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# DIFFERENTIAL EQUATIONS FOR ALGEBRAIC FUNCTIONS 

ALIN BOSTAN, FRÉDÉRIC CHYZAK, GRÉGOIRE LECERF, BRUNO SALVY, AND ÉRIC<br>SCHOST


#### Abstract

It is classical that univariate algebraic functions satisfy linear differential equations with polynomial coefficients. Linear recurrences follow for the coefficients of their power series expansions. We show that the linear differential equation of minimal order has coefficients whose degree is cubic in the degree of the function. We also show that there exists a linear differential equation of order linear in the degree whose coefficients are only of quadratic degree. Furthermore, we prove the existence of recurrences of order and degree close to optimal. We study the complexity of computing these differential equations and recurrences. We deduce a fast algorithm for the expansion of algebraic series.


## 1. Introduction

A power series $\alpha(X) \in \mathbb{K}[[X]]$ is called algebraic over the field $\mathbb{K}$ when there is a nonzero polynomial $P \in \mathbb{K}[X, Y]$ such that $P(X, \alpha(X))=0$.

The existence of a linear differential equation with polynomial coefficients satisfied by algebraic series seems to have been observed first in 1827 by Abel [1, p. 287] who did not publish it. Cockle gave an algorithm for the computation of such a differential equation, that he called a differential resolvent [11]. The method was then presented by Harley [15] and rediscovered by Tannery [21] (see also [14]).

One of the applications of these differential equations is the computation of power series solutions: a linear differential equation translates into a linear recurrence, with the consequence that the number of operations required to compute the first $N$ coefficients grows only linearly with $N$. This is useful in combinatorial applications, since the enumeration sequences of context-free languages have algebraic generating series. This method has been described by Comtet [12] and studied from the complexity point of view by Chudnovsky and Chudnovsky [8]. In practice, this procedure will be all the more efficient as its dependency on the size of $P$ is small. This in turn is related to the order and degree of the coefficients of the differential equation.

We say that a linear differential equation or operator vanishing at all the roots $\alpha(X)$ of $P$ is associated to $P$. In this work, we derive precise bounds for the order and degree of coefficients of several such equations. For brevity and clarity, the following hypothesis will be assumed throughout.
Hypothesis (H). The characteristic of the field $\mathbb{K}$ is 0 . The polynomial $P \in$ $\mathbb{K}[X, Y]$ is separable (wrt $Y$ ), i.e., coprime with its derivative $P_{Y}$ wrt $Y$. The total degree of $P$ is denoted $D$ and its degrees in $X$ and $Y$ are written $D_{X}$ and $D_{Y}$. We assume that $D_{Y} \geq 1$.

The first work on bounds that we know of in this area is due to Cormier et alii [13]. They consider the monic linear differential operator of minimal order
associated to $P$, which we call the differential resolvent of $P$. A bound in $\mathcal{O}\left(D_{X} D_{Y}^{4}\right)$ for the degrees of its coefficients can be deduced from their work. Nahay [17, 18] obtained degree bounds in $\mathcal{O}\left(D^{3}\right)$ for the coefficients of an analogous equation satisfied by $\alpha^{\lambda}$ with $\lambda$ transcendent over $\mathbb{K}$. In $\S 2$, we prove the following cubic bound ( $\partial_{X}$ denotes differentiation wrt $X$ ).

Theorem 1. Under (H), the differential resolvent of $P$ can be written

$$
\partial_{X}^{r}+\frac{M_{r-1}}{M_{r}} \partial_{X}^{r-1}+\cdots+\frac{M_{1}}{M_{r}} \partial_{X}+\frac{M_{0}}{M_{r}}
$$

where $r \leq D_{Y}$ and $M_{0}, \ldots, M_{r}$ in $\mathbb{K}[X]$ have degrees at most

$$
\eta=\left((2 r-1) D_{Y}+2 r^{2}-4 r+3\right) D_{X}-r(r-1) / 2 \in \mathcal{O}\left(r D_{X} D_{Y}\right)
$$

Experiments show that a cubic degree is to be expected (see $\S 4.2$ ). This theorem improves and extends the results of $[13,17,18]$. It leads to a linear recurrence of order $\mathcal{O}\left(D^{3}\right)$ with coefficients of degree $\mathcal{O}(D)$. Using fast multipoint evaluation on an arithmetic progression then leads to a computation of the first $N$ coefficients in $\mathcal{O}\left(D^{2} \mathrm{M}(D) N\right)$ arithmetic operations (see $\left.\S 3.5\right)$. As usual, $\mathrm{M}(D)$ denotes a bound on the number of arithmetic operations in the ring of coefficients needed to multiply two polynomials of degree at most $D$. We make the standard assumption that $\mathrm{M}\left(d_{1}\right)+\mathrm{M}\left(d_{2}\right) \leq \mathrm{M}\left(d_{1}+d_{2}\right)$ for all $d_{1}, d_{2}$ and that $D \log D \in \mathcal{O}(\mathrm{M}(D))$.

One reason for this cubic degree growth is the presence of numerous apparent singularities. These are points where the differential resolvent is singular (its leading coefficient vanishes) while none of its solutions is. However, the number of apparent singularities decreases dramatically if one does not insist on the order of the equation being minimal. It is actually possible to remove all apparent singularities by Weyl closure [22], but we do not know how to get bounds from this algorithm.

In $\S 3.2$, we show that a moderate increase of the order leads to the following quadratic bound.

Theorem 2. Under (H), there exists a nonzero linear differential operator in $\mathbb{K}[X]\left\langle\partial_{X}\right\rangle$ associated to $P$ whose degree in $\partial_{X}$ is at most $6 D_{Y}$ and with coefficients of degree at most $3 D_{X} D_{Y}$.

This operator then translates into a linear recurrence of order at most $3 D_{Y}\left(D_{X}+\right.$ 2) with coefficients of degree at most $6 D_{Y}$. A natural question is then the order of a minimal linear recurrence, or equivalently the minimal degree in $X$ of linear differential operators in $X$ and $\theta_{X}=X \partial_{X}$ associated to $P$. In $\S 4.2$, we explain why quadratic order is optimal. We prove in $\S 3$ the following precise bound.

Theorem 3. Under $(H)$, there exists a non-zero differential operator in $\mathbb{K}[X]\left\langle\theta_{X}\right\rangle$ associated to $P$ whose degrees in $X$ and in $\theta_{X}$ are bounded above by

$$
2 D_{X} D_{Y}+D_{Y}-\Delta^{2}-\Delta+1 \quad \text { with } \quad \Delta=D_{X}+D_{Y}-D
$$

Experiments showing that this bound is not far from optimal are reported in §4.2. This result implies the existence of a linear recurrence of order $\mathcal{O}\left(D^{2}\right)$ with coefficients of the same degree, whence a computation of the first $N$ coefficients of series solutions in $\mathcal{O}\left(\mathrm{M}\left(D^{2}\right) N\right)$, with a small constant in the $\mathcal{O}(\cdot)$ term. In $\S 3.4$, we go a bit further: with a degree in $X$ slightly larger, but still quadratic, we obtain a degree in $\theta_{X}$ that is only linear; this results in an algorithm for algebraic series in $\mathcal{O}(D \mathrm{M}(D) N)$ operations only.

## 2. Differential Resolvent

We write $\alpha_{1}, \ldots, \alpha_{D_{Y}}$ for the zeroes of $P(X, \cdot)$ in the algebraic closure $\overline{\mathbb{K}(X)}$ of $\mathbb{K}(X)$. Since $P$ is separable wrt $Y$, the derivation $\partial_{X}$ uniquely extends to the splitting field $\mathbb{K}(X)\left(\alpha_{1}, \ldots, \alpha_{D_{Y}}\right)$ of $P$ by $\partial_{X} \alpha_{i}=-P_{X}\left(X, \alpha_{i}\right) / P_{Y}\left(X, \alpha_{i}\right)$.
2.1. Cockle's Algorithm. Cockle's algorithm [11] computes the differential resolvent. It reduces the computation to linear algebra in the quotient ring $\mathbb{K}(X)[Y] /(P)$, that is a vector space of dimension $D_{Y}$ over $\mathbb{K}(X)$. For the sake of its presentation, we introduce two sequences of polynomials.

