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# Integration of UML and B Specification Techniques: Systematic Transformation from OCL Expressions into B

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## Abstract

*In the continuity of our research on integration of UML and B, we address in this paper the transformation from OCL (Object Constraint Language), which is an integral part of UML, into B. Our derivation schemes allow to automatically derive into B not only the complementary class invariants, the guard conditions in state-charts (in OCL) but also the OCL specifications for class operations.*

**Keywords:** UML, OCL, OCL operation, B expression, B generalised substitution.

## 1 Introduction

The Unified Modelling Language (UML) [19] and the B language [1] are two specification techniques well recognised in software engineering for their application capability in industry. Their integration is motivated by the hope to be able to use them jointly in a practice, unified and rigorous software development. The practicality comes from UML as the specification technique largely practised and accepted in software industry. It also comes from B as the formal technique whose industrial application are effective [9, 5]. The rigour comes from B as a formal method. The unification comes at the same time from UML and B since they are used during the whole software development from requirements expressions until the design and programming.

The transformation from UML specifications into B aims at a two-fold goal. On the one hand, one can use B powerful support tools like AtelierB [20], B-Toolkit [2] to analyse and detect inconsistencies within UML specifications (see further discussions in [12]). On the other hand, we can also use UML specifications as the starting point to develop B specifications which can then be refined automatically to an executable code [10].

Meyer and Souquière [17] and Nguyen [18], based on the previous work of Lano [11], have proposed the derivation schemes from UML structural concepts into B. Each class, attribute, association and state is modelled as a B variable. The properties of those concepts are modelled as B invariants. The inheritance relationship between classes is also modelled as B invariant predicates between B variables for the classes in question.

In [13] we have proposed approaches for modelling UML operations (operations declared in class diagrams). Each UML operation is firstly modelled by a B abstract operation in which the expected effects of such an operation on related data is specified directly on the derived data. If a UML operation is realised by an interaction or activity diagram then the B operation corresponding is refined to give rise a B implementation operation.

The UML-B derivation schemes for UML structural concepts and for UML operations are used in the derivation procedure which allow to integrate UML class and collaboration diagrams into one B specification. At this stage, only the architecture, data and the operations' signature of the derived B specification are generated automatically. For the invariant within B specification, only the part that reflects the properties of UML structural concepts expressed graphically in UML diagrams is generated. Therefore, the B specification should be completed with invariants for supplementary class invariant as well as B operations' body.

As cited in the UML literature [19], OCL (Object Constraint Language) is often used to specify supplementary class invariant as well as pre- and post-conditions of UML operations; in the continuity of our research on integration of UML and B, we address in this paper the transformation from OCL expressions into B. This OCL-B translation is applied for generating supplementary invariant and the abstract operations' body of the derived B specification.

Section 2 outlines the principles of the transformation

from OCL expressions into B. The derivation schemes for OCL types and their operations are presented in Section 3. The derivation schemes specific for postconditions are presented in Section 4. A case study is presented in Section 5. Discussions in Section 6 conclude our presentation.

## 2 From OCL expressions to B : overview

### 2.1 The OCL language

The Object Constraint Language (OCL) is now an integral part of UML [19]. One can use OCL to write constraints that contain extra information about, or restrictions to, UML diagrams. OCL is intended to be simple to read and write. Its syntax is similar to object-oriented programming languages. Most OCL expressions can be read left-to-right where the left part usually represents - in object-oriented terminology - the receiver of a message. Frequently used language features are attribute access of objects, navigation to objects that are connected via association links, and “is-Query” operation calls. OCL expressions are not only used to define invariants on classes and types, they also allow specification of guard conditions in UML state-charts and pre- and postconditions on UML operations.

### 2.2 The B language and method

B [1] is a formal software development method that covers the software process from specifications to implementations. The B notation is based on Zermelo-Frankel set theory and first order logic. Specifications are composed of abstract machines similar to modules or classes; they consist of variables, invariance properties related to those variables and operations. The state of the system, i.e. the set of variable values, is only modifiable by operations. The means by which B operations specifies state transitions is the *generalised substitution language* whose semantics is defined by means of predicate transformers [8] and the weakest precondition [7]. A generalised substitution is an abstract mathematical programming construct, built up from basic substitution  $x := e$ , corresponding to assignments to state variables, via a set of operators like No-op (*skip*), bounded choice (**choice**  $S_1$  **or**  $S_2$  ...), preconditioning (**pre**  $P$  **then**  $S$  **end**), unbounded non-determinism (**var**  $v$  **in**  $S$  **end**, **any**  $v$  **where**  $P$  **then**  $S$  **end**), guarding, sequential composition, multiple generalised composition and looping.

The abstract machine can be composed in various ways. Thus, large systems can be specified in a modular way, possibly reusing parts of other specifications. B refinement can be seen as an implementation technique but also as a specification technique to progressively augment a specification with more details until an implementation that can then be

translated into a programming language like ADA, C or C++. At every stage of the specification, proof obligations ensure that operations preserve the system invariant. A set of proof obligations that is sufficient for correctness must be discharged when a refinement is postulated between two B components.

### 2.3 Principles to translate OCL expressions into B

The core of OCL is given by an expression language. OCL expressions can be used in various contexts, for example, to define constraints such as class invariants and pre- and postconditions of UML operations. Our derivation schemes from OCL to B are therefore defined for concepts related to OCL expressions: (i) the OCL types and the associated operations and (ii) the postconditions of UML operations.

