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Abstract

The minimum distance computation problem between two surfaces is very important in many applications such as robotics, CAD/CAM and computer graphics. Given two implicit algebraic surfaces, a new method based on the offset technique is presented to compute the minimum distance and a pair of points where the minimum distance occurs. The new method also works where there are an implicit algebraic surface and a parametric surface. Quadric surfaces, tori and canal surfaces are used to demonstrate our new method. When the two surfaces are a general quadric surface and a surface which is a cylinder, a cone or an elliptic paraboloid, the new method can produce two bivariate equations where the degrees are lower than those of any existing method.

Keywords: Minimum distance; Offset; Canal surface; Implicit algebraic surface; Parametric surface

1. Introduction

Computation of the minimum distance between two objects is very important in many applications such as robotics [15,16,20,25,34,42], CAD/CAM [9,17,32,33,45], and computer graphics [27,39]. In robot motion planning, distance information is needed for computing interaction force and penalty function [20]. Distance information is useful for interference avoidance in CAD/CAM and NC verification [17]. It is also very important for collision detection and dynamic simulation in computer graphics [15,16,27]. Many papers [12, 13,16,24,36] are devoted to the computation of the minimum distance between two polyhedra. When the objects concerned are curved surfaces, they are converted into polyhedra to compute the minimum distance. However, the number of polyhedral faces approximating the given curved surfaces is usually very large when high precision is desired. The space cost would become a considerable problem, and the methods based on polyhedra may fail for the real time requirement.

Recently, some papers [11,17,22,39,40,45] have considered methods that work directly with curved surfaces. When the two surfaces intersect, the minimum distance between them is zero. There are many algorithms [2,10,14,18,19,21,23,28–30, 35,41,44] designed to test whether two given surfaces intersect. There are three dominant techniques used in the intersection test: the subdivision method, the lattice method, and the analytic method. For some simple cases, the analytic methods may yield good results [11,45]. In [11], Fu gives a necessary and sufficient condition to judge whether a bicubic Bézier patch and a plane intersect or not. In [45], Wang and Kim show that two ellipsoids are separated if and only if their degree 4 characteristic polynomial has two distinct positive roots. In practice, these techniques are often combined. Moreover, other techniques are also used in the intersection test. For instance, Tanaka et al. [40] use Monte Carlo simulation based on a stochastic differential equation to sample implicit surfaces for intersection test. Some other papers on collision detection can be found in [5,26].

The computation of the minimum distance between two curved surfaces usually involves a system of non-linear equations. There are many references for solving non-linear equation systems. Some of the common techniques used to find local solutions are the conjugate direction method, the steepest descent method, the Newton–Raphson method, and
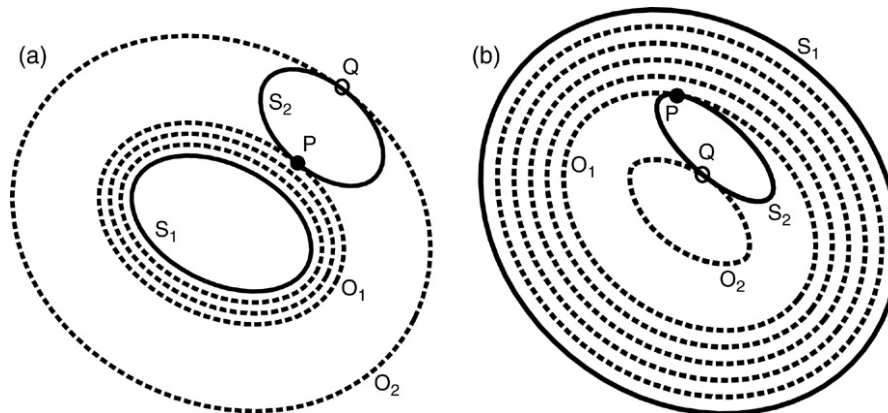
Fig. 1. Offset surfaces and touching point (a) growing or (b) shrinking.

the gradient flow method [4,8,38,43]. To compute all roots in a given area of interest, several global solution methods such as algebraic and hybrid methods, homotopy methods and subdivision methods [31] are useful. The resultant based methods [1,3,6,7,46,47] seem to be very promising in providing solutions for low degree algebraic equations.

In this paper, we only consider the case where the two surfaces do not intersect. There are a number of efficient ways for computing the minimum distance between two surfaces such as planes, quadrics, tori, or canal surfaces [17,22,39]. Lennerz and Schömer [22] discuss the minimum distance computation problem between two quadric surfaces. They use the Lagrange technique to obtain two bivariate polynomial equations in Lagrange multipliers. When the two surfaces are a cone and a general quadric surface, [22] obtains two bivariate polynomial equations both of degree 4 or a univariate equation of degree 12. Sohn et al. [39] use line geometry to convert the minimum distance problem into an intersection problem between the normal congruences of the two surfaces. Each point on a surface is associated with the line passing through the point and spanned by the normal of the surface at that point. The normal congruence of a surface is defined as a set of lines associated with the points on the surface. They describe one of the two normal congruences in parametric form and the other normal congruence in implicit form, and obtain two bivariate equations. Usually one of the two surfaces in [39] is in implicit form and the other is in parametric form. To describe a normal congruence for a general implicit algebraic surface, one needs to eliminate seven unknown variables from 10 polynomial equations (or eliminate four variables from seven polynomial equations by simplification), which is a non-trivial task. When the two surfaces are a cone and a general quadric surface, [39] produces two bivariate polynomial equations of degrees 4 and 8. In [17], Kim computes the minimum distance between a canal surface and a simple surface such as a plane, a sphere, a cylinder, a cone, or a torus. The fact that the normals at the closest points between the two surfaces are collinear is exploited, and characteristic circles are introduced to convert the minimum distance problem into a problem of finding the roots of a univariate function.

