



HAL
open science

Fully Interval Integer Transshipment Problem- A Solution Approach

G. Sudha, G. Ramesh, D. Datta, K. Ganesan

► **To cite this version:**

G. Sudha, G. Ramesh, D. Datta, K. Ganesan. Fully Interval Integer Transshipment Problem- A Solution Approach. 4th International Conference on Intelligence Science (ICIS), Feb 2021, Durgapur, India. pp.196-206, 10.1007/978-3-030-74826-5_17 . hal-03741707

HAL Id: hal-03741707

<https://inria.hal.science/hal-03741707>

Submitted on 1 Aug 2022

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.



Distributed under a Creative Commons Attribution 4.0 International License

Fully interval integer transshipment problem- A solution approach

G. Sudha¹, G. Ramesh², D. Datta³, K. Ganesan⁴

^{1, 2, 4}Department of Mathematics, SRM Institute of science and Technology,
Kattankulathur, Tamil Nadu, India.

³Adjunct faculty, Department of Mathematics, SRM Institute of science and
Technology, (former) Health Physics Division, Bhabha Atomic Research Centre,
Trombay, Mumbai-400085, India

E-Mail: sudhag@srmist.edu.in;

rameshapg@gmail.com; dbbrt_datta@yahoo.com; ganesank@srmist.edu.in

Abstract—

We present an improved interval Vogel's approximation algorithm for the initial feasible solution of a transshipment problem involving fully interval integer numbers in this paper. The interval variant of MODI algorithm is used to measure the optimality of the interval feasible solution. In a transportation problem, shipments are only permitted between source-sink pairs, while in a transshipment problem, shipments can be permitted both between sources and sinks. Furthermore, there may be points where units of a substance may be transferred from a source to a sink. To demonstrate the effectiveness of the proposed algorithm, a numerical example is given.

Keywords— Interval Numbers, New interval arithmetic, Ranking, Interval initial feasible solution, Improved interval Vogel's approximation algorithm.

1. INTRODUCTION

Goods are transported from specific sources to specific destinations in classic transportation problems (TP). However, in some situations, products can be delivered to a specific location either directly or through one or more intermediary nodes. Either of these nodes has the ability to function as an origin or a destination. Transshipment issues are models with these additional features. Thus, transshipment issues are a type of transportation problem in which merchandise are elated from a source to a destination via various intermediary nodes (sources/destinations), with the possibility of changing modes of transportation or consolidating or de-consolidating shipments. A transshipment problem's main goal is to find the lowest transportation cost while

satisfying supply and demand at each node. These issues have found a lot of use in the e-commerce era, when the online shopping business is thriving thanks to its customers' convenience. To solve the transshipment problem, translate it to an equivalent TP first, and then use existing techniques to find an initial simple feasible solution. Furthermore, an optimal solution of equivalent TP can be obtained using the modified distribution (MODI) process. Orden [17] introduced the concept of transshipment problem in the year 1956. When the decision parameters are known vaguely, the classical mathematical techniques are successfully applied to model and solve the transshipment problems. But actually there are problems that the charge coefficients and the supply and demand quantities may be uncertain due to many reasons. Hence we must deal with vague in order in formulating and solving real earth troubles. Hurt et.al [5] have developed an alternative formulations of the transshipment problem.

Ganesan et.al [3] discussed on arithmetic operations of interval numbers and some properties of interval matrices. Pandian et.al [11] studied a new method for finding an optimal solution of fully interval integer transportation problems. NagoorGani et.al [8] developed a transshipment problem in fuzzy environment. Rajendran et.al [12] discussed fully interval transshipment problems. Akilbasha et.al [1] formulated an innovative exact method for solving fully interval integer transportation problems.

Sophia Porchelvi et.al [14] studied a comparative study of optimum solution between interval transportation and interval transshipment problem. In this paper, we propose a new interval arithmetic and an improved interval Vogel's approximation algorithm to solve the problem of interval transshipment without translating into crisp form.

