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# Algorithms and Arithmetic Operators for Computing the $\eta_T$ Pairing in Characteristic Three – Appendices

Jean-Luc Beuchat, Nicolas Brisebarre, Jérémie Detrey, Eiji Okamoto, Masaaki Shirase, and Tsuyoshi Takagi

We describe here how to implement the arithmetic operations over  $\mathbb{F}_{3^{2m}}$ ,  $\mathbb{F}_{3^{3m}}$ , and  $\mathbb{F}_{3^{6m}}$  involved in the  $\eta_T$  pairing calculation. In order to compute the number of operations over  $\mathbb{F}_{3^m}$ , we assume that the ALU is able to compute  $u \cdot v$ ,  $\pm u \pm v$  and  $\pm u^3$ , where  $u$  and  $v \in \mathbb{F}_{3^m}$ .

## I. MULTIPLICATION OVER $\mathbb{F}_{3^{2m}}$

Let  $U = u_0 + u_1\sigma$  and  $V = v_0 + v_1\sigma$ , where  $u_0, u_1, v_0$ , and  $v_1 \in \mathbb{F}_{3^m}$ . The product  $UV$  is carried out according to Karatsuba-Ofman's algorithm:

$$U \cdot V = (u_0v_0 - u_1v_1) + ((u_0 + u_1)(v_0 + v_1) - u_0v_0 - u_1v_1)\sigma.$$

It requires 3 multiplications and 5 additions over  $\mathbb{F}_{3^m}$ .

## II. MULTIPLICATION OVER $\mathbb{F}_{3^{3m}}$

Assume that  $U = u_0 + u_1\rho + u_2\rho^2$  and  $V = v_0 + v_1\rho + v_2\rho^2$ , where  $u_i, v_i \in \mathbb{F}_{3^m}$ ,  $0 \leq i \leq 2$ . The product  $W = U \cdot V$  is then given by

$$\begin{aligned} w_0 &= b(bu_1 + u_2)(v_1 + bv_2) + u_0v_0 - u_1v_1 - u_2v_2, \\ w_1 &= (u_0 + u_1)(v_0 + v_1) + (bu_1 + u_2)(v_1 + bv_2) \\ &\quad - u_0v_0 - (b+1)u_1v_1, \text{ and} \\ w_2 &= (u_0 + u_2)(v_0 + v_2) - u_0v_0 + u_1v_1. \end{aligned}$$

Multiplication over  $\mathbb{F}_{3^{3m}}$  involves 6 multiplications and 12 additions over  $\mathbb{F}_{3^m}$  (Algorithm 1).

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### Algorithm 1 Multiplication over $\mathbb{F}_{3^{3m}}$ .

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**Input:**  $U = u_0 + u_1\rho + u_2\rho^2$  and  $V = v_0 + v_1\rho + v_2\rho^2 \in \mathbb{F}_{3^{3m}}$ .

**Output:**  $W = U \cdot V \in \mathbb{F}_{3^{3m}}$ .

1.  $a_0 \leftarrow u_0 + u_1$ ;  $a_1 \leftarrow u_0 + u_2$ ;  $a_2 \leftarrow bu_1 + u_2$ ; (3A)
  2.  $a_3 \leftarrow v_0 + v_1$ ;  $a_4 \leftarrow v_0 + v_2$ ;  $a_5 \leftarrow v_1 + bv_2$ ; (3A)
  3.  $m_0 \leftarrow u_0 \cdot v_0$ ;  $m_1 \leftarrow u_1 \cdot v_1$ ;  $m_2 \leftarrow u_2 \cdot v_2$ ; (3M)
  4.  $m_3 \leftarrow a_0 \cdot a_3$ ;  $m_4 \leftarrow a_1 \cdot a_4$ ;  $m_5 \leftarrow a_2 \cdot a_5$ ; (3M)
  5.  $a_6 \leftarrow m_0 - m_1$ ; (1A)
  6.  $w_0 \leftarrow a_6 - m_2 + bm_5$  (2A)
  7. **if**  $b = 1$  **then**
  8.  $w_1 \leftarrow -a_6 + m_3 + m_5$ ; (2A)
  9. **else**
  10.  $w_1 \leftarrow -m_0 + m_3 + m_5$ ; (2A)
  11. **end if**
  12.  $w_2 \leftarrow -a_6 + m_4$ ; (1A)
  13. **return**  $w_0 + w_1\rho + w_2\rho^2$ ;
- 

