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## Lvalue closures

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# Lvalue closures

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<https://modernc.gforge.inria.fr/>



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# Policy

- *extend* the standard
  - valid code remains valid
  - new feature integrates syntactically and semantically
- fix as much requirements as possible through constraints
  - specific syntax
  - explicit constraints
- avoid new undefined behavior
  - only, if property is not (or hardly) detectable at translation time
  - or we leave design space to implementations
- don't mess with ABI
  - no changes
  - no extensions

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# Example: simple lvalue capture

```
auto const λ0 = [ ](void) { printf("%d\n", var); }; // invalid
auto const λ1 = [ var](void) { printf("%d\n", var); }; // freeze var
auto const λ2 = [&var](void) { printf("%d\n", var); }; // current
```

- It is a conscient user decision if and when ‘var’ is evaluated
- these policies can even be combined:

```
auto const λ3 = [var0 = var, &var](void) {
    printf("var0 previously: %d, now: %d\n", var0, var);
};
```



# Example: access a lambda from another one

```

void matmult(size_t k0, size_t l0, size_t m0,
             double A0[k0][l0], double B0[l0][m0], double C0[k0][m0]) {
    // dot product, stride of m0 for B0, constant prop of l0 and m0
    auto const λ1 =
        [l0, m0](double v1[l0], double B1[l0][m0], size_t m1) {
            double ret = 0.0;
            for (size_t i = 0; i < l0; ++i) ret += v1[i]*B1[i][m1];
            return ret;
        };
    // vector matrix product, ensure accessibility of λ1
    auto const λ2 =
        [l0, m0, &λ1](double v2[l0], double B2[l0][m0], double r[m0]) {
            for (size_t m = 0; m < m0; ++m) r[m] = λ1(v2, B2, m);
        };
    for (size_t k = 0; k < k0; ++k) λ2(A[k], B, C[k]);
}

```

- $\lambda_1$  is accessed by  $\lambda_2$ , but the address is not taken
- dimensions of  $B_1$  and  $B_2$  are fixed, once  $\lambda_1$  and  $\lambda_2$  are evaluated



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# Terminology

C does not have C++' references

- *lvalue capture* and *lvalue closure*

# Syntax

- a capture with a leading `&` is an lvalue capture

```
& identifier
```

- object is accessible (fresh and modifiable) throughout lambda body
- no support for C++ aliasing syntax

```
& identifier = unary-expression
```

- support for default capture mechanisms
  - `[=]` or `[=, &id, ...]`
    - default is value capture, `id` lvalue capture
  - `[&]` or `[&, id, ...]`
    - default is lvalue capture, `id` value capture

# Semantics

- lvalue capture  $\Leftrightarrow$  access to variable of an outerscope
- lvalue closure  $\lambda_1$  can be returned from another lvalue closure  $\lambda_2$

```

size_t n;
double R[10];
... // after a long decision process n is now fixed
auto const  $\lambda_2$  = [&R, n](void) {
    size_t m = really_complicated_computation(n);
    auto const  $\lambda_1$  = [&R, m](void) {
        // returns a double*
        return &R[m];
    };
    // returns a lambda value
    return  $\lambda_1$ ;
};
auto const Fun $\lambda$  =  $\lambda_2$ ();
for (i = 0; i < k; ++i) *Fun $\lambda$ () += i;

```

- **only if** lvalue closures of  $\lambda_1$  are a subset of those for  $\lambda_2$
- safe access of outer scope

# Semantics

- lvalue capture is **not** taking the address
- **register** variables possible  
(take some precautions)
- it is up to the implementation how to implement this  
(function, **goto**, **set jmp**, ...)
- potential access conflict with another thread  
(when we have wide function pointers)

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# Existing extensions

## C++ and widely used gcc extensions

	language	value capture	lvalue capture
nested function	gcc' C	no	always
statement expression	gcc' C	no	always
blocks	objective C gcc' C and C++	default	object property
lambda	C++	explicit	explicit

possible mechanical transformation of existing code in many cases



# New C23 feature `defer` ?

In other programming languages `defer` allows to freeze values.

## Use with lambda

If `defer` sees a lambda value

```
defer [ var](void){ printf("%d\n", var); }; // freeze var
defer [&var](void){ printf("%d\n", var); }; // current
```

when leaving the current block for whatever reason:  
execute the lambda



# New C23 feature `defer` ?

## Use without lambda

A `defer` with just an expression or compound statement

```
double* array = malloc(sizeof(double[n]));
defer free(array);           // or equivalent
defer { free(array); }
```

is equivalent to

```
double* array = malloc(sizeof(double[n]));
defer [&](void){ free(array); };
```

is equivalent to

```
double* array = malloc(sizeof(double[n]));
defer [&array](void){ free(array); };
```