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Farkas Lemma made easy

Tool Demonstration

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Abstract
In this paper, we present FKCC, a scripting tool to prototype program analyses and transformations exploiting the affine form of Farkas lemma. Our language is general enough to prototype in a few lines sophisticated termination and scheduling algorithms. The tool is freely available and may be tried online via a web interface. We believe that FKCC is the missing chain to accelerate the development of program analyses and transformations exploiting the affine form of Farkas lemma.

ACM Reference Format:

1 Introduction
Many program analyses require to handle universally quantified constraints such as \( \forall x \in D : \Phi(x) \), where \( D \) is a convex polyhedron and \( \Phi \) is a conjunction of affine constraints. For instance, this occurs in loop scheduling [5, 6], loop tiling [3], program termination [2] or generation of loop invariants [4]. Farkas lemma – affine form – provides a way to get rid of that universal quantification, at the price of introducing quadratic terms. It is even possible to use Farkas lemma to design the Pluto-style tiling [3] for the Jacobi 1D kernel.

In this tool demonstration, we present a scripting tool, FKCC [1] which makes it possible to manipulate easily Farkas lemma to benefit from those nice properties. Specifically, we will discuss the following points:

- A general formulation for the resolution of equations \( \forall x : S(x) = 0 \) where \( S \) is summation of affine forms including Farkas terms. So far, this resolution was applied for specific instances of Farkas summation. This result is the basic engine of the FKCC scripting language.
- A scripting language to apply and exploit Farkas lemma; among polyhedra, affine functions and affine forms.
- Our tool, FKCC, implementing these principles, is available at http://foobar.ens-lyon.fr/fkcc. FKCC may be downloaded and tried online via a web interface. FKCC comes with many examples, making it possible to adopt the tool easily.

This tool demonstration is structured as follows. Section 2 presents the affine form of Farkas lemma, our resolution theorem, and explains how it applies to compute scheduling functions. Then, Section 3 defines the syntax and outlines informally the semantics of the FKCC language. Finally, Section 4 concludes this paper and draws future research perspectives, then Annex A gives a real-life example of FKCC script to compute a Pluto-style tiling [3] for the Jacobi 1D kernel.

2 Farkas lemma in polyhedral compilation
This section presents the theoretical background of this tool demonstration. We first introduce the affine form of Farkas lemma. Then, we present our theorem to solve equations \( \forall x : S(x) = 0 \) where \( S \) is a summation of affine forms including Farkas terms. This formalization will then be exploited to design the FKCC language.

Lemma 2.1 (Farkas Lemma, affine form). Consider a non-empty convex polyhedron \( P = \{ x \mid A^T x + b \geq 0 \} \subseteq \mathbb{R}^n \) and an affine form \( \phi : \mathbb{R}^n \rightarrow \mathbb{R} \) such that \( \phi(x) \geq 0 \ \forall x \in P \). Then:
\[
\exists \lambda \geq 0, \lambda_0 \geq 0 \text{ such that: } \phi(x) = \sum_i \lambda_i \cdot x_i + \lambda_0 \ \forall x \in P.
\]

Hence, Farkas lemma makes it possible to remove the quantification \( \forall x \in P \) by encoding directly the positivity over \( P \) into the definition of \( \phi \), thanks to the Farkas multipliers \( \lambda \) and \( \lambda_0 \). In the remainder, Farkas terms will be denoted by: \( \bar{\phi}(\lambda_0, \lambda, A, b)(x) = \sum_i \lambda_i (A^T x + b) + \lambda_0 \). We now recall our theorem [1] to solve equations \( \forall x : S(x) = 0 \) where \( S \) involves Farkas terms. The result is expressed as a conjunction of affine constraints, which is suited for integer linear programming:

