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On Combinatorial Proofs for Logics of Relevance and Entailment

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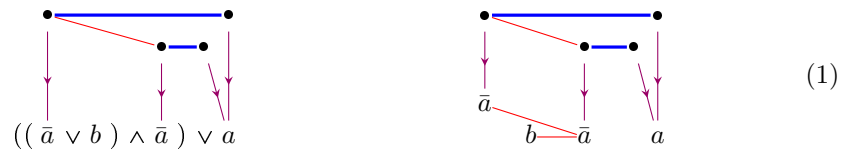
Abstract. Hughes’ combinatorial proofs give canonical representations for classical logic proofs. In this paper we characterize classical combinatorial proofs which also represent valid proofs for relevant logic with and without the mingle axiom. Moreover, we extend our syntax in order to represent combinatorial proofs for the more restrictive framework of entailment logic.

Keywords: Combinatorial Proofs · Relevant Logic · Entailment Logic · Skew fibrations · Proof Theory

1 Introduction

Combinatorial proofs have been conceived by Hughes [11] as a way to write *proofs* for classical propositional logic *without syntax*. Informally speaking, a combinatorial proof consists of two parts: first, a purely linear part, and second, a part that handles duplication and erasure. More formally, the first part is a variant of a proof net of multiplicative linear logic (MLL), and the second part is given by a skew fibration (or equivalently, a contraction-weakening map) from the cograph of the conclusion of the MLL proof net to the cograph of the conclusion of the whole proof. For the sake of a concise presentation, the MLL proof net is given as a cograph together with a perfect matching on the vertices of that graph. An important point is that in order to represent correct proofs, the proof nets have to obey a connectedness- and an acyclicity-condition.

To give an example, we show here the combinatorial proof of Pierce’s law $((a \rightarrow b) \rightarrow a) \rightarrow a$, which can be written in negation normal form (NNF) as $(\bar{a} \vee b) \vee \bar{a} \vee a$:



On the left above, we have written the conclusion of the proof as formula, and on the right as cograph, whose vertices are the atom occurrences of the formula, and whose edges are depicted as regular (red) lines. The linear part of the proof is given by the cograph determined by the four vertices and the regular (red)

edge in the upper half of the diagram. The perfect matching is indicated by the bold (blue) edges. Finally, the downward arrows describe the skew fibration. In our example there is one atom in the conclusion (the a) that is the image of two vertices, indicating that it is the subject of a *contraction* in the proof. Then there is an atom (the b) that is not the image of any vertex above, indicating that it is coming from a *weakening* in the proof.

Relevance logics have been studied by philosophers [2,3] to investigate when an implication is *relevant*, i.e., uses all its premises. In particular, in relevance logic, the implication $A \rightarrow (B \rightarrow A)$ is rejected because the B is not used to draw the conclusion A . In other words, we can no longer deduce A from $A \wedge B$. Put in proof theoretic terms, this corresponds to disallowing the *weakening* rule in a proof system. Carrying this observation to our combinatorial proofs mentioned above, this says that the skew fibration, that maps the linear part to the conclusion, must be a *relevant*, i.e. surjective with respect of vertices and edges of the cographs. The first contribution of this paper is to show that the converse also holds. a classical combinatorial proof is a proof of relevance logic if and only if its skew fibration is surjective.

The mingle axiom is in its original form $A \rightarrow (A \rightarrow A)$ [2, p.97]. In the implication-negation fragment of relevant logic, it can be derived from the more primitive form $A \rightarrow (B \rightarrow (\bar{B} \rightarrow A))$ (see also [2, p.148]), which is equivalent to $(A \wedge B) \rightarrow (A \vee B)$, which is known as *mix* in the linear logic community. When mix is added to MLL, the connectedness-condition has to be dropped. This leads to the second result of this paper: adding mingle to relevance logic corresponds to dropping the connectedness condition from the combinatorial proofs.

Interestingly, Hughes' original version of combinatorial proofs included *mix* (i.e., there was no connectedness-condition). If *weakening* $\perp \rightarrow A$ is present, then mix is derivable, so that the presence or absence of mix does not have an effect on provability. However, when weakening is absent, as it is the case with many substructural logics, then mix/mingle has an impact on provability, and for this reason, we present combinatorial proofs in their basic form without mix, and follow the presentation in [17], using the notion of RB-cographs due to Retoré [15].

Entailment logic is a further refinement of relevance logic, insisting not only on the relevance of premises but also on their necessity (in the sense of the modal logic S4)³. More precisely, the logic rejects the implication $A \rightarrow ((A \rightarrow A) \rightarrow A)$. In terms of the sequent calculus, this means that the two sequents

$$\Gamma \vdash A \rightarrow B \quad \text{and} \quad \Gamma, A \vdash B \quad (2)$$

are only equivalent if all formulas in the context Γ are of shape $C \rightarrow D$ for some C and D . If we write $A \rightarrow B$ as $\bar{A} \vee B$, then \vee is not associative, as the rejected

³ We do not discuss the philosophical considerations leading to this logic. For this, the reader is referred to the Book by Anderson and Belnap [2]. We take here the logic as given and discuss its proofs.