## III. SQUARING OVER $\mathbb{F}_{3^{3m}}$

Let  $U = u_0 + u_1\rho + u_2\rho^2 \in \mathbb{F}_{3^{3m}}$ , with  $u_i \in \mathbb{F}_{3^m}$ ,  $0 \leq i \leq 2$ .  $V = U^2$  is given by

$$\begin{aligned} v_0 &= u_0^2 - bu_1u_2, \\ v_1 &= bu_2^2 - u_0u_1 - u_1u_2, \text{ and} \\ v_2 &= (u_0 + u_1) \cdot (u_0 + u_1 + u_2) - u_0^2 + u_0u_1 + u_1u_2. \end{aligned}$$

Thus, squaring over  $\mathbb{F}_{3^{3m}}$  requires 5 multiplications and 7 additions over  $\mathbb{F}_{3^m}$  (Algorithm 2).

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### Algorithm 2 Squaring over $\mathbb{F}_{3^{3m}}$ .

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**Input:**  $U = u_0 + u_1\rho + u_2\rho^2 \in \mathbb{F}_{3^{3m}}$ .

**Output:**  $V = U^2 \in \mathbb{F}_{3^{3m}}$ .

1.  $a_0 \leftarrow u_0 + u_1$ ;  $a_1 \leftarrow a_0 + u_2$ ; (2A)
  2.  $m_0 \leftarrow u_0^2$ ;  $m_1 \leftarrow u_0 \cdot u_1$ ;  $m_2 \leftarrow u_1 \cdot u_2$ ; (3M)
  3.  $m_3 \leftarrow u_2^2$ ;  $m_4 \leftarrow a_1^2$ ; (2M)
  4.  $a_2 \leftarrow m_1 + m_2$ ; (1A)
  5.  $v_0 \leftarrow m_0 - bm_2$ ; (1A)
  6.  $v_1 \leftarrow bm_3 - a_2$ ; (1A)
  7.  $v_2 \leftarrow m_4 + a_2 - m_0$ ; (2A)
  8. **return**  $v_0 + v_1\rho + v_2\rho^2$ ;
- 

## IV. INVERSION OVER $\mathbb{F}_{3^{3m}}$

Let  $V = v_0 + v_1\rho + v_2\rho^2 \in \mathbb{F}_{3^{3m}}$  be the multiplicative inverse of  $U = u_0 + u_1\rho + u_2\rho^2 \in \mathbb{F}_{3^{3m}}$ ,  $U \neq 0$ , where the  $u_i, v_i \in \mathbb{F}_{3^m}$ ,  $0 \leq i \leq 2$ . Since  $U \cdot V = 1$ , we obtain

$$\begin{cases} u_0v_0 + bu_2v_1 + bu_1v_2 = 1, \\ u_1v_0 + (u_0 + u_2)v_1 + (u_1 + bu_2)v_2 = 0, \\ u_2v_0 + u_1v_1 + (u_0 + u_2)v_2 = 0. \end{cases}$$

The solution of this system of equations is then given by

$$\begin{bmatrix} v_0 \\ v_1 \\ v_2 \end{bmatrix} = w^{-1} \begin{bmatrix} u_0^2 - (u_1^2 - u_2^2) - u_2(u_0 + bu_1) \\ bu_2^2 - u_0u_1 \\ u_1^2 - u_2^2 - u_0u_2 \end{bmatrix},$$

where  $w = u_0^2(u_0 - u_2) + u_1^2(-u_0 + bu_1) + u_2^2(-(-u_0 + bu_1) + u_2) \in \mathbb{F}_{3^m}$ . This operation involves 12 multiplications, 11 additions (or subtractions), and a single inversion over  $\mathbb{F}_{3^m}$  (Algorithm 3).

