

Enforce storage stability (slides)

Jens Gustedt

▶ To cite this version:

Jens Gustedt. Enforce storage stability (slides). 2021. hal-03328555

HAL Id: hal-03328555 https://inria.hal.science/hal-03328555

Preprint submitted on 30 Aug 2021

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers. L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Enforce storage stability

ISO/IEC JTC 1/SC 22/WG14 N2756

Jens Gustedt

INRIA - Camus

ICube - ICPS Université de Strasbourg







https://modernc.gforge.inria.fr/



- Introduction
- Policy
- Proposed terminology
- Normative change





our failure

C has no commonly understood strategy for initialization

- users often don't understand the corner cases
 - component-wise initialization by assignment is common
 - memcmp is the only tool in the C language or library to compare struct, but does it work with padding?
 - atomic_compare_exchange_strong is based on a combination of memcpy and memcmp
- implementations have different strategies for incomplete initialization
 - lazy (just use what's there)
 - freeze (produce a value and stick to it)
 - wobbly (unreliably produce a value)
 - crash
 - not even consistent for the same platform, e.g depending on the size of the storage instance

3/15

Let's do our homework

Regardless what strategy we decide (or not)

- we should provide reasonable definitions
 - for valid use cases of uninitialized memory (if any)
 - ensure that all C features are consistent
 - memcmp
 - atomic_compare_exchange_strong
 - provide clear guidelines for implementations that handle the common cases



Reasonable expectations

A store operation conveys an intent.

A use of an object after a store operation is intentional and should not be undefined per se.

Expect an object to be usable

- whenever an Ivalue is assigned to, even partially
- whenever a byte is written into storage

For a usable object expect

multiple evaluations lead to the same value



Systematic pitfalls

Completely uninitialized objects

- access may crash
- evaluation may unbehave

Large storage instances

- may be implemented page-wise
- uninitialized pages may unbehave



6/15

- Introduction
- Policy
- Proposed terminology
- Normative change





Proposed design principles

observation stability

Every observable change to the abstract state shall be either induced by a store operation, by an IO operation or an external events.

modification locality

A valid store operation shall have no observable effect to storage outside the storage instance that is operand to the operation.

storage stability

Consecutive valid calls to **memcpy** and **memcmp** without intermediate store shall be consistent.





- Introduction
- Policy
- Proposed terminology
- Normative change





Value initialization (semantic initialization)

An object is value-initialized if any of the following holds:

- It stems from an object definition with an explicit or implicit initializer.
- It is a function parameter or a value capture.
- It (or an object that contains it) has been the target of an assignment operation.
- It has an aggregate type and all its elements or members have been value-initialized.
- It is not an array type and a pointer to it has been argument to a call of one of the library functions that store into their pointed-to parameter.



Byte initialization (storage initialization)

It is byte-initialized if it is not atomic and if any of the following holds:

- Its storage instance is allocated with the calloc library function.
- It has scalar, structure or union type and at least one representation byte has a stored value.
- It has structure or union type and at least one of its members has been value-initialized.
- It is a member of a byte-initialized structure or union.
- It is the representation byte of a value-initialized or byte-initialized object.



Initialization

An object is initialized if it is value-initialized or byte-initialized, otherwise it is uninitialized.

Notes

- A structure can be initialized by writing one byte into its representation.
- An array with more than one element cannot be initialized by writing one byte into its representation.

Rationale

- Structures have always been regarded as a whole.
 - They have a compile time fixed size.
- Arrays have no value and size information is futile.
- Representations of padding bits and bytes shall be stable.

Jens Gustedt (INRIA) Enforce storage stability August 30, 2021

12/15

Stable storage instance

A storage instance ... is stable ... if at least one of the following conditions holds ...

- As a whole it represents an object that has been initialized.
- At least one representation byte has been initialized and the macro

```
___STDC_STORAGE_GRANULARITY___
```

- is undefined or expands to zero, or,
- is strictly positive and the size of the storage instance is smaller or equal.



August 30, 2021 13

- Introduction
- Policy
- Proposed terminology
- Normative change





Normative change

Stable storage instances represent stable values ...

As long as a stable storage instance is not the target of any further write operation, all read operations of a given Ivalue that is represented by the storage instance result in the same value ...

and values are consistent

and reads to different Ivalues that have overlapping representations by the storage instance result in consistent values.



August 30, 2021