$A \rightarrow ((A \rightarrow A) \rightarrow A)$ would be written as $\bar{A} \vee ((A \wedge \bar{A}) \vee A)$, and the accepted $(A \rightarrow A) \rightarrow (A \rightarrow A)$ as $(A \wedge \bar{A}) \vee (\bar{A} \vee A)$. The consequence of this is that in combinatorial proofs we can no longer use simple cographs to encode formulas, as these identify formulas up to associativity and commutativity of \wedge and \vee . We solve this problem by putting weights on the edges in the graphs. This leads us to our third contribution of this paper: combinatorial proofs for entailment logic.

Outline of the paper In this paper we study the *implication-negation-fragment* of relevance logic. For this, we recall in Section 2 the corresponding sequent calculi, following the presentation in [13] and [5]. Then, in Section 3 we introduce another set of sequent calculi, working with formulas in NNF, and we show the equivalence of the two presentations. The NNF presentation allows us to reuse standard results from linear logic. In Sections 4 and 5, we introduce cographs and skew fibrations, so that in section 6 we can finally define combinatorial proofs for relevance logic with and without mingle. Then, in Sections 7–8 we extend our construction to the logic of entailment.

2 Sequent Calculus, Part I

In this section we recall the sequent calculi for the implication-negation-fragment of relevance logic (denoted by $\mathbf{R}_\bar{\cdot}$), of relevance logic with mingle (denoted by $\mathbf{RM}_\bar{\cdot}$), of entailment logic (denoted by $\mathbf{E}_\bar{\cdot}$), and classical propositional logic (denoted by $\mathbf{CL}_\bar{\cdot}$).

For this, we consider the class \mathcal{I} of *formulas* (denoted by A, B, \dots) generated by a countable set $\mathcal{A} = \{a, b, \dots\}$ of *propositional variables* and the connectives \rightarrow and $(\bar{\cdot})$ by the following grammar:

$$A, B ::= a \mid \bar{A} \mid A \rightarrow B \quad (3)$$

A *sequent* Γ in \mathcal{I} is a multiset of occurrence of formulas, written as list and separated by commas: $\Gamma = A_1, \dots, A_n$. We denote by $\bar{\Gamma}^{\rightarrow}$ a sequent of formulas in \mathcal{I} of the form $\bar{A}_1 \rightarrow \bar{B}_1, \dots, \bar{A}_n \rightarrow \bar{B}_n$, and we write $\bar{\Gamma}$ for the sequent obtained from Γ by negating all its formulas, i.e., if $\Gamma = A_1, \dots, A_n$ then $\bar{\Gamma} = \bar{A}_1, \dots, \bar{A}_n$.

In Figure 1 we give the standard sequent systems for the logics $\mathbf{E}_\bar{\cdot}$, $\mathbf{R}_\bar{\cdot}$, $\mathbf{RM}_\bar{\cdot}$, and $\mathbf{CL}_\bar{\cdot}$ as given in [13,5].

Theorem 2.1 *A formula is a theorem of the logic $\mathbf{E}_\bar{\cdot}$ (resp. $\mathbf{R}_\bar{\cdot}$, $\mathbf{RM}_\bar{\cdot}$, $\mathbf{CL}_\bar{\cdot}$) iff it is derivable in the sequent calculus $\mathbf{LE}_\bar{\cdot}$, (resp. $\mathbf{LR}_\bar{\cdot}$, $\mathbf{LRM}_\bar{\cdot}$, $\mathbf{LK}_\bar{\cdot}$). [13]*

Observe that the system $\mathbf{LE}_\bar{\cdot}$ in Figure 1 does contain the cut-rule, whereas the other systems are cut-free. The reason is that due to the form of the \mathbf{EAX} , the cut cannot be eliminated.

In order to obtain cut-free systems for all four logics, we move to the negation normal form in the next section.