The offset technique employed in this paper generalizes some of the methods mentioned above. Given two surfaces,

when we offset one of the closest points in one of the two surfaces along or against its normal direction using the minimum distance, we get a point on the other surface, and the distance between these two points is exactly the minimal distance. The method presented in this paper is based on this simple but important fact. Suppose two surfaces S_1 and S_2 do not intersect. We build the dynamic offset surface of one of the two surfaces, for example S_1. As Fig. 1 shows, if S_1 contains S_2, the dynamic offset surface of S_1 is shrinking, i.e., offsetting inward, towards S_2; otherwise, the dynamic offset surface of S_1 is growing, i.e., offsetting outward, towards S_2. Eventually, a dynamic offset surface O_1 will "touch" S_2. The touching point P could be one of the two points where the minimum distance occurs. Together with the fact that the normals of the two given surfaces at the two closest points where the minimum distance occurs are collinear, we obtain the equations for calculating one of the two closest points. From these equations we may get several candidate points, e.g., in Fig. 1(b), both P and Q are candidate points which satisfy the equations. To determine the point we want, we define the corresponding distance of a candidate point. If the normal line at a candidate point intersects with S_1, the corresponding distance of the candidate point is defined as the distance from the point to the closest intersection point of the normal line and S_1; if not, it is defined as a positive infinite real number. The point we want is the candidate point whose corresponding distance is minimal. Once this point is identified, the other point in S_2 can be calculated along or against the unit normal vector of the first point with the minimum distance.

The remaining part of this paper is organized as follows: an algorithm based on the offset method to compute the minimum distance between two surfaces is given in Section 2; in Section 3, we illustrate our method for two implicit algebraic surfaces, using quadric surfaces and tori; Section 4 illustrates the new method for an implicit algebraic surface and a canal surface; whereas Section 5 compares our method with some known existing methods. Some concluding remarks are given in the last section.

2. Minimum distance between two surfaces

Given two surfaces S_1 and S_2, the minimum distance computation problem is to find a non-negative real number $d_{\min}$

such that

$$d_{\min} = \inf\{\|\mathbf{p} - \mathbf{q}\| \mid \mathbf{p} \in \mathbf{S}_1 \text{ and } \mathbf{q} \in \mathbf{S}_2\},$$

where $\|\mathbf{p} - \mathbf{q}\|$ is the Euclidean distance between the points $\mathbf{p}$ and $\mathbf{q}$. $d_{\min}$ is the minimum distance between the surfaces $\mathbf{S}_1$ and $\mathbf{S}_2$. Any pair of points $(\mathbf{p}, \mathbf{q})$ where the minimum distance occurs is called a "location of the minimum distance". If $\mathbf{S}_1$ and $\mathbf{S}_2$ intersect, $d_{\min}$ is equal to zero, and any intersection point is a location. Here, we assume that $\mathbf{S}_1$ and $\mathbf{S}_2$ do not intersect.

In this paper, we assume that at least one of the two surfaces is an implicit algebraic surface. In this case, without loss of generality, let $\mathbf{S}_2$ be an implicit algebraic surface. The other surface $\mathbf{S}_1$ can be either an implicit algebraic surface or a parametric surface. Suppose that $\mathbf{S}_2$ is given by

$$f_2(\mathbf{q}) = 0,$$

where $\mathbf{q}$ is a point in $\mathbb{R}^3$. A point $\mathbf{q}$ is on the surface $\mathbf{S}_2$ if and only if $f_2(\mathbf{q}) = 0$. We write the coordinates of $\mathbf{q}$ as a column vector $(x, y, z)^{\mathrm{T}}$. The gradient at point $\mathbf{q} \in \mathbf{S}_2$ is

$$(f_{2x}(\mathbf{q}), f_{2y}(\mathbf{q}), f_{2z}(\mathbf{q}))^{\mathrm{T}}, \tag{1}$$

where $f_{2x}(\mathbf{q}), f_{2y}(\mathbf{q}), f_{2z}(\mathbf{q})$ are the first partial derivatives of $f_2(x, y, z)$ with respect to x, y and z.

For any real number d, we define the dynamic d-offset surface of the surface $\mathbf{S}_1$ to be

$$\mathbf{S}_1^d = \left\{ \begin{pmatrix} x + dn_x \\ y + dn_y \\ z + dn_z \end{pmatrix} \middle| \begin{pmatrix} n_x \\ n_y \\ n_z \end{pmatrix} \text{ is a unit normal vector of } \mathbf{S}_1 \text{ at } \begin{pmatrix} x \\ y \\ z \end{pmatrix} \in \mathbf{S}_1 \right\},$$

which depends completely on the surface $\mathbf{S}_1$. In fact, the new method only uses the information of $\mathbf{S}_1$ and d, and is not affected by the singularity of the d-offset surface $\mathbf{S}_1^d$.

Suppose that the minimum distance between $\mathbf{S}_1$ and $\mathbf{S}_2$ is a positive real number $d_{\min}$, and point $\mathbf{p}$ is one point of a location of $d_{\min}$ on the surface $\mathbf{S}_1$. Let $\mathbf{N}_p$ be the gradient of $\mathbf{S}_1$ at $\mathbf{p}$. If $\mathbf{p}$ is regular, i.e., $\mathbf{N}_p$ is not degenerated into a zero vector, there exists $d = d_{\min}$ or $d = -d_{\min}$, such that a dynamic d-offset surface $\mathbf{S}_1^d$ is tangent to the surface $\mathbf{S}_2$ at point

$$\mathbf{q} = \mathbf{p} + \mu \mathbf{N}_p, \tag{2}$$

where $\mu = \frac{d}{\|\mathbf{N}_p\|}$. The relationship between the minimum distance and the dynamic offset surface is given by

$$d_{\min} = \inf\{d \geq 0 \mid \mathbf{S}_1^d \cap \mathbf{S}_2 \neq \emptyset \text{ or } \mathbf{S}_1^{-d} \cap \mathbf{S}_2 \neq \emptyset\}.$$