The structure of this article is given as follows: Section 2 deals with the basic definitions. Section 3 deals with the mathematical formulation of fully interval transshipment problem. Section 4 introduces improved interval Vogel's approximation algorithm under generalized interval arithmetic. Section 5 provides an example to illustrate the theory developed in this paper. Section 6 concludes this manuscript.

2. PRELIMINARIES

The purpose of this segment is to provide some observations, ideas and results which are useful in our further consideration .

2.1. Interval numbers

Let $\tilde{a} = [a_1, a_2] = \{x \in \mathbb{R} : a_1 \leq x \leq a_2 \text{ and } a_1, a_2 \in \mathbb{R}\}$ be an interval on the real line \mathbb{R} . If $\tilde{a} = [a_1, a_2] = [a, a]$, then $\tilde{a} = [a, a] = a$ is a real number (or a degenerate interval). The terms interval and interval number are used interchangeably. The mid-point and width (or half-width) of an interval number $\tilde{a} = [a_1, a_2]$ are defined as $m(\tilde{a}) = \left(\frac{a_1 + a_2}{2}\right)$ and $w(\tilde{a}) = \left(\frac{a_2 - a_1}{2}\right)$.

The interval number \tilde{a} can also be expressed in terms of its midpoint and width as $\tilde{a} = [a_1, a_2] = \langle m(\tilde{a}), w(\tilde{a}) \rangle$. We use \mathbb{IR} to denote the set of all interval numbers defined on the real line \mathbb{R} .

2.2. Ranking of Interval Numbers

Sengupta and Pal [13] suggested a easy and powerful index to compare any two intervals on \mathbb{IR} through the satisfaction of decision-makers .

Definition 2.2.1. Let \preceq be an extended order relation between the interval numbers $\tilde{a} = [a_1, a_2]$, $\tilde{b} = [b_1, b_2]$ in \mathbb{IR} , then for $m(\tilde{a}) < m(\tilde{b})$, we construct a premise $(\tilde{a} \preceq \tilde{b})$ which implies that \tilde{a} is inferior to \tilde{b} (or \tilde{b} is superior to \tilde{a}).

An *acceptability function* $A_{\preceq} : \mathbb{IR} \times \mathbb{IR} \rightarrow [0, \infty)$ is defined as:

$$A_{\preceq}(\tilde{a}, \tilde{b}) = A(\tilde{a} \preceq \tilde{b}) = \frac{m(\tilde{b}) - m(\tilde{a})}{w(\tilde{b}) + w(\tilde{a})}, \text{ where } w(\tilde{b}) + w(\tilde{a}) \neq 0.$$

A_{\preceq} may be interpreted as the grade of acceptability of the first interval number to be inferior to the second interval number. For any two interval numbers \tilde{a} and \tilde{b} in \mathbb{IR} either $A(\tilde{a} \preceq \tilde{b}) \geq 0$ (or) $A(\tilde{b} \succeq \tilde{a}) \geq 0$ (or) $A(\tilde{a} \preceq \tilde{b}) = 0$ (or) $A(\tilde{b} \succeq \tilde{a}) = 0$ (or) $A(\tilde{a} \preceq \tilde{b}) + A(\tilde{b} \succeq \tilde{a}) = 0$. If $A(\tilde{a} \preceq \tilde{b}) = 0$ and $A(\tilde{b} \succeq \tilde{a}) = 0$, then we say that the interval numbers \tilde{a} and \tilde{b} are equivalent (non-inferior to each other) and we denote it by $\tilde{a} \approx \tilde{b}$. Also if $A(\tilde{a} \preceq \tilde{b}) \geq 0$, then $\tilde{a} \preceq \tilde{b}$ and if $A(\tilde{b} \succeq \tilde{a}) \geq 0$, then $\tilde{b} \succeq \tilde{a}$.

2.3. A New Interval Arithmetic

Ming Ma et al. [7] suggested a new fuzzy arithmetic based on the index of locations and the index function of fuzziness. The location index number follows the usual arithmetic where as the fuzziness index functions are assumed to obey the lattice law which is the least upper bound and the greatest lower bound. That is for $a, b \in L$ we define $a \vee b = \max\{a, b\}$ and $a \wedge b = \min\{a, b\}$.

