



HAL
open science

Factorially Many Maximum Matchings Close to the Erdős-Gallai Bound

Stéphane Bessy, Johannes Pardey, Lucas Picasarri-Arrieta, Dieter Rautenbach

► **To cite this version:**

Stéphane Bessy, Johannes Pardey, Lucas Picasarri-Arrieta, Dieter Rautenbach. Factorially Many Maximum Matchings Close to the Erdős-Gallai Bound. *The Electronic Journal of Combinatorics*, 2022, 29 (2), 10.37236/10610 . hal-03701284

HAL Id: hal-03701284

<https://inria.hal.science/hal-03701284v1>

Submitted on 21 Jun 2022

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.



Distributed under a Creative Commons Attribution - NoDerivatives 4.0 International License

Factorially many maximum matchings close to the Erdős-Gallai bound*

Stéphane Bessy¹ Johannes Pardey² Lucas Picasarri-Arrieta³
Dieter Rautenbach²

¹ LIRMM, Univ Montpellier, CNRS, Montpellier, France
stephane.bessy@lirmm.fr

² Institute of Optimization and Operations Research, Ulm University, Ulm, Germany
{johannes.pardey,dieter.rautenbach}@uni-ulm.de

³ Université Côte d'Azur, Inria, CNRS, I3S, France
lucas.picasarri-arrieta@inria.fr

Abstract

A classical result of Erdős and Gallai determines the maximum size $m(n, \nu)$ of a graph G of order n and matching number νn . We show that G has factorially many maximum matchings provided that its size is sufficiently close to $m(n, \nu)$.

Keywords: Matching

1 Introduction

We consider finite, simple, and undirected graphs. A *matching* in a graph G is a set of pairwise disjoint edges, and the *matching number* $\nu(G)$ of G is the largest size of a matching in G . For a matching M in G , let $V(M)$ be the set of vertices of G that are incident with an edge in M ; the set $V(M)$ contains the vertices of G that are *saturated* by M .

A classical result of Erdős and Gallai, Theorem 4.1 in [3], states that a graph G of order n , size m , and matching number $\nu(G)$ such that $\nu(G) = \nu n$ for some $\nu \in [0, \frac{1}{2}]$ satisfies

$$m \leq m(n, \nu) := \begin{cases} \nu n(n - \nu n) + \binom{\nu n}{2} & , \text{ if } \nu \leq \frac{2}{5} - \frac{3}{5n}, \text{ and} \\ \binom{2\nu n + 1}{2} & , \text{ if } \frac{2}{5} - \frac{3}{5n} \leq \nu \leq \frac{1}{2}. \end{cases} \quad (1)$$

Furthermore, they showed that equality holds in (1) if and only if

- (i) the complement \overline{G} of G is $K_{n-\nu n} \cup \overline{K_{\nu n}}$ for $\nu \leq \frac{2}{5} - \frac{3}{5n}$, and
- (ii) G is $K_{2\nu n + 1} \cup \overline{K_{n-2\nu n - 1}}$ for $\frac{2}{5} - \frac{3}{5n} \leq \nu \leq \frac{1}{2}$.

Recall that, for positive integers n and k with $k \leq n$, the *falling factorial* $n^{\underline{k}}$ is $n(n-1)\dots(n-k+1)$.

The starting point here was the observation that the two extremal graphs in (i) and (ii) have

$$(n - \nu n)^{\nu n} \quad \text{and} \quad \frac{(2\nu n + 1)!}{(\nu n)! 2^{\nu n}}$$

*Research supported by research grant DIGRAPHS ANR-19-CE48-0013.

maximum matchings, respectively. Estimating quite roughly, it follows that, for positive ν , the extremal graphs for (1) have between $\lceil 0.4n \rceil^{\lceil 0.5\nu n \rceil}$ and $n^{\nu n}$ maximum matchings. We show that G still has factorially many maximum matchings provided that $m(n, \nu) - m$ is sufficiently small. Since $m(n, \nu) = \Theta(\nu n^2)$, it is natural to bound $m(n, \nu) - m$ in terms of ν and n^2 .

The following is our first main result; all proofs are given in the next section.

Theorem 1. *For every real ν with $\nu \in (0, \frac{1}{2}]$, the following holds: If G is a graph of order n , size m , and matching number νn such that $n \geq (\frac{50}{\nu})^2$ and $m \geq m(n, \nu) - (\frac{\nu}{50})^2 n^2$, then G has at least $\lceil 0.1n \rceil^{\lceil 0.1\nu n \rceil}$ maximum matchings.*

For the sake of simplicity, we did not try to optimize the constants that appear in this statement, which works over the full range $(0, \frac{1}{2}]$ of ν . Our purpose here is rather to illustrate the effect and present arguments and tools that allow to capture it. In particular, the exact dependence of the minimum number of maximum matchings on the difference $m(n, \nu) - m$ remains a natural yet challenging open problem.