It is natural to model an OCL type by a B type, which would be a B predefined type such as  $\mathbb{Z}$ , *BOOL* etc, or a B user-defined type such as sets or relations. In addition, the formalisation in B of OCL types is guided and motivated by the wish to facilitate the formalisation in B of operations on OCL types. Intuitively, an OCL expression for class invariants, for guard conditions or for preconditions on behavioural concepts should be modelled by a B expression; meaning that every OCL operation (except *oclsNew*, which is used in postconditions of UML operations) should be represented by a B expression.

The derivation schemes from OCL to B for types and associated operations are sufficient to derive a B expression from an OCL expression of class invariants on class diagrams, guard conditions on state-charts or preconditions of UML operations. To model postconditions of UML operations, the use of B generalised substitutions is necessary. The OCL expressions involving values after the execution of a UML operation are translated into B substitutions.

### 2.4 Related work

The transformation from OCL into other formal notations have been discussed in several works [3, 4], however our choice of B as the target notation is motivated by the fact that B is a stable language with powerful support tools that have been advocated in industrial applications [9, 5].

The transformation from OCL into B has been previously discussed by Marcano and Lévy [15], in which the authors presented the derivation schemes from OCL expressions to B expressions. However there are several shortcomings in this proposal:

1. the postconditions of behavioural concepts have not been considered;

2. the fact to formalise an OCL operation by a B corresponding operation seems to be ambiguous due to the mismatch between OCL type system and B type system. For example, the OCL type `String` does not correspond to the B type `STRING`; B does not support the real type;
3. the fact to model OCL collection types by B sets is only appropriate if the collection type is a set; consequently the modelling in B of collection operations `collect`, `select` etc. is not appropriate.

Dealing with the three shortcoming above is essentially our contributions in this paper.

### 3 Modelling OCL types and their operations

The types in OCL can be classified as follows. The group of predefined basic types includes Integer, Real, Boolean and String. Enumeration types are user-defined. An object type corresponds to a classifier in an object model.

Collections of values can be described by the collection types `Set(T)`, `Sequence(T)` and `Bag(T)`. These are the classical types for bulk data, namely sets, lists and multi-sets respectively. The parameter T denotes the type of the elements. Notice that types at the meta-level such as `OclExpression` are not considered in the translation from OCL expressions into B.

#### 3.1 Predefined basic types

**Derivation 1 (Integer)** In B there are two predefined types  $\mathcal{Z}$  and `INT` which correspond to the OCL type Integer.  $\mathcal{Z}$  is chosen as the formalisation of Integer since  $\mathcal{Z}$  is more abstract than `INT`. An integer value `nn` in OCL is modelled in B as a value `nn` of  $\mathcal{Z}$ . As shown in Table 1, all Integer OCL operations but “/” can be expressed by a B expression on  $\mathcal{Z}$ .

**Remark 1 (Modelling the operation “/”)**

1. In OCL, the operation “/” between two integers `a` and `b`, gives as result a real value. Since B does not define the data type for real values, we propose to model the ratio `a/b` by the pair  $a \mapsto b$ , where  $a$  and  $b$  denote respectively the B formalisation of `a` and `b`.
2. The fact of using a ratio to express the real division “/” between two integers implies to define the formalisation in B for operations on ratios. Operations between an integer and a ratio can be converted into operations on ratios. As an example, the operation “+” between two ratio `a/b` and `c/d` can be modelled by  $a \cdot d + c \cdot b \mapsto b \cdot d$ ; details of a such formalisation can be found in [14].

Operations OCL	Semantics in B
<code>a=b</code>	$a = b$
<code>a&lt;&gt;b</code>	$\neg(a = b)$
<code>a+b</code>	$a+b$
<code>a-b</code>	$a-b$
<code>-a</code>	$-a$
<code>a*b</code>	$a \times b$
<code>a div b</code>	$a/b$
<code>a mod b</code>	$a \text{ mod } b$
<code>a&lt;b</code>	$a < b$
<code>a&lt;=b</code>	$a \leq b$
<code>a&gt;b</code>	$b < a$
<code>a&gt;=b</code>	$b \leq a$
<code>a.min(b)</code>	$\min(\{a, b\})$
<code>a.max(b)</code>	$\max(\{a, b\})$
<code>a.abs</code>	$\max(-a, a)$
<code>a/b</code>	$a \mapsto b$

Table 1. Operations on integers

**Derivation 2 (Boolean)** The OCL type Boolean is modelled in B by its correspondence `BOOL`. Table 2 shows the B formalisation of OCL Boolean operations, where `a`, `b` and `c` are three booleans and  $a$ ,  $b$  and  $c$  are their formalisation in B. An OCL boolean variable or constant `x` is modelled in B by  $x$  (variable or constante). An OCL boolean expression `exp` is modelled in B as a predicate `exp` or as a boolean expression `bool(exp)`, where `exp` is derived from `exp` according to the rules in Table 2 and `bool` is a B predefined function to convert a predicate into a boolean value.

Operations OCL	Semantics in B
<code>a=b</code>	$\text{bool}(a) = \text{bool}(b)$
<code>a&lt;&gt;b</code>	$\neg(\text{bool}(a) = \text{bool}(b))$
<code>a or b</code>	$a \vee b$
<code>a xor b</code>	$\neg(\text{bool}(a) = \text{bool}(b))$
<code>a and b</code>	$a \wedge b$
<code>not a</code>	$\neg a$
<code>a implies b</code>	$\neg a \vee b$
<code>if a then b else c endif</code>	$\neg(\text{bool}(a \wedge b) = \text{bool}(\neg a \wedge c))$

Table 2. Operations on booleans

**Derivation 3 (String)** The B predefined type `STRING` cannot be used to model the OCL type String due to restrictions of operations on `STRING` (only “=” and “<>” are defined for `STRING`). Our solution is to use a B user-defined type `seq(0..255)` to model String. The background idea is that a string can be considered as a sequence of characters and the range `0..255` models the set of possible characters. A string `str` in OCL is therefore modelled

by an element  $str$  of  $seq(0..255)$ . The OCL operations on String (except two operations `toUpper` and `toLower`) can be expressed by B expressions on  $seq(0..255)$  (cf. Table 3).