For any two intervals $\tilde{a} = [a_1, a_2]$, $\tilde{b} = [b_1, b_2] \in \mathbb{IR}$ and for $*$ $\in \{+, -, \cdot, /, \dots\}$ the arithmetic operations on \tilde{a} and \tilde{b} are defined as :

$$\begin{aligned} \tilde{a} * \tilde{b} &= [a_1, a_2] * [b_1, b_2] = \langle m(\tilde{a}), w(\tilde{a}) \rangle * \langle m(\tilde{b}), w(\tilde{b}) \rangle \\ &= \langle m(\tilde{a}) * m(\tilde{b}), \max\{w(\tilde{a}), w(\tilde{b})\} \rangle. \end{aligned}$$

In particular, here we should use following interval arithmetic operations like addition, subtraction, multiplication and division.

3. MAIN RESULTS

3.1. Mathematical formulation for using fully interval transshipment problem

In the transshipment table $O_1, O_2, O_3, \dots, O_1, \dots, O_m$ are sources from where goods are to be transported to destinations $D_1, D_2, D_3, \dots, D_j, \dots, D_n$. Any of the sources can transport to any of the destinations. \tilde{c}_{ij} is per unit transporting cost of goods from i^{th} source O_i to j^{th} destination D_j for all $i = 1, 2, \dots, m$ and $j = 1, 2, \dots, n$. \tilde{x}_{ij} is the amount of goods transporting from i^{th} source O_i to j^{th} destination D_j . \tilde{a}_i be the amount of goods available at the origins O_i and \tilde{b}_j the demand at the destination D_j . The mathematical model of fully interval transportation problem is as follows

$$\begin{aligned} \text{Minimize } \tilde{Z} &\approx \sum_{i=1}^m \sum_{j=1}^n \tilde{c}_{ij} \tilde{x}_{ij} \\ \text{subject to } &\sum_{j=1}^n \tilde{x}_{ij} \approx \tilde{a}_i, \quad i = 1, 2, 3, \dots, m \\ &\sum_{i=1}^m \tilde{x}_{ij} \approx \tilde{b}_j, \quad j = 1, 2, 3, \dots, n \\ &\sum_{i=1}^m \tilde{a}_i \approx \sum_{j=1}^n \tilde{b}_j, \quad \text{where } i = 1, 2, 3, \dots, m; \quad j = 1, 2, 3, \dots, n \text{ and} \\ &\tilde{x}_{ij} \succeq \tilde{0} \quad \text{for all } i \text{ and } j \end{aligned}$$

Table:1 Interval transshipment table

		1	2	...	i	...	m	$m+1$	$m+2$...	$m+j$...	$m+n$	
		O_1	O_2	...	O_i	...	O_m	D_1	D_2	...	D_j	...	D_n	supply
1	O_1	x_{11}	x_{12}	...	x_{1i}	...	x_{1m}	$x_{1,m+1}$	$x_{1,m+2}$...	$x_{1,m+j}$...	$x_{1,m+n}$	a_1
2	O_2	x_{21}	x_{22}	...	x_{2i}	...	x_{2m}	$x_{2,m+1}$	$x_{2,m+2}$...	$x_{2,m+j}$...	$x_{2,m+n}$	a_2
...
i	O_i	x_{i1}	x_{i2}	...	x_{ii}	...	x_{im}	$x_{i,m+1}$	$x_{i,m+2}$...	$x_{i,m+j}$...	$x_{i,m+n}$	a_i
...
m	O_m	x_{m1}	x_{m2}	...	x_{mi}	...	x_{mm}	$x_{m,m+1}$	$x_{m,m+2}$...	$x_{m,m+j}$...	$x_{m,m+n}$	a_m
$m+1$	D_1	$x_{m+1,1}$	$x_{m+1,2}$...	$x_{m+1,i}$...	$x_{m+1,m}$	$x_{m+1,m+1}$	$x_{m+1,m+2}$...	$x_{m+1,m+j}$...	$x_{m+1,m+n}$...
$m+2$	D_2	$x_{m+2,1}$	$x_{m+2,2}$...	$x_{m+2,i}$...	$x_{m+2,m}$	$x_{m+2,m+1}$	$x_{m+2,m+2}$...	$x_{m+2,m+j}$...	$x_{m+2,m+n}$...
...
$m+j$	D_j	$x_{m+j,1}$	$x_{m+j,2}$...	$x_{m+j,i}$...	$x_{m+j,m}$	$x_{m+j,m+1}$	$x_{m+j,m+2}$...	$x_{m+j,m+j}$...	$x_{m+j,m+n}$...
...
$m+n$	D_n	$x_{m+n,1}$	$x_{m+n,2}$...	$x_{m+n,i}$...	$x_{m+n,m}$	$x_{m+n,m+1}$	$x_{m+n,m+2}$...	$x_{m+n,m+j}$...	$x_{m+n,m+n}$...
Demand		b_1	b_2	...	b_j	...	b_n	

