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# Comparison of coupled solvers for FEM/BEM linear systems arising from discretization of aeroacoustic problems 

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## Résumé

When discretization of an aeroacoustic physical model is based on the application of both the Finite Elements Method (FEM) and the Boundary Elements Method (BEM), this leads to coupled FEM/BEM linear systems combining sparse and dense parts. In this work, we propose and compare a set of implementation schemes relying on the coupling of the open-source sparse direct solver MUMPS with the proprietary direct solvers from Airbus Central R\&T, i.e. the scalapack-like dense solver SPIDO and the hierarchical $\mathcal{H}$-matrix compressed solver HMAT. For this preliminary study, we limit ourselves to a single 24-core computational node.

Mots-clés : aeroacoustics, finite elements, boundary elements, solver comparison, coupled FEM/BEM linear systems

## 1. Introduction

In the industrial context of Airbus, we are interested in models of aeroacoustic phenomena such as the propagation of acoustic waves emitted by an aircraft on take-off, landing and taxiing (see Figure 1. left). Such physical models are typically expressed using partial differential equations. Prior to computing the numerical model, an approximation of its original physical expression is made over a limited domain using a suitable discretization technique.


FIGURE 1 - An acoustic wave (blue arrow) emitted by the aircraft's engine, reflected on the wing and crossing the jet of exhaust gas. Real-life case (left) [22] and a numerical model example (right). The red 2D mesh represents the aircraft's surface and the surface of the green 3D cylinder volume mesh representing the exhaust flow.

Aeroacoustic waves may interfere with various media. In our model, we take into account the air around the aircraft, considered as a homogeneous medium, and jet exhaust flow, considered as heterogeneous. The discretization relies on a coupling of two different techniques, the Finite Elements Method (FEM) [10, 14, 19, 26] applied to the exhaust flow and the Boundary Elements Method (BEM) [9, 20, 23] applied to the surfaces (see Figure 1, right). This leads to a coupled FEM/BEM linear system having two groups of unknowns, $x_{s}$ associated to a surface mesh $s$ resulting from BEM and $x_{v}$ associated to a volume mesh $v$ resulting from FEM (see Equation 11. Ultimately, we target systems counting around $10^{9}$ FEM-related unknowns and $10^{6}$ BEM-related unknowns. Note that, $R_{1}$ and $R_{2}$ denote respectively the first and the second row of the linear system.

Here, $\mathcal{A}$ is a $2 \times 2$ symmetric coefficient matrix where $A_{v v}$ is a large sparse submatrix representing the action of the volume part on itself, $A_{s s}$ is a smaller dense submatrix representing the action of the exterior surface on itself, $A_{s v}$ is a sparse submatrix representing the action of the volume part on the exterior surface and $A_{v s}$ is the transpose of $A_{s v}$ and represents the action of the exterior surface on the volume part (see Figure 2).
In this work, we focus on a direct method for solving the FEM/BEM linear system in Equation 1 . We distinguish two categories of implementation schemes for this method, single-stage and two-stage. The single-stage schemes rely on a single solver which takes the entire coefficient matrix $A$ and the right-hand side on input and gives directly the final solution of the whole system on output. Ideally, the solver should be aware of the dif-


Figure 2 Internal dimensions and sparsity of $A$ in Equation 1 . ferent sparsity levels of $A$ and capable of applying appropriate sparse or dense operations. The two-stage schemes are based on a coupling of a sparse direct solver and a dense direct solver. Our goal here is to present and compare the two-stage schemes between them. In this case, we consider the coupling of the open-source sparse direct solver MUMPS [7, 12] with either the scalapack-like dense direct solver SPIDO or the hierarchical $\mathcal{H}$-matrix compressed solver HMAT. Both are proprietary solvers from Airbus Central R\&T. However, an open-source sequential implementation of HMAT is available as HMAT-OSS [18, 2].
The document is organized as follows : in Section 2, we explain the direct solution method for the target linear system; in Section 3, we detail the two-stage implementation schemes; in Section 4, we describe our experimental evaluation of the latter and present the comparative results. We conclude in Section 5 .

