



Lvalue closures

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Lvalue closures

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<https://modernc.gforge.inria.fr/>



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Policy

- *extend* the standard
 - valid code remains valid
 - new feature integrates syntactically and semantically
- fix as much requirements as possible through constraints
 - specific syntax
 - explicit constraints
- avoid new undefined behavior
 - only, if property is not (or hardly) detectable at translation time
 - or we leave design space to implementations
- don't mess with ABI
 - no changes
 - no extensions



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Example: simple lvalue capture

```
auto const  $\lambda_0$  = [ ] (void) { printf("%d\n", var); }; // invalid
auto const  $\lambda_1$  = [ var ] (void) { printf("%d\n", var); }; // freeze var
auto const  $\lambda_2$  = [ &var ] (void) { printf("%d\n", var); }; // current
```

- It is a conscient user decision if and when 'var' is evaluated
- these policies can even be combined:

```
auto const  $\lambda_3$  = [ var_0 = var, &var ] (void) {
    printf("var_previously: %d, now: %d\n", var_0, var);
};
```



Example: access a lambda from another one

```

void matmult(size_t k0, size_t l0, size_t m0,
             double A0[k0][l0], double B0[l0][m0], double C0[k0][m0]) {
    // dot product, stride of m0 for B0, constant prop of l0 and m0
    auto const  $\lambda_1$  =
        [l0, m0](double v1[l0], double B1[l0][m0], size_t m1) {
            double ret = 0.0;
            for (size_t i = 0; i < l0; ++i) ret += v1[i]*B1[i][m1];
            return ret;
        };
    // vector matrix product, ensure accessibility of  $\lambda_1$ 
    auto const  $\lambda_2$  =
        [l0, m0, & $\lambda_1$ ](double v2[l0], double B2[l0][m0], double r[m0]) {
            for (size_t m = 0; m < m0; ++m) r[m] =  $\lambda_1$ (v2, B2, m);
        };
    for (size_t k = 0; k < k0; ++k)  $\lambda_2$ (A[k], B, C[k]);
}

```

- λ_1 is accessed by λ_2 , but the address is not taken
- dimensions of B_1 and B_2 are fixed, once λ_1 and λ_2 are evaluated



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Terminology

C does not have C++' references

- *lvalue capture* and *lvalue closure*



Syntax

- a capture with a leading **&** is an lvalue capture

```
& identifier
```

- object is accessible (fresh and modifiable) throughout lambda body
- no support for C++ aliasing syntax

```
& identifier = unary-expression
```

- support for default capture mechanisms
 - **[=]** or **[=, &id, ...]**
 - default is value capture, **id** lvalue capture
 - **[&]** or **[&, id, ...]**
 - default is lvalue capture, **id** value capture



Semantics

- lvalue capture $\Leftrightarrow$ access to variable of an outerscope
- lvalue closure λ_1 can be returned from another lvalue closure λ_2

```

size_t n;
double R[10];
... // after a long decision process n is now fixed
auto const  $\lambda_2$  = [&R, n](void) {
    size_t m = really_complicated_computation(n);
    auto const  $\lambda_1$  = [&R, m](void) {
        // returns a double*
        return &R[m];
    };
    // returns a lambda value
    return  $\lambda_1$ ;
};
auto const Fun $\lambda$  =  $\lambda_2$ ();
for (i = 0; i < k; ++i) *Fun $\lambda$ () += i;

```

- **only if** lvalue closures of λ_1 are a subset of those for λ_2
- safe access of outer scope



Semantics

- lvalue capture is **not** taking the address
- **register** variables possible
(take some precautions)
- it is up to the implementation how to implement this
(function, **goto**, **setjmp**, ...)
- potential access conflict with another thread
(when we have wide function pointers)



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Existing extensions

C++ and widely used gcc extensions

	language	value capture	lvalue capture
nested function	gcc' C	no	always
statement expression	gcc' C	no	always
blocks	objective C gcc' C and C++	default	object property
lambda	C++	explicit	explicit

possible mechanical transformation of existing code in many cases



New C23 feature `defer` ?

In other programming languages `defer` allows to freeze values.

Use with lambda

If `defer` sees a lambda value

```
defer [ var](void){ printf("%d\n", var); }; // freeze var
defer [&var](void){ printf("%d\n", var); }; // current
```

when leaving the current block for whatever reason:
execute the lambda



New C23 feature `defer` ?

Use without lambda

A `defer` with just an expression or compound statement

```
double* array = malloc(sizeof(double[n]));
defer free(array);           // or equivalent
defer { free(array); }
```

is equivalent to

```
double* array = malloc(sizeof(double[n]));
defer [&](void) { free(array); };
```

is equivalent to

```
double* array = malloc(sizeof(double[n]));
defer [&array](void) { free(array); };
```