Since any source or destination may ship to any other source or destination in a transshipment problem, it would be convenient to number the origins and destinations sequentially, with the origins numbered 1 to m and the destinations numbered $m+1$ to $m+n$.

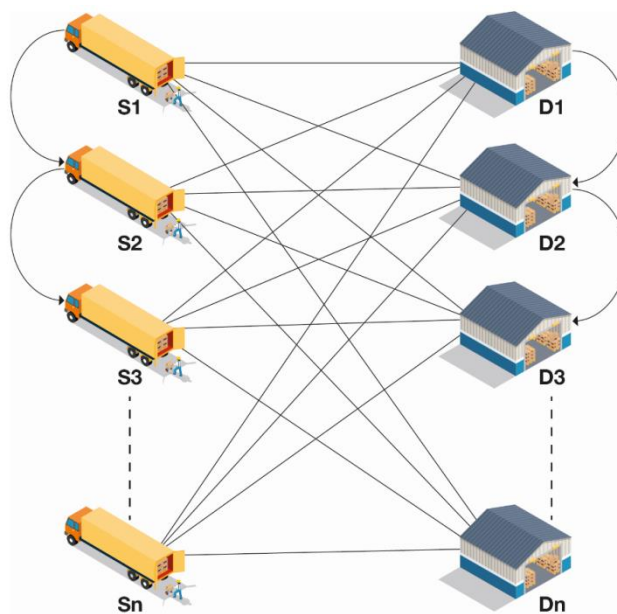


Figure 1

The unit cost of transport from a point chosen as a shipper to the same location designated as a receiver is set to zero.

$$\begin{aligned} \tilde{x}_{i1} + \tilde{x}_{i2} + \dots + \tilde{x}_{i,m+n} &= \tilde{a}_i + (\tilde{x}_{i1} + \tilde{x}_{i2} + \dots + \tilde{x}_{i,m+n,i}) \\ \tilde{x}_{i1} + \tilde{c}_{i2} + \dots + \tilde{c}_{i,i-1} \dots \tilde{x}_{i,i+1} + \dots + \tilde{x}_{i,m+n} &= \tilde{a}_i + (\tilde{x}_{i1} + \tilde{x}_{i2} + \dots + \tilde{x}_{i-1,i} + \tilde{x}_{i+1,i} + \dots + \tilde{x}_{m+n,i}) \\ \sum_{\substack{j=1 \\ j \neq i}}^n \tilde{x}_{ij} &\approx \tilde{a}_i + \sum_{\substack{j=1 \\ j \neq i}}^{m+n} \tilde{x}_{ji}, \quad i = 1, 2, 3, \dots, m \\ \sum_{\substack{j=1 \\ j \neq i}}^{m+n} \tilde{x}_{ij} - \sum_{\substack{j=1 \\ j \neq i}}^{m+n} \tilde{x}_{ji} &= \tilde{a}_i, \quad i = 1, 2, 3, \dots, m \end{aligned}$$

Similarly the total amount received at a destination D_j must be equal to its demand plus what it transships.