			$\frac{}{A, \bar{A}} \text{AX}$	$\frac{}{\bar{A}, A \rightarrow B, B} \text{EAX}$	$\frac{}{\Gamma, \bar{\Gamma}} \text{mAX}$
\mathbf{E}_ε	\mathbf{LE}_ε	AX, EAX, \rightarrow^E , \neg , C, cut	$\frac{\Gamma, A}{\Gamma, \bar{A}} \neg$	$\frac{\Gamma, A \quad \bar{B}, \Delta}{\Gamma, A \rightarrow B, \Delta} \Rightarrow$	$\frac{\Gamma, \bar{A}, B}{\Gamma, A \rightarrow B} \rightarrow$
\mathbf{R}_ε	\mathbf{LR}_ε	AX, \rightarrow , \Rightarrow , \neg , C			
\mathbf{RM}_ε	\mathbf{LRM}_ε	AX, \rightarrow , \Rightarrow , \neg , C, mAX	$\frac{\Gamma, A, A}{\Gamma, A} \text{C}$	$\frac{\bar{\Gamma}^\Rightarrow, \bar{A}, B}{\bar{\Gamma}^\Rightarrow, A \rightarrow B} \rightarrow^E$	$\frac{\Gamma}{\Gamma, A} \text{W}$
\mathbf{CL}_ε	\mathbf{LK}_ε	AX, \rightarrow , \Rightarrow , \neg , C, W		$\frac{\Gamma, A \quad \bar{A}, \Delta}{\Gamma, \Delta} \text{cut}$	

Fig. 1. Rules for the standard sequent systems for the logics \mathbf{E}_ε , \mathbf{R}_ε , \mathbf{RM}_ε , and \mathbf{CL}_ε

3 Sequent Calculus, Part II

In this section we consider formulas in negation normal form (NNF), i.e., the class \mathcal{L} of *formulas* (also denoted by A, B, \dots) generated by the countable set $\mathcal{A} = \{a, b, \dots\}$ of *propositional variables*, their duals $\bar{\mathcal{A}} = \{\bar{a}, \bar{b}, \dots\}$, and the binary connectives \wedge and \vee , via the following grammar:

$$A, B ::= a \mid \bar{a} \mid A \wedge B \mid A \vee B \quad (4)$$

An *atom* is a formula a or \bar{a} with $a \in \mathcal{A}$. As before, a sequent Γ is a multiset of formulas separated by comma. We define *negation* as a function on all formulas in NNF via the De Morgan laws:

$$\bar{\bar{a}} = a \quad \overline{A \wedge B} = \bar{A} \vee \bar{B} \quad \overline{A \vee B} = \bar{A} \wedge \bar{B} \quad (5)$$

There is a correspondence between the class \mathcal{I} defined in the previous section and the class \mathcal{L} of formulas in NNF, defined via the two translations $[\cdot]_{\mathcal{L}}: \mathcal{I} \rightarrow \mathcal{L}$ and $[\cdot]_{\mathcal{I}}: \mathcal{L} \rightarrow \mathcal{I}$:

$$[a]_{\mathcal{L}} = a, \quad [\bar{a}]_{\mathcal{L}} = \overline{[a]_{\mathcal{L}}}, \quad [A \rightarrow B]_{\mathcal{L}} = \overline{[A]_{\mathcal{L}}} \vee [B]_{\mathcal{L}} \quad (6)$$

and

$$[a]_{\mathcal{I}} = a, \quad [\bar{a}]_{\mathcal{I}} = \bar{a}, \quad [A \vee B]_{\mathcal{I}} = \overline{[A]_{\mathcal{I}}} \rightarrow [B]_{\mathcal{I}}, \quad [A \wedge B]_{\mathcal{I}} = \overline{[A]_{\mathcal{I}} \rightarrow [B]_{\mathcal{I}}} \quad (7)$$

Proposition 3.1 *If A is a formula in NNF, then $[[A]_{\mathcal{I}}]_{\mathcal{L}} = A$.*

The proof is straightforward, but in general we do not have $[[B]_{\mathcal{L}}]_{\mathcal{I}} = B$ for formulas in \mathcal{I} , since we can have arbitrary nestings of negation and $[\bar{\bar{B}}]_{\mathcal{L}} = [B]_{\mathcal{L}}$. For this reason, we use here the NNF notation, as it is more concise and carries less redundancy.

We can use this correspondence to translate the sequent systems in Figure 1 into the NNF notation. We go one step further and give cut-free systems \mathbf{LE}' , \mathbf{LR}' , \mathbf{LRM}' , and \mathbf{LK} . They are given in Figure 2, where we denote by Γ^\wedge a sequent of the form $A_1 \wedge B_1, \dots, A_n \wedge B_n$ (i.e., all formulas in Γ^\wedge are conjunctions).

MLL	ax, \vee, \wedge			
MLL _{mix}	ax, \vee, \wedge, mix	$\frac{}{a, \bar{a}} \text{ax}$	$\frac{\Gamma^\wedge, A, B}{\Gamma^\wedge, A \vee B} \vee_E$	$\frac{\Gamma, A, B}{\Gamma, A \vee B} \vee$
MLL _E	ax, \wedge, \vee_E			
LE'	ax, \wedge, \vee_E, C	$\frac{\Gamma, A, A}{\Gamma, A} C$	$\frac{\Gamma, A \quad \Delta, B}{\Gamma, \Delta, A \wedge B} \wedge$	
LR'	ax, \vee, \wedge, C			
LRM'	ax, $\vee, \wedge, C, \text{mix}$	$\frac{\Gamma}{\Gamma, A} W$	$\frac{\Gamma \quad \Delta}{\Gamma, \Delta} \text{mix}$	
LK	ax, \vee, \wedge, C, W			