Since $\mathbf{q}$ is also a point on the surface $\mathbf{S}_2$, we have

$$f_2(\mathbf{p} + \mu \mathbf{N}_p) = 0. \tag{3}$$

Let $\mathbf{N}_q$ be the gradient of $\mathbf{S}_2$ at the point $\mathbf{q}$. Then, $\mathbf{N}_q$ and $\mathbf{N}_p$ are two collinear vectors. From (1), $\mathbf{N}_q$ is $(f_{2x}(\mathbf{q}), f_{2y}(\mathbf{q}), f_{2z}(\mathbf{q}))^{\mathrm{T}}$. Therefore, we obtain

$$\lambda \mathbf{N}_p + (f_{2x}(\mathbf{q}), f_{2y}(\mathbf{q}), f_{2z}(\mathbf{q}))^{\mathrm{T}} = \mathbf{0}, \tag{4}$$

where $\lambda \in \mathbb{R}$. Substituting Eq. (2) into Eq. (4), we have

$$\lambda \mathbf{N}_p + (f_{2x}(\mathbf{p} + \mu \mathbf{N}_p), f_{2y}(\mathbf{p} + \mu \mathbf{N}_p), f_{2z}(\mathbf{p} + \mu \mathbf{N}_p))^{\mathrm{T}} = \mathbf{0}. \tag{5}$$

The resultant method [1,3,6,7,46,47] is useful for eliminating one variable from two polynomial equations. Given two polynomial equations

$$F_a = a_0 + a_1 \mu + a_2 \mu^2 + \cdots + a_n \mu^n, \quad a_n \neq 0$$

and

$$F_b = b_0 + b_1 \mu + b_2 \mu^2 + \cdots + b_m \mu^n, \quad b_m \neq 0,$$

F_a and F_b have common roots if and only if

$$R_a^b = \begin{vmatrix} a_0 & a_1 & \cdots & a_n & & & \\ & a_0 & a_1 & \cdots & a_n & & \\ & & \cdots & \cdots & \cdots & \cdots & \cdots \\ & & & a_0 & a_1 & \cdots & a_n \\ b_0 & b_1 & \cdots & b_m & & & \\ & b_0 & b_1 & \cdots & b_m & & \\ & & \cdots & \cdots & \cdots & \cdots & \cdots \\ & & & b_0 & b_1 & \cdots & b_m \end{vmatrix} = 0.$$

R_a^b is called the resultant of F_a and F_b with variable μ. Since $f_2(x, y, z)$ is an implicit algebraic function about x, y and z, the equations in equation systems (3) and (5) are all polynomial equations in λ and μ. So we are able to eliminate λ and μ from equation systems (3) and (5) with the resultant method, and thus obtain

$$\begin{cases} g_1(x, y, z) = 0, \\ g_2(x, y, z) = 0. \end{cases}$$

Both of the above equations have three variables which are the coordinates of the point $\mathbf{p}$ on the surface $\mathbf{S}_1$. Therefore, one more equation is required to determine the point $\mathbf{p}$.

If the first surface $\mathbf{S}_1$ is an implicit surface given by $f_1(x, y, z) = 0$, then we have the equation system

$$\begin{cases} g_1(x, y, z) = 0, \\ g_2(x, y, z) = 0, \\ f_1(x, y, z) = 0. \end{cases} \tag{6}$$

If the first surface $\mathbf{S}_1$ is a parametric surface given by

$$(x, y, z)^{\mathrm{T}} = (X(u, v), Y(u, v), Z(u, v))^{\mathrm{T}},$$

where u, v are parameters, then we have the equation system

$$\begin{cases} g_1(X(u, v), Y(u, v), Z(u, v)) = 0, \\ g_2(X(u, v), Y(u, v), Z(u, v)) = 0. \end{cases} \tag{7}$$

In either case, from the equation system (6) or (7), by using the methods in [3,31], we can find a solution for the coordinates of the point $\mathbf{p}$. Substituting the coordinates of $\mathbf{p}$ into Eq. (3), we obtain the value of d. The other point $\mathbf{q}$ of the location is obtained from Eq. (2).

If $\mathbf{p}$ is singular, i.e., $\mathbf{N}_p$ is degenerated into a zero vector, $\mathbf{p}$ can be found by solving

$$\mathbf{N}_p = \mathbf{0}. \tag{8}$$

For each singular point $\mathbf{p}_1(x_1, y_1, z_1)^T \in \mathbf{S}_1$, the closest point $\mathbf{q}(x, y, z)$ on $\mathbf{S}_2$ can be obtained by solving

$$\begin{cases} f_2(x, y, z) = 0, \\ (x - x_1, y - y_1, z - z_1)^T \times (f_{2x}, f_{2y}, f_{2z})^T = \mathbf{0}. \end{cases} \quad (9)$$

The corresponding distance d is just $\|\mathbf{p}_1 - \mathbf{q}\|$. Thus, whether the closest point $\mathbf{p} \in \mathbf{S}_1$ is singular or not, from both the equation systems (6) or (7), (2) and (3) and the equation systems (8) and (9), we are able to find the minimum distance as well as the corresponding closest points $\mathbf{p} \in \mathbf{S}_1$ and $\mathbf{q} \in \mathbf{S}_2$.

The algorithm is as follows.

Input: Two surfaces $\mathbf{S}_1$ and $\mathbf{S}_2$. $\mathbf{S}_1$ is either an implicit surface or a parametric surface. $\mathbf{S}_2$ is an implicit algebraic surface.

Output: The minimum distance $d_{\min}$ between $\mathbf{S}_1$ and $\mathbf{S}_2$, and a pair of points $(\mathbf{p}, \mathbf{q})$ with the properties $\|\mathbf{p}-\mathbf{q}\| = d_{\min}$, $\mathbf{p} \in \mathbf{S}_1$ and $\mathbf{q} \in \mathbf{S}_2$.