$$\begin{aligned} \tilde{x}_{1,m+j} + \tilde{c}_{2,m+j} + \dots + \tilde{c}_{m+j-1,m+j} + \tilde{c}_{m+j+1,m+j} + \dots + \tilde{c}_{m+n,m+j} \\ = \tilde{b}_{m+j} + \tilde{x}_{m+j,1} + \tilde{x}_{m+j,2} + \dots + \tilde{x}_{m+j,m+j-1} + \tilde{x}_{m+j,m+j+1} + \dots + \tilde{x}_{m+j,m+n} \\ \tilde{x}_{1,m+j} + \tilde{c}_{2,m+j} + \dots + \tilde{c}_{m+j-1,m+j} + \tilde{c}_{m+j+1,m+j} + \dots + \tilde{c}_{m+n,m+j} \\ = \tilde{b}_{m+j} + \tilde{x}_{m+j,1} + \tilde{x}_{m+j,2} + \dots + \tilde{x}_{m+j,m} + \tilde{x}_{m+j,m+j} + \dots + \tilde{x}_{m+j,m+n} \\ \sum_{\substack{i=1 \\ i \neq j}}^{m+n} \tilde{x}_{i,m+j} &\approx \tilde{b}_{m+j} + \sum_{\substack{i=1 \\ i \neq j}}^{m+n} \tilde{x}_{m+j,i}, \quad j = 1, 2, \dots, n \\ \sum_{\substack{i=1 \\ i \neq j}}^{m+n} \tilde{x}_{i,m+j} - \sum_{\substack{i=1 \\ i \neq j}}^{m+n} \tilde{x}_{m+j,i} &= \tilde{b}_{m+j}, \quad j = 1, 2, \dots, n \\ \sum_{\substack{i=1 \\ i \neq j}}^{m+n} \tilde{x}_{i,j} - \sum_{\substack{i=1 \\ i \neq j}}^{m+n} \tilde{x}_{j,i} &= \tilde{b}_j, \quad j = m+1, m+2, \dots, m+n. \\ \text{and } \tilde{x}_{ij} &\geq \tilde{0}, i = 1, 2, \dots, m+n, j \neq i. \end{aligned}$$

Thus the transshipment difficulty may be written as

$$\begin{aligned} \text{Minimize } \tilde{Z} &\approx \sum_{i=1}^m \sum_{\substack{j=1 \\ j \neq i}}^{m+n} \tilde{c}_{ij} \tilde{x}_{ij} \\ \text{subject to } \sum_{\substack{j=1 \\ j \neq i}}^{m+n} \tilde{x}_{ij} + \sum_{\substack{j=1 \\ j \neq i}}^{m+n} \tilde{x}_{ji} &\approx \tilde{a}_i, \quad i = 1, 2, 3, \dots, m \\ \sum_{\substack{i=1 \\ i \neq j}}^{m+n} \tilde{x}_{ij} + \sum_{\substack{i=1 \\ i \neq j}}^{m+n} \tilde{x}_{ji} &\approx \tilde{b}_j, \quad j = m+1, m+2, \dots, m+n \\ \sum_{i=1}^m \tilde{a}_i &\approx \sum_{j=m+1}^{m+n} \tilde{b}_j, \quad \text{and } \tilde{x}_{ij} \geq \tilde{0}, i = 1, 2, \dots, m+n, j \neq i. \end{aligned}$$

The beyond formulation is a linear indocrination trouble, which is similar to but not identical to a transportation problem since x_{ij} 's coef-

coefficients are -1. The problem, on the other hand, can be easily transformed into a normal transportation problem.

