



## Lightweight symmetric primitives

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Submitted on 3 Feb 2025

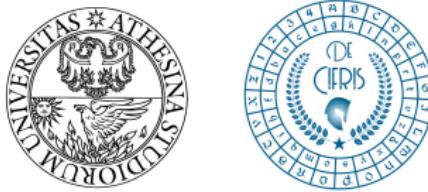
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# De Cifris Trends in Modern Cryptography: *the French Magisterium*



## Lecture 3



# Lightweight Symmetric Primitives

Anne Canteaut

Inria Paris, France

The Inria logo is a red, stylized, handwritten-style word "Inria".



# INTRODUCTION



20+ years later, AES is still secure



The banner features the NIST logo at the top left, followed by the text "Information Technology Laboratory" and "COMPUTER SECURITY RESOURCE CENTER". Below the banner are two green buttons labeled "UPDATES" and "2023".

## NIST Updates FIPS 197, Advanced Encryption Standard (AES)

May 09, 2023



Today, NIST has published an update of Federal Information Processing Standards Publication (FIPS) 197, [Advanced Encryption Standard \(AES\)](#). This update makes no technical changes to the algorithm specified in the standard, which was originally published in 2001.

... [View full post >](#)



# New implementation constraints

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## Questions raised about Oyster card security

Its RFID chip is cracked by researchers

By Network World and Computerworld UK staff | Published 12:20, 07 March 08

0 0 0 0 17

Smartcards with encrypted RFID chips, including London's Oyster fare card, might not be as secure as previously thought.

New research at the University of Virginia is causing a major stir in Boston, because it raises question over the smart "CharlieCards" used by commuters on the city's 'T' metro system.

However, London's Oyster card uses similar RFID technology – the Mifare Classic made by Philips spinoff NXP Semiconductors.

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# New implementation constraints

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# NEW YORK POST

NEWS

## Cheney feared terrorists would 'hack' pacemaker

By Bob Fredericks      October 19, 2013 | 4:11am

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PLAY CBS NEWS VIDEO



# Lightweight Competitions

## CAESAR for authenticated encryption (2014-2019):

<https://competitions.cr.yp.to/caesar.html>

Use case 1: Lightweight (resource-constrained) applications

- ① Ascon [Dobraunig, Eichlseder, Mendel, Schläffer 16]
- ② Acorn [Wu 14]

## NIST Lightweight Cryptography standardization (2015-23):

Ascon family (announced in Feb. 2023)



# Why is minimalism interesting?



Tretjakov Gallery

Besides (niche) application needs, it helps us **understand where security comes from.**



# Outline

- ① Symmetric encryption
- ② How to make it lightweight?
- ③ Ascon
- ④ Possible weaknesses coming from “minimal” building-blocks:
  - Simple key-schedule
  - ...



# Symmetric encryption



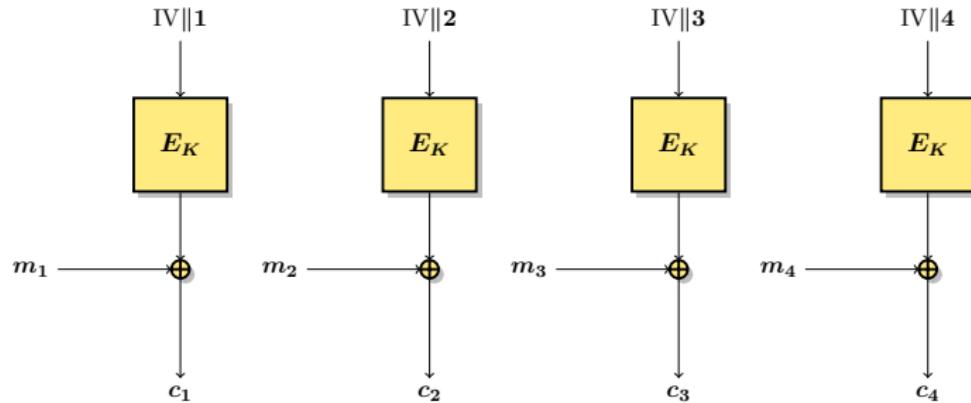
# Encryption scheme

Two steps for encrypting plaintexts of an arbitrary length:

- ① design a permutation, or a family of permutations, operating on  $\{0, 1\}^n$ .
- ② design a mode of operation describing how this primitive can be used for encrypting messages of any length (e.g. CTR, CBC).