## 2. Direct solution of FEM/BEM systems

The first step of a direct solution of Equation 1 consists in reducing the problem on boundaries and simplifying the system to solve. Based on its first row $R_{1}$, we express $x_{v}$ as :

$$
\begin{equation*}
x_{v}=A_{v v}^{-1}\left(b_{v}-A_{v s} x_{s}\right) \tag{2}
\end{equation*}
$$

Then, substituting $x_{v}$ in $R_{2}$ by Equation 2 yields in $R_{2}$ a reduced system without $x_{v}$. This represents one step of Gaussian elimination, i.e. subtracting $A_{s v} A_{v v}^{-1}$ times $R_{1}$ from $R_{2}$ :

$$
\underset{\mathrm{R}_{2} \leftarrow \mathrm{R}_{2}-\mathrm{A}_{s v} \mathrm{~A}_{v v}^{-1} \times \mathrm{R}_{1}}{\mathrm{R}_{1}}\left[\begin{array}{cc}
\mathrm{A}_{v v} & A_{v s}  \tag{3}\\
0 & \underbrace{A_{s s}-A_{s v} A_{v v}^{-1} A_{v s}}_{S}
\end{array}\right] \times\left[\begin{array}{l}
\mathrm{x}_{v} \\
\mathrm{x}_{s}
\end{array}\right]=\left[\begin{array}{c}
\mathrm{b}_{v} \\
\mathrm{~b}_{s}-A_{s v} A_{v v}^{-1} \mathrm{~b}_{v}
\end{array}\right] .
$$

The expression $A_{s s}-A_{s v} A_{v v}^{-1} A_{v s}$ that appeared in $R_{2}$ in Equation 3 is often referred to as the Schur complement [24] noted $S$ and associated with the partitioning, $v$ and $s$, of the variables in the system (see Section 17. When solving the problem numerically, we factorize $A_{v v}$ into a product of matrices making the system easier to solve rather than actually performing the inverse $A_{v v}^{-1}$. In summary, we have to compute $S$ and solve the reduced Schur complement system matching $R_{2}$ in Equation 3 after eliminating $x_{v}$ :

$$
\begin{equation*}
S x_{s}=b_{s}-A_{s v} A_{v v}^{-1} b_{v} \tag{4}
\end{equation*}
$$

Once we have computed $x_{s}$, we use its value to determine $x_{v}$ according to Equation 2 .

## 3. Two-stage implementation schemes and related work

The principle of a two-stage implementation scheme is to use a sparse direct solver to compute the Schur complement $S$ (see Equation 3) in the first stage and then, call a dense solver to factorize $S$ and solve the reduced system (see Equation 4) in the second stage. Some sparse direct solvers, such as MUMPS, expose in their Application Program Interface (API) a Schur complement computation function. Provided with $A_{v v}, A_{v s}, A_{s v}$ and $A_{s s}$ they can compute $S$ and yield the result as a dense matrix in $A_{\text {ss }}$. Although the sparse solver supports out-ofcore computation ${ }^{1}$. $S$ still has to be entirely stored in RAM. In terms of memory constraints, this quickly becomes a significant limiting factor in solving large systems. To the best of our knowledge, the approaches proposed in the literature for solving coupled FEM/BEM or more generally coupled sparse and dense linear systems based on direct methods, as considered here, make the assumption that the factors fit in memory [13, 16, 25, 15, 21, 17]. Here, we rely on an algorithm allowing us to store $S$ out-of-core, i.e. on disk, and keep using the community solvers. The idea is to compute the Schur complement $S$ block-wise and load into memory only the block necessary for the current computation step. We propose two variants of such an algorithm, multi-solve and multi-factorization.