Our second main result gives a better bound provided that ν is sufficiently small.

Theorem 2. *There are two functions $h_\nu : (0, 1) \rightarrow (0, \frac{1}{2}]$ and $h_\delta : (0, 1) \times (0, \frac{1}{2}] \rightarrow (0, 1)$ with the following property: If $\epsilon \in (0, 1)$, $\nu \in (0, h_\nu(\epsilon))$, and G is a graph of order n , size m , and matching number νn such that $n \geq 1/h_\delta(\epsilon, \nu)$ and $m \geq m(n, \nu) - h_\delta(\epsilon, \nu)n^2$, then G has at least $\lceil (1 - \epsilon)n \rceil^{\lceil (1 - \epsilon)\nu n \rceil}$ maximum matchings.*

Matchings in graphs are among the most well studied topics in graph theory [7], and we would like to mention only few related results. Computing the permanent of a matrix, and, hence, counting the perfect matchings of a given bipartite graph, is a well known #P-complete problem [9]. Van der Waerden's proved conjecture on the permanent of a doubly stochastic matrix [2, 5, 6, 8] allows to show that d -regular bipartite graphs have exponentially many perfect matchings for $d \geq 3$, and Brègman's [1] upper bound on the permanent allows to derive an exponential upper bound. Another famous related result, establishing a conjecture of Lovász and Plummer, is due to Esperet, Kardoš, King, Král, and Norine [4] who showed that cubic bridgeless graphs have exponentially many perfect matchings.

2 Auxiliary results and proofs

Throughout this section, let G be a graph of order n , size m , and matching number νn .

A key tool for our approach is the *Gallai-Edmonds decomposition* $D \cup A \cup C$ of a graph G , cf. [7], where

$$D = \{u \in V(G) : \nu(G - u) = \nu(G)\}, \quad A = \bigcup_{u \in D} N_G(u) \setminus D, \quad \text{and} \quad C = V(G) \setminus (D \cup A),$$

that is, the set D contains the vertices that are not saturated by some maximum matching in G , the set A is the set of neighbors of the vertices in D outside of D , and C contains the remaining vertices. Let the components of $G[D]$ be G_1, \dots, G_k . Each G_i is *factor-critical*, that is, for every vertex u of G_i , the graph $G_i - u$ has a perfect matching. Let G_i have order n_i for $i \in [k]$, $d = |D|$, $a = |A|$, and $c = |C|$.

Every maximum matching of G consists of, cf. [7],

- a matching of size $(n_i - 1)/2$ in G_i for every $i \in [k]$,

- a matching of size a in the bipartite subgraph of G with the partite sets A and D formed by the edges between these two sets, and
- a perfect matching in $G[C]$.

In particular, such matchings are guaranteed to exist. Note that

$$n - 2\nu n = k - a \tag{2}$$

is the number of vertices that are not saturated by maximum matchings in G .

If G^* arises from G by adding all missing edges

- within $V(G_i)$ for each $i \in [k]$,
- between D and A , and
- within $A \cup C$,

then G^* has the same Gallai-Edmonds decomposition, and, hence, also the same matching number as G , and

$$m(G^*) = \sum_{i=1}^k \binom{n_i}{2} + da + \binom{n-d}{2}. \tag{3}$$

Using $d = n_1 + \dots + n_k$, $n_1, \dots, n_k \geq 1$, and the convexity of $x \mapsto x^2$, we obtain

$$m(G^*) \leq m^* := \binom{d-k+1}{2} + da + \binom{n-d}{2}. \tag{4}$$

Since (4) holds with equality if each but at most one component of G^* is an isolated vertex, the integer m^* is the size of a graph of order n and matching number νn , and (1) implies $m(n, \nu) \geq m^*$.

We introduce the useful variables x and y :

$$x = \nu - \frac{a}{n} \quad \text{and} \quad y = \frac{d-k}{n}. \tag{5}$$

Since $0 \leq a \leq \nu n$, we have $x \in [0, \nu]$.

Note that

$$k \stackrel{(2)}{=} n - 2\nu n + a = n - 2\nu n + \nu n - xn = (1 - \nu - x)n, \tag{6}$$

and, hence,

$$y = \frac{d-k}{n} \leq \frac{n-a-k}{n} = 1 - \nu + x - 1 + \nu + x = 2x,$$

that is, $y \in [0, 2x]$.