(1) If $\mathbf{S}_1$ and $\mathbf{S}_2$ have any intersection point $\mathbf{p}$;

Output $d_{\min} = 0$ and the pair of points $(\mathbf{p}, \mathbf{p})$, then go to Step 5;

(2) Solve the equation systems (6) or (7); Substitute each solution $\mathbf{p} \in \mathbf{S}_1$ into Eq. (3) to obtain values of d; Compute the corresponding point $\mathbf{q} \in \mathbf{S}_2$ with Eq. (2);

Add all the solutions $(\mathbf{p}, \mathbf{q}, d)$ into the set $\mathbf{E}_d$;

(3) Solve the equation systems (8) to obtain the singular points on $\mathbf{S}_1$; For each singular point $\mathbf{p} \in \mathbf{S}_1$, compute the closest point $\mathbf{q} \in \mathbf{S}_2$ with equation system (9), and calculate $d = \|\mathbf{p} - \mathbf{q}\|$;

Add all the solutions $(\mathbf{p}, \mathbf{q}, d)$ into the set $\mathbf{E}_d$;

(4) Find $d_{\min}$ such that

$$d_{\min} = \min\{|d| \mid (\mathbf{p}, \mathbf{q}, d) \in \mathbf{E}_d, |d| = \|\mathbf{p} - \mathbf{q}\|\};$$

Output $d_{\min}$ and $(\mathbf{p}, \mathbf{q})$;

(5) **End** of the algorithm.

3. Minimum distance between a quadric surface and an implicit algebraic surface

Quadric surfaces and tori are widely used [22,39,45]. In this section, the first surface $\mathbf{S}_1$ is an implicit algebraic surface defined by $f_1(x, y, z) = 0$, and the second surface $\mathbf{S}_2$ is a quadric surface such as a cylinder, cone, elliptic paraboloid, ellipsoid or torus. The singular points of $\mathbf{S}_1$ can be found by solving $(f_{1x}, f_{1y}, f_{1z})^T = \mathbf{0}$, and it is easy to compute the distance between a point and $\mathbf{S}_2$. Here we omit the discussion on the singular points of $\mathbf{S}_1$.

For the convenience of degree analysis, we provide the following theorem.

Theorem 1. *Suppose that an equation system has three polynomial equations of degree 2, m and r. After eliminating one variable, one can obtain two bivariate equations of degree $\min\{2m, 2r\}$ and $m + r - 1$.*

Proof. Without loss of generality, assume $m \leq r$, and the first equation of degree 2 is of the form (if necessary, with coordinate transformation)

$$x^2 = F_2(y, z) + F_1(y, z)x, \quad (10)$$

where $F_i(y, z)$ is a polynomial of degree i in y and z, $i = 1, 2$. Substituting Eq. (10) into the other two equations of the equation system, we get two new equations

$$\begin{cases} G_{m-1}(y, z)x + G_m(y, z) = 0, \\ \bar{G}_{r-1}(y, z)x + \bar{G}_r(y, z) = 0, \end{cases} \quad (11)$$

where $G_i(y, z)$ and $\bar{G}_j(y, z)$ are polynomial of degree i and j in y and z, $i = m - 1, m$, and $j = r - 1, r$.

Combining Eq. (10) with the first equation in the equation system, we have

$$G_m(y, z)^2 - F_2(y, z)G_{m-1}(y, z)^2$$
$$- F_1(y, z)G_{m-1}(y, z)G_m(y, z) = 0. \quad (12)$$

From Eq. (11), we have

$$G_{r-1}(y, z)\bar{G}_m(y, z) - G_r(y, z)\bar{G}_{m-1}(y, z) = 0. \quad (13)$$

Eqs. (12) and (13) are of degree $2m$ and $m + r - 1$. Thus, we complete the proof. $\square$

When the three polynomial equations in an equation system are of degrees 2, 2 and r, respectively, where $r \geq 2$ is a positive integer, we can obtain two bivariate polynomial equations of degrees 4 and $r + 1$ by Theorem 1, or a univariate polynomial equation of degree $4r$ with Levin's method [23,43,48].

That is, in general detail: given two distinct quadrics $\mathbb{A} : \mathbf{X}^T \mathbf{A} \mathbf{X} = 0$ and $\mathbb{B} : \mathbf{X}^T \mathbf{B} \mathbf{X} = 0$, where $\mathbf{A}$, $\mathbf{B}$ are two distinct 4×4 matrices, and $\mathbf{X} = (x, y, z, 1)^T$, the pencil generated by $\mathbb{A}$ and $\mathbb{B}$ is the set $R(\lambda)$ with equations $\mathbf{X}^T(\mathbf{A} + \lambda \mathbf{B})\mathbf{X} = 0, \lambda \in \mathbb{R}$.

Levin [23] proves that there exists a real ruled quadric of the form

$$\mathbf{q}(u, v) = \mathbf{b}(u) + v\mathbf{d}(u)$$

in the pencil of two distinct quadrics, where the curve $\mathbf{b}(u)$ and ruling direction $\mathbf{d}(u)$ are at most quadratic rational functions in u [43,48]. Any intersection curve of the two quadrics (**QSIC**) can be parameterized in the form

$$\mathbf{p}(u) = \mathbf{a}(u) \pm \mathbf{d}(u)\sqrt{s(u)},$$

where $c_2(u) = \mathbf{d}(u)^T \mathbf{A} \mathbf{d}(u), c_1(u) = \mathbf{b}(u)^T \mathbf{A} \mathbf{d}(u), c_0(u) = \mathbf{b}(u)^T \mathbf{A} \mathbf{b}(u), \mathbf{a}(u) = c_2(u)\mathbf{b}(u) - c_1(u)\mathbf{d}(u), s(u) = c_1^2(u) - c_2(u)c_0(u)$ (Refer to [23,43,48] for more details). Substituting the intersection curve $\mathbf{p}(u)$ into the third equation of the equation system, and removing the radical expression, we obtain a univariate polynomial equation of degree $4r$.

