



Validation of Real Time Applications

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Validation of Real Time Applications

Françoise Simonot - Lion

**Zhejiang University
China**

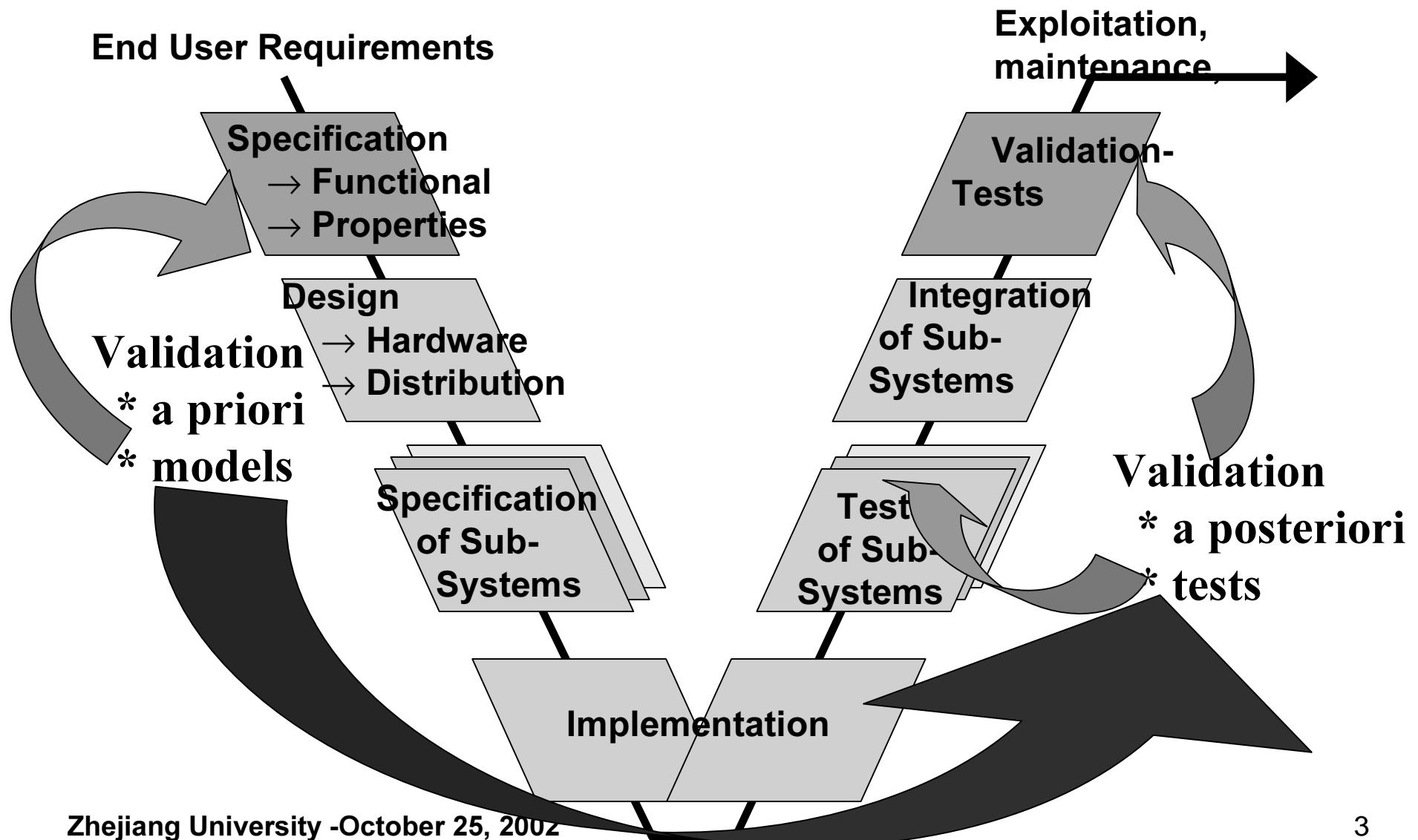
October 25, 2002



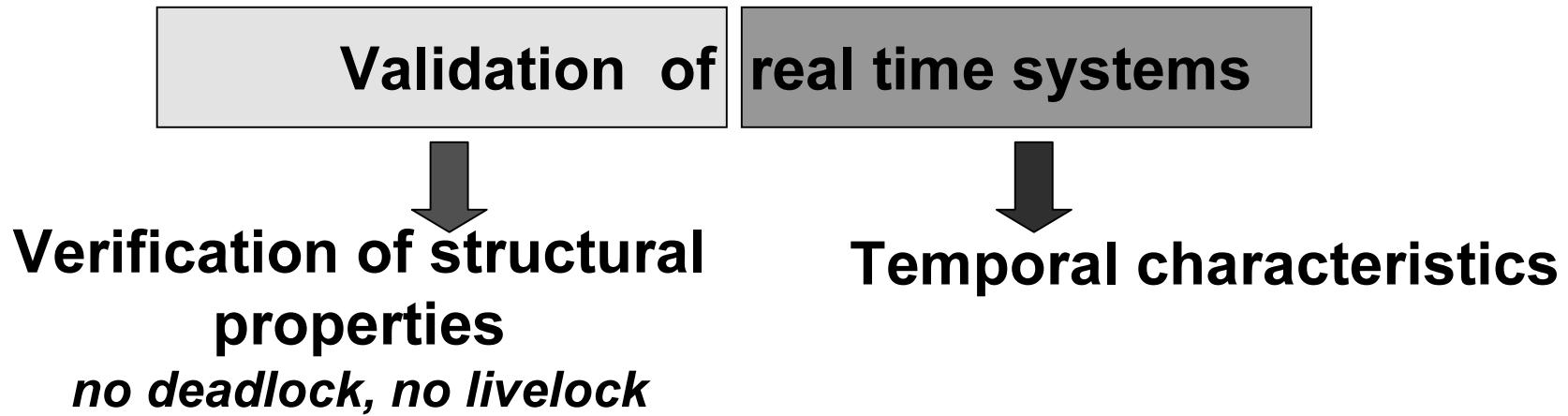
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- 2 ➔ Timed Input Output State Machine (TIOSM)
an adequate formalism
- 3 ➔ Validation a posteriori (Integration test)
how to be efficient
- 4 ➔ Validation a priori
how to master the complexity
- 5 ➔ Conclusions

Validation : when - what - how



Validation : when - what - how



Model

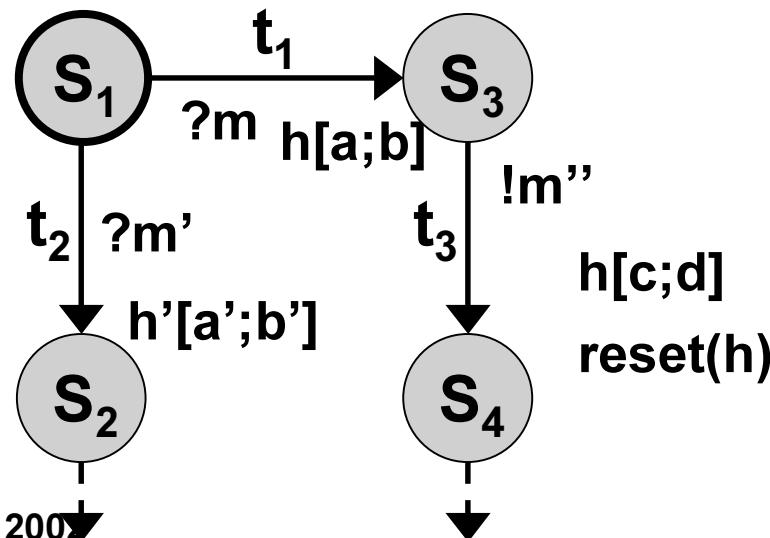
- state machines / temporal attributes
- deterministic approach

“Timed automata”

2 - Timed Input Output State Machine (TIOSM) *an adequate formalism*

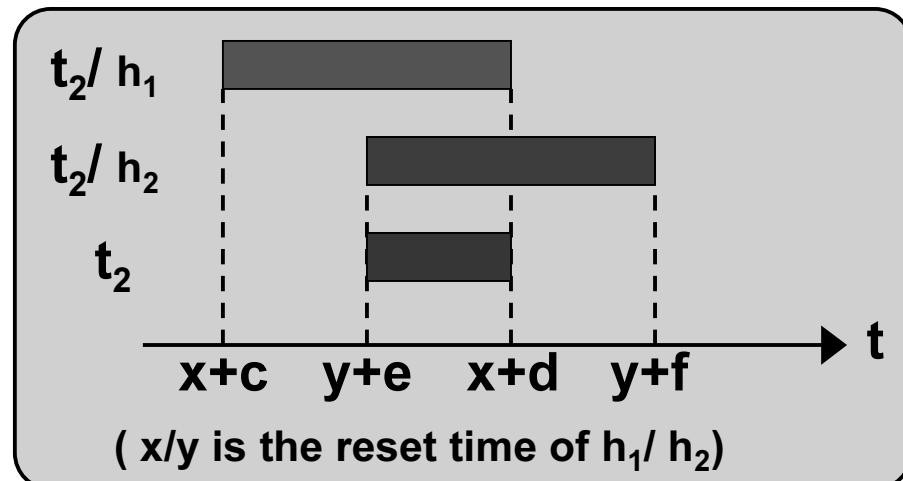
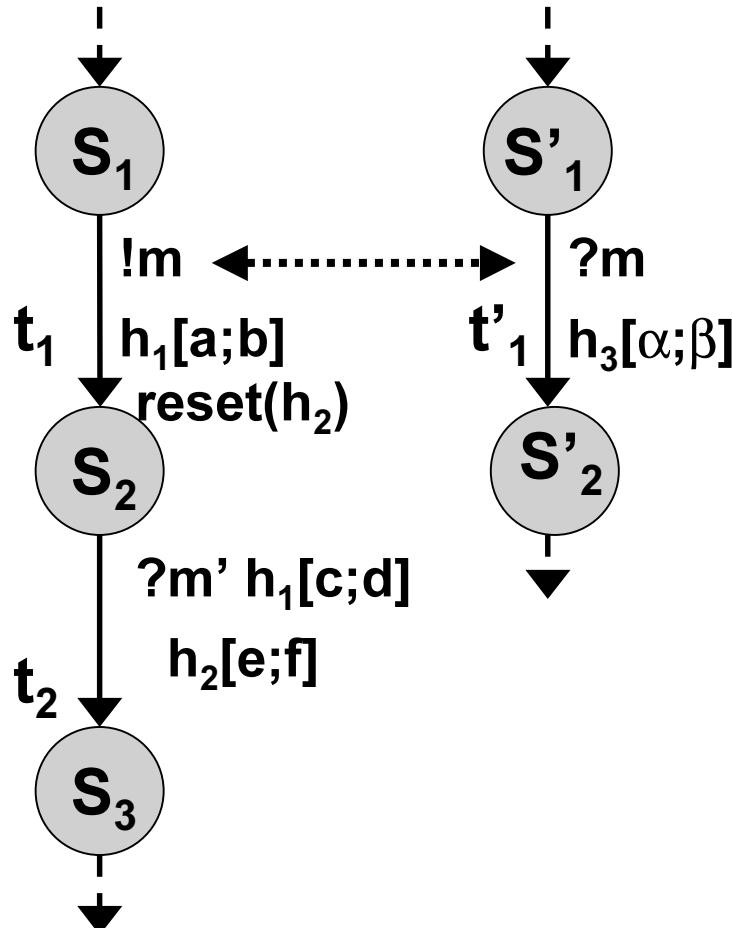
> TIOSM $T=(S, L, C, s_0, \varepsilon)$