We provide some intuition for x and y . The variable x measures which proportion of every maximum matching consists of edges between A and D . If $x = 0$, then $a = \nu n$, which means that every maximum matching lies completely between A and D . If $x = \nu$, then $a = 0$, which means that every maximum matching lies completely within $D \cup C$. The variable y relates to the orders of the components of $G[D]$. Each of these k components contains at least one vertex, that is, $d - k$ is the number of additional vertices of G distributed among the components of $G[D]$ on top of this one vertex. Therefore, y measures the proportion of the vertices of G that fill up the k components of

$G[D]$. If $y = 0$, then all components of $G[D]$ have just one vertex. If $y = 2x$, then $n = a + d$, which means that C is empty.

Our first lemma expresses the quadratic part of $m(n, \nu) - m^*$ in terms of x and y .

Lemma 3. $m(n, \nu) - m^* \geq g(x, y)n^2 - n$ for

$$g(x, y) := \begin{cases} x \left(1 - \nu - \frac{3}{2}x\right) + y(2x - y) & , \text{ if } \nu \leq \frac{2}{5} - \frac{3}{5n} \text{ and} \\ (\nu - x) \left(\frac{5}{2}\nu + \frac{3}{2}x - 1\right) + y(2x - y) & , \text{ if } \frac{2}{5} - \frac{3}{5n} < \nu \leq \frac{1}{2}. \end{cases}$$

Proof. If $\nu \leq \frac{2}{5} - \frac{3}{5n}$, then

$$\begin{aligned} & m(n, \nu) - m^* \\ \stackrel{(1),(4)}{=} & \left(\nu n(n - \nu n) + \binom{\nu n}{2} \right) - \left(\binom{d - k + 1}{2} + da + \binom{n - d}{2} \right) \\ = & \left(\nu n(n - \nu n) + \frac{(\nu n)^2}{2} \right) - \left(\frac{(d - k)^2}{2} + da + \frac{(n - d)^2}{2} \right) - \underbrace{\left(\frac{\nu n}{2} + \frac{d - k}{2} - \frac{n - d}{2} \right)}_{\leq n} \\ \geq & \left(\nu n(n - \nu n) + \frac{(\nu n)^2}{2} \right) - \left(\frac{(d - k)^2}{2} + da + \frac{(n - d)^2}{2} \right) - n \\ \stackrel{(5)}{=} & \left(x \left(1 - \nu - \frac{3}{2}x\right) + y(2x - y) \right) n^2 - n, \end{aligned}$$

where the last equality requires a tedious yet straightforward calculation.

Similarly, if $\frac{2}{5} - \frac{3}{5n} < \nu \leq \frac{1}{2}$, then

$$\begin{aligned} & m(n, \nu) - m^* \\ \stackrel{(1),(4)}{=} & \binom{2\nu n + 1}{2} - \left(\binom{d - k + 1}{2} + da + \binom{n - d}{2} \right) \\ = & 2(\nu n)^2 - \left(\frac{(d - k)^2}{2} + da + \frac{(n - d)^2}{2} \right) - \underbrace{\left(-\nu n + \frac{d - k}{2} - \frac{n - d}{2} \right)}_{\leq n} \\ \geq & 2(\nu n)^2 - \left(\frac{(d - k)^2}{2} + da + \frac{(n - d)^2}{2} \right) - n \\ \stackrel{(5)}{=} & \left((\nu - x) \left(\frac{5}{2}\nu + \frac{3}{2}x - 1\right) + y(2x - y) \right) n^2 - n, \end{aligned}$$

where the last equality again requires a tedious yet straightforward calculation. \square

The next lemma allows to identify the values of x and y for which $m(n, \nu) - m^*$ is small. Before we formulate it, we motivate its statement using plots of $g(x, y)$ for different values of ν , cf. Figure 1.

For $\nu = 0.4$, the function $g(x, y)$ is small in three cases:

- (1) Either x , and, hence, also y are small,
- (2) or x is close to ν , and y is small,
- (3) or x is close to ν , and y is close to 2ν .

For smaller values of ν , for instance, $\nu = 0.3$, only Case (1) leads to small values of $g(x, y)$, and for larger values of ν , for instance, $\nu = 0.5$, only Cases (2) and (3) lead to small values of $g(x, y)$. From the

