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Performance Evaluation of MongoDB I/O access patterns

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NoSQL databases

- Principle category of Big Data Storage
- Overcome the scalability issues of SQL DBs
- Almost Cloud-friend systems
- Different architectures and data-models



MongoDB

- A document-based NoSQL database (JSON documents)
- It uses sharding for distributing data
- Different tools for easy-deployment and migration (Ex. mongorestore, mongodump)
- Simple to use and fit well with various kinds of apps.

Querying data

- NoSQL flexibility: store data; think how to use after!
- \blacksquare No schema \Rightarrow not only one way to query the data

frequent need to create indexes on demand, on large-scale data!

Querying data

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Problem statement

 MongoDB takes unjustified amount of time for making indexes on a pre-existed data

this problem could affect other NoSQL databases!

Benchmarking?

- ⇒ Tools such as Yahoo! Cloud Serving Benchmarks (YCSB) could:
 - \blacksquare report high level metrics (Ex. throughput, latency, ...) \checkmark
 - \blacksquare reduce the effort for comparing different systems \checkmark

But

- could not reflect the complexity of production-like data ②
- are not suitable for explaining performance issues (2)

Research objectives

- Identifying the performance issue(s) of MongoDB while indexing data
- Introducing new experimental tools for explaining the causes of performance issue(s)

Datasets

Dataset type	(min, avrg, max) in Kb	N. of data units	Size (Gb)
SDU_1	(1, 3, 6)	20,000,000 docs	71

Every doc x (int, date, 2 x string[min, max], array[1..4] x string[min, max])

MongoDB's storage engine: Wiredtiger

- The default storage engine for persistent data.
- __id_s for holding the collection's records
- collection is a separate file
- Btree data structure on disks
- Contiguous allocation on disk, but unspanned
- Default page size is 4 KB



FIGURE - a) Spanned & b) Unspanned allocation

Experimental setup

- Experiments are performed on Grid'5000 testbed Machines x 2 CPUs Intel Xeon E5-2630 v3, 8 cores/CPU, 128GB RAM, 2x558GB HDD, 10Gbps ethernet
- Ubuntu 14.04, Linux 4.9.13, MongoDB v3.4
- wiredTiger cache = 2Gb, replicas are disabled, sharding key is *_id*_s
- Considering only hash sharding in this talk \Rightarrow load balancing



FIGURE - Experimental scenarios











- \Rightarrow Indexing is a 100% read workload
- \Rightarrow I/O bounded operations
- \Rightarrow Data distribution issue
- ⇒ Benchmarking is not sufficient to look beyond!





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- ⇒ Benchmarking is not sufficient to look beyond!

 \Rightarrow How to get the main reason behind these results?

extended Berkeley Packet Filter (eBPF)

- Connect to all Linux data sources : (Kprobes, Uprobes, tracepoints, ...)
- Almost no overhead (4 ns per syscall) [A. Starovoitov]
- In-kernel verifier, JIT-compilation, mapping (kernel-userspace exchange)
- Processing, tracing, filtering inside Linux kernel.
- Suitable to work with system in-production & real data.

BPF compiler collection (BCC)

- Like-standard frontend project for eBPF tools
- No more byte code! \Rightarrow write your BPF code in restricted C
- Write your frontend code using Python or Lua
- Towards script-driven tracing



Tracing impact on our experiments

■ Overhead cost is lower than 0.8% of execution time in all Experiments. ☺

Basically, it depends on the number issued I/O requests.



 ${\rm I}/{\rm O}$ access pattern and latency on standalone config.

Experiments : eBPF results with SDU_1



 ${\sf I}/{\sf O}$ access pattern and latency on standalone config.

Experiments : eBPF results with SDU_1



 ${\rm I}/{\rm O}$ access pattern and latency on standalone config.



First shard results in two-shards config (50.9% of data)



Second shard results in two-shards config (49.1% of data)

- $\Rightarrow \mbox{ Acceptable access pattern on standalone}$
- ⇒ Almost random access on shards!
- \Rightarrow Data distribution is poorly done

Mismatch between the scanning table vs data stored on disk

Collection _ids: used by MongoDB process and WiredTiger

_id 1	_id 2	_id 3	_id 5
_id 7	_id 8	_id 10	_id 11
_id 19	_id 20	_id 27	

Collection file on disk: allocated regarding the key sharding (hashed id)

_id 1	_id 3	_id 7	_id 5	_id 8	_id 10	_id 2	_id 19	_id 11	_id 20	_id 30	_id 27
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 $\ensuremath{\operatorname{FIGURE}}$ – MongoDB scanning table Vs records order on the disk

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- Rebuilding the _id index \Rightarrow consider the pre-existed ordering. ③

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Our ad-hoc solution

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The worst case example!

- Data on one shard in a four conf. experiment (1/4 out of data)
- indexing time : Doubling the indexing time of the whole data on standalone config (3000 sec).
- indexing time after : six time faster! (about 500 sec)



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 $\label{eq:FIGURE-I/O} \begin{array}{c} FIGURE-I/O \text{ access pattern a} \end{array} \text{ before \& b) after applying } \\ \\ the solution \end{array}$

Conclusion

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- A performance study on MongoDB I/O access pattern is done
- A generic tool for testing I/O access patterns is introduced
- \blacksquare An ad hoc solution is proposed to correct MongoDB I/O access patterns
- We showed how a trivial data management issue could affect the performance
- We demonstrated how it is worthy to use dynamic tracing to go beyond benchmarking results

Future works

- Including other solutions like Cassandra in those tests.
- Trying to enrich eBPF tools which explain specific issues.

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Questions are welcome !