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INSTITUT NATIONAL DE RECHERCHE EN INFORMATIQUE ET EN AUTOMATIQUE

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when the robber can hide and ride*

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## Cop and robber games when the robber can hide and ride

Jérémy Chalopin\* , Victor Chepoi\* , Nicolas Nisse† , Yann Vaxès\*

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**Abstract:** In the classical cop and robber game, two players, the cop  $\mathcal{C}$  and the robber  $\mathcal{R}$ , move alternatively along edges of a finite graph  $G = (V, E)$ . The cop captures the robber if both players are on the same vertex at the same moment of time. A graph  $G$  is called *cop win* if the cop always captures the robber after a finite number of steps. Nowakowski, Winkler (1983) and Quilliot (1983) characterized the cop-win graphs as graphs admitting a dismantling scheme. In this paper, we characterize in a similar way the cop-win graphs in the game in which the cop and the robber move at different speeds  $s'$  and  $s$ ,  $s' \leq s$ . We also investigate several dismantling schemes necessary or sufficient for the cop-win graphs in the game in which the robber is visible only every  $k$  moves for a fixed integer  $k > 1$ . We characterize the graphs which are cop-win for any value of  $k$ .

**Key-words:** Cops and Robber games, dismantling ordering, hyperbolicity

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\* Laboratoire d'Informatique Fondamentale, Université d'Aix-Marseille and CNRS, CMI and Faculté des Sciences de Luminy, F-13288 Marseille Cedex 9, France `firstname.lastname@lif.univ-mrs.fr`

† MASCOTTE, INRIA, I3S, CNRS, Univ. Nice Sophia, Sophia Antipolis, France. `firstname.lastname@sophia.inria.fr`

## Un gendarme face à un voleur véloce et discret

**Résumé :** Dans le jeu classique du gendarme et du voleur, deux joueurs, le gendarme  $\mathcal{C}$  et le voleur  $\mathcal{R}$ , se déplacent alternativement le long des arêtes d'un graphe fini  $G = (V, E)$ . Le gendarme capture le voleur si les deux joueurs occupent le même sommet simultanément. Un graphe  $G$  est dit *cop win* si le gendarme peut toujours capturer le voleur après un nombre fini de mouvements. Nowakowski, Winkler (1983) et Quilliot (1983) ont caractérisé les graphes *cop win* comme étant les graphes dont les sommets admettent un ordre de *démentèlement*. Dans ce papier, nous caractérisons de façon similaire les graphes *cop win* dans le jeu dans lequel le gendarme et le voleur se déplacent à des vitesses différentes  $s'$  et  $s$ ,  $s' \leq s$ . Nous nous intéressons également à plusieurs ordres de démentèlement nécessaires ou suffisants pour caractériser les graphes *cop win* dans le jeu dans lequel le voleur n'est visible que toutes les  $k$  étapes, pour  $k > 1$  fixé. Nous caractérisons les graphes *cop win* pour tous  $k$ .

**Mots-clés :** Jeux des gendarmes et du voleur, ordre de démentèlement, hyperbolicité

## 1 Introduction

### 1.1 The cop and robber game(s)

The cop and robber game originated in the 1980's with the work of Nowakowski, Winkler [30], Quilliot [31], and Aigner, Fromme [2], and since then has been intensively investigated by numerous authors and under different names (e.g., hunter and rabbit game [27]). Cop and robber is a pursuit-evasion game played on finite undirected graphs. Player cop  $\mathcal{C}$  has one or several cops who attempt to capture the robber  $\mathcal{R}$ . At the beginning of the game,  $\mathcal{C}$  occupies vertices for the initial position of his cops, then  $\mathcal{R}$  occupies another vertex. Thereafter, the two sides move alternatively, starting with  $\mathcal{C}$ , where a move is to slide along an edge or to stay at the same vertex, i.e. pass. Both players have full knowledge of the current positions of their adversaries. The objective of  $\mathcal{C}$  is to capture  $\mathcal{R}$ , i.e., to be at some moment of time, or *step*, at the same vertex as the robber. The objective of  $\mathcal{R}$  is to continue evading the cop. A *cop-win graph* [2, 30, 31] is a graph in which a *single cop* captures the robber after a finite number of moves for all possible initial positions of  $\mathcal{C}$  and  $\mathcal{R}$ . Denote by  $\mathcal{CW}$  the set of all cop-win graphs. The cop-number of a graph  $G$ , introduced by Aigner and Fromme [2], is the minimum number of cops necessary to capture the robber. Different combinatorial (lower and upper) bounds on the cop number for different classes of graphs were given in [2, 4, 9, 17, 22, 32, 33, 34] (see also the survey paper [3] and the annotated bibliography [21]).

In this paper, we investigate the cop-win graphs for three basic variants of the classical cop and robber game (for continuous analogous of these games, see [21]). In the *cop and fast robber game*, introduced by Fomin, Golovach, and Kratochvil [19] and further investigated in [29] (see also [20]), the cop is moving at unit speed while the speed of the robber is an integer  $s \geq 1$  or is unbounded ( $s \in \mathbb{N} \cup \{\infty\}$ ), i.e., at his turn,  $\mathcal{R}$  moves along a path of length at most  $s$  which does not contain vertices occupied by  $\mathcal{C}$ . Let  $\mathcal{CWF}\mathcal{R}(s)$  denote the class of all graphs in which a single cop having speed 1 captures a robber having speed  $s$ . Obviously,  $\mathcal{CWF}\mathcal{R}(1) = \mathcal{CW}$ . In a more general version, we will suppose that  $\mathcal{R}$  moves with speed  $s$  and  $\mathcal{C}$  moves with speed  $s' \leq s$  (if  $s' > s$ , then the cop can always capture the robber by strictly decreasing at each move his distance to the robber). We will denote the class of cop-win graphs for this version of the game by  $\mathcal{CWF}\mathcal{R}(s, s')$ . A *witness version* of the cop and robber game was recently introduced by Clarke [18]. In this game, the robber has unit speed and moves by having perfect information about cop positions. On the other hand, the cop no longer has full information about robber's position but receives it only occasionally, say every  $k$  units of time, in which case, we say that  $\mathcal{R}$  is *visible* to  $\mathcal{C}$ , otherwise,  $\mathcal{R}$  is *invisible* (this kind of constraint occurs, for instance, in the "Scotland Yard" game [14]). Following [18], we call a graph  $G$  *k-winnable* if a single cop can guarantee a win with such witness information and denote by  $\mathcal{CWW}(k)$  the class of all *k-winnable* graphs. Notice that  $\mathcal{CWF}\mathcal{R}(s) \subset \mathcal{CWW}(s)$  because the first game can be viewed as a particular version of the second game in which  $\mathcal{C}$  moves only at the turns when he receives the information about  $\mathcal{R}$ . Finally, the game of *distance k cop and robber* introduced by Bonato and Chiniforooshan [11] is played in the same way as classical cop and robber, except that the cop wins if a cop

is within distance at most  $k$  from the robber (following the name of an analogous game in continuous spaces [21], we will refer to this game as *cop and robber with radius of capture  $k$* ). We denote by  $\mathcal{CWR}\mathcal{C}(k)$  the set of all cop-win graphs in this game.

## 1.2 Cop-win graphs

Cop-win graphs have been characterized by Nowakowski and Winkler [30], and Quillot [32] (see also [2]) as dismantlable graphs (see Section 1.4 for formal definitions). Let  $G = (V, E)$  be a graph and  $u, v$  two vertices of  $G$  such that any neighbor of  $v$  (including  $v$  itself) is also a neighbor of  $u$ . Then there is a retraction of  $G$  to  $G \setminus \{v\}$  taking  $v$  to  $u$ . Following [25], we call this retraction a *fold* and we say that  $v$  is *dominated* by  $u$ . A graph  $G$  is *dismantlable* if it can be reduced, by a sequence of folds, to a single vertex. In other words, an  $n$ -vertex graph  $G = (V, E)$  is dismantlable if its vertices can be ordered  $v_1, \dots, v_n$  so that for each vertex  $v_i, 1 \leq i < n$ , there exists another vertex  $v_j$  with  $j > i$ , such that  $N_1(v_i) \cap X_i \subseteq N_1(v_j) \cap X_i$ , where  $X_i := \{v_i, v_{i+1}, \dots, v_n\}$  and  $N_1(v)$  denotes the closed neighborhood of  $v$ . For a simple proof that dismantlable graphs are the cop-win graphs, see the book [25]. An alternative (more algorithmic) proof of this result is given in [27]. Dismantlable graphs comprise bridged graphs (graphs in which all isometric cycles have length 3) and Helly graphs (absolute retracts) [6, 25] which occur in several other contexts in discrete mathematics. Except the cop and robber game, dismantlable graphs are used to model physical processes like phase transition [13], while bridged graphs occur as 1-skeletons of systolic complexes in the intrinsic geometry of simplicial complexes [15, 24, 26]. Dismantlable graphs are closed under retracts and direct products, i.e., they constitute a variety [30].

## 1.3 Our results

In this paper, we characterize the graphs of the class  $\mathcal{CWR}\mathcal{R}(s, s')$  for all speeds  $s, s'$  in the same vein as cop-win graphs, by using a specific dismantling order. Our characterization allows to decide in polynomial time if a graph  $G$  belongs to any of considered classes  $\mathcal{CWR}\mathcal{R}(s, s')$ . In the particular case  $s' = 1$ , we show that  $\mathcal{CWR}\mathcal{R}(2)$  is exactly the well-known class of dually chordal graphs. Then we show that, for all  $s \geq 3$ , the classes  $\mathcal{CWR}\mathcal{R}(s)$  coincide and that the graphs  $G$  of these classes have the following structure: the block-decomposition of  $G$  can be rooted in such a way that any block has a dominating vertex and that for each non-root block, this dominating vertex can be chosen to be the articulation point separating the block from the root. Then we characterize the graphs that are  $s$ -winnable for all  $s$  (i.e., graphs in  $\bigcap_{s \geq 1} \mathcal{CWW}(s)$ ) using a similar decomposition. On the other hand, we show that for each  $s$ ,  $\mathcal{CWW}(s) \setminus \mathcal{CWW}(s+1)$  is non-empty, contrary to the classes  $\mathcal{CWR}\mathcal{R}(s)$ . We show that all graphs of  $\mathcal{CWW}(2)$ , i.e., the 2-winnable graphs, have a special dismantling order (called bidismantling), which however does not ensure that a graph belongs to  $\mathcal{CWW}(2)$ . We present a stronger version of bidismantling and show that it is sufficient for ensuring that a graph is 2-winnable. We extend bidismantling to any  $k \geq 3$  and prove that for all odd  $k$ , bidismantling is sufficient to ensure that  $G \in \mathcal{CWR}\mathcal{R}(k)$ .

Finally, we characterize the bipartite members of  $CWRC(1)$  via an appropriate dismantling scheme. We also formulate several open questions.

## 1.4 Preliminaries

For a graph  $G = (V, E)$  and a subset  $X$  of its vertices, we denote by  $G(X)$  the subgraph of  $G$  induced by  $X$ . We will write  $G \setminus \{x\}$  and  $G \setminus \{x, y\}$  instead of  $G(V \setminus \{x\})$  and  $G(V \setminus \{x, y\})$ . The *distance*  $d(u, v) := d_G(u, v)$  between two vertices  $u$  and  $v$  of a graph  $G$  is the length (number of edges) of a shortest  $(u, v)$ -path. An induced subgraph  $H$  of  $G$  is *isometric* if the distance between any pair of vertices in  $H$  is the same as that in  $G$ . The *ball* (or disk)  $N_r(x)$  of center  $x$  and radius  $r \geq 0$  consists of all vertices of  $G$  at distance at most  $r$  from  $x$ . In particular, the unit ball  $N_1(x)$  comprises  $x$  and the neighborhood  $N(x)$ . The *punctured ball*  $N_r(x, G \setminus \{y\})$  of center  $x$ , radius  $r$ , and puncture  $y$  is the set of all vertices of  $G$  which can be connected to  $x$  by a path of length at most  $r$  avoiding the vertex  $y$ , i.e., this is the ball of radius  $r$  centered at  $x$  in the graph  $G \setminus \{y\}$ . Denote by  $N_\infty(v, G)$  the connected component of  $G$  that contains the vertex  $v$ . A *retraction*  $\varphi$  of a graph  $H = (W, F)$  is an idempotent nonexpansive mapping of  $H$  into itself, that is,  $\varphi^2 = \varphi : W \rightarrow W$  with  $d(\varphi(x), \varphi(y)) \leq d(x, y)$  for all  $x, y \in W$ . The subgraph of  $H$  induced by the image of  $H$  under  $\varphi$  is referred to as a *retract* of  $H$ .

A *strategy* for the cop is a function  $\sigma$  which takes as an input the first  $i$  moves of both players and outputs the  $(i + 1)$ th move  $c_{i+1}$  of the cop. A strategy for the robber is defined in a similar way. A cop's strategy  $\sigma$  is *winning* if for any sequence of moves of the robber, the cop, following  $\sigma$ , captures the robber after a finite sequence of moves. Note that if the cop has a winning strategy  $\sigma$  in a graph  $G$ , then there exists a winning strategy  $\sigma'$  for the cop that only depends of the last positions of the two players (such a strategy is called *positional*). This is because cop and robber games are parity games (by considering the directed graph of configurations) and parity games always admit positional strategies for the winning player [28]. A strategy for the cop is called *parsimonious* if at his turn, the cop captures the robber (in one move) whenever he can. For example, in the cop and fast robber game, at his move, the cop following an parsimonious strategy always captures a robber located at distance at most  $s'$  from his current position. It is easy to see that in the games investigated in this paper, if the cop has a (positional) winning strategy, then he also has a parsimonious (positional) winning strategy.

## 2 Cop-win graphs for game with fast robber: class $CWFR(s, s')$

In this section, first we characterize the graphs of  $CWFR(s, s')$  via a specific dismantling scheme, allowing to recognize them in polynomial time. Then we show that any  $\delta$ -hyperbolic graph belongs to the class  $CWFR(2r, r + 2\delta)$  for any  $r \geq 1$ . We conjecture that the converse is true, i.e., any graph from  $CWFR(s, s')$  with  $s' < s$  is  $\delta$ -hyperbolic for some value of  $\delta$  depending only of  $s$ , and we confirm this conjecture in several particular cases.



## 2.1 Graphs of $\mathcal{CWF}\mathcal{R}(s, s')$

For technical convenience, we will consider a slightly more general version of the game: given a subset of vertices  $X$  of a graph  $G = (V, E)$ , the  $X$ -restricted game with cop and robber having speeds  $s'$  and  $s$ , respectively, is a variant in which  $\mathcal{C}$  and  $\mathcal{R}$  can pass through any vertex of  $G$  but can stand only at vertices of  $X$  (i.e., the beginning and the end of each move are in  $X$ ). A subset of vertices  $X$  of a graph  $G = (V, E)$  is  $(s, s')$ -winnable if the cop captures the robber in the  $X$ -restricted game. In the following, given a subset  $X$  of admissible positions, we say that a sequence of vertices  $S_r = (a_1, \dots, a_p, \dots)$  of a graph  $G = (V, E)$  is  $X$ -valid for a robber with speed  $s$  (respectively, for a cop with speed  $s'$ ) if for any  $k$  we have  $a_k \in X$  and  $d(a_{k-1}, a_k) \leq s$  (respectively,  $d(a_{k-1}, a_k) \leq s'$ ). We will say that a subset of vertices  $X$  of a graph  $G = (V, E)$  is  $(s, s')$ -dismantlable if the vertices of  $X$  can be ordered  $v_1, \dots, v_m$  in such a way that for each vertex  $v_i, 1 \leq i < m$ , there exists another vertex  $v_j$  with  $j > i$ , such that  $N_s(v_i, G \setminus \{v_j\}) \cap X_i \subseteq N_{s'}(v_j)$ , where  $X_i := \{v_i, v_{i+1}, \dots, v_m\}$  and  $X_m = \{v_m\}$ . A graph  $G = (V, E)$  is  $(s, s')$ -dismantlable if its vertex-set  $V$  is  $(s, s')$ -dismantlable.

**Theorem 1.** *For any  $s, s' \in \mathbb{N} \cup \{\infty\}$ ,  $s' \leq s$ , a graph  $G = (V, E)$  belongs to the class  $\mathcal{CWF}\mathcal{R}(s, s')$  if and only if  $G$  is  $(s, s')$ -dismantlable.*

*Proof.* First, suppose that  $G$  is  $(s, s')$ -dismantlable and let  $v_1, \dots, v_n$  be an  $(s, s')$ -dismantling ordering of  $G$ . By induction on  $n - i$  we will show that for each level-set  $X_i = \{v_i, \dots, v_n\}$  the cop captures the robber in the  $X_i$ -restricted game. This is obviously true for  $X_n = \{v_n\}$ . Suppose that our assertion is true for all sets  $X_n, \dots, X_{i+1}$  and we will show that it still holds for  $X_i$ . Let  $N_s(v_i, G \setminus \{v_j\}) \cap X_i \subseteq N_{s'}(v_j)$  for a vertex  $v_j \in X_i$ . Consider a parsimonious positional winning strategy  $\sigma_{i+1}$  for the cop in the  $X_{i+1}$ -restricted game. We build a parsimonious winning strategy  $\sigma_i$  for the cop in the  $X_i$ -restricted game: the intuitive idea is that if the cop sees the robber in  $v_i$ , he plays as in the  $X_{i+1}$ -restricted game when the robber is in  $v_j$ . Let  $\sigma_i$  be the strategy for the  $X_i$ -restricted game defined as follows. For any positions  $c \in X_i$  of the cop and  $r \in X_i$  of the robber, set  $\sigma_i(c, r) = r$  if  $d(c, r) \leq s'$ , otherwise  $\sigma_i(c, r) = \sigma_{i+1}(c, r)$  if  $c, r \neq v_i$ ,  $\sigma_i(c, v_i) = \sigma_{i+1}(c, v_j)$  if  $c \notin \{v_i, v_j\}$ , and  $\sigma_i(v_i, r) = v_j$  if  $r \neq v_i$  (in fact, if the cop plays  $\sigma_i$  he will never move to  $v_i$  except to capture the robber there). By construction, the strategy  $\sigma_i$  is parsimonious; in particular,  $\sigma_i(v_j, v_i) = v_i$ , because  $d(v_i, v_j) \leq s'$ . We now prove that  $\sigma_i$  is winning.

Consider any  $X_i$ -valid sequence  $S_r = (r_1, \dots, r_p, \dots)$  of moves of the robber and any trajectory  $(\pi_1, \dots, \pi_p, \dots)$  extending  $S_r$ , where  $\pi_p$  is a simple path of length at most  $s$  from  $r_p$  to  $r_{p+1}$  along which the robber moves. Let  $S'_r = (r'_1, \dots, r'_p, \dots)$  be the sequence obtained by setting  $r'_k = r_k$  if  $r_k \neq v_i$  and  $r'_k = v_j$  if  $r_k = v_i$ . For each  $p$ , set  $\pi'_p = \pi_p$  if  $v_i \notin \{r_p, r_{p+1}\}$ . If  $v_i = r_{p+1}$  (resp.  $v_i = r_p$ ), set  $\pi'_p$  be a shortest path from  $r_p$  to  $v_j$  (resp. from  $v_j$  to  $r_{p+1}$ ) if  $\pi_p$  does not contain  $v_j$  and set  $\pi'_p$  be the subpath of  $\pi_p$  between  $r_p$  and  $v_j$  (resp. between  $v_j$  and  $r_{p+1}$ ) otherwise. Since  $N_s(v_i, G \setminus \{v_j\}) \cap X_i \subseteq N_{s'}(v_j)$ , we infer that  $S'_r$  is a  $X_{i+1}$ -valid sequence of moves for the robber. By induction hypothesis, for any initial location of  $\mathcal{C}$  in  $X_{i+1}$ , the strategy  $\sigma_{i+1}$  allows the cop to capture the robber which moves according to  $S'_r$  in the  $X_{i+1}$ -restricted game. Let  $c'_{m+1}$  be the position of the cop after his last move and  $S'_c = (c'_1, \dots, c'_{m+1})$  be the sequence of positions of the cop in the  $X_{i+1}$ -restricted game

against  $S'_r$  using  $\sigma_{i+1}$ . Let  $S_c = (c_1, \dots, c_p, \dots)$  be the sequence of positions of the cop in the  $X_i$ -restricted game against  $S_r$  using  $\sigma_i$ . From the definition of  $S'_r$  and  $\sigma_i$ ,  $S_c$  and  $S'_c$  coincide at least until step  $m$ , i.e.,  $c'_k = c_k$  for  $k = 1, \dots, m$ . Moreover, if  $c'_{m+1} \neq c_{m+1}$  then  $c_{m+1} = r_m = v_i$  and  $c'_{m+1} = r'_m = v_j$ . In the  $X_{i+1}$ -restricted version of the game, the robber is captured, either (i) because after his last move, his position  $r'_m$  is at distance at most  $s'$  from cop's current position  $c'_m$ , or (ii) because his trajectory  $\pi'_m$  from  $r'_m$  to  $r'_{m+1}$  passes via  $c'_{m+1}$ .

In case (i), since  $d(r'_m, c'_m) \leq s'$  and the strategy  $\sigma_{i+1}$  is parsimonious, we conclude that  $c'_{m+1} = r'_m$ . If  $c'_{m+1} = r'_m \neq v_j$ , then from the definition of  $S'_r$  and  $\sigma_i$ , we conclude that  $c_{m+1} = c'_{m+1} = r'_m = r_m$ , whence  $c_{m+1} = r_m$  and  $\mathcal{C}$  captures  $\mathcal{R}$  using  $\sigma_i$ . Now suppose that  $c'_{m+1} = r'_m = v_j$ . If  $r_m = v_j$ , then  $d(c_m, r_m) \leq s'$  because  $c_m = c'_m$  and thus  $\mathcal{C}$  captures  $\mathcal{R}$  at  $v_j$  using  $\sigma_i$ . On the other hand, if  $r_m = v_i$ , either  $c_{m+1} = v_i$  and we are done, or  $c_{m+1} = v_j$  and since  $N_s(v_i, G \setminus \{v_j\}) \cap X_i \subseteq N_{s'}(v_j)$ , the robber is captured at the next move of the cop, i.e.,  $c_{m+2} = r_{m+1}$  holds.