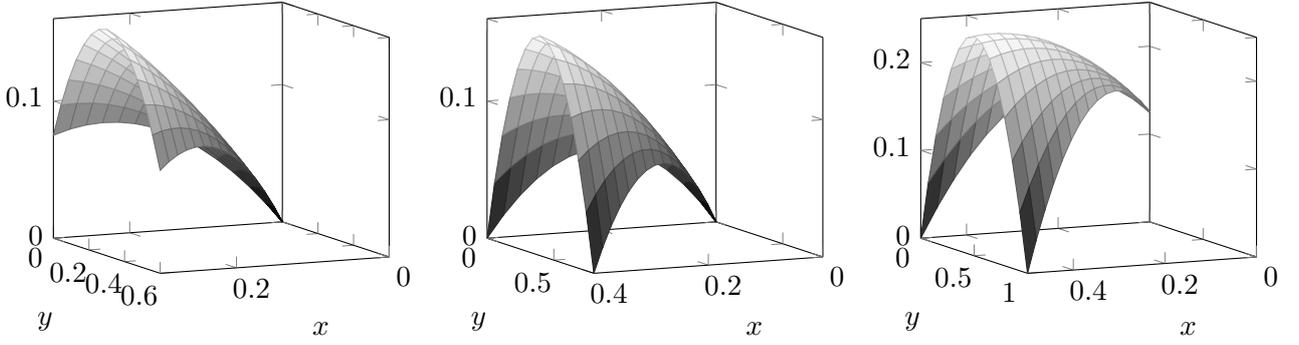


Figure 1: The function $g(x, y)$ for $\nu = 0.3$ (left), $\nu = 0.4$ (middle), and $\nu = 0.5$ (right), assuming that n is large enough.

intuition provided for x and y before Lemma 3, one can derive the structural implications of each case on the Gallai-Edmonds decomposition: In Case (1), the Gallai-Edmonds decomposition degenerates towards the bipartite graph between A and D , which is complete bipartite in G^* . In Case (2), most of G is in C , which is complete in G^* . Finally, in Case (3), most of G is in D , which is the union of k complete graphs in G^* . Within the proof of Theorem 1, we consider exactly these three cases.

Lemma 4. *Let δ be such that $\delta \geq \frac{3}{n}$.*

If

$$\left((x \geq \delta) \vee \left(\nu \geq \frac{2}{5} + \frac{2}{5}\delta \right) \right) \wedge \left((x \leq \nu - \delta) \vee \left(\nu \leq \frac{2}{5} - \frac{2}{5}\delta \right) \vee (\delta \leq y \leq 2\nu - 3\delta) \right), \quad (7)$$

then $g(x, y) \geq \delta^2$.

Proof. First, we assume that $\nu \leq \frac{2}{5} - \frac{3}{5n}$, which implies $x \geq \delta$.

If $x \leq \nu - \delta$, then

$$g(x, y) = x \left(1 - \nu - \frac{3}{2}x \right) + \underbrace{y(2x - y)}_{\geq 0} \geq \underbrace{x}_{\geq \delta} \underbrace{\left(1 - \nu - \frac{3}{2}x \right)}_{\geq \frac{3}{2}\delta} \geq \delta^2,$$

if $x > \nu - \delta$ and $\nu \leq \frac{2}{5} - \frac{2}{5}\delta$, then

$$g(x, y) = x \left(1 - \nu - \frac{3}{2}x \right) + \underbrace{y(2x - y)}_{\geq 0} \geq \underbrace{x}_{\geq \delta} \underbrace{\left(1 - \nu - \frac{3}{2}x \right)}_{\geq \delta} \geq \delta^2,$$

and, if $x > \nu - \delta$ and $\nu > \frac{2}{5} - \frac{2}{5}\delta$, then

$$g(x, y) = \underbrace{x \left(1 - \nu - \frac{3}{2}x \right)}_{\geq 0} + y(2x - y) \geq \underbrace{y}_{\geq \delta} \underbrace{(2x - y)}_{\geq \delta} \geq \delta^2.$$

Next, we assume that $\frac{2}{5} - \frac{3}{5n} < \nu \leq \frac{1}{2}$, which implies $\nu > \frac{2}{5} - \frac{2}{5}\delta$.

Note that

$$\frac{5}{2}\nu + \frac{3}{2}x - 1 \geq \min \left\{ \frac{5}{2} \left(\frac{2}{5} - \frac{3}{5n} \right) + \frac{3}{2}\delta - 1, \frac{5}{2} \left(\frac{2}{5} + \frac{2}{5}\delta \right) - 1 \right\} \geq \delta.$$

If $x \leq \nu - \delta$, then

$$g(x, y) = (\nu - x) \left(\frac{5}{2}\nu + \frac{3}{2}x - 1 \right) + \underbrace{y(2x - y)}_{\geq 0} \geq \underbrace{(\nu - x)}_{\geq \delta} \underbrace{\left(\frac{5}{2}\nu + \frac{3}{2}x - 1 \right)}_{\geq \delta} \geq \delta^2,$$

and, if $x > \nu - \delta$, then

$$g(x, y) = \underbrace{(\nu - x) \left(\frac{5}{2}\nu + \frac{3}{2}x - 1 \right)}_{\geq 0} + y(2x - y) \geq \underbrace{y}_{\geq \delta} \underbrace{(2x - y)}_{\geq \delta} \geq \delta^2.$$

□

The next two lemmas establish the existence of many maximum matchings in graphs that are close to complete bipartite graphs or complete graphs, respectively.