$$t_i = \sum_{\substack{j=1 \\ j \neq i}}^{m+n} \tilde{x}_{ij}, \quad i = 1, 2, 3, \dots, m, \quad t_i + \tilde{x}_{ii} = \sum_{j=1}^{m+n} \tilde{x}_{ji}, \quad i = 1, 2, 3, \dots, m$$

$$\text{and } t_j = \sum_{\substack{i=1 \\ i \neq j}}^{m+n} \tilde{x}_{ji}, \quad j = m+1, m+2, m+3, \dots, m+n$$

$$t_i + \tilde{x}_{ij} = \sum_{i=1}^{m+n} \tilde{x}_{ji}, \quad j = m+1, m+2, m+3, \dots, m+n$$

where t_i represents the total amount of transshipment through the i^{th} origin and t_j represents the total amount shipped put from the j^{th} destination as transshipment. Let $T > 0$ be sufficiently large number so that $t_i \leq T$, for all i and $t_j \leq T$ for all j . We now write $t_i + x_{ii} = T$ then the non negative slack variable \tilde{x}_{ii} represents the difference between T and the actual amount of transshipment through the i^{th} origin. Similarly, if we let $t_j + x_{jj} = T$ then the non negative slack variable \tilde{x}_{jj} represents the difference between T and the actual amount of transshipment through the j^{th} destination. It is self-evident that the total volume of products transshipped at any given time cannot exceed the total amount generated or obtained . and hence $T = \sum_{i=1}^m \tilde{a}_i$.

The transshipment problem then reduces to

$$\text{Minimize } \tilde{Z} \approx \sum_{i=1}^{m+n} \sum_{\substack{j=1 \\ j \neq i}}^{m+n} \tilde{c}_{ij} \tilde{x}_{ij}$$

$$\text{subject to } \sum_{j=1}^{m+n} \tilde{x}_{ij} = \tilde{a}_i + T, \quad i = 1, 2, 3, \dots, m$$

$$\sum_{j=1}^{m+n} \tilde{x}_{ij} = T, \quad i = m+1, m+2, m+3, \dots, m+n,$$

$$\sum_{i=1}^{m+n} \tilde{x}_{ij} = T, \quad j = 1, 2, \dots, m$$

$$\sum_{i=1}^{m+n} \tilde{x}_{ij} = \tilde{b}_j + T, \quad j = m+1, m+2, m+3, \dots, m+n$$

$$\text{and } \tilde{x}_{ij} \geq \tilde{0}, i = 1, 2, \dots, m+n, j = 1, 2, \dots, m+n.$$

$$\text{where } \tilde{C}_{ii} = 0, i = 1, 2, \dots, m+n.$$

A typical transportation problem with $(m+n)$ roots and $(m+n)$ destinations is described by the mathematical model above. The problem's solution includes $2m+2n-1$ basic variables. However, the remaining buff-

er stock is represented by $m+n$ of these variables appearing in the diagonal cells if they are omitted. Our interest has $m+n-1$ basic variables.

4. Improved Interval Vogel's Approximation algorithm

Row opportunity interval cost matrix

Each element of the same row is subtracted from the row's smallest interval expense.

Column opportunity interval cost matrix

The smallest cost of each column in the original interval transshipment cost matrix is subtracted from each segment of the same column. (Note: the diagonal zero is ignored when determining the smallest interval cost.) On the TOC matrix, the proposed algorithm is used. The following are the steps in detail:

Step 1: If either (full amount supply > full amount demand) or (total supply total demand), balance the given interval transshipment problem.

Step 2: Convert the interval transshipment problem in to interval transportation problem.

Step 3: state the all interval parameters supply, demand and unit transportation cost in terms of midpoint and half width form.

Step 4: Obtain the ROC, COC and then TOC matrix

Step 5: Identify the boxes in each row with the lowest and next-lowest interval transportation costs, and write the difference (penalty) along the table's side against the corresponding row.

Step 6: In each column, find the boxes with the lowest and next-lowest interval transportation costs, and write the difference (penalty) against the corresponding column.

Step 7: Determine the most severe penalty possible. If it's on the table's side, give the most space to the box with the lowest transportation interval cost in that row. If the box with the lowest interval cost of transportation in that column is below the table, give it the most money.

Step 8: Select the topmost row and the farthest left column if the penalties for two or more rows or columns are equal.

Step 9: The satisfied row or column does not need any further thought. If both the row and column are fulfilled at the same time, delete only one of them, and assign zero supply to the remaining row or column (or demand).

Step 10: Continue until all of the rows and columns are satisfied.

