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Strong Shift Equivalence and Bialgebras in traced monoidal categories

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E. Jeandel

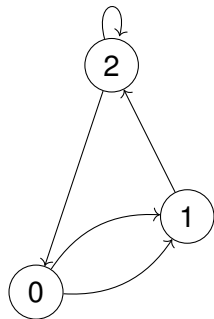
Université de Lorraine, France

Plan

- 1 Symbolic Dynamics
- 2 Categories for symbolic dynamics
- 3 Applications
- 4 Conclusion

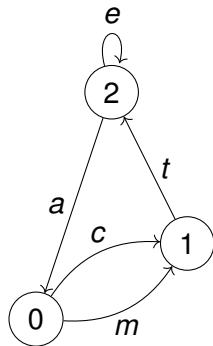
Edge shift

An edge shift is the set of all biinfinite walks (on edges) in a finite graph



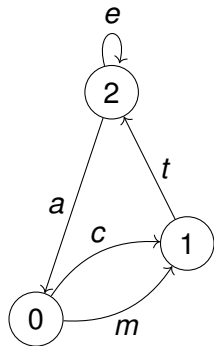
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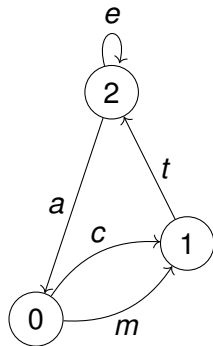
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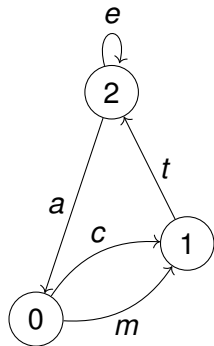
Conjugacy

Edge shifts are conjugate if they're isomorphic via local transformations



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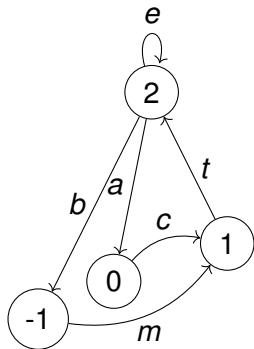
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$ac \rightarrow ac$

$am \rightarrow bm$

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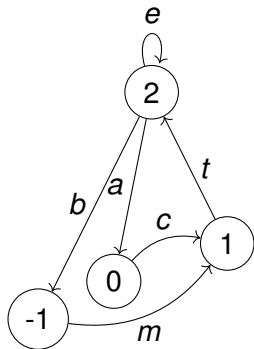
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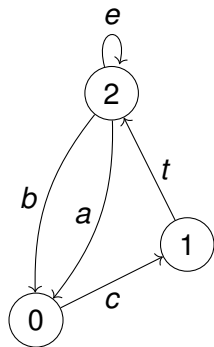
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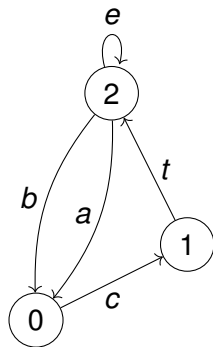
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Conjugacy

Edge shifts are conjugate if they're isomorphic via local transformations



More generally, choose a vertex to duplicate

- Duplicate the inputs
- Redistribute the outputs

or vice-versa

SSE

This can be defined directly on adjacency matrices:

Definition

Two matrices M and N are 1-step equivalent if $M = RS$ and $N = SR$ for (nonnecessarily square) nonnegative integral matrices R, S

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Two matrices M and N are 1-step equivalent if $M = RS$ and $N = SR$ for (nonnecessarily square) nonnegative integral matrices R, S

Example:

$$M = \begin{pmatrix} 0 & 2 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \end{pmatrix} \quad N = \begin{pmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 1 \end{pmatrix}$$

$$R = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \quad S = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \end{pmatrix}$$

SSE

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Definition

Strong shift equivalence (SSE) is the transitive closure of 1-step equivalence.

Graphs G_1 and G_2 represent conjugate edge shifts iff their adjacency matrices are SSE.

Main open problem of symbolic dynamics: decide conjugacy/SSE

- Open since the 70s
- Decidable for matrices in \mathbb{Z} (Krieger, 1980)
 - (almost the) same as conjugacy in $GL_n(\mathbb{Z})$
- Decidable for one-sided edge-shifts (Williams, 1973)
 - The rewriting system on graphs is confluent.

Definition

An invariant is a function ψ s.t. $\psi(M) = \psi(N)$ if M and N are SSE.

- To prove that M and N are NOT SSE, just find an invariant that distinguishes them

Examples:

- Trace of the matrix
- More generally, number of cycles of size k
- The entropy (exponential growth of the number of paths of size n)
- ...