## CTR mode for encryption



where  $E_K$  = family of permutations of  $\{0, 1\}^n$  indexed by the key.



# Practical Pseudo-Random Permutation

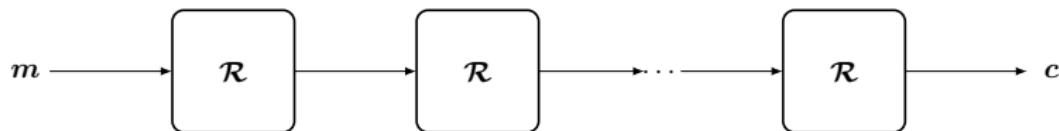
$$E_K : \{0, 1\}^n \longrightarrow \{0, 1\}^n$$

- indistinguishable from randomly chosen bijections of  $\{0, 1\}^n$  with  $n \in \{64, 128\}$
- implementable

→ Contradiction!

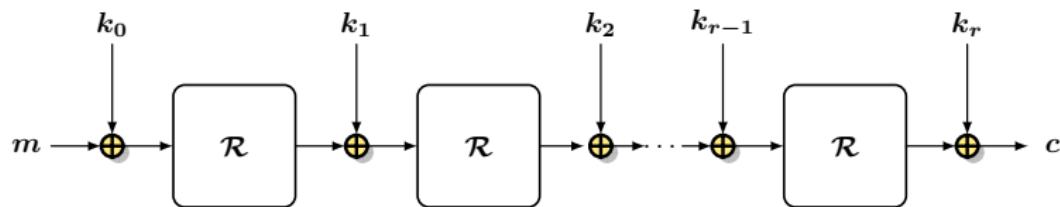


# Iterated construction



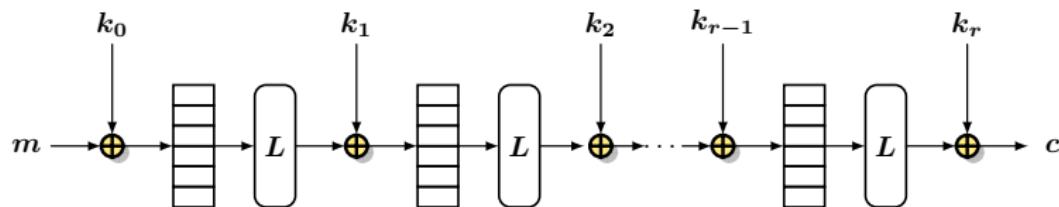


# Iterated construction





# Iterated construction





- blocksize: 128 bits
- **10 rounds** for the 128-bit key version
- Sbox operates on 8 bits
- diffusion layer is linear over  $F_{2^8}$
- nonlinear key schedule.



