Differential properties of functions $x \mapsto x^{2^t-1}$ – extended version*–

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Abstract

We provide an extensive study of the differential properties of the functions $x \mapsto x^{2^t-1}$ over \mathbb{F}_{2^n} , for 1 < t < n. We notably show that the differential spectra of these functions are determined by the number of roots of the linear polynomials $x^{2^t} + bx^2 + (b+1)x$ where b varies in \mathbb{F}_{2^n} . We prove a strong relationship between the differential spectra of $x \mapsto x^{2^t-1}$ and $x \mapsto x^{2^s-1}$ for s = n - t + 1. As a direct consequence, this result enlightens a connection between the differential properties of the cube function and of the inverse function. We also determine the complete differential spectra of $x \mapsto x^{2^t-1}$ for $t \in \{\lfloor n/2 \rfloor, \lfloor n/2 \rfloor + 1, n-2\}$.

Keywords. Differential cryptanalysis, block cipher, S-box, power function, monomial, differential uniformity, APN function, permutation, linear polynomial, Kloosterman sum, cyclic codes.

1 Introduction

Differential cryptanalysis is the first statistical attack proposed for breaking iterated block ciphers. Its publication [4] then gave rise to numerous works which investigate the security offered by different types of functions regarding differential attacks. This security is quantified by the so-called *differential uniformity* of the Substitution box used in the cipher [22]. Most

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notably, finding appropriate S-boxes which guarantee that the cipher using them resist differential attacks is a major topic for the last twenty years, see e.g. [11, 16, 9, 6, 8].

Power functions, *i.e.*, monomial functions, form a class of suitable candidates since they usually have a lower implementation cost in hardware. Also, their particular algebraic structure makes the determination of their differential properties easier. However, there are only a few power functions for which we can prove that they have a low differential uniformity. Up to equivalence, there are two large families of such functions: a subclass of the quadratic power functions (a.k.a. Gold functions) and a subclass of the socalled Kasami functions. Both of these families contain some permutations which are APN over \mathbb{F}_{2^n} for odd *n* and differentially 4-uniform for even *n*. The other known power functions with a low differential uniformity correspond to "sporadic" cases in the sense that the corresponding exponents vary with n [17] and they do not belong to a large class: they correspond to the exponents defined by Welch [14, 10], by Niho [13, 18], by Dobbertin [15], by Bracken and Leander [7], and to the inverse function [21]. It is worth noticing that some of these functions seem to have different structures because they do not share the same differential spectrum. For instance, for a quadratic power function or a Kasami function, the differential spectrum has only two values, *i.e.*, the number of occurrences of each differential belongs to $\{0, \delta\}$ for some δ [5]. The inverse function has a very different behavior since its differential spectrum has three values, namely 0, 2 and 4 and, for each input difference, there is exactly one differential which is satisfied four times.

However, when classifying all functions with a low differential uniformity, it can be noticed that the family of all power functions $x \mapsto x^{2^{t}-1}$ over \mathbb{F}_{2^n} , with 1 < t < n, contains several functions with a low differential uniformity. Most notably, it includes the cube function and the inverse function, and also $x \mapsto x^{2^{(n+1)/2}-1}$ for n odd, which is the inverse of a quadratic function. At a first glance, this family of exponents may be of very small relevance because the involved functions have distinct differential spectra. Then, they are expected to have distinct structures. For this reason, one of the motivations of our study was to determine whether some link could be established between the differential properties of the cube function and of the inverse function. Our work then answers positively to this question since it exhibits a general relationship between the differential spectra of $x \mapsto x^{2^{t-1}}$ and $x \mapsto x^{2^{n-t+1}-1}$ over \mathbb{F}_{2^n} . We also determine the complete differential spectra of some other exponents in this family.

The rest of the paper is organized as follows. Section 2 recalls some defi-

nitions and some general properties of the differential spectrum of monomial functions. Section 3 then focuses on the differential spectra of the monomials $x \mapsto x^{2^{t}-1}$. First, the differential spectrum of any such function is shown to be determined by the number of roots of a family of linear polynomials. Then, we exhibit a symmetry property for the exponents in this family: it is proved that the differential spectra of $x \mapsto x^{2^{t}-1}$ and $x \mapsto x^{2^{n-t+1}-1}$ over $\mathbb{F}_{2^{n}}$ are closely related. In Section 5, we determine the whole differential spectrum of $x \mapsto x^{7}$ over $\mathbb{F}_{2^{n}}$. It is expressed by means of some Kloosterman sums, and explicitly computed using the work of Carlitz [12]. We then derive the differential spectra of $x \mapsto x^{2^{n-2}-1}$. Further, we study the functions $x \mapsto x^{2^{\lfloor n/2 \rfloor}-1}$ and $x \mapsto x^{2^{\lfloor n/2 \rfloor+1}-1}$. We finally end up with some conclusions.

2 Preliminaries

2.1 Functions over \mathbb{F}_{2^n} and their derivatives

Any function F from \mathbb{F}_{2^n} into \mathbb{F}_{2^n} can be expressed as a univariate polynomial in $\mathbb{F}_{2^n}[X]$. The *univariate degree* of the polynomial F is, as usual, the maximal integer value of its exponents. The *algebraic degree* of F is the maximal Hamming weight of its exponents:

$$\deg \left(\sum_{i=0}^{2^n-1} \lambda_i X^i\right) = \max \{wt(i) \mid \lambda_i \neq 0\},\$$

where $\lambda_i \in \mathbb{F}_{2^n}$ and the *Hamming weight* is calculated as follows :

$$i = \sum_{j=0}^{n-1} i_j 2^j$$
 with $i_j \in \{0,1\}, wt(i) = \sum_{j=0}^{n-1} i_j$.

In this paper, we will identify a polynomial of $\mathbb{F}_{2^n}[X]$ with the corresponding function over \mathbb{F}_{2^n} . For instance, $F \in \mathbb{F}_{2^n}[X]$ is called a *permutation polynomial* of \mathbb{F}_{2^n} if the function $x \mapsto F(x)$ is a permutation of \mathbb{F}_{2^n} .

Boolean functions are also involved in this paper and are generally of the form

$$x \in \mathbb{F}_{2^n} \quad \mapsto \quad Tr(P(x)) \in \mathbb{F}_2,$$

where P is any function from \mathbb{F}_{2^n} into \mathbb{F}_{2^n} and where Tr denotes the absolute trace on \mathbb{F}_{2^n} , *i.e.*,

$$Tr(\beta) = \beta + \beta^2 + \dots + \beta^{2^{n-1}}, \ \beta \in \mathbb{F}_{2^n}$$

In the whole paper, #E denotes the cardinality of any set E.

The resistance of a cipher to differential attacks and to its variants is quantified by some properties of the *derivatives* of its S(ubstitution)-box, in the sense of the following definition. It is worth noticing that this definition is general: it deals with functions from \mathbb{F}_{2^n} into \mathbb{F}_{2^m} for any $m \geq 1$.

Definition 1 Let F be a function from \mathbb{F}_{2^n} into \mathbb{F}_{2^m} . For any $a \in \mathbb{F}_{2^n}$, the derivative of F with respect to a is the function D_aF from \mathbb{F}_{2^n} into \mathbb{F}_{2^m} defined by

$$D_a F(x) = F(x+a) + F(x), \quad \forall x \in \mathbb{F}_{2^n}.$$

The resistance to differential cryptanalysis is related to the following quantities, introduced by Nyberg and Knudsen [22, 21].

Definition 2 Let F be a function from \mathbb{F}_{2^n} into \mathbb{F}_{2^n} . For any a and b in \mathbb{F}_{2^n} , we denote

$$\delta(a,b) = \#\{x \in \mathbb{F}_{2^n}, \ D_a F(x) = b\}.$$

Then, the differential uniformity of F is

$$\delta(F) = \max_{a \neq 0, \ b \in \mathbb{F}_{2^n}} \delta(a, b).$$

Those functions for which $\delta(F) = 2$ are said to be almost perfect nonlinear (APN).

2.2 Differential spectrum of power functions

In this paper, we focus on the case where the S-box is a power function, *i.e.*, a monomial function on \mathbb{F}_{2^n} . In other words, $F(x) = x^d$ over \mathbb{F}_{2^n} , which will be denoted by F_d when necessary. In the case of such a power function, the differential properties can be analyzed more easily since, for any nonzero $a \in \mathbb{F}_{2^n}$, the equation $(x + a)^d + x^d = b$ can be written

$$a^d \left(\left(\frac{x}{a}+1\right)^d + \left(\frac{x}{a}\right)^d \right) = b,$$

implying that

$$\delta(a,b) = \delta(1,b/a^d)$$
 for all $a \neq 0$.

Then, when $F: x \mapsto x^d$ is a monomial function, the differential characteristics of F are determined by the values $\delta(1, b), b \in \mathbb{F}_{2^n}$. From now on, this quantity $\delta(1, b)$ is denoted by $\delta(b)$. Since

$$\#\{b \in \mathbb{F}_{2^n} | \delta(a, b) = i\} = \#\{b \in \mathbb{F}_{2^n} | \delta(b) = i\} \quad \forall a \neq 0,$$

the differential spectrum of F can be defined as follows.

Definition 3 Let $F(x) = x^d$ be a power function on \mathbb{F}_{2^n} . We denote by ω_i the number of output differences b that occur i times:

$$\omega_i = \#\{b \in \mathbb{F}_{2^n} | \delta(b) = i\}.$$

$$\tag{1}$$

 \diamond

The differential spectrum of F_d is the set of ω_i :

$$\mathbb{S} = \{\omega_0, \omega_2, ..., \omega_{\delta(F)}\}.$$

With same notation, we have the following equalities. They are well-known but we indicate the proof for clarity.