**Remark 2** Two operations `toUpper` and `toLower` involve a repetitive computation which is very sophisticated such that they cannot be expressed by an expression B at the abstract machine level.

Operations OCL	Semantics in B
<code>a=b</code>	$a = b$
<code>a&lt;&gt;b</code>	$\neg(a = b)$
<code>a.size</code>	$size(a)$
<code>a.concat(b)</code>	$a \hat{\ } b$
<code>a.substring(lower,upper)</code>	$(a \uparrow upper) \downarrow lower$
<code>a.toUpper</code>	<i>no definition</i>
<code>a.toLower</code>	<i>no definition</i>

**Table 3. Operations on strings**

**Derivation 4 (Real)** There is no B predefined type for real values, a definitive solution for modelling the OCL type Real in B has not been yet achieved. In waiting for a such solution, a temporary solution, inspired from Remark 1, can be used. It is to approximate a real value by a ratio. Hence the type Real can be modelled in B by relation  $\mathcal{Z} \leftrightarrow \mathcal{Z}$ . The OCL operations on Real are modelled in B in a similar manner of the operations on ratios.

**Derivation 5 (Enumeration types)** Each enumeration type  $Enum = \{val_1, \dots, val_n\}$  is modelled in B by an enumerated set serving as a user-defined type  $Enum = \{val_1, \dots, val_n\}$ . Each element  $val_i$  in Enum is modelled by an element  $val_i$  in Enum. The modelling in B of operations on an enumeration type is shown in Table 4, where  $a\#, b\#$  are two values of type Enum and  $a, b$  denote their B formalisation.

Operations OCL	Semantics in B
<code>a#=b#</code>	$a = b$
<code>a#&lt;&gt;b#</code>	$\neg(a = b)$

**Table 4. Operations on enumerations**

**Derivation 6 (Object types)** According to Meyer and Souquières [17], for each class Class, the B constant *CLASS* models the possible instance set and the B variable *class*, which is defined as a subset of *CLASS*, models the effective instance set of Class. Therefore, the object type Class is modelled in B as *CLASS*, whereas the operation `Class.allInstances` is modelled as *class*. An

object *cc* of Class is modelled in B by an element *cc* of *class*.

### 3.2 Collection types

**Derivation 7 (Collection types)** Given T an OCL type and T its B formalisation:

1. the collection type `Set(T)`, which denotes all subsets of T, is modelled in B by  $\mathcal{P}(T)$ ;
2. the collection type `Bag(T)`, which denotes all multi-sets on T, is modelled in B by  $T \rightarrow \mathcal{N}$ . An element *bb* of `Bag(T)` can be modelled as  $bb \in T \rightarrow \mathcal{N}$  and for each element *tt* of *bb*,  $bb(tt)$  denotes the occurrence number of *tt* in *bb*;
3. the collection type `Sequence(T)`, which denotes all lists of elements of T, is directly modelled in B by  $seq(T)$ .

Using Derivation 7, almost predefined operations on collection types (except `asSequence` on sets or bags and `excluding` on sequences<sup>1</sup>) can be expressed by a B expression. The semantics of OCL operations `select`, `reject`, `collect`, `forAll`, exists on collection types can also be interpreted by B expressions. Details of those formalisations can be found in [14], Derivation 8 shows some examples.

**Derivation 8 (OCL operations on collection types)**

Given T an OCL type on which the collection types are defined and  $ss : Set(T)$ ;  $bb : Bag(T)$ ;  $se, se_2 : Sequence(T)$ ;  $tt : T$ ; *boolexprrt* is a boolean expression on *tt* and *exprrt* is an expression on *tt*. Let's call *T*, *ss*, *bb*, *se*, *se<sub>2</sub>*, *tt*, *boolexprrt* and *exprrt* respectively the B formalisation of T *ss*, *bb*, *se*, *se<sub>2</sub>*, *tt*, *boolexprrt* and *exprrt*:

1. the OCL expression `bb->includes(tt)`, which checks whether *tt* is an element of the bag *bb*, is modelled in B by  $tt \in dom(bb)$ ;
2. the OCL sequence union expression `se->union(se2)` is modelled in B by the sequence concatenation expression  $se \hat{\ } se_2$ ;
3. the OCL expression `ss->select(tt|boolexprrt)`, which extracts elements *tt* of *ss* satisfying *boolexprrt*, is modelled in B by  $\{tt | tt \in ss \wedge boolexprrt\}$ ;
4. the OCL expression `ss->collect(tt|exprrt)`, which derives a collection (a bag) of values *exprrt* computed on every element *tt* of the set *ss*, is modelled in B by  $\{tt, nn | tt \in exprrt[ss] \wedge nn \in \mathcal{N} \wedge nn = card(\{xx | xx \in ss \wedge exprrt(xx) = tt\})\}$ .

<sup>1</sup>The semantics of the operation `asSequence` on sets or bags has not been defined in OCL therefore we cannot model it in B. It is the same for the operation `excluding` on a sequence.

### 3.3 Property access operations

OCL expressions can refer to attributes, association ends and “is-Query” operations thanks to property access operations. A property access operation on an object might return a single value/object, a set of values/objects, a multi-set of values/objects or a sequence of values/objects. It is also possible to apply a property access operation on the result of another property access operation. Hence the target of a property access operation might be an object, a set of objects, a multi-set of objects or even a sequence of objects. Our derivation schemes for property access operations are based on the derivation schemes for UML structural concepts and the derivation schemes for collection types (cf. Derivation 7). The derivation schemes for UML structural concepts are detailed in [17, 18, 16]; in Derivation 9, we recall only essential points which facilitate the presentation afterwards.