In case (ii), either the path  $\pi'_m$  from  $r'_m$  to  $r'_{m+1}$  is a subpath of  $\pi_m$ , or  $v_i \in \{r_m, r_{m+1}\}$  and  $\pi_m$  does not go via  $v_j$ . In the first case, note that  $c_{m+1} = c'_{m+1}$ , otherwise  $c_{m+1} = v_i = r_m$  by construction of  $\sigma_i$  and thus the robber has been captured before. Therefore the trajectory  $\pi_m$  of the robber in the  $X_i$ -game traverses the position  $c_{m+1}$  of the cop and we are done. Now suppose that  $\pi_m$  does not go via  $v_j$  and  $v_i \in \{r_m, r_{m+1}\}$ . Note that in this case,  $c_{m+1} = c'_{m+1}$  holds; otherwise,  $c'_{m+1} = r'_m = v_j$  and  $c_{m+1} = r_m = v_i$  and therefore, the robber is caught at step  $m+1$ . If  $c_{m+1}$  belongs to  $\pi_m$ , then we are done as in the first case. So suppose that  $c_{m+1} \notin \pi_m$ . If  $r_{m+1} = v_i$ , then  $r_m \in N_s(v_i, G \setminus \{v_j\}) \subseteq N_{s'}(v_j)$ . Since,  $\pi'_m$  is a shortest path and  $c'_{m+1}$  belongs to this path,  $d(c'_{m+1}, v_j) \leq s'$  and thus either  $c_{m+2} = v_i = r_{m+1}$  if  $d(c'_{m+1}, v_i) \leq s'$ , or  $c_{m+2} = v_j$  since  $\sigma_{i+1}$  is parsimonious. In the latter case, since  $N_s(v_i, G \setminus \{v_j\}) \subseteq N_{s'}(v_j)$ ,  $r_{m+1} = v_i$ , and  $c_{m+2} = v_j$ , the robber will be captured at the next move. Finally, suppose that  $r_m = v_i$ . Then  $r'_m = v_j$ . Since  $\pi_m$  is a path of length at most  $s$  avoiding  $v_j$ , we conclude that  $r_{m+1} \in N_s(v_i, G \setminus \{v_j\}) \subseteq N_{s'}(v_j)$ . Since  $\pi'_m$  is a shortest path from  $v_j$  to  $r_{m+1}$  containing the vertex  $c'_{m+1} = c_{m+1}$ , we have  $d(c_{m+1}, r_{m+1}) \leq d(v_j, r_{m+1}) \leq s'$ . Therefore, the cop captures the robber in  $r_{m+1}$  at his next move, i.e.,  $c_{m+2} = r_{m+1}$ . This shows that a  $(s, s')$ -dismantlable graph  $G$  belongs to  $\mathcal{CWF}\mathcal{R}(s, s')$ .

Conversely, suppose that for a  $X$ -restricted game played on a graph  $G = (V, E)$  there is a positional winning strategy  $\sigma$  for the cop. We assert that  $X$  is  $(s, s')$ -dismantlable. Consider a  $X$ -valid sequence of moves of the robber having a maximum number of steps before the capture of the robber. Let  $u \in X$  be the position occupied by the cop before the capture of  $\mathcal{R}$  and let  $v \in X$  be the position of the robber at this step. Since wherever the robber moved next in  $X$  (including remaining in  $v$  or passing via  $u$ ), the cop would capture him, necessarily  $N_s(v, G \setminus \{u\}) \cap X \subseteq N_{s'}(u)$  holds. Set  $X' := X \setminus \{v\}$ .

We assert that  $X'$  is  $(s, s')$ -winnable as well. In this proof, we use a strategy that is not positional but uses one bit of memory. A strategy using one bit memory can be presented as follows: it is a function which takes as input the current positions of the two players and a boolean (the current value of the memory) and that outputs the next position of the cop

and a boolean (the new value of the memory). Using the positional winning strategy  $\sigma$ , we define  $\sigma'(c, r, m)$  for any positions  $c \in X'$  of the cop and  $r \in X'$  of the robber and for any value of the memory  $m \in \{0, 1\}$ . The intuitive idea for defining  $\sigma'$  is that the cop plays using  $\sigma$  except when he is in  $u$  and his memory contains 1; in this case, he uses  $\sigma$  as if he was in  $v$ . If  $m = 0$  or  $c \neq u$ , then we distinguish two cases: if  $\sigma(c, r) = v$  then  $\sigma'(c, r, m) = (u, 1)$  (this is a valid move since  $N_{s'}(v) \cap X \subseteq N_{s'}(u)$ ) and  $\sigma'(c, r, m) = (\sigma(c, r), 0)$  otherwise. If  $m = 1$  and  $c = u$ , we distinguish two cases: if  $\sigma(v, r) = v$ , then  $\sigma'(u, r, 1) = (u, 1)$  and  $\sigma'(u, r, 1) = (\sigma(v, r), 0)$  otherwise (this is a valid move since  $N_{s'}(v) \cap X \subseteq N_{s'}(u)$ ). Let  $S_r = (r_1, \dots, r_p, \dots)$  be any  $X'$ -valid sequence of moves of the robber. Since  $X' \subset X$ ,  $S_r$  is also a  $X$ -valid sequence of moves of the robber. Let  $S_c := (c_1, \dots, c_p, \dots)$  be the corresponding  $X$ -valid sequence of moves of the cop following  $\sigma$  against  $S_r$  in  $X$  and let  $S'_c = (c'_1, \dots, c'_p, \dots)$  be the  $X'$ -valid sequence of moves of the cop following  $\sigma'$  against  $S_r$ . Note that the sequences of moves  $S_c$  and  $S'_c$  differ only if  $c_k = v$  and  $c'_k = u$ . Finally, since the cop follows a winning strategy for  $X$ , there is a step  $j$  such that  $c_j = r_j \in X \setminus \{v\}$  (note that  $r_j \neq v$  because we supposed that  $S_r \subseteq X'$ ). Since  $c_j \neq v$ , we also have  $c'_j = r_j$ , thus  $\mathcal{C}$  captures  $\mathcal{R}$  in the  $X'$ -restricted game. Starting from a positional strategy for the  $X$ -restricted game, we have constructed a winning strategy using memory for the  $X'$ -restricted game. As mentioned in the introduction, it implies that there exists a positional winning strategy for the  $X'$ -restricted game.

Applying induction on the number of vertices of the cop-winning set  $X$ , we conclude that  $X$  is  $(s, s')$ -dismantlable. Applying this assertion to the vertex set  $V$  of cop-win graph  $G = (V, E)$  from the class  $\mathcal{CWFR}(s, s')$ , we will conclude that  $G$  is  $(s, s')$ -dismantlable.  $\square$

**Corollary 1.** *Given a graph  $G = (V, E)$  and the integers  $s, s' \in \mathbb{N} \cup \{\infty\}$ ,  $s' \leq s$ , one can recognize in polynomial time if  $G$  belongs to  $\mathcal{CWFR}(s, s')$ .*

*Proof.* By Theorem 1,  $G \in \mathcal{CWFR}(s, s')$  if and only if  $G$  is  $(s, s')$ -dismantlable. Moreover, from the last part of the proof of Theorem 1 we conclude that if a subset  $X$  of vertices of  $G$  is  $(s, s')$ -winnable and  $u, v \in X$  such that  $N_s(v, G \setminus \{u\}) \cap X \subseteq N_{s'}(u)$  holds, then the set  $X' = X \setminus \{v\}$  is  $(s, s')$ -winnable as well. Therefore it suffices to run the following algorithm. Start with  $X := V$  and as long as possible find in  $X$  two vertices  $u, v$  satisfying the inclusion  $N_s(v, G \setminus \{u\}) \cap X \subseteq N_{s'}(u)$ , and set  $X := X \setminus \{v\}$ . If the algorithm ends up with a set  $X$  containing at least two vertices, then  $G$  is not  $(s, s')$ -winnable, otherwise, if  $X$  contains a single vertex, then  $G$  is  $(s, s')$ -dismantlable and therefore  $G \in \mathcal{CWFR}(s, s')$ .  $\square$

## 2.2 Graphs of $\mathcal{CWFR}(s, s')$ and hyperbolicity

Introduced by Gromov [23],  $\delta$ -hyperbolicity of a metric space measures, to some extent, the deviation of a metric from a tree metric. A graph  $G$  is  $\delta$ -hyperbolic if for any four vertices  $u, v, x, y$  of  $G$ , the two larger of the three distance sums  $d(u, v) + d(x, y)$ ,  $d(u, x) + d(v, y)$ ,  $d(u, y) + d(v, x)$  differ by at most  $2\delta \geq 0$ . Every 4-point metric  $d$  has a canonical representation in the rectilinear plane as illustrated in Fig. 1. The three distance sums are ordered from small to large, thus implying  $\xi \leq \eta$ . Then  $\eta$  is half the difference of the

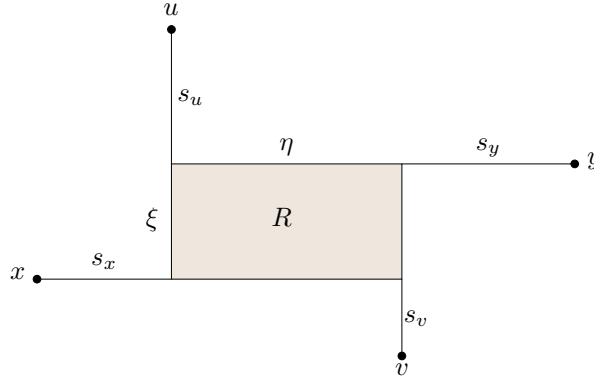


Figure 1: Realization of a 4-point metric in the rectilinear plane.

largest and the smallest sum, while  $\xi$  is half the largest minus the medium sum. Hence, a graph  $G$  is  $\delta$ -hyperbolic iff  $\xi$  does not exceed  $\delta$  for any four vertices  $u, v, w, x$  of  $G$ . Many classes of graphs are known to have bounded hyperbolicity [6, 16]. Our next result, based on Theorem 1 and a result of [16], establishes that in a  $\delta$ -hyperbolic graph a “slow” cop captures a faster robber provided that  $s' > s/2 + 2\delta$  (in the same vein, Benjamini [8] showed that in the competition of two growing clusters in a  $\delta$ -hyperbolic graph, one growing faster than the other, the faster cluster not necessarily surround the slower cluster).

**Proposition 1.** *Given  $r > 2\delta \geq 0$ , any  $\delta$ -hyperbolic graph  $G = (V, E)$  is  $(2r, r + 2\delta)$ -dismantlable and therefore  $G \in \mathcal{CWF}\mathcal{R}(2r, r + 2\delta)$ .*

*Proof.* The second assertion follows from Theorem 1. To prove the  $(2r, r + 2\delta)$ -dismantlability of  $G$ , we will employ Lemma 2 of [16]. According to this result, in a  $\delta$ -hyperbolic graph  $G$  for any subset of vertices  $X$  there exist two vertices  $x \in X$  and  $c \in V$  such that  $d(c, y) \leq r + 2\delta$  for any vertex  $y \in X \cap N_{2r}(x)$ , i.e.,  $N_{2r}(x) \cap X \subseteq N_{r+2\delta}(c)$ . The proof of [16] shows that the vertices  $x$  and  $c$  can be selected in the following way: pick any vertex  $z$  of  $G$  as a basepoint, construct a breadth-first search tree  $T$  of  $G$  rooted at  $z$ , and then pick  $x$  to be the furthest from  $z$  vertex of  $X$  and  $c$  to be vertex located at distance  $s$  on the unique path between  $x$  and  $z$  in  $T$ . Using this result, we will establish a slightly stronger version of dismantlability of a  $\delta$ -hyperbolic graph  $G$ , in which the inclusion  $N_s(v_i, G \setminus \{v_j\}) \cap X_i \subseteq N_{s'}(v_j)$  is replaced by  $N_s(v_i) \cap X_i \subseteq N_{s'}(v_j)$  with  $s := 2r$  and  $s' := r + 2\delta$ . We recursively construct the ordering of  $V$ . By previous result, there exist two vertices  $v_1 \in X_1 := V$  and  $c \in X_2 := V \setminus \{x_1\}$  such that  $N_{2r}(v_1) \cap X_1 \subseteq N_{r+2\delta}(c)$ . At step  $i \geq 1$ , suppose by induction hypothesis that  $V$  is the disjoint union of the sets  $\{v_1, \dots, v_i\}$  and  $X_{i+1}$ , so that, for any  $j \leq i$ , there exists a vertex  $c \in X_{j+1}$  such that  $N_{2r}(v_j) \cap X_j \subseteq N_{r+2\delta}(c)$  with  $X_j = \{v_j, \dots, v_i\} \cup X_{i+1}$ . We assert that this ordering can be extended. Applying the previous result to the set  $X := X_{i+1}$  we can define two vertices  $v_{i+1} \in X_{i+1}$  and  $c \neq v_{i+1}$  such that  $N_{2r}(v_{i+1}) \cap X_{i+1} \subseteq N_{r+2\delta}(c)$ .

The choice of the vertices  $x \in X$  and  $c \in V$  provided by [16] and the definition of the sets  $X_1, X_2, \dots$  ensure that if a vertex of  $G$  is closer to the root than another vertex, then the first vertex will be labeled later than the second one. Since by construction  $c$  is closer to  $z$  than  $v_{i+1}$ , necessarily  $c$  belongs to the set  $X_{i+1} \setminus \{v_{i+1}\}$ .  $\square$

In general, dismantlable graphs do not have bounded hyperbolicity because they are universal in the following sense. As we noticed in the introduction, any finite Helly graph is dismantlable. On the other hand, it is well known that an arbitrary connected graph can be isometrically embedded into a Helly graph (see for example [6, 31]). However, dismantlable graphs without some short induced cycles are 1-hyperbolic:

**Corollary 2.** *Any dismantlable graph  $G = (V, E)$  without induced 4-, 5-, and 6-cycles is 1-hyperbolic, and therefore  $G \in \mathcal{CWFR}(2r, r + 2)$  for any  $r > 0$ .*

*Proof.* A dismantlable graph  $G$  not containing induced 4- and 5-cycles does not contain 4-wheels and 5-wheels as well (a  $k$ -wheel is a cycle of length  $k$  plus a vertex adjacent to all vertices of this cycle), therefore  $G$  is bridged by a result of [1]. Since  $G$  does not contain 6-wheels as well,  $G$  is 1-hyperbolic by Proposition 11 of [16]. Then the second assertion immediately follows from Proposition 1.  $\square$

**Open question 1:** Is it true that the converse of Proposition 1 holds? More precisely, is it true that if  $G \in \mathcal{CWFR}(s, s')$  for  $s' < s$ , then the graph  $G$  is  $\delta$ -hyperbolic, where  $\delta$  depends only of  $s$ ?

We give some confidences in the truth of this conjecture by showing that for  $s \geq 2s'$  all graphs  $G \in \mathcal{CWFR}(s, s')$  are  $(s - 1)$ -hyperbolic. On the other hand, since  $\mathcal{CWFR}(s, s') \subset \mathcal{CWFR}(s, s' + 1)$ , to answer our question for  $s' < s < 2s'$  it suffices to show its truth for the particular case  $s' = s - 1$ . We give a positive answer to our question for Helly and bridged graphs by showing that if such a graph  $G$  belongs to the class  $\mathcal{CWFR}(s, s - 1)$ , then  $G$  is  $s^2$ -hyperbolic.

In the following results, for an  $(s, s')$ -dismantling order  $v_1, \dots, v_n$  of a graph  $G \in \mathcal{CWFR}(s, s')$  and a vertex  $v$  of  $G$ , we will denote by  $\alpha(v)$  the rank of  $v$  in this order (i.e.,  $\alpha(v) = i$  if  $v = v_i$ ). For two vertices  $u, v$  with  $\alpha(u) < \alpha(v)$  and a shortest  $(u, v)$ -path  $P(u, v)$ , an  $s$ -net  $N(u, v)$  of  $P(u, v)$  is an ordered subset of vertices of  $P(u, v)$ ,  $(u = x_0, x_1, \dots, x_k, x_{k+1} = v)$ , such that  $d(x_i, x_{i+1}) = s$  for any  $i = 0, \dots, k - 1$  and  $0 < d(x_k, x_{k+1}) \leq s$ .

**Proposition 2.** *If  $G \in \mathcal{CWFR}(s, s - 1)$  and  $u, v$  are two vertices of  $G$  such that  $\alpha(u) < \alpha(v)$  and  $d(u, v) > s^2$ , then for any shortest  $(u, v)$ -path  $P(u, v)$ , the vertex  $x_1$  of its  $s$ -net  $N(u, v) = (u = x_0, x_1, \dots, x_k, x_{k+1} = v)$  satisfies the condition  $\alpha(u) < \alpha(x_1)$ .*

*Proof.* Suppose by way of contradiction that  $\alpha(u) > \alpha(x_1)$ . Let  $x_i$  ( $1 \leq i \leq k$ ) be a vertex of  $N(u, v)$  having a locally minimal index  $\alpha(x_i)$ , i.e.,  $\alpha(x_{i-1}) > \alpha(x_i) < \alpha(x_{i+1})$ . Let  $y_i$  be the vertex eliminating  $x_i$  in the  $(s, s - 1)$ -dominating order. We assert that  $d(y_i, x_{i-1}) \leq s - 1$  and  $d(y_i, x_{i+1}) \leq s - 1$ . Indeed, if  $y_i$  does not belong to the portion of the path  $P(u, v)$  comprised between  $x_{i-1}$  and  $x_{i+1}$ , then  $x_{i-1}, x_{i+1} \in X_{\alpha(x_i)} \cap N_s(x_i, G \setminus \{y_i\})$ , and therefore  $x_{i-1}, x_{i+1} \in N_{s-1}(y_i)$  by the dismantling condition. Now suppose that  $y_i$  belongs to one of

the segments of  $P(u, v)$ , say to the subpath between  $x_{i-1}, x_i$ . Since  $y_i \neq x_i$  we conclude that  $d(x_{i-1}, y_i) \leq s - 1$ . On the other hand, since  $x_{i+1} \in X_{\alpha(x_i)} \cap N_s(x_i, G \setminus \{y_i\})$ , by dismantling condition we conclude that  $d(y_i, x_{i+1}) \leq s - 1$ . Hence, indeed  $d(y_i, x_{i-1}) \leq s - 1, d(y_i, x_{i+1}) \leq s - 1$ , whence  $d(x_{i-1}, x_{i+1}) \leq 2s - 2$ . Since  $d(x_{i-1}, x_{i+1}) = 2s$  for any  $1 \leq i \leq k - 1$ , we conclude that  $i = k$ . Therefore the indices of the vertices of  $N(u, v)$  satisfy the inequalities  $\alpha(u) = \alpha(x_0) > \dots > \alpha(x_{k-1}) > \alpha(x_k) < \alpha(x_{k+1}) = \alpha(v)$ .

Denote by  $N$  the ordered sequence of vertices  $x_0 = u, x_1, \dots, y_k, x_{k+1} = v$  obtained from the  $s$ -net  $N(u, v)$  by replacing the vertex  $x_k$  by  $y_k$ . We say that  $N$  is obtained from  $N(u, v)$  by an *exchange*. Call two consecutive vertices of  $N$  a *link*:  $N$  has  $k + 1$  links,  $k - 1$  links of length  $s$  and two links of length at most  $s - 1$ . If  $\alpha(y_k) < \alpha(x_{k-1})$ , then we perform with  $y_k$  the same exchange operation as we did with  $x_k$ . After several such exchanges, we will obtain a new ordered set  $x_0 = u, x_1, \dots, z_k, x_{k+1} = v$  (denote it also by  $N$ ) having  $k - 1$  links of length  $s$  and two links of length  $\leq s - 1$  and  $\alpha(x_{k-1}) < \alpha(z_k)$ . Since  $\alpha(x_{k-2}) > \alpha(x_{k-1})$ , using the  $(s, s - 1)$ -dismantling order we can exchange in  $N$  the vertex  $x_{k-1}$  by a vertex  $y_{k-1}$  to get an ordered set (denote it also by  $N$ ) having  $k - 3$  links of length  $s$  and 3 links of length  $s - 1$ . Repeating the exchange operation with each occurring local minimum (different from  $u$ ) of  $N$  with respect to the total order  $\alpha$ , after a finite number of exchanges we will obtain an ordered set  $N = (u, z_1, z_2, \dots, z_k, v)$  consisting of  $k + 1$  links of length at most  $s - 1$  each and such that  $\alpha(u) < \alpha(z_i)$  for any  $i = 1, \dots, k$ . By triangle inequality,  $d(u, v) \leq d(u, z_1) + d(z_1, z_2) + \dots + d(z_k, v) \leq (k + 1)(s - 1)$ . On the other hand, from the definition of  $N(u, v)$  we conclude that  $d(u, v) = ks + \gamma$ , where  $0 < \gamma = d(x_k, v) \leq s$ . Hence  $(k + 1)(s - 1) \geq ks + \gamma$ , yielding  $k \leq s - \gamma - 1$ . But then  $d(u, v) = ks + \gamma \leq (s - \gamma - 1)s + \gamma = s^2 - s\gamma - s + \gamma < s^2$ , contrary to the assumption that  $d(u, v) \geq s^2$ . This contradiction shows that indeed  $\alpha(x_1) > \alpha(u)$ .  $\square$

We call a graph  $G \in \mathcal{CWF}\mathcal{R}(s, s - 1)$   $(s, s - 1)$ -*dismantlable* if for any  $(s, s - 1)$ -dismantling order  $v_1, \dots, v_n$  of  $G$ , for each vertex  $v_i, 1 \leq i < n$ , there exists another vertex  $v_j$  adjacent to  $v_i$  such that  $N_s(v_i, G \setminus \{v_j\}) \cap X_i \subseteq N_{s-1}(v_j)$ , where  $X_i := \{v_i, v_{i+1}, \dots, v_n\}$  and  $X_n = \{v_n\}$ . The difference between  $(s, s - 1)$ -dismantlability and  $(s, s - 1)$ -*dismantlability* is that in the second case the vertex  $v_j$  dominating  $v_i$  is necessarily adjacent to  $v_i$  but not necessarily eliminated after  $v_i$ .

**Proposition 3.** *If a graph  $G \in \mathcal{CWF}\mathcal{R}(s, s - 1)$  is  $(s, s - 1)$ -dismantlable, then  $G$  is  $s^2$ -hyperbolic.*

*Proof.* Pick any quadruplet of vertices  $u, v, x, y$  of  $G$ , consider its representation as in Fig. 1 where  $\xi \leq \eta$ , and proceed by induction on the total distance sum  $S(u, v, x, y) = d(u, v) + d(u, x) + d(u, y) + d(v, x) + d(v, y) + d(x, y)$ . From Fig. 1 we immediately conclude that if one of the distance between the vertices  $u, v, x, y$  is at most  $s^2$ , then  $\xi \leq s^2$  and we are done. So suppose that the distance between any two vertices of our quadruplet is at least  $s^2$ .