Lemma 1

$$\sum_{k=0}^{2^{n}} \omega_{k} = 2^{n} \text{ and } \sum_{k=2}^{2^{n}} (k \times \omega_{k}) = 2^{n},$$

where $\omega_i = 0$ for *i* odd.

Proof. The first equality is obviously deduced from (1). And, for k > 0, $k \times \omega_k$ equals the number of $x \in \mathbb{F}_{2^n}$ such that

$$x^d + (x+1)^d = b$$
 and $\delta(b) = k$

for some b. Thus, any x is counted in the second sum.

Remark 1 The differential spectrum of the power function $F(x) = x^d$ over \mathbb{F}_{2^n} is also related to the weight enumerator of the cyclic code of length $(2^n - 1)$ with defining set $\{1, s\}$ [11]. In particular, the number of codewords with Hamming weight 3 and 4 in this cyclic code can be derived from the differential spectrum of F (see e.g. Corollary 1 in [5]).

A power function F is said to be *differentially* 2-valued if and only if for any $b \in \mathbb{F}_{2^n}$, we have $\delta(b) \in \{0, \kappa\}$ (and then only two ω_i in \mathbb{S} do not vanish). It is known that $\kappa = 2^r$ for some r > 1 (see an extensive study in [5, Section 5]). Note that APN functions are differentially 2-valued with $\kappa = 2$.

There are some basic transformations which preserve S.

Lemma 2 Let $F_d(x) = x^d$ and $F_e(x) = x^e$ over \mathbb{F}_{2^n} . If there exists k such that $e = 2^k d \mod 2^n - 1$ or if $ed = 1 \mod 2^n - 1$, then F_e has the same differential spectrum as F_d .

2.3 General properties on the differential spectrum

In this section, $F_d(x) = x^d$ and notation is as in Section 2.2. Studying $\delta(b)$ for special values of b may give us as least a lower bound on $\delta(F_d)$. So we first focus on $\delta(0)$.

Lemma 3 Let d be such that $gcd(d, 2^n - 1) = s$. Then $F_d : x \mapsto x^d$ is such that $\delta(0) = s - 1$. In particular s = 1 if and only if $\delta(0) = 0$.

Proof. Note that s = 1 if and only if F_d is a permutation. Obviously, x is a solution of $x^d + (x+1)^d = 0$ if and only if

$$\left(\frac{x+1}{x}\right)^d = 1$$
 that is $x+1 = xz$ with $z^d = 1$,

since $x \mapsto (x+1)/x$ is a permutation over $\mathbb{F}_{2^n} \setminus \{0, 1\}$. As there are exactly s-1 such nonzero z, the proof is completed.

There is an immediate consequence of Lemma 3 for specific values of d.

Proposition 1 Let $d \ge 3$ such that d divides $2^n - 1$. Then $\delta(F_d) = \delta(0) = d - 1$.

In particular, if $d = 2^t - 1$ with gcd(t, n) = t then $\delta(F_d) = \delta(0) = 2^t - 2$.

Proof. Since $gcd(d, 2^n - 1) = d$, $\delta(0) = d - 1$ from Lemma 3. But the polynomial $x^d + (x+1)^d + b$ has degree d-1 for any b, so that $\delta(b) \leq d-1$. We conclude that $\delta(F_d) = d - 1$.

Now, let $d = 2^t - 1$ with gcd(t, n) = t. Then $gcd(d, 2^n - 1) = 2^t - 1$ so that $\delta(0) = 2^t - 2$. As previously we conclude that $\delta(F_d) = 2^t - 2$.

Example 1 If d = 3 then $\delta(F_d) = \delta(0) = 2$ for any even n. If d = 5 then $\delta(F_d) = \delta(0) = 4$ for n = 4k for all k > 1. If d = 7 then $\delta(F_d) = \delta(0) = 6$ for n = 3k for all k > 1.

The previous remarks combined with our simulation results point out that $\delta(0)$ and $\delta(1)$ play a very particular role in the differential spectra of power functions. This leads us to investigate the properties of the differential spectrum restricted to the values $\delta(b)$ with $b \notin \mathbb{F}_2$.

Definition 4 Let F be a power function on \mathbb{F}_{2^n} . We say that F has the same restricted differential spectrum as an APN function when

$$\delta(b) \leq 2 \text{ for all } b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2.$$

For the sake of simplicity, we will say that F is locally-APN.

This definition obviously generalizes the APN property. For instance, the *inverse* function over \mathbb{F}_{2^n} is locally-APN for any n, while it is APN for odd n only. Another infinite class of locally-APN functions is exhibited in Section 5.2.

3 The differential spectrum of $x \mapsto x^{2^t-1}$

From now on, we investigate the differential spectra of the following specific monomial functions

$$G_t: x \mapsto x^{2^t-1}, \ 2 \le t \le n-1, \text{ over } \mathbb{F}_{2^n}.$$

Note that such a function has algebraic degree t.

3.1 Link with linear polynomials

In this section, we first give some general properties.

Theorem 1 Let $G_t(x) = x^{2^t-1}$ over \mathbb{F}_{2^n} with $2 \le t \le n-1$. Then,

$$G_t(x+1) + G_t(x) + 1 = \frac{(x^{2^{t-1}} + x)^2}{x^2 + x}.$$
(2)

Consequently, for any $b \in \mathbb{F}_{2^n} \setminus \{1\}$, $\delta(b)$ is the number of roots in $\mathbb{F}_{2^n} \setminus \mathbb{F}_2$ of the linear polynomial

$$P_b(x) = x^{2^t} + bx^2 + (b+1)x$$
.

And we have

$$\begin{aligned} \delta(0) &= 2^{\gcd(t,n)} - 2\\ \delta(1) &= 2^{\gcd(t-1,n)} \end{aligned}$$
for any $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2, \ \delta(b) &= 2^r - 2 \end{aligned}$

for some r with $1 \le r \le \min(t, n - t + 1)$.

Proof. To prove (2) we simply compute

$$(x+x^{2})(1+x^{2^{t}-1}+(1+x)^{2^{t}-1}) = x+x^{2}+x^{2^{t}}+x^{2^{t}+1}+x(1+x)^{2^{t}} = x^{2}+x^{2^{t}}.$$

Thus, $\delta(1)$ is directly deduced and it corresponds to the number of roots of $P_1(x) = (x^{2^{t-1}} + x)^2$. Let $b \in \mathbb{F}_{2^n} \setminus \{1\}$. Then $x \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2$ is a solution of

$$(x+1)^d + x^d = b, \ d = 2^t - 1,$$

if and only if it is a solution of

$$(x^{2^{t-1}} + x)^2 = (b+1)x(x+1),$$

or equivalently if it is a root of the linear polynomial

$$P_b(x) = x^{2^t} + bx^2 + (b+1)x.$$

The values x = 0 and x = 1 are counted in $\delta(1)$ (as solutions of $(x+1)^d + x^d = 1$), while $P_b(0) = P_b(1) = 0$ for any b. So, we get that, if $b \neq 1$, the number of roots of P_b in \mathbb{F}_{2^n} is equal to $(\delta(b) + 2)$. Because the set of all roots of a linear polynomial is a linear space, we deduce that

$$\forall b \in \mathbb{F}_{2^n} \setminus \{1\}, \ \delta(b) = 2^r - 2 \text{ with } r \leq t.$$

Moreover, by raising P_b to the 2^{n-t} -th power, we get that any root of P_b is also a root of

$$b'x^{2^{n-t+1}} + (b'+1)x^{2^{n-t}} + a$$

with $b' = b^{2^{n-t}}$. This then implies that $\delta(b) = 2^r - 2$ with $r \leq n - t + 1$. Finally, for b = 0, $P_0(x) = x^{2^t} + x$, implying that $\delta(0) = 2^{\gcd(t,n)} - 2$, which naturally corresponds to Lemma 3.

Remark 2 As a first easy corollary, we recover the following well-known form of the differential spectrum of the inverse function, $G_{n-1} : x \mapsto x^{2^{n-1}-1}$ over \mathbb{F}_{2^n} . Actually, the previous theorem applied to t = n-1 leads to $\delta(0) = 0$ and $\delta(1) = 2$ when n is odd and $\delta(1) = 4$ when n is even. For all $b \notin \mathbb{F}_2$, $\delta(b) \in \{0, 2\}$. Therefore, we have

- if n is odd, $\delta(G_{n-1}) = 2$ and $\omega_0 = 2^{n-1}$, $\omega_2 = 2^{n-1}$;
- if n is even, $\delta(G_{n-1}) = 4$ and $\omega_0 = 2^{n-1} + 1$, $\omega_2 = 2^{n-1} 2$, $\omega_4 = 1$.

Clearly G_{n-1} is locally-APN for any n, as we previously noticed (see Definition 4).

The following corollary is a direct consequence of Theorem 1.

Corollary 1 Let $G_t(x) = x^{2^t-1}$ over \mathbb{F}_{2^n} with $2 \leq t \leq n-1$. Then, its differential uniformity is of the form either $2^r - 2$ or 2^r for some $2 \leq r \leq n$. Moreover, if $\delta(G_t) = 2^r$ for some r > 1, then this value appears only once in the differential spectrum, i.e., $\omega_{2^r} = 1$, and it corresponds to the value of $\delta(1)$, implying $\delta(G_t) = 2^{\gcd(t-1,n)}$.