Theorem 2.2 (solve). Consider a summation \( S(x) = \bar{u} \cdot x + v + \sum_i \delta(\lambda_{i0}, \lambda_i, A_i, b_i)(x) \) of affine forms, including Farkas terms. Then:
\[
\forall x : S(x) = 0 \text{ iff } \left\{ \begin{array}{l}
\bar{u} + \sum_i \lambda_i \cdot A_i = 0 \wedge \\
v + \sum_i (\lambda_i \cdot b_i + \lambda_{i0}) = 0
\end{array} \right.
\]
Application to scheduling  Figure 1 depicts an example of a program (a) computing the product of two polynomials specified by their arrays of coefficients $a$ and $b$, and the iteration domain with the data dependence across iterations (b) and an example schedule $\theta(i,j) = i$ prescribing a parallel execution by vertical waves. This paragraph reformulates the technique presented in [5] with our theorem 2.2. This formulation will directly inspire the \texttt{FKCC} syntax presented in the next section.

A schedule must be positive everywhere on its iteration domain:

$$\theta(i,j,N) \geq 0 \quad \forall (i,j) \in D_N$$  \hfill (1)

Applying Farkas lemma, this translates to:

$$\exists \lambda_0 \geq 0, \bar{x} \geq 0 \quad \text{s.t.} \quad \theta(i,j,N) = \bar{\theta}(\lambda_0, \bar{x}, A, \bar{b})(i,j,N)$$  \hfill (2)

Moreover, a schedule must satisfy the data dependences $(i,j) \rightarrow (i',j')$, abstracted by a dependence polyhedron $\Delta_N$:

$$\theta(i',j',N) > \theta(i,j,N) \quad \forall (i,j,i',j') \in \Delta_N$$  \hfill (3)

This is equivalently written as the positivity of the affine form $(i,j,i',j',N) \rightarrow \theta(i',j',N) - \theta(i,j,N) - 1$ over the convex polyhedron $\Delta_N$. Applying Farkas lemma:

$$\exists \mu_0 \geq 0, \bar{\mu} \geq 0 \quad \text{s.t.} \quad \theta(i',j',N) - \theta(i,j,N) - 1 = \bar{\theta}(\mu_0, \bar{\mu}, C, \bar{d})(i,j,i',j',N)$$

Substituting $\theta$ using Equation (2), this translates to solving:

$$\forall (i,j,i',j',N) : S(i,j,i',j',N) = 0$$

where $S(i,j,i',j',N)$ is the summation:

$$\bar{\theta}(\lambda_0, \bar{x}, A, \bar{b})(i',j',N) - \bar{\theta}(\lambda_0, \bar{x}, A, \bar{b})(i,j,N) - 1 - \bar{\theta}(\mu_0, \bar{\mu}, C, \bar{d})(i,j,i',j',N)$$

Since $-\bar{\theta}(\lambda_0, \bar{x}, A, \bar{b}) = \bar{\theta}(\lambda_0, -\bar{x}, A, \bar{b})$, we may apply theorem 2.2 to obtain a system of affine constraints with $\lambda_0, \bar{x}, \mu_0, \bar{\mu}$. Linear programming may then be applied to find out the desired schedule [3, 6].

3 FKCC at a glance

This section outlines briefly the input syntax of FKCC on our motivating example. For a detailed description, the reader is referred to [1].

Program, instructions, polyhedra  An FKCC program consists of a sequence of instructions. There is no other control structure than the sequence. An instruction may assign an FKCC object (polyhedron, affine form or affine function) to an FKCC identifier, or may be an FKCC object alone. In the latter case, the FKCC object is streamed out to the standard output. FKCC objects are expressed with the same syntax as ICC7:

- \textbf{parameters} := \{M, eps\};
- \textbf{parametrized_iterations} := \{M\} \rightarrow \{[i] : 0 \leq i \leq M\};

Parameters \textbf{must} be declared with the parameters construct. The parameters of a polyhedron \textbf{may} optionally be declared on preceding brackets $[M] \rightarrow \ldots$. The set intersection of two polyhedra $P$ and $Q$ is obtained with $P \cap Q$.