- S set of states / ε set of transitions
- s_0 is the initial state
- L set of messages
- C set of clocks



2 - Timed Input Output State Machine (TIOSM) *an adequate formalism*

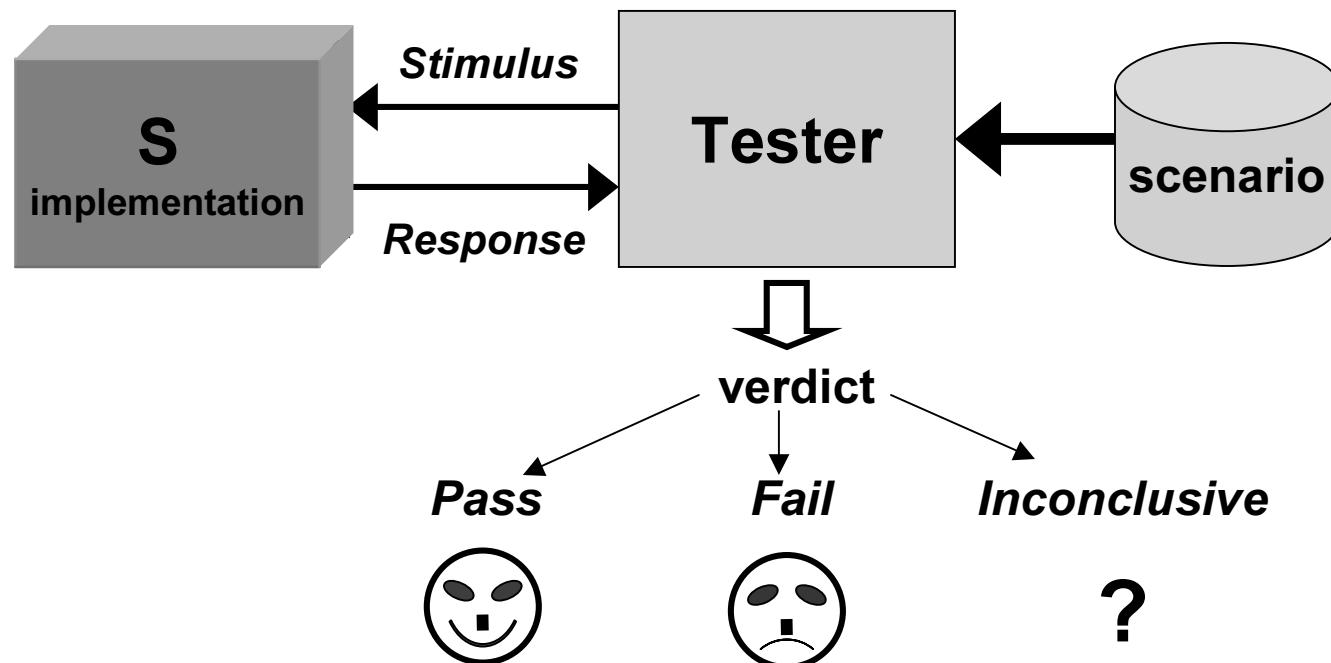
> Temporal behavior : an exemple



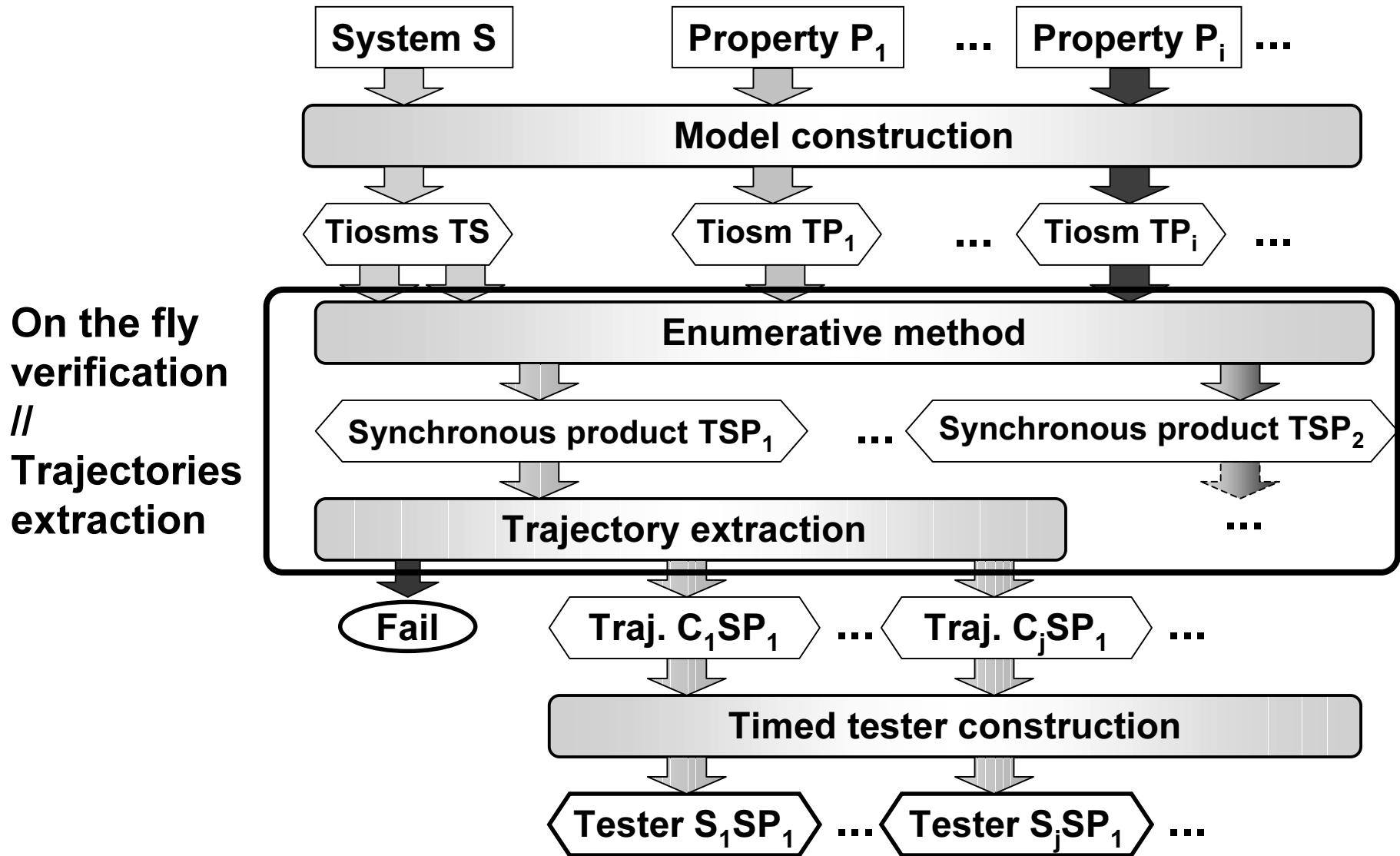
3 - Validation a posteriori (Integration test)

Implementation of S

Verification of Temporal Interoperability Properties
Testing approach

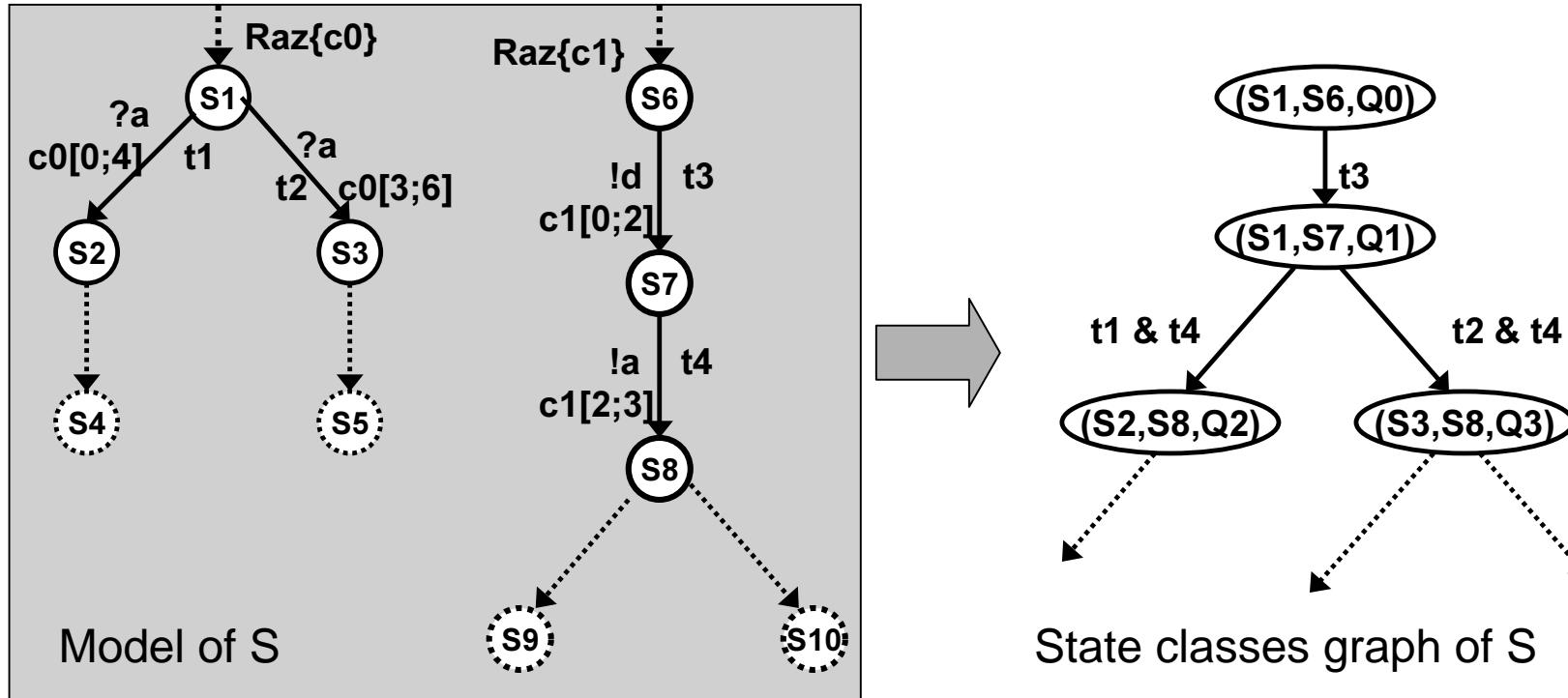


3 - Validation a posteriori (Integration test)