3.2 Equivalent formulations

In Theorem 1, we exhibited some tools for the computation of the differential spectra of functions $x \mapsto x^{2^{t-1}}$. The problem boils down to the determination of the roots of a linear polynomial whose coefficients depend on $b \in \mathbb{F}_{2^n}$. There are equivalent formulations that we are going to develop now. The first one is obtained by introducing another class of linear polynomials over \mathbb{F}_{2^n} . For any subspace E of \mathbb{F}_{2^n} (where \mathbb{F}_{2^n} is identified with \mathbb{F}_2^n), we define its dual as follows:

$$E^{\perp} = \{ x \mid Tr(xy) = 0, \forall y \in E \}.$$

Also, we denote by $\mathcal{I}m(F)$ the image set of any function F.

Lemma 4 Let $t, s \ge 2$ and s = n - t + 1. Let us consider the linear applications

$$P_{t,b}(x) = x^{2^t} + bx^2 + (b+1)x, \quad b \in \mathbb{F}_{2^n}$$

Then the dual of $\mathcal{I}m(P_{t,b})$ is the set of all α satisfying $P_{t,b}^*(\alpha) = 0$ where

$$P_{t,b}^*(x) = x^{2^s} + (b+1)^2 x^2 + bx.$$

Note that $P_{t,b}^*$ is called the adjoint application of $P_{t,b}$.

Proof. By definition, $\mathcal{I}m(P_{t,b})^{\perp}$ consists of all α such that $Tr(\alpha P_{t,b}(x)) = 0$ for all $x \in \mathbb{F}_{2^n}$. We have

$$Tr(\alpha P_{t,b}(x)) = Tr(\alpha x^{2^{t}}) + Tr(b\alpha x^{2}) + Tr(\alpha(b+1)x)$$

= $Tr(\alpha^{2^{n-t+1}}x^{2}) + Tr(b\alpha x^{2}) + Tr(\alpha^{2}(b+1)^{2}x^{2})$
= $Tr(x^{2}(\alpha^{2^{s}} + \alpha^{2}(b+1)^{2} + \alpha b)).$

Hence α belongs to the dual of the image of $P_{t,b}$ if and only if $\alpha^{2^s} + \alpha^2(b+1)^2 + \alpha b = 0$, *i.e.*, α is a root of $P_{t,b}^*$, completing the proof.

The following theorem gives an equivalent formulation of the quantity r which is presented in Theorem 1.

Theorem 2 Notation is as in Lemma 4. Then

$$\dim Ker(P_{t,b}) = \dim Ker(P_{t,b}^*).$$

Consequently, this dimension can be determined by solving $P_{t,b}(x) = 0$ or equivalently by solving

$$x^{2^{s}} + (b+1)^{2}x^{2} + bx = 0$$
, where $s = n - t + 1$.

Proof. Let κ be the dimension of the image set of $P_{t,b}$. It is well-known that $n = \kappa + \dim Ker(P_{t,b})$. On the other hand, Lemma 4 shows that α is in the dual of the image of $P_{t,b}$ if and only if $P_{t,b}^*(\alpha) = 0$. We deduce that

$$n - \kappa = \dim Ker(P_{t,b}^*) = \dim Ker(P_{t,b}) ,$$

completing the proof.

 \diamond

Now, we discuss a different point of view, using an equivalent linear system.

Theorem 3 For any $2 \le t < n$, we define the following equations:

$$E_b$$
 : $x^{2^{\iota}} + bx^2 + (b+1)x = 0, \ b \in \mathbb{F}_{2^n}.$

Let N_b be the number of solutions of E_b in $\mathbb{F}_{2^n} \setminus \mathbb{F}_2$. Let M_b be the number of solutions in $\mathbb{F}_{2^n}^*$ of the system

$$y^{2^{t-1}} + \dots + y^2 + y(b+1) = 0 Tr(y) = 0$$

Then $N_b = 2 \times M_b$.

Proof. We simply write

$$x^{2^{t}} + bx^{2} + (b+1)x = x^{2^{t}} + x + b(x^{2} + x)$$

which is equal to

$$= (x^{2} + x)^{2^{t-1}} + (x^{2} + x)^{2^{t-2}} + \dots + (x^{2} + x) + b(x^{2} + x)$$

= $y^{2^{t-1}} + y^{2^{t-2}} + \dots + y^{2} + y(b+1)$, with $y = x^{2} + x$.

We are looking at the number of solutions of E_b which are not in \mathbb{F}_2 . So, it is equivalent to compute the number of nonzero solutions y of

$$y^{2^{t-1}} + y^{2^{t-2}} + \dots + y^2 + y(b+1) = 0$$

such that the equation $x^2 + x + y = 0$ has solutions. This last condition holds if and only if Tr(y) = 0, providing two distinct solutions $x_1, x_2 = x_1 + 1$ such that $x_i^2 + x_i = y$, completing the proof.

Remark 3 In Theorem 3, b takes any value while P_b is defined for $b \neq 1$ in Theorem 1. For all $b \neq 1$, we have clearly $N_b = \delta(b)$. If b = 1, $P_1(x) = x^{2^t} + x^2$ and the number of roots of P_1 in \mathbb{F}_{2^n} is equal to

$$N_1 + 2 = 2^{\gcd(t-1,n)} = \delta(1).$$

Therefore, we have $M_1 = \delta(1)/2 - 1$.

4 A property of symmetry

Recall that $G_t(x) = x^{2^t-1}$. Now, we are going to examine some symmetries between the differential spectra of G_t and G_s where $t, s \ge 2$ and s = n-t+1. In the list of properties below, notation is conserved as soon it is defined. Recall that

$$P_{t,b}^*(x) = x^{2^s} + x^2(b+1)^2 + xb$$

is the adjoint polynomial of $P_{t,b}(x) = x^{2^t} + bx^2 + (b+1)x$. Thus, both polynomials have a kernel with the same dimension (see Lemma 4 and Theorem 2). It is worth noticing that this dimension is at least 1 since $P_{t,b}(0) = P_{t,b}(1) = 0$. In this section we want to prove the following theorem.

Theorem 4 For any ν with $2 \leq \nu \leq n-1$, we define

$$S^{i}_{\nu} = \{ b \mid \dim Ker(P_{\nu,b}) = i \} \text{ with } 1 \le i \le \nu .$$

Then, for any $s, t \ge 2$ with t = n - s + 1 and for any i, we have $\#S_s^i = \#S_t^i$.

We begin by some lemmas. The next one will not be used for the proof of Theorem 4 but clarifies some arguments.

Lemma 5 Let $a \in \mathbb{F}_{2^n}^*$ and $2 \leq t \leq n-1$. Then there are exactly two elements, b_1 and b_2 with $b_2 = b_1 + a^{-1}$, such that $P_{t,b_i}^*(a) = 0$ for i = 1, 2. In particular, $P_{t,b}^*(1) = 0$ for $b \in \{0, 1\}$.

Proof. Let a be fixed and let us consider the equation $P_{t,b}^*(a) = 0$ for some b:

$$b^{2}a^{2} + ba + a^{2^{s}} + a^{2} = a^{2}\left(b^{2} + \frac{b}{a} + \frac{a^{2^{s}} + a^{2}}{a^{2}}\right) = 0$$

There is b such that this equation is satisfied if and only if

$$Tr\left(\frac{a^{2^{s}}+a^{2}}{a^{2}}\times a^{2}\right) = Tr(a^{2^{s}}+a^{2}) = 0,$$

which holds for any a. Thus, for any nonzero a there are exactly two solutions, say b_1 and b_2 whose sum equals a^{-1} . To complete the proof, we observe that $P_{t,b}^*(1) = b^2 + b$.

Lemma 6 Let $s,t \geq 2$ with t = n - s + 1. Let π be the permutation of $\mathbb{F}_{2^n}^* \times \mathbb{F}_{2^n}$ defined by

$$\pi(a,b) = \left(a^{2^s}, \frac{ab}{a^{2^s}} + 1\right) \;.$$

Then, for any (a, b) in $\mathbb{F}_{2^n}^* \times \mathbb{F}_{2^n}$, $(\alpha, \beta) = \pi(a, b)$ satisfies

$$P^*_{s,\beta}(\alpha) = P^*_{t,b}(\alpha) \; .$$

Proof. First, we clearly have that π is a permutation of $\mathbb{F}_{2^n}^* \times \mathbb{F}_{2^n}$. Indeed, $\pi(\mathbb{F}_{2^n}^* \times \mathbb{F}_{2^n}) \subset \mathbb{F}_{2^n}^* \times \mathbb{F}_{2^n}$ and one can define the inverse of π as follows:

$$\pi^{-1}(\alpha,\beta) = \left(\alpha^{2^{n-s}}, \frac{\alpha(\beta+1)}{\alpha^{2^{n-s}}}\right) \ .$$

Actually, $(\alpha^{2^{n-s}})^{2^s} = \alpha$ and it can be checked that

$$\pi(\pi^{-1}(\alpha,\beta)) = \left(\alpha, \frac{\alpha^{2^{n-s}}\alpha(\beta+1)}{\alpha\alpha^{2^{n-s}}} + 1\right) = (\alpha,\beta).$$

Then, by using that $(\beta + 1)^2 = \frac{a^2b^2}{a^{2^{s+1}}}$ and s + t = n + 1, we deduce that

$$P_{s,\beta}^*(\alpha) = (a^{2^s})^{2^t} + (a^{2^s})^2(\beta+1)^2 + (a^{2^s})\beta$$

= $a^2 + a^2b^2 + ab + a^{2^s}$
= $P_{t,b}^*(a)$.