## V. CUBING OVER $\mathbb{F}_{3^{6m}}$

### A. General Algorithm

Let  $U = u_0 + u_1\sigma + u_2\rho + u_3\sigma\rho + u_4\rho^2 + u_5\sigma\rho^2 \in \mathbb{F}_{3^{6m}}$ .  $U^3 \in \mathbb{F}_{3^{6m}}$  is defined as follows:

$$U^3 = u_0^3 + u_1^3\sigma^3 + u_2^3\rho^3 + u_3^3(\sigma\rho)^3 + u_4^3(\rho^2)^3 + u_5^3(\sigma\rho^2)^3.$$

**Algorithm 3** Inversion over  $\mathbb{F}_{3^{3m}}$ .**Input:**  $U = u_0 + u_1\rho + u_2\rho^2 \in \mathbb{F}_{3^{3m}}$ ,  $U \neq 0$ .**Output:**  $V = U^{-1} \in \mathbb{F}_{3^{3m}}$ .

1.  $a_0 \leftarrow u_0 + bu_1$ ;  $a_1 \leftarrow u_0 - u_2$ ; (2A)
2.  $a_2 \leftarrow -u_0 + u_1$ ;  $a_3 \leftarrow -a_2 + u_2$ ; (2A)
3.  $m_0 \leftarrow u_0^2$ ;  $m_1 \leftarrow u_1^2$ ;  $m_2 \leftarrow u_2^2$ ; (3M)
4.  $m_3 \leftarrow u_0 \cdot u_1$ ;  $m_4 \leftarrow u_0 \cdot u_2$ ;  $m_5 \leftarrow u_2 \cdot a_0$ ; (3M)
5.  $m_6 \leftarrow m_0 \cdot a_1$ ;  $m_7 \leftarrow m_1 \cdot a_2$ ;  $m_8 \leftarrow m_2 \cdot a_3$ ; (3M)
6.  $w \leftarrow m_6 + m_7 + m_8$ ; (2A)
7.  $i \leftarrow w^{-1}$ ; (1I)
8.  $a_4 \leftarrow m_1 - m_2$ ;  $a_5 \leftarrow -a_4 + m_0 - m_5$ ; (3A)
9.  $a_6 \leftarrow bm_2 - m_3$ ;  $a_7 \leftarrow a_4 - m_4$ ; (2A)
10.  $v_0 \leftarrow i \cdot a_5$ ;  $v_1 \leftarrow i \cdot a_6$ ;  $v_2 \leftarrow i \cdot a_7$ ; (3M)
11. **return**  $v_0 + v_1\rho + v_2\rho^2$ ;

Since

$$\begin{cases} \rho^3 & = \rho + b, \\ (\rho^2)^3 & = \rho^2 - b\rho + 1, \end{cases} \quad \begin{cases} \sigma^3 & = -\sigma, \\ (\sigma\rho)^3 & = -\sigma\rho - b\sigma, \\ (\sigma\rho^2)^3 & = -\sigma\rho^2 + b\sigma\rho - \sigma, \end{cases}$$

we obtain the following coefficients for  $V = U^3$ :

$$\begin{aligned} v_0 &= u_0^3 + bu_1^3 + u_2^3, & v_1 &= -u_1^3 - bu_2^3 - u_3^3, \\ v_2 &= u_2^3 - bu_4^3, & v_3 &= -u_3^3 + bu_5^3, \\ v_4 &= u_4^3, & v_5 &= -u_5^3. \end{aligned}$$

As our unified operator computes  $-u_5^3$  in one clock cycle, cubing over  $\mathbb{F}_{3^{6m}}$  requires 6 cubings and 6 additions over  $\mathbb{F}_{3^m}$ .**B. Computation of  $(-t^2 + u\sigma - t\rho - \rho^2)^3$** Let  $U = -t^2 + u\sigma - t\rho - \rho^2$ . According to the previous formula for cubing over  $\mathbb{F}_{3^{6m}}$ , we have

$$V = U^3 = v_0 + v_1\sigma + v_2\rho - \rho^2,$$

where

$$\begin{aligned} v_0 &= -t^6 - bt^3 - 1, \\ v_1 &= -u^3, \text{ and} \\ v_2 &= -t^3 + b. \end{aligned}$$

Therefore,  $U^3$  is as sparse as  $U$  and this specific cubing involves a single multiplication, 2 cubings, and 3 additions over  $\mathbb{F}_{3^m}$  (Algorithm 4).This operation is usually followed by a multiplication which is optimized to take advantage of  $v_4 = -1$  (see for instance Appendix VI-C). Thus, our coprocessor does not explicitly compute  $v_4 \leftarrow -1$ .**Algorithm 4** Computation of  $(-t^2 + u\sigma - t\rho - \rho^2)^3$ .**Input:**  $t$  and  $u \in \mathbb{F}_{3^m}$ .**Output:**  $V = (-t^2 + u\sigma - t\rho - \rho^2)^3 \in \mathbb{F}_{3^{6m}}$ .