First, polynomials $\left(W_{k}\right)_{k \geq 1}$ in $\mathbb{K}[X, Y]$ satisfying

$$
\alpha_{i}^{(k)}=\frac{W_{k}\left(X, \alpha_{i}\right)}{P_{Y}\left(X, \alpha_{i}\right)^{2 k-1}}, \quad i \in\left\{1, \ldots, D_{Y}\right\}
$$

are defined by induction: $W_{1}=-P_{X}$, and for all $k \geq 1, W_{k+1}=\left(P_{Y} \partial_{X} W_{k}-\right.$ $\left.P_{X} \partial_{Y} W_{k}\right) P_{Y}-(2 k-1)\left(P_{Y} \partial_{X} P_{Y}-P_{X} \partial_{Y} P_{Y}\right) W_{k}$. They satisfy

$$
\begin{aligned}
\operatorname{deg}_{X}\left(W_{k}\right) & \leq\left(2 D_{X}-1\right) k-D_{X} \\
\operatorname{deg}_{Y}\left(W_{k}\right) & \leq 2\left(D_{Y}-1\right) k-D_{Y}+2
\end{aligned}
$$

Cockle's algorithm computes a second sequence of polynomials, $\left(V_{k}\right)_{k \geq 0}$, such that $V_{k}(\alpha)$ is the image of $W_{k} / P_{Y}^{2 k-1}$ in $\mathbb{K}(X)[Y] /(P)$ for $k \geq 1$, while $V_{0}(\alpha)=\alpha$ is that of $Y$. The computation is as follows:
(1) $V_{1}$ is computed by first inverting $P_{Y}$ modulo $P$ and then reducing the product with $-P_{X}$ modulo $P$;
(2) for $k=0,1, \ldots$, the polynomial $V_{k+1}$ is computed by taking the remainder of $\partial_{X} V_{k}+V_{1} \partial_{Y} V_{k}$ modulo $P$;
(3) the computation stops with the first integer $r \leq D_{Y}$ such that $V_{r}$ belongs to the $\mathbb{K}(X)$-vector space spanned by $V_{0}, \ldots, V_{r-1}$. Then, together with $r$, Cockle's algorithm computes the relation

$$
\begin{equation*}
V_{r}=A_{r-1} V_{r-1}+\cdots+A_{0} V_{0} \tag{1}
\end{equation*}
$$

and returns the unique monic differential operator of minimal order $M=$ $\partial_{X}^{r}-A_{r-1} \partial_{X}^{r-1}+\cdots+A_{1} \partial_{X}+A_{0}$ associated to $P$.
2.2. Degree Bound. We are now going to prove Thm. 1 by establishing a formula for the differential resolvent in terms of the elementary symmetric polynomials of the $\alpha_{i}$. This idea originates in Nahay's PhD thesis [16, Thm. 46, p. 96] (published in [18, Thm. 9.1]).

Let $Y_{1}, \ldots, Y_{D_{Y}}$ be new indeterminates and $\mathcal{N}$ the matrix:

$$
\left(\begin{array}{cccc}
Y_{1} P_{Y}\left(X, Y_{1}\right)^{2 r-1} & \cdots & Y_{r} P_{Y}\left(X, Y_{r}\right)^{2 r-1} & 1 \\
W_{1}\left(X, Y_{1}\right) P_{Y}\left(X, Y_{1}\right)^{2 r-2} & \cdots & W_{1}\left(X, Y_{r}\right) P_{Y}\left(X, Y_{r}\right)^{2 r-2} & \partial_{X} \\
W_{2}\left(X, Y_{1}\right) P_{Y}\left(X, Y_{1}\right)^{2 r-4} & \cdots & W_{2}\left(X, Y_{r}\right) P_{Y}\left(X, Y_{r}\right)^{2 r-4} & \partial_{X}^{2} \\
\vdots & & \vdots & \vdots \\
W_{r}\left(X, Y_{1}\right) & \cdots & W_{r}\left(X, Y_{r}\right) & \partial_{X}^{r}
\end{array}\right)
$$

By subtracting the $j$ th column of $\mathcal{N}$ to the $i$ th one $(i, j \leq r)$ we observe that $Y_{i}-Y_{j}$ divides the determinant $\operatorname{det}(\mathcal{N})$ of $\mathcal{N}$, hence $\delta:=\prod_{1 \leq k<l \leq r}\left(Y_{k}-Y_{l}\right) \operatorname{divides} \operatorname{det}(\mathcal{N})$. The operator $L=\operatorname{det}(\mathcal{N}) / \delta$ is thus in $\mathbb{K}[X]\left[Y_{1}, \ldots, Y_{r}\right]\left\langle\partial_{X}\right\rangle$ and is symmetric in $Y_{1}, \ldots, Y_{r}$.

By definition of the $W_{k}$ 's, for all $\sigma$ in the permutation group $S_{D_{Y}}$ of $\left\{1, \ldots, D_{Y}\right\}$ and all $i \in\{1, \ldots, r\}$, we have

$$
\begin{aligned}
L\left(\alpha_{\sigma(1)}, \ldots, \alpha_{\sigma(r)}, \partial_{X}\right) \alpha_{i} & = \\
& \left(\prod_{j=1}^{r} P_{Y}\left(X, \alpha_{\sigma(j)}\right)\right)^{2 r-1} \quad \frac{\mathrm{Wr}\left(\alpha_{\sigma(1)}, \ldots, \alpha_{\sigma(r)}, \alpha_{i}\right)}{\delta\left(\alpha_{\sigma(1)}, \ldots, \alpha_{\sigma(r)}\right)}=0,
\end{aligned}
$$

where $\mathrm{Wr}(\cdot)$ is the Wronskian determinant. In other words $L\left(\alpha_{\sigma(1)}, \ldots, \alpha_{\sigma(r)}, \partial_{X}\right)$ is associated to $P$ for all $\sigma \in S_{D_{Y}}$.

For all $k \in\{0, \ldots, r\}$, the entries of the $(k+1)$ th row of $\mathcal{N}$, have degree in $X$ at most $m_{k}=(2 r-1) D_{X}-k$, except in the last column. Therefore, for all $i \in\{0, \ldots, r\}$, the degree of the coefficient $L_{i}$ of $\partial_{X}^{i}$ in $L$ is bounded by:

$$
\operatorname{deg}_{X}\left(L_{i}\right) \leq \sum_{k=0, k \neq i}^{r} m_{k} \leq r(2 r-1) D_{X}-r(r-1) / 2=: m_{X}
$$

On the other hand, for all $l \in\{1, \ldots, r\}$, the entries of the $l$ th column of $\mathcal{N}$ have degree in $Y_{l}$ at most $(2 r-1) D_{Y}-2 r+2$. Therefore, for all $i \in\{0, \ldots, r\}$, all $j \geq 0$, and all $l \in\{1, \ldots, r\}$, the degree of the coefficient $L_{i, j}$ of $X^{j}$ in $L_{i}$ is bounded by:

$$
\operatorname{deg}_{Y_{l}}\left(L_{i, j}\right) \leq(2 r-1) D_{Y}-2 r+2-(r-1)=: m_{Y}
$$

By definition of $r$, there exists a permutation $\sigma \in S_{D_{Y}}$ such that the operator $L\left(\alpha_{\sigma(1)}, \ldots, \alpha_{\sigma(r)}, \partial_{X}\right)$ has order $r$ exactly. Therefore, Lemma 1 below provides us with a nonzero operator $\bar{L} \in \mathbb{K}[X]\left[Y_{1}, \ldots, Y_{D_{Y}}\right]\left\langle\partial_{X}\right\rangle$ symmetric in $Y_{1}, \ldots, Y_{D_{Y}}$, such that $\bar{L}\left(\alpha_{1}, \ldots, \alpha_{D_{Y}}, \partial_{X}\right)$ has order $r$, is associated to $P$, and $\operatorname{deg}_{Y_{i}}(\bar{L}) \leq$ $\max \left(m_{Y}, D_{Y}-1\right)=m_{Y}$, for all $i \in\left\{1, \ldots, D_{Y}\right\}$.

Let $p_{i}$ denote the coefficient of $Y^{i}$ in $P$ seen in $\mathbb{K}[X][Y]$. By Lemma 2 below, there exists $\tilde{L} \in \mathbb{K}[X]\left[Z_{1}, \ldots, Z_{D_{Y}}\right]\left\langle\partial_{X}\right\rangle$ of total degree at most $m_{Y}$ in $Z_{1}, \ldots, Z_{D_{Y}}$, of partial degree in $X$ at most $m_{X}$, and such that

$$
\bar{L}\left(\alpha_{1}, \ldots, \alpha_{D_{Y}}, \partial_{X}\right)=\tilde{L}\left(p_{0} / p_{D_{Y}}, \ldots, p_{D_{Y}-1} / p_{D_{Y}}, \partial_{X}\right)
$$

Finally, $\tilde{M}=p_{D_{Y}}^{m_{Y}} \tilde{L}\left(p_{0} / p_{D_{Y}}, \ldots, p_{D_{Y}-1} / p_{D_{Y}}, \partial_{X}\right)$ is an operator in $\mathbb{K}[X]\left\langle\partial_{X}\right\rangle$ of order $r$ and degree at most $m_{X}+D_{X} m_{Y}$ in $X$, which is associated to $P$. The conclusion of Thm. 1 thus follows from the uniqueness of the minimal operator.