### 3.1. Multi-solve

In this approach, we compute the Schur complement $S$ by blocks of columns (see Figure 3). Let us note $A_{s v_{i}}^{\top}$ and $A_{s s_{i}}$ blocks of $n_{c}$ columns of $A_{s v}^{\top}$ and $A_{s s}$ respectively. Then, based on the definition of $S$ in Equation $3, S_{i}$ is a block of $n_{c}$ columns of $S$ defined as :

$$
\begin{equation*}
S_{i}=A_{s s_{i}}-\underbrace{A_{s v} \overbrace{\left(\mathrm{~L}_{v v} L_{v v}^{\top}\right)^{-1} A_{s v_{i}}^{\top}}^{Y}}_{Z} . \tag{5}
\end{equation*}
$$

The multi-solve algorithm (see Algorithm 1 in the appendix) begins by the $L L^{\top}$ factorization of $A_{v v}$ (line 3) into the product of matrices $\mathrm{L}_{v v} \mathrm{~L}_{v v}^{\top}$ where $\mathrm{L}_{v v}$ is a lower triangular matrix and $\mathrm{L}_{v v}^{\top}$ its transpose. Then, we loop (line 4) over the blocks of $A_{s v}^{\top}$ and $A_{s s}$ to compute all the blocks $S_{i}$ of $S$ while overwriting $A_{s s}$. The first step of this computation (line 5 ) is a triangular solve done by a sparse solver for determining $\mathrm{Y}=\left(\mathrm{L}_{v v} \mathrm{~L}_{v v}^{\top}\right)^{-1} A_{s v_{i}}^{\top}$. It is possible to benefit from the sparsity of the right-hand side matrix $A_{s v_{i}}^{\top}$ during the solve operation [6]. However, independently from the sparsity of the input right-hand side, the resulting $Y$ is a dense matrix [12] with only a few zero entries. Also, the block $\mathcal{A}_{s v_{i}}^{\top}$ needs to be stored explicitly (see Figure 3). The number of columns $n_{c}$ in $S_{i}$ depends on $n_{\text {BEM }}$, the count of unknowns associated with the formulation of

1. Using out-of-core feature the solver can store on disk (out of the core memory) the parts of matrices not being actively used for computation, reduce memory consumption and process bigger problems with the same amount of memory.

BEM in the global linear system, and the amount of available RAM. If $n_{c}$ is too small, a great number of solve operations are performed on a small number of right-hand sides inside of the main loop which may degrade the performance. On the other hand, if it is too high, data (such as $Y$ and $Z$ ) may not fit in RAM.

### 3.2. Multi-factorization

In this case, the $A_{s s}$ submatrix receiving the Schur complement $S$ is split into square blocks of equal size (see Figure 4). Let us note $A_{s v_{i}}$ a block of $n_{b}$ rows of $A_{s v}$ and $A_{s v_{j}}^{\top}$ a block of $n_{b}$ columns of $A_{s v}^{\top}$. Then, based on the definition of $S$ in Equation 3 . $S_{i j}$ is a block of $n_{b}$ rows and columns of $S$ such as: $\quad S_{i j}=A_{s s_{i j}}-A_{s v_{i}}\left(L_{v v} \mathrm{U}_{v v}\right)^{-1} A_{s v_{j}}^{\top}$.
In the loop (line 2) of the multi-factorization algorithm (see Algorithm 2 in the appendix), we construct a submatrix $W$ from $A_{v v}, A_{s v_{i}}$ and $A_{s v_{j}}^{\top}$ :

$$
W \leftarrow\left[\begin{array}{cc}
A_{\nu v} & A_{s v_{j}}^{\top}  \tag{7}\\
A_{s v_{i}} & 0
\end{array}\right]
$$

Then, we call the SchurComplement function on $W$ (line 3) to compute the associated block $S_{i j}$ of $S$ relying on the Schur complement computation function provided by the sparse solver, e. g. MUMPS. Here, we decompose $A_{v v}$ using the LU factorization as $W$ is not symmetric except when $\mathfrak{i}=\mathfrak{j}$. This yields a duplicated storage of $A_{v v}, A_{s v_{i}}$ and $A_{s v_{j}}^{\top}$ in $W$ (see Figure 4).


Figure 3 - Computation loop of $S$ in multisolve (see Algorithm 1). $A_{s v_{i}}^{\top}$ is explicitly stored and $Y$ is a dense matrix.