\textbf{for} $i := 0$ \textbf{to} $N$

\textbf{for} $j := 0$ \textbf{to} $N$

$c[i+j] := c[i+j] + a[i]*b[j];$

(a) Product of polynomials

Affine forms  An affine form may be defined as a Farkas term:

- \textbf{iterations} := \{[] \rightarrow \{[i,j,N] : 0 \leq i \leq N \text{ and} \ 0 \leq j \leq N\};
- \textbf{theta} := \text{positive on iterations};

If \textbf{iterations} is \{$x \mid Ax + b \geq 0\}, then \textbf{theta} is defined as $\bar{\theta}(\lambda_0, \bar{x}, A, \bar{b})$ where $\lambda_0$ and $\bar{x}$ are fresh positive variables. In this case, the polyhedron is \textbf{never} parametrized: the \textbf{parameters must be handled as variables}. Affine forms may be summed, scaled and composed with \textbf{affine functions}, typically to adjust the input dimension:

- $\textbf{dependence} := \{[] \rightarrow \{[i,j,i',j',N] : 0 \leq i \leq N \text{ and} \ 0 \leq j \text{ and} \ j' \leq N \text{ and} \ 0 \leq i' \leq N \text{ and} \ 0 \leq j' \leq N \text{ and} \ i+j' = i'+j \text{ and} \ i \leq i'\};$
- $\textbf{to_target} := \{[i,j,i',j',N] \rightarrow [i',j',N];$
- $\textbf{to_source} := \{[i,j,i',j',N] \rightarrow [i,j,N];$
- $\textbf{sum} := \langle \text{theta . to_target} \rangle \rightarrow \langle \text{theta . to_source} \rangle - 1 - \text{positive on dependence};$

In a summation of affine forms, affine forms must have the same input dimension. Also, a constant $(-1)$ is automatically interpreted as an affine form $\langle [i,j,i',j',N] \rightarrow -1\rangle$. The terms of the summation are simply separated with $+$ and $-$, no parenthesis are allowed. Affine forms may also be stated explicitly:

- $\textbf{sum_eps} := \langle \text{theta . to_target} \rangle \rightarrow \langle \text{theta . to_source} \rangle + \langle [i,j,i',j',N] \rightarrow \text{-eps} \rangle - \text{positive on dependence};$

Resolution  The main feature of FKCC is the resolution of equations $\forall x : S(x) = 0$ where $S$ is a summation of affine forms including Farkas terms. This is obtained with the instruction $\textbf{solve}$:

- $\textbf{solve sum} := 0;$

The result is a polyhedron with Farkas multipliers (obtained after applying Theorem 2.2 (solve)),

\text{Solve constraints!};
When the summation contains affine forms with parameters (as \( \text{sum} \cdot \text{eps} \)), the resolution interprets parameters as constants. In particular, this makes it possible to tune dependence satisfaction: \( \theta(i', j') \geq \theta(i, j) + \epsilon_d \) with \( 0 \leq \epsilon_d \leq 1 \).

At this point, we need to recover the coefficients of our affine form \( \theta \) in terms of \( \lambda \) (\( \lambda_0, \ldots, \lambda_3 \)) and \( \lambda_0 \) (\( \lambda_4 \)). Observe that \( \theta(x) = \lambda_0 \lambda \lambda x + \lambda \lambda \lambda x + \lambda \lambda \lambda x + \lambda \lambda \lambda \lambda x \). If the coefficients of \( \theta \) are expressed as Farkas terms in a summation: \( \theta(x) = \lambda \cdot \lambda \lambda x + \lambda \lambda \lambda x + \lambda \lambda \lambda \lambda x \). We may filter the coefficients for several affine forms expressed as Farkas terms in a summation: \( \theta(x) = \lambda \cdot \lambda \lambda x + \lambda \lambda \lambda x + \lambda \lambda \lambda \lambda x \).