3 - Validation a posteriori (Integration test)

Construction of the state classes graph (cf. Berthomieu&Diaz) :

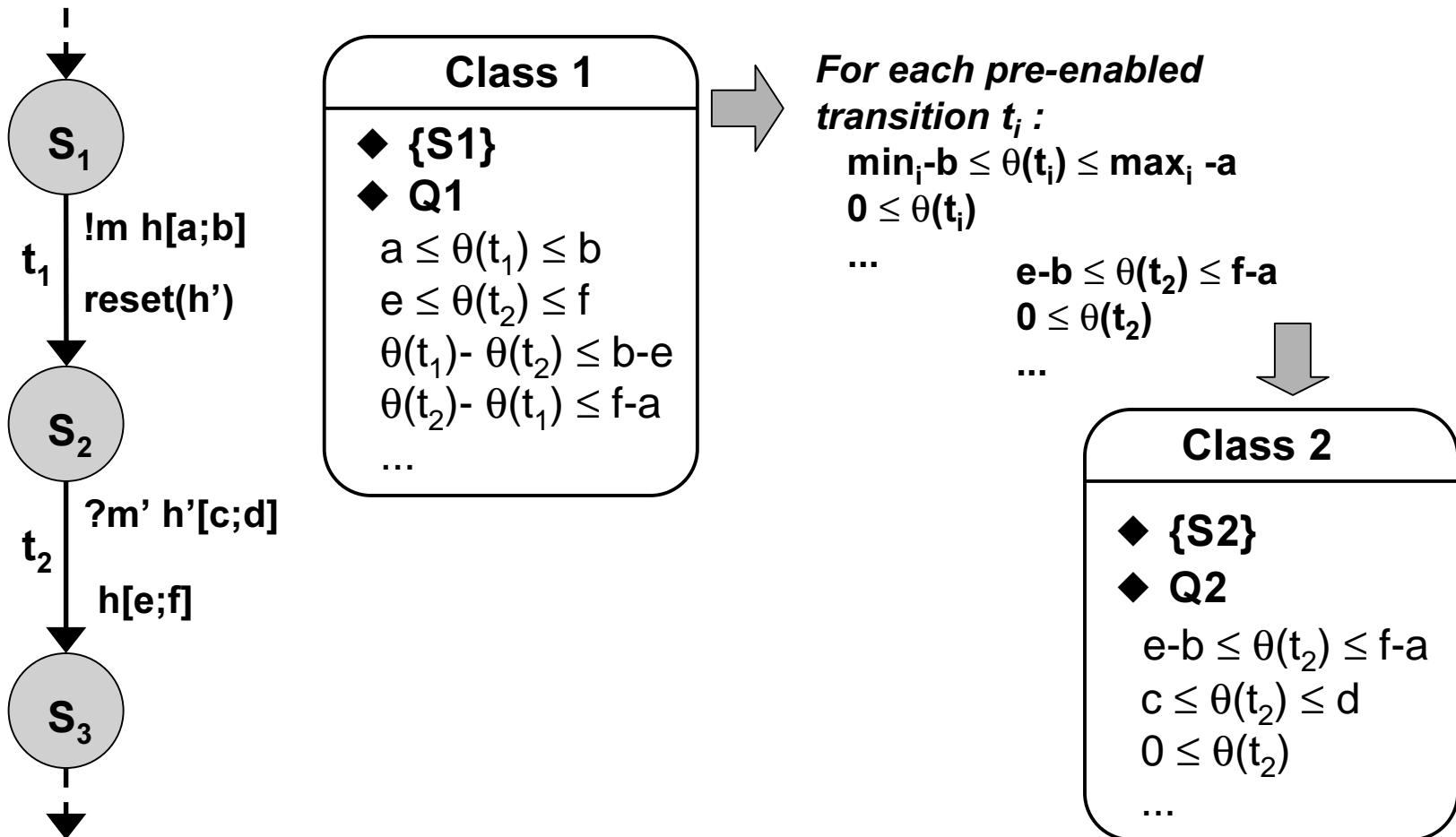


S_i, S_j : current states of the TIOSM

Q_k : inequations system

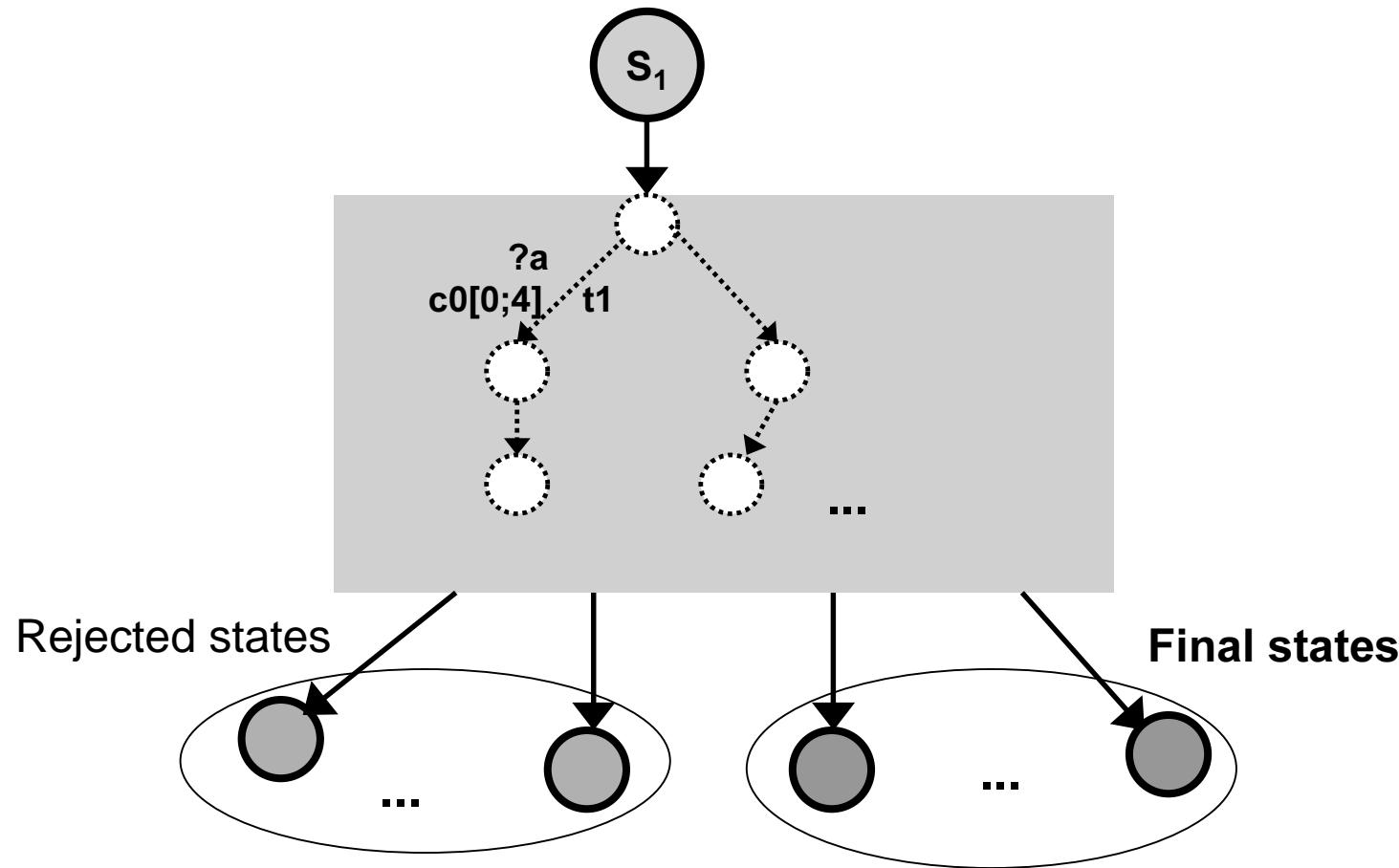
3 - Validation a posteriori (Integration test)

Construction of the inequations system / Pre-enabled transitions



3 - Validation a posteriori (Integration test)

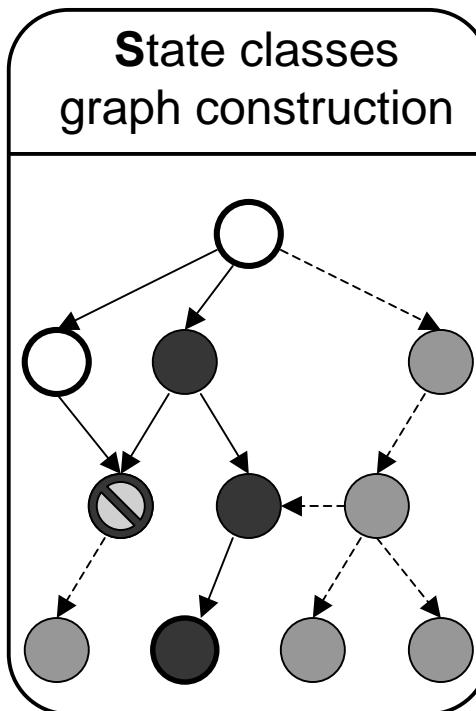
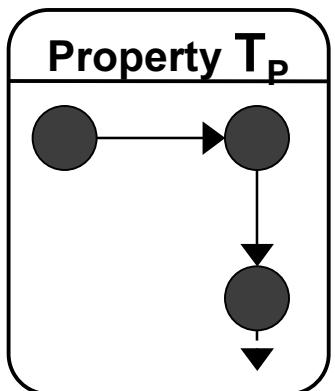
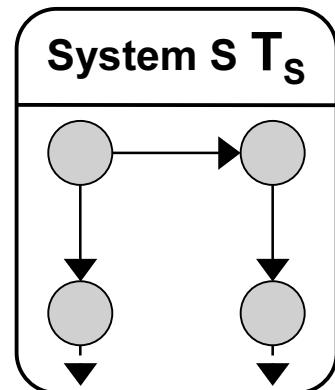
> Model of a Property T_p (TIOSM)



3 - Validation a posteriori (Integration test)