1.  $c_0 \leftarrow t^3$ ;  $c_1 \leftarrow -u^3$ ; (2C)
2.  $m_0 \leftarrow c_0^2$ ; (1M)
3.  $v_0 \leftarrow -m_0 - bc_0 - 1$ ; (2A)
4.  $v_1 \leftarrow c_1$ ;
5.  $v_2 \leftarrow b - c_0$ ; (1A)
6. **return**  $v_0 + v_1\sigma + v_2\rho - \rho^2$ ;

VI. MULTIPLICATION OVER  $\mathbb{F}_{3^{6m}}$ **A. General Algorithm**

Elements of  $\mathbb{F}_{3^{6m}}$  can be represented as degree-2 polynomials over  $\mathbb{F}_{3^{2m}}$ . Gorla *et al.* introduced an evaluation-interpolation scheme to perform multiplication over  $\mathbb{F}_{3^{6m}}$  by means of five multiplications over  $\mathbb{F}_{3^{2m}}$  [1]. Then, Karatsuba-Ofman's algorithm allows one to compute each multiplication over  $\mathbb{F}_{3^{2m}}$  by means of three multiplications over  $\mathbb{F}_{3^m}$  (see Appendix I). Thus, the scheme proposed by Gorla *et al.* to multiply two elements of  $\mathbb{F}_{3^{6m}}$  involves 15 multiplications over  $\mathbb{F}_{3^m}$  (Algorithm 5).

**B. Multiplication by a Sparse Operand**

The last multiplication over  $\mathbb{F}_{3^{6m}}$  of the  $\eta_T$  pairing algorithms is cheaper: it consists in computing the product  $(u_0 + u_1\sigma + u_2\rho) \cdot (v_0 + v_1\sigma + v_2\rho + v_3\sigma\rho + v_4\rho^2 + v_5\sigma\rho^2)$  and requires 12 multiplications and 51 additions over  $\mathbb{F}_{3^m}$  (Algorithm 6).

The first multiplication of the  $\eta_T$  pairing algorithms based on the reversed-loop approach also benefits from this optimization. Since

$$(u_0 + u_1\sigma + u_2\rho - \rho^2) \cdot V = (u_0 + u_1\sigma + u_2\rho) \cdot V - \rho^2 \cdot V, \quad (1)$$

it suffices to subtract  $\rho^2 V$  from the element of  $\mathbb{F}_{3^{6m}}$  returned by Algorithm 6. Recall that  $\rho^3 = \rho + b$  and note that Algorithm 6 requires two intermediate variables  $r_1 = v_0 + v_4$  and  $r_2 = v_1 + v_5$ . We then have

$$\begin{aligned} -\rho^2 \cdot V &= -v_2b - bv_3\sigma - (v_2 + bv_4)\rho \\ &\quad - (v_3 + bv_5)\sigma\rho - (v_0 + v_4)\rho^2 - (v_1 + v_5)\sigma\rho^2 \\ &= -v_2b - bv_3\sigma - (v_2 + bv_4)\rho \\ &\quad - (v_3 + bv_5)\sigma\rho - r_1\rho^2 - r_2\sigma\rho^2. \end{aligned}$$

Therefore, subtracting  $\rho^2 \cdot V$  involves 8 additions over  $\mathbb{F}_{3^m}$  and the total cost of Equation (1) is 12 multiplications and 59 additions over  $\mathbb{F}_{3^m}$ .

**C. Computation of  $(u_0 + u_1\sigma + u_2\rho - \rho^2) \cdot (v_0 + v_1\sigma + v_2\rho - \rho^2)$** 

The multiplication of  $U = u_0 + u_1\sigma + u_2\rho - \rho^2$  by  $V = v_0 + v_1\sigma + v_2\rho - \rho^2$ , where both  $U$  and  $V$  are in  $\mathbb{F}_{3^{6m}}$ , requires 6 multiplications and 21 additions over  $\mathbb{F}_{3^m}$  (Algorithm 7).