The rest of the proof consists of elementary properties of symmetric polynomials. If $r \leq n$ then the group $S_{r}$ of the permutations of $\{1, \ldots, r\}$ is naturally embedded into $S_{n}$. We write $\left(S_{n} / S_{r}\right)_{l}$ for the left quotient of $S_{n}$ by $S_{r}$.

Lemma 1. Let $\mathbb{E}$ be a field, let $r \leq n$ be nonnegative integers, and let $\alpha_{1}, \ldots, \alpha_{n}$ be pairwise distinct elements in $\mathbb{E}$. If $E \in \mathbb{E}\left[Y_{1}, \ldots, Y_{r}\right]$ is symmetric and does not vanish at $\left(\alpha_{1}, \ldots, \alpha_{r}\right)$, then there exist nonnegative integers $e_{r+1} \leq r, \ldots, e_{n} \leq n-1$ such that

$$
\bar{E}=\sum_{\sigma \in\left(S_{n} / S_{r}\right)_{l}} E\left(Y_{\sigma(1)}, \ldots, Y_{\sigma(r)}\right) Y_{\sigma(r+1)}^{e_{r+1}} \cdots Y_{\sigma(n)}^{e_{n}}
$$

is a symmetric polynomial in $\mathbb{E}\left[Y_{1}, \ldots, Y_{n}\right]$ that does not vanish at $\left(\alpha_{1}, \ldots, \alpha_{n}\right)$.
Proof. For fixed $r$, the proof proceeds by induction on $n$. If $n=r$, then the lemma trivially holds. We now assume that the lemma holds for some $n \geq r$, and prove that it holds for $n+1$. The induction hypothesis provides us with nonnegative integers $e_{r+1} \leq r, \ldots, e_{n} \leq n-1$ such that $\bar{E}$ (as defined in the lemma) is a
symmetric polynomial in $Y_{1}, \ldots, Y_{n}$ that does not vanish at $\left(\alpha_{1}, \ldots, \alpha_{n}\right)$. Since the Vandermonde matrix of $\alpha_{1}, \ldots, \alpha_{n+1}$ is invertible, there exists $e_{n+1} \leq n$ such that

$$
\sum_{\rho \in\left(S_{n+1} / S_{n}\right)_{l}} \bar{E}\left(Y_{\rho(1)}, \ldots, Y_{\rho(n)}\right) Y_{\rho(n+1)}^{e_{n+1}}
$$

is symmetric in $Y_{1}, \ldots, Y_{n+1}$ and moreover does not vanish at $\left(\alpha_{1}, \ldots, \alpha_{n+1}\right)$. Plugging the expression of $\bar{E}$ into the latter sum concludes the proof for $n+1$.

Lemma 2. [24, Thm 6.21] Let $\mathbb{A}$ be a commutative ring, and let $E \in \mathbb{A}\left[Y_{1}, \ldots, Y_{n}\right]$ be symmetric of partial degree $D$ in each of the $Y_{i}$ 's. Then there exists a unique polynomial $F \in \mathbb{A}\left[Z_{1}, \ldots, Z_{n}\right]$ of total degree $D$ such that $E\left(Y_{1}, \ldots, Y_{n}\right)=F\left(\Sigma_{1}, \ldots, \Sigma_{n}\right)$, where $\Sigma_{i}$ represents the $i$-th elementary symmetric polynomial in $Y_{1}, \ldots, Y_{n}$.
2.3. Algorithm and Complexity. With the bounds of the previous section, we now turn to the computation of the differential resolvent. As in [8, 13], the idea is to compute power series expansions of the coefficients and recover them by Padé approximants. For this, we have to choose a point where expansion is easy and determine a bound on the order of the series expansions.

Definition 1. A point $a \in \mathbb{K}$ that annihilates neither the leading coefficient $p_{D_{Y}}$ nor the discriminant $\Delta_{Y}$ of $P$ seen in $\mathbb{K}[X][Y]$, and such that $V_{0}(a), \ldots, V_{r-1}(a)$ are linearly independent in $\mathbb{K}[Y] /(P(a, Y))$ is called a lucky point.

If $a$ is lucky, then the $\alpha_{i}$ 's can be seen as power series in $\overline{\mathbb{K}}[[X-a]]$, and the $V_{i}$ 's as elements of $\mathbb{K}[[X-a]][Y] /(P)$. Here $\overline{\mathbb{K}}$ represents the algebraic closure of $\mathbb{K}$.

Lemma 3. There exists $\Phi \in \mathbb{K}[X] \backslash\{0\}$ of total degree at most $\eta$ (as defined in Thm. 1) such that $a$ is lucky whenever $p_{D_{Y}}(a) \Delta_{Y}(a) \Phi(a) \neq 0$.

Proof. For $\Phi$ we take the leading coefficient of the operator $\tilde{M}$ introduced at the end of the proof of Thm. 1. If $p_{D_{Y}}(a) \Delta_{Y}(a) \neq 0$ then we can define $A(a)$ as the $r \times D_{Y}$ matrix $\left(\alpha_{j}^{(i-1)}(a)\right)_{i, j}, B(a)$ as the $D_{Y} \times r$ matrix in $\mathbb{K}[[X-a]]$ whose columns are $V_{0}, \ldots, V_{r-1}$, and $C(a)$ as the invertible Vandermonde matrix of the $\alpha_{i}$. If $\Phi(a) \neq 0$ then $A(a)$ has rank $r$ by construction, and so has $B(a)$ since $A(a)=C(a) B(a)$.

Lemma 3 implies that all but a finite set of elements in $\mathbb{K}$ are lucky. A lucky point $a$ can thus be obtained either deterministically by computing the rank of $V_{0}(a), \ldots, V_{r-1}(a)$ for more than $\operatorname{deg}_{X}\left(p_{D_{Y}} \Delta_{Y} \Phi\right) \in \mathcal{O}\left(D_{X} D_{Y}^{2}\right)$ points, or probabilistically by means of the Schwartz-Zippel zero test [23, Lemma 6.44]. Once a lucky point is determined (which is very fast by the probabilistic approach) then the computation of the differential resolvent of $P$ can be done efficiently.

Proposition 1. Given a lucky point $a \in \mathbb{K}$ and the order $r$ of $M$, the computation of $M$ can be done with $\mathcal{O}\left(\left(r^{\omega}+r \mathrm{M}\left(D_{Y}\right)+r \log \left(D_{X}\right)\right) \mathrm{M}(\eta)\right)$ arithmetic operations in $\mathbb{K}$.

Here $2<\omega \leq 3$ is the exponent of matrix multiplication, see, e.g., [23].
Proof. First of all, shifting the variable $X$ by $a$ in $P$ takes $\mathcal{O}\left(D_{Y} \mathrm{M}\left(D_{X}\right)\right)$ operations in $\mathbb{K}$ by means of $[4, \mathrm{Ch} .1, \S 2]$ since the characteristic is zero. In order to reconstruct the coefficients of $M$ from their series expansions via Padé approximants [23, §5.9], a precision $(X-a)^{2 \eta+1}$ is sufficient. The idea is to follow the steps of Cockle's algorithm in $\mathbb{K}[[X-a]][Y] /(P)$ instead of $\mathbb{K}(X)[Y] /(P)$. The precision of the series
is initially set to $2 \eta+1+r \in \mathcal{O}(\eta)$ in order to accommodate the precision loss due to differentiation.

In Step 1, the computation of $V_{1}$ takes $\mathcal{O}\left(\mathrm{M}\left(D_{Y}\right) \mathrm{M}(\eta)\right)$ operations in $\mathbb{K}$ by [23, Thms. $15.10 \& 15.11]$. In Step 2 , each $V_{k+1}$ can be obtained from $V_{k}$ and $V_{1}$ by means of $\mathcal{O}\left(\mathrm{M}\left(D_{Y}\right) \mathrm{M}(\eta)\right)$ operations. In Step 3 , we start by computing the approximation of the relation (1) to precision $(X-a)$. This boils down to linear algebra over $\mathbb{K}$ in size $(r+1) \times D_{Y}$, hence a cost in $\mathcal{O}\left(r^{\omega-1} D_{Y}\right)$. From this we use Hensel-Newton lifting in order to recover relation (1) to precision $(X-a)^{2 \eta+1}$ with $\mathcal{O}\left(r^{\omega} \mathrm{M}(\eta)\right)$ operations in $\mathbb{K}$. The rational reconstruction of each $A_{i}$ and its shift back of $X$ by $-a \operatorname{cost} \mathcal{O}(\mathrm{M}(\eta) \log (\eta))$.