The result is a conjunction of definition equalities, gathered in a polyhedron:

\[
[] \Rightarrow \{ \langle \lambda_0, \lambda_1, \lambda_2, \lambda_3, \lambda_4 \rangle \mid \\
\quad \text{tau}_0 \cdot \text{tau}_1 \cdot \text{tau}_2 \cdot \text{tau}_3 = \text{tau}_0 \cdot \text{tau}_1 \cdot \text{tau}_2 \cdot \text{tau}_3 \} \\
\quad \text{tau}_0 \cdot \text{tau}_1 \cdot \text{tau}_2 \cdot \text{tau}_3 = \text{tau}_1 \cdot \text{tau}_2 \cdot \text{tau}_3 \\
\quad \text{tau}_0 \cdot \text{tau}_1 \cdot \text{tau}_2 \cdot \text{tau}_3 = \text{tau}_2 \cdot \text{tau}_3 \\
\quad \text{tau}_0 \cdot \text{tau}_1 \cdot \text{tau}_2 \cdot \text{tau}_3 = \text{tau}_3 \\
\quad \text{tau}_0 \cdot \text{tau}_1 \cdot \text{tau}_2 \cdot \text{tau}_3 = 0 \}
\]

The first coefficients \( \text{tau}_k \) define \( \tau \), the last one defines the constant \( \tau_0 \). On our example, \( \theta(i, j, N) = \text{tau}_0 \cdot i + \text{tau}_1 \cdot j + \text{tau}_2 \cdot N + \text{tau}_3 \). Now we may gather the results and eliminate the \( \lambda \) to keep only \( \tau \) and \( \tau_0 \):

\[
\text{keep tau}_0, \text{tau}_1, \text{tau}_2, \text{tau}_3 \text{ in}
\]

\[
\text{solvem sum} = 0 \} \text{define theta with tau)}
\]

The result is a polyhedron with all the valid schedules:

\[
[] \Rightarrow \{ \langle \text{tau}_0, \text{tau}_1, \text{tau}_2, \text{tau}_3 \rangle \mid \\
\quad \text{tau}_3 \geq 0 \text{ and (tau}_0 \cdot \text{tau}_1 \cdot \text{tau}_2 \cdot \text{tau}_3 > 0 \text{ and} \\
\quad \text{tau}_1 \cdot \text{tau}_2 \cdot \text{tau}_3 > 0 \text{ and } \text{tau}_2 \cdot \text{tau}_3 > 0 \}
\]

All these steps may be applied at once with the \text{find} command:

\[
\text{find theta s.t. sum} = 0 \\
\]

The coefficients are automatically named \( \text{theta}_0, \text{theta}_1, \text{theta}_2, \text{theta}_3 \) etc with the same convention as \( \text{define} \). We point out that \text{find} fresh names for coefficients (e.g. \( \text{tau}_4, \text{tau}_5 \) on the second time with ‘\( \text{tau}_4 \)’) whereas \text{find} always chooses the same names. Hence \text{find} would be preferred when deriving separately constraints on the coefficients of \( \theta \). We may filter the coefficients for several affine forms expressed as Farkas terms in a summation:

\[
\text{find theta_s,theta_t s.t.} \\
\text{theta_T.to_target - theta_S.to_source - 1} \\
\text{= positive_on dependences_from_S_to_T} = 0
\]

This is typically used to compute schedules for programs with multiple assignments (here \( S \) and \( T \) with dependences from iterations of \( S \) to iterations of \( T \)). Finally, note that keep \( \text{tau}_0, \text{tau}_1, \text{tau}_2, \text{tau}_3 \) in \( P \); projects \( P \) on variables \( \text{tau}_0, \text{tau}_1, \text{tau}_2, \text{tau}_3 \); the result is a polyhedron with integral points of coordinates (\( \text{tau}_0, \text{tau}_1, \text{tau}_2, \text{tau}_3 \)). This way, the order in which \( \text{tau}_0, \text{tau}_1, \text{tau}_2, \text{tau}_3 \) are specified to keep impacts directly a further lexicographic optimization. As in iscc, the lexicographic minimum of a polyhedron is obtained with the command \text{lexmin}. This example, however, does not admit a lexico-minimum solution. We can find a solution by forcing positive coefficients:

\[
\text{positive_theta := [ ] \rightarrow \{ \langle \text{theta}_0, \text{theta}_1, \text{theta}_2, \text{theta}_3 \rangle \mid \\
\quad \text{theta}_0 > 0 \text{ and } \text{theta}_1 > 0 \text{ and } \text{theta}_2 > 0 \text{ and } \text{theta}_3 > 0 \} \\
\quad \text{lexmin (find theta s.t. sum = 0) * positive_theta}
\]