On the fly trajectories extraction

- The property is searched during the state classes graph construction



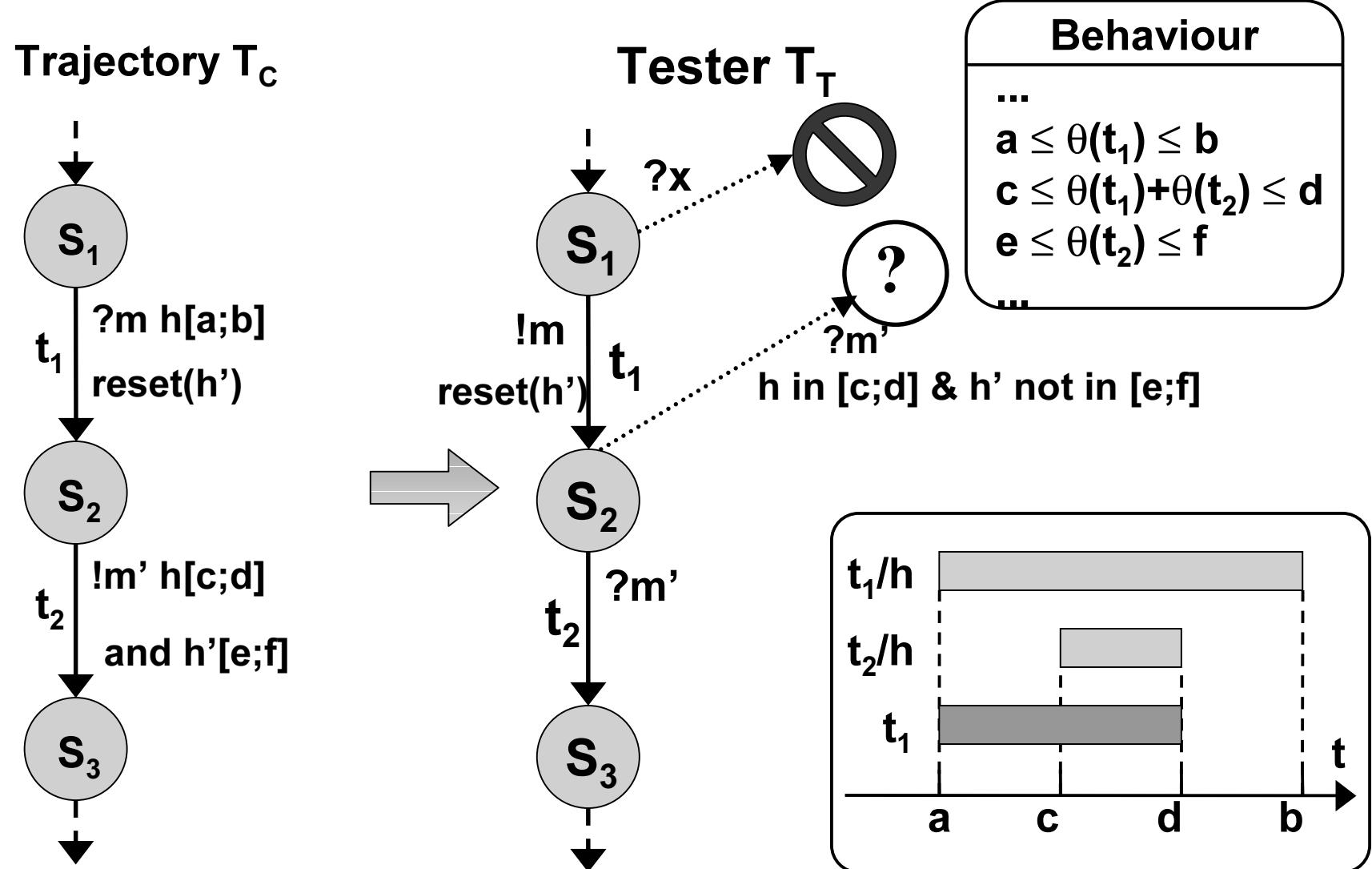
- Stop as soon as the property is founded or ...

No Trajectory

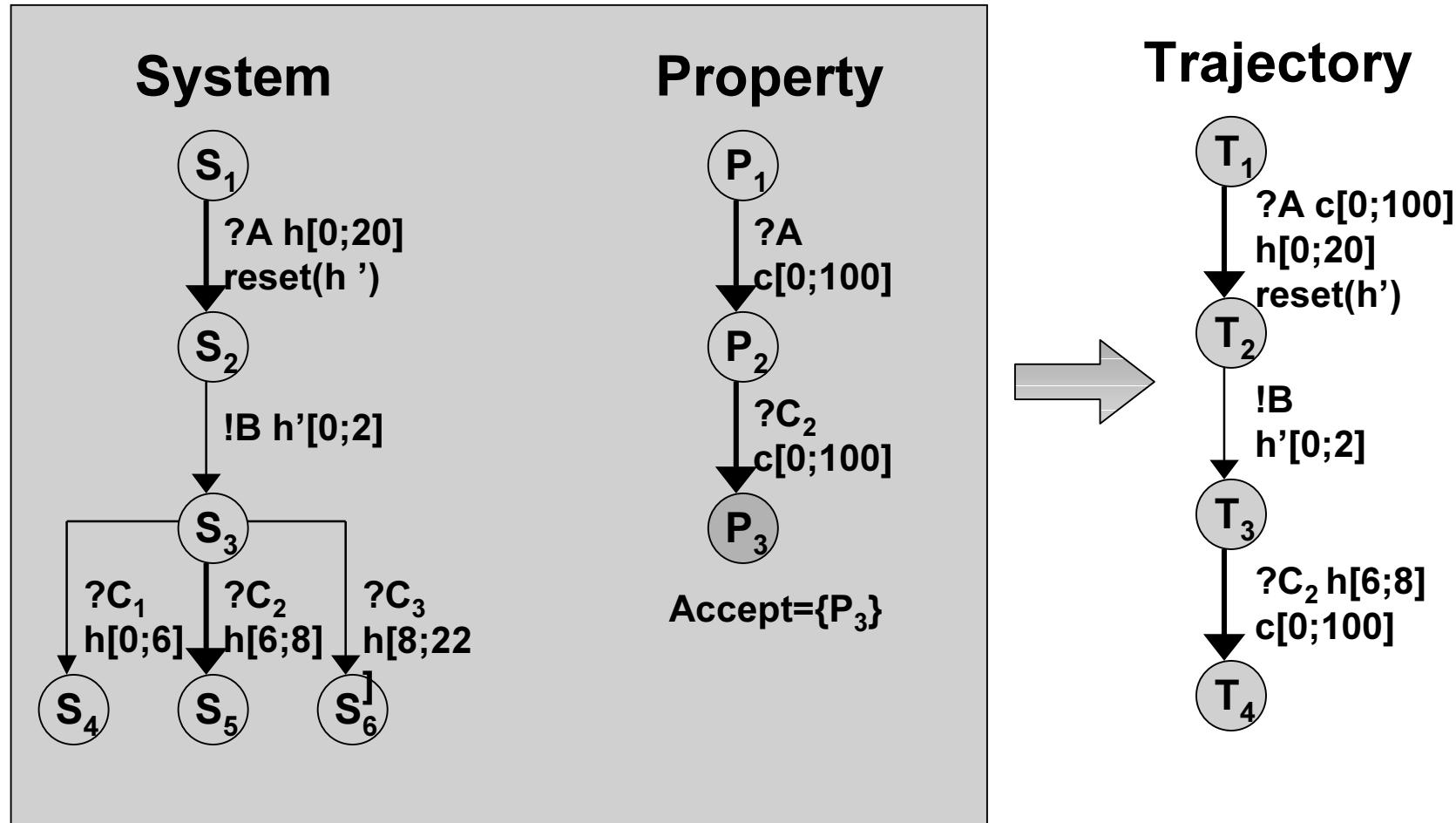
Trajectory T_c

verify P

3 - Validation a posteriori (Integration test)

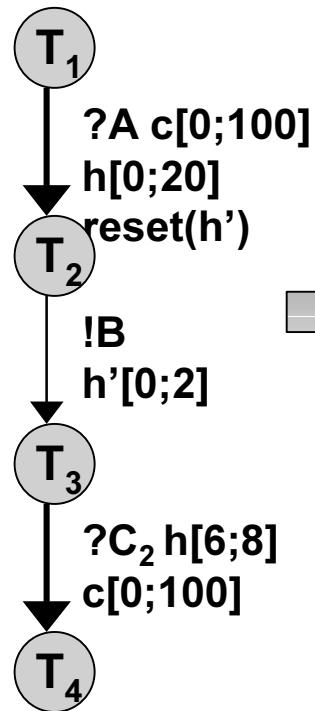


3 - Validation a posteriori (Integration test)

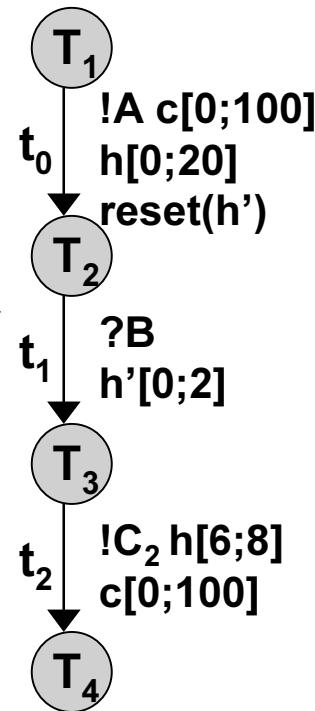


3 - Validation a posteriori (Integration test)

Trajectory



Tester



$$\begin{aligned}
 & 0 \leq \theta(t_0) \leq 100 \quad \& 0 \leq \theta(t_0) \leq 20 \quad \& \\
 & 0 \leq \theta(t_1) \leq 2 \quad \& \\
 & 0 \leq \theta(t_0) + \theta(t_1) + \theta(t_2) \leq 100 \quad \& \\
 & 6 \leq \theta(t_0) + \theta(t_1) + \theta(t_2) \leq 8
 \end{aligned}$$

Normal Tester

Fire of t_0

$[0;20]$

9

Inconclusive

Reception of B
Fire of t_1

$[0;2]$

OK

2

Inconclusive

Adaptative Tester

$[0;8]$

7

$[0;1]$

OK

$[1;2]$

Inconclusive

1

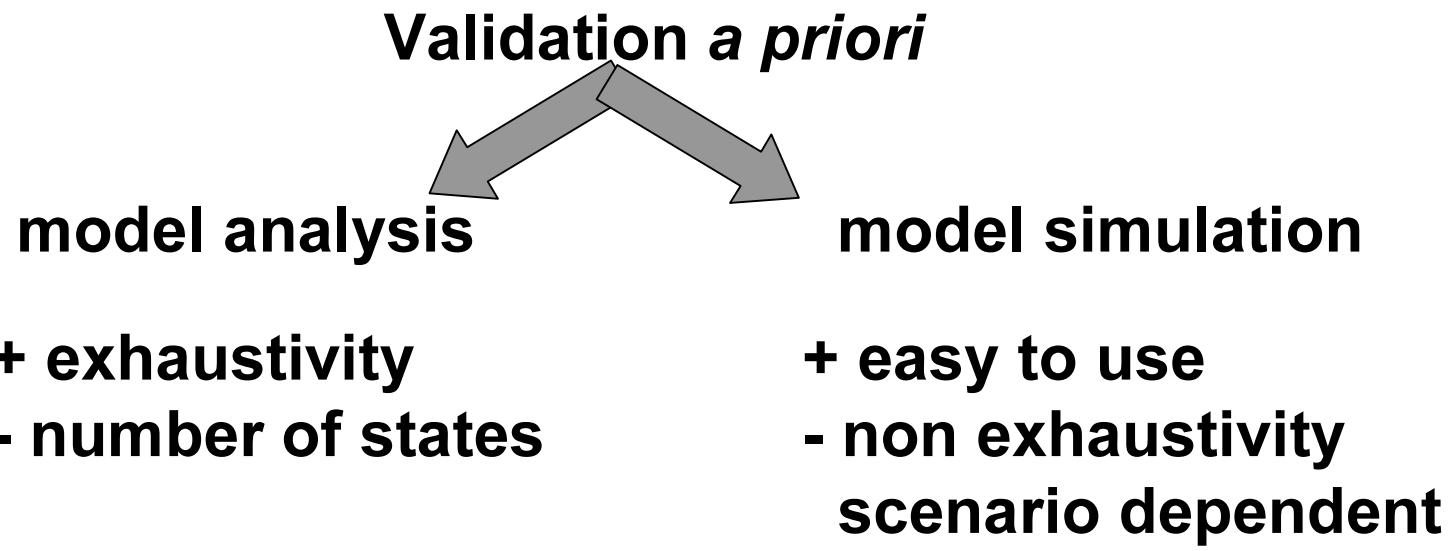
Pass

3 - Validation a posteriori (Integration test)