3.1. Minimum distance between a cylinder and an implicit algebraic surface

Suppose the first surface $\mathbf{S}_2$ is a cylinder given by

$$x^2 + y^2 - r^2 = 0, \quad r > 0.$$

The equation system (6) becomes

$$\begin{cases} f_{1z} = 0, \\ x f_{1y} - y f_{1x} = 0, \\ f_1(x, y, z) = 0. \end{cases} \tag{14}$$

When $\mathbf{S}_1$ is a quadric surface, the equations in the equation system (14) are of degree 1, 2 and 2, and we can obtain two bivariate polynomial equations of degree 2, or a univariate polynomial of degree 4.

3.2. Minimum distance between a cone and an implicit algebraic surface

Suppose the first surface $\mathbf{S}_2$ is a cone given by

$$x^2 + y^2 - mz^2 = 0,$$

where m is a positive real number.

The equation system (6) becomes

$$\begin{cases} x f_{1y} - y f_{1x} = 0, \\ \lambda^2 (m f_{1x}^2 + m f_{1y}^2 - f_{1z}^2) = 0, \\ f_1(x, y, z) = 0, \end{cases} \tag{15}$$

where $\lambda = 0$ means that the candidate closest points on the cone are singular. So the equation system (15) is equivalent to two new equation systems

$$\begin{cases} x f_{1y} - y f_{1x} = 0, \\ m f_{1x}^2 + m f_{1y}^2 - f_{1z}^2 = 0, \\ f_1(x, y, z) = 0, \end{cases} \tag{16}$$

and

$$\begin{cases} (x, y, z)^T \times (f_{1x}, f_{1y}, f_{1z})^T = \mathbf{0}, \\ f_1(x, y, z) = 0. \end{cases} \tag{17}$$

When $\mathbf{S}_1$ is a quadric surface, the equations in the equation system (16) and (17) are all of degree 2. From each new equation system, we can obtain two bivariate polynomial equations of degree 3 and 4, or a univariate polynomial of degree 8.

3.3. Minimum distance between an elliptic paraboloid and an implicit algebraic surface

Suppose the first surface $\mathbf{S}_2$ is an elliptic paraboloid given by

$$x^2 + y^2 - 2pz = 0,$$

where p is a positive real number.

The equation system (6) becomes

$$\begin{cases} f_1(x, y, z) = 0, \\ x f_{1y} - y f_{1x} = 0, \\ p(f_{1x}^2 + f_{1y}^2) f_{1x} - 2 f_{1z}^2(z f_{1x} - x f_{1z} - p f_{1x}) = 0. \end{cases} \tag{18}$$

When $\mathbf{S}_1$ is a quadric surface, the three equations in the equation system (18) are of degrees 2, 2 and 4, and we can obtain two bivariate polynomials of degrees 4 and 5, or a univariate polynomial of degree 16.

3.4. Minimum distance between an ellipsoid and an implicit algebraic surface

Suppose the first surface $\mathbf{S}_2$ is an ellipsoid given by

$$\frac{x^2}{a^2} + \frac{y^2}{b^2} + \frac{z^2}{c^2} - 1 = 0.$$

If $a = b = c$, then we obtain

$$\begin{cases} f_1(x, y, z) = 0, \\ (x, y, z)^T \times (f_{1x}, f_{1y}, f_{1z})^T = \mathbf{0}. \end{cases}$$

Otherwise, suppose that $a \neq b$, the equation system (6) becomes

$$\begin{cases} f_1(x, y, z) = 0, \\ a^2 f_{1x}(y f_{1z} - z f_{1y}) + b^2 f_{1y}(z f_{1x} - x f_{1z}) \\ \quad + c^2 f_{1z}(x f_{1y} - y f_{1x}) = 0, \\ (x f_{1y} - y f_{1x})^2 (a^2 f_{1x}^2 + b^2 f_{1y}^2 + c^2 f_{1z}^2) \\ \quad - (b^2 - a^2)^2 f_{1x}^2 f_{1y}^2 = 0. \end{cases} \tag{19}$$

When $\mathbf{S}_1$ is a quadric surface, the three equations in the equation system (19) are of degrees 2, 3 and 6. By using the transformation $(x', y', z')^T = (x/f_{1x}, y/f_{1y}, z/f_{1z})^T$, we can obtain two bivariate polynomials of degree 6, or a univariate polynomial of degree 36, which can be factorized into two polynomials of degrees 12 and 24. We show that the equation system (19) will produce the same formulas as the method in [39]. The details are as follows: the normal congruence of $\mathbf{S}_1$, i.e., $(L_1, L_2, L_3, L_4, L_5, L_6)^T$, can be described as

$$(f_{1x}, f_{1y}, f_{1z}, y f_{1z} - z f_{1y}, z f_{1x} - x f_{1z}, x f_{1y} - y f_{1x})^T.$$

From the equation system (19), we obtain

$$\begin{cases} a^2 L_1 L_4 + b^2 L_2 L_5 + c^2 L_3 L_6 = 0, \\ L_6^2 (a^2 L_1^2 + b^2 L_2^2 + c^2 L_3^2) - L_1^2 L_2^2 (a^2 - b^2)^2 = 0. \end{cases} \tag{20}$$

Since $f_L = (L_1 L_4 + L_2 L_5 + L_3 L_6) = 0$, by adding $-c^2 f_L$ to the first equation of equation system (20), we obtain the same formulae as the method in [39].

3.5. Minimum distance between a torus and an implicit algebraic surface

Suppose the first surface $\mathbf{S}_2$ is a torus given by

$$(x^2 + y^2 + z^2 + R_2^2 - R_1^2)^2 - 4R_2^2(x^2 + y^2) = 0.$$

The equation system (6) becomes

$$\begin{cases} f_1(x, y, z) = 0, \\ x f_{1y} - y f_{1x} = 0, \\ (x f_{1z} - z f_{1x})^2 + (y f_{1z} - z f_{1y})^2 - R_2^2 f_{1z}^2 = 0. \end{cases} \tag{21}$$

When $\mathbf{S}_1$ is a quadric surface, the degrees of the three above equations are 2, 2 and 4, and we can obtain two bivariate polynomials of degrees 4 and 5, or a univariate polynomial of degree 16.