Using the ‘\text{-pretty}’ option, we obtain:

\[
\theta_0 = 1 \\
\theta_1 = 0 \\
\theta_2 = 0 \\
\theta_3 = 0
\]

Which corresponds to the schedule \( \theta(i, j, N) = i \). We point out that it is also possible to write in a few lines of \text{fkcc} the latency minimization as in [5]. A complete description may be found in [1].

4 Conclusion

\text{FKCC} is a scripting tool to prototype program analyses using the affine form of Farkas lemma. The script language of \text{FKCC} is powerful enough to write in a few lines sophisticated scheduling algorithms. The object representation (polyhedral, affine functions) is compatible with iscc, a widespread polyhedral tool featuring manipulation of affine relations. \text{FKCC} provides features to generate iscc code, and conversely, the output of iscc might be injected in \text{FKCC}. This will allow to take profit of both worlds.

We believe that scripting tools are mandatory to evaluate rapidly research ideas. So far, Farkas lemma-based approaches were locked by two facts: applying by hand Farkas Lemma is a pain; and implementing an analysis with Farkas lemma is notoriously time consuming and bug prone. With \text{FKCC}, computer scientists are now freed from these constraints.

Let the power of \text{FKCC} be with you!

References


A Pluto-style tiling with \texttt{fkcc}

We give here the \texttt{fkcc} script discussed in the demo session to compute a Pluto-style tiling \cite{3} for the Jacobi-1D kernel. Note the sections expressing successively the data dependences (→), the correctness condition ($⟨S,i⟩ → ⟨T,j⟩ ⇒ φ_T(j) − φ_S(i) ≥ 0$), the laziness (minimal dependence distance $δ_{ST} = φ_T(j) − φ_S(i)$), and finally the computation of non-null and linearly independent affine tiling hyperplanes in $ϕ_S$ and $ϕ_T$. Applying \texttt{fkcc} with the \texttt{-pretty} option, we obtain:

```bash
$ fkcc -pretty < pluto.fk
phi_S_0 = 1
phi_S_0 = 2
phi_S_1 = 0
phi_S_1 = 1
phi_S_2 = 0
phi_S_2 = 0
phi_S_3 = 0
phi_S_3 = 0
phi_S_4 = 0
phi_S_4 = 0
phi_T_0 = 1
phi_T_0 = 2
phi_T_1 = 0
phi_T_1 = 1
phi_T_2 = 0
phi_T_2 = 0
phi_T_3 = 0
phi_T_3 = 0
phi_T_4 = 0
phi_T_4 = 1
latency_0 = 0
latency_0 = 0
latency_1 = 0
latency_1 = 0
latency_2 = 1
latency_2 = 2
```