Consider any  $(s, s - 1)$ -dismantling order  $v_1, \dots, v_n$  of  $G$  and suppose that  $u$  is the vertex of our quadruplet occurring first in this order. Pick three shortest paths  $P(u, v), P(u, x)$ , and  $P(u, y)$  between the vertex  $u$  and the three other vertices of the quadruplet. Denote by  $v_1, x_1$ , and  $y_1$  the vertices of the paths  $P(u, v), P(u, x)$ , and  $P(u, y)$ , respectively, located at distance



$s$  from  $u$ . From Proposition 2 we infer that  $u$  is eliminated before each of the vertices  $v_1, x_1, y_1$ . Let  $u'$  be the neighbor of  $u$  eliminating  $u$  in the  $(s, s-1)^*$ -dismantling order associated with the  $(s, s-1)$ -dismantling order  $v_1, \dots, v_n$ . From the  $(s, s-1)^*$ -dismantling condition we infer that each of the distances  $d(u', v_1), d(u', x_1), d(u', y_1)$  is at most  $s-1$ . Since  $u$  is adjacent to  $u'$  and  $u$  is at distance  $s$  from  $v_1, x_1, y_1$ , necessarily  $d(u', v_1), d(u', x_1), d(u', y_1)$  are all equal to  $s-1$ . Therefore, if we will replace in our quadruplet the vertex  $u$  by  $u'$ , we will obtain a quadruplet with a smaller total distance sum:  $S(u', v, x, y) = S(u, v, x, y) - 3$ . Therefore, by induction hypothesis, the two largest of the distance sums  $d(u', v) + d(x, y)$ ,  $d(u', x) + d(v, y)$ ,  $d(u', y) + d(v, x)$  differ by at most  $2s^2$ . On the other hand,  $d(u, v) + d(x, y) = d(u', v) + d(x, y) + 1$ ,  $d(u, x) + d(v, y) = d(u', x) + d(v, y) + 1$ , and  $d(u, y) + d(v, x) = d(u', y) + d(v, x) + 1$ , whence the two largest distance sums of the quadruplet  $u, v, x, y$  also differ by at most  $2s^2$ . Hence  $G$  is  $s^2$ -hyperbolic.  $\square$

A graph  $G$  is called a *Helly graph* if its family of balls satisfies the Helly property: any collection of pairwise intersecting balls has a common vertex. A graph  $G$  is called a *bridged graph* if all isometric cycles of  $G$  have length 3. Equivalently,  $G$  is a bridged graph if all balls around convex sets are convex (a subset  $S$  of vertices is convex if together with any two vertices  $u, v$ , the set  $S$  contains the *interval*  $I(u, v) = \{x \in V : d(u, v) = d(u, x) + d(x, v)\}$  between  $u$  and  $v$ ). For a comprehensive survey of results and bibliography on Helly and bridged graphs, see [6].

**Proposition 4.** *If  $G \in \mathcal{CWFR}(s, s-1)$  is a Helly or a bridged graph, then  $G$  is  $(s, s-1)^*$ -dismantlable and therefore  $G$  is  $s^2$ -hyperbolic.*

*Proof.* First, let  $G$  be an  $(s, s-1)$ -dismantlable Helly graph. Let  $v_i$  be the  $i$ th vertex in an  $(s, s-1)$ -dismantling order and let  $y_i$  be the vertex eliminating  $v_i$ . Suppose that  $k := d(v_i, y_i) \geq 2$ . We assert that we can always eliminate  $v_i$  with a vertex  $y'_i$  adjacent to  $y_i$  and located at distance  $k-1$  from  $v_i$ . Then repeating the same reasoning with  $y'_i$  instead of  $y_i$ , we will eventually arrive at a vertex of  $I(v_i, y_i)$  adjacent to  $v_i$  which still eliminate  $v_i$ . Set  $A := (X_i \cap N_s(v_i)) \setminus \{v_i, y_i\}$ . For each vertex  $x \in A$ , consider the ball  $N_{s-1}(x)$  of radius  $s-1$  centered at  $x$ . Consider also the balls  $N_{k-1}(v_i)$  and  $N_1(y_i)$ . We assert that the balls of the resulting collection pairwise intersect. Indeed, any two balls centered at vertices of  $A$  intersect in  $y_i$ . The ball  $N_1(y_i)$  intersects any ball centered at  $A$  in  $y_i$ . The ball  $N_{k-1}(v_i)$  intersects any ball centered at a vertex  $x \in A$  because  $d(v_i, x) \leq s \leq k-1 + s-1$ . Finally,  $N_{k-1}(v_i)$  and  $N_1(y_i)$  intersect because  $d(v_i, y_i) = k = k-1 + 1$ . By Helly property, the balls of this collection intersect in a vertex  $y'_i$ . Since  $y'_i$  is at distance at most  $k-1$  from  $v_i$  and at distance at most 1 from  $y_i$ , from the equality  $d(v_i, y_i) = k$  we immediately deduce that  $y'_i$  is a neighbor of  $y_i$  located at distance  $k-1$  from  $v_i$ . This establishes the  $(s, s-1)^*$ -dismantling property for Helly graphs.

Now, suppose that  $G$  is a bridged graph and let the vertices  $v_i, y_i$  and the set  $A$  be defined as in previous case. Since  $G$  is bridged, the convexity of the ball  $N_{k-1}(v_i)$  implies that the set  $C$  of neighbors of  $y_i$  in the interval  $I(v_i, y_i)$  induces a complete subgraph. Pick any vertex  $x \in A$ . Clearly,  $d(x, y_i) \leq s-1$  and  $d(x, v_i) \leq s$ . If  $d(x, v_i) \leq s-1$ , then  $v_i, y_i \in N_{s-1}(x)$  and from the convexity of the ball  $N_{s-1}(x)$  we conclude that  $I(v_i, y_i) \subset N_{s-1}(x)$ . Hence, in this

case,  $d(x, y) \leq s - 1$  for any  $y \in I(v_i, y_i)$ , in particular, for any vertex of  $C$ . Analogously, if  $d(x, y_i) < s - 1$ , then  $d(x, y) \leq s - 1$  for any vertex  $y \in C$ . Therefore the choice of the vertex  $y'_i$  in  $C$  depends only of the vertices of the set  $A_0 = \{x \in A : d(x, v_i) = s \text{ and } d(x, y_i) = s - 1\}$ . Pick any vertex  $x \in A_0$ . If  $I(x, y_i) \cap I(y_i, v_i) \neq \{y_i\}$ , then  $y_i$  has a neighbor  $y'$  in this intersection located at distance  $s - 2$  from  $x$ . Since  $y' \in C$  and  $C$  is a complete subgraph, then  $d(y, x) \leq s - 1$  for any  $y \in C$ . Therefore we can discard all such vertices of  $A_0$  from our future analysis and suppose without loss of generality that  $I(x, y_i) \cap I(y_i, v_i) = \{y_i\}$  for any  $x \in A_0$ . For  $x \in A_0$ , let  $x_0$  be a furthest from  $x$  vertex of  $I(x, y_i) \cap I(x, v_i)$ . Let  $v_0$  be a furthest from  $v_i$  vertex of  $I(v_i, x_0) \cap I(v_i, y_i)$ . Since  $I(x, y_i) \cap I(y_i, v_i) = \{y_i\}$  and  $G$  is bridged, the vertices  $y_i, x_0, v_0$  define an equilateral metric triangle sensu [5, 6]:  $d(y_i, x_0) = d(x_0, v_0) = d(v_0, y_i) =: m$ . Moreover, any vertex of  $I(v_0, y_i)$  is located at distance  $m$  from  $x_0$  and therefore at distance  $s - 1$  from  $x$ . From the definition of  $x_0$  and  $v_0$  we conclude that  $m + d(x_0, x) = s - 1$ ,  $d(x, x_0) + m + d(v_0, v_i) = s$ , and  $d(v_i, v_0) + m \leq s - 1$ . Whence  $d(v_i, v_0) = 1$ , yielding  $d(v_i, y_i) = m$ .

Pick in  $C$  a vertex  $y$  belonging to a maximum number of balls  $N_{s-1}(x)$  centered at  $x \in A_0$ . Suppose by way of contradiction that  $A_0$  contains a vertex  $x'$  such that  $y \notin N_{s-1}(x')$ . Since  $d(x', y_i) = s - 1$  and  $y$  is adjacent to  $y_i$ , we have  $d(x', y) = s$ . Let  $y'$  be a vertex of  $C$  belonging to  $N_{s-1}(x')$ . Let  $v'_0$  be the neighbor of  $v_i$  defined with respect to  $x'$  in the same way as  $v_0$  was defined for  $x$ . Then all vertices of  $I(v'_0, y')$  are located at distance  $s - 1$  from  $x'$ . We can suppose that there exists a vertex  $x \in A_0$  such that  $y \in N_{s-1}(x)$  but  $y' \notin N_{s-1}(x)$ , otherwise we will obtain a contradiction with the choice of  $y$ . Since the balls  $N_{s-1}(x)$  and  $N_{s-1}(x')$  are convex, the intervals  $I(v_0, y_i)$  and  $I(v'_0, y_i)$  belong to these balls, respectively, whence  $d(v_0, y) = d(v'_0, y') = m - 1$  but  $d(v_0, y') = d(v'_0, y) = m$ . Let  $z$  be a neighbor of  $y$  in  $I(v_0, y)$ . Since  $z, y' \in I(y, v'_0)$  and  $G$  is bridged, the vertices  $z$  and  $y'$  are adjacent. Now, since  $y' \notin N_{s-1}(x)$ , we conclude that  $d(x, y') = s$ . On the other hand,  $d(z, x) = d(y_i, x) = s - 1$  because  $z \in I(v_0, y_i) \subset N_{s-1}(x)$ . Hence  $z, y_i \in I(y', x)$ , and therefore  $y_i$  and  $z$  must be adjacent because the ball  $N_{s-1}(x)$  is convex. But this is impossible because  $z$  is a neighbor of  $y$  in  $I(y, v_0) \subset I(y_i, v_0)$ . This contradiction shows that  $C$  contains a vertex belonging to all balls  $N_{s-1}(x)$  centered at vertices of  $A_0$ , thus establishing the  $(s, s - 1)^*$ -dismantling property for bridged graphs.  $\square$

**Proposition 5.** *If  $s \geq 2s'$ , then any graph  $G$  of  $CWFR(s, s')$  is  $(s - 1)$ -hyperbolic.*

*Proof.* First, similarly to Proposition 2, we prove that if  $d(u, v) \geq s$  and  $\alpha(u) < \alpha(v)$ , then the vertex  $x_1$  of the  $s$ -net  $N(u, v)$  of any shortest  $(u, v)$ -path satisfies the inequality  $\alpha(x_1) > \alpha(u)$ . Suppose by way of contradiction that  $\alpha(u) > \alpha(x_1)$ . Then as in proof of Proposition 2 we conclude that  $x_k$  is the unique local minimum of  $\alpha$  on  $N(u, v) : \alpha(x_{k-1}) > \alpha(x_k) < \alpha(x_{k+1})$ . Let  $y_k$  be the vertex eliminating  $x_k$  in the  $(s, s')$ -dominating order. If  $y_k$  does not belong to the segment of  $P(u, v)$  between  $x_{k-1}$  and  $x_k$ , then  $d(x_{k-1}, x_{k+1}) \leq d(x_{k-1}, y_k) + d(y_k, x_{k+1}) \leq 2s'$ , contrary to the assumption that  $d(x_{k-1}, x_{k+1}) > s \geq 2s'$ . So  $y_k$  belongs to the subpath of  $P(u, v)$  between  $x_{k-1}$  and  $x_{k+1}$ . If  $y_k$  belongs to the subpath comprised between  $x_k$  and  $x_{k+1}$ , then the dismantling condition implies that  $d(y_k, x_{k-1}) \leq s'$ , which is impossible because  $d(y_k, x_{k-1}) = d(y_k, x_k) + s > 2s'$ . The same contradiction



is obtained if  $y_k$  belongs to the second half of the subpath between  $x_{k-1}$  and  $x_k$ . Finally, if  $y_k$  belong to the first half of this subpath, then  $d(y_k, x_{k+1}) \leq s'$  by the dismantling condition, contradicting the fact that the location of  $y_k$  on this subpath of  $P(u, v)$  implies that  $d(y_k, x_{k+1}) > s'$ . This shows that indeed  $\alpha(x_1) > \alpha(u)$ .

To establish  $(s - 1)$ -hyperbolicity of  $G$ , as in the proof of Proposition 3 we pick any quadruplet of vertices  $u, v, x, y$  of  $G$  and proceed by induction on the total distance sum  $S(u, v, x, y) = d(u, v) + d(u, x) + d(u, y) + d(v, x) + d(v, y) + d(x, y)$ . Again, we can suppose that the distances between any two vertices of this quadruplet is at least  $s$ , otherwise we are done. Let  $u$  be the vertex of our quadruplet occurring first in some  $(s, s')$ -dismantling order of  $G$ . Pick three shortest paths  $P(u, v)$ ,  $P(u, x)$ , and  $P(u, y)$  and denote by  $v_1, x_1$ , and  $y_1$  their respective vertices located at distance  $s$  from  $u$ . From first part of our proof we infer that  $u$  is eliminated before  $v_1, x_1$ , and  $y_1$ . Let  $u'$  be the vertex eliminating  $u$ . From the  $(s, s')$ -dismantling condition we infer that  $d(u, u') \leq s'$ . Moreover, either  $d(u', v_1) \leq s'$  or  $v_1 \notin N_s(u, G \setminus \{u'\})$ . Since  $d(u, v_1) = s \geq 2s'$ , in both cases we conclude that  $u'$  belongs to a shortest  $(u, v_1)$ -path of  $G$ . Analogously, we conclude that  $u'$  lie on a shortest  $(u, x_1)$ -path and on a shortest  $(u, y_1)$ -path. Therefore, if we replace in our quadruplet  $u$  by  $u'$ , we will get a quadruplet with total distance sum  $S(u', v, x, y) = S(u, v, x, y) - 3d(u, u') < S(u, v, x, y)$ . By induction hypothesis, the two largest distance sums of this quadruplet differ by at most  $2(s - 1)$ . On the other hand, since  $d(u, v) + d(x, y) = d(u', v) + d(x, y) + d(u, u')$ ,  $d(u, x) + d(v, y) = d(u', x) + d(v, y) + d(u, u')$ , and  $d(u, y) + d(v, x) = d(u', y) + d(v, x) + d(u, u')$ , the two largest distance sums of the quadruplet  $u, v, x, y$  also differ by at most  $2(s - 1)$ . Hence  $G$  is  $(s - 1)$ -hyperbolic.  $\square$

### 3 Cop-win graphs for game with fast robber: class $\mathcal{CWFR}(s)$

In this section, we specify the dismantling scheme provided by Theorem 1 in order to characterize the graphs in which one cop with speed 1 captures a robber with speed  $s \geq 2$ . First we show that the graphs from  $\mathcal{CWFR}(2)$  are precisely the dually chordal graphs [12]. Then we show that for  $s \geq 3$  the classes  $\mathcal{CWFR}(s)$  coincide with  $\mathcal{CWFR}(\infty)$  and we provide a structural characterization of these graphs.

#### 3.1 $\mathcal{CWFR}(2)$ and dually chordal graphs

We start by showing that when the cop has speed 1 and the robber has speed  $s \geq 1$ , then the dismantling order in Theorem 1 can be defined using the subgraphs  $G_i = G(X_i)$ .

**Proposition 6.** *A graph  $G$  is  $(s, 1)$ -dismantlable if and only if the vertices of  $G$  can be ordered  $v_1, \dots, v_n$  in such a way that for each vertex  $v_i \neq v_n$  there exists a vertex  $v_j$  with  $j > i$  such that  $N_s(v_i, G_i \setminus \{v_j\}) \subseteq N_1(v_j, G_i)$ .*

*Proof.* First, note that for any  $i \leq j$ ,  $N_1(v_j, G) \cap X_i = N_1(v_j, G_i)$ . Thus, if a graph  $G$  is  $(s, 1)$ -dismantlable, then any  $(s, 1)$ -dismantling order satisfies the requirement  $N_s(v_i, G_i \setminus \{v_j\}) \subseteq N_1(v_j, G_i)$ . Conversely, consider an order  $v_1, \dots, v_n$  on the vertices of  $G$  satisfying this

condition. If  $s = 1$ , then  $N_1(v_i, G_i \setminus \{v_j\}) = N_1(v_i, G \setminus \{v_j\}) \cap X_i$  and thus our assertion is obviously true. We now suppose that  $s \geq 2$ . By induction on  $i$ , we will show that  $N_s(v_i, G \setminus \{v_j\}) \cap X_i \subseteq N_1(v_j)$ . For  $i = 1$ ,  $G_i = G$  and thus the property holds. Consider  $i$  such that for any  $i' < i$ , the property is satisfied. Pick any vertex  $u \in N_s(v_i) \cap X_i$ . If the distance in  $G_i \setminus \{v_j\}$  between  $v_i$  and  $u$  is at most  $s$ , then  $u \in N_s(v_i, G_i \setminus \{v_j\}) \subseteq N_1(v_j)$  and we are done. Otherwise, we can find a unique index  $i_0 < i$  such that the distance between  $v_i$  and  $u$  in the graph  $G_{i_0} \setminus \{v_j\}$  is at most  $s$  and in the graph  $G_{i_0+1} \setminus \{v_j\}$  is larger than  $s$ . Consider a shortest path  $\pi$  between  $v_i$  and  $u$  in  $G_{i_0} \setminus \{v_j\}$ . From the choice of  $i_0$ , necessarily  $v_{i_0}$  is a vertex of  $\pi$ . Since the length of  $\pi$  is at most  $s$ , we deduce that  $d_{G_{i_0}}(u, v_{i_0}) \leq s$  and  $d_{G_{i_0}}(v_i, v_{i_0}) \leq s$ . By the induction hypothesis, there exists  $j_0 > i_0$  such that  $N_s(v_{i_0}, G_{i_0} \setminus \{v_{j_0}\}) \subseteq N_1(v_{j_0})$ . If  $j_0 \neq j$ , then there exists a path  $(u, v_{j_0}, v_i)$  of length 2 between  $u$  and  $v_i$  in  $G_{j_0}$ . Since  $j_0 > i_0$ , we obtain a contradiction with the definition of  $i_0$ . Hence  $j_0 = j$ , and, by our induction hypothesis,  $u \in N_s(v_{i_0}, G_{i_0} \setminus \{v_j\}) \subseteq N_1(v_j)$ , and we are done.  $\square$

Analogously to Theorem 3 of Clarke [18] for the witness version of the game, it can be easily shown that, for any  $s$ , the class  $\mathcal{CWFR}(s)$  is closed under retracts:

**Proposition 7.** *If  $G \in \mathcal{CWFR}(s)$  and  $G'$  is a retract of  $G$ , then  $G' \in \mathcal{CWFR}(s)$ .*

Recall that a graph  $G$  is called *dually chordal* [12] if its clique hypergraph (or, equivalently, its ball hypergraph) is a hypertree, i.e., it satisfies the Helly property and its line graph is chordal (see the Berge's book on hypergraphs [10] for these two definitions). Dually chordal graphs are equivalently defined as the graphs  $G$  having a spanning tree  $T$  such that any maximal clique or any ball of  $G$  induces a subtree of  $T$ . Finally, dually chordal graphs are exactly the graphs  $G = (V, E)$  admitting a maximum neighborhood ordering of its vertices. A vertex  $u \in N_1(v)$  is a *maximum neighbor* of  $v$  if for all  $w \in N_1(v)$  the inclusion  $N_1(w) \subseteq N_1(u)$  holds. The ordering  $\{v_1, \dots, v_n\}$  is a *maximum neighborhood ordering* (mno for short) of  $G$  [12], if for all  $i < n$ , the vertex  $v_i$  has a maximum neighbor in the subgraph  $G_i$  induced by the vertices  $X_i = \{v_i, v_{i+1}, \dots, v_n\}$ . Dually chordal graphs comprise strongly chordal graphs, doubly chordal, and interval graphs as subclasses and can be recognized in linear time. Any graph  $H$  can be transformed into a dually chordal graph by adding a new vertex  $c$  adjacent to all vertices of  $H$ .

**Theorem 2.** *For a graph  $G = (V, E)$ , the following conditions are equivalent:*

- (i)  $G \in \mathcal{CWFR}(2)$ ;
- (ii)  $G$  is  $(2, 1)$ -dismantlable;
- (iii)  $G$  admits an mno ordering;
- (iv)  $G$  is dually chordal.

*Proof.* Since  $\mathcal{CWFR}(2) = \mathcal{CWFR}(2, 1)$ , the equivalence (i) $\Leftrightarrow$ (ii) follows from Theorem 1. The equivalence (iii) $\Leftrightarrow$ (iv) is a result of [12]. Notice that  $u$  is a maximum neighbor of  $v$  in

$G$  iff  $N_2(v) = N_1(u)$ . Therefore,  $\{v_1, \dots, v_n\}$  is a maximum neighborhood ordering of  $G$  iff for all  $i < n$ ,  $N_2(v_i, G_i) = N_1(v_j, G_i)$  for some  $v_j, j > i$ . Hence any mno ordering is a  $(2, 1)$ -dismantling ordering, establishing (iii) $\Rightarrow$ (ii). Finally, by induction on the number of vertices of  $G$  we will show that any  $(2, 1)$ -dismantling ordering  $\{v_1, \dots, v_n\}$  of the vertex set of  $G$  is an mno, thus (ii) $\Rightarrow$ (iii). Suppose that  $N_2(v_1, G \setminus \{u\}) \subset N_1(u)$  for some  $u := v_j, j > 1$ . Then  $u$  is adjacent to  $v_1$  and to all neighbors of  $v_1$ . Since for any neighbor  $w \neq u$  of  $v_1$  the ball  $N_1(w)$  is contained in the punctured ball  $N_2(v_1, G \setminus \{u\})$ , we conclude that  $N_1(w) \subseteq N_1(u)$ , i.e.,  $u$  is a maximum neighbor of  $v_1$ . The graph  $G'$  obtained from  $G$  by removing the vertex  $v_1$  is a retract, and therefore an isometric subgraph of  $G$ . Thus for any vertex  $v_i, i > 1$ , by what has been noticed above, the intersection of a ball (or of a punctured ball) of  $G$  centered at  $v_i$  with the set  $X_2 = \{v_2, \dots, v_n\}$  coincides with the corresponding ball (or punctured ball) of the graph  $G' = G(X_2)$  centered at the same vertex  $v_i$ . Therefore  $\{v_2, \dots, v_n\}$  is a  $(2, 1)$ -dismantling ordering of the graph  $G'$ . By induction assumption,  $\{v_2, \dots, v_n\}$  is an mno of  $G'$ . Since  $v_1$  has a maximum neighbor in  $\{v_2, \dots, v_n\}$ , we conclude that  $\{v_1, v_2, \dots, v_n\}$  is a maximum neighborhood ordering of  $G$ .  $\square$

### 3.2 $CWFR(k), k \geq 3$ , and big brother graphs

A *block* of a graph  $G$  is a maximal by inclusion vertex two-connected subgraph of  $G$  (possibly reduced to a single edge). Two blocks of  $G$  are either disjoint or share a single vertex, called an *articulation point*. Any graph  $G = (V, E)$  admits a block-decomposition in the form of a rooted tree  $T$ : each vertex of  $T$  is a block of  $G$ , pick any block  $B_1$  as a root of  $T$ , label it, and make it adjacent in  $T$  to all blocks intersecting it, then label that blocks and make them adjacent to all nonlabeled blocks which intersect them, etc. A block  $B$  of  $G$  is *dominated* if it contains a vertex  $u$  (called the *big brother* of  $B$ ) which is adjacent to all vertices of  $B$ . A graph  $G$  is a *big brother graph*, if its block-decomposition can be represented in the form of a rooted tree  $T$  is such a way that (1) each block of  $G$  is dominated and (2) for each block  $B$  distinct from the root  $B_1$ , the articulation point between  $B$  and its father-block dominates  $B$ . Equivalently,  $G$  is a big brother graph if its blocks can be ordered  $B_1, \dots, B_r$  such that  $B_1$  is dominated and, for any  $i > 1$ , the block  $B_i$  is a leaf in the block-decomposition of  $\cup_{j < i} B_j$  and is dominated by the articulation point connecting  $B_i$  to  $\cup_{j < i} B_j$  (we will call such a decomposition a *bb-decomposition* of  $G$ ); see Fig. 2(a) for an example.