**D. Computation of  $(\lambda y_P t - \lambda y_Q \sigma - \lambda y_P \rho) \cdot (-t^2 + y_P y_Q \sigma - t\rho - \rho^2)$** 

We consider here the first multiplication over  $\mathbb{F}_{3^{6m}}$  of the  $\eta_T$  pairing calculation based on the reversed-loop approach. Noting

$$\begin{aligned} W &= w_0 + w_1\sigma + w_2\rho + w_3\sigma\rho + w_4\rho^2 + w_5\sigma\rho^2 \\ &= (\lambda y_P t - \lambda y_Q \sigma - \lambda y_P \rho) \cdot (-t^2 + y_P y_Q \sigma - t\rho - \rho^2), \end{aligned}$$

we easily check that

$$\begin{aligned} w_0 &= -\lambda y_P t^3 + \lambda y_P y_Q^2 + b\lambda y_P, \\ w_1 &= \lambda y_P^2 y_Q t + \lambda y_Q t^2, \\ w_2 &= \lambda y_P, \\ w_3 &= \lambda y_Q t - \lambda y_P^2 y_Q, \\ w_4 &= 0, \text{ and} \\ w_5 &= \lambda y_Q. \end{aligned}$$

These equations involve a single cubing, 6 additions, and 6 multiplications over  $\mathbb{F}_{3^m}$ .

## REFERENCES

- [1] E. Gorla, C. Puttmann, and J. Shokrollahi, "Explicit formulas for efficient multiplication in  $\mathbb{F}_{3^6m}$ ," in *Selected Areas in Cryptography – SAC 2007*, ser. Lecture Notes in Computer Science, C. Adams, A. Miri, and M. Wiener, Eds., no. 4876. Springer, 2007, pp. 173–183.

**Algorithm 5** Multiplication over  $\mathbb{F}_{3^6m}$  [1].

**Input:**  $U, V \in \mathbb{F}_{3^6m}$  with  $U = u_0 + u_1\sigma + u_2\rho + u_3\sigma\rho + u_4\rho^2 + u_5\sigma\rho^2$  and  $V = v_0 + v_1\sigma + v_2\rho + v_3\sigma\rho + v_4\rho^2 + v_5\sigma\rho^2$ .