Fig. 2. The cut-free sequent systems for formulas in NNF

$$\frac{\Gamma\{(A \vee B) \wedge (C \vee D)\}}{\Gamma\{(A \vee C) \vee (B \vee D)\}} m^\downarrow \quad \frac{\Gamma\{A \vee A\}}{\Gamma\{A\}} C^\downarrow \quad \frac{\Gamma\{a \vee a\}}{\Gamma\{a\}} ac^\downarrow \quad \frac{\Gamma\{B\}}{\Gamma\{B \vee A\}} W^\downarrow$$

Fig. 3. The *deep* rules for *medial*, *contraction*, *atomic contraction* and *weakening*.

That figure also defines the linear logic systems MLL_E and MLL that we will need in the course of this paper.

We make also use of the *deep inference* rules in Figure 3 (see also [10,6]), where a *context* $\Gamma\{\}$ is a sequent or a formula, where a hole $\{\}$ takes the place of an atom. We write $\Gamma\{A\}$ when we replace the hole with a formula A .

If S is a sequent system and Γ a sequent, we write $\vdash^S \Gamma$ if Γ is derivable in S . Moreover, if S is a set of inference rules with exactly one premise, we write $\Gamma' \vdash^S \Gamma$ whenever there is a derivation from Γ' to Γ using only rules in S .

Theorem 3.2 *If Γ is a sequent in \mathcal{L} , then*

$$\frac{}{\vdash^{\text{LE}' \cup \{\text{cut}\}} \Gamma} \iff \frac{}{\vdash^{\text{LE}_\exists} [\Gamma]_{\mathcal{I}}}$$

Proof The proof follows the definitions of the two translations $[\cdot]_{\mathcal{I}}$ and $[\cdot]_{\mathcal{L}}$. In fact, C- and cut-rules are the same in the two systems and \rightarrow -rule is equivalent to \vee_E -rule. Moreover, it is trivial to prove by induction that $\vdash^{\text{LE}'} A, \bar{A}$. Finally:

$$\frac{}{\bar{A}, A \rightarrow B, B} \text{EAX} \rightsquigarrow \frac{\frac{}{[\bar{A}]_{\mathcal{L}}, [A]_{\mathcal{L}}} \text{LE}' \parallel \quad \frac{}{[\bar{B}]_{\mathcal{L}}, [B]_{\mathcal{L}}} \text{LE}' \parallel}{[\bar{A}]_{\mathcal{L}}, ([A]_{\mathcal{L}} \wedge [\bar{B}]_{\mathcal{L}}), [B]_{\mathcal{L}}} \wedge_E$$

and

$$\frac{\frac{}{\Gamma, A \quad B, \Delta} \text{LE}' \parallel \quad \frac{}{\Gamma, A \wedge B, \Delta} \wedge_E}{\frac{\frac{}{[\bar{A}]_{\mathcal{I}}, [A]_{\mathcal{I}} \rightarrow [\bar{B}]_{\mathcal{I}}, [B]_{\mathcal{I}}} \text{EAX} \quad \frac{}{[A]_{\mathcal{I}}, [\Gamma]_{\mathcal{I}}} \parallel \text{LE}_\exists}{[\Gamma]_{\mathcal{I}}, [A]_{\mathcal{I}} \rightarrow [\bar{B}]_{\mathcal{I}}, [B]_{\mathcal{I}}} \text{cut} \quad \frac{}{[B]_{\mathcal{I}}, [\Delta]_{\mathcal{I}}} \parallel \text{LE}_\exists}{[\Gamma]_{\mathcal{I}}, [A]_{\mathcal{I}} \rightarrow [\bar{B}]_{\mathcal{I}}, [\Delta]_{\mathcal{I}}} \text{cut}}}$$

□

One important property of the systems in Figure 2 is cut admissibility.

Theorem 3.3 (Cut admissibility) *Let Γ be a sequent in \mathcal{L} , and let S be one of the systems MLL , MLL_{mix} , MLL_{E} , LE' , LR' , LRM' , LK . Then*

$$\frac{S \cup \{\text{cut}\}}{\vdash \Gamma} \iff \frac{S}{\vdash \Gamma}$$

Proof The proof is a standard cut permutation argument. For LK it can already be found in [9] and for all other systems it is the same proof, observing that no reduction step introduces a rule that is not present in the system. \square

The following lemma relates the mix-rule from linear logic to the mingle axiom rule mAX :

Lemma 3.4 *Let S be a sequent system, if Γ is a sequent in \mathcal{L} then*

$$\frac{S \cup \{\text{mAX}\}}{\vdash F} \iff \frac{S \cup \{\text{mix}\}}{\vdash F}$$