#### Derivation 9 (Structural concepts)

1. An attribute *attr* of type *typeAttr* in a class *Class* is modelled by a B variable *attr* defined as  $attr \in class \leftrightarrow typeAttr$ , where the variable *class* models the effective instance set of *Class* and *typeAttr* is defined as a B set to model *typeAttr*. The relation defining *attr* might be further refined according to additional properties of *attr*.
2. A binary association *assos* between two classes *Class* and *Class*<sub>2</sub> is modelled by a B variable *assos* defined as  $assos \in class \leftrightarrow class_2$ . If there are eventual qualifiers  $q_1 : Q_1, \dots, q_n : Q_n$  at the role end of *Class*, they are modelled in a similar manner to an attribute:  $q_1 \in class_2 \leftrightarrow Q_1, \dots, q_n \in class_2 \leftrightarrow Q_n$ . We add also a B invariant linking *assos* and  $q_1, \dots, q_n$  as follows:  $(assos^{-1} \otimes q_1 \otimes \dots \otimes q_n)^{-1} \in class \times Q_1 \times \dots \times Q_n \leftrightarrow class_2$ <sup>2</sup>. As for attributes, the relation defining *assos* could be further refined according to additional properties of *assos*.

**Derivation 10 (Attribute operations)** Given *attr*, *cc*, *sc*, *bc*, *seqc* an attribute, an object, a set of objects, a bag of objects and a sequence of objects of a class *Class*. Let’s call *attr*, *cc*, *sc*, *bc* and *seqc* their B formalisation according to Derivation 6, Derivation 7 and Derivation 9:

1. the expression *cc.attr*, which denotes the value(s) of the attribute *attr* associated to the object *cc*, is generally modelled in B by  $attr[\{cc\}]$ . If the cardinality of *attr* is equal to 1, *cc.attr* can be modelled by  $attr(cc)$ ; otherwise and if *attr* is ordered,  $attr[\{cc\}]$  is interpreted as a sequence;

<sup>2</sup>  $\otimes$  is the direct product between two relations.

2. the expression *sc.attr* denotes a collection of values for *attr* associated with elements in *sc*. If the cardinality of *attr* is equal to 1 then *sc.attr* denotes a set and is modelled by  $attr[sc]$ . If the cardinality of *attr* is multiple but *attr* is not ordered then *sc.attr* denotes a bag and is modelled by  $\{vv, nn | vv \in attr[sc] \wedge nn \in \mathcal{N} \wedge nn = card(attr^{-1}[\{vv\}] \cap sc)\}$ . Otherwise there is no semantics for *sc.attr* and there is no therefore corresponding B formalisation;
3. the expression *bc.attr* has no semantics if *attr* is multiple and ordered; otherwise *bc.attr* denotes a bag of values of *attr* associated to the bag *bc* and is modelled by  $\{vv, nn | vv \in attr[dom(bc)] \wedge nn \in \mathcal{N} \wedge nn = \Sigma(cc).(cc \in attr^{-1}[\{vv\}] \cap dom(bc) | bc(cc))\}$ ;
4. the expression *seqc.attr* has only semantics if the cardinality of *attr* is equal to 1 and in that case it denotes a sequence of values for *attr* and is modelled by  $\lambda(ii).(ii \in dom(seqc) | attr(seqc(ii)))$ .

**Remark 3** The navigation operations without qualifiers are modelled in a similar manner to the attribute operations.

#### Derivation 11 (Navigation operations with qualifiers)

Given a binary association *assos* from the class *Class* to *Class*<sub>2</sub>. The association *assos* is qualified by attributes  $q_1 : Q_1, \dots, q_n : Q_n$  at the role end of *Class*. Given values and objects  $cc : Class, v_1 : Q_1, \dots, v_n : Q_n$ . Let’s call *roleClass*<sub>2</sub> the role end of *Class*<sub>2</sub> in *assos*. The expression  $cc.roleClass_2[v_1, \dots, v_n]$  is modelled in B by  $(assos^{-1} \otimes q_1 \otimes \dots \otimes q_n)^{-1}[\{cc \mapsto v_1 \mapsto \dots \mapsto v_n\}]$ , where *cc*,  $v_1, \dots, v_n$  are the B formalisation of *cc*,  $v_1, \dots, v_n$  and  $q_1, \dots, q_n$  are defined according to Derivation 9. Furthermore, if the multiplicity property of *Class*<sub>2</sub> in *assos* is equal to 1,  $cc.roleClass_2[v_1, \dots, v_n]$  can be expressed in B by  $(assos^{-1} \otimes q_1 \otimes \dots \otimes q_n)^{-1}(cc \mapsto v_1 \mapsto \dots \mapsto v_n)$ .

**Remark 4** It is always possible to define the B semantics for navigation operations with qualifiers on a set or a bag of objects. However those situations are rarely encountered and the corresponding derivation schemes are omitted here.