# How to make it lightweight?



# Lightweight block ciphers

## AES [Daemen-Rijmen 98][FIPS PUB 197]

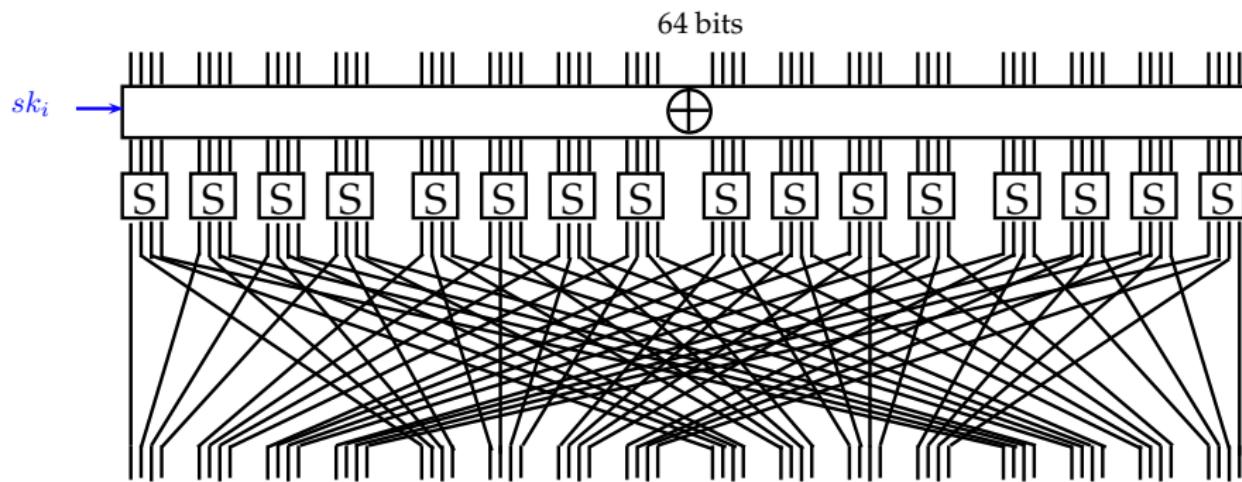
- blocksize: 128 bits
- Sbox operates on 8 bits
- diffusion layer is linear over  $\mathbb{F}_{2^8}$

## To make it smaller in hardware:

- blocksize: 64 bits
- smaller Sbox, on 3 or 4 bits
- linear diffusion layer over a smaller alphabet
- simplified key-schedule



# PRESENT [Bogdanov et al. 07]



31 rounds (+ a key addition)



## Lightweight but secure...

Increase the number of rounds!

- PRESENT [Bogdanov et al. 07]: 31 rounds
- LED [Guo et al. 11]:  
LED-64: 32 rounds, LED-128: 48 rounds
- SPECK [Beaulieu et al. 13]:  
SPECK64/128: 27 rounds, SPECK128/256: 34 rounds
- SIMON [Beaulieu et al. 13]:  
SIMON64/128: 44 rounds, SIMON128/256: 72 rounds

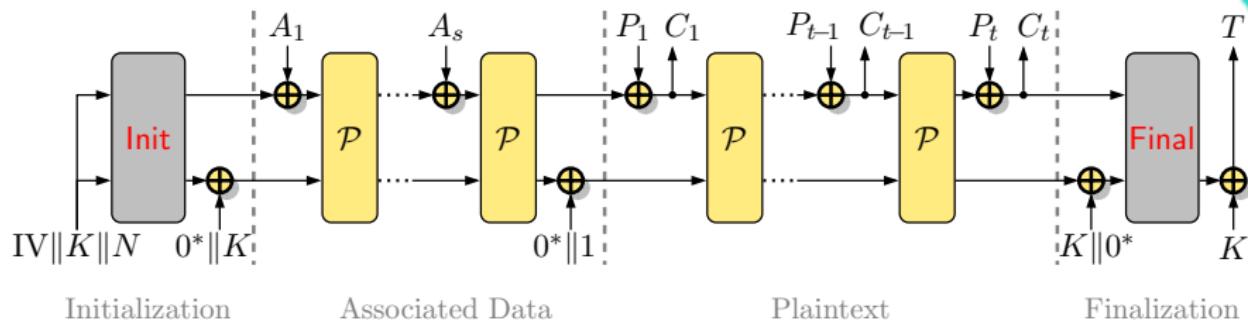


Does lightweight mean “light + wait” ?  
[Knežević et al. 12]



# Duplex-Sponge mode for AEAD encryption

[Bertoni et al. 12]



where  $\mathcal{P}$  is a permutation of  $\{0, 1\}^n$ .