Figure 4 - Computation loop of $S$ in multifactorization (see Algorithm 2). Constructing $W$ requires a duplicated storage of $A_{v v}, A_{s v_{i}}$, $A_{s v_{j}}^{\top}$.
$n_{b}$ represents the count of blocks per row or column of $A_{s s}$. The SchurComplement function operating on $W$ implies a re-factorization of $A_{\nu v}$ in $W$ at each iteration, yet it does not change during the computation. Consequently, the more blocks $A_{s s}$ is split into, the more superfluous factorizations of $A_{v v}$ are performed. In Section 4, we show that too much blocks can substantially degrade the performance of the multi-factorization scheme. Similarly, too small values of $n_{c}$ in the case of multi-solve (see Section 3.1) imply a high number of solve operations. Despite being considerably less expensive than factorizations, a sufficiently high number of solves may make the multi-solve method perform worse than multi-factorization. We highlight the fact that both methods are designed to compute the Schur complement S, then solve the reduced system (see Section [2], but differ only in the way they make use of the available API provided by the community solvers. In the Airbus Central R\&T solver framework, both variants are implemented using MUMPS as sparse solver, and SPIDO or HMAT as dense solver.

## 4. Experimental study

Within this preliminary study, we compare to each other the performance of the two-stage implementation schemes, multi-solve and multi-factorization, for solving the target coupled FEM/BEM linear system in Equation 1. Moreover, for each scheme, we consider both of the available solver couplings, i.e. using MUMPS as sparse solver and either HMAT or SPIDO as dense solver. In this section, we use the notations MUMPS/HMAT and MUMPS/SPIDO respectively to refer to these couplings. Our experiments rely on a short pipe test case (see Figure 6 in the appendix) yielding linear systems close enough to those arising from real life models (see Figure 1) while providing the scientific community with a reproducible example [5]. Note that MUMPS and HMAT both provide an out-of-core feature which we disable in our experiments. Also, these solvers rely on data compression and expose a precision parameter $\epsilon$ set to $10^{-3}$ for all the computations. We design the test case so as to know the expected result in advance. This way, we can determine the relative error of the computed solution. The relative error of the results we obtained was approximately $10^{-3}$, which is a satisfactory accuracy in our context. We do not report further on that in the following.
We conduct our experiments on a single node of miriel on the PlaFRIM [3] high-performance computing platform. A miriel node has a total of 24 processor cores running each at 2.5 GHz and 126 GiB of RAM (see Figure 7 in the appendix). As for the software, we use the Airbus Central R\&T solver test suite test_FEMBEM [5]. Table 1 in the appendix details the version and commit numbers of some of the core packages of our experimental software environment. We consider coupled FEM/BEM linear systems with $N$, the total unknown count, ranging from 250,000 up to $4 \times 10^{6}$. Table 2 in the appendix details the proportions of FEM and BEM unknowns for each value of $N$. In addition, we evaluate multiple configurations for each implementation scheme. Regarding multi-solve, we vary the number of columns $n_{c}$ in a Schur complement block $S_{i}$ (see Section 3.1). In case of MUMPS/HMAT, the value of $n_{c}$ ranges from 512 to 4,096 columns. As of MUMPS/SPIDO, it varies between 32 and 256. When it comes to multi-factorization, we vary the count $n_{b}$ of Schur complement blocks $S_{i j}$ per row or column of the submatrix $A_{s s}$ (see Section 3.2). Regardless the solver coupling, the tested values of $n_{b}$ are between 1 and 12 blocks per row or column.
In Figure 5. for each solver coupling, we show the best computation times of multi-solve and multi-factorization among all of the evaluated configurations. The worst performing variant seems to be multi-factorization using SPIDO as dense solver. Using HMAT instead, we can benefit from data compression and lower the memory footprint. This allows us to consider fewer Schur complement blocks and consequently speed-up the computation by reducing the number of sparse factorizations performed by the routine CreateSchurComplement (see Section 3.2). Up to $N=10^{6}$, this configuration can even outperform the multi-solve scheme. However, in case of larger systems, the memory constraint prevents us from lowering $n_{b}$ enough despite using data compression. It may seem the compression is necessary for multi-factorization to be competitive at all. However, in the current implementation of SPIDO, the total size of a block in bytes is encoded using a 32-bit integer type. In double complex arithmetic, the size of one element is 16 bytes which limits the count of rows and columns of one block to $m$ such that $16 m^{2}=\frac{2^{32}}{2}-1$ allowing $m$ to be at most 11,585 . This restricts the size of a Schur complement block in the case of the MUMPS/SPIDO coupling. Consequently, we can not always consider $n_{b}$ as low as possible even if there is enough of memory available. Although we would not be able to lower $n_{b}$ as much as in the case of the MUMPS/HMAT coupling due to the significantly hihger memory consumption of SPIDO compared to HMAT, a better performance for the MUMPS/SPIDO coupling might be achieved without this limitation. The multi-solve im-
plementation is not the best performing one, but its lower memory requirements allow us to process systems with up to $4 \times 10^{6}$ unknowns in the case of MUMPS/SPIDO. We have recently implemented multi-solve using HMAT as dense solver too.