Lemma 5. *For positive integers k and a with $k > 1.1a$, let the bipartite graph H arise from $K_{a,k}$ with partite sets A and K of orders a and k , respectively, by removing up to $0.08a^2$ edges in such a way that H has a matching saturating all vertices in A .*

The graph H has at least $\lceil k - 1.1a \rceil^{\min\{\lceil 0.2a \rceil, \lceil k - 1.1a \rceil\}}$ matchings saturating all vertices in A .

Proof. Let M be a matching in H saturating all vertices in A . If A contains a subset A' of at least $0.2a$ vertices u with $|N_H(u) \setminus V(M)| \geq k - 1.1a$, then there are at least $\lceil k - 1.1a \rceil^{\min\{\lceil 0.2a \rceil, \lceil k - 1.1a \rceil\}}$ matchings that connect the vertices in A' to vertices in $K \setminus V(M)$. Since each of these matchings can be extended to a matching saturating all vertices in A by using edges from M , the desired statement follows. Hence, for a contradiction, we suppose that a set A' as above does not exist. Since H has at most a^2 edges within $V(M)$, this implies that

$$m(H) < a^2 + 0.2a(k - a) + 0.8a(k - a - 0.1a) = ak - 0.08a^2,$$

which is a contradiction. □

Lemma 6. *For a positive integer p with $p \geq 10^3$, let the graph K arise from K_{2p} by removing up to $0.01\binom{p}{2}$ many edges in such a way that K has a perfect matching.*

The graph K has at least $\lceil 0.447p \rceil!$ perfect matchings.

Proof. Let $M = \{x_1y_1, \dots, x_py_p\}$ be a perfect matching in K . Let the graph H with vertex set $\{z_1, \dots, z_p\}$ be such that, for distinct indices i and j from $[p]$, the two vertices z_i and z_j are adjacent in H if and only if the two edges x_iy_j and x_jy_i belong to K . Since every non-edge in K leads to at most one non-edge in H , we have

$$m(H) \geq \binom{p}{2} - 0.01\binom{p}{2} \geq 0.99\binom{p}{2} \stackrel{p \geq 10^3}{\geq} 0.989\frac{p^2}{2},$$

which implies that H has average degree at least $0.989p$. This implies that H contains at least $0.895p$ vertices of degree at least $0.895p$. Hence, the graph H contains a set X of at least $0.447p$ vertices such that each vertex in X has at least $0.447p$ neighbors outside of X . This immediately implies that H contains at least $\lceil 0.447p \rceil!$ distinct matchings.

For a matching N in H , let $M(N)$ arise from the perfect matching M in K by replacing, for every edge z_iz_j in N , the two edges x_iy_i and x_jy_j in M with the two edges x_iy_j and x_jy_i . Clearly, the set

$M(N)$ is a perfect matching in K , and $M(N)$ is distinct from $M(N')$ for distinct matchings N and N' in H . Therefore, the graph K has at least $\lceil 0.447p \rceil!$ perfect matchings. \square

Proof of Theorem 1. Let $\nu \in (0, \frac{1}{2}]$.

Let

$$\epsilon = \left(\frac{\nu}{50}\right)^2 \leq 10^{-4}. \quad (8)$$

Let G be a graph of order n , size m , and matching number νn such that $\epsilon n \geq 1$ and $m \geq m(n, \nu) - \epsilon n^2$. We have

$$\epsilon n^2 \geq m(n, \nu) - m = \underbrace{(m(n, \nu) - m^*)}_{\geq 0} + \underbrace{(m^* - m(G^*))}_{\geq 0} + \underbrace{(m(G^*) - m)}_{\geq 0}, \quad (9)$$

where we use the notation introduced earlier in this section.

Let $\delta = 2\sqrt{\epsilon}$.

Using (8), $\nu \leq \frac{1}{2}$, and $n \geq \frac{1}{\epsilon}$, we obtain

$$0 < \delta \leq \frac{\nu}{25}, \text{ and } n \geq \max\left\{\frac{3}{\delta}, 10^4\right\}. \quad (10)$$

If (7) holds, then Lemma 3 and Lemma 4 imply

$$\epsilon n^2 \stackrel{(9)}{\geq} m(n, \nu) - m^* \geq g(x, y)n^2 - n \geq \underbrace{\delta^2 n^2}_{\geq 4\epsilon n^2} - \underbrace{n}_{\leq \epsilon n^2} \geq 3\epsilon n^2,$$

which is a contradiction. Hence (7) fails, which leads to the following three cases:

- $x < \delta$ and $\nu < \frac{2}{5} + \frac{2}{5}\delta$.
- $x > \nu - \delta$, $\nu > \frac{2}{5} - \frac{2}{5}\delta$, and $y < \delta$.
- $x > \nu - \delta$, $\nu > \frac{2}{5} - \frac{2}{5}\delta$, and $y > 2\nu - 3\delta$.