Step 11: Using the original balanced interval transshipment cost matrix, calculate the cumulative interval carrying cost for the possible allocations.

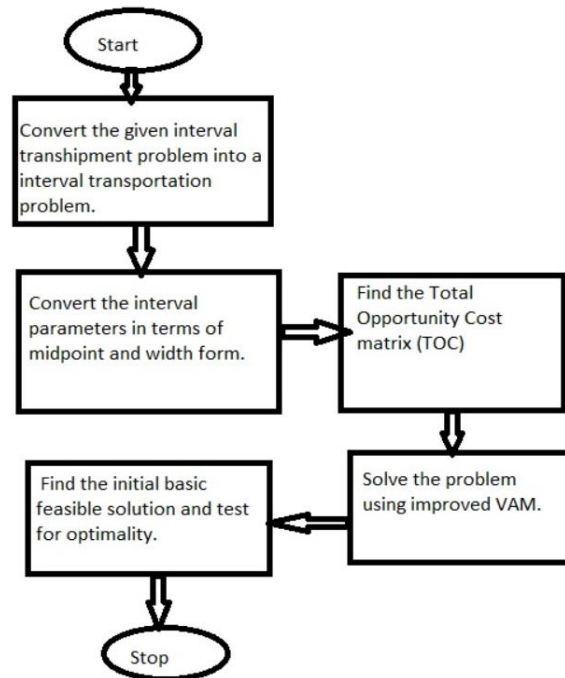


Figure 2

5. NUMERICAL EXAMPLE

To solve the following fully interval integer transshipment problem with 2 source and 2 destinations. The supply values of the source S1 and S2 are [4,6], [5,7] units respectively. The demand values of destination D1 and D2 are [3,5], [6,8] units respectively. Transportation cost per unit between various defined source and destination are given in the following table [12].

Table 2. Interval transshipment problem

	S1	S2	D1	D2	Supply
S1	-	[1,3]	[6,9]	[2,4]	[4,6]
S2	[1,3]	-	[8,10]	[2,4]	[5,7]
D1	[0,1]	[7,9]	-	[3,5]	-
D2	[1,3]	[0,2]	[1,3]	-	-
Demand	-	-	[3,5]	[6,8]	

Table 3. Equivalent interval transportation problem of the interval transshipment problem

	S1	S2	D1	D2	Supply
S1	[0,0]	[1,3]	[6,9]	[2,4]	[13,19]
S2	[1,3]	[0,0]	[8,10]	[2,4]	[14,20]
D1	[0,1]	[7,9]	[0,0]	[3,5]	[9,13]
D2	[1,3]	[0,2]	[1,3]	[0,0]	[9,13]
Demand	[9,13]	[9,13]	[12,18]	[15,21]	[45,65]

By applying the proposed algorithm, the initial basic feasible solution is

Table 4. Initial basic feasible solution

	S1	S2	D1	D2	Supply
S1	<0,0> <11,2>	<1,1>	<11,1.5>	<1,1> <5,3>	<16,3>
S2	<1.5,1>	<0,0> <11,2>	<14,1>	<1,1> <6,3>	<17,3>
D1	<0,0.5>	<14.5,1>	<0,0> <11,2>	<4.5,1>	<11,2>
D2	<2.5,1>	<0,1>	<1,1> <4,3>	<0,0> <7,3>	<11,2>
Demand	<11,2>	<11,2>	<15,3>	<18,3>	<55,3>

The route from source to destination is given as S1 to D2 = <5,3>, S2 to D2 = <6,3>, D2 to D1 = <4,3>

Then the initial solution of the given interval transshipment problem is

$$\text{Min } z = \langle 3,1 \rangle \langle 5,3 \rangle + \langle 3,1 \rangle \langle 6,3 \rangle + \langle 2,1 \rangle \langle 4,3 \rangle = \langle 41,3 \rangle = [38,44]$$