Proof First, mAX can be derived using mix :

$$\frac{}{A_1, \dots, A_n, \bar{A}_1, \dots, \bar{A}_n} \text{mAX} \rightsquigarrow \frac{\frac{\frac{S \parallel}{A_1, \bar{A}_1} \quad \frac{S \parallel}{A_2, \bar{A}_2}}{A_1, \bar{A}_1, A_2, \bar{A}_2} \text{mix}}{\frac{\frac{\parallel S \cup \{\text{mix}\}}{A_1, \bar{A}_1, \dots, A_{n-1}, \bar{A}_{n-1}} \quad \frac{S \parallel}{A_n, \bar{A}_n}}{A_1, \bar{A}_1, \dots, A_n, \bar{A}_n} \text{mix}}$$

Conversely, if Γ, Δ is the conclusion of a mix inference,

$$\frac{\frac{\frac{}{a_1, \bar{a}_1} \text{ax} \quad \dots \quad \frac{}{a_n, \bar{a}_n} \text{ax}}{S} \quad \frac{\frac{}{a_{n+1}, \bar{a}_{n+1}} \text{ax} \quad \dots \quad \frac{}{a_{n+m}, \bar{a}_{n+m}} \text{ax}}{S}}{\Gamma, \Delta} \text{mix}$$

it suffices to replace one axiom of the derivation of Γ and one axiom of the derivation of Δ by a single mAX , that is

$$\frac{\frac{\frac{}{a_2, \bar{a}_2} \text{ax} \quad \dots \quad \frac{}{a_{n+m-1}, \bar{a}_{n+m-1}} \text{ax}}{S} \quad \frac{}{a_1, \bar{a}_1, a_{n+m}, \bar{a}_{n+m}} \text{mAX}}{\Gamma, \Delta} \text{S}$$

\square

This is enough to show the equivalence between the systems in Figures 1 and 2.

Theorem 3.5 *If Γ is a sequent in \mathcal{L} , then*

$$\begin{array}{ccc} \frac{}{\vdash \Gamma} \text{LE}' \iff \frac{}{\vdash [\Gamma]_{\mathcal{I}}} \text{LE}_{\mathfrak{s}} & \frac{}{\vdash \Gamma} \text{LR}' \iff \frac{}{\vdash [\Gamma]_{\mathcal{I}}} \text{LR}_{\mathfrak{s}} \\ \frac{}{\vdash \Gamma} \text{LRM}' \iff \frac{}{\vdash [\Gamma]_{\mathcal{I}}} \text{LRM}_{\mathfrak{s}} & \frac{}{\vdash \Gamma} \text{LK} \iff \frac{}{\vdash [\Gamma]_{\mathcal{I}}} \text{LK}_{\mathfrak{s}} \end{array}$$

Proof This follows from Theorems 3.2 and 3.3 and Lemma 3.4, using the definitions of $[\cdot]_{\mathcal{I}}$ and $[\cdot]_{\mathcal{L}}$. \square

Finally, the most important reason to use the systems in Figure 2 instead of the ones in Figure 1 is the following decomposition theorem:

Theorem 3.6 *If Γ is a sequent in \mathcal{L} , then*

$$\begin{array}{ccc} \frac{}{\vdash \Gamma} \text{LE}' \iff \frac{}{\vdash \Gamma'} \text{MLLE} \frac{}{\vdash \Gamma} \text{C}^{\downarrow} & \frac{}{\vdash \Gamma} \text{LR}' \iff \frac{}{\vdash \Gamma'} \text{MLL} \frac{}{\vdash \Gamma} \text{C}^{\downarrow} \\ \frac{}{\vdash \Gamma} \text{LRM}' \iff \frac{}{\vdash \Gamma'} \text{MLL}_{\text{mix}} \frac{}{\vdash \Gamma} \text{C}^{\downarrow} & \frac{}{\vdash \Gamma} \text{LK} \iff \frac{}{\vdash \Gamma'} \text{MLL} \frac{}{\vdash \Gamma} \text{C}^{\downarrow, \text{W}^{\downarrow}} \end{array}$$

Proof The proof is given by rules permutation. It suffices to consider all W- and C-rules as their deep counterpart and move their instance as down as possible in the derivation. Conversely, it suffices to move up all occurrences of C^{\downarrow} and W^{\downarrow} until the context is shallow and then replace them by C and W instances. \square

4 Cographs

A *graph* $\mathcal{G} = \langle V_{\mathcal{G}}, \overset{\mathcal{G}}{\mathcal{E}} \rangle$ is a set $V_{\mathcal{G}}$ *vertices* and a set $\overset{\mathcal{G}}{\mathcal{E}}$ of *edges*, which are two-element subsets of $V_{\mathcal{G}}$. We write $v \overset{\mathcal{G}}{\sim} w$ for $\{v, w\} \in \overset{\mathcal{G}}{\mathcal{E}}$, and we write $v \not\overset{\mathcal{G}}{\sim} w$ if $\{v, w\} \notin \overset{\mathcal{G}}{\mathcal{E}}$. We omit the index/superscript \mathcal{G} when it is clear from the context. When drawing a graph we use $v \text{---} w$ for $v \overset{\mathcal{G}}{\sim} w$. If $v \not\overset{\mathcal{G}}{\sim} w$ and $v \neq w$ we either draw no edge or use $v \cdots w$.