When $r=D_{Y}$ the cost of the computation of $M$ drops to $\tilde{\mathcal{O}}\left(D_{X} D_{Y}^{\omega+2}\right)$ when fast multiplication is used. (Here the tilde indicates that the factors polynomial in $\log D_{X}, \log D_{Y}$ and their logarithms have been omitted.) Note that the results presented in this section extend to analyze the bit-complexity via $p$-adic approximation when $\mathbb{K}=\mathbb{Q}$.

## 3. Creative Telescoping for Algebraic Functions

In this section we search for differential operators associated to $P$ in $\mathbb{K}(X)\left\langle\partial_{X}\right\rangle$ with non-minimal orders but with degrees in $X$ much smaller than in the resolvent of $P$.
3.1. Principle. Let $\alpha(X)$ be a solution of $P(X, \alpha(X))=0$. If $P \in \mathbb{C}[X, Y]$, then by the residue theorem,

$$
\alpha(X)=\frac{1}{2 \pi i} \oint F(X, Y) d Y, \text { with } F(X, Y)=\frac{Y P_{Y}(X, Y)}{P(X, Y)}
$$

where the contour is for instance a circle containing $\alpha(X)$ and no other root of $P$. In this setting, creative telescoping can be seen as an algorithm based on differentiating under the integral sign and integrating by parts [2]. It consists in computing a linear differential operator of the form

$$
\begin{equation*}
\Lambda=A\left(X, \partial_{X}\right)+\partial_{Y} B\left(X, \partial_{X}, Y, \partial_{Y}\right) \tag{2}
\end{equation*}
$$

that cancels the rational function $F$. There, $\partial_{X}$ and $\partial_{Y}$ denote differentiation with respect to $X$ and $Y$. If such a $\Lambda$ exists, integrating $\Lambda\left(Y P_{Y} / P\right)=0$ shows that $A$ cancels $\alpha$. We now give an algebraic proof of this property.

Proposition 2. Under (H), suppose that an element in $\Lambda \in \mathbb{K}[X, Y]\left\langle\partial_{X}, \partial_{Y}\right\rangle$ is of the form (2) and that $\partial_{Y}^{k} \Lambda$ cancels the rational function $F=Y P_{Y} / P$ for some integer $k \geq 0$. Then, the operator $A$ is associated to $P$.

Proof. The element $F=Y P_{Y} / P$ of $\mathbb{K}(X, Y)$ can be expanded as a Laurent series in $\mathbb{K}(X, \alpha)((Y-\alpha))$ :

$$
F=\frac{\alpha}{Y-\alpha}+\sum_{i \geq 0} \gamma_{i}(Y-\alpha)^{i}
$$

with coefficients $\gamma_{i}$ in $\mathbb{K}(X, \alpha)$. On the other hand the action of derivations in $X$ and $Y$ on a summand $\gamma(Y-\alpha)^{i}$, for $\gamma$ in $\mathbb{K}(X, \alpha)$, is given for any $i \in \mathbb{Z}$ by

$$
\begin{gathered}
\partial_{Y}\left(\gamma(Y-\alpha)^{i}\right)=i \gamma(Y-\alpha)^{i-1} \\
\partial_{X}\left(\gamma(Y-\alpha)^{i}\right)=\gamma^{\prime}(Y-\alpha)^{i}-i(Y-\alpha)^{i-1} \alpha^{\prime} .
\end{gathered}
$$

Thus no term in $(Y-\alpha)^{-1}$ remains after action of $\partial_{Y}$. In this way, applying $\partial_{Y}^{k} \Lambda$ to $F$ and considering the resulting Laurent series yields

$$
0=\partial_{Y}^{k} \Lambda(F)=\partial_{Y}^{k} \Lambda\left(\frac{\alpha}{Y-\alpha}\right)=(-1)^{k} \frac{A\left(X, \partial_{X}\right)(\alpha)}{(Y-\alpha)^{k+1}}
$$

which proves that $A$ annihilates $\alpha$.
3.2. Quadratic Bound. We now prove Thm. 2 by considering the degrees of the numerators and denominators of the successive derivatives of the rational function $F=Y P_{Y} / P$. By Leibniz's rule, for any nonnegative integers $N_{X}$ and $N_{\partial}$, the $r_{N_{X}, N_{\partial}}:=\left(N_{X}+1\right)\binom{N_{\partial}+2}{2}$ rational functions

$$
X^{i} \partial_{X}^{j} \partial_{Y}^{k} F, \quad 0 \leq i \leq N_{X}, \quad 0 \leq j+k \leq N_{\partial}
$$

are contained in the $\mathbb{K}$-vector space of dimension at most $s_{N_{X}, N_{\partial}}:=\left(D_{Y}\left(N_{\partial}+1\right)+\right.$ 1) $\left(N_{X}+1+D_{X}\left(N_{\partial}+1\right)\right)$ spanned by the elements

$$
\frac{X^{i} Y^{j}}{P^{N_{\partial}+1}}, \quad 0 \leq i \leq N_{X}+D_{X}\left(N_{\partial}+1\right), 0 \leq j \leq D_{Y}\left(N_{\partial}+1\right)
$$

With the values of $N_{X}$ and $N_{\partial}$ given in the theorem, the difference $r_{N_{X}, N_{\partial}}-s_{N_{X}, N_{\partial}}$ becomes

$$
\left(12 D_{Y}^{2}-7 D_{Y}-1\right) D_{X}+12 D_{Y}^{2}+8 D_{Y}
$$

which is positive for positive $D_{Y}$.
This proves the existence of a nonzero differential operator $\Lambda\left(X, \partial_{X}, \partial_{Y}\right)$ with degree in $X$ at most $N_{X}$ and total degree in $\partial_{X}$ and $\partial_{Y}$ at most $N_{\partial}$ that annihilates $F$. Let now $k \leq N_{\partial}$ be such that $\Lambda$ rewrites $\partial_{Y}^{k}\left(A\left(X, \partial_{X}\right)+\partial_{Y} B\left(X, \partial_{X}, \partial_{Y}\right)\right)$ with $A \neq 0$. The conclusion follows from Prop. 2.
3.3. Refined Bound. We now prove Thm. 3. The bound of the previous section relies on the elimination of $Y$ in the left ideal annihilating $F$ in the Weyl algebra $W_{X, Y}(\mathbb{K})$. In the context of creative telescoping, this is known as Zeilberger's slow algorithm, which relies on holonomic D-modules [25]. There is also a fast algorithm for hyperexponential integrands [2] that consists in looking for an operator of the form (2). Letting $B$ in (2) depend on $Y$ also allows for the existence of operators of smaller order. In this context, the only good bounds [3] we know of do not apply to rational integrands. Instead, we give a construction inspired by one of Takayama [20, Thm. 2.4 and Appendix 6] showing that any differentiably finite function is holonomic.

For each integer $d \geq 0$, we consider the vector space $V_{d}$ of rational functions $G / P^{d+1}$, with polynomials $G \in \mathbb{K}[X, Y]$ satisfying

$$
\begin{gathered}
\operatorname{deg} G \leq(d+1) D+d, \quad \operatorname{deg}_{X} G \leq(d+1) D_{X}+d \\
\operatorname{deg}_{Y} G \leq(d+1) D_{Y}
\end{gathered}
$$

This is defined in such a way that an induction on $d$ shows that the vector space $F_{d}$ generated by the rational functions

$$
\begin{equation*}
X^{i} \theta_{X}^{j} \cdot F, \quad 0 \leq i, j \leq d \tag{3}
\end{equation*}
$$

is a subspace of $V_{d}$, and that so is $\partial_{Y} \cdot V_{d-1}$. We are going to establish that for $d$ large enough,

$$
\begin{equation*}
\operatorname{dim}_{\mathbb{K}} F_{d}+\operatorname{dim}_{\mathbb{K}}\left(V_{d} \cap \partial_{Y} \cdot V_{d+1}\right)>\operatorname{dim}_{\mathbb{K}} V_{d} \tag{4}
\end{equation*}
$$

From here the result follows: both $F_{d}$ and $V_{d} \cap \partial_{Y} \cdot V_{d+1}$ are subspaces of $V_{d}$ whose dimensions are such that they have a nonzero intersection. Any element of this intersection is of the form $A F=\partial_{Y}\left(G / P^{d+2}\right)$ with $A F \in F_{d}$. By the same arguments as in the proof of Prop. 2, we deduce that $A$ is associated to $P$.