It appears that using HMAT, we can speed up


Figure 5 - Best computation times of multisolve and multi-factorization for both of the solver couplings MUMPS/HMAT and MUMPS/SPIDO. Parallel runs using 24 threads on single miriel node. the computations but there is still a limitation in our implementation preventing us from running cases with $\mathrm{N}>10^{6}$ independently from the memory constraint. As of multisolve, reducing the number of Schur complement blocks impacts each of the couplings differently. For MUMPS/SPIDO, it means solve operations with more right-hand sides treated simultaneously by MUMPS, hence improving performance, although not to the same extent as in multi-factorization. Moreover, increasing the block size too much would prevent us from solving systems with up to $4 \times 10^{6}$ unknowns due to the memory constraint. For MUMPS/HMAT, we need to take into account an additional computational cost of transforming the Schur complement blocks into $\mathcal{H}$-matrices. To minimize this cost, we use substantially bigger blocks, i.e. blocks of 2,048 columns compared to 256 in case of MUMPS/SPIDO. For now, setting $n_{c}$ to values bigger than 2,048 seems causing MUMPS to produce incorrect results and does not allow us to further increase the block size regardless the memory constraint.

## 5. Conclusion

We have proposed and evaluated the multi-solve and multi-factorization two-stage implementation schemes for solving the target FEM/BEM linear system using both of the available solver couplings relying on MUMPS as sparse solver and on HMAT or on SPIDO as dense solver. According to our findings, replacing the usage of SPIDO by the HMAT solver based on the lowrank data compression can speed-up the computation in case of both implementation schemes. As of the overall comparison between the two schemes, we are not able to make a clear statement for now. Multi-solve using the coupling with SPIDO allowed us to process the largest coupled systems we have considered. On the other hand, due to the current limitations in our implementation of multi-solve with HMAT as dense solver, we could not confirm yet that this variant performs better on systems with more than $10^{6}$ unknowns in total.
In the future, we should address the implementation limitations, perform a more complex study of the two-stage implementation schemes and try to establish a criterion determining when to prefer multi-solve over multi-factorization and vice-versa. Then, we shall the study of the impact of out-of-core computation. Eventually, we shall extend our benchmarks to multiple computational nodes. We also plan to include other solver alternatives such as qr_mumps [11, 4] (version 3 now includes a sparse Gaussian elimination) and further explore the possibilities of the alternative single-stage implementation schemes relying on one single solver optimized for both sparse and dense matrix operations.

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## Appendix

```
Algorithm 1: Multi-solve algorithm for computing the Schur complement \(S\) based on Equa-
tion 5 and solving the target system (see Equation 1).
Function Multi-solve \((A, b)\) :
    \(\triangleright\) USING A SPARSE SOLVER :
    \(A_{\nu v} \leftarrow L^{\top}\)-Factorization \(\left(A_{\nu v}\right)\)
    for \(i=1\) to \(n_{B E M} / n_{c}\) do
        \(\triangleright\) Using the \(i^{\text {th }}\) block of columns of \(A_{s v}^{\top}\) as right-hand side :
        \(\mathrm{Y} \leftarrow\) TriangularSolve \(\left(A_{v v}, A_{s v_{i}}^{\top}\right)\)
        \(\mathrm{Z} \leftarrow \mathrm{A}_{s v} \times \mathrm{Y}\)
        \(A_{s s_{i}} \leftarrow A_{s s_{i}}-Z\)
        \(\triangleright\) AXPY
    \(\mathrm{b}_{v} \leftarrow\) TriangularSolve \(\left(\mathrm{A}_{v v}, \mathrm{~b}_{v}\right)\)
    \(\triangleright\) USING A DENSE SOLVER:
    \(A_{\text {ss }} \leftarrow\) LL \(^{\top}\)-Factorization \(\left(A_{s s}\right)\)
    \(\chi_{s} \leftarrow \operatorname{TriangularSolve}\left(A_{s s}, b_{s}-A_{s v} b_{v}\right)\)
```