**Output:**  $W = U \cdot V$ . The algorithm requires 15 multiplications and 67 additions over  $\mathbb{F}_{3^m}$ .

1.  $r_0 \leftarrow u_0 + u_4; a_0 \leftarrow r_0 + u_2; a_{12} \leftarrow r_0 - u_2;$  (3A)
2.  $r_0 \leftarrow v_0 + v_4; a_3 \leftarrow r_0 + v_2; a_{15} \leftarrow r_0 - v_2;$  (3A)
3.  $r_0 \leftarrow u_0 - u_4; a_6 \leftarrow r_0 - u_3; a_{18} \leftarrow r_0 + u_3;$  (3A)
4.  $r_0 \leftarrow v_0 - v_4; a_9 \leftarrow r_0 - v_3; a_{21} \leftarrow r_0 + v_3;$  (3A)
5.  $r_0 \leftarrow u_1 + u_5; a_1 \leftarrow r_0 + u_3; a_{13} \leftarrow r_0 - u_3;$  (3A)
6.  $r_0 \leftarrow v_1 + v_5; a_4 \leftarrow r_0 + v_3; a_{16} \leftarrow r_0 - v_3;$  (3A)
7.  $r_0 \leftarrow u_1 - u_5; a_7 \leftarrow r_0 + u_2; a_{19} \leftarrow r_0 - u_2;$  (3A)
8.  $r_0 \leftarrow v_1 - v_5; a_{10} \leftarrow r_0 + v_2; a_{22} \leftarrow r_0 - v_2;$  (3A)
9.  $a_2 \leftarrow a_0 + a_1; a_5 \leftarrow a_3 + a_4; a_8 \leftarrow a_6 + a_7;$  (3A)
10.  $a_{11} \leftarrow a_9 + a_{10}; a_{14} \leftarrow a_{12} + a_{13}; a_{17} \leftarrow a_{15} + a_{16};$  (3A)
11.  $a_{20} \leftarrow a_{18} + a_{19}; a_{23} \leftarrow a_{21} + a_{22};$  (2A)
12.  $a_{24} \leftarrow u_4 + u_5; a_{25} \leftarrow v_4 + v_5;$  (2A)
13.  $m_0 \leftarrow a_0 \cdot a_3; m_1 \leftarrow a_2 \cdot a_5; m_2 \leftarrow a_1 \cdot a_4;$  (3M)
14.  $m_3 \leftarrow a_6 \cdot a_9; m_4 \leftarrow a_8 \cdot a_{11}; m_5 \leftarrow a_7 \cdot a_{10};$  (3M)
15.  $m_6 \leftarrow a_{12} \cdot a_{15}; m_7 \leftarrow a_{14} \cdot a_{17}; m_8 \leftarrow a_{13} \cdot a_{16};$  (3M)
16.  $m_9 \leftarrow a_{18} \cdot a_{21}; m_{10} \leftarrow a_{20} \cdot a_{23}; m_{11} \leftarrow a_{19} \cdot a_{22};$  (3M)
17.  $m_{12} \leftarrow u_4 \cdot v_4; m_{13} \leftarrow a_{24} \cdot a_{25}; m_{14} \leftarrow u_5 \cdot v_5;$  (3M)
18. **if**  $b = 1$  **then**
19.  $t_0 \leftarrow m_0 + m_4 + m_{12}; t_1 \leftarrow m_2 + m_{10} + m_{14};$  (4A)
20.  $t_2 \leftarrow m_6 + m_{12}; t_3 \leftarrow -m_8 - m_{14};$  (2A)
21.  $t_4 \leftarrow m_7 + m_{13}; t_5 \leftarrow t_3 + m_2;$  (2A)
22.  $t_6 \leftarrow t_2 - m_0; t_7 \leftarrow t_3 - m_2 + m_5 + m_{11};$  (4A)
23.  $t_8 \leftarrow t_2 + m_0 - m_3 - m_9;$  (3A)
24.  $w_0 \leftarrow -t_0 + t_1 - m_3 + m_{11};$  (3A)
25.  $w_1 \leftarrow t_0 + t_1 - m_1 + m_5 + m_9 - m_{13};$  (5A)
26.  $w_2 \leftarrow t_5 + t_6;$  (1A)
27.  $w_3 \leftarrow t_5 - t_6 + t_4 - m_1;$  (3A)
28.  $w_4 \leftarrow t_7 + t_8;$  (1A)
29.  $w_5 \leftarrow t_7 - t_8 + t_4 + m_1 - m_4 - m_{10};$  (5A)
30. **else**
31.  $t_0 \leftarrow m_4 + m_8 + m_{14}; t_1 \leftarrow m_6 + m_{12};$  (3A)
32.  $t_2 \leftarrow t_1 + m_{10}; t_3 \leftarrow m_2 + m_{14};$  (2A)
33.  $t_4 \leftarrow t_3 - m_8; t_5 \leftarrow -m_0 + m_6 - m_{12};$  (3A)
34.  $t_6 \leftarrow -t_3 + m_5 - m_8 + m_{11};$  (3A)
35.  $t_7 \leftarrow t_1 + m_0 - m_3 - m_9; t_8 \leftarrow m_1 + m_{13};$  (4A)
36.  $w_0 \leftarrow t_0 - t_2 + m_5 - m_9;$  (3A)
37.  $w_1 \leftarrow t_0 + t_2 + m_3 - m_7 + m_{11} - m_{13};$  (5A)
38.  $w_2 \leftarrow t_4 + t_5;$  (1A)
39.  $w_3 \leftarrow t_4 - t_5 - t_8 + m_7;$  (3A)
40.  $w_4 \leftarrow t_6 + t_7;$  (1A)
41.  $w_5 \leftarrow t_6 - t_7 + t_8 - m_4 + m_7 - m_{10};$  (5A)
42. **end if**
43. **return**  $w_0 + w_1\sigma + w_2\rho + w_3\sigma\rho + w_4\rho^2 + w_5\sigma\rho^2;$

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**Algorithm 6** Computation of  $(u_0 + u_1\sigma + u_2\rho) \cdot (v_0 + v_1\sigma + v_2\rho + v_3\sigma\rho + v_4\rho^2 + v_5\sigma\rho^2)$ .