**Theorem 3.** *For a graph  $G = (V, E)$  the following conditions are equivalent:*

- (i)  $G \in CWFR(3)$ ;
- (i')  $G$  is  $(3, 1)$ -dismantlable;
- (ii)  $G \in CWFR(\infty)$ ;
- (ii')  $G$  is  $(\infty, 1)$ -dismantlable;
- (iii)  $G$  is a big brother graph.

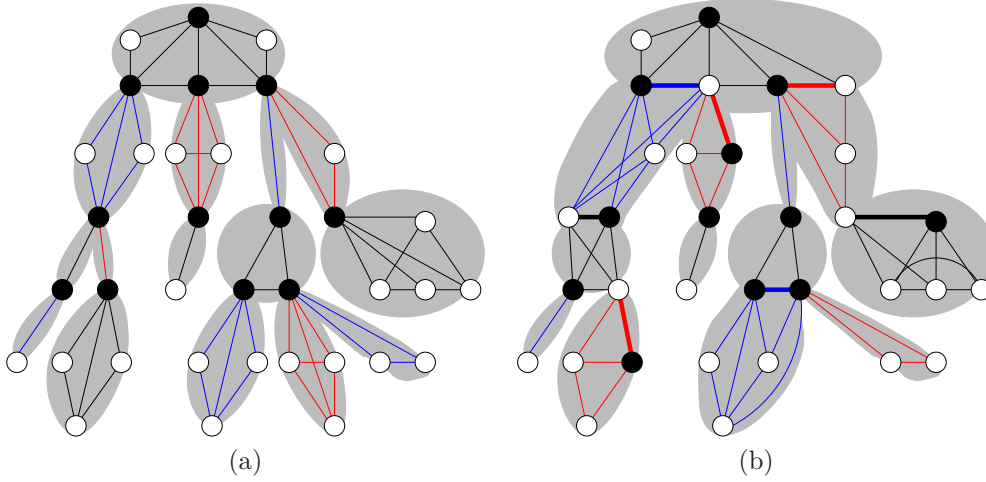


Figure 2: (a) A big brother graph. (b) A big two-brother graph.

In particular, the classes of graphs  $\mathcal{CWFR}(s)$ ,  $s \geq 3$ , coincide.

*Proof.* The equivalences (i) $\Leftrightarrow$ (i') and (ii) $\Leftrightarrow$ (ii') are particular cases of Theorem 1. Next we will establish (iii) $\Rightarrow$ (i)&(ii), i.e., that any big brother graph  $G$  belongs to  $\mathcal{CWFR}(s)$  for all  $s \geq 3$ . Let  $B_1, \dots, B_r$  be a bb-decomposition of  $G$ . We consider the following strategy for the cop. At the beginning of the game, we locate the cop at the big brother of the root-block  $B_1$ . Now, at each subsequent step, the cop moves to the neighbor of his current position that is closest to the position of the robber. Notice the following invariant of the strategy: the position of the cop will always be at the articulation point of a block  $B$  on the path of  $T$  between the previous block hosting  $\mathcal{C}$  and the current block hosting  $\mathcal{R}$ . This means that, since  $\mathcal{R}$  cannot traverse this articulation point without being captured,  $\mathcal{R}$  is restricted to move only in the union of blocks in the subtree rooted at  $B$ . Now, if before the move of the cop,  $\mathcal{C}$  and  $\mathcal{R}$  occupy their positions in the same block, then  $\mathcal{C}$  captures  $\mathcal{R}$  at the next move. Otherwise, the next move will increase the distance in  $T$  between the root and the block hosting  $\mathcal{C}$ . Therefore after at most diameter of  $T$  rounds,  $\mathcal{R}$  and  $\mathcal{C}$  will be located in the same block, and thus the cop captures the robber at next move. This shows that (iii) $\Rightarrow$ (i)&(ii).

The remaining part of the proof is devoted to the implication (i)&(i') $\Rightarrow$ (iii). Let  $G$  be a graph of  $\mathcal{CWFR}(3)$ . Notice first that for any articulation point  $u$  of  $G$ , and any connected component  $C$  of  $G \setminus \{u\}$ , the graph induced by  $C \cup \{u\}$  also belongs to  $\mathcal{CWFR}(3)$ . Indeed, this follows by noticing that  $G(C \cup \{u\})$  is a retract of  $G$  (this retraction is obtained by mapping all vertices outside  $C$  to  $u$ ) and that  $\mathcal{CWFR}(3)$  is closed under retracts by Proposition 7. To prove that a graph  $G = (V, E) \in \mathcal{CWFR}(3)$  is a big brother graph, we will proceed by induction on the number of vertices of  $G$ . If  $G$  has one or two vertices, the result is obviously

true. For the inductive step, we distinguish two cases, depending if  $G$  is two-connected or not.

**Case 1:**  $G$  is not two-connected.

Since each block of  $G$  has strictly less vertices than  $G$ , by induction hypothesis each block is a big brother graph, i.e., it has a dominating vertex. First suppose that the block-decomposition of  $G$  has a leaf  $B$  such that the articulation point  $a$  of  $B$  separating  $B$  from the rest of  $G$  is a big brother of  $B$ . Let  $G'$  be the subgraph of  $G$  induced by all blocks of  $G$  except  $B$ , i.e.,  $G' = G(V \setminus (B \setminus \{a\}))$ . Since  $G' \in \mathcal{CWR}(3)$  by what has been shown above, from the induction hypothesis we infer that  $G'$  is a big brother graph. Consequently, there exists a bb-decomposition  $B_1, \dots, B_r$  of  $G'$ . Then,  $B_1, \dots, B_r, B$  is a bb-decomposition of  $G$  and thus,  $G$  is a big brother graph. Suppose now that for any leaf in the block-decomposition of  $G$ , the articulation point of the corresponding block does not dominate it. Pick two leaves  $B_1$  and  $B_2$  in the block-decomposition of  $G$  and consider their unique articulation points  $a_1$  and  $a_2$  ( $a_i$  disconnects  $B_i$  from the rest of  $G$ ). We claim that in this case, a robber that moves at speed 3 can always escape, which will contradict the assumption that  $G \in \mathcal{CWR}(3)$ . Let  $b_i$  be the dominating vertex of the block  $B_i$ ,  $i = 1, 2$  (by assumption,  $b_i \neq a_i$ ). Consider now a vertex  $c_i \in B_i \setminus \{b_i\}$  which can be connected with  $a_i$  by a 2-path  $(c_i, g_i, a_i)$  avoiding  $b_i$  (such a vertex exists because  $B_i$  is two-connected and, by assumption,  $a_i$  is not a dominating vertex of  $B_i$ ). Let  $\pi$  be a shortest path from  $a_1$  to  $a_2$  in  $G$  and let  $h_1$  and  $h_2$  be the neighbors in  $\pi$  of  $a_1$  and  $a_2$ , respectively. Note that  $h_i$  does not belong to  $B_i$ , thus  $a_i$  is the only neighbor of  $h_i$  in  $B_i$ . We now describe a strategy that enables the robber to escape. Initially, if the cop is not in  $B_1$ , then the robber starts in  $c_1$ ; otherwise, he starts in  $c_2$ . Then the robber stays in  $c_i$ , as long as the cop is at distance  $\geq 2$  from  $c_i$ . When the cop moves to a neighboring vertex of  $c_i$ , then the robber goes to  $h_i$  (either via the path  $(c_i, b_i, a_i, h_i)$  or via the path  $(c_i, g_i, a_i, h_i)$ ) and then, no matter how the cop moves, he goes to  $c_{3-i}$  using the shortest path  $\pi$ . Now notice that when  $\mathcal{R}$  is in  $h_i$ ,  $\mathcal{C}$  is in  $B_i \setminus \{a_i\}$  and thus he cannot capture the robber. When the robber is moving from  $h_i$  to  $c_{3-i}$ , he uses a shortest path  $\pi$  of  $G$ : the cop cannot capture him either because he is initially at distance 2 from the robber and he moves slower than the robber. Consequently, the cop cannot capture the robber, contrary with the assumption  $G \in \mathcal{CWR}(3)$ .

**Case 2:**  $G$  is two-connected.

We must show that  $G$  has a dominating vertex. Consider a  $(3, 1)$ -dismantling order  $v_1, \dots, v_n$  of the vertices of  $G$ . Let  $u$  be a vertex such that  $N_3(v_1, G \setminus \{u\}) \subseteq N_1(u)$ . Since  $u$  is a maximum neighbor of  $v_1$ , the isometric subgraph  $G' := G(V \setminus \{v_1\})$  of  $G$  also belongs to  $\mathcal{CWR}(3)$  because  $v_2, \dots, v_n$  is a  $(3, 1)$ -dismantling ordering of  $G'$ . By induction hypothesis,  $G'$  is a big brother graph. Again, we distinguish two subcases, depending on the two-connectivity of  $G'$ . First suppose that  $G'$  is two-connected. Since  $G'$  is a big brother graph, it contains a dominating vertex  $t$ . If  $t$  is adjacent to  $v_1$ , then  $t$  dominates  $G$  and we are done. Otherwise, consider a neighbor  $w \neq u$  of  $v_1$ . Any vertex  $x \neq u$  of  $G$  can be connected to  $v_1$  by the path  $(v_1, w, t, x)$  of length 3 avoiding  $u$ , thus  $x$  belongs to the punctured ball  $N_3(v_1, G \setminus \{u\})$ . As a consequence,  $x$  is a neighbor of  $u$ , thus  $u$  dominates  $G$ . Now suppose that  $G'$  is not two-connected. We assert that  $u$  is the only articulation point of  $G'$ . Assume by

way of contradiction that  $w \neq u$  is an articulation point of  $G'$  and let  $x$  and  $y$  be two vertices of  $G'$  such that all paths connecting  $x$  to  $y$  go through  $w$ . In  $G$ ,  $x$  and  $y$  can be connected by two vertex-disjoint paths  $\pi_1$  and  $\pi_2$ . Assume without loss of generality that  $w \notin \pi_1$ . Since  $\pi_1$  cannot be a path of  $G'$ , the vertex  $v_1$  belongs to  $\pi_1$ . Let  $\pi_1 = (x, x_1, \dots, x_k, v_1, y_l, \dots, y_1)$ . Since  $x_k, y_l \in N_1(v_1) \subseteq N_3(v_1, G \setminus \{u\}) \cup \{u\} \subseteq N_1(u)$ , necessarily  $x_k, y_l \in N_1(u)$ . If  $x_k = u$  or  $y_l = u$ , then  $(x, x_1, \dots, x_k, y_l, \dots, y_1)$  is a path between  $x$  and  $y$  in  $G' \setminus \{w\}$ , which is impossible. Thus  $u$  is different from  $x_k$  and  $y_l$  but adjacent to these vertices. But then  $(x, x_1, \dots, x_k, u, y_l, \dots, y_1)$  is a path from  $x$  to  $y$  in  $G' \setminus \{w\}$ , leading again to a contradiction. This shows that  $w$  cannot be an articulation point of  $G'$ . Since  $G'$  is not two-connected, we conclude that  $u$  is the only articulation point of  $G'$ . By the induction hypothesis, any block  $B$  of  $G'$  is dominated by some vertex  $b$ . Suppose that  $u$  does not dominate  $G'$ , for instance,  $u$  is not adjacent to some vertex  $t$  of  $B$ . Since  $u$  is the unique articulation point of  $G'$  but is not an articulation point of  $G$ ,  $v_1$  necessarily has a neighbor  $w \neq u$  in  $B$ . Hence, there is a path  $(v_1, w, b, t)$  of length 3 in  $G \setminus \{u\}$  and thus  $t$  is a neighbor of  $u$ , because  $t \in N_3(v_1, G \setminus \{u\}) \subseteq N_1(u)$ . Thus  $u$  dominates  $G' = G \setminus \{v_1\}$ , and, since  $v_1 \in N_1(u)$ ,  $u$  dominates  $G$  as well. This concludes the analysis of Case 2 and the proof of the theorem.  $\square$

## 4 Cop-win graphs for game with witness: class $\bigcap_{k \geq 1} \mathcal{CWW}(k)$

In this and next sections, we investigate the structure of  $k$ -winnable graphs. In analogy with big brother graphs, we characterize here the graphs  $G$  that are  $k$ -winnable for all  $k \geq 1$ , i.e., the graphs from the intersection  $\bigcap_{k \geq 1} \mathcal{CWW}(k)$ .

### 4.1 Game with witness: preliminaries

In the  $k$ -witness version of the game, the cop first selects his initial position and then the robber selects his initial position which is visible to the cop. As in the classical cop and robber game, the players move alternatively along an edge or pass. However, the robber is visible to the cop only every  $k$  moves. After having seen the robber, the cop decides a sequence of his next  $k$  moves (the first move of such a sequence is called a *visible* move). The cop captures the robber if they both occupy the same vertex at the same step (even if the robber is invisible). In particular, the cop can capture the visible robber if after the robber shows up, they occupy two adjacent vertices of the graph. Since we are looking for winning strategies for the cop, we may assume that the robber knows the cop's strategy, i.e., after each visible move, the robber knows the next  $k - 1$  moves of the cop. In the  $k$ -witness version of the game, a *strategy* for the cop is a function  $\sigma$  which takes as an input the  $i$  first visible positions of the robber and the  $ik$  first moves of the cop and outputs the next  $k$  moves of the cop. A winning strategy is defined as before and in any  $k$ -winnable graph, the cop has a positional winning strategy. We will call a *phase* of the game the movements of the two players comprised between two consecutive visible moves. We will call the behavior of the cop during several consecutive moves of the same phase  $\{a, b\}$ -*oscillating* if his moves alternate between the adjacent vertices  $a$  and  $b$ . In a  $k$ -winnable graph  $G$ , given a winning

cop's strategy  $\sigma$ , any trajectory  $S_r$  of the robber ends up in a vertex  $r_p$  at which the robber is captured. We will say that the trajectory  $S_r = (r_1, \dots, r_p)$  is *maximal* if  $(r_1, \dots, r_{p-1})$  cannot be extended to a longer trajectory for which the robber is not captured by the cop. Notice that the last vertex  $r_p$  in a maximal trajectory  $S_r$  corresponds to an invisible move if and only if it is a leaf of  $G$ . Indeed, otherwise let  $r_{p-1}$  be the previous position of the robber. If  $r_{p-1} \neq r_p$ , the robber could have stayed in  $r_{p-1}$  to avoid being captured. Thus  $r_{p-1} = r_p$  and if  $r_p$  has at least two neighbors, the robber can safely move to one of the neighbors of  $r_p$  not occupied by the cop, and survive for an extra unit of time. We continue with two simple observations, the first shows that during a phase an invisible robber can always safely move around a cycle, while the second shows that an invisible robber should always avoid an  $\{a, b\}$ -oscillating cop.

**Lemma 1.** *Suppose that at his move, the robber  $\mathcal{R}$  occupies a vertex  $v$  of a cycle  $C$  of a graph  $G$  and is not visible after this move. Then  $\mathcal{R}$  has a move (either staying at  $v$  or going to a neighbor of  $v$ ) such that the cop does not capture the robber during his next move.*

*Proof.* Let  $u$  be a neighbor of  $v$  in  $C$  which is not occupied by the cop. Since the robber will not be visible after his next move, the strategy of the cop is defined *a priori*. Let  $z$  be the next vertex to be occupied by the cop. Then the robber can stay at  $v$  if  $v \neq z$  or can move to  $u$  if  $u \neq z$ .  $\square$

**Lemma 2.** *If during one phase, the cop is performing  $\{a, b\}$ -oscillating moves and the robber moves to one of the vertices  $a$  or  $b$ , then the robber is captured either immediately or at the next move of the cop.*

*Proof.* Suppose that  $\mathcal{R}$  moves to the vertex  $a$ . If  $\mathcal{C}$  is located at  $a$ , then the robber is captured immediately. If  $\mathcal{C}$  is located at  $b$  and this is not the last vertex of the phase, then  $\mathcal{C}$  will move to  $a$  and will capture there the robber. Finally, if  $a$  and  $b$  are the positions of  $\mathcal{R}$  and  $\mathcal{C}$  at the end of the phase, then the robber will be visible at  $a$  and with the next visible move of  $\mathcal{C}$  from  $b$  to  $a$ , the robber will be caught at  $a$ .  $\square$

## 4.2 On the inclusion of $CWW(k+1)$ in $CWW(k)$

Clarke [18] noticed that for any  $k \geq 2$ , the inclusion  $CWFR(k) \subseteq CWW(k)$  holds. Contrary to the classes considered in the previous section which collapses for  $k \geq 3$ , we present now, for each  $k$ , an example of a graph in  $CWW(k) \setminus CWW(k+1)$ .

**Proposition 8.** *For any  $k \geq 2$ ,  $CWFR(k)$  is a proper subclass of  $CWW(k)$ . For any  $k \geq 1$ , there exists a graph contained in  $CWW(k) \setminus CWW(k+1)$ .*

*Proof.* To see the inclusion  $CWFR(k) \subseteq CWW(k)$  (which was also mentioned in [18]), it suffices to note that we can interpret the moves at speed  $k$  of the robber as if the cop moves only when the robber is visible (i.e., each  $k$ th move). Now, let  $S_3$  be the 3-sun, the graph on 6 vertices obtained by gluing a triangle to each of the three edges of another triangle (see Fig. 3(a)). Since no vertex of  $S_3$  has a maximum neighbor, the 3-sun is not dually chordal,



thus  $S_3 \notin \mathcal{CWF}\mathcal{R}(2)$  by Theorem 2. Then clearly,  $S_3$  is not a big brother graph either. On the other hand,  $S_3 \in \mathcal{C}\mathcal{W}\mathcal{W}(k)$  for any  $k \geq 2$ . Indeed, initially the cop is placed at a vertex  $u$  of degree 4. Then, the robber shows himself at the unique vertex  $v$  which is not adjacent to  $u$ . Let  $x$  and  $y$  be the two neighbors of  $v$  in  $S_3$ . The strategy of the cop consists in oscillating between  $x$  and  $y$  until the robber becomes visible again. Suppose without loss of generality that the cop's sequence of moves is  $x, y, x, y, \dots, y$ . Then from Lemma 2 we infer that  $\mathcal{R}$  is jammed at vertex  $v$ . At the end, when the robber shows his position again, then either he is at  $v$  or he desperately moves to  $x$ . In both cases, he is caught by  $\mathcal{C}$  at the next move. This shows that  $\mathcal{CWF}\mathcal{R}(k)$  is a proper subclass of  $\mathcal{C}\mathcal{W}\mathcal{W}(k)$

Now we will establish the second assertion. Let  $k \geq 1$  and  $G_k$  be the graph defined as follows. The vertex set of  $G_k$  is  $\{x, y, u, v, u_1, \dots, u_k, v_1, \dots, v_k\}$ . The vertex  $x$  is adjacent to any vertex except  $v$ , while  $y$  is adjacent to any vertex except  $u$ . For any  $i < k$ , the couples  $\{u_i, u_{i+1}\}, \{u_i, v_{i+1}\}, \{v_i, v_{i+1}\}, \{v_i, u_{i+1}\}$  are edges of  $G_k$ . Finally,  $u$  is adjacent to  $x, u_1$ , and  $v_1$ , while  $v$  is adjacent to  $y, u_k$ , and  $v_k$  ( $G_4$  is depicted in Figure 3(b)). To prove that  $G_k \in \mathcal{C}\mathcal{W}\mathcal{W}(k)$ , consider the following strategy for one cop. Initially, the cop occupies  $x$ . To avoid being caught immediately, the robber must show up at  $v$ . The cop occupies alternatively  $x$  and  $y$  in such a way that after  $k$  moves he is at  $y$  (if  $k$  is odd, then the cop passes his first move). Therefore, after  $k$  steps, the robber shows up at a vertex of  $N_k(v, G \setminus \{x, y\}) \cup \{x\} \subseteq N_1(y)$ , and at the next move the cop caught him. On the other hand, we assert that in  $G_k$  a robber with witness  $k + 1$  can evade against any strategy of the cop. Indeed, assume without loss of generality (in view of symmetry) that the initial position of the cop belongs to the set  $L = \{x, u, u_1, \dots, u_{\lceil k/2 \rceil}, v_1, \dots, v_{\lfloor k/2 \rfloor}\}$ . Then the robber chooses  $v$  (or  $v_1$  if  $k = 1$  and the cop is occupying  $u_k$ ) as his initial position. Let  $z$  be the vertex occupied by the cop after  $k + 1$  steps. If  $z \in L$ , then by Lemma 1 the robber can move in the triangle  $\{v, v_k, y\}$  in order to avoid the cop during the  $k + 1$  steps and to finish at a vertex of the triangle that is not adjacent to  $z$ . If  $z \notin L$ , then the robber uses the  $k + 1$  steps to reach  $u$  (or  $u_1$  if  $k = 1$  and  $z = v_1$ ). At any step, there is some  $i \leq k$ , such that the two vertices  $u_i$  and  $v_i$  allow the robber to decrease his distance to  $u$  (or to  $u_1$ ) by one; the robber chooses one of these vertices that is not occupied and will not be occupied by the cop after his move.  $\square$

**Open question 2:** Is it true that  $\mathcal{C}\mathcal{W}\mathcal{W}(k + 1) \subset \mathcal{C}\mathcal{W}\mathcal{W}(k)$ ?