4. Minimum distance of a canal surface and an implicit algebraic surface

Our method is able to compute the minimum distance between an implicit algebraic surface and a canal surface. This section, therefore, also considers the case between two canal surfaces as well. It appears that this new method generalizes the method in [17].

A canal surface is a swept surface along the central trajectory curve $\mathbf{C}(t) = (x(t), y(t), z(t))^T$ and with the radius function $r(t)$. In this section, a canal surface is the same as that defined in [17], which is the set of all $\mathbf{x}$ satisfying

$$\begin{cases} \|\mathbf{x} - \mathbf{C}(t)\|^2 - r(t)^2 = 0, \\ \langle \mathbf{x} - \mathbf{C}(t), \mathbf{C}'(t) \rangle + r(t)r'(t) = 0. \end{cases} \quad (22)$$

In [17], the characteristic circle $\mathbf{K}(t)$ is the circle embedded in the moving sphere with center $\mathbf{C}_k(t) = \mathbf{C}(t) - r(t)r'(t)\frac{\mathbf{C}'(t)}{\|\mathbf{C}'(t)\|}$ and radius $r_k(t) = r(t)\frac{\sqrt{\|\mathbf{C}'(t)\| - r'^2(t)}}{\|\mathbf{C}'(t)\|}$. The normal of the main plane of the characteristic circle is $\mathbf{N}_k(t) = \mathbf{C}'(t)$. $\mathbf{K}(t)$ lies on the canal surface. Let $\mathbf{p}$ be a point on $\mathbf{K}(t)$. If

$$\mathbf{d}(t) = \mathbf{p} - \mathbf{C}(t)$$

is parallel to the normal of the canal surface at $\mathbf{p}$, then Kim [17] provides an equation for solving the property of the canal surface

$$\frac{\langle \mathbf{d}(t), \mathbf{C}'(t) \rangle}{\|\mathbf{d}(t)\|\|\mathbf{C}'(t)\|} + \frac{r'(t)}{\|\mathbf{C}'(t)\|} = 0. \quad (23)$$

Let $\mathbf{S}_1$ be an implicit algebraic surface defined by $f_1(x, y, z) = 0$, and the second surface $\mathbf{S}_2$ a canal surface defined by Eq. (22). Suppose the pair $(\mathbf{p}, \mathbf{q})$ is a location of the minimum distance. The line $\overline{\mathbf{pq}}$ intersects the central trajectory curve at a point $\mathbf{C}(t)$ and the vector $\mathbf{V} = \mathbf{p} - \mathbf{C}(t)$ is a normal of the canal surface at $\mathbf{p}$. Thus, according to Eq. (23), we have

$$\frac{\langle \mathbf{p} - \mathbf{C}(t), \mathbf{C}'(t) \rangle}{\|\mathbf{p} - \mathbf{C}(t)\|\|\mathbf{C}'(t)\|} + \frac{r'(t)}{\|\mathbf{C}'(t)\|} = 0. \quad (24)$$

$\mathbf{V}$ is also a normal of $\mathbf{S}_1$ at $\mathbf{p}$. We obtain

$$\begin{cases} \lambda(f_{1x}(\mathbf{p}), f_{1y}(\mathbf{p}), f_{1z}(\mathbf{p}))^T + (\mathbf{p} - \mathbf{C}(t))^T = \mathbf{0}, \\ f_1(x, y, z) = 0. \end{cases} \quad (25)$$

In particular, when $\mathbf{S}_1$ is a plane, cylinder, cone, or torus, the point $\mathbf{p}$ can be solved geometrically from (25). Substituting them into Eq. (24) yields the same result as in [17].

When $\mathbf{S}_1$ is a general quadric surface, the first three equations in the equation system (25) are linear in the coordinates of point $\mathbf{p}$ (which can be rewritten in matrix form $\mathbf{M}(\lambda)\mathbf{p} = \mathbf{C}(t)$) and we thus obtain $\mathbf{p} = \mathbf{M}(\lambda)^{-1}\mathbf{C}(t)$. Substituting it into Eq. (24) and the last equation in the equation system (25), we obtain two bivariate equations in λ and t. When $\mathbf{C}(t)$ and $r(t)$ are polynomials of degree n and m, the degrees of the two equations are $\max\{2n+4, 6\}$ and $\max\{4n+2, 2n+2m+2\}$, respectively. When $r(t)$ is a constant function, the degrees of the two equations are $\max\{2n+4, 6\}$ and $2n+1$, respectively.

When $\mathbf{S}_1$ is a canal surface, suppose that the two canal surfaces $\mathbf{S}_1$ and $\mathbf{S}_2$ are defined by Eq. (22) with their central trajectory curves and radius functions being $\mathbf{C}_1(u)$, $r_1(u)$ and $\mathbf{C}_2(v)$, $r_2(v)$, respectively. Suppose $\mathbf{p}$ and $\mathbf{q}$ are the closest points on $\mathbf{S}_1$ and $\mathbf{S}_2$, respectively. The line $\overline{\mathbf{pq}}$ intersects the two central trajectory curves at points $\mathbf{C}_1(u)$ and $\mathbf{C}_2(v)$, respectively. According to Eq. (23), we obtain

$$\begin{cases} \dfrac{\langle \mathbf{C}_1(u) - \mathbf{C}_2(v), \mathbf{C}_1'(u) \rangle}{\|\mathbf{C}_1(u) - \mathbf{C}_2(v)\|\|\mathbf{C}_1'(u)\|} + \dfrac{r_1'(u)}{\|\mathbf{C}_1'(u)\|} = 0, \\ \dfrac{\langle \mathbf{C}_1(u) - \mathbf{C}_2(v), \mathbf{C}_2'(v) \rangle}{\|\mathbf{C}_1(u) - \mathbf{C}_2(v)\|\|\mathbf{C}_2'(v)\|} + \dfrac{r_2'(v)}{\|\mathbf{C}_2'(v)\|} = 0. \end{cases} \quad (26)$$

Suppose $\mathbf{C}_1(u)$, $\mathbf{C}_2(v)$, $r_1(u)$ and $r_2(v)$ are polynomials of degrees $m_1 \leq m_2$, n_1 and n_2, respectively. Then the two equations in the equation system (26) are of degrees $\max\{2m_2 + 2m_1 - 2, 2m_2 + 2n_1 - 2\}$ and $\max\{4m_2 - 2, 2m_2 + 2n_2 - 2\}$, respectively. When $r_1(u)$ and $r_2(v)$ are constant functions, the degrees of the two equations are $m_2 + m_1 - 1$ and $2m_2 - 1$, respectively.