Which corresponds to the affine tiling $ϕ_S(t,i) = (t,2t + i)$ and $ϕ_T(t,i) = (t,2t + i + 1)$ with a maximum dependence distance 1 for the first tiling hyperplane and 2 for the second tiling hyperplane.

```bash
# # jacobi-1D non-perfect
# #
# # for (t = 1; t <= T; t++)
# # {
# #   # for (i = 1; i < N - 1; i++)
# #   # for (i = 1; i < N - 1; i++)
# #   #   A[i] = B[i]; //T
# # # }

D_S := [] -> { [t,i,T,N]: 1 <= t and t <= T and 1 <= i and i <= N-1};
D_T := [] -> { [t,i,T,N]: 1 <= t and t <= T and 1 <= i and i <= N-1};
phi_S := positive_on D_S;
phi_T := positive_on D_T;

#S --> T
Delta_ST := [] -> { [t,i,t',i',T,N]: 1 <= t and t <= T and 1 <= i and i <= N-1 and t'=t and i'=i+1};

#S --> T (anti read(a[i-1]) --> write(a[i]))
Delta_ST_anti_1 := [] -> { [t,i,t',i',T,N]: 1 <= t and t <= T and 1 <= i and i <= N-1 and 1 <= t' and t' <= T and 1 <= i' and i' <= N-1 and t+1=t' and i=i'-1};

#S --> T (anti read(a[i+1]) --> write(a[i]))
Delta_ST_anti_2 := [] -> { [t,i,t',i',T,N]: 1 <= t and t <= T and 1 <= i and i <= N-1 and 1 <= t' and t' <= T and 1 <= i' and i' <= N-1 and t+1=t' and i=i'+1};

#T --> S, read a[i-1]
Delta_TS_1 := [] -> { [t,i,t',i',T,N]: 1 <= t and t <= T and 1 <= i and i <= N-1 and 1 <= t' and t' <= T and 1 <= i' and i' <= N-1 and t+1=t' and i=i-1'};

#T --> S, read a[i]
Delta_TS_2 := [] -> { [t,i,t',i',T,N]: 1 <= t and t <= T and 1 <= i and i <= N-1 and 1 <= t' and t' <= T and 1 <= i' and i' <= N-1 and t+1=t' and i=i'};

#T --> S, read a[i+1]
Delta_TS_3 := [] -> { [t,i,t',i',T,N]: 1 <= t and t <= T and 1 <= i and i <= N-1 and 1 <= t' and t' <= T and 1 <= i' and i' <= N-1 and t+1=t' and i=i'+1'};
```
# Correctness: s -> t ==> phi(s) <= phi(t)
#

to_target := \{[t, i, t', i', T, N] -> [t', i', T, N]\};
to_source := \{[t, i, t', i', T, N] -> [t, i, T, N]\};

phi_correct :=
# S --> T
(find phi_S, phi_T s.t. (phi_T . to_target) - (phi_S . to_source) - positive_on Delta_ST = 0) *
# S --> T, (anti read(a[i-1]) --> write(a[i]))
(find phi_S, phi_T s.t. (phi_T . to_target) - (phi_S . to_source) - positive_on Delta_ST_anti_1 = 0) *
# S --> T, (anti read(a[i+1]) --> write(a[i]))
(find phi_S, phi_T s.t. (phi_T . to_target) - (phi_S . to_source) - positive_on Delta_ST_anti_2 = 0) *
# T --> S, read a[i-1]
(find phi_S, phi_T s.t. (phi_S . to_target) - (phi_T . to_source) - positive_on Delta_TS_1 = 0) *
# T --> S, read a[i]
(find phi_S, phi_T s.t. (phi_S . to_target) - (phi_T . to_source) - positive_on Delta_TS_2 = 0) *
# T --> S, read a[i+1]
(find phi_S, phi_T s.t. (phi_S . to_target) - (phi_T . to_source) - positive_on Delta_TS_3 = 0);

# Efficiency: s --> t ==> phi(t) - phi(s) <= latency(N), then min latency(N)
#

# L(N) >= 0 on the parameter domain
latency := positive_on ([] -> \{[T, N]: T >= 0 and N >= 0\});
latency_def := define latency with latency;
to_param := \{[t, i, t', i', T, N] -> [T, N]\};

# phi(t) - phi(s) <= L(N) \forall s --> t
phi_bounded :=
# S --> T
(find latency, phi_S, phi_T s.t. latency . to_param - (phi_T . to_target) + (phi_S . to_source) - positive_on Delta_ST = 0) *
# S --> T, (anti read(a[i-1]) --> write(a[i]))
(find latency, phi_S, phi_T s.t. latency . to_param - (phi_T . to_target) + (phi_S . to_source) - positive_on Delta_ST_anti_1 = 0) *
# S --> T, (anti read(a[i+1]) --> write(a[i]))
(find latency, phi_S, phi_T s.t. latency . to_param - (phi_T . to_target) + (phi_S . to_source) - positive_on Delta_ST_anti_2 = 0) *
# T --> S, read a[i-1]
(find latency, phi_S, phi_T s.t. latency . to_param - (phi_S . to_target) + (phi_T . to_source) - positive_on Delta_TS_1 = 0) *
# T --> S, read a[i]
(find latency, phi_S, phi_T s.t. latency . to_param - (phi_S . to_target) + (phi_T . to_source) - positive_on Delta_TS_2 = 0) *
# T --> S, read a[i+1]
(find latency, phi_S, phi_T s.t. latency . to_param - (phi_S . to_target) + (phi_T . to_source) - positive_on Delta_TS_3 = 0);