### 4.3 $\bigcap_{k \geq 1} \mathcal{C}\mathcal{W}\mathcal{W}(k)$ and big two-brother graphs

In analogy to the big brother graphs, we say that a graph  $G$  is called a *big two-brother graph*, if  $G$  can be represented as an ordered union of subgraphs  $G_1, \dots, G_r$  in the form of a tree  $T$  rooted at  $G_1$  such that (1)  $G_1$  has a dominating vertex and (2) any  $G_i, i > 1$ , contains one or two adjacent vertices disconnecting  $G_i$  from its father and one of these two vertices dominates  $G_i$ . Note that if  $G_i$  and its father intersect in an articulation point  $x$ , then  $x$  is not necessarily the vertex which dominates  $G_i$ . Equivalently,  $G$  is a big two-brother graph if  $G$  can be represented as a union of its subgraphs  $G_1, \dots, G_r$ , labeled in such a way that  $G_1$  has a dominating vertex, and for any  $i > 1$ , either the subgraph  $G_i$  intersects  $\bigcup_{j < i} G_j$



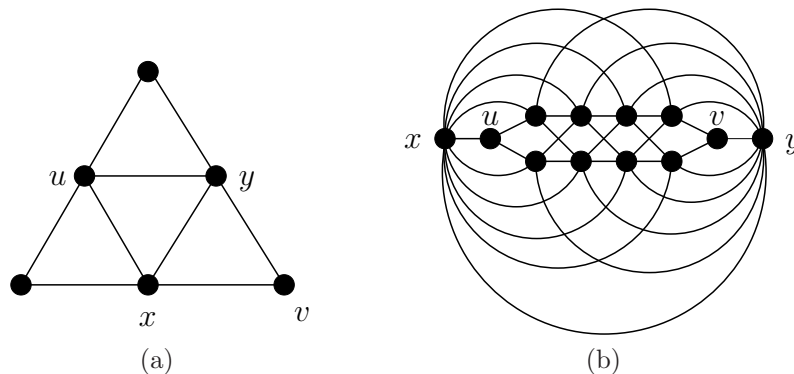


Figure 3: Two graphs in (a)  $\mathcal{CWW}(k) \setminus \mathcal{CWFR}(k)$ ,  $k \geq 2$  and (b)  $\mathcal{CWW}(4) \setminus \mathcal{CWW}(5)$ .

in two adjacent vertices  $x_i, y_i$  so that  $y_i$  dominates  $G_i$ , or  $G_i$  has a dominating vertex  $y_i$  and intersects  $\cup_{j < i} G_j$  in a single vertex  $x_i$  (that may coincide with  $y_i$ ); we will call such a decomposition  $G_1, \dots, G_r$  a *btb-decomposition* of  $G$ . The vertices  $y_i$  and  $x_i$  are the big and the small brothers of  $G_i$ . Let  $\mathcal{CWW}$  be the class of all big two-brother graphs. See Fig. 2(b) for an example of a big two-brother graph. As for big brother graphs, one can associate a rooted tree  $T$  with the decomposition  $G_1, \dots, G_r$  of a big two-brother graph  $G$ . Obviously any big brother graph  $G$  is also a big two-brother graph because the required union of subgraphs is provided by the block decomposition of  $G$  and  $x_i = y_i$  is the articulation point of the block  $G_i = B_i$  relating it with its father. The 2-trees and, more generally, the chordal graphs in which all minimal separators are vertices or edges are examples of big two-brother graphs which are not big brother graphs.

**Theorem 4.** *A graph  $G = (V, E)$  is  $k$ -winnable for all  $k \geq 1$  if and only if  $G$  is a big two-brother graph, i.e.,  $\mathcal{CWW} = \bigcap_{k \geq 1} \mathcal{CWW}(k)$ .*

*Proof.* First we show that any big two-brother graph  $G$  is  $k$ -winnable for any  $k \geq 1$ . Let  $G_1, \dots, G_r$  be a btb-decomposition of  $G$ . We consider the following strategy for the cop. The cop starts the game in the big brother of the root graph  $G_1$  and, more generally, at the beginning of each phase, we have the following property: the cop is located in the big brother  $y_i$  of some subgraph  $G_i$  such that the robber is located in a subgraph  $G_k$  that is a descendent of  $G_i$  in the decomposition tree  $T$  of  $G$ . If  $G_i = G_k$ , then the cop will capture the robber at the first move of the phase. Otherwise, let  $G_j$  be the son of  $G_i$  on the unique path of  $T$  between  $G_i$  and  $G_k$ . If  $G_i$  and  $G_j$  intersect in an articulation point  $x_j$ , then the cop moves from  $y_i$  to  $x_j$ , stays there during  $k - 2$  steps, and then, at the last step of the phase, if  $x_j$  is not the big brother  $y_j$  of  $G_j$ , he moves to  $y_j$ . If  $G_i$  and  $G_j$  intersect in an edge  $x_j y_j$  where  $y_j$  is the big brother of  $G_j$ , then the cop moves from  $y_i$  to one of the vertices  $x_j, y_j$  and then oscillate between  $x_j$  and  $y_j$  in such a way that when  $\mathcal{R}$  becomes visible again  $\mathcal{C}$  occupies the vertex  $y_j$  (the decision to move first to  $x_j$  or to  $y_j$  depends only on the parity of  $k$ ).

During this phase, the robber cannot leave the subgraph induced by the descendants of  $G_j$ , otherwise, he has to go from  $G_j$  to  $G_i$ . In the first case, the cop stays during the whole phase in the unique vertex  $x_j$  which cannot be traversed by the robber. In the second case, the cop oscillates between  $x_j$  and  $y_j$ ; therefore, by Lemma 2 the robber cannot traverse  $\{x_j, y_j\}$ . Therefore, after this phase, the invariant is preserved and the distance in  $T$  between the root and the subgraph  $G_j$  hosting the cop has strictly increased. Thus after at most diameter of  $T$  phases,  $\mathcal{R}$  and  $\mathcal{C}$  will be located in the same subgraph  $G_k$ , and the cop captures the robber.

Conversely, let  $G \in \mathcal{CWW}(k)$  for any  $k \geq 1$ . If  $G$  has a vertex  $z$  of degree 1, then  $G' = G \setminus \{z\}$  is a retract of  $G$ , thus  $G' \in \mathcal{CWW}(k)$  for any  $k \geq 1$ . Hence  $G'$  has a btb-decomposition  $G_1, \dots, G_{r-1}$  by induction hypothesis. If  $w$  is the unique neighbor of  $z$ , then setting  $G_r$  to be the edge  $zw$  and  $y_r = x_r := w$ , we will conclude that  $G$  is a big two-brother graph as well. So, we can suppose that  $G$  does not contain vertices of degree 1.

Since  $G \in \mathcal{CWW}(n^2)$ , applying Proposition 9 below for  $k = n$ , where  $n$  is the number of vertices of  $G$ , we deduce that  $G$  contains a vertex  $v$  and two adjacent neighbors  $x, y$  of  $v$  such that  $N_n(v, G \setminus \{x, y\}) \subseteq N_1(y)$ . This means that the connected component  $C$  of  $G \setminus \{x, y\}$  containing the vertex  $v$  is dominated by  $y$ . The graph  $G' := G(V \setminus C)$  is a retract of  $G$ , thus by Theorem 3 of [18]  $G' \in \mathcal{CWW}(k)$  for any  $k \geq 1$ . By induction assumption, either  $G'$  is empty or  $G'$  has a btb-decomposition  $G_1, \dots, G_{r-1}$ . If  $G'$  is empty, then, since  $y$  dominates  $C$ , we conclude that  $G$  has a btb-decomposition consisting of a single subgraph. Otherwise, setting  $G_r := G(C \cup \{x, y\})$ ,  $y_r := y$  and  $x_r := x$ , one can easily see that  $G_1, \dots, G_{r-1}, G_r$  is a btb-decomposition of  $G$ .  $\square$

**Proposition 9.** *Let  $G \in \mathcal{CWW}(k^2)$  for  $k \geq 1$ . If the minimum degree of a vertex of  $G$  is at least 2, then  $G$  contains a vertex  $v$  and an edge  $xy$  such that  $N_k(v, G \setminus \{x, y\}) \subseteq N_1(y)$ .*

*Proof.* If  $G$  contains a dominating vertex  $y$ , then the result follows by taking as  $x$  any vertex of  $G$  different from  $y$ . Assume thus that  $G$  does not have any dominating vertex. Consider a parsimonious winning strategy of the cop and suppose that the robber uses a strategy to avoid being captured as long as possible. Since  $G$  does not contain leaves, the robber is caught immediately after having been visible, i.e., at step  $pk^2 + 1$ . Since  $G$  does not have dominating vertices, the robber is visible at least twice, i.e.  $p \geq 1$ . Let  $y$  be the vertex occupied by the cop when the robber becomes visible for the last time before his capture. Let  $v$  be the next-to-last visible vertex occupied by the robber, i.e., his position at step  $(p-1)k^2 + 1$ , and let  $c_0$  be the vertex occupied by the cop at that moment. Finally, let  $S_c^p = (c_0, c_1, \dots, c_{k^2} = y)$  be the trajectory of the cop between the steps  $(p-1)k^2 + 1$  and  $pk^2 + 1$  (repetitions are allowed). Note that  $v \notin N_1(c_0)$ , otherwise the robber would have been caught immediately at step  $(p-1)k^2 + 1$ . We distinguish two cases depending on whether or not the cop occupies  $y$  at least once every two consecutive steps.

**Case 1:** There exists an index  $(p-1)k^2 + 1 \leq i < pk^2 - 1$  such that  $y \notin \{c_i, c_{i+1}\}$ .

Let  $i$  be the largest index satisfying the condition of Case 1 and set  $x := c_{i+1}$ . We will use the following assertion.

**Claim 1.** *If  $G$  contains a cycle  $C$  and a vertex  $w \in C$  such that  $d(v, w) < d(c_1, w)$ , then  $G \setminus \{x, y\}$  has a connected component that is dominated by  $y$ .*

*Proof.* Let us consider the vertex  $w$  satisfying the condition of the claim that is closest to  $v$ . If the claim is not satisfied, we will exhibit a strategy allowing the robber to escape the cop during more steps, contradicting the choice of the strategy of the robber. Suppose that at the beginning of the  $p$ th phase the robber move from  $v$  to  $w$  along a shortest  $(v, w)$ -path. Since  $d(v, w) < d(c_1, w)$ , the robber cannot be intercepted by the cop during these moves.

First assume that the cop reaches  $c_i$  before the step when the robber reaches  $w$ . We assert that  $y$  dominates its connected component  $C_y$  in  $G \setminus \{x\}$ . Let  $z$  be the current position of the robber at the step when the cop reaches  $c_i$ . Then  $z \in N_1(y)$ , otherwise the robber could have remain at  $z$  without being caught because starting with this step the cop moves only on vertices of  $N_1(y)$ . By way of contradiction, assume that there exists  $z' \in C_y \setminus N_1(y)$ . Consider such a vertex  $z'$  that is at distance exactly 2 from  $y$  in  $C_y$ , and let  $r \in C_y$  be a common neighbor of  $y$  and  $z'$  (note that  $r \neq x$ ). In this case the robber can go from  $z$  to  $y$  (the cop goes from  $c_i$  to  $x$ ), then from  $y$  to  $r$  (the cop goes from  $x$  to  $y$ ), and finally from  $r$  to  $z'$  where the robber can stay until he becomes visible. This contradicts the maximality of the strategy of the robber. Hence  $y$  dominates its connected component  $C_y$  in  $G \setminus \{x\}$ .

Now, assume that the robber reaches the vertex  $w$  before the  $i$ th step when the cop arrives at  $c_i$ . Then by Lemma 1 the robber can safely move on  $C$  until the cop reaches the vertex  $c_i$ . Let  $u \in C$  be the vertex occupied by the robber at the  $i$ th step when the cop arrives in  $c_i \neq u$ . If  $u \notin N_1(y)$ , then the robber could remain at  $u$  until step  $pk^2$  and would not be caught because the cop, after moving to  $c_i$ , at remaining moves of the phase will remain in the neighborhood of  $y$  and finishes the phase in  $y$ . Therefore  $u \in N_1(y)$ . If  $u = y$ , we assert that  $y$  dominates its connected component in the graph  $G \setminus \{x\}$  (such a component is not reduced to  $\{y\}$  because the degree of  $y$  is at least 2). If  $u \neq y$ , then we assert that  $y$  dominates the connected component of the graph  $G \setminus \{x, y\}$  containing  $u$ .

First suppose that  $u = y$  and assume by way of contradiction that  $G \setminus \{x\}$  has a vertex  $z$  at distance 2 from  $y$ . Let  $r \neq x$  be a common neighbor of  $z$  and  $y$ . Then the following sequence of moves is valid for the robber: after the cop went to  $x = c_{i+1} \neq u = y$ , the robber goes from  $u$  to  $r$ . Finally, the cop goes to  $y$  and the robber goes to  $z$ . After this step, by definition of  $c_i$ , the cop occupies only vertices in  $N_1(y)$  and finishes in  $y$ . Hence, the robber can remain at  $z$  and will not be captured at step  $k^2p + 1$ , a contradiction. Now suppose that  $u \neq y$  and assume by way of contradiction that  $G \setminus \{x, y\}$  contains a vertex  $z$  at distance 2 from  $y$ . Again, let  $r \neq x$  be a common neighbor of  $y$  and  $z$ . Then, the following sequence of moves is valid for the robber: when the cop arrives at  $c_i$ , the robber goes to  $y$ , then the cop goes to  $x = c_{i+1}$  and the robber goes to  $r$ . Finally, the cop goes to  $y$  and the robber goes to  $z$ . After this step, by definition of  $c_i$ , the cop occupies only vertices in  $N_1(y)$  and finishes in  $y$ . Hence, the robber can remain at  $z$  and will not be captured at step  $k^2p + 1$ , a contradiction. This concludes the proof of the claim.  $\square$

If the vertex  $v$  belongs to a cycle  $C$ , then setting  $w := v$  and applying Claim 1 we conclude that  $y$  dominates the connected component of  $G \setminus \{x, y\}$  containing  $v$ , establishing thus the assertion of Proposition 9. So, suppose that  $v$  is an articulation point of  $G$  not contained in

a cycle. Since the minimum degree of  $G$  is at least 2,  $G \setminus \{v\}$  has a connected component  $D$  that does not contain  $c_0$  (nor  $c_1$ ). Necessarily  $D$  contains a cycle  $C$ , otherwise we will find in  $D$  a vertex of degree 1 in  $G$ . Since any path from  $c_1$  to a vertex  $w$  of  $C$  passes via  $v$ , we obtain  $d(v, w) < d(c_1, w)$ . Then, the result follows from the claim. This concludes the analysis of Case 1.

**Case 2:** For any  $(p-1)k^2 \leq i \leq pk^2$  we have  $y \in \{c_i, c_{i+1}\}$ , i.e.,  $C$  occupies  $y$  at least once every 2 steps.

First, assume that there exists a vertex  $x$  (possibly  $x = y$ ) and indices  $q \geq k$  and  $(p-1)k^2 \leq i \leq pk^2 - q$  such that  $c_i, \dots, c_{i+q} \in \{y, x\}$ , i.e., that there are at least  $k$  consecutive steps when the cop remains at  $x$  or  $y$ . Then, we claim that  $N_q(v, G \setminus \{x, y\}) \subseteq N_1(y)$ . Indeed, pick  $z \in N_q(v, G \setminus \{x, y\})$  and let  $P = (v = p_1, \dots, p_q = z)$  be a shortest path in  $G \setminus \{x, y\}$  between  $v$  and  $z$ . Until the  $i$ th step of the phase, the robber may progress "slowly" along  $P$ : either by staying at his current position, or moving to the next vertex of  $P$  toward  $z$ , depending on the moves of the cop. The cop starts oscillating between  $x$  and  $y$  at step  $i$ . Then during the next  $q$  steps, the robber can follow  $P$  until he reaches  $z$  (since the length of  $P$  is at most  $q$ ). Therefore, if  $z$  is not a neighbor of  $y$ , then the robber can remain at  $z$  until step  $k^2p$  without being captured. Since by our assumption the robber is caught at step  $k^2p$ , necessarily  $z \in N_1(y)$ . Hence  $N_q(v, G \setminus \{x, y\}) \subseteq N_1(y)$  in this case. In particular, if  $q = k$ , the assertion of Proposition 9 holds.

Therefore, we may assume that between the steps  $(p-1)k^2$  and  $pk^2$ , for all  $k$  consecutive steps, the cop occupies at least three distinct vertices (and one of which is  $y$ ). We assert in this case that  $N_k(v, G \setminus \{y\}) \subseteq N_1(y)$ . Pick  $z \in N_k(v, G \setminus \{y\})$  and let  $P$  be a shortest path between  $v$  and  $z$  in  $G \setminus \{y\}$ . Then for any vertex  $w$  of  $P$ , among any sequence of  $k$  moves of the cop we can find three consecutive moves during which the cop does not occupy  $w$ . Therefore, for any sequence of  $k$  consecutive steps the robber can reduce by one his distance to  $z$  by moving on  $P$  towards  $z$  without being captured. Hence, he will reach  $z$  before step  $pk^2$ . If  $z$  is not adjacent to  $y$ , then staying at  $z$  the robber will not be captured, a contradiction. This concludes the proof of Proposition 9 and Theorem 4.  $\square$

## 5 Cop-win graphs for game with witness: classes $\mathcal{CWW}(k)$

In this section we investigate the dismantling orders related to  $k$ -winnable graphs. We provide a dismantling order which must be satisfied by all graphs of the class  $\mathcal{CWW}(2)$ . We show that this order is not sufficient but some its reinforcement is. Then we continue with similar results about  $k$ -winnable graphs for odd values of  $k \geq 3$ .

### 5.1 Class $\mathcal{CWW}(2)$

We continue with the definition of a dismantling ordering which seems to be intimately related with the witness variant of the cop and robber game. Again, we will consider a slightly more general version of the game: given a subset of vertices  $X$  of a graph  $G = (V, E)$ ,

the  $X$ -restricted  $k$ -witness game with cop and robber, is a variant in which  $\mathcal{R}$  can pass through any vertex of  $G$ ,  $\mathcal{C}$  can move only inside  $X$ , and all visible positions of the robber are at vertices of  $X$ . Then  $X$  is called  $k$ -winnable if for any starting positions of  $\mathcal{C}$  and  $\mathcal{R}$ , the cop wins in the  $X$ -restricted variant of the  $k$ -witness version of the game. We will say that a subset of vertices  $X$  of a graph  $G = (V, E)$  is  $k$ -bidismantlable if the vertices of  $X$  can be ordered  $v_1, \dots, v_m$  in such a way that for each vertex  $v_i, 1 \leq i < m$ , there exist two adjacent or coinciding vertices  $x, y$  with  $y = v_j, x = v_\ell$  and  $j, \ell > i$  such that  $N_k(v_i, G \setminus \{x, y\}) \cap X_i \subseteq N_1(y)$ , where  $X_i := \{v_i, v_{i+1}, \dots, v_m\}$  and  $X_m = \{v_m\}$ . We say that a graph  $G = (V, E)$  is  $k$ -bidismantlable if its vertex-set  $V$  is  $k$ -bidismantlable. In case  $k = 2$ , the inclusion  $N_2(v_i, G \setminus \{x, y\}) \cap X_i \subseteq N_1(y)$ , can be equivalently written as  $N_2(v_i, G \setminus \{x\}) \cap X_i \subseteq N_1(y)$ . Any  $(k, 1)$ -dismantlable graph is  $k$ -bidismantlable but the converse is not true: for any  $k \geq 2$ , the 3-sun  $S_3$  presented in Fig. 3 is  $k$ -bidismantlable but not  $(k, 1)$ -dismantlable. In some proofs, we will denote by  $x(v)$  and  $y(v)$  the vertices eliminating a vertex  $v$  in a  $k$ -bidismantling order.

**Proposition 10.** *Any graph  $G = (V, E)$  of  $CWW(2)$  is 2-bidismantlable.*

*Proof.* Suppose that a subset  $X$  of vertices (inducing a connected subgraph  $G(X)$ ) of  $G$  is 2-winnable. We will show that  $X$  is 2-bidismantlable. Consider a parsimonious winning strategy of the cop and a maximal trajectory of the robber. If the capture happened when  $\mathcal{R}$  is invisible, then, analogously to what we noticed above, the robber is located at a leaf  $v_1$  of  $G(X)$ . In this case, set  $x = y$  to be the unique neighbor of  $v_1$  in  $G(X)$ . Then  $N_2(v_1, G \setminus \{x\}) \cap X = \{v_1\} \subseteq N_1(y)$  holds (otherwise, the robber can live longer). Now suppose that  $\mathcal{C}$  captures  $\mathcal{R}$  at the next visible move. This means that when  $\mathcal{C}$  sees  $\mathcal{R}$ , the cop is located in some vertex  $y \in X$  and the robber is located in some vertex  $w \in X$  and  $w \in N_1(y)$  holds. Then the cop moves from  $y$  to  $w$  and captures  $\mathcal{R}$  there. Denote by  $v_1$  the vertex of  $X$  where  $\mathcal{R}$  is visible for the next-to-last time. Suppose that after having seen the robber in  $v_1$ , the cop moves first to a vertex of  $X$  which we denote by  $x$  and then to vertex  $y$  (notice that  $y$  is either a neighbor of  $x$  or it coincides with  $x$ ). In response to these moves of the cop, the robber first moves to some vertex  $u \in N_1(v_1) \setminus \{x\}$  and then to the vertex  $w \in N_1(u) \cap X \subseteq (N_2(v_1, G \setminus \{x\}) \cup \{x\})$ . By the definition of the vertices  $y$  and  $w$ , in  $y$  the cop sees (for the last time) the robber which is located at  $w$  and with the next move captures him. Since  $\mathcal{R}$  follows a maximal sequence of moves before his capture, any vertex of  $N_2(v_1, G \setminus \{x\}) \cap X$  must be adjacent to  $y$ , otherwise, if there exists  $z \in N_2(v_1, G \setminus \{x\}) \cap X$  not adjacent to  $y$ , instead of moving to  $w$ , in two moves the robber can safely reach  $z$  and survive for a longer time. Thus indeed  $N_2(v_1, G \setminus \{x\}) \cap X \subseteq N_1(y)$  holds.