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**Input:**  $U, V \in \mathbb{F}_{3^{6m}}$  with  $U = u_0 + u_1\sigma + u_2\rho$  and  $V = v_0 + v_1\sigma + v_2\rho + v_3\sigma\rho + v_4\rho^2 + v_5\sigma\rho^2$ .

**Output:**  $W = U \cdot V$ . The algorithm requires 12 multiplications and 51 additions over  $\mathbb{F}_{3^m}$ .

1.  $a_0 \leftarrow u_0 + u_2; a_{10} \leftarrow u_0 - u_2;$  (2A)
  2.  $r_1 \leftarrow v_0 + v_4; a_2 \leftarrow r_1 + v_2; a_{12} \leftarrow r_1 - v_2;$  (3A)
  3.  $r_0 \leftarrow v_0 - v_4; a_7 \leftarrow r_0 - v_3; a_{17} \leftarrow r_0 + v_3;$  (3A)
  4.  $r_2 \leftarrow v_1 + v_5; a_3 \leftarrow r_2 + v_3; a_{13} \leftarrow r_2 - v_3;$  (3A)
  5.  $a_5 \leftarrow u_1 + u_2; a_{15} \leftarrow u_1 - u_2;$  (2A)
  6.  $r_0 \leftarrow v_1 - v_5; a_8 \leftarrow r_0 + v_2; a_{18} \leftarrow r_0 - v_2;$  (3A)
  7.  $a_1 \leftarrow a_0 + u_1; a_4 \leftarrow a_2 + a_3; a_6 \leftarrow u_0 + a_5;$  (3A)
  8.  $a_9 \leftarrow a_7 + a_8; a_{11} \leftarrow a_{10} + u_1; a_{14} \leftarrow a_{12} + a_{13};$  (3A)
  9.  $a_{16} \leftarrow u_0 + a_{15}; a_{19} \leftarrow a_{17} + a_{18};$  (2A)
  10.  $m_0 \leftarrow a_0 \cdot a_2; m_1 \leftarrow a_1 \cdot a_4; m_2 \leftarrow u_1 \cdot a_3;$  (3M)
  11.  $m_3 \leftarrow u_0 \cdot a_7; m_4 \leftarrow a_6 \cdot a_9; m_5 \leftarrow a_5 \cdot a_8;$  (3M)
  12.  $m_6 \leftarrow a_{10} \cdot a_{12}; m_7 \leftarrow a_{11} \cdot a_{14}; m_8 \leftarrow u_1 \cdot a_{13};$  (3M)
  13.  $m_9 \leftarrow u_0 \cdot a_{17}; m_{10} \leftarrow a_{16} \cdot a_{19}; m_{11} \leftarrow a_{15} \cdot a_{18};$  (3M)
  14. **if**  $b = 1$  **then**
  15.  $t_0 \leftarrow m_0 + m_4; t_1 \leftarrow m_2 + m_{10};$  (2A)
  16.  $t_2 \leftarrow -m_8 + m_2; t_3 \leftarrow m_6 - m_0;$  (2A)
  17.  $t_4 \leftarrow -m_8 - m_2 + m_5 + m_{11};$  (3A)
  18.  $t_5 \leftarrow m_6 + m_0 - m_3 - m_9;$  (3A)
  19.  $w_0 \leftarrow -t_0 + t_1 - m_3 + m_{11};$  (3A)
  20.  $w_1 \leftarrow t_0 + t_1 - m_1 + m_5 + m_9;$  (4A)
  21.  $w_2 \leftarrow t_2 + t_3;$  (1A)
  22.  $w_3 \leftarrow t_2 - t_3 + m_7 - m_1;$  (3A)
  23.  $w_4 \leftarrow t_4 + t_5;$  (1A)
  24.  $w_5 \leftarrow t_4 - t_5 + m_7 + m_1 - m_4 - m_{10};$  (5A)
  25. **else**
  26.  $t_0 \leftarrow m_4 + m_8; t_1 \leftarrow m_6 + m_{10};$  (2A)
  27.  $t_2 \leftarrow m_2 - m_8; t_3 \leftarrow -m_0 + m_6;$  (2A)
  28.  $t_4 \leftarrow -m_2 + m_5 - m_8 + m_{11};$  (3A)
  29.  $t_5 \leftarrow m_6 + m_0 - m_3 - m_9;$  (3A)
  30.  $w_0 \leftarrow t_0 - t_1 + m_5 - m_9;$  (3A)
  31.  $w_1 \leftarrow t_0 + t_1 + m_3 - m_7 + m_{11};$  (4A)
  32.  $w_2 \leftarrow t_2 + t_3;$  (1A)
  33.  $w_3 \leftarrow t_2 - t_3 - m_1 + m_7;$  (3A)
  34.  $w_4 \leftarrow t_4 + t_5;$  (1A)
  35.  $w_5 \leftarrow t_4 - t_5 + m_1 - m_4 + m_7 - m_{10};$  (5A)
  36. **end if**
  37. **return**  $w_0 + w_1\sigma + w_2\rho + w_3\sigma\rho + w_4\rho^2 + w_5\sigma\rho^2;$
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**Algorithm 7** Computation of  $(u_0 + u_1\sigma + u_2\rho - \rho^2) \cdot (v_0 + v_1\sigma + v_2\rho - \rho^2)$ .