Each of these three cases corresponds to a different degeneration of the Gallai-Edmonds decomposition of G . We now consider these cases separately.

Case 1 $x < \delta$ and $\nu < \frac{2}{5} + \frac{2}{5}\delta$.

Note that $\nu < \frac{2}{5} + \frac{2}{5}\delta \stackrel{(10)}{\leq} \frac{2}{5} + \frac{2}{125}\nu$, which implies

$$\nu < 0.407. \quad (11)$$

By (5) and (6), we have

$$\begin{aligned} \nu n \geq a &= (\nu - x)n > (\nu - \delta)n \stackrel{(10)}{\geq} 0.96\nu n, \text{ and} \\ (1 - \nu)n \geq k &\stackrel{(6)}{=} (1 - \nu - x)n > (1 - \nu - \delta)n \stackrel{(10), (11)}{>} 0.576n. \end{aligned} \quad (12)$$

Let H_0 be the bipartite subgraph of G with the partite sets A and D formed by the edges between these two sets. Let M be a matching of size a in H_0 that is a subset of some maximum matching in

G . In particular, for every $i \in [k]$, there is at most one edge in M between A and $V(G_i)$. Let H arise from H_0 by removing, for every component G_i of $G[D]$, all but exactly one vertex in such a way that $V(M) \subseteq V(H)$. The properties of the Gallai-Edmonds decomposition imply that every matching in H that saturates all vertices in A can be extended to a maximum matching in G . By (9), we have $m(G^*) - m \leq \epsilon n^2$, and, hence, the graph H arises from $K_{a,k}$ with partite sets A and K of orders a and k , respectively, by removing up to ϵn^2 edges in such a way that $\nu(G) = a$. Since $k \stackrel{(11),(12)}{>} 1.1a$ and $0.08a^2 \stackrel{(12)}{\geq} 0.08 \cdot 0.96^2 \nu^2 n^2 \stackrel{(8)}{\geq} \epsilon n^2$, Lemma 5 implies that H has at least $\lceil k - 1.1a \rceil^{\min\{[0.2a], [k-1.1a]\}}$ matchings saturating all vertices in A . Since $k - 1.1a \stackrel{(11),(12)}{>} 0.1n \geq 0.1\nu n$ and $0.2a \stackrel{(12)}{>} 0.1\nu n$, the desired statement follows in this case.

Case 2 $x > \nu - \delta$, $\nu > \frac{2}{5} - \frac{2}{5}\delta$, and $y < \delta$.

We have $\nu > \frac{2}{5} - \frac{2}{5}\delta \stackrel{(10)}{\geq} \frac{2}{5} - \frac{2}{125}\nu$, which implies

$$\nu > 0.393. \quad (13)$$

Furthermore,

$$\begin{aligned} |C| &= n - d - a = n - (d - k) - k + (\nu n - a) - \nu n \\ &\stackrel{(5),(6)}{=} (1 - y - 1 + \nu + x + x - \nu)n \\ &= (2x - y)n > (2\nu - 3\delta)n \stackrel{(10)}{\geq} \left(2 - \frac{3}{25}\right)\nu n \stackrel{(13)}{\geq} 0.738n. \end{aligned}$$

With $|C| = 2p$, then $p \geq 0.369n \stackrel{(10)}{\geq} 10^3$, and, hence, $p - 1 \geq 0.368n$. Note that $G[C]$ arises from K_{2p} by removing at most $m(G^*) - m \leq \epsilon n^2$ many edges. Since $0.01 \binom{p}{2} \geq \frac{0.01 \cdot 0.368^2}{2} n^2 \stackrel{(8)}{\geq} \epsilon n^2$, and every perfect matching in $G[C]$ can be extended to a maximum matching in G , Lemma 6 implies that G has at least $\lceil 0.447p \rceil! \geq \lceil 0.1n \rceil!$ maximum matchings, and the desired statement follows in this case.

Case 3 $x > \nu - \delta$, $\nu > \frac{2}{5} - \frac{2}{5}\delta$, and $y > 2\nu - 3\delta$.