We show that the set  $X' := X \setminus \{v_1\}$  is 2-winnable as well. Consider a positional parsimonious winning strategy  $\sigma$  of the cop in  $X$ . For any positions  $c$  of the cop and  $r$  of the robber in  $X'$ , we note  $\sigma(c, r) = (c_1, c_2)$ . As in the proof of Theorem 1, we construct a strategy that uses one bit of memory  $m$ : it is a function that associates to each  $(c, r, m)$  a couple  $((c'_1, c'_2), m)$ . As in the proof of Theorem 1, the intuitive idea is that the cop plays using  $\sigma$ , except when he is in  $y$  and his memory contains 1; in that case, he plays using  $\sigma$  as if he was in  $v_1$ .

If  $m = 0$  or  $c \neq y$ , let  $(c_1, c_2) = \sigma(c, r)$ . If  $c_1 = v_1$ , then  $c'_1 = y$  and  $c'_2 = c_1$  otherwise. If  $c_2 = v_1$ , then  $\sigma'(c, r, m) = ((c'_1, y), 1)$  and  $\sigma'(c, r, m) = ((c'_1, c_2), 0)$  otherwise. If  $m = 1$  and  $c = y$ , let  $(c_1, c_2) = \sigma(v_1, r)$ . If  $c_1 = v_1$ , then  $c'_1 = y$  and  $c'_2 = c_1$  otherwise. If  $c_2 = v_1$ , then  $\sigma'(y, r, 1) = ((c'_1, y), 1)$  and  $\sigma'(y, r, 1) = ((c'_1, c_2), 0)$  otherwise. Since  $N_1(v_1) \cap X \subseteq N_1(y)$ , one can easily check that  $\sigma'$  is a valid strategy for the  $X'$ -restricted game.

By way of contradiction, suppose now that there exists an infinite  $X'$ -valid sequence  $S'_r$  of moves of the robber in the  $X'$ -restricted game allowing him to escape forever against a cop using the strategy  $\sigma'$ . First note that the sequence of moves  $S_c$  of the cop playing  $\sigma$  against  $S'_r$  differs from the sequence of moves  $S'_c$  of the cop playing  $\sigma'$  against  $S'_r$  only in the positions where the cop is in  $v_1$  in  $S_c$ .

We show that there exists an infinite sequence  $S_r$  in the  $X$ -restricted game enabling the robber to escape forever against a cop using the strategy  $\sigma$ . The visible positions  $\mathcal{R}$  of in  $S_r$  will coincide with the visible positions of  $\mathcal{R}$  in  $S'_r$  (thus the cop's strategies  $\sigma$  and  $\sigma'$  behave in the same way against both sequences). It is sufficient to show that if during a phase of  $S'_r$ , the robber goes from  $r'_0 \in X'$  to  $r'_2 \in X'$  via  $r'_1 \in V(G)$ , then in the  $X$ -restricted game where the cop plays with strategy  $\sigma$  (going first to  $c_1$  and then to  $c_2$ ), there exists  $r_1$  such that  $\mathcal{R}$  can go from  $r'_0$  to  $r'_2$  via  $r_1$  without being captured in  $r_1$ .

If  $r'_1 \neq v_1$ , or if  $v_1 \notin \{c_1, c_2\}$  one can choose  $r_1 = r'_1$  (since  $r'_0, r'_2 \in X'$ , they are different from  $v_1$ ). Thus, we may assume that  $r'_1 = v_1$  and that  $c_1 = v_1$  or  $c_2 = v_1$ . If  $c_2 \in \{v_1, y\}$ , then  $c'_2 = y$ . Since  $r'_1 = v_1$ ,  $r'_2 \in N_1(v_1) \cap X \subseteq N_1(y)$  and thus the robber is captured when he shows up in  $r'_2$ , i.e.,  $S'_r$  does not enable the robber to escape forever. Consequently,  $c_2 \notin \{v_1, y\}$ ; suppose  $c_1 = v_1$  and let  $r_1 = y$ . In this case,  $(r'_0, y, r'_2)$  is a  $X$ -valid sequence since  $r'_0, r'_2 \in N_1(v_1) \cap X \subseteq N_1(y)$  and moreover  $y \notin \{c_1, c_2\}$  (since  $c_1 = v_1$  and  $y \neq c_2$ ). It implies that there exists an infinite  $X$ -valid sequence  $S_r$  enabling the robber to escape forever, a contradiction.

Starting from a positional strategy for the  $X$ -restricted game, we have constructed a winning strategy using memory for the  $X'$ -restricted game. As mentioned in the introduction, it implies that there exists a positional winning strategy for the  $X'$ -restricted game. Consequently, the set  $X' := X \setminus \{v_1\}$  is 2-winnable as well. By induction assumption,  $X'$  admits a 2-bidismantling order  $v_2, \dots, v_m$ . Then clearly  $v_1, v_2, \dots, v_m$  is a 2-bidismantling of  $X$ . If  $G$  is 2-winnable, then its set of vertices is 2-winnable and therefore 2-bidismantlable, showing that  $G$  is 2-bidismantlable.  $\square$

We continue with two examples. The first one shows that we cannot replace in the definition of 2-bidismantlability the condition  $N_2(v_i, G \setminus \{x\}) \cap X_i \subseteq N_1(y)$  by a weaker condition  $N_2(v_i, G_i \setminus \{x\}) \subseteq N_1(y)$  (i.e., instead of all vertices of  $X_i$  reachable from  $v_i$  by paths of length 2 avoiding  $x$  of the whole graph  $G$  to consider only the vertices reachable by such paths of the subgraph  $G_i$ ). The second example shows that unfortunately 2-bidismantlability is not a sufficient condition.

**Proposition 11.** *Let  $G$  be the graph from Fig. 4. Then  $G$  admit a dismantling order satisfying the condition  $N_2(v_i, G_i \setminus \{x\}) \subseteq N_1(y)$ , however  $G$  is not 2-bidismantlable nor 2-winnable.*

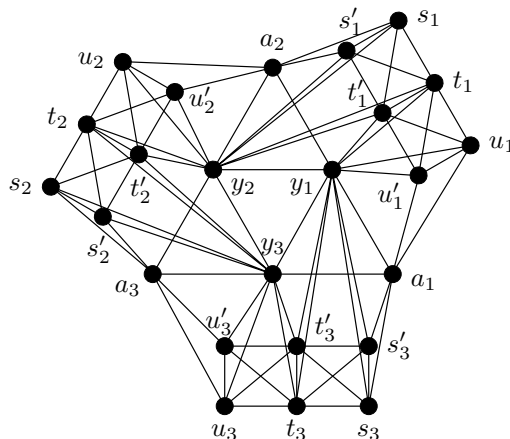


Figure 4: A weakly 2-bidismantlable graph that is not 2-bidismantlable

*Proof.* Consider the following order on the vertices of  $G$ :  $a_1, a_2, a_3, u_1, u_1', u_2, u_2', u_3, u_3', s_1, s_1', s_2, s_2', s_3, s_3', t_1, t_1', t_2, t_2', t_3, t_3', y_1, y_2, y_3$ . For each vertex  $v \in V(G) \setminus \{y_3\}$ , we give below two adjacent vertices  $x(v), y(v)$  that are eliminated later than  $v$  and such that  $N_2(v, G_i \setminus \{x(v), y(v)\}) \subseteq N_1(y(v))$  (when  $N_2(v, G_i \setminus \{y(v)\}) \subseteq N_1(y(v))$ , the vertex  $x(v)$  is not defined).

$v$	$a_1$	$a_2$	$a_3$	$u_1$	$u_1'$	$u_2$	$u_2'$	$u_3$	$u_3'$	$s_1$	$s_1'$	$s_2$
$y(v)$	$y_1$	$y_2$	$y_3$	$t_1$	$t_1'$	$t_2$	$t_2'$	$t_3$	$t_3'$	$y_2$	$y_2$	$y_3$
$x(v)$	$y_3$	$y_1$	$y_2$	$y_1$	$y_1$	$y_2$	$y_2$	$y_3$	$y_3$	—	—	—
$v$	$s_2'$	$s_3$	$s_3'$	$t_1$	$t_1'$	$t_2$	$t_2'$	$t_3$	$t_3'$	$y_1$	$y_2$	$y_3$
$y(v)$	$y_3$	$y_1$	$y_1$	$y_2$	$y_2$	$y_3$	$y_3$	$y_1$	$y_1$	$y_2$	$y_3$	—
$x(v)$	—	—	—	$y_1$	$y_1$	—	—	—	—	—	—	—

We prove now that  $G$  is not 2-bidismantlable. Note that for  $a_1$  (resp.  $a_2, a_3$ ), there exists  $y(a_1) = y_1$  (resp.  $y(a_2) = y_2, y(a_3) = y_3$ ) and  $x(a_1) = y_3$  (resp.  $x(a_2) = y_1, x(a_3) = y_2$ ) such that  $N_2(a_1, G \setminus \{y_1, y_2\}) \subseteq N_1(y_1)$  (resp.  $N_2(a_2, G \setminus \{y_2, y_3\}) \subseteq N_1(y_2), N_2(a_3, G \setminus \{y_3, y_1\}) \subseteq N_1(y_3)$ ). Consequently, any 2-bidismantling order of  $G$  can start with  $a_1, a_2, a_3$ . In fact, one can check that any 2-bidismantling of  $G$  must start with a permutation of  $a_1, a_2, a_3$ . We will show now that it is impossible to extend a 2-bidismantling order starting with  $a_1, a_2, a_3$ . To prove this, it suffices to show that for any  $v \in V(G) \setminus \{a_1, a_2, a_3\}$  and for all adjacent vertices  $x(v), y(v) \in N_1(v)$ , there exists a vertex  $z(v) \in N_2(v, G \setminus \{x(v), y(v)\}) \setminus (\{a_1, a_2, a_3\})$  such that  $z(v) \notin N_1(y(v))$ . In view of symmetry of  $G$ , it is sufficient to check this property for  $v \in \{u_1, t_1, y_1\}$ .



If  $v = u_1$ , then  $y(u_1), x(u_1) \in \{a_1, u'_1, t_1, t'_1, y_1\}$ . If  $y(u_1) \in \{a_1, u'_1, y_1\}$ , then either  $t_1 \neq x(v_1)$ , or  $t'_1 \neq x(v_1)$ ; in both cases,  $s_1 \in N_2(v, G \setminus \{x(u_1), y(u_1)\})$  and  $s_1 \notin N_1(y(u_1))$ . By symmetry, we can suppose that  $y(u_1) = t_1$ . Since  $a_1 \notin N_1(t_1)$ , we must have  $x(u_1) \neq a_1$  and consequently,  $s_3 \in N_2(v, G \setminus \{x(u_1), y(u_1)\})$  and  $s_3 \notin N_1(y(u_1))$ .

If  $v = t_1$ , then  $y(t_1), x(t_1) \in \{u_1, u'_1, s_1, s'_1, t'_1, y_1, y_2\}$ . If  $y(t_1) \in \{y_1, u_1, u'_1\}$  (resp.  $y(t_1) = \{y_2, s_1, s'_1\}$ ), then set  $z(v) = s_1$  (resp.  $z(v) = u_1$ ); in all cases,  $z(v) \in N_2(v, G \setminus \{x(t_1), y(t_1)\})$  and  $u \notin N_1(y(t_1))$ . If  $y(t_1) = t'_1$ , then either  $y_2 \neq x(t_1)$  or  $y_1 \neq x(t_1)$ ; in both cases,  $y_3 \in N_2(v, G \setminus \{x(t_1), y(t_1)\})$  and  $y_3 \notin N_1(y(t_1))$ .

If  $v = y_1$ , then, since  $N_1(y_1) \subseteq N_1(y(y_1))$ , the vertex  $y(y_1)$  must belong to  $N_1(y_1) \cap N_1(y_2) \cap N_1(y_3)$ . Consequently, by symmetry, we can assume  $y(y_1) = y_2$ . However, since  $u_1 \in N_1(y_1) \setminus N_1(y_2)$ , we obtain  $u_1 \in N_2(y_1, G \setminus \{x(y_1), y(y_1)\})$  and  $u_1 \notin N_1(y(y_1))$ . This completes the proof that  $G$  is not 2-bidismantlable. Since any graph  $G \in \mathcal{CWW}(2)$  is 2-bidismantlable, it also implies that  $G \notin \mathcal{CWW}(2)$ .  $\square$

**Proposition 12.** *Let  $G$  be the graph from Fig. 5. Then  $G$  is 2-bidismantlable, however  $G \notin \mathcal{CWW}(2)$ .*

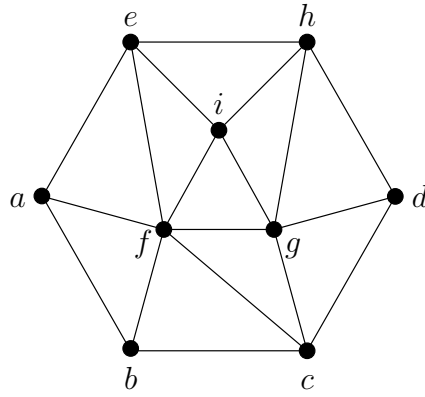


Figure 5: A 2-bidismantlable graph  $G \notin \mathcal{CWW}(2)$

*Proof.* The graph presented in Fig. 5 is 2-bidismantlable with the following 2-bidismantling order  $a, b, c, d, e, f, g, h, i$ , where each vertex  $v$  is eliminated by the vertices  $x(v), y(v)$  defined as follows:

$v$	$a$	$b$	$c$	$d$	$e$	$f$	$g$	$h$	$i$
$y(v)$	$f$	$f$	$g$	$g$	$i$	$i$	$i$	$i$	$-$
$x(v)$	$e$	$c$	$f$	$h$	$-$	$-$	$-$	$-$	$-$

However, one can show that for any vertex  $c$  there exists a vertex  $r$  such that, if at step  $i \geq 0$  the cop moves to (or starts in)  $c$  (going through any intermediate vertex), then the



robber can move to (or starts in)  $r$  without being caught. Since for any such couple  $(c, r)$  the vertices  $c$  and  $r$  are not adjacent, it means that the cop cannot catch the robber in this graph. The definition of the pairs  $(c, r)$  is given in the following table:

$c$	$a$	$b$	$c$	$d$	$e$	$f$	$g$	$h$	$i$
$r$	$d$	$d$	$e$	$e$	$b$	$d$	$e$	$b$	$b$

Note that if the robber wants to go from  $d$  to  $e$ , (resp., from  $e$  to  $d$ ), then this means that the cop is in  $a, b$  or  $f$  (resp., wants to go to  $a, b$  or  $f$ ). Since  $h \notin N(a) \cup N(b) \cup N(f)$ ,  $h$  cannot be the intermediate vertex used by the cop. Thus, the robber can always go from  $d$  to  $e$  (resp. from  $e$  to  $d$ ) via  $h$ .

If the robber wants to go from  $b$  to  $d$  (resp., from  $d$  to  $b$ ), then this implies that the cop is in  $e, h, i$  (resp., wants to go to  $e, h, i$ ). Since  $c \notin N(e) \cup N(h) \cup N(i)$ ,  $c$  cannot be the intermediate vertex used by the cop. Thus, the robber can always go from  $b$  to  $d$  (resp. from  $d$  to  $b$ ) through  $c$ .

If the robber wants to go from  $b$  to  $e$  (resp. from  $e$  to  $b$ ), then this means that the cop neither starts in  $a$  nor  $f$  (because in this case the robber would have been in  $d$ ), nor goes to  $a$  or  $f$  (since in this case, the robber wants to go in  $d$ ). Moreover, the intermediate vertex used by the cop is different from  $a$  or  $f$ . In the first case (resp. second case), the robber can go from  $b$  to  $e$  via  $a$  (resp.  $f$ ).  $\square$

We continue with a condition on 2-bidismantling which turn out to be sufficient for 2-winability. We say that a graph  $G$  is *strongly 2-bidismantlable* if  $G$  admits a 2-bidismantling order such that for any vertex  $v_i, i < n$ ,  $y(v_i) = x(v_i)$  or  $N_2(v_i, G \setminus \{y(v_i)\}) \cap X_i \subseteq N_2(x(v_i), G \setminus \{y(v_i)\})$  (recall that  $x(v)$  and  $y(v)$  denote the vertices eliminating a vertex  $v$  in a 2-bidismantling order).

**Proposition 13.** *If a graph  $G$  is strongly 2-bidismantlable, then  $G \in CWW(2)$ .*

*Proof.* Suppose that a subset  $X$  of vertices of  $G$  admits a strong 2-bidismantling order  $v_1, \dots, v_m$ . Assume by induction assumption that the set  $X' = \{v_2, \dots, v_n\}$  is 2-winnable and we will establish that the set  $X$  itself is 2-winnable. Let  $N_2(v_1, G \setminus \{x\}) \cap X \subseteq N_1(y)$ .

Let  $\sigma'$  be a parsimonious positional winning strategy for  $\mathcal{C}$  in  $X'$ . We define the strategy  $\sigma$  for  $\mathcal{C}$  in  $X$  as follows:  $\sigma(c, r) = r$  if  $r \in N_1(c)$ ,  $\sigma(c, v_1) = (x, y)$  if  $c \in N_1(x)$  (in this case, the robber will be caught during the next move because  $N_2(v_1, G \setminus \{x\}) \cup X \subseteq N_1(y)$ ) and  $\sigma(c, v_1) = \sigma'(c, x)$  otherwise, and  $\sigma(c, v) = \sigma'(c, v)$  in all other cases. We now prove that  $\sigma$  is winning. Let  $S_r = (r_1, r_2, \dots)$  be any  $X$ -valid sequence of moves of the robber. We will transform  $S_r$  into a  $X'$ -valid sequence  $S'_r = (r'_1, r'_2, \dots)$  of moves of the robber and prove that, since  $\sigma'$  eventually captures  $\mathcal{R}$  following  $S'_r$ , then  $\mathcal{C}$  playing  $\sigma$  captures  $\mathcal{R}$  following  $S_r$ .

Let  $r'_1 := x$  if  $r_1 = v_1$ , and  $r'_1 := r_1$  otherwise. Suppose that  $r'_1, \dots, r'_{2j-1}$  ( $j \geq 1$ ) have been already defined and we wish to define  $r'_{2j}$  and  $r'_{2j+1}$ . We set  $r'_{2j+1} := r_{2j+1}$  if  $r_{2j+1} \neq v_1$ , and  $r'_{2j+1} := x$  otherwise (indeed, when the cop sees the robber in the vertex  $v_1$ , then  $\mathcal{C}$  will play against  $\mathcal{R}$  as like the latter was in  $x$ ). We set  $r'_{2j} := r_{2j}$  in all cases

unless  $v_1 \in \{r_{2j+1}, r_{2j-1}\}$  and  $r_{2j} \notin N_1(x)$  (in particular  $r_{2j} \neq y$ ). If  $r_{2j-1} = v_1$  (resp., if  $r_{2j+1} = v_1$ ) and  $r_{2j} \notin N_1(x)$ , then there exists a common neighbor  $u$  of  $r_{2j-1}$  (resp.,  $r_{2j+1}$ ) and  $x$  different from  $y$ . The choice of  $r'_{2j}$  depends of the current position  $c_{2j}$  of the cop pursuing  $\mathcal{R}$ . We set  $r'_{2j} := u$  if  $c_{2j} \neq u$  and  $r'_{2j} := y$  otherwise (this is to avoid to artificially create a move where the robber goes to a vertex occupied by the cop). It can be easily seen that  $S'_r$  is a  $X'$ -valid sequence of moves of the robber.

Let  $S'_c = (c'_1, c'_2, \dots)$  be the  $X'$ -valid sequence of moves of the cop playing  $\sigma'$  against a robber  $\mathcal{R}'$  moving according to  $S'_r$ , and let  $S_c = (c_1, c_2, \dots)$  be the  $X$ -valid sequence of moves of the cop playing  $\sigma$  against the robber  $\mathcal{R}$  following  $S_r$ . It is easy to check that  $S'_c$  and  $S_c$  are similar except one or two steps before the capture of the robber. Moreover, since  $\sigma'$  is a winning strategy in  $X'$ , there is  $j > 0$  such that  $c'_j = r'_j$ .

First suppose that  $\mathcal{C}$  captures the robber  $\mathcal{R}'$  when he is visible, say  $\mathcal{R}'$  is located in  $r'_{2j+1}$ . If  $r'_{2j+1} = r_{2j+1}$ , then we are done. So, suppose that  $r'_{2j+1} \neq r_{2j+1}$ , i.e.,  $r_{2j+1} = v_1$  and  $r'_{2j+1} = x$ . Therefore, when he sees  $\mathcal{R}$  in  $v_1$ ,  $\mathcal{C}$  is located in a neighbor of  $x$ . According to  $\sigma$ ,  $\mathcal{C}$  will move to  $x$  and then to  $y$ , while  $\mathcal{R}$  can only reach a vertex in  $N_2(v_1, G \setminus \{x\}) \cup X$ . Since  $N_2(v_1, G \setminus \{x\}) \cap X \subseteq N_1(y)$ , the cop will capture the visible robber at his next move.

Now suppose that  $\mathcal{C}$  captures  $\mathcal{R}'$  when the latter is invisible, say  $\mathcal{R}'$  is located in  $r'_{2j}$ . Again, if  $r'_{2j} = r_{2j}$ , then we are done. Otherwise, according to the definition of  $S'_r$ , we conclude that  $r_{2j}$  is a common neighbor of  $r_{2j-1}$  and  $r_{2j+1}$  different from  $y$  with either  $v_1 = r_{2j+1}$  or  $v_1 = r_{2j-1}$ . Suppose that  $v_1 = r_{2j+1}$  (the other case is analogous),  $r'_{2j}$  is either  $y$  or a common neighbor  $u$  of  $r_{2j-1}$  and  $x$  provided by the strong 2-bidismantling order. Since, between  $r_{2j-1}$  and  $r_{2j+1} = v_1$  the trajectory of  $\mathcal{R}'$  avoids the cop if possible, we deduce that  $\{c_{2j-1}, c_{2j}\} = \{u, y\}$  or  $\{c_{2j}, c_{2j+1}\} = \{u, y\}$ . If  $\{c_{2j-1}, c_{2j}\} = \{u, y\}$ , then, when  $\mathcal{C}$  sees  $\mathcal{R}$  in  $r_{2j-1}$ , the cop is located in a neighbor of  $r_{2j-1}$ . By the definition of  $\sigma$ ,  $\mathcal{C}$  will move to  $r_{2j-1}$  and captures  $\mathcal{R}$ . Otherwise, if  $\{c_{2j}, c_{2j+1}\} = \{u, y\}$ , and, when he sees  $\mathcal{R}$  in  $v_1$ ,  $\mathcal{C}$  is located in a neighbor of  $x$ . By the definition of  $\sigma$ , as before,  $\mathcal{C}$  will move to  $x$  and then to  $y$ , while  $\mathcal{R}$  can only reach a vertex in  $N_2(v_1, G \setminus \{x\}) \cup X$ . Since  $N_2(v_1, G \setminus \{x\}) \cap X \subseteq N_1(y)$ , the cop will capture the visible robber at his next move.  $\square$

We conclude this section by showing that the existence of a strong 2-bidismantling order is not necessary.