- Conclusions
 - Formal generation of timed test sequences
 - Reduction of *inconclusive* cases during the test process

4 - Validation a priori

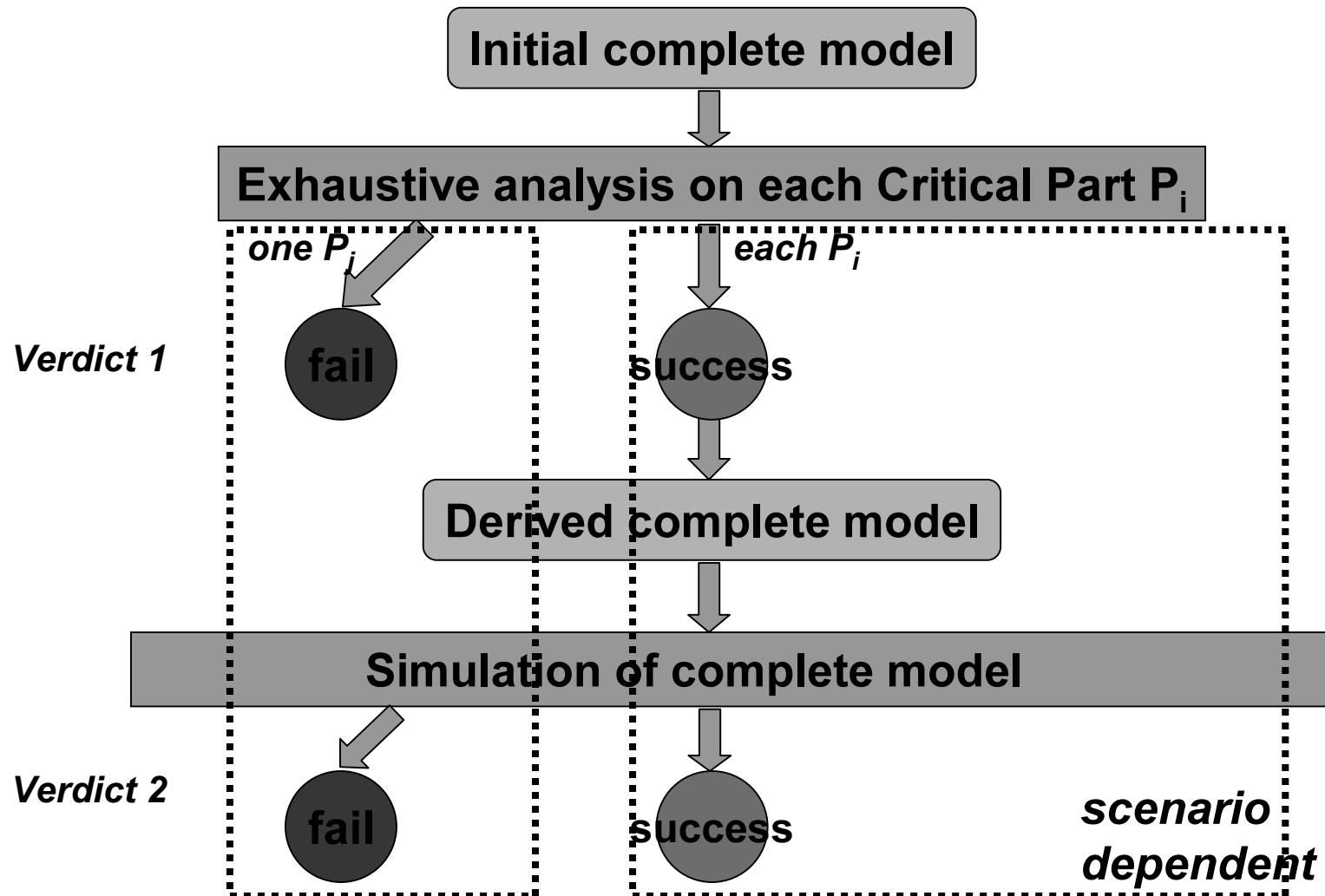
how to master the complexity



An hybrid validation method

{ exhaustive analysis on critical parts } + { simulation on the complete system }

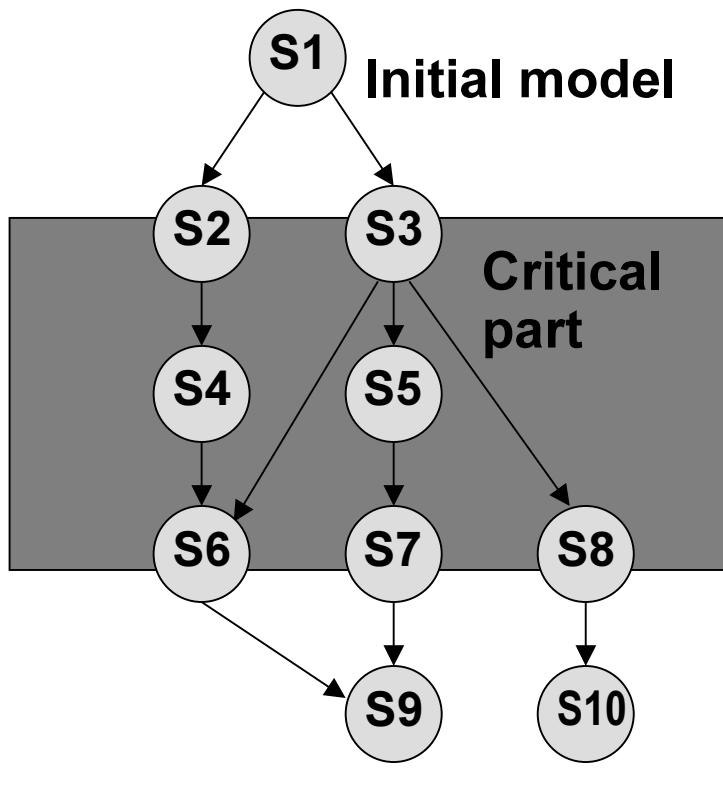
4 - Validation a priori - Hybrid Method



4 - Validation a priori - Hybrid Method

Critical parts - partial models definition

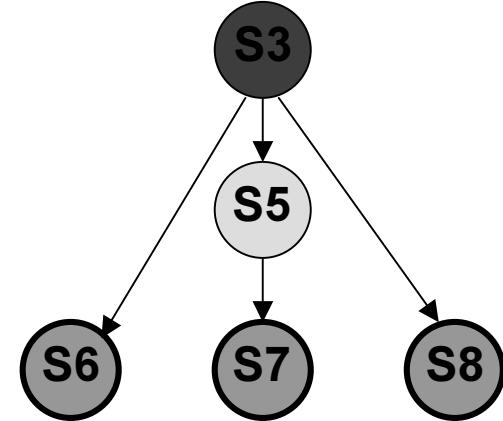
- 1 and only 1 input state
- at least one output states
- 2 partial models cannot be overlapped