#### Derivation 12 (Navigation to association classes)

Given *assos* a binary association class between two classes *Class* and *Class*<sub>2</sub>. Given *cc*, *sc*, *bc*, *seqc* an object, a set of objects, a bag of objects and a sequence of objects of *Class*. Let’s call *assos*, *cc*, *sc*, *bc* and *seqc* the B formalisation of *assos*, *cc*, *sc*, *bc* and *seqc* according to Derivation 9:

1. the expression *cc.assos*, which denotes the instance(s) of *assos* associated to *cc*, is modelled by  $\{cc\} \triangleleft assos$ ; if the cardinality of *Class*<sub>2</sub> in *assos* is equal to 1, *cc.assos* can also be modelled by  $cc \mapsto assos(cc)$ ;

2. the expression  $sc.assos$ , which denotes the instances of  $assos$  associated to the elements of  $sc$ , is modelled by  $sc \triangleleft assos$ ;
3. the expression  $bc.assos$ , which denotes a bag of instances of  $assos$  associated to the elements of  $bc$ , is modelled by  $\{cc, nn \mid cc \in dom(bc) \triangleleft assos \wedge nn \in \mathcal{N} \wedge nn = bc(dom(\{cc\}))\}$ ;
4. the expression  $seqc.assos$  has only semantics if the cardinality of  $Class_2$  in  $assos$  is equal to 1 and in that case it denotes a sequence of instances of  $assos$  associated to the elements of  $seqc$  and is modelled by  $\lambda(ii).(ii \in dom(seqc) \mid seqc(ii) \mapsto assos(seqc(ii)))$ .

**Derivation 13 (“is-Query” operations)** The call to “is-Query” class operations can appear in OCL expressions, however, we can not call B operations in B expressions. Therefore, we propose to create for each “is-Query” operation  $Class::oper(p_1:P_1, \dots, p_n:P_n):P$ , which appears in OCL expressions, a B variable  $oper$  defined by  $oper \in class \otimes P_1 \otimes \dots \otimes P_n \rightarrow P$ . A call  $oper(v_1, \dots, v_n)$  to an object  $cc$  of Class in OCL expressions is therefore modelled by  $oper(cc \mapsto v_1 \mapsto \dots \mapsto v_n)$  in the corresponding B expressions. Note equally that the call  $oper$  to a set, a bag or a sequence of objects of Class is rarely encountered, hence the corresponding derivation schemes are omitted here.

**Remark 5 (Let expressions)** All the variables declared by let expressions should be replaced by their values before the transformation.

## 4 Modelling postconditions

This section presents the modelling of OCL expressions on postconditions of an UML operation. As said earlier (cf. Section 2.3), the OCL postconditions of an UML operation  $oper$  are modelled in B by generalised substitutions in the body of the B abstract operation  $oper$  for  $oper$ . First of all are some definitions.

### 4.1 Definitions

Given an operation  $oper$ , the postconditions of  $oper$  can be considered as a constraint  $P(out_1, \dots, out_n, in_1, \dots, in_m)$  which links the potential “outputs” (cf. Definition 2) and potential “inputs” (cf. Definition 1) of  $oper$ .

**Definition 1 (Operation potential inputs)** The set  $Input = \{in_1, \dots, in_m\}$  of potential inputs of an operation  $oper$  consists of: (i) the possible parameters stereotyped by “in” or “inout” whose value is provided upon every call to  $oper$  and (ii) the objects, the attributes and the associations available upon the operation call.

**Definition 2 (Operation potential output)** The set  $Output = \{out_1, \dots, out_n\}$  of potential outputs of an operation  $oper$  consists of: (i) the possible return parameter, which is referenced by the name  $result$  in OCL, of  $oper$ ; (ii) the possible parameters stereotyped by “out” or “inout” of  $oper$ ; (iii) the possible newly created objects during the execution of  $oper$  and (iv) the possible updated attributes and associations.

Definition 3 presents a standard style of the constraint  $P(out_1, \dots, out_n, in_1, \dots, in_m)$ . In our opinion, the definition is enough generalised to be able to cover almost class operations. Our derivation schemes in the sequel are defined in reference to this definition.

### Definition 3 (Well-formed postconditions)

1. Every potential output  $out_i$  is defined by an elementary constraint  $P_i(out_i, [Input], [NewObject])$ , which defines  $out_i$  according to elements of  $Input$  as well as the newly created objects (the elements of  $NewObject$ , which is a subset of  $Output$ ):
  - (a)  $P_i$  states the creation of an object ( $out_i$ ) by  $oper$ ;
  - (b)  $P_i$  is a comparison between the value of  $out_i$  and the values of elements in  $Input \cup NewObject$ . Two cases should be distinguished: (i)  $P_i$  is represented by  $out_i = \text{expr}(Input \cup NewObject)$ , meaning that  $out_i$  is defined deterministically in terms of elements of  $Input \cup NewObject$ ; (ii)  $P_i$  is represented by a boolean expression but not an equality on  $out_i$  and possible elements of  $Input \cup NewObject$ , meaning that  $out_i$  is defined non deterministically in terms of elements of  $Input \cup NewObject$ ;
  - (c)  $P_i$  might represent the application of the operation  $forall$  on a set of objects/values to be updated by  $oper$ .
2. The constraint  $P$  is a combination between the elementary constraints  $P_1, \dots, P_n$  and the operations  $and$  and  $if \dots then \dots else \dots endif$ :
  - (a) the condition part in an expression  $if \dots then \dots else \dots endif$  refers to potential inputs;
  - (b) the elementary constraints  $P_1, \dots, P_n$  are linked by  $and$  in order to compose the body of expressions  $if \dots then \dots else \dots endif$ ;
  - (c) the expressions  $if \dots then \dots else \dots endif$  can be nested;
  - (d) two body expressions in an expression  $if \dots then \dots else \dots endif$  contain either another expression  $if \dots then \dots else \dots endif$  or an expression  $and$  on elementary constraints;