**Proposition 14.** *The graph  $G$  from Fig. 6 belongs to  $CWW(2)$ , however  $G$  does not admit a strong 2-bidismantling order.*

*Proof.* We first show that the graph  $G$  from Fig. 6 is in  $CWW(2)$ , by exhibiting a strategy for the cop. The cop starts in  $u$ . Hence, if the robber starts in  $x, x', y_1$  or  $y_2$ , he is immediately caught. If the robber starts in  $s_1$  (or  $s_2$ ), then the cop moves to  $y_1$  (resp., to  $y_2$ ) and since  $N_2(s_1) \subseteq N_1(y_1)$  (resp.,  $N_2(s_2) \subseteq N_1(y_2)$ ), the robber is caught the next time he shows up. If the robber starts in  $v_1$  (the cases  $v'_1, v_2, v'_2$  are similar), then the cop first moves to  $x$  and then to  $y_1$ . Then the robber has to show up in a vertex of  $\{v_1, s_1, v'_1, x\} \subseteq N_1(y_1)$  and the cop can catch him.

Consider now any 2-bidismantling order of  $G$ . Consider the first vertex  $v$  different from  $s_1, s_2$  that appears in the 2-bidismantling order of  $G$ . We may assume without loss of

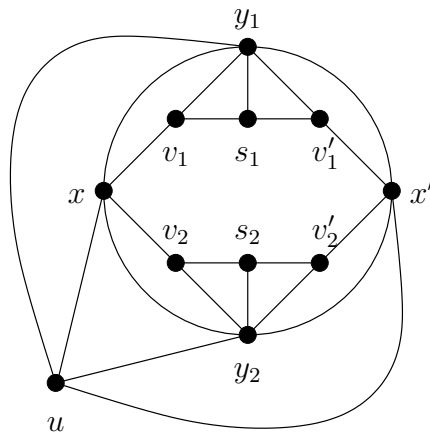


Figure 6: A graph  $G \in \mathcal{CWW}(2)$  which is not strongly 2-bidismantlable.

generality that  $v \in \{v_1, x, y_1, u\}$ . Let  $X = V(G) \setminus \{s_1, s_2\}$ . Since there does not exist  $t$  such that  $N_1(x) \cap X$  (resp.  $N_1(y_1) \cap X$  or  $N_1(u) \cap X$ ) is included in  $N_1(t)$ , it implies that  $v \in \{v_1, v'_1, v_2, v'_2\}$ . By symmetry, we can assume that  $v = v_1$ . We know that  $x(v_1), y(v_1) \in N_1(v_1)$  and that  $N_1(v_1) \cap X \subseteq N_1(y(v_1))$ . Consequently,  $y(v_1) \in \{x, y_1\}$ . If  $y(v_1) = x$ , then  $x(v_1) \neq s_1$  (since  $s_1$  and  $x$  are not adjacent) and thus  $v'_1 \in N_2(v_1, G \setminus \{x(v_1), y(v_1)\}) \setminus N_1(y(v_1))$ , which is impossible. Thus,  $y(v_1) = y_1$ . If  $x(v_1) \neq x$ , then  $v_2 \in N_2(v_1, G \setminus \{x(v_1), y(v_1)\}) \setminus N_1(y_1)$ . Consequently,  $y(v_1) = y_1$  and  $x(v_1) = x$ . However,  $v'_1 \in N_2(v_1) \cap X$  but  $v'_1 \notin N_2(x, G \setminus \{y_1\})$  and thus  $G$  is not strongly 2-bidismantlable.  $\square$

## 5.2 Classes $\mathcal{CWW}(k)$ for $k \geq 3$

In this subsection we show that  $k$ -bidismantlable graphs are  $k$ -winnable for any odd  $k \geq 3$ . We also show that for any  $k \geq 3$ , there exist graphs in  $\mathcal{CWW}(k)$  that are not  $k$ -bidismantlable, i.e., for  $k \geq 3$ ,  $k$ -bidismantlability of a graph is not a necessary condition to be  $k$ -winnable.

**Theorem 5.** *For any odd integer  $k \geq 3$ , if a graph  $G$  is  $k$ -bidismantlable, then  $G \in \mathcal{CWW}(k)$ .*

*Proof.* Suppose that  $X \subseteq V$  is a  $k$ -bidismantlable set of vertices of a graph  $G$ . We prove that there is a winning strategy for the cop in the  $X$ -restricted  $k$ -witness game on  $G$ . To do so, we proceed as in the papers [27, 30] and use the  $k$ -bidismantling order to mark all  $X$ -configurations  $(c, r)$ . A  $X$ -configuration of  $X$ -restricted game is a couple  $(c, r)$  that consists of a position of the cop  $c \in X$  and a position of the robber  $r \in X$ , with  $r \neq c$ . A  $X$ -configuration  $(c, r)$  is called *terminal* if  $r \in N_1(c)$ .

To mark the  $X$ -configurations, we use the following procedure  $\text{Mark}(X)$ .

1. Initially, all  $X$ -configurations are unmarked.
2. Any *terminal*  $X$ -configuration  $(c, r)$  is marked with label 1.
3. While it is possible, mark an unmarked  $X$ -configuration  $(c, r)$  with the smallest possible integer  $\ell + 1$  such that there exist vertices  $y_{(c,r)} \in N_1(c) \cap X$  and  $x_{(c,r)} \in (N_1(y_{(c,r)}) \setminus \{r\}) \cap X$  such that for all  $z \in N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\}) \cap X$ , the  $X$ -configuration  $(y_{(c,r)}, z)$  is marked with a label at most  $\ell$ .

**Claim 2.** *If all  $X$ -configurations are marked by  $\text{Mark}(X)$ , then there is a winning strategy for the cop in the  $X$ -restricted  $k$ -witness game on  $G$ .*

Indeed, pick any initial positions  $c \in X$  of the cop and  $r \in X$  of the robber. If the configuration  $(c, r)$  is terminal, then  $r \in N_1(c)$  and the robber is captured at the next move. Otherwise, the cop first moves to  $y_{(c,r)}$  and then oscillates between  $x_{(c,r)}$  and  $y_{(c,r)}$  during  $k - 1$  steps, i.e., the cop ends in  $y_{(c,r)}$  since  $k$  is odd. If during one of his invisible moves the robber goes to  $x_{(c,r)}$  or  $y_{(c,r)}$ , then he will be captured immediately. Otherwise, in  $k$  moves the robber goes from  $r$  to a vertex  $z \in N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\}) \cap X$ . According to  $\text{Mark}(X)$ , the label of  $(y_{(c,r)}, z)$  is strictly less than that of  $(c, r)$ . Therefore, by repeating the same process, after a finite number of steps either the cop captures the robber during an invisible move or the cop and the robber arrive at a terminal configuration.

**Claim 3.** *If  $X$  is  $k$ -bidismantlable, then  $\text{Mark}(X)$  marks all  $X$ -configurations.*

The general idea of our proof follows the proof of Theorem 12 of [27]. Let  $\{v_1, \dots, v_t\}$  be a  $k$ -bidismantling ordering of  $X$ . We prove by induction on  $t - i$  that  $\text{Mark}(X_i)$  marks all  $X_i$ -configurations, where  $X_i = \{v_i, \dots, v_t\}$ . The assertion trivially holds for  $X_{t-1}$ . Let  $i < t - 1$ , assuming that all  $X_{i+1}$ -configurations are marked by  $\text{Mark}(X_{i+1})$ , we prove that  $\text{Mark}(X_i)$  marks all  $X_i$ -configurations.

By definition of the  $k$ -bidismantling ordering, there exist two adjacent or coinciding vertices  $x, y \in X_{i+1}$  such that  $N_k(v_i, G \setminus \{x, y\}) \cap X_i \subseteq N_1(y)$ . Roughly speaking,  $\text{Mark}(X_i)$  marks the  $X_i$ -configurations in the same order as  $\text{Mark}(X_{i+1})$  marks the  $X_{i+1}$ -configurations, but once a configuration  $(c, y)$  ( $c \in X_{i+1}$ ) is marked,  $\text{Mark}(X_i)$  also marks the configuration  $(c, v_i)$ . Once  $\text{Mark}(X_i)$  has marked all  $X_i$ -configurations  $(c, r) \in X_{i+1} \times X_i$ , the remaining  $X_i$ -configurations  $(v_i, r)$ ,  $r \in X_{i+1}$ , can also be marked by  $\text{Mark}(X_i)$ .

Let  $\ell \geq 1$ . By induction on  $\ell$ , we prove that *any  $X_{i+1}$ -configurations  $(c, r)$  that is marked by  $\text{Mark}(X_{i+1})$  with label at most  $\ell$  will be also marked by  $\text{Mark}(X_i)$* . Moreover, if  $r = y$ , we prove that *once  $\text{Mark}(X_i)$  has marked  $(c, r)$ , then it can mark  $(c, v_i)$* . Let us first prove this assertion for  $\ell = 1$ . For any  $(c, r) \in X_i \times X_i$  with  $r \in N_1(c)$ ,  $(c, r)$  is marked by  $\text{Mark}(X_i)$  with label 1. If  $(c, y)$  is marked with label 1 (i.e.,  $y \in N_1(c) \cap X_i$ ), then  $(c, v_i)$  can be marked with 2. Indeed, for all  $z \in N_k(v_i, G \setminus \{x, y\}) \cap X_i$ , we have  $z \in N_1(y)$  (by definition of the  $k$ -bidismantling order), and thus the  $X_i$ -configuration  $(y, z)$  is marked with label 1. Hence, by setting  $(x_{(c,v_i)}, y_{(c,v_i)}) = (x, y)$ , the procedure  $\text{Mark}(X_i)$  marks  $(c, v_i)$  with label 2.

Assume now that the induction hypothesis holds for some  $\ell \geq 1$  and we will show that it still holds for  $\ell + 1$ . Let  $(c, r)$  be a  $X_{i+1}$ -configuration marked by  $\text{Mark}(X_{i+1})$  with label  $\ell + 1$ . We first prove that  $(c, r)$  is eventually marked by  $\text{Mark}(X_i)$ . By definition of  $\text{Mark}(X_{i+1})$ , there exist  $y_{(c,r)} \in N_1(c) \cap X_{i+1}$  and  $x_{(c,r)} \in (N_1(y_{(c,r)}) \setminus \{r\}) \cap X_{i+1}$  such that for all  $z \in N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\}) \cap X_{i+1}$ , the  $X_{i+1}$ -configuration  $(y_{(c,r)}, z)$  is marked with label at most  $\ell$  by  $\text{Mark}(X_{i+1})$ . By the induction hypothesis, this implies that for all  $z \in N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\}) \cap X_{i+1}$ , the  $X_{i+1}$ -configuration  $(y_{(c,r)}, z)$  is marked by  $\text{Mark}(X_i)$ . If  $v_i \notin N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\})$ , then clearly  $(c, r)$  is marked by  $\text{Mark}(X_i)$ . Let us assume that  $v_i \in N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\})$ . We aim at proving that  $(y_{(c,r)}, v_i)$  is eventually marked by  $\text{Mark}(X_i)$ . We distinguish three cases.

- If  $y_{(c,r)} = y$ , then  $(y_{(c,r)}, v_i)$  is marked with label 1 since  $y_{(c,r)} = y \in N_1(v_i)$ .
- If  $x_{(c,r)} = y$ , then  $(y_{(c,r)}, v_i)$  is marked with label 1 or 2 by setting  $(x_{(y_{(c,r)}, v_i)}, y_{(y_{(c,r)}, v_i)}) = (x, y)$ . Indeed, for all  $z \in N_k(v_i, G \setminus \{x, y\}) \cap X_i$ , we have  $z \in N_1(y)$  (by definition of the  $k$ -bidismantling order), and thus the  $X_i$ -configuration  $(y, z)$  is marked with label 1.
- Otherwise, we assert that  $(y_{(c,r)}, y)$  has already been marked by  $\text{Mark}(X_i)$ . By the induction hypothesis, this implies that  $(y_{(c,r)}, v_i)$  was also marked.

If  $y \in N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\}) \cap X_{i+1}$  and since  $(c, r)$  is marked with label  $\ell + 1$  by the marking procedure in  $X_{i+1}$ , then  $(y_{(c,r)}, y)$  must be marked by  $\text{Mark}(X_{i+1})$  with label at most  $\ell$ . By the induction hypothesis, this implies that  $(y_{(c,r)}, y)$  has been marked by  $\text{Mark}(X_i)$ . Hence, it remains to show that  $y \in N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\}) \cap X_{i+1}$ .

Let  $P$  be a path of length at most  $k$  between  $r$  and  $v_i$  in  $G \setminus \{x_{(c,r)}, y_{(c,r)}\}$ . If  $x$  or  $y$  belongs to  $P$ , then we trivially get that  $y \in N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\}) \cap X_{i+1}$ . Otherwise, this means that  $r \in N_k(v_i, G \setminus \{x, y\}) \cap X_i$  and  $r \in N_1(y)$  holds by definition of the bidismantling order. Hence,  $y \in N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\}) \cap X_{i+1}$ .

In all three cases, the pair  $(y_{(c,r)}, v_i)$  is marked by  $\text{Mark}(X_i)$ . Thus, for all  $z \in N_k(r, G \setminus \{x_{(c,r)}, y_{(c,r)}\}) \cap X_i$ , the  $X_i$ -configuration  $(y_{(c,r)}, z)$  has been marked. Therefore, this is also the case for the  $X_i$ -configuration  $(c, r)$ . To conclude the proof, we need to show that, once a  $X_i$ -configuration  $(c, y)$  ( $c \neq v_i$ ) is marked by  $\text{Mark}(X_i)$ , then  $(c, v_i)$  can be marked as well. Since  $(c, y)$  has been marked, there exist  $y_{(c,y)} \in N_1(c) \cap X_i$  and  $x_{(c,y)} \in (N_1(y_{(c,y)}) \setminus \{y\}) \cap X_i$  such that for all  $z \in N_k(y, G \setminus \{x_{(c,y)}, y_{(c,y)}\}) \cap X_i$ , the  $X_i$ -configuration  $(y_{(c,y)}, z)$  is marked. Let  $z' \in N_k(v_i, G \setminus \{x_{(c,y)}, y_{(c,y)}\}) \cap X_i$ . We prove that  $z' \in N_k(y, G \setminus \{x_{(c,y)}, y_{(c,y)}\}) \cap X_i$ , which shows that  $(y_{(c,y)}, z')$  has been already marked. Let  $P$  be a shortest path between  $v_i$  and  $z'$  in  $G \setminus \{x_{(c,y)}, y_{(c,y)}\}$ . Note that  $|P| \leq k$ . If  $y \in P$ , clearly  $z' \in N_k(y, G \setminus \{x_{(c,y)}, y_{(c,y)}\}) \cap X_i$ . Else, if  $x \in P$ , then let  $P'$  be the subpath of  $P$  from  $z'$  to  $x$ . Then  $P' \cup \{x, y\}$  is a path of length at most  $k$  between  $z'$  and  $y$  in the graph  $G \setminus \{x_{(c,y)}, y_{(c,y)}\}$ . Otherwise,  $z' \in N_k(v_i, G \setminus \{x, y\}) \cap X_i$  and thus  $z' \in N_1(y)$ . Therefore, for any  $z' \in N_k(v_i, G \setminus \{x_{(c,y)}, y_{(c,y)}\}) \cap X_i$ ,  $(y_{(c,y)}, z')$  is marked and thus the pair  $(c, v_i)$  can be marked as well.

Summarizing, for all  $c, r \in X_{i+1}$ , the configurations  $(c, r)$  and  $(c, v_i)$  are marked by the procedure  $\text{Mark}(X_i)$ . To conclude the proof, note that any configuration  $(v_i, r)$  can be marked as well: either with 1 if  $r \in N_1(v_i)$  or by setting  $(x_{(v_i, r)}, y_{(v_i, r)}) = (y, y)$  otherwise.  $\square$

From Theorem 4 and by noticing that if a graph  $G = (V, E)$  with  $n$  vertices is  $n$ -bidismantlable, then there are two vertices  $x, y$  such that  $y$  dominates a connected component of  $G \setminus \{x, y\}$ , we obtain the following observation:

**Proposition 15.** *CWW is the class of graphs which are  $k$ -bidismantlable for all  $k \geq 1$ .*

We continue with an example showing that  $k$ -bidismantlability is not a necessary condition for any  $k \geq 3$ .

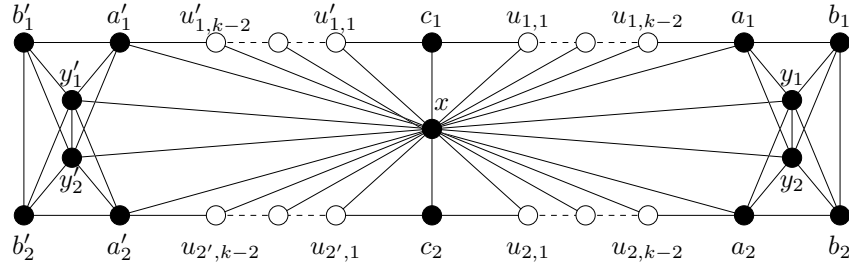


Figure 7: A graph  $G \in \text{CWW}(k)$  that is not  $k$ -bidismantlable

**Proposition 16.** *Let  $G$  be the graph from Fig. 7. Then  $G \in \text{CWW}(k)$ , however  $G$  is not  $k$ -bidismantlable.*

*Proof.* To show that  $G \in \text{CWW}(k)$ , we exhibit a strategy for  $\mathcal{C}$ . The cop starts at the vertex  $x$ . To avoid being captured immediately, the robber starts in  $b_1, b_2, b'_1$ , or  $b'_2$ , say in  $b_1$  (the other cases are similar). Then the cop moves to  $a_1$ , goes to  $x$  and stays there during  $k - 2$  steps, and finally goes to  $y_1$ . Once  $\mathcal{C}$  is in  $a_1$ ,  $\mathcal{R}$  can go to  $b_1, b_2, y_1$  or  $y_2$ . Then, while  $\mathcal{C}$  is in  $x$ ,  $\mathcal{R}$  can move to a vertex in  $\{a_1, a_2, b_1, b_2, y_1, y_2\} \cup \{u_{1,i}, u_{2,i} : 2 \leq i \leq k - 2\}$ . Finally, when the cop moves to  $y_1$ ,  $\mathcal{R}$  can go to a vertex in  $\{a_1, a_2, b_1, b_2, y_2, x\} \cup \{u_{1,i}, u_{2,i} : i \leq k - 2\}$ . Thus, if the robber is not caught the second time he shows up, he is in a vertex of  $\{u_{1,i}, u_{2,i} : i \leq k - 2\}$ . If the robber shows up in a vertex of  $\{u_{1,i} : i \leq k - 2\}$  while the cop is in  $y_1$  (the other case is similar), the cop oscillates between  $a_1$  and  $x$  during  $k$  steps and finishes in  $x$  (if  $k$  is odd,  $\mathcal{C}$  first moves to  $x$  and if  $k$  is even,  $\mathcal{C}$  first moves to  $a_1$ ). Thus, the next time  $\mathcal{R}$  shows up, he is in a vertex of  $\{c_1, a'_1\} \cup \{u_{1,i}, u'_{1,i} : i \leq k - 2\}$ : in any case,  $\mathcal{C}$  (positioned at  $x$ ) catches  $\mathcal{R}$  when he shows up.

Now we show that  $G$  is not  $k$ -bidismantlable. We can eliminate a vertex  $v$  if there exist two neighbors  $x(v), y(v)$  such that  $N_k(v, G \setminus \{x, y\}) \cap X \subseteq N_1(y)$ , where  $X$  is the set of vertices that have not been yet eliminated. First note that for any  $i \leq k - 2$ , we can

eliminate  $u_{1,i}$  with  $x(u_{1,i}) = a_{1,i}$  and  $y(u_{1,i}) = x$ . By symmetry, we can also eliminate  $u_{2,i}$ ,  $u'_{1,i}$  and  $u'_{2,i}$  for any  $i \leq k-2$  (these vertices are colored white in Fig. 7). We show that no other vertex can be eliminated after that a set  $Y \subseteq \{u_{1,i}, u'_{1,i}, u_{2,i}, u'_{2,i} : i \leq k-2\}$  of vertices has been eliminated ( $Y$  can be empty or contain all these vertices).

Let  $X = V(G) \setminus \{u_{1,i}, u_{2,i}, u'_{1,i}, u'_{2,i} : i \leq k-2\}$ . By symmetry, it is sufficient to show that for any  $v \in \{a_1, b_1, c_1, x, y_1\}$ , and any adjacent vertices  $x(v), y(v) \in V(G)$ , there exists  $z(v) \in (N_k(v, G \setminus \{x(v), y(v)\}) \cap X) \setminus N_1(y(v))$ . For any  $v$  and  $y(v)$ , this condition is true as soon as  $N_1(v) \cap X \not\subseteq N_1(y(v))$ .

If  $v = x$ , then there does not exist any  $y(v)$  such that  $N_1(v) \cap X \subseteq N_1(y(v))$ . If  $v = a_1$  and  $N_1(v) \cap X \subseteq N_1(y(v))$ , then  $y(v) \in \{y_1, y_2\}$ ; in both cases,  $x(v) \notin \{u_{1,i} : i \leq k-2\}$  and thus  $z(v) = c_1$  satisfies the condition. If  $v = b_1$  and  $N_1(v) \cap X \subseteq N_1(y(v))$  then  $y(v) \in \{y_1, y_2\}$ ; by symmetry, we assume  $y(b_1) = y_1$ . Note that there exist two vertex-disjoint paths  $(b_1, a_1, u_{1,k-2}, \dots, u_{1,1}, c_1)$  and  $(b_1, y_2, x, c_1)$  of length at most  $k$  from  $b_1$  to  $c_1$  avoiding  $y_1$ . Consequently, for any choice of  $x(b_1)$ , the vertex  $z(v) = c_1$  satisfies the condition. If  $v = y_1$  and  $N_1(v) \cap X \subseteq N_1(y(v))$ , then  $y(v) = y_2$ . Again, there exist two vertex-disjoint paths  $(y_1, a_1, u_{1,k-2}, \dots, u_{1,1}, c_1)$  and  $(y_1, x, c_1)$  of length at most  $k$  from  $y_1$  to  $c_1$  avoiding  $y_2$ . Hence, for any choice of  $x(y_1)$ , the vertex  $z(v) = c_1$  satisfies the condition. Finally, suppose that  $v = c_1$  and  $N_1(v) \cap X \subseteq N_1(y(v))$  then  $y(v) \in \{u_{1,1}, u'_{1,1}, x\}$ . If  $y(v) = u_{1,1}$  (the case  $y(v) = u'_{1,1}$  is similar), then  $x(v) \notin \{a'_1\} \cup \{u'_{1,i} : i \leq k-2\}$  and thus  $z(v) = b_1$  satisfies the condition. If  $y(v) = x$ , then either  $x(v) \notin \{a'_1\} \cup \{u'_{1,i} : i \leq k-2\}$ , or  $x(v) \notin \{a_1\} \cup \{u_{1,i} : i \leq k-2\}$ . By symmetry, we assume  $x(v) \notin \{a'_1\} \cup \{u'_{1,i} : i \leq k-2\}$ ; in this case,  $z(v) = b_1$  satisfies the condition.  $\square$

**Open question 3:** Characterize the  $k$ -winnable graphs for  $k = 2, 3$  and, more generally, for all  $k$ .

### 5.3 Cop-win graphs for game with fast robber and witness.

We now consider a variant of the game where the robber is visible every  $k$  moves and has speed  $s$  while the cop has speed 1. It means that at each step, the robber can move to a vertex at distance at most  $s$  from his current position, and that the cop can see the robber only every  $k$  steps. We denote by  $CWFRW(s, k)$  the class of graphs where a single cop with speed 1 can catch a robber with speed  $s$  that is visible every  $k$  moves. By definition, we have  $CWFRW(1, k) = CWW(k)$  and  $CWFRW(s, 1) = CWFR(s)$ .