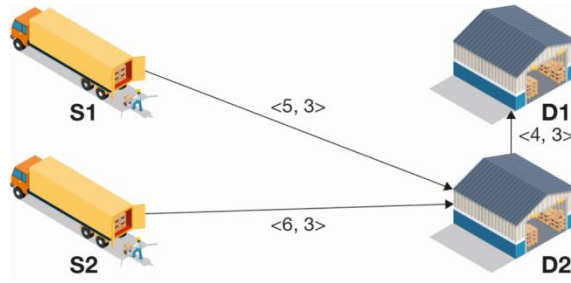


Figure 3

Table 5. Comparison of solutions

S.No	Solution by the proposed method	Solution by P.Rajendran and P.Pandian [12]	Solution by R.Sophiaporchelvi and M.Anitha[14]
1	[38,44]	[20,67]	[1,126]

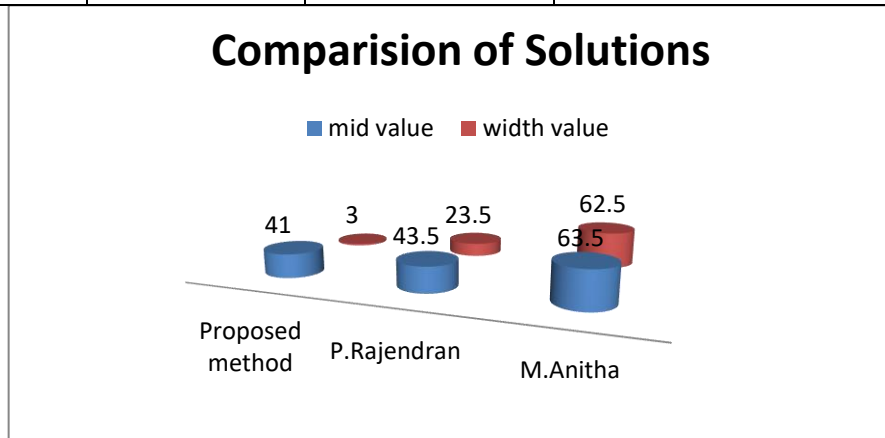


Figure 4

6. CONCLUSION

To discover the primary interval possible explanation to the interval shipment problem, we proposed an improved interval Vogel's approximation algorithm. The interval possible solution's optimality has been checked. A mathematical model is solved to demonstrate the suggested technique, and the outcome is compared to the results of other obtainable methods.

References

- [1]. Akilbasha, P.Pandian and G.Natarajan, An innovative exact method for solving fully interval integer transportation problems, *Informatics in Medicine Unlocked*, 11, 95-99,(2018).
- [2]. K. Ganesan and P. Veeramani, On Arithmetic Operations of Interval Numbers, *International Journal of Uncertainty, Fuzziness and Knowledge - Based Systems*, 13 (6), 619 – 631,(2005).
- [3]. V.Hurt and T.Tramel, Alternative formulations of the transshipment problem. *Journal of Farm Economics*, Vol.47 ,763–773,(1965).
- [4]. Ming Ma, Menahem Friedman and Abraham Kandel. A new fuzzy arithmetic. *Fuzzy sets and systems*, 108: 83-90,(1999).
- [5]. A.NagoorGani, R.Baskaran, and S.N.Mohamed Assarudeen, Transshipment problem in fuzzy environment, *International J. of Math. Sci. &Engg. Applns.*, 5, 57-74,(2011).
- [6]. P. Pandian and G. Natrajan, A new method for finding an optimal solution of fully interval integer transportation problems, *applied mathematical sciences*. *Applied* 4(37), 1819 – 1830(2010).
- [7]. Rajendran, P., Pandian, P., Solving fully Interval Transshipment problems, *International Mathematical Forum*, 41, 2027-2035(2012).
- [8]. A. Sengupta and T.K. Pal, Interval-valued transportation problem with multiple penalty factors, *VU Journal of Physical Sciences*, 9,71 – 81,(2003).
- [9]. Sophia Porchelvi and M. Anitha, Comparative study of optimum solution between interval transportation and interval transshipment problem, *International Journal of Advanced Science and Engineering*, 4 (4),764-767,(2018).
- [10]. A.Orden, Transshipment problem, *Management Science*, 3, 276–285(1956)