5. Comparisons

This section compares our method with the methods in [17, 22, 39]. Let $\mathbf{M}_1$, $\mathbf{M}_2$, $\mathbf{M}_3$ be the corresponding methods in [17, 22, 39] and $\mathbf{M}_n$ be the corresponding method in this paper respectively. Firstly, $\mathbf{M}_n$ is able to produce the same formulae as $\mathbf{M}_1$, and is also able to provide the solution for the case between a canal surface and an ellipsoid, or the case between two canal surfaces (see Fig. 2).

Secondly, regardless of whether a case is degenerated in $\mathbf{M}_2$ or not, $\mathbf{M}_n$ is able to deal with it uniformly. In principle, $\mathbf{M}_n$ is able to deal with the cases between two implicit algebraic surfaces of degree more than 2.

Thirdly, $\mathbf{M}_3$ presents formulae for the cases between an ellipsoid and a cylinder, a cone, a torus or a parametric surface, and $\mathbf{M}_n$ is able to produce the same formulae as $\mathbf{M}_3$. Given two surfaces, $\mathbf{M}_3$ needs to describe two normal congruences for the two surfaces, one in parametric form, and the other in implicit form. Usually $\mathbf{M}_3$ deals with the cases between an implicit algebraic surface and a parametric surface. To obtain a normal congruence for a general implicit surface, one needs to eliminate seven unknown variables from 10 polynomial equations (or eliminate four variables from seven polynomial equations by simplification), which is a non-trivial task. On the other hand, in $\mathbf{M}_n$, λ is easily eliminated by turning the equation system (5) into

$$\mathbf{N}_p \times (f_{2x}(\mathbf{p} + \mu\mathbf{N}_p), f_{2y}(\mathbf{p} + \mu\mathbf{N}_p), f_{2z}(\mathbf{p} + \mu\mathbf{N}_p))^T = \mathbf{0}, \quad (27)$$

and one only needs to eliminate μ from the equation system (3) and the above equation system (27), which is much easier. Furthermore, the parametric expression of an ellipsoid excludes one point on the ellipsoid, which should be considered additionally. For example, given two ellipsoids $\mathbf{S}_1$: $f_1(x, y, z) = x^2 + y^2/4 + (z + 3)^2/9 - 1 = 0$ and $\mathbf{S}_2$: $f_2(x, y, z) = x^2/4 + y^2/9 + z^2/49 - 1 = 0$. The
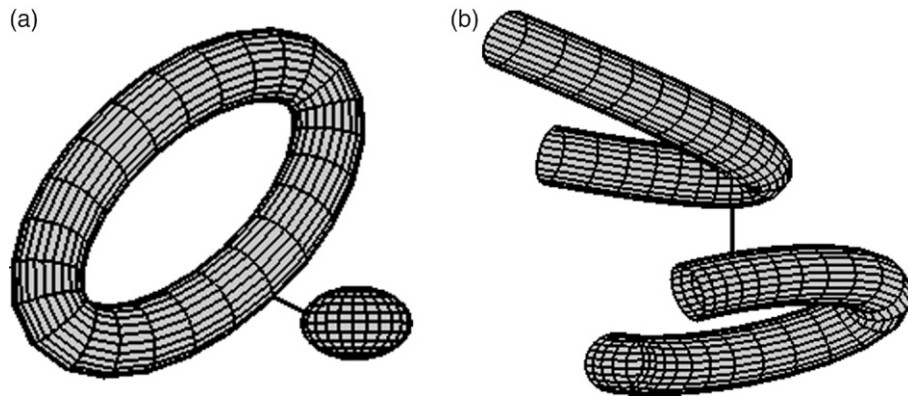
Fig. 2. The minimum distance between two surfaces. The first surface is a canal surface, and the second one is (a) an ellipsoid or (b) a canal surface.
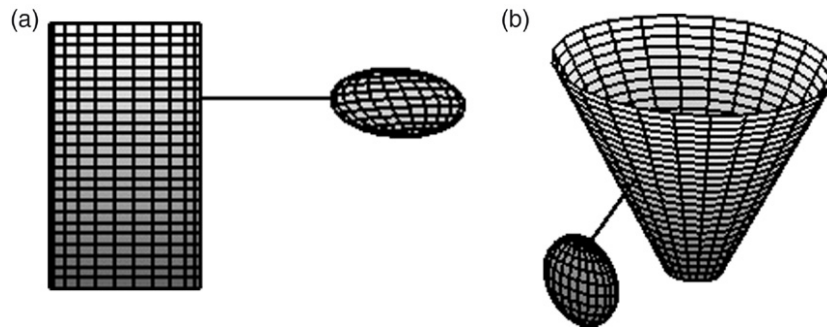

Fig. 3. The minimum distance between two implicit surfaces. The first surface is an ellipsoid, and the second one is (a) a cylinder or (b) a cone.

minimum distance is 1.0, and the location of the minimum distance is $((0, 0, -6)^T, (0, 0, -7)^T)$. $\mathbf{S}_1$ can be parameterized as $(\frac{2u}{1+u^2+v^2}, \frac{4v}{1+u^2+v^2}, \frac{-6(u^2+v^2)}{1+u^2+v^2})$, which excludes the closest point $(0, 0, -6)^T$.