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Lemma 7 Let $s, t \ge 2$ with t = n - s + 1. Let $b \in \mathbb{F}_{2^n}$ and let $a \in \mathbb{F}_{2^n}^*$ such that $P_{t,b}^*(a) = 0$. Then dim $Ker(P_{t,b}^*) = \dim Ker(P_{s,\beta}^*)$, where $\beta = 1 + ab/a^{2^s}$.

Proof. Recall that $P_{t,b}^*(x) = x^{2^s} + x^2(b+1)^2 + xb$. We know that for any $b \notin \mathbb{F}_2$ there is $a \in \mathbb{F}_{2^n} \setminus \{0,1\}$ such that $P_{t,b}^*(a) = 0$. This is because dim $Ker(P_{t,b}) = \dim Ker(P_{t,b}^*)$ (see Theorem 2) and $\{0,1\}$ is included in the kernel of $P_{t,b}$. Moreover, $P_{t,b}^*(1) = b^2 + b = 0$ if and only if $b \in \mathbb{F}_2$.

We treat the case a = 1 separately, a case where $P_{t,b}(a) = 0$ for $b \in \mathbb{F}_2$ only. In this case, Lemma 6 leads to $P^*_{s,\beta}(1) = 0$ too where $\beta = b + 1$, since $\pi(1,b) = (1,b+1)$. And we have for b = 0

$$P_{t,0}^*(x) = x^{2^s} + x^2 = P_{s,1}(x)$$

and for b = 1

$$P_{t,1}^*(x) = x^{2^s} + x = P_{s,0}(x).$$

Thus, we conclude: for a = 1, if b is such that $P^*_{t,b}(1) = 0$ then $\beta = b + 1$ and

$$\dim Ker(P_{t,b}^*) = \dim Ker(P_{s,\beta}) = \dim Ker(P_{s,\beta}^*)$$

where the last equality comes from Theorem 2.

Now, we suppose that $a \notin \mathbb{F}_2$. With x = ay, the equation $P_{t,b}^*(x) = 0$ is equivalent to

$$a^{2^{s}}y^{2^{s}} + a^{2}y^{2}(b+1)^{2} + ayb = 0$$

which is

$$a^{2^{s}}\left(y^{2^{s}} + \frac{a^{2}(b+1)^{2}}{a^{2^{s}}}y^{2} + y\frac{ab}{a^{2^{s}}}\right) = 0.$$

We can set

$$\beta = \frac{a^2(b+1)^2}{a^{2^s}}$$
 and $\beta + 1 = \frac{ab}{a^{2^s}}$,

since

$$\frac{a^2(b+1)^2}{a^{2^s}} + 1 = \frac{ab}{a^{2^s}}$$

is equivalent to

$$a^{2^s} + a^2(b+1)^2 + ab = 0, \ i.e., P^*_{t,b}(a) = 0.$$

We have proved that $P_{t,b}^*(x) = 0$ is equivalent to

$$P_{s,\beta}(y) = y^{2^s} + \beta y^2 + (\beta + 1)y = 0.$$

Then, dim $Ker(P_{s,\beta}) = \dim Ker(P_{t,b}^*)$. But dim $Ker(P_{s,\beta}) = \dim Ker(P_{s,\beta}^*)$ by Theorem 2, completing the proof.

Proof of Theorem 4. Recall that

$$S_{\nu}^{i} = \{ b \in \mathbb{F}_{2^{n}} \mid \dim Ker(P_{\nu,b}) = i \}$$

Then, we want to show that, for any i, $\#S_t^i = \#S_s^i$. For any $2 \le \nu \le n-1$ and for any $1 \le i \le \nu$, we define

$$\mathcal{E}_{\nu}^{i} = \{(a,b) \in \mathbb{F}_{2^{n}}^{*} \times \mathbb{F}_{2^{n}} \mid P_{\nu,b}^{*}(a) = 0 \text{ and } \dim Ker(P_{\nu,b}) = i \}.$$

From Theorem 2, we know that dim $Ker(P_{\nu,b}) = \dim Ker(P_{\nu,b}^*)$. Then,

$$\mathcal{E}_{\nu}^{i} = \{(a,b) \in \mathbb{F}_{2^{n}}^{*} \times \mathbb{F}_{2^{n}} \mid P_{\nu,b}^{*}(a) = 0 \text{ and } \dim Ker(P_{\nu,b}^{*}) = i \}.$$

For any $b \in S^i_{\nu}$ there are $2^i - 1$ nonzero a in $Ker(P^*_{\nu,b})$ and then $2^i - 1$ pairs (a, b), for a fixed b, in \mathcal{E}^i_{ν} so that

$$#\mathcal{E}^{i}_{\nu} = (2^{i} - 1) #S^{i}_{\nu} .$$
(3)

We use Lemma 6. Recall that π is the permutation of $\mathbb{F}_{2^n}^* \times \mathbb{F}_{2^n}$ defined by

$$\pi(a,b) = \left(a^{2^s}, \frac{ab}{a^{2^s}} + 1\right).$$

Then, we have

$$\begin{aligned} \mathcal{E}_t^i &= \{(a,b) \in \mathbb{F}_{2^n}^* \times \mathbb{F}_{2^n} \mid P_{t,b}^*(a) = 0 \text{ and } \dim Ker(P_{t,b}^*) = i \}, \\ \mathcal{E}_s^i &= \{(\alpha,\beta) \in \mathbb{F}_{2^n}^* \times \mathbb{F}_{2^n} \mid P_{s,\beta}^*(\alpha) = 0 \text{ and } \dim Ker(P_{s,\beta}^*) = i \} \\ &= \{(\alpha,\beta) = \pi(a,b), (a,b) \in \mathcal{E}_t^i\}. \end{aligned}$$

Indeed, any (α, β) is as follows specified from (a, b). We have $P_{s,\beta}^*(\alpha) = P_{t,b}^*(a)$ from Lemma 6. Moreover, according to Lemma 7, dim $Ker(P_{t,b}^*) = \dim Ker(P_{s,\beta}^*)$, where β is calculated from a and b, for any a such that $P_{t,b}^*(a) = 0$.

In other terms, to any pair $(a, b) \in \mathcal{E}_t^i$ corresponds a unique pair $(\alpha, \beta) \in \mathcal{E}_s^i$. We finally get that $\#\mathcal{E}_s^i = \#\mathcal{E}_t^i$ and it directly follows from (3) that $\#S_s^i = \#S_t^i$, completing the proof. \diamond

Now we are going to explain Theorem 4, in terms of the differential spectra of G_t and G_s , $s, t \ge 2$ with t = n - s + 1. Actually, we can deduce from the previous theorem that both functions G_t and G_s have the same restricted differential spectrum, *i.e.* the multisets $\{\delta(b), b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2\}$ are the same for both functions.

Corollary 2 We denote by $\delta_{\nu}(b)$, $b \in \mathbb{F}_{2^n}$, the quantities $\delta(b)$ corresponding to $G_{\nu} : x \mapsto x^{2^{\nu}-1}$. Then, for any $s, t \geq 2$ with t = n - s + 1, we have

$$\begin{aligned} \delta_s(0) &= \delta_t(1) - 2 &= 2^{\gcd(t-1,n)} - 2 \\ \delta_s(1) &= \delta_t(0) + 2 &= 2^{\gcd(t,n)} \end{aligned}$$

and we have equality between both following multisets:

$$\{\delta_s(b), b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2\} = \{\delta_t(b), b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2\}.$$
(4)

Moreover, G_t and G_s have the same differential spectrum if and only if

$$gcd(s,n) = gcd(t,n) = 1$$

which can hold for odd n only. In any case, G_t is locally-APN if and only if G_s is locally-APN.

Proof. Since s = n - t + 1, we clearly have

$$gcd(s,n) = gcd(t-1,n)$$
 and $gcd(s-1,n) = gcd(t,n)$.

Thus, applying Theorem 1, we get

$$\delta_s(0) = 2^{\gcd(s,n)} - 2 = 2^{\gcd(t-1,n)} - 2 = \delta_t(1) - 2$$

and

$$\delta_s(1) = 2^{\gcd(s-1,n)} = 2^{\gcd(t,n)} = \delta_t(0) + 2$$

Moreover, we have

$$(P_{t,1}(x))^{2^{s-1}} = (x^{2^t} + x^2)^{2^{s-1}} = x + x^{2^s} = P_{s,0}(x)$$

$$(P_{t,0}(x))^{2^s} = (x^{2^t} + x)^{2^s} = x^{2^s} + x^2 = P_{s,1}(x) ,$$

implying that

$$\{\dim KerP_{t,0}, \dim KerP_{t,1}\} = \{\dim KerP_{s,0}, \dim KerP_{s,1}\}.$$

We deduce from Theorem 4 that

$$\#\{b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2 \mid \dim Ker(P_{t,b}) = i\} = \#\{b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2 \mid \dim Ker(P_{s,b}) = i\}.$$

Equality (4) is then a direct consequence of Theorem 1, since

$$\{\delta_{\nu}(b), b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2\} = \{2^{\kappa(b)} - 2, \ \kappa(b) = \dim Ker(P_{\nu,b})\}.$$

Now, we note that $\delta_s(0) = \delta_t(0)$ if and only if $\delta_s(1) = \delta_t(1)$. Thus, G_t and G_s have the same differential spectrum if and only if $\delta_s(0) = \delta_t(0)$. Since

$$\delta_s(0) = 2^{\gcd(s,n)} - 2$$
 and $\delta_t(0) = 2^{\gcd(t,n)} - 2$,

this holds if and only if gcd(t, n) = gcd(s, n) = 1. It cannot hold when n is even, because in this case either s or t is even too.