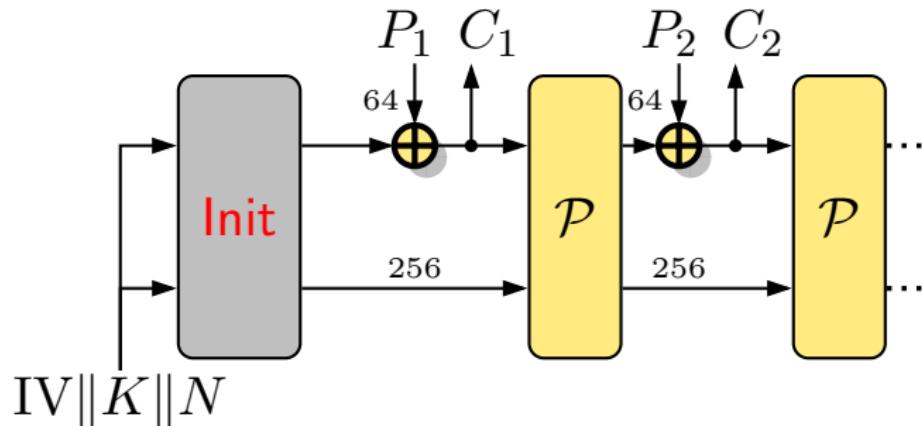


# Ascon



# Duplex-Sponge mode in Ascon

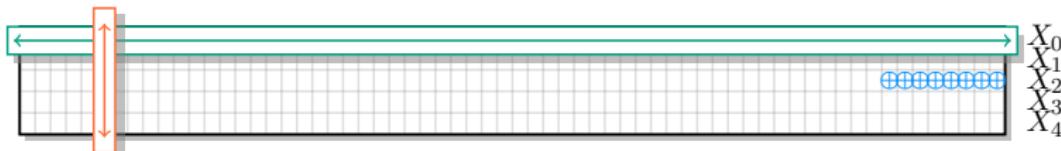
[Dobraunig, Eichlseder, Mendel, Schläffer 16]



where  $\mathcal{P}$  is a permutation on 320 bits  
of which 64 are known/controlled.



## $\mathcal{P}$ in Ascon



Permutation operating on a 320-bit state:

- 8-bit constant addition;
- Nonlinear Sbox on 5 bits of degree 2 (on the 64 columns);
- 5 simple linear transformations on 64 bits

$$\Sigma_i(X_i) = X_i \oplus (X_i \ggg a_i) \oplus (X_i \ggg b_i)$$

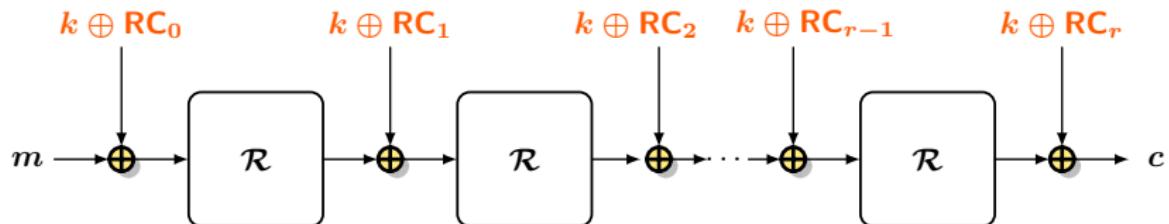
→ 6 rounds



# Possible weaknesses coming from minimal building-blocks



## Lightweight key schedules



where  $RC_0, RC_1, \dots, RC_r$  are fixed round-constants.