In this talk: A systematic way to find invariants

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Goal

Build a category where arrows can represent matrices and equivalence classes of SSE.

Starting point


Folklore

The bialgebra PROP is exactly the PROP of nonnegative integer matrices.

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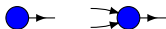
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The bialgebra PROP is exactly the PROP of nonnegative integer matrices.

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A cocommutative comonoid $(\epsilon : 1 \rightarrow 0, \Delta : 1 \rightarrow 2)$

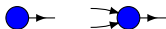


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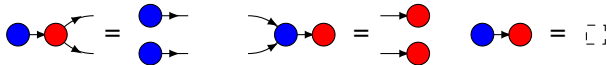
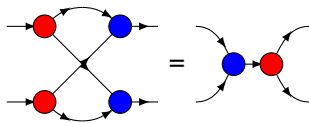
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Bialgebra rules



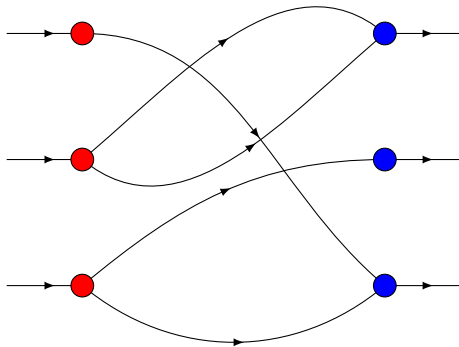
Starting point

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The bialgebra PROP is exactly the PROP of nonnegative integer matrices.

arrows $n \rightarrow m$ are exactly matrices of size $m \times n$.

$$M = \begin{pmatrix} 0 & 2 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \end{pmatrix}$$



Technically, our matrices represent graph, so with inputs = outputs

SSE states that $RS \equiv SR$, how can we simulate that ?

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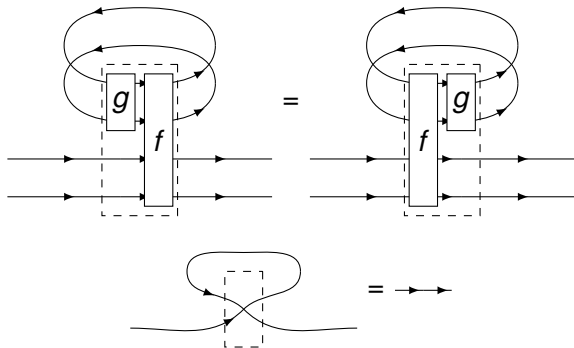
With a trace!

Trace

Consider the traced bialgebra prop (the previous prop + a trace).

Trace

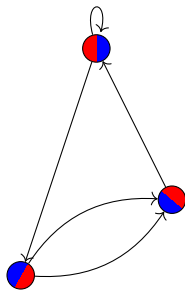
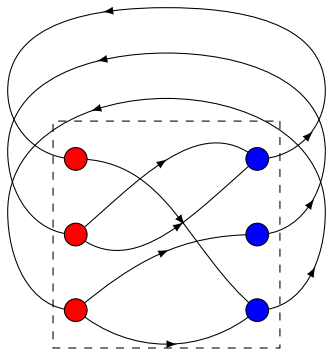
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Trace

In this prop, we can associate a scalar $t(M)$ to each matrix M .

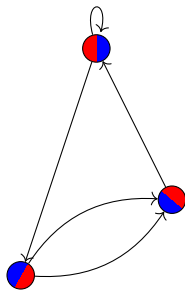
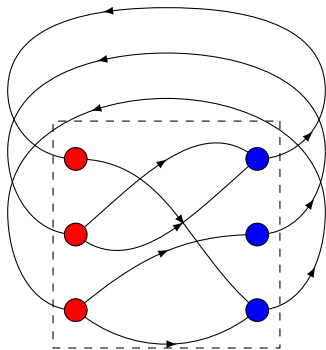
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Trace

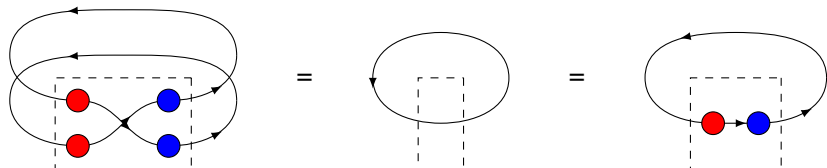
In this prop, we can associate a scalar $t(M)$ to each matrix M .
Is it true that $t(M) = t(N)$ iff M and N are SSE ?

$$M = \begin{pmatrix} 0 & 2 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \end{pmatrix}$$



Flow equivalence

$M = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ and $N = (1)$ satisfy $t(M) = t(N)$ but are not SSE.