## 4.2 Elementary constraints

### Derivation 14 (Return parameter)

1. Given  $P_i$  in the form  $\text{result}=\text{expr}(\text{Input}[\cup\text{NewObject}])$  to define deterministically the return value of  $\text{oper}$ . We add in the B operation  $\text{oper}$  the following substitution:  $\text{out} = \text{expr}(\text{Input}[\cup\text{NewObject}])$ , where  $\text{out}$  represents the return parameter of  $\text{oper}$  and  $\text{expr}$  is the B formalisation of  $\text{expr}$ .
2. Given  $P_i$  in the form  $\text{expr}(\text{result},\text{Input}[\cup\text{NewObject}])$  to define non-deterministically the return parameter of  $\text{oper}$ . We can rewrite the constraint  $P_i$  in the following manner:
  - we introduce a temporary variable  $\text{res}$  which takes the place of  $\text{result}$  in the old  $P_i$ ;
  - we rewrite  $P_i$  as:  $\text{expr}(\text{res},\text{Input}[\cup\text{NewObject}])$  and  $\text{result} = \text{res}$ ,

the new form of  $P_i$  enables us to update  $\text{oper}$  in the following manner:

- we create a clause **any...where...** if it has not been created (cf. Derivation 18);
- we declare  $\text{res}$ , which models  $\text{res}$ :  
**any..., res where**  
 $\dots \text{expr}(\text{res}, \text{Input}[\cup\text{NewObject}])$
- we add the substitution  $\text{out} := \text{res}$  in the body of **any...where...**

### Remark 6

1. The B formalisation of  $\text{expr}$  is done using derivation schemes in Section 3; all the possible occurrences of  $\text{@pre}$  are omitted without losing the semantics of  $\text{@pre}$ . To justify this point, let's take an exemple: in the B assignment statement  $x := a$ , the value of  $a$  is always referred to be the value before execution of the corresponding operation.
2. Derivation 14 can be extended to apply for possible parameters stereotyped by  $\text{out}$  or  $\text{inout}$  of  $\text{oper}$ .

**Derivation 15 (Object creation)** Given  $P_i$  a constraint specifying that an object  $\text{cc}$  of a class  $\text{Class}$  is created by  $\text{oper}$ . We create in  $\text{oper}$ :

- a clause **any...where...** if this clause has not been created (cf. Derivation 18);
- a temporary variable  $\text{cc}$ , which models  $\text{cc}$ :  
**any ..., cc where**  
 $\dots \wedge \text{cc} \in \text{CLASS} - \text{class}$

- a B substitution, which models the updating of a set of effective instances of  $\text{Class}$  by the object  $\text{cc}$ , in the body **any...where...**  
 $\text{class} := \text{class} \cup \{\text{cc}\}$

### Derivation 16 (Updating attribute value of an object)

Given an attribute  $\text{attr}$  and an object  $\text{cc}$  of a class  $\text{Class}$ :

1. given  $P_i$  in the form  $\text{cc.attr}=\text{expr}(\text{Input}[\cup\text{NewObject}])$  to define deterministically the new value of  $\text{cc.attr}$  in which the cardinality of  $\text{attr}$  is equal to 1. We model the constraint  $P_i$  by the substitution B  
 $\text{attr} := \text{attr} \triangleleft \{\text{cc} \rightarrow \text{expr}\}$ ;
2. given  $P_i$  in the form  $\text{cc.attr}=\text{expr}(\text{Input}[\cup\text{NewObject}])$  to define deterministically  $\text{cc.attr}$  in which the cardinality of  $\text{attr}$  is multiple. We model  $P_i$  by the substitution:  
 $\text{attr} := \text{attr} \cup \{\text{cc}\} \times \text{expr}$ .

In the two cases above,  $\text{attr}$ ,  $\text{cc}$  and  $\text{expr}$  are the B formalisation of  $\text{attr}$ ,  $\text{cc}$  and  $\text{expr}$ .

**Remark 7** Derivation 16 can be extended for updating associations.

**Derivation 17 (forAll operation on an object set)** Given a set  $\text{sc}$  of objects and an attribute  $\text{attr}$  of the cardinality 1 of a class  $\text{Class}$ :

1. given  $P_i$  in the form  $\text{sc} \rightarrow \text{forAll}(p | p.\text{attr} = \text{expr}(\text{Input}[\cup\text{NewObject}]))$  to define deterministically the value of attribute  $\text{attr}$  of all elements in  $\text{sc}$ . We model  $P_i$  by the following substitution B:  $\text{attr} := \text{attr} \triangleleft \text{sc} \times \{\text{expr}\}$ ,
2. given  $P_i$  in the form  $\text{sc} \rightarrow \text{forAll}(p | \text{expr}(p.\text{attr}, \text{Input}[\cup\text{NewObject}]))$  to define non deterministically the value of attribute  $\text{attr}$  of all elements in  $\text{sc}$ . We model  $P_i$  in creating in  $\text{oper}$ :
  - a clause **any...where...** if this clause has not been created (cf. Derivation 18);
  - a temporary variable  $\text{at}$  defined as:  
**any..., at where**  
 $\text{at} \in \text{class} \rightarrow \text{type Attr} \wedge$   
 $\text{dom}(\text{at}) = \text{sc} \wedge$   
 $\forall (tt). (tt \in \text{ran}(\text{at}) \Rightarrow$   
 $\text{expr}(\text{at}, \text{Input}[\cup\text{NewObject}]))$
  - the substitution:  
 $\text{attr} := \text{attr} \triangleleft \text{at} \dots ||$   
in the body of the clause **any...where...**

In the two cases above,  $\text{attr}$  and  $\text{sc}$  are the B formalisation of  $\text{attr}$  and  $\text{sc}$ ,  $\text{expr}$  is the B formalisation of  $\text{expr}$  in which  $\text{tt}$  replaces  $p.\text{attr}$ .