# First hyperplane: avoid the null solution
#
phi_filter_level_1 := [] -> \{[phi_S_0, phi_S_1, phi_S_2, phi_S_3, phi_S_4, phi_T_0, phi_T_1, phi_T_2, phi_T_3, phi_T_4]\):

#phi_S: positive coefficients + no parameters
phi_S_0 := 0 and
phi_S_1 := 0 and
phi_S_2 := 0 and
phi_S_3 := 0 and
phi_S_4 := 0 and

#phi_T: positive coefficients + no parameters
phi_T_0 := 0 and
phi_T_1 := 0 and
phi_T_2 := 0 and
phi_T_3 := 0 and
phi_T_4 := 0 and
\texttt{#phi_S} \neq 0 \text{ and } \phi_T \neq 0
\phi_S_0 + \phi_S_1 + \phi_S_2 + \phi_S_3 \geq 1 \text{ and }
\phi_S_0 + \phi_T_1 + \phi_T_2 + \phi_T_3 \geq 1
};
\text{all_level}_1 := \text{keep latency}_0, \text{latency}_1, \text{latency}_2, \phi_S_0, \phi_S_1, \phi_S_2, \phi_S_3, \phi_S_4, \phi_T_0, \phi_T_1, \phi_T_2, \phi_T_3, \phi_T_4
\text{in phi_correct} \ast \phi_{\text{bounded}} \ast \phi_{\text{filter_level}_1};
\text{lexmin all_level}_1;

\#
\# \text{Second hyperplane: avoid the null solution} + \text{linear independence with the first hyperplane}
#
\phi_{\text{filter_level}_2} := [] \rightarrow \{[\phi_S_0, \phi_S_1, \phi_S_2, \phi_S_3, \phi_S_4, \phi_T_0, \phi_T_1, \phi_T_2, \phi_T_3, \phi_T_4]:
\#
\#\text{phi}_S: \text{positive coefficients} + \text{no parameters} + \text{lin. ind (coef(i) > 0)}
\phi_S_0 \geq 0 \text{ and }
\phi_S_1 > 0 \text{ and }
\phi_S_2 = 0 \text{ and }
\phi_S_3 = 0 \text{ and }
\phi_S_4 \geq 0 \text{ and }
\#
\#\text{phi}_T: \text{positive coefficients} + \text{no parameters} + \text{lin. ind (coef(i) > 0)}
\phi_T_0 \geq 0 \text{ and }
\phi_T_1 > 0 \text{ and }
\phi_T_2 = 0 \text{ and }
\phi_T_3 = 0 \text{ and }
\phi_T_4 \geq 0 \text{ and }
\#
\text{phi}_S \neq 0 \text{ and } \phi_T \neq 0
\phi_S_0 + \phi_S_1 + \phi_S_2 + \phi_S_3 \geq 1 \text{ and }
\phi_S_0 + \phi_T_1 + \phi_T_2 + \phi_T_3 \geq 1
}\};
\text{all_level}_2 := \text{keep latency}_0, \text{latency}_1, \text{latency}_2, \phi_S_0, \phi_S_1, \phi_S_2, \phi_S_3, \phi_S_4, \phi_T_0, \phi_T_1, \phi_T_2, \phi_T_3, \phi_T_4
\text{in phi_correct} \ast \phi_{\text{bounded}} \ast \phi_{\text{filter_level}_2};
\text{lexmin all_level}_2;