 \diamond

Using Definition 4, the last statement is obviously derived.

The previous result implies that, if G_t is APN over \mathbb{F}_{2^n} , then G_s is locally-APN. Moreover, the differential spectrum of G_s can be completely determined as shown by the following corollary.

Corollary 3 Let n and t < n be two integers such that $G_t : x \mapsto x^{2^t-1}$ is APN over \mathbb{F}_{2^n} . Let s = n - t + 1. Then,

- if n is odd, both G_t and G_s are APN permutations;
- if n is even, G_t is not a permutation and G_s is a differentially 4uniform permutation (locally-APN) with the following differential spectrum: $\omega_4 = 1$, $\omega_2 = 2^{n-1} - 2$ and $\omega_0 = 2^{n-1} + 1$.

Proof. From Theorem 1, we deduce that, if F is APN, then $\delta_t(b) \in \{0, 2\}$ for all $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2$; moreover, gcd(n, t - 1) = 1 and $gcd(n, t) \in \{1, 2\}$ since $\delta_t(1) = 2$ and $\delta_t(0) \in \{0, 2\}$.

If n is odd, gcd(n,t) = 1 is then the only possible value, implying that $\delta_t(0) = 0$. It follows that $\delta_s(0) = 0$, $\delta_s(1) = 2$ and $\delta_s(b) \in \{0,2\}$ for all $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2$. In other words, both G_t and G_s are APN permutations.

If n is even, it is well-known that G_t is not a permutation (see e.g. [2]). More precisely, we have here gcd(n,t) = 2 since t and t-1 cannot be both coprime with n. Then, we deduce that $\delta_s(0) = 0$ and $\delta_s(1) = 4$. The differential spectrum of G_s directly follows from Corollary 2.

Example 2 Notation is as in Corollary 3. For t = 2, we have $G_t(x) = x^3$. It is well-known that G_2 is an APN function over \mathbb{F}_{2^n} for any n. Since s = n - 1, $G_s(x)$ is equivalent to the inverse function and it is also well-known that the inverse function is APN for odd n. For even n, $\delta(G_{n-1}) = 4$ and the differential spectrum is computed in Remark 2.

Corollary 4 Let n and t < n be two integers such that $G_t : x \mapsto x^{2^t-1}$ is differentially 4-uniform. Then, n is even and G_t is a permutation with the following differential spectrum: $\omega_4 = 1$, $\omega_2 = 2^{n-1} - 2$ and $\omega_0 = 2^{n-1} + 1$. Moreover, for s = n - t + 1, G_s is APN.

Proof. From Corollary 1, we deduce that $\delta(G_t) = 4$ implies gcd(n, t-1) = 2 and $\omega_4 = 1$. In particular, *n* is even. Since gcd(n, t-1) and gcd(n, t) cannot be both equal to 2, we also deduce that that G_t is a permutation. Its differential spectrum is then derived from Lemma 1.

Moreover, we have $\delta_s(0) = 2$ and $\delta_s(1) = 2$, implying that G_s is APN. \diamond

5 Specific classes

In this section, we apply the results of Section 3 to the study of the differential spectrum of $G_t: x \mapsto x^{2^t-1}$, for special values of t.

5.1 The function $x \mapsto x^7$

We first focus on $G_3 : x \mapsto x^7$ over \mathbb{F}_{2^n} , *i.e.*, t = 3. In this case, we determine the complete differential spectrum of the function. Moreover, thanks to the work of Carlitz [12], we emphasize that this spectrum is related to some Kloosterman sums defined as follows.

Proposition 2 [12, Formula (6.8)] Let K(1) be the Kloosterman sum

$$K(1) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{Tr(x^{-1} + x)}$$

extended to 0 assuming that $(-1)^{Tr(x^{-1})} = 1$ for x = 0. Then,

$$K(1) = 1 + \frac{(-1)^{n-1}}{2^{n-1}} \sum_{i=0}^{\lfloor \frac{n}{2} \rfloor} (-1)^i \binom{n}{2i} 7^i.$$

Theorem 5 Let $G_3 : x \mapsto x^7$ over \mathbb{F}_{2^n} with $n \ge 4$. Then, its differential spectrum is given by:

• if n is odd,

$$\begin{aligned}
\omega_6 &= \frac{2^{n-2}+1}{6} - \frac{K(1)}{8} \\
\omega_4 &= 0 \\
\omega_2 &= 2^{n-1} - 3\omega_6 \\
\omega_0 &= 2^{n-1} + 2\omega_6;
\end{aligned}$$

• if n is even,

$$\begin{aligned}
\omega_6 &= \frac{2^{n-2}-4}{6} + \frac{K(1)}{8} \\
\omega_4 &= 1 \\
\omega_2 &= 2^{n-1} - 3\omega_6 - 2 \\
\omega_0 &= 2^{n-1} + 2\omega_6 + 1.
\end{aligned}$$

where K(1) is the Kloosterman sum defined as in Proposition 2. In particular, G_3 is differentially 6-uniform for all $n \ge 6$.

To prove this theorem, we need some preliminary results. We first recall some basic results on cubic equations.

Lemma 8 [3] The cubic equation $x^3 + ax + b = 0$, where $a \in \mathbb{F}_{2^n}$ and $b \in \mathbb{F}_{2^n}^*$ has a unique solution in \mathbb{F}_{2^n} if and only if $Tr(a^3/b^2) \neq Tr(1)$. In particular, if it has three distinct roots in \mathbb{F}_{2^n} , then $Tr(a^3/b^2) = Tr(1)$.

Proposition 3 [19, Appendix] Let $f_a(x) = x^3 + x + a$ and

$$M_i = \#\{ a \in \mathbb{F}_{2^n}^* \mid f_a(x) = 0 \text{ has precisely } i \text{ solutions in } \mathbb{F}_{2^n} \}.$$

Then, we have for odd n

$$M_0 = \frac{2^n + 1}{3}, \ M_1 = 2^{n-1} - 1, \ M_3 = \frac{2^{n-1} - 1}{3}$$

and for even n

$$M_0 = \frac{2^n - 1}{3}, \ M_1 = 2^{n-1}, \ M_3 = \frac{2^{n-1} - 2}{3}.$$

Now we are going to solve the equations $P_b(x) = 0$ (see Theorem 1) by solving a system of equations, including a cubic equation, thanks to the equivalence presented in Theorem 3.

Theorem 6 Let

$$P_b(x) = x^8 + bx^2 + (b+1)x, \ b \in \mathbb{F}_{2^n} \setminus \{1\}$$

The number ν_0 of $b \in \mathbb{F}_{2^n} \setminus \{1\}$ such that P_b has no roots in $\mathbb{F}_{2^n} \setminus \{0,1\}$ is given by

$$\nu_0 = \frac{2^n + (-1)^{n+1}}{3} + 2^{n-2} + (-1)^n \frac{K(1)}{4}$$

where K(1) is the Kloosterman sum defined as in Proposition 2.

Proof. Let $b \in \mathbb{F}_{2^n} \setminus \{1\}$. According to Theorem 3 we know that the number (denoted by N_b) of roots in $\mathbb{F}_{2^n} \setminus \mathbb{F}_2$ of P_b is twice the number of roots in $\mathbb{F}_{2^n}^*$ of the following system where $\beta = b + 1$:

$$\begin{cases} Q_{\beta}(y) = y^{3} + y + \beta = 0\\ Tr(y) = 0. \end{cases}$$
(5)

Since $\beta \neq 0$, $Q_{\beta}(y) \neq 0$ for $y \in \mathbb{F}_2$. Then, for any $\beta \neq 0$, the following situations may occur:

• Q_{β} has no root in \mathbb{F}_{2^n} . In this case, $N_b = 0$.

- Q_{β} has a unique root $y \in \mathbb{F}_{2^n}$. From Lemma 8, this occurs if and only if $Tr(\beta^{-1}) \neq Tr(1)$. In this case, $N_b = 0$ if Tr(y) = 1 and $N_b = 2$ if Tr(y) = 0.
- Q_{β} has three roots $y_1, y_2, y_3 \in \mathbb{F}_{2^n}$. Since these roots are roots of a linear polynomial of degree 4 then $y_3 = y_1 + y_2$, implying $Tr(y_3) = Tr(y_1) + Tr(y_2)$. Then, at least one y_i is such that $Tr(y_i) = 0$. It follows that, in this case, N_b is either 6 or 2.