Algorithm 2: Multi-factorization algorithm for computing the Schur complement $S$ based on Equation 6 and solving the target system (see Equation 1 ).

```
Function Multi-factorization ( \(A, b\) ) :
    \(\triangleright\) USING A SPARSE SOLVER:
    for \(i=1\) to \(n_{b}\) do
        for \(\mathrm{j}=1\) to \(\mathrm{n}_{\mathrm{b}}\) do
            \(W \leftarrow\left[\begin{array}{cc}A_{v v} & A_{s v_{j}}^{\top} \\ A_{s v_{i}} & 0\end{array}\right]\)
                \(A_{s s_{i j}} \leftarrow A_{s s_{i j}}+\) SchurComplement (W)
    \(\mathrm{A}_{\nu v} \leftarrow\) LL \(^{\mathrm{T}}\)-Factorization \(\left(\mathrm{A}_{v v}\right)\)
    \(\mathrm{b}_{v} \leftarrow\) TriangularSolve \(\left(\mathrm{A}_{v v}, \mathrm{~b}_{v}\right)\)
    \(\triangleright\) Using a Dense solver :
    \(A_{\text {ss }} \leftarrow\) LL \(^{\mathrm{T}}\)-Factorization \(\left(A_{\text {ss }}\right)\)
    \(\chi_{s} \leftarrow\) TriangularSolve \(\left(A_{s s}, b_{s}-A_{s v} b_{v}\right)\)
```



Figure 6 - Short pipe test case (length : 2 m , radius : 4 m ) with BEM surface mesh in red and FEM volume mesh in green.

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FIGURE 7 - Hardware configuration scheme of miriel nodes determined using the Linux Shell command lstopo.

| Software | Commit/Version | Branch | Licensing |
| :--- | :---: | :---: | :---: |
| HMAT | 912aa93 | master | proprietary |
| MPF | 432 b 53 b | $\mathrm{mf} /$ devel | proprietary |
| SCAB (test_FEMBEM) | $0 \mathrm{0c} 09 \mathrm{~d} 1 \mathrm{f}$ | $\mathrm{mf} /$ devel | proprietary |
| gcc [1] | 9.3 .0 | $\mathrm{~N} / \mathrm{A}$ | open-source |
| OpenMPI | 4.1 .0 | $\mathrm{~N} / \mathrm{A}$ | open-source |
| Intel(R) MKL | 2019.1 .144 | $\mathrm{~N} / \mathrm{A}$ | proprietary |
| MUMPS [7] | 5.2 .1 | $\mathrm{~N} / \mathrm{A}$ | open-source |
| StarPU [8] | 1.3 .7 | $\mathrm{~N} / \mathrm{A}$ | open-source |

TABLE 1-Commit and version numbers of selected packages used in the experimental software environment of the study.

| Total unknowns (N) | \# BEM unknowns (n $n_{\text {BEM }}$ ) | \# FEM unknowns |
| :--- | :--- | :--- |
| 250,000 | 14,835 | 235,165 |
| 500,000 | 23,577 | 476,423 |
| $1,000,000$ | 37,169 | 962,831 |
| $1,200,000$ | 41,992 | $1,158,008$ |
| $1,400,000$ | 46,482 | $1,353,518$ |
| $1,500,000$ | 48,750 | $1,451,250$ |
| $2,000,000$ | 58,910 | $1,941,090$ |
| $4,000,000$ | 93,593 | $3,906,407$ |

TABLE 2 - Counts of BEM and FEM unknowns in the target coupled FEM/BEM systems according to the overall unknown count.