We now consider the dimensions in (4) more precisely.
The generators (3) are linearly independent since $P_{Y} \neq 0$ : the reduced form of $\theta_{X}^{i} \cdot F$ has denominator $P^{i+1}$, which cannot be eliminated by multiplication by a polynomial in $X$ only. Therefore $\operatorname{dim}_{\mathbb{K}} F_{d}=(d+1)^{2}$.

Write $P$ in the form $X^{D_{X}} \phi+\tilde{P}$ where $\phi$ depends only on $Y$. For $d \geq 0$, define the vector space $W_{d}$ of rational functions $F / P^{d}$, with polynomials $F \in \mathbb{K}[X, Y]$ of the form $F=c X^{\left(D_{X}+1\right) d+1} \phi^{d}+H$ where $c$ is in $\mathbb{K}$ and $H$ satisfies

$$
\begin{gather*}
\operatorname{deg} H \leq(D+1) d+1, \quad \operatorname{deg}_{X} H \leq\left(D_{X}+1\right) d  \tag{5}\\
\operatorname{deg}_{Y} H \leq D_{Y} d \tag{6}
\end{gather*}
$$

The decomposition of $F$ induces $W_{d} \subseteq V_{d+1}$ directly. Then, from

$$
\begin{aligned}
P^{d+1} \partial_{Y}\left(\frac{F}{P^{d}}\right)= & c d \phi^{d-1} X^{\left(D_{X}+1\right) d+1}\left(\tilde{P} \partial_{Y} \phi-\phi \partial_{Y} \tilde{P}\right) \\
& +P \partial_{Y} H-d H \partial_{Y} P
\end{aligned}
$$

we obtain the inclusion $\partial_{Y} \cdot W_{d} \subseteq V_{d}$. It follows that $\partial_{Y} \cdot W_{d} \subseteq V_{d} \cap \partial_{Y} \cdot V_{d+1}$, and that Eq. (4) holds as soon as

$$
\begin{equation*}
\operatorname{dim}_{\mathbb{K}} F_{d}+\operatorname{dim}_{\mathbb{K}} \partial_{Y} \cdot W_{d}>\operatorname{dim}_{\mathbb{K}} V_{d} \tag{7}
\end{equation*}
$$

In order to compute the dimension of $\partial_{Y} \cdot W_{d}$, we first consider the kernel of $\partial_{Y} \mid W_{d}$, which is precisely $\mathbb{K}[X] \cap W_{d}$. In view of the degree constraints (5-6), this corresponds to numerators $F=P^{d} Q$ with $Q \in \mathbb{K}[X]$ and

$$
\operatorname{deg}_{X} Q \leq\left(D_{X}+1\right) d+1-D_{X} d=d+1
$$

so that $\operatorname{dim}_{\mathbb{K}}\left(\partial_{Y} \cdot W_{d}\right)=\operatorname{dim}_{\mathbb{K}} W_{d}-(d+2)$.
The dimensions of $V_{d}$ and $W_{d}$ are given by the possible supports of the numerators of their elements. Lemma 4 below thus provides us with:

$$
\begin{gathered}
\operatorname{dim}_{\mathbb{K}} V_{d}=(d+1)\left(D_{X}+1\right)\left((d+1) D_{Y}+1\right)-\binom{d \Delta+\Delta+1}{2} \\
\operatorname{dim}_{\mathbb{K}} W_{d}=\left(\left(D_{X}+1\right) d+1\right)\left(D_{Y} d+1\right)-\binom{d \Delta}{2}+1
\end{gathered}
$$

Injecting the three dimensions in Eq. (7) finally gives the following condition:

$$
\begin{aligned}
d^{2}-\left(2 D_{X} D_{Y}+D_{Y}-\Delta^{2}-\Delta-1\right) d & \\
& +\frac{(\Delta-1) \Delta}{2}-\left(D_{X}+1\right)\left(D_{Y}+1\right)>0
\end{aligned}
$$

Call $\psi$ the value obtained on substituting the bound of the theorem for $d$ in the left-hand side:

$$
\psi=(3 D y-1) D x+D y-3 \Delta(\Delta+1) / 2+2
$$

From $D_{Y} \geq 1$ and $0 \leq \Delta \leq \min \left(D_{X}, D_{Y}\right)=: m \geq 1$ it follows that

$$
\psi \geq(3 m-1) m+m-3 m(m+1) / 2+2=3 m(m-1) / 2+2>0 .
$$

Lemma 4. The number of monomials in $X$ and $Y$ of total degree at most $\delta$, degree in $X$ at most $\delta_{X}$, and degree in $Y$ at most $\delta_{Y}$ is $\left(\delta_{X}+1\right)\left(\delta_{Y}+1\right)-\binom{\delta_{X}+\delta_{Y}-\delta+1}{2}$, whenever $\delta_{X}+\delta_{Y}+1 \geq \delta \geq \max \left(\delta_{X}, \delta_{Y}\right)$.
3.4. Construction of Differential Operators. Given a polynomial $P \in \mathbb{K}[X, Y]$ satisfying Hypothesis (H), we have proved in the previous sections the existence of several operators in $\mathbb{K}[X]\left\langle\partial_{X}\right\rangle$ associated to $P$. The aim of this section is to study the efficient construction of such operators. This is done under an additional, but mild, hypothesis (see ( $H^{\prime}$ ) below). As a by-product, we obtain in Th. 4 an alternative unified proof of slightly weaker versions of Thms. 1, 2 and 3.
Direct Algorithms. The proof of Thm. 2 boils down to a kernel computation of a matrix of size $\sim D_{X} D_{Y}^{3}$. Even using FFT for (univariate and bivariate) polynomial multiplication and a Wiedemann-like approach for the linear algebra step, this computation requires $\sim D_{X}^{2} D_{Y}^{6}$ operations in $\mathbb{K}$.

In the sequel we provide an algorithm with better cost, which is based on PadéHermite approximation of series expansions of an algebraic solution and its derivatives.
Notation and Assumptions. We assume in the rest of this section that the following hypothesis holds.

Hypothesis (H'). Besides Hypothesis (H),
$\left(H_{a}\right) D_{Y} \geq 2$ and $P$ is absolutely irreducible in $\mathbb{K}[X, Y]$ (which means that $P$ is irreducible when seen over the algebraic closure $\overline{\mathbb{K}}$ of $\mathbb{K}$ ), and
$\left(H_{b}\right) x=0$ annihilates neither the discriminant $\Delta_{Y}(X)$ nor the leading coefficient $P_{D_{Y}}(X)$ of $P$, seen in $\mathbb{K}[X][Y]$.

To clarify the status of these assumptions, note that $\left(H_{a}\right)$ is a global property of $P$ which is crucial in the subsequent proof that our algorithm works correctly. In contrast, the local property $\left(H_{b}\right)$ is purely of technical nature, being similar with that on lucky points in Section 2.3. $\left(H_{b}\right)$ can be ensured deterministically and in good complexity after a change of coordinates $X \leftarrow X+a$ with $a \in \mathbb{K}$ not a zero of $P_{D_{Y}} \Delta_{Y}$.
Principle of the Algorithm. By Hypothesis $\left(H_{b}\right)$, the polynomial $p(Y):=P(0, Y)$ is separable of degree $D_{Y}$. Hensel's lemma ensures that the $D_{Y}$ roots $\alpha_{i}$ of $P\left(X, \alpha_{i}\right)=$ 0 are all in $\overline{\mathbb{K}}[[X]]$. To avoid the factorization of $p(Y)$, we will work in the quotient algebra $\mathbb{A}:=\mathbb{K}[Y] /(p(Y))$. Let $y$ denote the residue class of $Y$ in $\mathbb{A}$. By Hensel's lemma again, there exists a unique series $\varphi \in \mathbb{A}[[X]]$ such that $\varphi(0)=y$ and $P(X, \varphi(X))=0$. Moreover, the expansion of $\varphi$ to any prescribed precision can be efficiently computed by means of the Newton operator.

To give a unified presentation, let us denote by $\partial$ one of the derivation operators $\partial_{X}$ and $\theta_{X}$. The principle of our algorithm is to compute enough terms of $\varphi(X)$, so that the coefficients of an associated operator can be recovered by Padé-Hermite (abbrev. PH) approximation of the vector $\left(\varphi, \partial \varphi, \partial^{2} \varphi, \ldots\right)^{t}$ of the first derivatives of $\varphi$ with respect to $\partial$. Thus, we obtain first an operator with coefficients in $\mathbb{A}[X]$, from which we extract an operator in $\mathbb{K}[X]\langle\partial\rangle$ associated to $P$. We give in Figure 1 a detailed presentation of the algorithm. The main difficulty of this approach is to provide suitable bounds on the orders of truncation. This is done in the next theorem, which encapsulates the main results of this section.