Let us now define

$$B = \# \{ \beta \in \mathbb{F}_{2^n}^*, \ Q_\beta \text{ has a unique root } y \in \mathbb{F}_{2^n} \text{ and } Tr(y) = 1 \}.$$

From the previous discussion, we have

$$\nu_0 = \#\{\beta \in \mathbb{F}_{2^n}^*, \ Q_\beta \text{ has no root in } \mathbb{F}_{2^n}\} + B$$
$$= \frac{2^n + (-1)^{n+1}}{3} + B$$

where the last equality comes from Proposition 3. Let us now compute the value of B.

$$B = \#\{\beta \in \mathbb{F}_{2^{n}}^{*}, \ Q_{\beta} \text{ has a unique root } y \in \mathbb{F}_{2^{n}} \text{ and } Tr(y) = 1\}$$
$$= \#\{(y^{3} + y) \in \mathbb{F}_{2^{n}}^{*}, \ Tr\left(\frac{1}{y^{3} + y}\right) \neq Tr(1) \text{ and } Tr(y) = 1\},$$

by using that $\beta = y^3 + y$. But, we have

$$\frac{1}{y^3 + y} = \frac{1 + y^2}{y^3 + y} + \frac{y^2 + y}{y^3 + y} + \frac{y}{y^3 + y} = \frac{1}{y} + \frac{1}{y + 1} + \frac{1}{y^2 + 1},$$

implying that

$$Tr\left(\frac{1}{y^3+y}\right) = Tr\left(\frac{1}{y}\right).$$

Therefore,

$$B = \#\{(y^3 + y) \in \mathbb{F}_{2^n}^*, \ Tr\left(\frac{1}{y}\right) \neq Tr(1) \text{ and } Tr(y) = 1\}.$$

Now, we clearly have that $(y^3 + y) = 0$ if and only if $y \in \mathbb{F}_2$. Moreover, two distinct elements y_1 and y_2 in $\mathbb{F}_{2^n} \setminus \mathbb{F}_2$ with $Tr(y_1^{-1}) \neq Tr(1)$ and

 $Tr(y_2^{-1}) \neq Tr(1)$ satisfy $(y_1^3+y_1) \neq (y_2^3+y_2)$ (otherwise, Q_β with $\beta = y_1^3+y_1$ has at least 2 roots in \mathbb{F}_{2^n}). Therefore, we deduce that

$$B = \#\{y \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2, \ Tr\left(\frac{1}{y}\right) \neq Tr(1) \text{ and } Tr(y) = 1\}.$$

If n is odd, we deduce that

$$B = \#\{y \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2, \ Tr\left(\frac{1}{y}\right) = 0 \text{ and } Tr(y) = 1\}.$$

If n is even, we deduce that

$$B = \#\{y \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2, \ Tr\left(\frac{1}{y}\right) = 1 \text{ and } Tr(y) = 1\}$$

= $\#\{y \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2, \ Tr(y) = 1\}$
 $-\#\{y \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2, \ Tr\left(\frac{1}{y}\right) = 0 \text{ and } Tr(y) = 1\}$
= $2^{n-1} - \#\{y \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2, \ Tr\left(\frac{1}{y}\right) = 0 \text{ and } Tr(y) = 1\}.$

On the other hand, by definition of the Kloosterman sum K(1), we have

$$K(1) - 2 = \sum_{x \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2} (-1)^{Tr(x^{-1} + x)}$$

= $-2\#\{x \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2, Tr(x^{-1} + x) = 1\} + 2^n - 2$
= $-4\#\{x \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2, Tr(x^{-1}) = 0 \text{ and } Tr(x) = 1\} + 2^n - 2.$

Thus,

$$#\{x \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2, Tr(x^{-1}) = 0 \text{ and } Tr(x) = 1\} = 2^{n-2} - \frac{K(1)}{4}.$$

We then deduce that, for any n,

$$B = 2^{n-2} + (-1)^n \frac{K(1)}{4}.$$

It follows that

$$\nu_0 = \frac{2^n + (-1)^{n+1}}{3} + 2^{n-2} + (-1)^n \frac{K(1)}{4}.$$

\diamond	
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Proof. (Proof of Theorem 5) In accordance with Lemma 1, we obtain the differential spectrum of G_3 as soon as we are able to solve the following system:

$$\frac{\omega_0 + \omega_2 + \omega_4 + \omega_6 = 2^n}{2\omega_2 + 4\omega_4 + 6\omega_6 = 2^n}$$
(6)

Now, we apply Theorem 1 and we recall first that $\delta(b) \in \{0, 2, 6\}$ for any $b \in \mathbb{F}_{2^n} \setminus \{1\}$. Moreover, we know that $\omega_0 = \nu_0$ as defined in Theorem 6.

Since t = 3, gcd(t - 1, n) equals 1 for odd n and 2 otherwise. Then, if n is even then $\delta(1) = 4$ else $\delta(1) = 2$. Thus, $\omega_4 = 1$ for even n and $\omega_4 = 0$ otherwise. From the second equation of (6), we get

$$\omega_2 = 2^{n-1} - 3\omega_6 - 2\omega_4$$

and using the first equation of (6)

$$\omega_6 = 2^n - \omega_0 - \omega_2 - \omega_4 = 2^{n-1} - \omega_0 + \omega_4 + 3\omega_6 + \frac{1}{2}\omega_6 + \frac{1}{2}\omega_$$

leading to

$$\omega_6 = -2^{n-2} + \frac{\omega_0 - \omega_4}{2}$$

Finally, we deduce from Theorem 6 that, for odd n,

$$\omega_6 = -2^{n-2} + \frac{\omega_0}{2} = -2^{n-3} + \frac{2^n + 1}{6} - \frac{K(1)}{8}$$
$$= \frac{2^{n-2} + 1}{6} - \frac{K(1)}{8}$$

and for even n

$$\omega_6 = -2^{n-2} + \frac{\omega_0 - 1}{2} = -2^{n-3} + \frac{2^n - 1}{6} + \frac{K(1)}{8} - \frac{1}{2}$$
$$= \frac{2^{n-2} - 4}{6} + \frac{K(1)}{8}.$$

Finally, it can be proved that $\omega_6 \ge 1$ for any $n \ge 6$, implying that G_3 is differentially 6-uniform. Actually, it has been proved in [20, Th. 3.4] that

$$-2^{\frac{n}{2}+1} + 1 \le K(1) \le 2^{\frac{n}{2}+1} + 1$$

implying that $\omega_6 > 0$ when n > 5. It is worth noticing that G_3 is APN when n = 5 since its inverse is the quadratic APN permutation $x \mapsto x^9$. When n = 4, G_3 is locally-APN, and not APN, since it corresponds to the inverse function over \mathbb{F}_{2^4} .

By combining the previous theorem and Corollary 2, we deduce the differential spectrum of $G_{n-2}: x \mapsto x^{2^{n-2}-1}$ over \mathbb{F}_{2^n} .

Corollary 5 Let $G_{n-2}: x \mapsto x^{2^{n-2}-1}$ over \mathbb{F}_{2^n} with $n \ge 6$. Then, we have:

 if gcd(n,3) = 1, G_{n-2} is differentially 6-uniform and for any b ∈ 𝔽_{2ⁿ}, δ(b) ∈ {0,2,6}. Moreover, its differential spectrum is given by:

$$\omega_{6} = \begin{cases} \frac{2^{n-2}+1}{6} - \frac{K(1)}{8} & \text{for odd } n \\ \frac{2^{n-2}-4}{6} + \frac{K(1)}{8} & \text{for even } n \end{cases}$$

$$\omega_{2} = 2^{n-1} - 3\omega_{6}$$

$$\omega_{0} = 2^{n-1} + 2\omega_{6};$$

 if 3 divides n, G_{n-2} is differentially 8-uniform and for any b ∈ 𝔽_{2ⁿ}, δ(b) ∈ {0,2,6,8}. Moreover, its differential spectrum is given by:

$$\begin{split} \omega_8 &= 1\\ \omega_6 &= \begin{cases} \frac{2^{n-2}-5}{6} - \frac{K(1)}{8} & \text{for odd } n\\ \frac{2^{n-2}-10}{6} + \frac{K(1)}{8} & \text{for even } n \end{cases}\\ \omega_2 &= 2^{n-1} - 3\omega_6 - 4\\ \omega_0 &= 2^{n-1} + 2\omega_6 + 3 ; \end{split}$$

Proof. Let $(\omega'_0, \omega'_2, \omega'_4, \omega'_6)$ denote the differential spectrum of G_3 over \mathbb{F}_{2^n} . We apply Corollary 2 (with s = 3). Then, if gcd(3, n) = 1, $\delta_3(0) = 0$ and $\delta_{n-2}(1) = 2$. Otherwise, $\delta_3(0) = 6$ and $\delta_{n-2}(1) = 8$. Moreover, in both cases, $\delta_3(1) = 4$ for n even and $\delta_3(1) = 2$ for n odd. It follows that,

- for gcd(3, n) = 1, *n* odd, we have $(\delta_3(0), \delta_3(1)) = (0, 2)$ and $(\delta_{n-2}(0), \delta_{n-2}(1)) = (0, 2)$. Then, $\omega_i = \omega'_i$ for all *i*;
- for gcd(3,n) = 1, *n* even, we have $(\delta_3(0), \delta_3(1)) = (0,4)$ and $(\delta_{n-2}(0), \delta_{n-2}(1)) = (2,2)$. Then, $\omega_0 = \omega'_0 1$, $\omega_4 = \omega'_4 1$ and $\omega_2 = \omega'_2 + 2$.
- for gcd(3,n) = 3, n odd, we have $(\delta_3(0), \delta_3(1)) = (6,2)$ and $(\delta_{n-2}(0), \delta_{n-2}(1)) = (0,8)$. Then, $\omega_8 = 1$, $\omega_6 = \omega'_6 1$, $\omega_2 = \omega'_2 1$ and $\omega_0 = \omega'_0 + 1$.
- for gcd(3,n) = 3, n even, we have $(\delta_3(0), \delta_3(1)) = (6,4)$ and $(\delta_{n-2}(0), \delta_{n-2}(1)) = (2,8)$. Then, $\omega_8 = 1$, $\omega_6 = \omega'_6 1$, $\omega_4 = \omega'_4 1$, $\omega_2 = \omega'_2 + 1$ and $\omega_0 = \omega'_0$.