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**Input:**  $U = (u_0 + u_1\sigma + u_2\rho - \rho^2)$  and  $V = v_0 + v_1\sigma + v_2\rho - \rho^2 \in \mathbb{F}_{3^{6m}}$ .

**Output:**  $W = U \cdot V \in \mathbb{F}_{3^{6m}}$ .

1.  $a_0 \leftarrow u_0 + u_1; a_1 \leftarrow u_0 + u_2; a_2 \leftarrow u_1 + u_2;$  (3A)
  2.  $a_3 \leftarrow v_0 + v_1; a_4 \leftarrow v_0 + v_2; a_5 \leftarrow v_1 + v_2;$  (3A)
  3.  $a_6 \leftarrow u_2 + v_2;$  (1A)
  4.  $m_1 \leftarrow u_0 \cdot v_0; m_2 \leftarrow u_1 \cdot v_1; m_3 \leftarrow u_2 \cdot v_2;$  (3M)
  5.  $m_4 \leftarrow a_0 \cdot a_3; m_5 \leftarrow a_1 \cdot a_4; m_6 \leftarrow a_2 \cdot a_5;$  (3M)
  6.  $w_0 \leftarrow m_1 - m_2 - ba_6;$  (2A)
  7.  $w_1 \leftarrow m_4 - m_1 - m_2;$  (2A)
  8.  $w_2 \leftarrow m_5 - m_1 - m_3 - a_6 + b;$  (4A)
  9.  $w_3 \leftarrow m_6 - m_2 - m_3;$  (2A)
  10.  $w_4 \leftarrow 1 + m_3 - u_0 - v_0;$  (3A)
  11.  $w_5 \leftarrow -u_1 - v_1;$  (1A)
  12. **return**  $w_0 + w_1\sigma + w_2\rho + w_3\sigma\rho + w_4\rho^2 + w_5\sigma\rho^2;$
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**Algorithm 8** First multiplication of the reversed-loop  $\eta_T$  pairing calculation.

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**Input:**  $U = \lambda y_P t - \lambda y_Q \sigma - \lambda y_P \rho$  and  $V = -t^2 + y_P y_Q \sigma - t\rho - \rho^2$ .

**Output:**  $W = U \cdot V \in \mathbb{F}_{3^{6m}}$ .

1.  $m_0 \leftarrow y_Q \cdot t; m_1 \leftarrow y_P \cdot y_Q; m_2 \leftarrow y_P \cdot m_1;$  (3M)
  2.  $a_0 \leftarrow \lambda m_0 + \lambda m_2;$  (1A)
  3.  $c_0 \leftarrow t^3;$  (1C)
  4.  $m_3 \leftarrow a_0 \cdot t; m_4 \leftarrow y_P \cdot c_0; m_5 \leftarrow y_Q \cdot m_1;$  (3M)
  5.  $w_0 \leftarrow -\lambda m_4 + \lambda m_5 + b\lambda y_P;$  (2A)
  6.  $w_1 \leftarrow m_3;$  (1A)
  7.  $w_2 \leftarrow \lambda y_P;$  (1A)
  8.  $w_3 \leftarrow \lambda m_0 - \lambda m_2;$  (1A)
  9.  $w_5 \leftarrow \lambda y_Q;$  (1A)
  10. **return**  $w_0 + w_1\sigma + w_2\rho + w_3\sigma\rho + w_5\sigma\rho^2;$
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