Theorem 4. Let $P \in \mathbb{K}[X, Y]$ be of bi-degree $\left(D_{X}, D_{Y}\right)$ and satisfy $\left(H^{\prime}\right)$.

```
    \(\operatorname{AlgToDiff}\left(P, B_{X}, B_{\partial}, \partial\right)\)
Input: \(P \in \mathbb{K}[X, Y], B_{X}, B_{\partial} \in \mathbb{N}, \partial \in\left\{\partial_{X}, \theta_{X}\right\}\)
Output: an operator \(L \in \mathbb{K}[X]\langle\partial\rangle \backslash\{0\}\) associ-
ated to \(P\), with \(\operatorname{deg}_{X}(L) \leq B_{X}\) and \(\operatorname{deg}_{\theta}(L) \leq B_{\partial}\).
1. Set \(\Sigma:=B_{X} B_{\partial}+B_{X}+B_{\partial}\).
2. Expand \(\varphi(X) \in \mathbb{A}[[X]]\) up to precision \(\Sigma+B_{\partial}\).
3. Compute \(Z_{i}=\partial^{i} \varphi\), for \(i=0,1, \ldots, B_{\partial}\) to precision \(\Sigma\).
4. Compute a PH approximant \(\left(\ell_{0}(X), \ldots, \ell_{B_{\partial}}(X)\right) \neq 0\)
    in \(\mathbb{A}[X]^{B_{\partial}+1}\) of \(\left(Z_{0}, \ldots, Z_{B_{\partial}}\right)\), of type \(\left(B_{X}, \ldots, B_{X}\right)\).
5. Write \(L=\ell_{0}(X)+\ell_{1}(X) \partial+\cdots+\ell_{B_{\partial}}(X) \partial^{B_{\partial}}\) as
    \(L_{0}+y L_{1}+\cdots+y^{D_{Y}-1} L_{D_{Y}-1}\), with \(L_{i} \in \mathbb{K}[X]\langle\partial\rangle\).
6. Return any \(L_{i} \neq 0\).
```

Figure 1. Construction of differential operators by Padé-Hermite approximation.
(1) $\operatorname{AlgToDiff}\left(P, 4 D_{X} D_{Y}^{2}, D_{Y}, \partial\right)$ outputs a non-zero differential operator in $\mathbb{K}[X]\langle\partial\rangle$ associated to $P$, using $\tilde{\mathcal{O}}\left(D_{X} D_{Y}^{\omega+3}\right)$ operations in $\mathbb{K}$.
(2) $\operatorname{AlgToDiff}\left(P, 5 D_{X} D_{Y}, 5 D_{Y}, \partial\right)$ outputs a non-zero differential operator in $\mathbb{K}[X]\langle\partial\rangle$ associated to $P$, using $\tilde{\mathcal{O}}\left(D_{X} D_{Y}^{\omega+2}\right)$ operations in $\mathbb{K}$.
(3) $\operatorname{AlgToDiff}(P, B, B, \partial)$ with $B=4 D_{X} D_{Y}+D_{Y}-2 D_{X}-2$ outputs a nonzero differential operator in $\mathbb{K}[X]\langle\partial\rangle$ associated to $P$, using $\tilde{\mathcal{O}}\left(D_{X}^{\omega+1} D_{Y}^{\omega+2}\right)$ operations in $\mathbb{K}$.

First we give a bound on the number of terms needed to reconstruct an operator.
Lemma 5. Under ( $H^{\prime}$ ), let $L \in \overline{\mathbb{K}}[X]\langle\partial\rangle \backslash\{0\}$ be of order at most $B_{\partial}$ in $\partial$, with coefficients of degrees at most $B_{X}$.
(a) If $\alpha_{0} \in \overline{\mathbb{K}}[[X]]$ is such that $P\left(X, \alpha_{0}\right)=0$ and $L \alpha_{0}=0$, then $L$ is associated to $P$.
(b) Let $\sigma=4 D_{X} D_{Y} B_{\partial}+B_{X} D_{Y}-2 D_{X} B_{\partial}$. If $\alpha_{0} \in \overline{\mathbb{K}}[[X]]$ is such that $P\left(X, \alpha_{0}\right)$ and $L \alpha_{0}$ are zero modulo $X^{\sigma}$, then $L$ is associated to $P$.

Proof. There exists $Q(X, Y) \in \overline{\mathbb{K}}[X, Y]$ with

$$
\begin{aligned}
& \operatorname{deg}_{X}(Q) \leq 2 B_{\partial} D_{X}+B_{X}-D_{X} \\
& \operatorname{deg}_{Y}(Q) \leq 2\left(D_{Y}-1\right) B_{\partial}-D_{Y}+2,
\end{aligned}
$$

such that $L \alpha=\frac{Q(X, \alpha)}{P_{Y}(X, \alpha)^{2 B_{\partial}-1}}$ for any root $\alpha \in \overline{\mathbb{K}(X)}$ of $P(X, \alpha)=0$. Indeed, if $L=$ $\ell_{0}(X)+\cdots+\ell_{B_{\partial}}(X) \partial^{B_{\partial}}$, we take $Q=\ell_{0} Y P_{Y}^{2 B_{\partial}-1}+\ell_{1} W_{1} P_{Y}^{2 B_{\partial}-2}+\cdots+\ell_{B_{\partial}} W_{B_{\theta}}$, where the $W_{i}(X, Y)$ are the polynomials introduced in Section 2.

Let $R(X) \in \overline{\mathbb{K}}[X]$ be the resultant of $P(X, Y)$ and $Q(X, Y)$ with respect to $Y$. To show that $L$ is associated to $P$, it is enough to prove that $R$ is identically zero. Indeed, since $P$ is irreducible in $\overline{\mathbb{K}}[X, Y]$, the nullity of $R$ implies that $P$ divides $Q[23$, Cor. 6.20], therefore $Q(X, \alpha)=0$, and thus $L \alpha=0$, for any $\alpha \in \overline{\mathbb{K}(X)}$ such that $P(X, \alpha)=0$.
(a) By definition, there exist polynomials $U, V \in \overline{\mathbb{K}}[X, Y]$ such that

$$
U(X, Y) P(X, Y)+V(X, Y) Q(X, Y)=R(X) .
$$

Evaluating this equality at $Y=\alpha_{0}$ yields $R(X)=0$, as wanted. The same reasoning shows that under the hypothesis of $(b)$ we have $R(X)=0 \bmod X^{\sigma}$. On the other hand, by Bézout's theorem, one can bound the degree of the resultant $R$ by $D_{X} \operatorname{deg}_{Y} Q+D_{Y} \operatorname{deg}_{X} Q$. Therefore $\operatorname{deg}(R)$ is at most

$$
\begin{aligned}
D_{X}\left(2 D_{Y} B_{\partial}-2 B_{\partial}-D_{Y}+2\right)+D_{Y}( & \left.2 B_{\partial} D_{X}+B_{X}-D_{X}\right) \\
& <4 D_{X} D_{Y} B_{\partial}+B_{X} D_{Y}-2 D_{X} B_{\partial}=\sigma
\end{aligned}
$$

since $D_{Y} \geq 2$. Finally $\operatorname{deg}(R) \leq \sigma$ and $R(X)=0 \bmod X^{\sigma}$ imply that $R$ is the zero polynomial.
of Thm. 4. In Step 4, the Padé-Hermite equality $\sum_{i} \ell_{i} Z_{i}=0 \bmod X^{\Sigma}$ can be viewed as an under-determined homogeneous linear system with $\Sigma$ equations in the $\Sigma+1$ unknown coefficients of the polynomials $\ell_{i}$. It thus admits at least one non-trivial solution in $\mathbb{A}^{\Sigma+1}$.

Next let $p=\prod_{i} p_{i}$ represent the irreducible factorization of $p$ in $\mathbb{K}[Y]$, and let $\mathbb{A}_{i}=\mathbb{K}[Y] /\left(p_{i}(Y)\right)$, so that $\mathbb{A}$ rewrites into the product of fields $\prod_{i} \mathbb{A}_{i}$. Let $\varphi^{(i)}$ be the image of $\varphi$ in $\mathbb{A}_{i}[[X]] \subseteq \overline{\mathbb{K}}[[X]]$ and let $L^{(i)}$ be the image of $L$ in $\mathbb{A}_{i}[X]\langle\partial\rangle \subseteq \overline{\mathbb{K}}[X]\langle\partial\rangle$. By construction $P\left(X, \varphi^{(i)}\right)=L^{(i)} \varphi^{(i)}=0 \bmod X^{\Sigma}$ holds in $\overline{\mathbb{K}}\left[[X]\right.$. Since all choices of $\left(B_{X}, B_{\partial}\right)$ as specified in Thm 4 ensure $\Sigma \geq \sigma$, it follows from Lemma $5(\mathrm{~b})$ that $L^{(i)}$ is associated to $P$ for all $i$. In other words, for all root $a \in \overline{\mathbb{K}}$ of $p(a)=0$, and all root $\alpha \in \overline{\mathbb{K}}[[X]]$ of $P(X, \alpha)=0$ we have that $L_{0} \alpha+a L_{1} \alpha+\cdots+a^{D_{Y}-1} L_{D_{Y}-1} \alpha=0$. We thus deduce that $L_{i}$ is associated to $P$ for all $i$, which concludes the proof of correctness.