### Remark 8

1. Derivation 17 did not consider the case where several attributes of the same set of objects have changed their values. However this situation can be solved by applying Derivation 17 several times.
2. We did not consider the case where `forall` is applied on a sequence or a bag of objects since those situations have no semantics in the context of postconditions.
3. Derivation 17 can be extended for the case where the body of `forall` is the navigation operation.
4. Derivation 17 can also be extended for the case where the cardinality of `attr` is greater than 1. However this situation is rarely encountered.

### 4.3 Substitution unification

In the previous section, we have presented the principles for modelling an elementary constraint. In general, each elementary constraint gives rise to a B substitution. This section discusses the way to unify the derived substitutions.

**Derivation 18 (Substitution unification)** Given “ $P_1$  and ... and  $P_k$ ” an OCL and expression on the elementary constraints  $P_1, \dots, P_k$  in the postconditions of `oper`. The B substitutions derived from this expression are unified in the following manner:

- all the possible temporary variables as well as the possible created objects are declared in the same clause **any...where...;**
- all the substitutions involving the same B variable are unified;
- the substitutions for the different B variables are placed on parallel (“||”) and they constitute the body of the clause **any...where...** if this clause has been created.

**Derivation 19 (The expression if ... then ... else ... endif)** Given `if cond then expr1 else expr2 endif` an expression of postconditions of an operation `oper`. The B formalisation of this expression gives rise to a clause:

- **if cond then subst(expr<sub>1</sub>) else subst(expr<sub>2</sub>) end**, if `expr2` is not an expression `if ... then ... else ... endif` or
- **if cond then subst(expr<sub>1</sub>) elsif cond<sub>2</sub> then subst(expr<sub>21</sub>) ... end**, if `expr2` is also in the form `if cond2 then expr21 ...;`

where `cond` is the B formalisation of `cond` and `subst(expr)` denotes the substitutions derived from `expr`.

**Remark 9** `skip` is the substitution of true.

**Derivation 20 (Operation)** Given {pre,post} the preconditions and the postconditions in OCL of an operation `oper`, where `post` is defined according to Definition 3. The B operation `oper` modelling `oper` is generated in the following manner:

- the signature and a part of precondition of `oper` for typing possible parameters are generated according to derivation schemes described in [17, 13];
- if `pre` is not true then the precondition in `oper` is augmented by predicates derived from `pre` in using derivation schemes in Section 3;
- the substitution part of `oper` is generated from `post` using derivation schemes previously presented in this section.

In the context of postconditions Remark 5 may be applied. However, there is also another alternative for `let` expressions as described in Derivation 21.

**Derivation 21 (let ... in postconditions)** The expression `let v1 : type1 = val1, ..., vn : typen = valn in expr` in postconditions of an operation `oper` gives rise to a clause `Let... in oper`:

```
let v1, ..., vn be  
  v1 = val1 ∧  
  .. ∧  
  vn = valn  
in  
  subst(expr)  
end
```

where `v1, ..., vn, val1, ..., valn` and `subst(expr)` denote the B formalisation of `v1, ..., vn, val1, ..., valn` and `expr`.

## 5 Transformation example

The class diagram in Figure 1 is extracted from the UML specification [13] for the pump component in a system controlling petrol dispensing, customer payment handling and petrol tank level monitoring as described in chapter 6 of [6].

The class diagram is composed of an aggregation of the classes `Pump`, `Clutch`, `Display`, `Gun`, `Holster` and `Motor` which model five pumps with included components. We defined the class `Delivery` to model deliverance. The attributes, the operations as well as data types used by attributes and operations are presented in [6].

### 5.1 The operation `Pump::enable.Pump` in OCL

To illustrate the application of derivation schemes, let's consider the OCL specification (Figure 2) of the class operation `Pump::enable.Pump`.

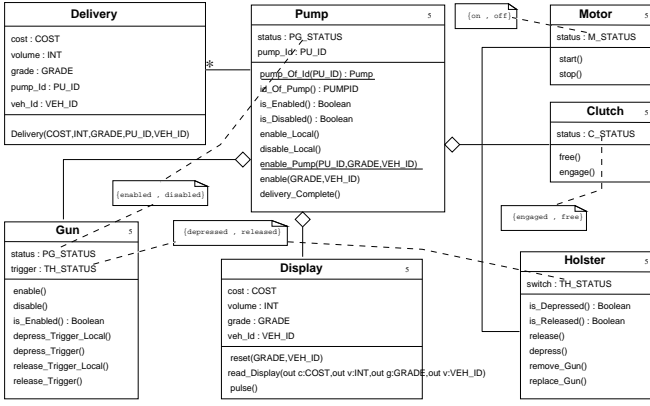


Figure 1. Pump class diagram

```

CONTEXT Pump::enable_Pump(pi : PU_ID, gg : GRADE,
                           vi : VEH_ID) : void
PRE
  Pump.allInstances → collect(pump_Id) → includes(pi)
POST
  let pp : Set(Pump) = Pump.allInstances → select(tt |
    tt.pump_Id@pre=pi and tt.status@pre=disabled)
  in
  if pp → notEmpty then
    pp → for all(p | p.status = enabled) and
    pp → for all(p | p.display.grade = gg) and
    pp → for all(p | p.display.cost = costOfGrade(gg)) and
    pp → for all(p | p.display.volume = 0) and
    pp → for all(p | p.display.veh_Id = vi) and
    pp → for all(p | p.motor.status = on) and
    pp → for all(p | p.clutch.status = free)
  else true endif

```

Figure 2. Operation enable\_Pump in OCL

The operation has three arguments: the identifier  $pi$  of the pump to be invoked; the category  $gg$  of the petrol to be distributed and the registration number  $vi$  of the vehicle. The precondition says that the pump with identifier  $pi$  exists. The postcondition is formatted according to Definition 3 and says that in case the pump with identifier  $pi$  is disabled, it should be enabled and its display is initialised, its motor is running and its clutch is free, otherwise nothing happens. Notice that although there is only one pump with the identifier  $pi$  but we can not extract it from the set of effective pumps (denoted by  $Pump.allInstances$ ) due to restrictions of OCL collection operations. That is why the postcondition is composed of expressions on a singleton  $pp$  whose unique element is the disabled pump with identifier  $pi$ .