A *cograph* \mathcal{G} is a P4-free graph, i.e. a graph \mathcal{G} with no $u, v, y, z \in V$ such that their induced subgraph has the following shape:⁴

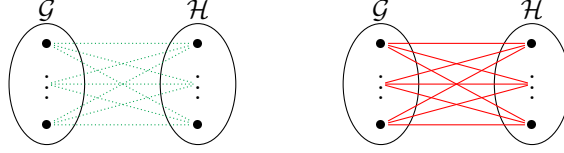


For two disjoint graphs \mathcal{G} and \mathcal{H} , we define their (*disjoint*) *union* $\mathcal{G} \vee \mathcal{H}$ and their *join* $\mathcal{G} \wedge \mathcal{H}$ as follows:

$$\begin{aligned} \mathcal{G} \vee \mathcal{H} &= \langle V_{\mathcal{G}} \cup V_{\mathcal{H}}, \overset{\mathcal{G}}{\mathcal{E}} \cup \overset{\mathcal{H}}{\mathcal{E}} \rangle \\ \mathcal{G} \wedge \mathcal{H} &= \langle V_{\mathcal{G}} \cup V_{\mathcal{H}}, \overset{\mathcal{G}}{\mathcal{E}} \cup \overset{\mathcal{H}}{\mathcal{E}} \cup \{\{u, v\} \mid u \in V_{\mathcal{G}}, v \in V_{\mathcal{H}}\} \rangle \end{aligned} \tag{8}$$

⁴ In the literature, this condition is also called Z-free or N-free.

which can be visualized as follows:



We say that a graph is \mathcal{A} -labeled if each vertex is marked with an atom in $\mathcal{A} \cup \bar{\mathcal{A}}$. We can associate to each formula F in \mathcal{L} an \mathcal{A} -labeled cograph $\llbracket F \rrbracket$ inductively:

$$\llbracket a \rrbracket = \bullet_a, \quad \llbracket \bar{a} \rrbracket = \bullet_{\bar{a}}, \quad \llbracket A \vee B \rrbracket = \llbracket A \rrbracket \vee \llbracket B \rrbracket, \quad \llbracket A \wedge B \rrbracket = \llbracket A \rrbracket \wedge \llbracket B \rrbracket$$

If $\Gamma = A_1, \dots, A_n$ is a sequent of formulas in \mathcal{L} , we define $\llbracket \Gamma \rrbracket = \llbracket A_1 \rrbracket \vee \dots \vee \llbracket A_n \rrbracket$.

The interest in cographs comes from the following two well-known theorems (see, e.g., [8,14]).

Theorem 4.1 *A \mathcal{A} -labeled graphs \mathcal{G} is a cograph iff there is a formula $F \in \mathcal{L}$ such that $\mathcal{G} = \llbracket F \rrbracket$.*

Theorem 4.2 *$\llbracket F \rrbracket = \llbracket F' \rrbracket$ iff F and F' are equivalent modulo associativity and commutativity of \wedge and \vee .*

5 Skew fibrations

Definition 5.1 Let \mathcal{G} and \mathcal{H} be graphs. A *skew fibration* $f: \mathcal{G} \rightarrow \mathcal{H}$ is a mapping from $V_{\mathcal{G}}$ to $V_{\mathcal{H}}$ that preserves \curvearrowright :

- if $u \xrightarrow{\mathcal{G}} v$ then $f(u) \xrightarrow{\mathcal{H}} f(v)$,

and that has the *skew lifting* property:

- if $w \xrightarrow{\mathcal{H}} f(v)$ then there is $u \in V_{\mathcal{G}}$ such that $u \xrightarrow{\mathcal{G}} v$ and $w \xrightarrow{\mathcal{H}} f(u)$.

A skew fibration $f: \mathcal{G} \rightarrow \mathcal{H}$ is *relevant* if it is surjective on vertices and on \curvearrowright :

- for every $w \in V_{\mathcal{H}}$ there is a $u \in V_{\mathcal{G}}$ such that $f(u) = w$, and
- if $w \xrightarrow{\mathcal{H}} t$ then there are $u, v \in V_{\mathcal{G}}$ such that $f(u) = w$ and $f(v) = t$ and $u \xrightarrow{\mathcal{G}} v$.

The purpose of skew fibrations in this setting is to give a combinatorial characterization of derivations containing only contractions and weakenings.