We are left with the detailed description and complexity study of each step of the algorithm. Operations $(+, \times)$ in $\mathbb{A}$ have $\operatorname{cost} \mathcal{O}\left(\mathrm{M}\left(D_{Y}\right)\right)$ operations in $\mathbb{K}$, while inversion (when possible) takes $\mathcal{O}\left(\mathrm{M}\left(D_{Y}\right) \log D_{Y}\right)$. For simplicity, in what follows we suppose that FFT is used for polynomial multiplication, thus $\mathrm{M}(d) \in \tilde{\mathcal{O}}(d)$ for all $d$.

For the lifting step 2, we use the Newton iteration. By [7, Algo. 2, Prop. 3], one can compute, under $\left(\mathrm{H}^{\prime}\right)$, the first $\Sigma$ terms of $\varphi \in \mathbb{A}[[X]]$ in $\mathcal{O}\left(\Sigma D_{Y}{ }^{(\omega+1) / 2}+\right.$ $\left.\sqrt{D_{Y}} \mathrm{M}(\Sigma) \mathrm{M}\left(D_{Y}\right)\right) \in \tilde{\mathcal{O}}\left(B_{X} B_{\partial} D_{Y}^{(\omega+1) / 2}\right)$ operations in $\mathbb{K}$.

Since differentiating has linear complexity, Step 3 requires $\mathcal{O}\left(B_{\partial}^{2} B_{X}\right)$ operations in $\mathbb{A}$, hence at most $\tilde{\mathcal{O}}\left(B_{\partial}^{2} B_{X} D_{Y}\right)$ operations in $\mathbb{K}$.

To explain and estimate the cost of Step 5 , we use the approximation algorithm in [19], which ensures that one can compute a nonzero Padé-Hermite approximant of type $\left(B_{X}, \ldots, B_{X}\right)$ for $\left(Z_{0}, \ldots, Z_{B_{\partial}}\right)$ within $\tilde{\mathcal{O}}\left(B_{\partial}^{\omega} B_{X}\right)$ operations in $\mathbb{A}$, provided that $\mathbb{A}$ is a field, or more generally, if no division by a non-invertible element in $\mathbb{A}$ occurs during the computation.

In the general case, since $\mathbb{A}$ is not a field, we use the following strategy. We begin by applying the approximation algorithm over $\mathbb{A}$. If no forbidden division occurs, this algorithm outputs the desired operator $L$ with coefficients in $\mathbb{A}[X]$. If we encounter a nonzero $q \in \mathbb{A}$ that is not invertible, then by a gcd computation we obtain a proper factor $\tilde{p}$ of $p$ of degree at most $D_{Y} / 2$ such that $q$ is either invertible or zero modulo $\tilde{p}$. This takes $\tilde{\mathcal{O}}\left(D_{Y}\right)$ operations in $\mathbb{K}[23$, Th. 11.5]. We replace $\mathbb{A}$ by $\mathbb{K}[Y] /(\tilde{p}(Y))$ and restart the approximation algorithm. This procedure is repeated at most $\log D_{Y}$ times, until all the divisions are doable. The total cost of this computation is $\tilde{\mathcal{O}}\left(B_{\partial}^{\omega} B_{X} D_{Y}\right)$ operations in $\mathbb{K}$.

Finally, Steps 1, 5 and 6 do not require any arithmetic operation. Adding together these costs, and under the assumption $D_{Y} \leq B_{\partial}^{2}$ (which is satisfied by the choices in Thm. 4), we get that $\operatorname{Alg} \operatorname{ToDiff}\left(P, B_{X}, B_{\partial}, \partial\right)$ uses $\tilde{\mathcal{O}}\left(B_{\partial}^{\omega} B_{X} D_{Y}\right)$ operations in $\mathbb{K}$.

Probabilistic and Heuristic Versions. We now give a probabilistic version of Algorithm AlgToDiff, called AlgToDiffP. The idea is to compute a Padé-Hermite approximant over $\mathbb{K}$ rather than over $\mathbb{A}$. The probabilistic algorithm AlgToDiffP is obtained by replacing Steps 3-6 in AlgToDiff by:
$3^{\prime}$. Let $\varphi(X)=\varphi_{0}(X)+y \varphi_{1}(X)+\cdots+y^{D_{Y}-1} \varphi_{D_{Y}-1}(X)$.
Choose random $a_{0}, a_{1}, \ldots, a_{D_{Y}-1} \in \mathbb{K} \backslash\{0\}$ and compute
$Z_{0}=\sum_{i} a_{i} \varphi_{i}(X) \bmod X^{\Sigma}, Z_{i}=\partial^{i} Z_{0}, i=1, \ldots, B_{\partial}$.
$4^{\prime}$. Compute a PH approximant $\left(\ell_{0}(X), \ldots, \ell_{B_{\partial}}(X)\right)$ of type $\left(B_{X}, \ldots, B_{X}\right)$ of $\left(Z_{0}, Z_{1}, \ldots, Z_{B_{\partial}}\right)$ in $\mathbb{K}[X]^{B_{\partial}+1}$.
$5^{\prime}$. Return $L=\ell_{0}(X)+\ell_{1}(X) \partial+\cdots+\ell_{B_{\partial}}(X) \partial^{B_{\partial}}$.
Algorithm $\operatorname{Alg} \operatorname{ToDiffP}\left(P, 5 D_{X} D_{Y}, 5 D_{Y}, \partial\right)$ has complexity in $\tilde{\mathcal{O}}\left(D_{X} D_{Y}^{(\omega+5) / 2}\right)$, while $\operatorname{Alg} \operatorname{ToDiffP}(P, B, B, \partial)$ with $B$ as in Thm. 4 uses $\tilde{\mathcal{O}}\left(D_{X}^{2} D_{Y}^{(\omega+5) / 2}+\left(D_{X} D_{Y}\right)^{\omega+1}\right)$ ops in $\mathbb{K}$. Compared to the deterministic algorithms, this improves the exponent of $D_{Y}$.

Another improvement comes from the remark that the algorithms derived from AlgToDiff and AlgToDiffP do not return operators of order and degree as small as predicted by Thms. 2 and 3. The reason for this is technical: we need the values $B_{X}, B_{\partial}$ it uses in order to prove that the operator constructed by PadéHermite approximation cancels $\varphi$ and not merely its truncation. Replacing arguments $B_{X}, B_{\partial}$ by the better bounds from Thms. 2 and 3 we obtain heuristic modifications $\operatorname{Alg} \operatorname{ToDiffP}\left(P, 3 D_{X} D_{Y}+6 D_{Y}, 6 D_{Y}, \partial\right)$ and $\operatorname{AlgToDiffP}(P, B, B, \partial)$ with $B$ as in Thm. 3. So far, we do not have a proof that these heuristic always return a correct output, but they behave very well in practice (see $\S 4$ ).
3.5. Unrolling the Recurrence. The operators constructed in Th. 4 can be used to compute power series expansions of algebraic functions asymptotically faster than Newton iteration. The following result quantifies the improvement.
Theorem 5. Let $\psi \in \mathbb{L}[[X]]$ be such that $P(X, \psi(X))=0$, where $\mathbb{L}=\mathbb{K}(\alpha)$ is a finite extension of $\mathbb{K}$. Then, under Hypothesis $\left(H^{\prime}\right)$, the first $N$ terms of $\psi$ can be computed in $\mathcal{O}\left(N D_{X} D_{Y}\left([\mathbb{L}: \mathbb{K}]+\mathrm{M}\left(D_{Y}\right) / D_{Y}\right)\right)$ operations in $\mathbb{K}$, as $N \rightarrow \infty$.

In the case when $\mathbb{L}=\mathbb{K}$, the complexity we get is therefore linear in $N$ and quadratic in the degree if fast multiplication is used. A better complexity $O(N D)$ was announced by Chudnovsky \& Chudnovsky in [8, 9]. While the dependency in $N$ is correct, they seem to have overlooked the degree growth of the coefficients of the differential resolvent, that gives a cubic dependency in $D$ for their algorithm.