**Theorem 6.** *If  $s \geq 3, k \geq 1$  or  $s \geq 2, k \geq 2$ , then  $CWFRW(s, k)$  is the class of big brother graphs.*

*Proof.* We know from Theorem 3 that if  $s \geq 3$  and  $k \geq 1$ ,  $CWFR(sk) = CWFR(s)$  is the class of big brother graphs. Consequently, since  $CWFR(sk) \subseteq CWFRW(s, k) \subseteq CWFR(s)$ ,  $CWFRW(s, k)$  is the class of big brother graphs for all  $s \geq 3$  and  $k \geq 1$ .

In the remaining of this proof, we show that when  $s = 2$  and  $k \geq 2$ ,  $CWFRW(2, k)$  also coincides with the class of big brother graphs. This proof follows closely the proof of Theorem 4. In particular, the following proposition is the counterpart of Proposition 9.



**Proposition 17.** *Let  $G \in \mathcal{CWFRW}(2, k)$  for  $k \geq 2$ . If the minimum degree of  $G$  is at least 2, then  $G$  contains two vertices  $v$  and  $y$  such that  $N_{2k}(v, G \setminus \{y\}) \subseteq N_1(y)$ .*

*Proof.* If  $G$  contains a dominating vertex  $y$ , then the result holds for any  $v \neq y$ . Assume thus that  $G$  has no dominating vertex. Consider a parsimonious winning strategy of the cop and suppose that the robber uses a strategy to avoid being captured as long as possible. Since  $G$  does not contain leaves, the robber is caught immediately after having been visible. Since  $G$  does not have any dominating vertex, the robber is visible at least twice. Let  $y$  be the vertex occupied by the cop when the robber becomes visible for the last time before his capture. Let  $v$  be the next-to-last visible vertex occupied by the robber. Finally, let  $S_c = (c_0, c_1, \dots, c_k = y)$  be the trajectory of the cop during the last  $k$  steps (repetitions are allowed). Note that  $v \notin N_1(c_0)$ , otherwise the robber would have been caught immediately. We distinguish two cases depending of whether or not  $c_i = y$  for all  $1 \leq i \leq k$ .

If for any  $1 \leq i \leq k$ , we have  $c_i = y$ , then  $\mathcal{R}$  could have move to any vertex  $w \in N_{2k}(v, G \setminus \{y\})$ . Since the trajectory of the robber is maximal, the robber is caught in any such  $w$  and thus  $N_{2k}(v, G \setminus \{y\}) \subseteq N_1(y)$ . Suppose now that there exists  $i$  such that  $c_i \neq y$ . Let  $i$  be the largest index such that  $c_i \neq y$ .

**Claim 4.** *If  $G$  contains a cycle  $C$  and a vertex  $w \in C$  such that  $d(v, w) < d(c_1, w)$ , then  $G \setminus \{y\}$  has a connected component that is dominated by  $y$ .*

*Proof.* Let  $w$  be the closest to  $v$  vertex satisfying the condition of the claim. If the assertion of the claim is not satisfied, we will exhibit a strategy allowing the robber to escape the cop during more steps, contradicting the choice of the trajectory  $\mathcal{R}$ . Suppose that at the beginning of the last phase, the robber moves from  $v$  to  $w$  along a shortest  $(v, w)$ -path. Since  $d(v, w) < d(c_1, w)$ , the robber cannot be intercepted by the cop during these moves. Suppose that the robber reaches the vertex  $w$  before the  $i$ th step when the cop arrives at  $c_i$ . Then by Lemma 1 (adapted to this game) the robber can safely move on  $C$  until the cop reaches the vertex  $c_i$ .

In both cases, let  $z$  be the current position of the robber when the cop reaches  $c_i$ . Then  $z \in N_1(y)$ , otherwise the robber can remain at  $z$  without being caught because starting with this step the cop remains in  $y$ . If  $z \neq y$ , then let  $u = z$ , and if  $z = y$ , let  $u$  be a neighbor of  $y$  different from  $c_i$  (it exists because the minimum degree of a vertex of  $G$  is at least 2). In both cases, let  $H$  be the connected component of  $G \setminus \{y\}$  that contains  $u$ . We assert that  $y$  dominates all the vertices of  $H$ . Suppose this is not the case and consider a vertex  $t$  in  $V(H) \setminus N(y)$  that is at a minimum distance from  $u$  in  $H$ . From our choice of  $t$ , we can find a common neighbor  $r \in V(H)$  of  $y$  and  $t$ . If  $r \neq c_i$ , then while the cop is in  $c_i$ , the robber can go from  $z$  ( $z$  is either  $u$  or  $y$ ) to  $r$  through  $y$  and then, when  $\mathcal{C}$  goes to  $y$ ,  $\mathcal{R}$  goes to  $t$  and stays there until he becomes visible. If  $r = c_i$ , then let  $s$  be a neighbor of  $r$  on a shortest  $(u, r)$ -path in  $G \setminus \{y\}$ . By our choice of  $t$ , necessarily  $s \in N(y)$ . Thus, when the cop is in  $c_i$ , the robber can go from  $z$  to  $s$  through  $y$ . And then, when the cop goes to  $y$ ,  $\mathcal{R}$  goes to  $t$  through  $r$  and stays there until he becomes visible. In both cases, by following such a strategy  $\mathcal{R}$  could avoid being caught. This contradicts the maximality of the trajectory of the robber. This concludes the proof of the claim.  $\square$



We now complete the proof of Proposition 17. If the vertex  $v$  belongs to a cycle  $C$ , then setting  $w := v$  and applying Claim 4, we conclude that  $y$  dominates a non-empty connected component  $H$  of  $G \setminus \{y\}$  establishing thus the assertion. So, suppose that  $v$  is an articulation point of  $G$  not contained in a cycle. Since the minimum degree of  $G$  is at least 2,  $G \setminus \{v\}$  has a connected component  $H$  that does not contain  $c_0$  (nor  $c_1$ ). Necessarily  $H$  contains a cycle  $C$ , otherwise we will find in  $H$  a vertex of degree 1 in  $G$ . Since any path from  $c_1$  to a vertex  $w$  of  $C$  goes through  $v$ , we obtain  $d(v, w) < d(c_1, w)$ . Then, the result again follows from Claim 4. This ends the proof of Proposition 17.  $\square$

Finally, we prove Theorem 6 when  $s = 2$  and  $k \geq 2$ . Consider a graph  $G \in \mathcal{CWF}\mathcal{RW}(2, k)$ . To establish that  $G$  is a big brother graph, in view of Theorem 3, it suffices to show that  $G$  is  $(2k, 1)$ -dismantlable. For this, by Proposition 6, we just have to show that there exists an ordering  $v_1, \dots, v_n$  of the vertices of  $G$  such that for each  $1 \leq i < n$  there exists  $y \in \{v_{i+1}, \dots, v_n\}$  such that  $N_{2k}(v_i, G_i \setminus \{y\}) \subseteq N_1(y, G_i)$ . We proceed by induction on the size of  $G$ . Suppose that  $G$  has at least two vertices, otherwise the result is trivial. If  $G$  has a vertex  $v$  of degree 1, then let  $y$  be the unique neighbor of  $v$  in  $G$ . In this case, then obviously  $N_{2k}(v, G \setminus \{y\}) \subseteq N_1(y, G)$ . Otherwise, by Proposition 17, we can find vertices  $v$  and  $y$  such that  $N_{2k}(v, G \setminus \{y\}) \subseteq N_1(y, G)$ .

We now show that  $G' = G \setminus \{v_1\}$  also belongs to  $\mathcal{CWF}\mathcal{RW}(2, k)$ . Consider a winning positional strategy  $\sigma$  for the cop in  $G$ . As in the proof of Theorem 1, we define a strategy  $\sigma'$  for the cop in  $G'$  using one bit of memory. Starting from  $\sigma$ , we define  $\sigma'(c, r, m)$  for any positions  $c, r \in V(G')$  of the cop and the robber and for any value of the memory  $m \in \{0, 1\}$ . The idea is that the cop plays using  $\sigma$  except when he is in  $y$  and his memory contains 1; in this case he uses  $\sigma$  as if he was in  $v$  (going via  $y$  instead of  $v$  if  $v$  appears in his sequence of moves).

If  $m = 0$  or  $c \neq y$ , let  $\sigma(c, r) = (c_1, \dots, c_k)$  and for each  $i$ , let  $c'_i = c_i$  if  $c_i \neq v$  and  $c'_i = y$  otherwise (this is possible since  $N_1(v) \subseteq N_1(y)$ ). If  $c_k = v$ , then  $\sigma'(c, r, m) = ((c'_1, \dots, c'_{k-1}, y), 1)$ , otherwise let  $\sigma'(c, r, m) = ((c'_1, \dots, c'_{k-1}, c_k), 0)$ .

If  $m = 1$  and  $c = y$ , let  $\sigma(v, r) = (c_1, \dots, c_k)$  and for each  $i$ , let  $c'_i = c_i$  if  $c_i \neq v$  and  $c'_i = y$  otherwise. If  $c_k = v$ , then  $\sigma'(y, r, 1) = ((c'_1, \dots, c'_{k-1}, y), 1)$ , otherwise let  $\sigma'(y, r, 1) = ((c'_1, \dots, c'_{k-1}, c_k), 0)$ .

Let  $S_r = (r_1, r_2, \dots, r_p, \dots)$  be a valid sequence of moves in  $G'$ . Since  $V(G') \subseteq V(G)$ ,  $S_r$  is also a valid sequence of moves in  $G$ . Let  $S_c = (c_1, \dots, c_p, \dots)$  be the corresponding valid sequence of moves of the cop playing  $\sigma$  against  $S_r$  in  $G$  and let  $S'_c = (c'_1, \dots, c'_p, \dots)$  be the valid sequence of moves of the cop playing  $\sigma'$  against  $S_r$  in  $G'$ . Note that the sequences of moves  $S_c$  and  $S'_c$  differ only if  $c_k = v$  and  $c'_k = y$ . Finally, since the cop follows a winning strategy for  $X$ , there exists a step  $j$  such that  $c_j = r_j \in V(G')$  (note that  $r_j \neq v$  because we supposed that  $S_r$  is a valid sequence of moves in  $G'$ ). Since  $c_j \neq v$ , we also have  $c'_j = r_j$ , thus  $\mathcal{C}$  captures  $\mathcal{R}$  in the game restricted to  $G'$ . In conclusion, starting from a positional strategy for the game in  $G$ , we have constructed a winning strategy using memory for the game in  $G'$ . As mentioned in the introduction, it implies that there exists a positional winning strategy for the game in  $G'$ . Consequently,  $G' \in \mathcal{CWF}\mathcal{RW}(2, k)$  and by induction hypothesis,  $G'$  has

$(2k, 1)$ -dismantling order  $(v_2, \dots, v_n)$ , whence  $(v, v_2, \dots, v_n)$  is a  $(2k, 1)$ -dismantling order of  $G$ . This concludes the proof of Theorem 6.  $\square$

## 6 Bipartite cop-win graphs for game with “radius of capture”

In this section we characterize bipartite graphs of the class  $CWRC(1)$ , i.e., the bipartite cop-win graphs in the cop and robber game with radius of capture 1. Recall that in this game introduced in [11],  $\mathcal{C}$  and  $\mathcal{R}$  move at unit speed and the cop wins if after his move he is within distance at most 1 from the robber. Notice that any graph of diameter 2 belongs to  $CWRC(1)$ : given the positions  $u$  and  $v$  of the cop and the robber, to capture the robber the cop simply moves from  $u$  to a common neighbor of  $u$  and  $v$ .

Following [7], a bipartite graph  $G$  is called *dismantlable* if its vertices can be ordered  $v_1, \dots, v_n$  so that  $v_{n-1}v_n$  is an edge of  $G$  and for each  $v_i, i < n-1$ , there exists a vertex  $y := v_j$  with  $j > i$  (necessarily not adjacent to  $v_i$ ) such that  $N(v_i, G_i) := N_1(v_i, G_i) \setminus \{v_i\} \subseteq N_1(y)$ . Note that for any  $i$ ,  $G_i$  is a retract of  $G$  and therefore an isometric subgraph of  $G$ .

**Theorem 7.** *A bipartite graph  $G$  belongs to  $CWRC(1)$  if and only if  $G$  is dismantlable.*

*Proof.* First suppose that  $G \in CWRC(1)$ . If  $G$  has diameter 2, then necessarily  $G$  is a complete bipartite graph, which is obviously dismantlable. Suppose now that  $G$  has diameter at least 3. As in previous proofs of similar results, we assume that  $\mathcal{C}$  uses a parsimonious strategy. Consider a maximal sequence of moves of the robber before he get caught. Let  $v$  be the next-to-last position of the robber and let  $y$  be the position of the cop at this step ( $y$  is not adjacent to  $v$ , otherwise  $\mathcal{C}$  captures  $\mathcal{R}$  in  $v$ ). This means that for any  $w \in N_1(v)$ , the cop can move in some vertex  $u \in N_1(y)$  such that  $w \in N_1(u)$ . This shows that  $N_1(v) \subseteq N_2(y)$ . Since  $G$  is bipartite, this means that  $d(v, y) = 2$  and all neighbors of  $v$  are adjacent to  $y$ , i.e.,  $N(v) \subseteq N(y)$ .

We now show that  $G' = G \setminus \{v\}$  also belongs to  $CWRC(1)$ . Consider a winning positional strategy  $\sigma$  for the cop in  $G$ . As in the proof of Theorem 1, we define a strategy  $\sigma'$  for the cop in  $G'$  using one bit of memory. Starting from  $\sigma$ , we define  $\sigma'(c, r, m)$  for any positions  $c, r \in V(G')$  of the cop and the robber and for any value of the memory  $m \in \{0, 1\}$ . The idea is that the cop plays using  $\sigma$  except when he is in  $y$  and his memory contains 1; in this case he uses  $\sigma$  as if he was in  $v$  (going to  $y$  instead of  $v$ ). If  $m = 0$  or  $c \neq y$ , if  $\sigma(c, r) = v$  then  $\sigma'(c, r, m) := (y, 1)$  (this is a valid move since  $N(v) \subseteq N(y)$  and  $c \neq v$ ), otherwise let  $\sigma'(c, r, m) := (\sigma(c, r), 0)$ . If  $m = 1$  and  $c = y$ , if  $\sigma(v, r) := v$ , then  $\sigma'(y, r, 1) = (y, 1)$ , otherwise let  $\sigma'(y, r, 1) := (\sigma(v, r), 0)$ .

Let  $S_r = (r_1, r_2, \dots, r_p, \dots)$  be a valid sequence of moves in  $G'$ . Since  $V(G') \subseteq V(G)$ ,  $S_r$  is also a valid sequence of moves in  $G$ . Let  $S_c = (c_1, \dots, c_p, \dots)$  be the corresponding valid sequence of moves of the cop playing  $\sigma$  against  $S_r$  in  $G$  and let  $S'_c = (c'_1, \dots, c'_p, \dots)$  be the valid sequence of moves of the cop playing  $\sigma'$  against  $S_r$  in  $G'$ . Note that the sequences of moves  $S_c$  and  $S'_c$  differ only if  $c_k = v$  and  $c'_k = y$ . Finally, since the cop follows a winning

strategy for  $X$ , there exists a step  $j$  such that  $c_j \in N_1(r_j) \in V(G')$  (note that  $r_j \neq v$  because we supposed that  $S_r$  is a valid sequence of moves in  $G'$ ). Since  $N(v) = N(y)$ , we also have  $c'_j \in N_1(r_j)$ , thus  $\mathcal{C}$  captures  $\mathcal{R}$  in the game restricted to  $G'$ . In conclusion, starting from a positional strategy for the game in  $G$ , we have constructed a winning strategy using memory for the game in  $G'$ . As mentioned in the introduction, it implies that there exists a positional winning strategy for the game in  $G'$ . Consequently,  $G' \in \mathcal{CWRC}(1)$  and by induction hypothesis,  $G'$  has a dismantling order  $(v_2, \dots, v_n)$ , whence  $(v, v_2, \dots, v_n)$  is a dismantling order of  $G$ .

Conversely, suppose that a bipartite graph  $G$  is dismantlable and let  $v = v_1, v_2, \dots, v_n$  be a dismantling order of  $G$ . If  $G$  has diameter 2, then  $G$  is a complete bipartite graph and thus,  $G \in \mathcal{CWRC}(1)$ . Suppose now that  $G$  has a diameter at least 3. By induction hypothesis,  $G' = G(\{v_2, \dots, v_n\})$  belongs to  $\mathcal{CWRC}(1)$ . Suppose that  $v$  is dominated by a vertex  $y$  at distance 2 from  $v$ .

Consider a parsimonious positional winning strategy  $\sigma'$  for the cop in  $G'$ . Using  $\sigma'$ , we build a parsimonious positional winning strategy  $\sigma$  for the cop in  $G$ . As in the previous proofs, the idea is that if  $\mathcal{C}$  sees  $\mathcal{R}$  in  $v$ , he plays as in the game on  $G'$  when the cop is in  $y$ . For any positions  $c, r \in V(G)$ , if  $r \in N_2(c)$ , then  $\sigma(c, r) := u \in N_1(c) \cap N_1(r)$ . Otherwise, if  $c \in N(y) \setminus N(v)$  and  $r = v$ , then  $\sigma(c, v) := y$ . Otherwise, if  $c, r \neq v$ , then  $\sigma(c, r) := \sigma'(c, r)$ ; if  $r = v$  and  $c \notin N_2(v)$ , then  $\sigma(c, v) := \sigma'(c, y)$ ; finally, if  $c = v$  and  $r \notin N_2(v)$ , then  $\sigma(v, r) := u \in N(v)$  (in fact, if the cop plays according to  $\sigma$ , he will never move to  $v$ , except to catch the robber). By construction,  $\sigma$  is parsimonious and positional; in particular,  $\sigma(y, v) \in N(v)$ .

We now prove that  $\sigma$  is a winning strategy. Consider any valid sequence  $S_r = (r_1, \dots, r_k, \dots)$  of moves of the robber in  $G$ , and let  $S'_r = (r'_1, \dots, r'_k, \dots)$  be the sequence obtained by setting  $r'_k = r_k$  if  $r_k \neq v$  and  $r'_k = y$  if  $r_k = v$ . Since  $N(v) \subseteq N(y)$ ,  $S'_r$  is a valid sequence of moves for the robber in  $G'$ . By induction hypothesis, for any initial position of  $\mathcal{C}$  in  $G'$ , the strategy  $\sigma'$  enables  $\mathcal{C}$ , following a trajectory  $S'_c = (c'_1, \dots, c'_k, \dots)$ , to catch  $\mathcal{R}$  which moves according to  $S'_r$ , i.e., there exists an index  $m$  such that  $r'_m \in N_1(c'_m)$ . Suppose that  $\mathcal{C}$  chooses his starting position  $c_1$  in  $G'$  and that the cop, following  $\sigma$ , plays in  $G$  the sequence  $S_c = (c_1, \dots, c_k, \dots)$  against the sequence  $S_r$  of the robber. Note that  $c_k = c'_k$  for any  $k < m$ , i.e., except when the robber is caught in  $G'$ . If  $r_m = r'_m$ , then  $r_m \in N_1(c_{m+1})$  and  $\mathcal{C}$  captures  $\mathcal{R}$ . Now suppose that  $r_m \neq r'_m$ . From the definition of  $S'_r$  we conclude that  $r'_m = y$  and  $r_m = v$ . If  $v \in N(c_{m+1})$ , the robber is caught at step  $m + 1$ . Otherwise, since  $c_{m+1} \in N(y) \setminus N(v)$ , then  $c_{m+2} = y$ . If at step  $r_{m+1}$  the robber moves to a neighbor of  $v$ , since  $N(v) \subseteq N(y)$ , we conclude that  $r_{m+1} \in N(c_{m+2})$  and the robber is caught. Finally, if the robber remains at  $v$  (i.e.,  $r_{m+1} = v$ ), then to avoid being caught while moving, at step  $m + 2$   $\mathcal{R}$  must also stay in  $v$  and then at step  $m + 3$  the cop moves from  $y$  to any common neighbor of  $y$  and  $v$  and catches the robber.  $\square$

Bonato and Chiniforooshan [11] asked for a characterization of graphs of  $\mathcal{CWRC}(1)$ . Theorem 7 answers the question in the case of bipartite graphs. On the other hand, characterizing the graphs of  $\mathcal{CWRC}(1)$  using a specific dismantling scheme seems to be quite

challenging. Two natural candidates for us were the total orders  $v_1, \dots, v_n$  satisfying the following conditions:

- (i) for each  $i$ , there exists  $j > i$  such that  $N_1(v_i, G_i) \subseteq N_2(v_j, G_{i+1})$ ;
- (ii) for each  $i$ , there exists  $j > i$  such that  $N_2(v_i, G_i) \subseteq N_2(v_j, G_i)$ .

It seems that the first condition is necessary, while the second condition is sufficient. However, we were not able to prove this. In fact, one can easily show that any graph  $G \in \mathcal{CWR}\mathcal{C}(1)$  contains two vertices  $v, y$  such that  $N_1(v, G_i) \subseteq N_2(y, G_{i+1})$ , but we cannot show that  $G \setminus \{v\}$  also belongs to  $\mathcal{CWR}\mathcal{C}(1)$ . A similar difficulty occurs while establishing the sufficiency of the second dismantling order.

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