Exactly as in Case 2, we obtain (13). Furthermore,

$$|D| \geq |D| - k = d - k \stackrel{(5)}{=} yn > (2\nu - 3\delta)n \stackrel{(10)}{\geq} \left(2 - \frac{3}{25}\right)\nu n \stackrel{(13)}{\geq} 0.738n.$$

Suppose, for a contradiction, that $\max\{n_1, \dots, n_k\} \leq 0.54n$. Using the convexity of $x \mapsto x^2$, we obtain

$$\begin{aligned} \epsilon n^2 &\stackrel{(9)}{\geq} m^* - m(G^*) \stackrel{(3),(4)}{=} \binom{d-k+1}{2} - \sum_{i=1}^k \binom{n_i}{2} > \frac{(d-k)^2}{2} - \frac{1}{2} \sum_{i=1}^k n_i^2 \\ &\geq \frac{1}{2} (0.738n)^2 - \frac{1}{2} \left(\frac{n}{0.54n}\right) (0.54n)^2 \geq 0.002n^2, \end{aligned}$$

contradicting (8). Hence, we may assume that $n_1 \geq 0.54n \stackrel{(10)}{\geq} 10^3$. For $p = \frac{n_1-1}{2}$, we obtain $p \geq 0.269n \stackrel{(10)}{\geq} 10^3$ and $p - 1 \stackrel{(10)}{\geq} 0.268n$. Note that removing any one vertex from G_1 , we obtain a graph K that arises from K_{2p} by removing at most $m(G^*) - m \leq \epsilon n^2$ many edges. Furthermore, by the properties of the Gallai-Edmonds decomposition, the graph K has a perfect matching, and every

perfect matching in K can be extended to a maximum matching in G . Since $0.01 \binom{p}{2} \geq \frac{0.01 \cdot 0.268^2}{2} n^2 \stackrel{(8)}{\geq} \epsilon n^2$, Lemma 6 implies that G has at least $\lceil 0.447p \rceil! \geq \lceil 0.1n \rceil!$ maximum matchings, and the desired statement follows in this case.

This completes the proof. \square

For the proof of Theorem 2, we need the following variant of Lemma 5.

Lemma 7. *Let $\epsilon, \gamma \in (0, 1)$. For positive integers k and a with $\frac{\epsilon k}{2} \geq a$, let the bipartite graph H arise from $K_{a,k}$ with partite sets A and K of orders a and k , respectively, by removing up to γak edges in such a way that H has a matching saturating all vertices in A .*

The graph H has at least $\lceil (1 - \sqrt{\gamma} - \frac{\epsilon}{2}) k \rceil^{\lceil \frac{(1 - \sqrt{\gamma} - \frac{\epsilon}{2}) a}{2} \rceil}$ matchings saturating all vertices in A .

Proof. Let M be a matching in H saturating all vertices in A . If A contains a subset A' of at least $(1 - \sqrt{\gamma}) a$ vertices of degree at least $(1 - \sqrt{\gamma}) k$, then each vertex u in A' satisfies $|N_H(u) \setminus V(M)| \geq (1 - \sqrt{\gamma}) k - a \geq (1 - \sqrt{\gamma} - \frac{\epsilon}{2}) k$, and there are at least $\lceil (1 - \sqrt{\gamma} - \frac{\epsilon}{2}) k \rceil^{\lceil \frac{(1 - \sqrt{\gamma} - \frac{\epsilon}{2}) a}{2} \rceil}$ matchings that connect the vertices in A' to vertices in $K \setminus V(M)$. Since each of these matchings can be extended to a matching saturating all vertices in A by using edges from M , the desired statement follows. Hence, for a contradiction, we suppose that a set A' as above does not exist. This implies that $m(H) < (1 - \sqrt{\gamma}) ak + \sqrt{\gamma} a (1 - \sqrt{\gamma}) k = ak - \gamma ak$, which is a contradiction. \square

Proof of Theorem 2. Let $\epsilon \in (0, 1)$. Let $\nu \in (0, \frac{1}{2}]$ and $\delta \in (0, 1)$ be such that

$$\nu \leq \min \left\{ \frac{1}{5}, \frac{\epsilon}{2 \left(1 + \frac{\epsilon}{2} \left(1 + \frac{\epsilon}{2}\right)\right)} \right\}, \quad \text{and} \quad \delta \leq \frac{\epsilon \nu}{8}. \quad (14)$$

Further restricting ν (in terms of ϵ) and δ (in terms of ϵ and ν), we assume that

$$1 - \left(1 + \frac{\epsilon}{2}\right) \nu \geq 1 - \frac{\epsilon}{2}, \quad \text{and} \quad \left(1 - \sqrt{\frac{\delta}{\left(1 - \frac{\epsilon}{2}\right) \nu \left(1 - \left(1 + \frac{\epsilon}{2}\right) \nu\right)}} - \frac{\epsilon}{2}\right) \left(1 - \frac{\epsilon}{2}\right) \geq 1 - \epsilon. \quad (15)$$