$\mathbf{M}_2$, $\mathbf{M}_3$ and $\mathbf{M}_n$ are all able to obtain two resulting bivariate equations, in the unknown variables Lagrange multipliers, parameters of one surface and the coordinates of the closest point on one surface respectively. Figs. 3 and 4 show the cases between an ellipsoid and a cylinder, a cone, an elliptic paraboloid, or a torus, and the corresponding degrees of the two resulting bivariate equations are indicated in Table 1. In Table 1, the first surface $\mathbf{S}_1$ is always a general quadric surface, the second surface $\mathbf{S}_2$ is a general quadric surface, an ellipsoid, a cylinder, a cone, an elliptic paraboloid, a hyperboloid, or a torus. The pair (m, n) of numbers denotes the degrees of the two bivariate equations, and the single number denotes the degree of the resulting univariate polynomial. As shown in Table 1, the two bivariate equations obtained by our method have lower degrees than those in other methods.

We have implemented the algorithm proposed in this paper on a 1.7 GHz Windows PC using C + +, including the case between an ellipsoid and a cylinder, a cone, an ellipsoid, or a torus. we first use the method in [39] to obtain the initial solution, and then apply a standard Newton method in two variables. The average computation time for the initial solution is about 20 μs for the cases between an ellipsoid and a cylinder, a cone, or a torus, and 40 μs for the case between two ellipsoids. Table 2 shows the corresponding average computation time of

Table 1
Degree analysis for the case between two quadric surfaces

$\mathbf{S}_2$	$\mathbf{M}_1$	$\mathbf{M}_2$	$\mathbf{M}_3$	$\mathbf{M}_n$
A general quadric	–	(6, 6) or 24	(8, 16)	(6, 6) or 24
An ellipsoid	–	(6, 6) or 24	(8, 16)	(6, 6) or 24
A cylinder	–	(4, 4) or 8	(4, 4)	(2, 2) or 4
A cone	–	(4, 6) or 12	(4, 8)	(3, 4) or 8
An elliptic paraboloid	–	(5, 6) or 18	–	(4, 5) or 16
A hyperboloid	–	(6, 6) or 24	(8, 16)	(6, 6) or 24
A torus	–	–	(4, 8)	(4, 5) or 16

Table 2
Average computation time for an ellipsoid and a quadric by iterative search

Quadric	Cylinder (μs)	Cone (μs)	Ellipsoid (μs)	Torus (μs)
T_2	4.69	12.81	19.86	23.84
T_3	10.69	21.53	28.75	21.32
T_n	1.87	6.47	16.25	11.72

$\mathbf{M}_2$, $\mathbf{M}_3$, and $\mathbf{M}_n$ by iterative search, and the total time should be the sum of the computation time for initial solution and for iterative search. As Table 2 shows, $\mathbf{M}_n$ is much faster than $\mathbf{M}_2$ and $\mathbf{M}_3$. Since $\mathbf{M}_n$ produces a resulting univariate polynomial equation of degree 4, we are able to directly compute all the discrete solutions in less than 5 μs, which is much less than the computation time for the initial solution.

For the case between an ellipsoid and a torus, the above iterative search methods may fall into a local minimum, and it may be necessary to compute all the discrete solutions. The
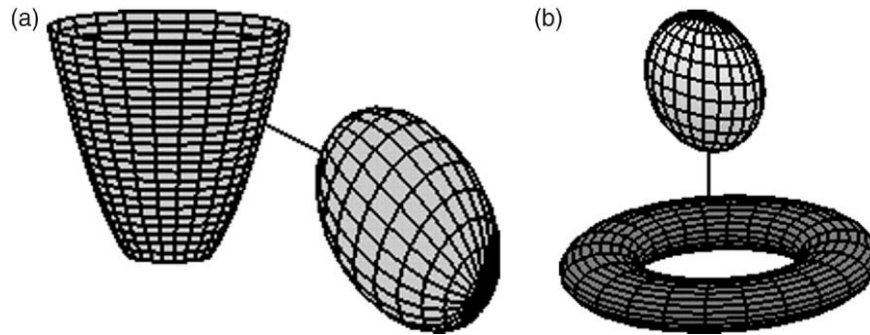
Fig. 4. The minimum distance between two implicit surfaces. The first surface is an ellipsoid, and the second one is (a) a paraboloid or (b) a torus.

algebraic and hybrid method explained in [31] is useful for computing all the discrete solutions, which turns the bivariate equation system into a univariate polynomial equation using the resultant method. The real root isolating method explained in [37] works with huge polynomials of degree 1000 and more, which is useful for solving the resulting univariate polynomial. To compute all the discrete solutions for the case between an ellipsoid and a torus, the average computation times for $\mathbf{M}_2$, $\mathbf{M}_3$ and $\mathbf{M}_n$ are 3235.0, 5417.2 and 2154.7 µs. Again, $\mathbf{M}_n$ is much faster than $\mathbf{M}_2$ and $\mathbf{M}_3$.

6. Conclusions

An offset method for computing the minimum distance between two implicit algebraic surfaces is presented in this paper. The offset method also works with the cases between an implicit algebraic surface and a parametric surface. Quadric surfaces, tori and canal surfaces are used to illustrate our method. Our method generalizes the method in [17]. As Table 1 shows, the two resulting bivariate equations in the new method have lower degrees than those in the methods explained in [22] and [39], which may lead to less computation time cost (see Table 2). Our method is able to produce the same (possibly simpler) formulae as the method in [39], but is much easier, for it only needs to eliminate one variable from three polynomial equations, while the method in [39] needs to eliminate four variables from seven polynomial equations.

The minimum distance computation problem usually involves non-linear equation systems. It is possible, with great difficulty, to compute all the roots for the corresponding general non-linear equation systems, but it is not necessary to do so. In the future, we will investigate an efficient method for finding a good candidate solution or for pruning redundant solutions. Another future task is to extend the offset method to the cases between two parametric surfaces.

Acknowledgements

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