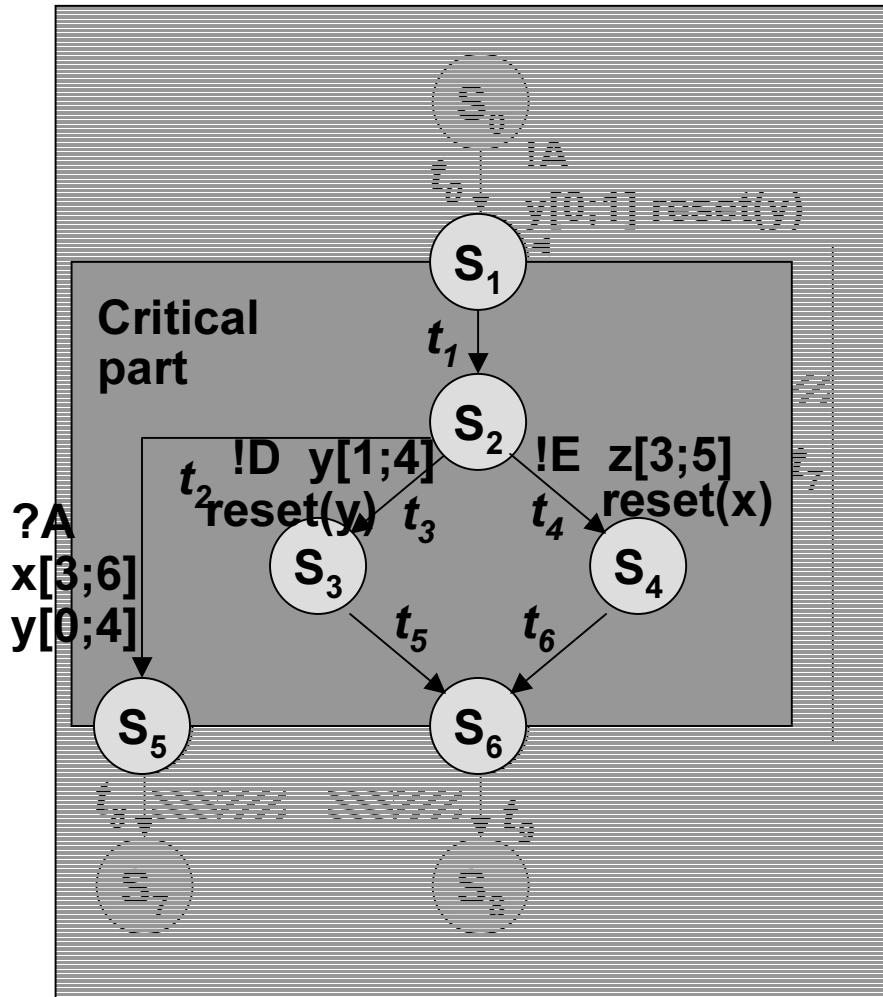
Partial model
 PM_1



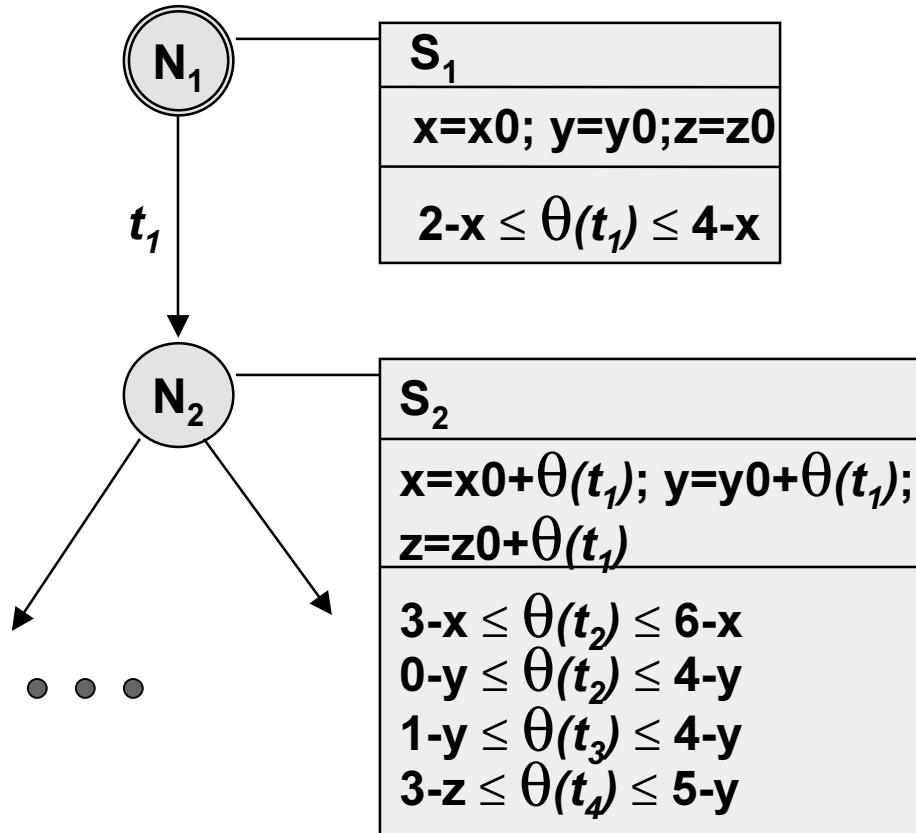
Partial model
 PM_2



4 - Validation a priori - Hybrid Method

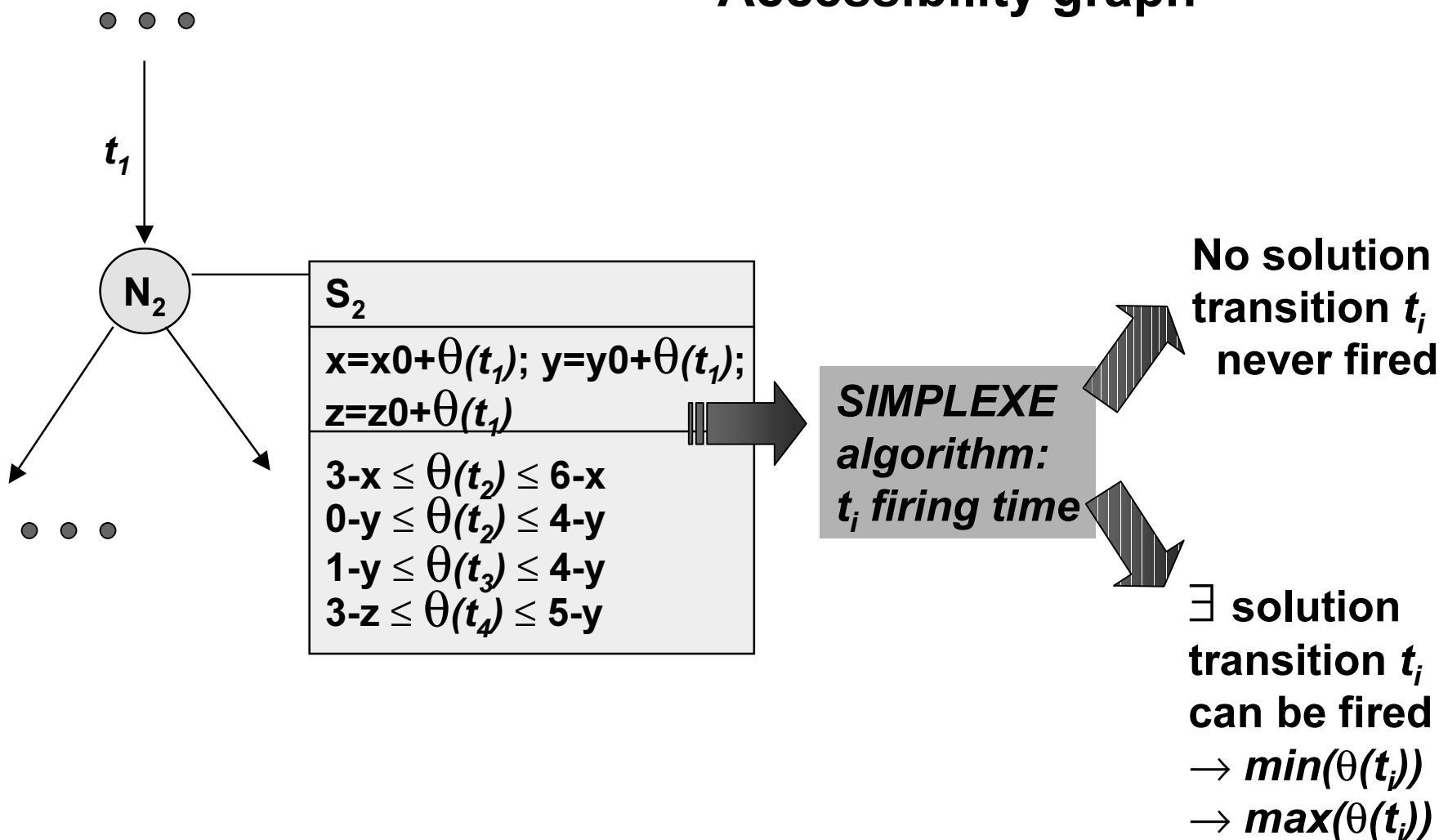


Accessibility graph
Set of clocks : $C = \{x,y,z\}$



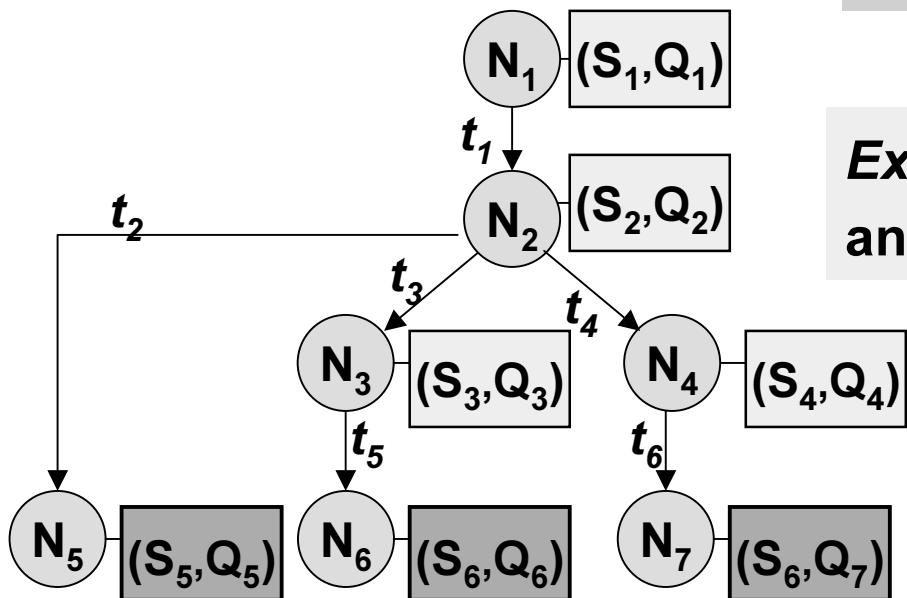
4 - Validation a priori - Hybrid Method

Accessibility graph



4 - Validation a priori - Hybrid Method

Accessibility graph



Exit states of partial model : {S5, S6}

Exit Nodes : node characterized with an exit state

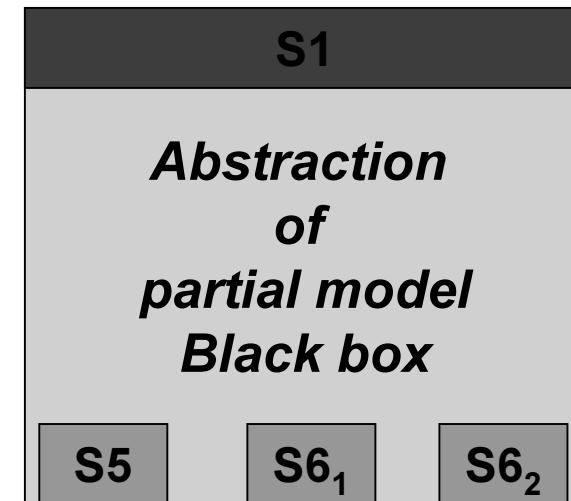
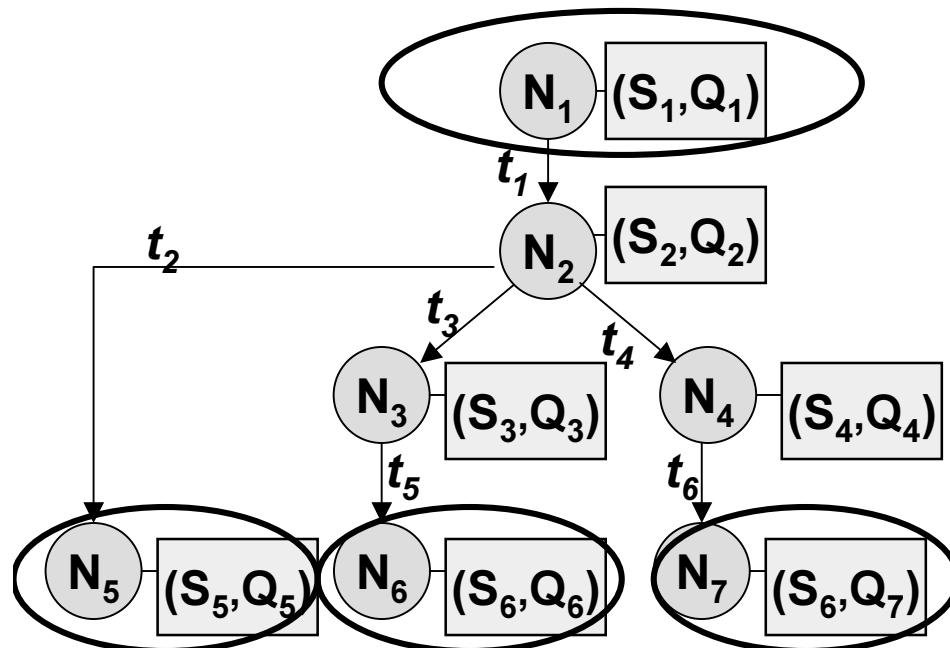
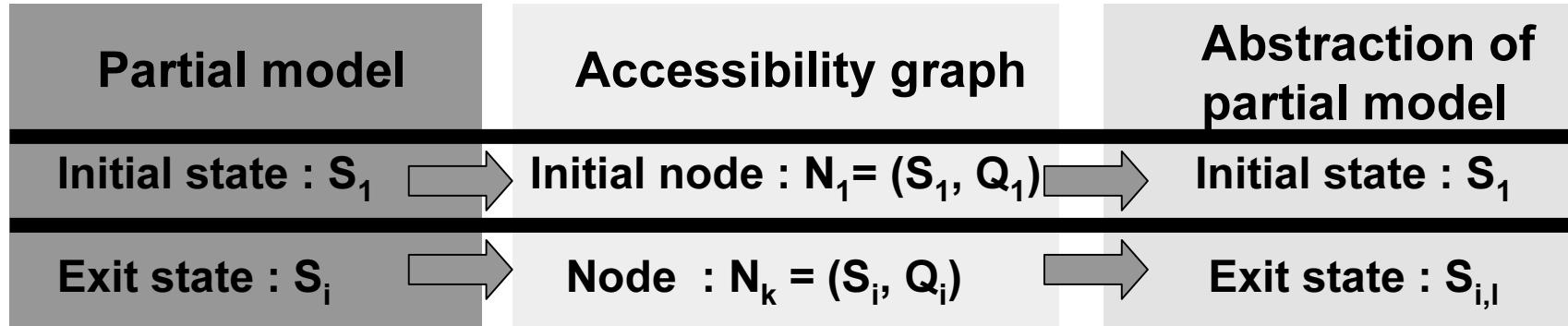
Every exit node can be reached