Theorem 5.2 *If Γ, Γ' are sequents in \mathcal{L} then*

1. $\Gamma' \xrightarrow{\text{c}^\downarrow, \text{w}^\downarrow} \Gamma$ iff there is a skew fibration $f: \llbracket \Gamma' \rrbracket \rightarrow \llbracket \Gamma \rrbracket$.
2. $\Gamma' \xrightarrow{\text{c}^\downarrow} \Gamma$ iff there is a relevant skew fibration $f: \llbracket \Gamma' \rrbracket \rightarrow \llbracket \Gamma \rrbracket$.

Proof The first statement has been proved independently in [12] and in [16].

The proof of the second statement is similar, but relevant condition rules out weakening. Let $\Gamma' = \Gamma_0, \Gamma_1, \dots, \Gamma_n = \Gamma$ such that $\frac{\Gamma_i \{A_i \vee A_i\}}{\Gamma_{i+1} = \Gamma_i \{A_i\}} \mathbf{C}\downarrow$. By induction over the size of A_i , there is a relevant skew fibration $f_i: \Gamma_i \rightarrow \Gamma_{i+1}$ for each $i \in \{0, \dots, n-1\}$ and the composition of such f_i is still a relevant skew fibration. Conversely, in case of f a relevant skew fibration, the lifting property becomes the following:

– if $w \xrightarrow{\mathcal{H}} f(v)$ then there is $u \in V_{\mathcal{G}}$ such that $u \xrightarrow{\mathcal{G}} v$ and $f(u) = w$.

which, by induction over Γ , allows to prove that:

- if $\llbracket \Gamma \rrbracket = \bullet_a$ then $\llbracket \Gamma' \rrbracket = \bullet_a \vee \dots \vee \bullet_a$;
- if $\llbracket \Gamma \rrbracket = \mathcal{G} \vee \mathcal{H}$ then $\llbracket \Gamma' \rrbracket = \mathcal{G}' \vee \mathcal{H}'$ with $f(\mathcal{G}') = \mathcal{G}$ and $f(\mathcal{H}') = \mathcal{H}$;
- if $\llbracket \Gamma \rrbracket = \mathcal{G} \wedge \mathcal{H}$ then either $\llbracket \Gamma' \rrbracket = \mathcal{G}' \wedge \mathcal{H}'$ with $f(\mathcal{G}') = \mathcal{G}$ and $f(\mathcal{H}') = \mathcal{H}$, or $\llbracket \Gamma' \rrbracket = (\mathcal{G}'_1 \wedge \mathcal{H}'_1) \vee \dots \vee (\mathcal{G}'_n \wedge \mathcal{H}'_n)$ with $f(\mathcal{G}'_i) = \mathcal{G}$ and $f(\mathcal{H}'_i) = \mathcal{H}$ for each $i \in \{1, \dots, n\}$.

These decompositions guide the definition of a derivation $\Gamma' \xrightarrow{\mathbf{C}\downarrow} \Gamma$. □

6 RB-cographs and combinatorial proofs

In this section we finally define combinatorial proofs. For this we use Retoré's RB-cographs [15]:

Definition 6.1 ([15]) An *RB-cograph* is a tuple $\mathcal{G} = \langle V_{\mathcal{G}}, \xrightarrow{\mathcal{G}}, \underset{\mathcal{G}}{\vee} \rangle$ where $\mathcal{G}^R = \langle V_{\mathcal{G}}, \xrightarrow{\mathcal{G}} \rangle$ is a cograph and $\underset{\mathcal{G}}{\vee}$ a irreflexive, symmetric binary relation such that for every $v \in V_{\mathcal{G}}$ there is a unique $w \in V_{\mathcal{G}}$ with $v \underset{\mathcal{G}}{\vee} w$.

As done in (1) in the introduction, we use $v \text{---} w$ for $v \xrightarrow{\mathcal{G}} w$, and $v \text{---} w$ for $v \underset{\mathcal{G}}{\vee} w$ when drawing an RB-cograph.

Definition 6.2 ([15]) If u and v are two vertices of a RB-cograph, an *alternating elementary path* (*æ-path*) from x_0 to x_n is a sequence of pairwise disjoint vertices $x_0, \dots, x_n \in V$ such that either $x_0 \xrightarrow{\mathcal{G}} x_1 \underset{\mathcal{G}}{\vee} x_2 \xrightarrow{\mathcal{G}} x_3 \underset{\mathcal{G}}{\vee} x_4 \dots x_n$ or $x_0 \underset{\mathcal{G}}{\vee} x_1 \xrightarrow{\mathcal{G}} x_2 \underset{\mathcal{G}}{\vee} x_3 \xrightarrow{\mathcal{G}} x_4 \dots x_n$. An *æ-cycle* is an æ-path of even length with $x_0 = x_n$. A *chord* of æ-path x_0, \dots, x_n is an edge $x_i \xrightarrow{\mathcal{G}} x_j$ with $i+1 < j$. The æ-path is *chordless* if it has no chord. A RB-cograph is *æ-connected* if there is a chordless æ-path between each pair of vertices \mathcal{G} and it is *æ-acyclic* if there are no chordless æ-cycle.