Let G be a graph of order n , size m , and matching number νn such that $\delta n \geq 1$ and $m \geq m(n, \nu) - \delta n^2$. Using the notation introduced before in this section, Lemma 3 implies

$$2\delta \geq \frac{1}{n^2} (m(n, \nu) - m^* + n) \geq g(x, y) \geq x \left(1 - \nu - \frac{3}{2}x\right) \geq x \left(1 - \frac{5}{2}\nu\right) \geq \frac{x}{2}.$$

Hence, $x \leq 4\delta$, and

$$\begin{aligned} \nu n \geq a &= (\nu - x)n \geq (\nu - 4\delta)n \stackrel{(14)}{\geq} \left(1 - \frac{\epsilon}{2}\right) \nu n, \text{ and} \\ k &\stackrel{(6)}{\geq} (1 - \nu - x)n \geq (1 - \nu - 4\delta)n \stackrel{(14)}{\geq} \left(1 - \left(1 + \frac{\epsilon}{2}\right) \nu\right) n. \end{aligned} \quad (16)$$

Using the second upper bound on ν in (14), this implies $\frac{\epsilon}{2} k \geq a$. Let the graph H be defined as in Case 1 of the proof of Theorem 1. For $\gamma = \frac{\delta}{\left(1 - \frac{\epsilon}{2}\right) \nu \left(1 - \left(1 + \frac{\epsilon}{2}\right) \nu\right)}$, we obtain $\gamma ak \stackrel{(16)}{\geq} \delta n^2$, and Lemma 7 implies that H , and hence, also G , has at least $\lceil (1 - \sqrt{\gamma} - \frac{\epsilon}{2}) k \rceil^{\lceil \frac{(1 - \sqrt{\gamma} - \frac{\epsilon}{2}) a}{2} \rceil}$ maximum matchings. Now,

$$\left(1 - \sqrt{\gamma} - \frac{\epsilon}{2}\right) k \stackrel{(16)}{\geq} \left(1 - \sqrt{\gamma} - \frac{\epsilon}{2}\right) \left(1 - \left(1 + \frac{\epsilon}{2}\right) \nu\right) n \stackrel{(15)}{\geq} \left(1 - \sqrt{\gamma} - \frac{\epsilon}{2}\right) \left(1 - \frac{\epsilon}{2}\right) n \stackrel{(15)}{\geq} (1 - \epsilon)n, \text{ and}$$

$$\left(1 - \sqrt{\gamma} - \frac{\epsilon}{2}\right) a \stackrel{(16)}{\geq} \left(1 - \sqrt{\gamma} - \frac{\epsilon}{2}\right) \left(1 - \frac{\epsilon}{2}\right) \nu n \stackrel{(15)}{\geq} (1 - \epsilon)\nu n,$$

which completes the proof. □

References

- [1] L.M. Brègman, Certain properties of nonnegative matrices and their permanents, *Doklady Akademii Nauk SSSR* 211 (1973) 27-30.
- [2] G.P. Egorychev, The solution of van der Waerden's problem for permanents, *Advances in Mathematics* 42 (1981) 299-305.
- [3] P. Erdős and T. Gallai, On maximal paths and circuits of graphs, *Acta Mathematica. Academiae Scientiarum Hungaricae* 10 (1959) 337-356.
- [4] L. Esperet, F. Kardoš, A.D. King, D. Král, and S. Norine, Exponentially many perfect matchings in cubic graphs, *Advances in Mathematics* 227 (2011) 1646-1664.
- [5] D.I. Falikman, Proof of the van der Waerden conjecture on the permanent of a doubly stochastic matrix, *Akademiya Nauk SSSR. Matematicheskie Zametki* 29 (1981) 931-938, 957.
- [6] L. Gurvits, Van der Waerden/Schrijver-Valiant like conjectures and stable (aka hyperbolic) homogeneous polynomials: one theorem for all, *Electronic Journal of Combinatorics* 15 (2008) Paper 66, 26.
- [7] L. Lovász and M.D. Plummer, *Matching theory*, *Annals of Discrete Mathematics*, 29, North-Holland Publishing Co., Amsterdam, 1986.
- [8] A. Schrijver and W.G. Valiant, On lower bounds for permanents, *Koninklijke Nederlandse Akademie van Wetenschappen. Indagationes Mathematicae* 42 (1980) 425-427.
- [9] L.G. Valiant, The complexity of computing the permanent, *Theoretical Computer Science* 8 (1979) 189-201.