→ **Success**

One exit node cannot be reached

→ **Fail**

4 - Validation a priori - Hybrid Method

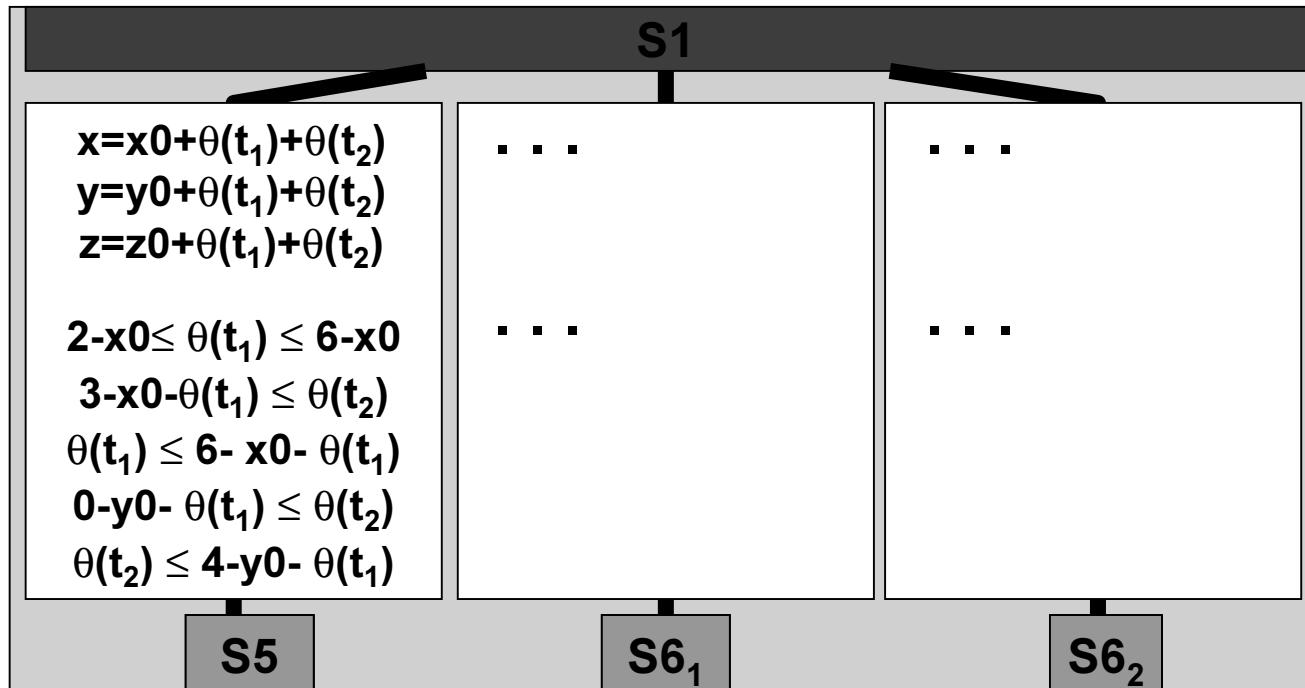


4 - Validation a priori - Hybrid Method

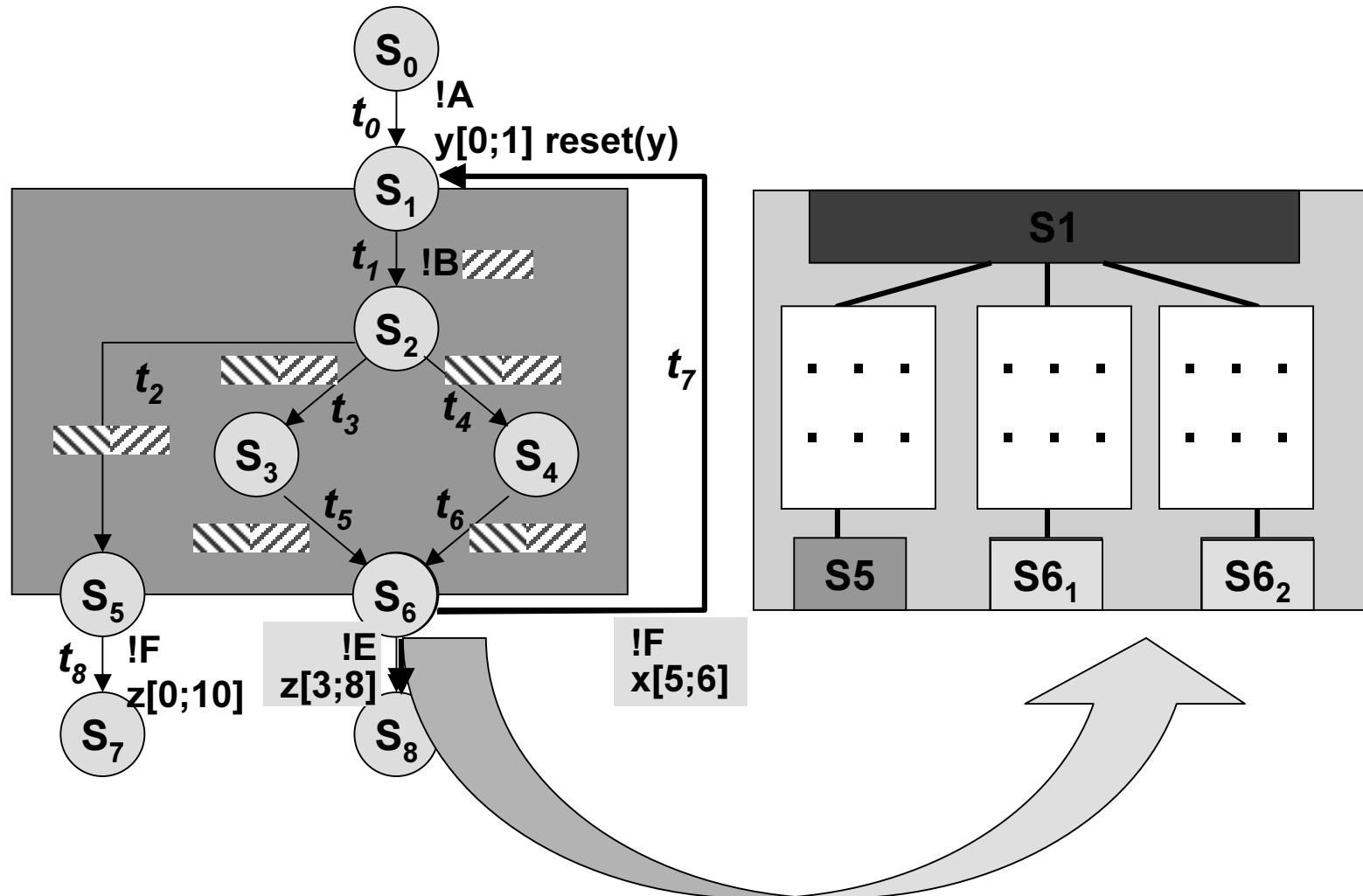
For each exit state S

Value of each clock when reaching this state

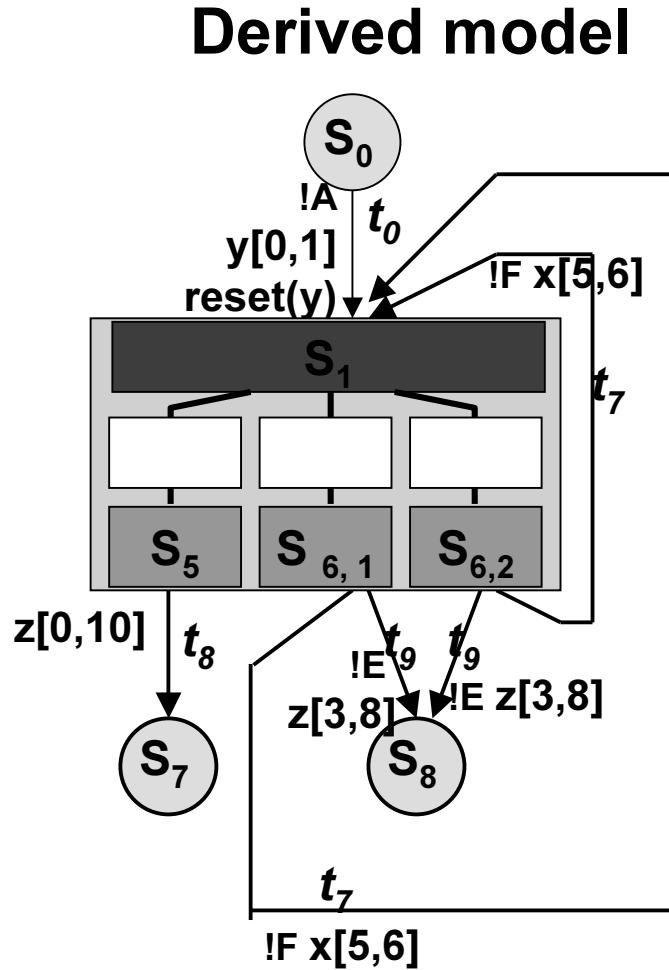
Constraints on clocks along the path from init state to S