The result finally follows from Theorem 5.

The minimum distance of the cyclic code of length $2^n - 1$ with defining set $\{1,7\}$ has been studied by van Lint and Wilson in [23]. More precisely, they have proved that this code has minimum distance at most 4 for $n \ge 6$. The previous corollary recovers this result and also provides the exact number of codewords of weight 3 and 4 in this code.

Corollary 6 Let B_3 (resp. B_4) denote the number of codewords of Hamming weight 3 (resp. of Hamming weight 4) in the binary cyclic code of length $2^n - 1$ with defining set {1,7}. Then, we have

• if n is odd

$$B_3 = 0$$

$$B_4 = (2^n - 1) \left(\frac{2^{n-2} + 1}{6} - \frac{K(1)}{8} \right);$$

• if n is even

$$B_3 = \frac{(2^n - 1)}{3}$$

$$B_4 = (2^n - 1)\left(\frac{2^{n-2} - 4}{6} + \frac{K(1)}{8}\right)$$

Proof. Let $F(x) = x^d$ over \mathbb{F}_{2^n} and let $\delta(b), b \in \mathbb{F}_{2^n}$, denote the number of solutions x of

$$D_1F(x) = F(x+1) + F(x) = b$$
.

It is known from Proposition 2 and Lemma 2 in [5] that the number of codewords of weight 3 and 4 in the cyclic code of length $(2^n - 1)$ with defining set $\{1, d\}$ is given by

$$B_3 = \frac{(2^n - 1)}{6} (\delta(1) - 2)$$

$$B_3 + B_4 = \frac{(2^n - 1)}{24} \left[\#\{(x, y) \in \mathbb{F}_{2^n} \times \mathbb{F}_{2^n} : D_1 F(x) = D_1 F(y) \} - 2^{n+1} \right]$$

Therefore, we have

$$B_3 + B_4 = \frac{(2^n - 1)}{24} \left(\sum_{b \in \mathbb{F}_{2^n}} \delta(b)^2 - 2^{n+1} \right)$$
$$= \frac{(2^n - 1)}{24} \left[\sum_{i=0}^{2^n} i^2 \omega_i - 2^{n+1} \right].$$

This formula was proved in Corollary 1 of [5], but only in the particular case where $gcd(d, 2^n - 1) = 1$. For d = 7, Theorem 5 implies that

$$B_3 + B_4 = (2^n - 1)\omega_6$$

Then, the values of B_3 and B_4 are deduced from the expression of ω_6 given in Theorem 5.

5.2 Exponents $2^{\lfloor n/2 \rfloor} - 1$

We are going to determine the differential uniformity of G_t for $t = \lfloor n/2 \rfloor$. We first consider the case where *n* is even. Note that in this case, G_t is not a permutation since $2^n - 1 = (2^t - 1)(2^t + 1)$.

Theorem 7 Let n be an even integer, n > 4 and $G_t(x) = x^{2^t-1}$ for $t = \frac{n}{2}$. Then G_t is locally-APN. More precisely

$$\delta(G_t) = 2^t - 2 \text{ and } \delta(b) \leq 2, \forall b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2.$$

Moreover, the differential spectrum of G_t is:

• if $n \equiv 0 \mod 4$ then

$$\begin{aligned}
\omega_{2^{t}-2} &= 1 \\
\omega_{i} &= 0, \ \forall \ i, \ 2 < i < 2^{t}-2 \\
\omega_{2} &= 2^{n-1} - 2^{t-1} + 1 \\
\omega_{0} &= 2^{n-1} + 2^{t-1} - 2;
\end{aligned}$$

• if $n \equiv 2 \mod 4$,

$$\begin{aligned} \omega_{2^{t}-2} &= 1\\ \omega_{i} &= 0, \ \forall \ i, \ 4 < i < 2^{t}-2\\ \omega_{4} &= 1\\ \omega_{2} &= 2^{n-1} - 2^{t-1} - 1\\ \omega_{0} &= 2^{n-1} + 2^{t-1} - 1. \end{aligned}$$

Proof. From Theorem 1, we obtain directly $\delta(0) = 2^t - 2$. Also, $\delta(1) = 2$ if t is even and $\delta(1) = 4$ otherwise.

Now, for all $b \notin \mathbb{F}_2$, we have to determine the number of roots in \mathbb{F}_{2^n} of $P_b(x) = x^{2^t} + bx^2 + (b+1)x$ or, equivalently, the number of roots of

$$(P_b(x))^{2^t} = x + b^{2^t} x^{2^{t+1}} + (b+1)^{2^t} x^{2^t}.$$
(7)

If x is a root of P_b then $x^{2^t} = bx^2 + (b+1)x$. So, $P_b(x) = 0$ implies

$$(P_b(x))^{2^t} = x + b^{2^t} (x^{2^t})^2 + (b^{2^t} + 1)x^{2^t}$$

= $x + b^{2^t} (bx^2 + (b+1)x)^2 + (b^{2^t} + 1)(bx^2 + (b+1)x)$
= $b^{2^t+2}x^4 + (b^{2^t+2} + b^{2^t+1} + b^{2^t} + b)x^2 + (b^{2^t+1} + b^{2^t} + b)x$
= $b^{2^t+2} (x^2 + x)^2 + (b^{2^t+1} + b^{2^t} + b)(x^2 + x).$

Thus, we get a linear polynomial of degree 4 which has at least the roots 0 and 1. Hence, this polynomial has τ roots where τ is either 4 or 2, including x = 0 and x = 1. Therefore, for any $b \notin \mathbb{F}_2$, $\delta(b) \leq 2$ since $\delta(b) \leq \tau - 2$. We deduce that G_t is locally-APN.

We also proved that $\omega_i = 0$ unless $i \in \{0, 2, 2^t - 2\}$ when t is even and $i \in \{0, 2, 4, 2^t - 2\}$ otherwise. Moreover $\omega_{2^t-2} = \omega_4 = 1$. According to Lemma 1, we have for t even :

$$2^{n} = \omega_{0} + \omega_{2} + \omega_{2^{t}-2} = \omega_{0} + \omega_{2} + 1$$

and

$$2^{n} = 2\omega_{2} + (2^{t} - 2)\omega_{2^{t} - 2} = 2\omega_{2} + (2^{t} - 2)$$

So, we get $\omega_2 = 2^{n-1} - 2^{t-1} + 1$ and conclude with $\omega_0 = 2^n - \omega_2 - 1$. We proceed similarly for odd t, with the following equalities derived from Lemma 1:

$$2^n = \omega_0 + \omega_2 + 2$$
 and $2^n = 2\omega_2 + 2^t + 2$.

 \diamond

And we directly deduce a property on the corresponding class of linear polynomials.

Corollary 7 Let n = 2t and let Tr_t denote the absolute trace on \mathbb{F}_{2^t} . Consider the polynomials over \mathbb{F}_{2^n} :

$$x^{2^{t}} + bx^{2} + (b+1)x$$
 and $x^{2^{t+1}} + bx^{2} + (b+1)x$.

Then, for any $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2$, these polynomials have either 2 or 4 roots in \mathbb{F}_{2^n} . The first one has 4 roots if and only if $Tr_t(b^{-(2^t+1)}) = 1$ with $(1+b) \notin \mathcal{G}$, where \mathcal{G} is the cyclic subgroup of \mathbb{F}_{2^n} of order $2^t + 1$.

Proof. Let $P_b(x) = x^{2^t} + bx^2 + (b+1)x$. We define

$$Q_b(x) = \frac{(P_b(x))^{2^t} + b^{2^t}(P_b(x))^2 + (b^{2^t} + 1)P_b(x)}{b^{2^t+2}(x^2 + x)}$$

Using (7), we get :

$$Q_b(x) = x^2 + x + A$$
, with $A = \frac{b^{2^t+1} + b^{2^t} + b}{b^{2^t+2}}$.

To be clear, we summarize the situation:

- if $P_b(x) = 0$, $x \notin \{0, 1\}$, then $Q_b(x) = 0$; - when $Q_b(x) = 0$, $x \notin \{0, 1\}$, one can have $P_b(x) \neq 0$; - if $Q_b(x) = 0$ for $x \in \{0, 1\}$ only, this holds for $P_b(x)$ too.

We consider the case where $P_b(x) = 0$ has more than the two solutions 0 and 1. The equation $Q_b(x) = 0$ has two solutions (not in $\{0, 1\}$) if and only if Tr(A) = 0 with $A \neq 0$. But

$$Tr(A) = Tr\left(\frac{1}{b} + \frac{1}{b^2} + \frac{1}{b^{2^t+1}}\right) = Tr\left(\frac{1}{b^{2^t+1}}\right) = 0,$$

for all b, since $b^{2^t+1} \in \mathbb{F}_{2^t}$. And,

$$A \neq 0 \iff b^{2^t+1} + b^{2^t} + b \neq 0 \iff (b+1)^{2^t+1} \neq 1,$$

that is : b+1 is not in the cyclic subgroup \mathcal{G} of order 2^t+1 of $\mathbb{F}_{2^n}^*$. On the other hand, if $Q_b(x) = 0$ then $x^2 + x = A$ and we get

$$P_b(x) = x^{2^t} + x + b(x^2 + x)$$

= $(x^2 + x)^{2^{t-1}} + (x^2 + x)^{2^{t-2}} + \dots + (x^2 + x) + b(x^2 + x)$
= $A^{2^{t-1}} + \dots + A + bA.$

We compute this last expression by replacing the value of A:

$$P_b(x) = \sum_{i=0}^{t-1} \left(\frac{1}{b} + \frac{1}{b^2}\right)^{2^i} + \sum_{i=0}^{t-1} \left(\frac{1}{b^{2^t+1}}\right)^{2^i} + 1 + \frac{1}{b} + \frac{1}{b^{2^t}}$$
$$= Tr_t \left(\frac{1}{b^{2^t+1}}\right) + 1,$$

where Tr_t is the absolute trace on \mathbb{F}_{2^t} . We conclude that $P_b(x) = 0$ if and only if $Tr_t(b^{-(2^t+1)}) = 1$, with $b+1 \notin \mathcal{G}$.