Flow equivalence

$t(M) = t(N)$ iff M and N are flow-equivalent, another equivalence notion coming from symbolic dynamics

(First observed by David Hillman, 1995)

We have lost a notion of time, that we need to recover

Solution

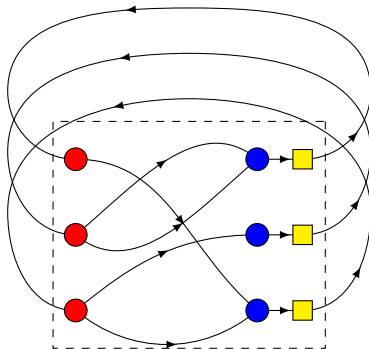
We add a generator that is a morphism for both the monoid and the comonoid.



Solution

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Note: this new prop represents matrices in $\mathbb{Z}_+[t]$ rather than \mathbb{Z}_+

Solution

We add a generator that is a morphism for both the monoid and the comonoid.

Theorem

In this prop, $t(M) = t(N)$ iff M and N are SSE.

Needs some new (easy) results in symbolic dynamics for the proof.

To recap: we need one bialgebra, a bialgebra morphism and a trace.

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What do we gain from it ?

- If one knows a concrete representation of this prop, one can decide conjugacy
- One can find invariants !

- Let \mathcal{C} be our category.
- Let \mathcal{D} be any traced category that contains a bialgebra
- By the universal property, there is a functor from \mathcal{C} to \mathcal{D}
- Using this functor, we can associate to each matrix M an arrow $\psi(M)$ in \mathcal{D} s.t. $\psi(M) = \psi(N)$ if M and N are SSE.

Good news: There are a lot of bialgebras in the wild

Bad news:

- Some of them are in categories that are not traced
- Some of them are Hopf algebras
 - Hopf algebras represent matrices with coefficients in \mathbb{Z} , not in \mathbb{Z}_+

Example from algebra 1/2

- If M is a finite abelian monoid, $\mathbb{K}[M]$ has a structure of a bialgebra
 - Product is the monoid product extended linearly
 - Coproduct is the copy extended linearly

Theorem

Let \mathcal{M} be an additive monoid, and h an homomorphism. Then the number of solutions in \mathcal{M} of the equation $h(Mx) = x$ is an invariant for SSE.

Example from algebra 2/2

- $\mathbb{K}[X]$ has a structure of a bialgebra
 - Product is the product of polynomials
 - Coproduct is defined by $\Delta(X) = 1 \otimes X + X \otimes 1$

Problem: no trace!

Example from algebra 2/2

- $\mathbb{K}[X]$ has a structure of a bialgebra
 - Product is the product of polynomials
 - Coproduct is defined by $\Delta(X) = 1 \otimes X + X \otimes 1$

Solution: Replace $\mathbb{K}[X]$ by formal series in a complete semiring (like $\mathbb{R}_+ \cup \{+\infty\}$).

If we take the morphism $h(X) = tX$ we get:

Theorem

The quantity $f_M(t) = \frac{1}{\det(I-tM)}$ is an invariant for SSE

This is the well-known Zeta function of a subshift.

Example from Category Theory

Let \mathcal{D} be a category with finite limits and colimits.

- Suppose one has a monoid X in a \mathcal{D} with a morphism h .
- It extends to a bialgebra using the diagonal map as a comonoid.

Problem: What if there is no trace in \mathcal{D} ?

Example from Category Theory

Let \mathcal{D} be a category with finite limits and colimits.

- Suppose one has a monoid X in a \mathcal{D} with a morphism h .
- It extends to a bialgebra using the diagonal map as a comonoid.

Solution: Replace \mathcal{D} with $\text{cospan}(\mathcal{D})$ which is a compact category.

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Theorem

Let M be a $n \times n$ matrix.

The object that is the coequalizer of $X^n \begin{array}{c} \xrightarrow{id^n} \\ \xrightarrow{h^n \circ M} \end{array} X^n$ is an invariant for SSE.

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- Starting with the monoid $(\mathbb{Z}, +)$ in the category of abelian groups, we obtain the *Bowen-Franks* group (1977)
- Starting with the monoid $(\mathbb{Z}[t], +)$ in the category of $\mathbb{Z}[t]$ -modules, we obtain the dimension group of Krieger (1977)

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Where to go from here ?

- We have a systematic way to find invariants
- Investigate new bialgebras to find new invariants!