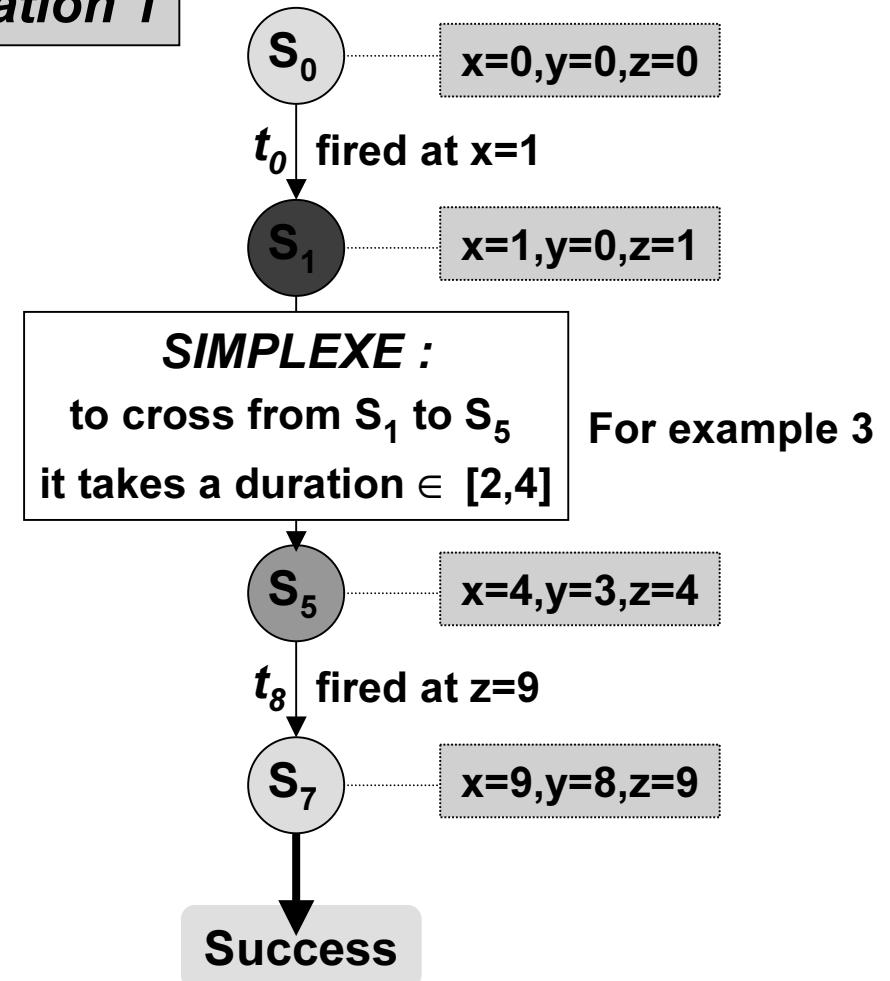
4 - Validation a priori - Hybrid Method



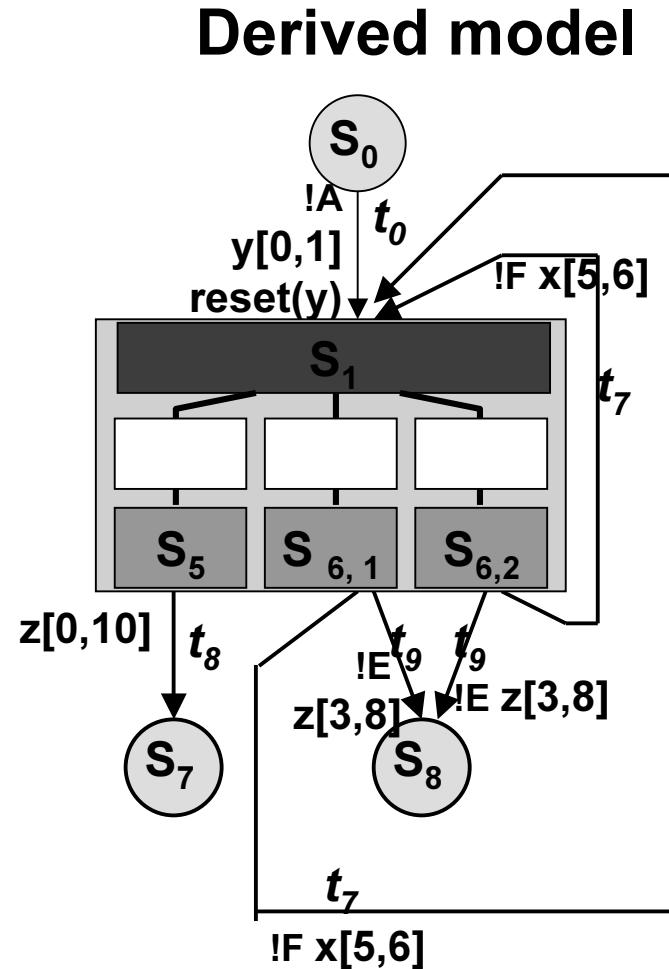
4 - Validation a priori - Hybrid Method



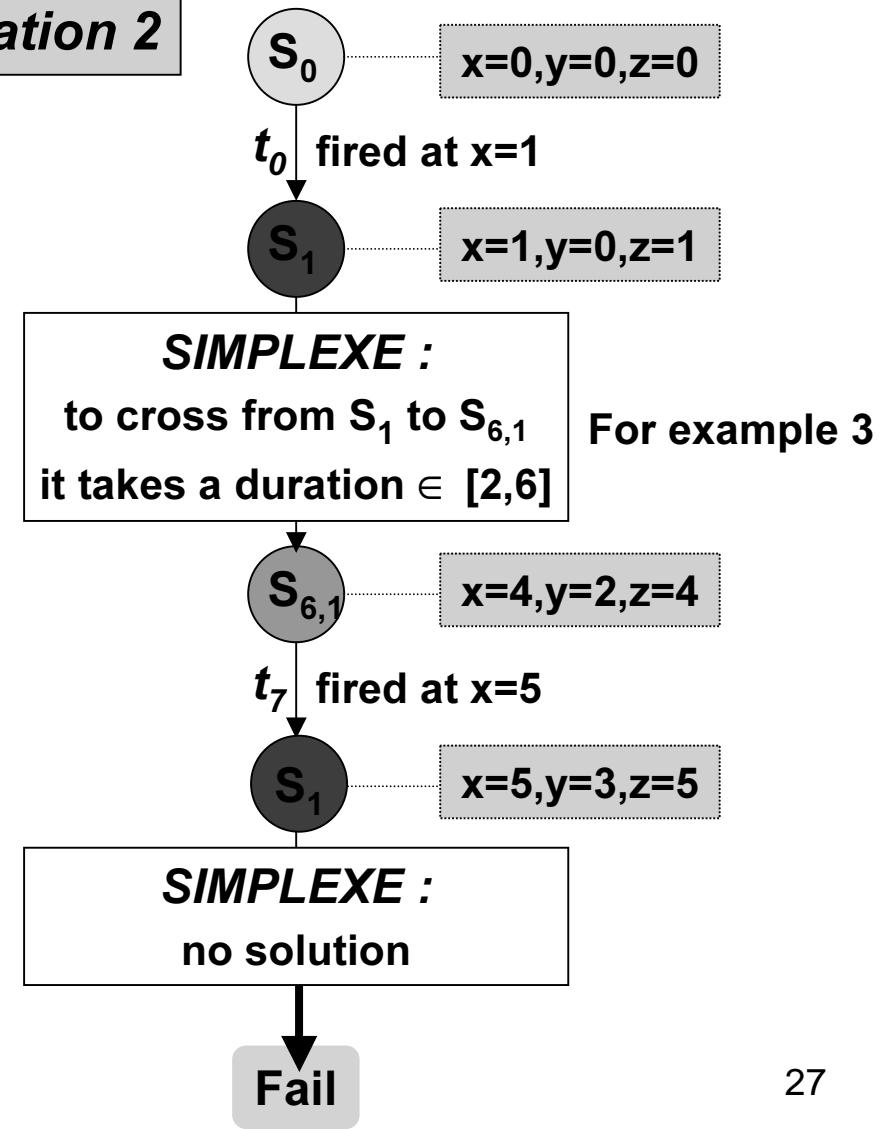
Simulation 1



4 - Validation a priori - Hybrid Method



Simulation 2

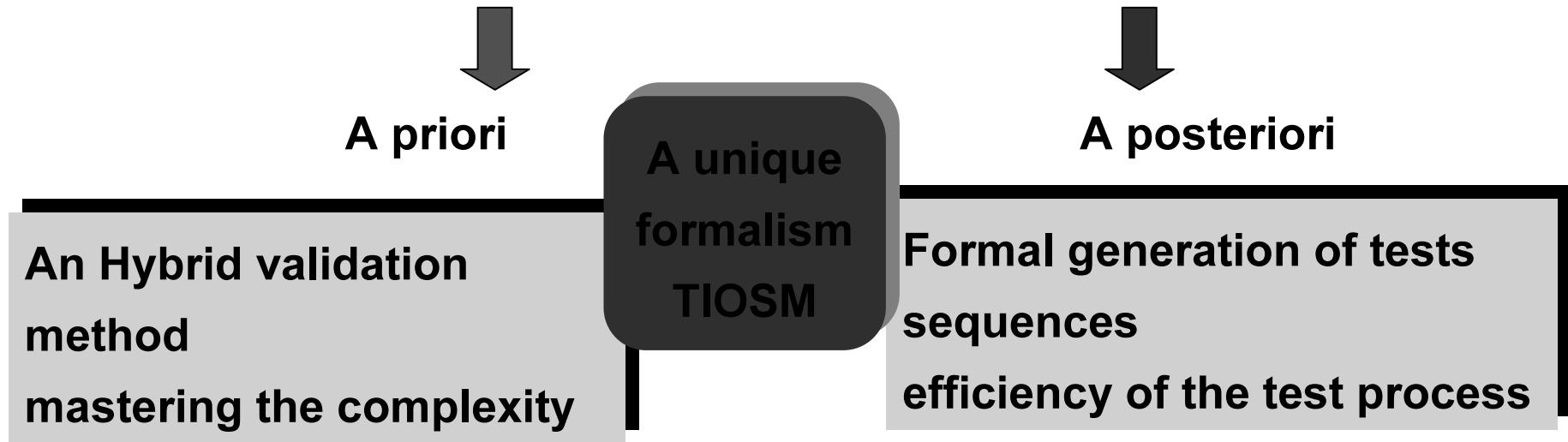


4 - Validation a priori - Hybrid Method

- Conclusions**

- A validation method decreasing the complexity**
- Simulation and Exhaustive analysis**

5 - Validation of Real Time Systems



Future trends

- ◆ *Composition / communication between critical parts*
- ◆ *COTS (Components off-the-shelf)*
- ◆ *Industrial Tester - CANoe (TTCN)*