According to Corollary 2, the differential spectrum of $x \mapsto x^{2^{\frac{n}{2}}-1}$ determines the differential spectrum of $x \mapsto x^{2^{\frac{n}{2}+1}-1}$

Theorem 8 Let n be an even integer n > 4 and $G_{t+1}(x) = x^{2^{t+1}-1}$ for $t = \frac{n}{2}$. Then, G_{t+1} is locally-APN. It is differentially 2^t -uniform and its differential spectrum is

$$\begin{array}{rcl} \omega_{2^t} &=& 1 \\ \omega_i &=& 0, \; \forall \; i, \; 2 < i < 2^t \\ \omega_2 &=& 2^{n-1} - 2^{t-1} \\ \omega_0 &=& 2^{n-1} + 2^{t-1} - 1 \; . \end{array}$$

Moreover, G_{t+1} is a permutation if and only if $n \equiv 0 \mod 4$.

Proof. First, since n = 2t, we have gcd(t + 1, n) = 1 if t is even (*i.e.*, $n \equiv 0 \mod 4$) and gcd(t + 1, n) = 2 if t is odd (*i.e.*, $n \equiv 2 \mod 4$). Here s = t + 1.

Let $(\omega'_i)_{0 \leq i \leq 2^n}$ (resp. $(\omega_i)_{0 \leq i \leq 2^n}$) denote the differential spectrum of G_t (resp. G_{t+1}) over \mathbb{F}_{2^n} .

- For $n \equiv 0 \mod 4$, we have $(\delta_t(0), \delta_t(1)) = (2^t 2, 2)$ and $(\delta_s(0), \delta_s(1)) = (0, 2^t)$. Thus, $\omega_0 = \omega'_0 + 1$, $\omega_2 = \omega'_2 1$, $\omega_{2^t-2} = \omega'_{2^t-2} 1$ and $\omega_{2^t} = 1$.
- For $n \equiv 2 \mod 4$, we have $(\delta_t(0), \delta_t(1)) = (2^t 2, 4)$ and $(\delta_s(0), \delta_s(1)) = (2, 2^t)$. Thus, $\omega_2 = \omega'_2 + 1$, $\omega_4 = \omega'_4 1$, $\omega_{2^t-2} = \omega'_{2^t-2} 1$ and $\omega_{2^t} = 1$.

The differential spectrum of G_{t+1} is then directly deduced by combining the previous formulas with the values of ω'_i computed in Theorem 7.

In the case where n is odd, the differential uniformity of G_t , with $t = \frac{n-1}{2}$, can also be determined.

Theorem 9 Let n be an odd integer, n > 3. Let $G_t(x) = x^{2^t-1}$ with t = (n-1)/2. Then, G_t is a permutation and for all $b \in \mathbb{F}_{2^n} \setminus \mathbb{F}_2$ we have $\delta(b) \in \{0,2,6\}$. Moreover

- if $n \equiv 0 \mod 3$, then $\delta(G_t) = 8$, and the differential spectrum satisfies $\omega_i = 0$ for all $i \notin \{0, 2, 6, 8\}$ and $\omega_8 = 1$.
- if $n \not\equiv 0 \mod 3$, then $\delta(G_t) \leq 6$ and the differential spectrum satisfies $\omega_i = 0$ for all $i \notin \{0, 2, 6\}$.

Proof. From Theorem 1, we have $\delta(0) = 0$; moreover, if 3 divides *n* then $\delta(1) = 8$ else $\delta(1) = 2$. Now, for all $b \notin \mathbb{F}_2$, we have to determine the number of roots in \mathbb{F}_{2^n} of

$$P_b(x) = x^{2^{t}} + bx^2 + (b+1)x$$

or, equivalently, the number of roots of

$$(P_b(x))^{2^{t+1}} = x + b^{2^{t+1}}x^{2^{t+2}} + (b+1)^{2^{t+1}}x^{2^{t+1}}$$

Set $c = b^{2^{t+1}}$ and $Q_b(x) = (P_b(x))^{2^{t+1}}$. If x is a root of P_b then $x^{2^t} =$ $bx^{2} + (b+1)x$. So, $P_{b}(x) = 0$ implies

$$Q_b(x) = x + c(x^{2^t})^4 + (c+1)(x^{2^t})^2$$

= $x + c(bx^2 + (b+1)x)^4 + (c+1)(bx^2 + (b+1)x)^2$
= $cb^4x^8 + (c(b+1)^4 + (c+1)b^2)x^4 + (c+1)(b^2+1)x^2 + x$.

Since Q_b has degree 8, it has either 8 or 4 or 2 solutions. In other terms, $\delta(b) \in \{0, 2, 6\}.$ \diamond

6 Conclusions

In this work, we point out that the family of all power functions

$$\{ G_t : x \mapsto x^{2^{\iota} - 1} \text{ over } \mathbb{F}_{2^n}, \ 1 < t < n \}$$

$$\tag{8}$$

has interesting differential properties. The study of these properties led us to introduce locally-APN functions, as a generalization of the differential spectrum of the inverse function.

In particular, we give several results about the functions with a low differential uniformity within family (8). There are classes of functions G_t such that $\delta(G_t) = 6$. It is the case for the functions G_3 over \mathbb{F}_{2^n} (see Theorem 5).

The functions such that $\delta(G_t) \leq 4$ can be differentially 4-uniform for even n only (see Corollary 4). We have shown that, for exponents of the form $2^t - 1$, the APN property imposes many conditions of the value of t.In particular, it is easy to prove, using Theorem 1 that such exponent must satisfy gcd(t,n) = 2 for even n and gcd(t,n) = gcd(t-1,n) = 1 for odd n. Another condition can be derived from the recent result by Aubry and Rodier [1] who proved the following theorem.

Theorem 10 [1, Theorem 9] Let $G_t : x \mapsto x^{2^t-1}$ over \mathbb{F}_{2^n} with $t \geq 3$. If $7 \leq 2^t - 1 < 2^{n/4} + 4.6$ then $\delta(G_t) > 4$.

Thanks to Corollary 2, we can extend this result as follows.

Corollary 8 Let $G_t : x \mapsto x^{2^{t-1}}$ over \mathbb{F}_{2^n} with $3 \leq t \leq n-2$. If $\delta(G_t) \leq 4$, then $(\frac{n}{4} + 5.6)$. 10

$$\log_2(2^{\overline{4}} + 5.6) \le t \le n + 1 - \log_2(2^{\overline{4}} + 5.6)$$

Proof. Let s = n - t + 1 so that $3 \le s \le n - 2$. In this proof, we denote by $\delta_t(b)$ (resp. $\delta_s(b)$) the quantities $\delta(b)$ corresponding to G_t (resp. G_s).

From Theorem 10, we know that $\delta(G_t) \leq 4$ implies

 $2^{n/4} + 4.6 \le 2^t - 1$, *i.e.*, $t \ge \log_2(2^{\frac{n}{4}} + 5.6)$.

We consider now the function G_s . Note that, from Theorem 1, $\delta(G_t) \leq 4$ implies $\delta_t(0) \in \{0, 2\}$ and $\delta_t(1) \in \{2, 4\}$. Moreover, we obtain directly from Corollary 2 :

- $\delta_s(b) \leq 4$, for any $b \notin \mathbb{F}_2$.
- $\delta_s(0) \in \{0, 2\}$ and $\delta_s(1) \in \{2, 4\}$.

Thus $\delta(G_s) \leq 4$ and, applying Theorem 10 again, we get

$$s \ge \log_2(2^{\frac{n}{4}} + 5.6), \quad i.e, \ n+1 - \log_2(2^{\frac{n}{4}} + 5.6) \ge t.$$

 \diamond

We now concentrate on APN functions belonging to the family (8). Some are well-known as the inverse permutation for n odd (t = n - 1) and the quadratic function $x \mapsto x^3$ (t = 2). There is also the function G_t for t = (n + 1)/2 with n odd, because this function is the inverse of the quadratic function $x \mapsto x^{2^{(n+1)/2}+1}$. Recall that x^{2^i+1} is an APN function over \mathbb{F}_{2^n} if and only if gcd(n, i) = 1 and we have obviously gcd(n, (n + 1)/2) = 1(for odd n). We conjecture that these three functions are the only APN functions within family (8).

Conjecture 1 Let $G_t(x) = x^{2^t-1}$, $2 \le t \le n-1$. If G_t is APN then either t = 2 or n is odd and $t \in \{\frac{n+1}{2}, n-1\}$.

If the previous conjecture holds then there are some consequences for the functions of (8) which are differentially 4-uniform. From Corollary 4, we can say that such a function G_t is a function over \mathbb{F}_{2^n} with n even. Moreover G_s , s = n - t + 1, is APN. If the conjecture holds then s = 2 (t = n - 1) is the only one possibility. So, in this case we could conclude that the inverse function is the only one differentially 4-uniform function of family (8).

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