Our algorithm is presented in the proof below.
Proof. The first part of the computation is independent of $N$. We begin by calling $\operatorname{AlgToDiff}\left(P, B_{X}, B_{\partial}, \partial\right)$ with $B_{X}=5 D_{X} D_{Y}, B_{\partial}=5 D_{Y}$ and $\partial=\theta_{X}$. We then convert the operator into a recurrence of the form

$$
r_{0}(n) u_{n}+r_{1}(n) u_{n+1}+\cdots+r_{B_{X}}(n) u_{n+B_{X}}=0, \quad \text { for } n \geq 0
$$

with coefficients $r_{i}(X) \in \mathbb{K}[X]$ of degree at most $B_{\partial}$. Let $\rho$ be the largest nonnegative integer root of the polynomial $r_{B_{X}}(X)$, or -1 if no such root exists. Thus
for any $n>\rho, u_{n+B_{X}}$ can be computed from the previous ones using the recurrence. Note that we also have at our disposal the initial segment $\sum_{n=0}^{B_{X}-1} u_{n} X^{n}$ of $\varphi(X)$. If $\rho \neq-1$, we use Newton iteration to extend this initial segment up to the index $\rho+B_{X}$ if necessary. Using the minimal polynomial of $\alpha$, we then deduce the corresponding segment of $\psi$.

Only the rest of the computation depends on $N$. It consists in unrolling the recurrence efficiently for $n$ from $\rho+1$ to $N-B_{X}-1$. Since the coefficients $r_{i}$ belong to $\mathbb{K}$, so do the values $r_{i}(n)$ and thus each new coefficient $u_{n}$ of $\psi$ can be obtained using $\mathcal{O}\left(B_{X}[\mathbb{L}: \mathbb{K}]\right) \in \mathcal{O}\left(D_{X} D_{Y}[\mathbb{L}: \mathbb{K}]\right)$ operations in $\mathbb{K}$, provided that the values taken by the coefficients $r_{i}(X)$ at the points $\rho+1, \ldots, N-B_{X}-1$ have been computed and stored.

This computation can be done separately for each coefficient $r_{i}(X)$ using fast multipoint evaluation on $\rho+1, \rho+2, \ldots, \rho+B_{\partial}$ followed by fast extrapolation on arithmetic sequences. The cost of all the fast multipoint evaluation steps is in $\mathcal{O}\left(B_{X} \mathrm{M}\left(B_{\partial}\right) \log B_{\partial}\right)$ operations in $\mathbb{K}$, while that of all the fast extrapolation steps is in $\mathcal{O}\left(B_{X} N \mathrm{M}\left(B_{\partial}\right) / B_{\partial}\right)$ operations in $\mathbb{K}$, see e.g. [23, 6].

We note that Hypothesis ( $H^{\prime}$ ) is not strictly necessary in practice: the same method handles singular cases as well, provided power series solutions exist at the point of interest.

## 4. Experiments

4.1. Timings. We have implemented the algorithm of $\S 2.3$ in the Magma computer algebra system V2.11-14 [5]. In Table 1 we report on our experiments with random dense bivariate polynomials over the prime field with 9973 elements on a Pentium (M) 1.8 GHz processor. The column "ser." indicates the time in seconds of our algorithm, while the column "rat." is for Magma's routine DifferentialOperator, where we removed the characteristic 0 requirement. The latter routine implements Cockle's algorithm in $\mathbb{K}(X)$ and has been contributed by A. van der Waall. The column " $\eta$ " indicates the value of $\eta$ and the last column " $\operatorname{deg}_{X}(M)$ " provides us with the actual degree in $X$ of the minimal operator. We observe that Cockle's algorithm is always faster when run over the series and that there is still room for improvement on our degree bound $\eta$.

In Table 2, we compare the practical behaviour of our heuristic algorithm from $\S 3.4$ for expanding algebraic series defined by polynomials of bi-degree $\left(1, D_{Y}\right)$ with that of Newton's iteration (column Newton). Column AlgToRec contains the cumulated timings of the heuristic $\operatorname{Alg} \operatorname{ToDiffP}(P, B, B, \partial)$ (with $B$ as in Th. 3) and of the conversion of a differential operator into a recurrence. Column $N$ gives the number of terms of the expansion, while column unroll shows the timings for unrolling the recurrence. One motivation for the computation when $D_{X}=1$ is its application in univariate polynomial root finding by homotopy [10, § 10].

### 4.2. On the Optimality of the Bounds.

Conjectures. Empirically, the degree of the differential resolvent seems to be bounded by $D\left(D^{2}-5 D / 2+5 / 2\right)$ for polynomials of total degree $D$ and $2\left(D^{2}-3 D+3\right)$ for polynomials of bi-degree $(D, D)$. This is to be compared with the degree bound $D\left(4 D^{2}-11 / 2 D+5 / 2\right)$ given by Thm. 1 .

Also, intensive experiments suggest that the bound from Thm. 3 is only slightly pessimistic: for a polynomial of total degree $D$, there often exists a recurrence

Table 1. Cockle's algorithm

| $\left(D_{X}, D_{Y}\right)$ | ser. | rat. | $\eta$ | $\operatorname{deg}_{X}(M)$ |
| :---: | :---: | :---: | :---: | :---: |
| $(1,1)$ | 0.002 | 0.002 | 2 | 2 |
| $(2,2)$ | 0.003 | 0.004 | 17 | 10 |
| $(3,3)$ | 0.02 | 0.03 | 69 | 36 |
| $(4,4)$ | 0.10 | 0.16 | 182 | 92 |
| $(5,5)$ | 0.47 | 0.98 | 380 | 190 |
| $(6,6)$ | 1.86 | 4.56 | 687 | 342 |
| $(7,7)$ | 5.80 | 16.5 | 1127 | 560 |
| $(8,8)$ | 15.5 | 49.9 | 1724 | 856 |
| $(9,9)$ | 38.0 | 138 | 2502 | 1242 |
| $(10,10)$ | 72.7 | 340 | 3485 | 1730 |

Table 2. Newton's algorithm vs. AlgToDiff

| $\left(D_{X}, D_{Y}\right)$ | $N$ | AlgToRec | unroll | Newton |
| :---: | :---: | :---: | :---: | :---: |
| $(1,8)$ | $2^{10}$ | 6.66 | 0.22 | 1.75 |
| $(1,8)$ | $2^{11}$ | 6.66 | 0.44 | 3.91 |
| $(1,8)$ | $2^{12}$ | 6.66 | 0.89 | 8.83 |
| $(1,8)$ | $2^{13}$ | 6.66 | 1.8 | 18.84 |
| $(1,8)$ | $2^{14}$ | 6.66 | 3.67 | 42.53 |
| $(1,8)$ | $2^{15}$ | 6.66 | 7.38 | 94.99 |
| $(1,9)$ | $2^{10}$ | 10.54 | 0.26 | 2.14 |
| $(1,9)$ | $2^{11}$ | 10.54 | 0.44 | 4.35 |
| $(1,9)$ | $2^{12}$ | 10.54 | 1.07 | 9.31 |
| $(1,9)$ | $2^{13}$ | 10.54 | 2.17 | 20.64 |
| $(1,9)$ | $2^{14}$ | 10.54 | 4.36 | 43.44 |
| $(1,9)$ | $2^{15}$ | 10.54 | 8.75 | 97.64 |

of order $D^{2}-2$ with coefficients of degree at most $D^{2}-1$. For a polynomial of bi-degree $\left(D_{X}, D_{Y}\right)$, experiments suggest the bound $2 D_{X} D_{Y}-2-\left(D_{X}-D_{Y}\right)$ if $D_{Y}>1$ and $D_{X}+1$ if $D_{Y}=1$.
Lower Bound. We note also that $D(D-1)$ is the degree of the discriminant of $P$. When this discriminant is square-free, the differential equation satisfied by all solutions of $P$ has to be singular at those points and therefore its leading coefficient cannot have a smaller degree. Thus our bound on the order of the recurrence is asymptotically optimal.

An example illustrating $D(D-1)$ as a lower bound is provided by $P(X, Y)=$ $Y^{D}-Y+X^{D}$. Then $\alpha=X^{D}+X^{D^{2}}+\mathcal{O}\left(X^{D^{2}}\right)$ is a root of of $P$ seen in $\mathbb{Q}[[X]][Y]$. It has a sequence of $D(D-1)-1$ zero coefficients before a non-zero one. Thus, the order of the recurrence satisfied by the coefficients of $\alpha$ is at least $D(D-1)$.

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