# Invariants for Midori-64

[Guo et al. 16][Todo et al. 16]

**Midori-64** [Banik et al. 15]:

$$E_K : (\{0, 1\}^4)^{16} \longrightarrow (\{0, 1\}^4)^{16}$$

with  $K = (k_0, k_1, k_0 \oplus \text{RC}_2, k_1 \oplus \text{RC}_3, k_0 \oplus \text{RC}_4, \dots)$

**Invariant set for Midori-64:**

If  $k_0, k_1 \in \{0x0, 0x1\}^{16}$ , then  $\{0x8, 0x9\}^{16}$  is invariant under  $E_K$ .

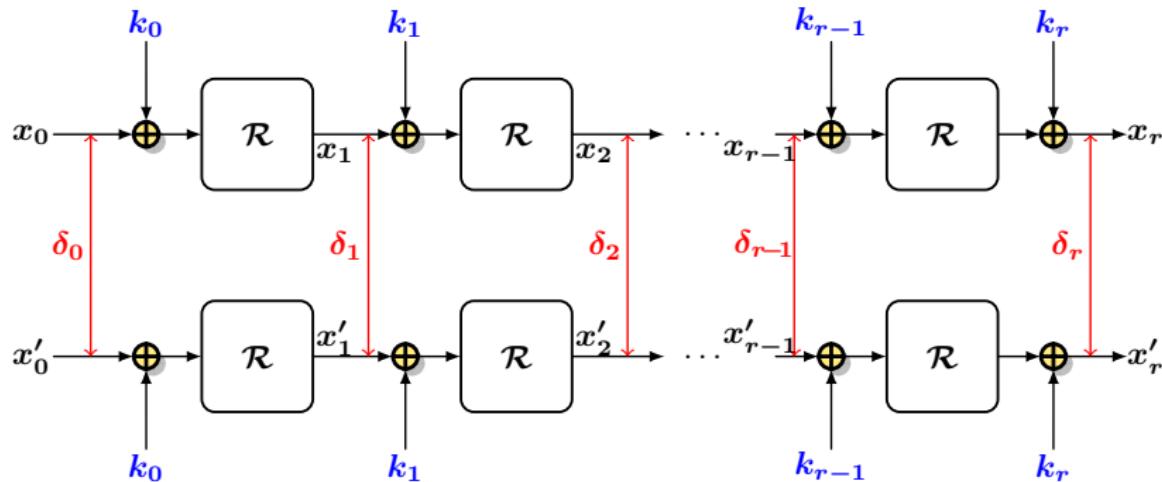
**Example :**

For  $(k_0, k_1) = (0x1100110011001100, 0x0011001100110011)$ ,  
 $m = 0x9999999999999999 \mapsto c = 0x8999999988988989$



# Probability of a differential path

[Biham, Shamir 90]





# Fixed-key differential paths for $r$ iterations of Midori-64

Differential path  $(\delta, \delta, \dots, \delta)$  with  $\delta = (0x1)^{16}$

- On average over all key sequences

$$\text{EDP}(\delta, \dots, \delta) = 2^{-48r}$$

- For a fixed key

$$\text{DP}(\delta, \dots, \delta) = \begin{cases} 2^{-48} & \text{if } \forall i, k_i \in \{0x0, 0x1\}^{16} \\ 0 & \text{otherwise.} \end{cases}$$



## Where does this come from?

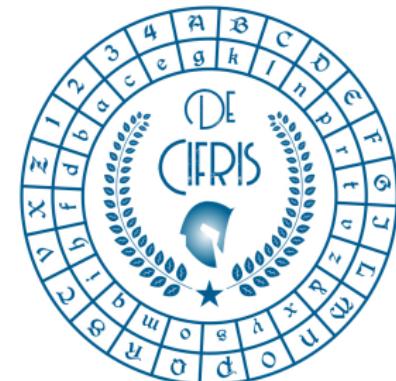
- The key schedule is **compatible** with the fact that each  $k_i$  lies in a **(affine) subspace**.
- The linear diffusion function is defined by an **orthogonal matrix**.



## Takeaway



Minimalism helps us understand where security comes from.  
It motivates new attacks, and new design criteria.



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