## 5.2 The operation Pump::enable\_Pump in B

Figure 3 represents the B specification of the operation Pump::enable\_Pump which is derived from the OCL specification in Figure 2 by applying systematically Deriva-

tion 20 and implied derivation schemes in Section 3 and Section 4.

```

...
OPERATIONS
  pump_enable_Pump(pi, gg, vi) =
  pre
    pi ∈ PU_ID ∧
    gg ∈ GRADE ∧
    vi ∈ VEH_ID ∧
    pi ∈ dom({tt, nn | tt ∈ pump_pump_Id[pump] ∧ nn ∈ N ∧
      nn = card({xx | xx ∈ pump ∧ pump_pump_Id(xx) = tt})})
  then
  let pp be
    pp = {tt | tt ∈ pump ∧ pump_pump_Id(tt) = pi ∧
      pump_status(tt) = pg_status_disabled}
  in
  if ¬(pp = φ) then
    pump_status := pump_status ◁
      pp × {pg_status_enabled} ||
    display_grade := display_grade ◁
      displayPump[pp] × {gg} ||
    display_cost := display_cost ◁
      displayPump[pp] × {costOfGrade(gg)} ||
    display_volume := display_volume ◁
      displayPump[pp] × {0} ||
    display_veh_Id := display_veh_Id ◁
      displayPump[pp] × {vi} ||
    motor_status := motor_status ◁
      motorPump[pp] × {m_status_on} ||
    clutch_status := clutch_status ◁
      clutchPump[pp] × {c_status_free}
  else skip end
  end
end;
...

```

Figure 3. Operation enable\_Pump in B

As an example, let's consider the expression  $Pump.allInstances \rightarrow collect(pump\_Id) \rightarrow includes(pi)$  in the OCL precondition of Pump::enable\_Pump. Analysing its sub-expressions we know that: the expression  $Pump.allInstances$  denotes the set of effective instances of the class Pump. Applying the operation  $collect(pump\_Id)$  on set  $Pump.allInstances$  we obtain a bag. The operation  $includes(pi)$  will check whether the identifier  $pi$  is an element of the bag  $Pump.allInstances \rightarrow collect(pump\_Id)$ . Let  $T(expr)$  denotes the B formalisation of the OCL expression  $expr$  and  $pump\_pump\_Id$  the B variable for the attribute  $pump\_Id$  of the class Pump. The transformation process from the precondition  $Pump.allInstances \rightarrow collect(pump\_Id) \rightarrow includes(pi)$  into B proceeds as follows:

$$\begin{aligned}
& T(Pump.allInstances \rightarrow collect(pump\_Id) \rightarrow includes(pi)) \\
& \quad \equiv (cf. \text{ point 1 in Derivation 8}) \\
& \quad pi \in dom(T(Pump.allInstances \rightarrow collect(pump\_Id))) \\
& \quad \equiv (cf. \text{ point 4 in Derivation 8}) \\
& \quad pi \in dom(\{tt, nn | tt \in pump\_pump\_Id[T(Pump.allInst- \\
& \quad \text{ances})] \wedge nn \in N \wedge nn = card(\{xx | xx \in T(Pump.allInst- \\
& \quad \text{ances}) \wedge pump\_pump\_Id(xx) = tt\})\}) \\
& \quad \equiv (cf. \text{ Derivation 6})
\end{aligned}$$

$$pi \in \text{dom}(\{tt, nn | tt \in \text{pump\_pump\_Id}[pump] \wedge nn \in \mathcal{N} \wedge nn = \text{card}(\{xx | xx \in \text{pump} \wedge \text{pump\_pump\_Id}(xx) = tt\})\})$$

**Remark 10** In order to avoid the eventual name conflicting in the derivation of attributes into B, we prefix the name of attributes by the the class name. Hence the attribute `pump_Id` is modelled by the variable `pump_pump_Id`. It is similar for constants in the enumeration type; the constant enabled of the enumeration type PG\_STATUS is modelled by `pg_status_enabled`.

Let's consider now the application of derivation schemes specific for postconditions. According to Derivation 21, we transform the let expression in the postcondition of the OCL operation `Pump::enable_Pump` into a clause **let...** in the B operation `pump_enable_pump`. The substitution body **if...** of **let...** is derived from the body expression `if ...` of `let ...` using Derivation 19. The B guard condition  $\neg(pp = \phi)$  is derived from the OCL guard condition `pp->notEmpty`. The body of the substitution **if...** is generated from the the expression `if ...` according to Derivation 18 and Derivation 17.

## 6 Conclusion

This paper presents a systematic way for transforming OCL expressions into B, which can be applied to generate B supplementary invariants as well as B abstract operations in B specifications generated according to the derivation procedure in our previous work [13], from the OCL specifications for the supplementary class invariants and for UML operations in UML specifications.

For the further work, the prototype `ArgoUML+B` has been developed from `ArgoUML`<sup>3</sup>, a platform for editing UML diagrams with the code in java and freely available. We have added to `ArgoUML` the possibility to transform a set of class and collaboration diagrams into a B specification according the derivation procedure in [13]. In `ArgoUML`, there is a component "ocl-argo", which is in charge to parse OCL expressions. We would like to extend this component by implementing our derivation schemes from OCL into B. So that the OCL constraints within UML class diagrams can be transformed into B.

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<sup>3</sup><http://www.ArgoUML.org>