Theorem 6.3 ([15]) *If Γ is a sequent over \mathcal{L} then*

1. $\xrightarrow{\text{MLL}} \Gamma \iff$ there is an æ-connected, æ-acyclic RB-cograph \mathcal{G} with $\mathcal{G}^R = \llbracket \Gamma \rrbracket$
2. $\xrightarrow{\text{MLL}_{\text{mix}}} \Gamma \iff$ there is an æ-acyclic RB-cograph \mathcal{G} with $\mathcal{G}^R = \llbracket \Gamma \rrbracket$

We say that a map f from an RB-cograph \mathcal{C} to a \mathcal{A} -labeled cograph is *axiom preserving* if for all u, v with $u \overset{\mathcal{C}}{\vee} v$ we have that $f(u)$ and $f(v)$ are labeled by two dual atoms.

Definition 6.4 Let Γ be a sequent over \mathcal{L} .

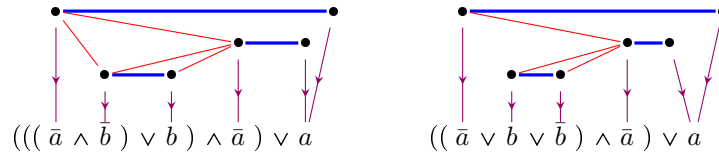
1. A *combinatorial LK-proof* of Γ is an axiom-preserving skew fibration $f: \mathcal{C} \rightarrow \llbracket \Gamma \rrbracket$ where \mathcal{C} is an \mathfrak{a} -connected, \mathfrak{a} -acyclic RB-cograph.
2. A *combinatorial LR'-proof* of Γ is an axiom-preserving relevant skew fibration $f: \mathcal{C} \rightarrow \llbracket \Gamma \rrbracket$ where \mathcal{C} is an \mathfrak{a} -connected, \mathfrak{a} -acyclic RB-cograph.
3. Finally, a *combinatorial LRM'-proof* of Γ is an axiom-preserving relevant skew fibration $f: \mathcal{C} \rightarrow \llbracket \Gamma \rrbracket$ where \mathcal{C} is an \mathfrak{a} -acyclic RB-cograph.

Theorem 6.5 Let Γ be a sequent over \mathcal{L} , and let $S \in \{\text{LR}', \text{LRM}', \text{LK}\}$. Then

$$\overset{S}{\vdash} \Gamma \iff \text{there is a combinatorial } S\text{-proof of } \Gamma .$$

Proof This follows from Theorems 3.6, 5.2 and 6.3. For LK this has already been shown in [11,12,16,17]. \square

Below are a combinatorial LR'-proof (on the left) and a combinatorial LRM'-proof (on the right):



Theorem 6.6 Let Γ be a sequent and \mathcal{G} a graph together with a perfect matching, and let f be a map from \mathcal{G} to $\llbracket \Gamma \rrbracket$. It can be decided in polynomial time in $|V_{\mathcal{G}}| + |\Gamma|$ whether $f: \mathcal{G} \rightarrow \llbracket \Gamma \rrbracket$ is a combinatorial LR'-proof (resp. a combinatorial LRM'-proof).

Proof All necessary properties can be checked in polynomial time. \square

7 Sequent Calculus, Part III

In the remainder of the paper, we extend our results to the *entailment logic* \mathbf{E}_5 . The reason why we need a separate treatment is due to some intrinsic technical drawbacks occurring in the LE' sequent calculus. The first is that commas used to separate formulas in a sequent can not be interpreted as disjunction, as we usually do in classical logic. Using the *display calculi* [4] terminology, in LE' the comma is extensional while \vee and \wedge are intensional. Moreover, \wedge and \vee are not associative and this give birth to unusual behaviors. For example $(A \vee A) \vee (\bar{A} \wedge \bar{A})$ is provable in LE' while $A \vee (A \vee (\bar{A} \wedge \bar{A}))$ is not.

$$\frac{\text{MLL}_{\dot{\mathcal{E}}}}{\text{LE}^{\bullet}} \left| \frac{\text{ax}, \wedge_{\dot{\mathcal{E}}}, \vee_{\dot{\mathcal{E}}}}{\text{ax}, \wedge_{\dot{\mathcal{E}}}, \vee_{\dot{\mathcal{E}}}, \mathbf{C}_{\dot{\mathcal{E}}}} \quad \frac{\text{ax}}{a, \bar{a}} \quad \frac{\Gamma, A^n \quad B^m, \Delta}{\Gamma, A^n \wedge B^m, \Delta} \wedge_{\dot{\mathcal{E}}} \quad \frac{\Gamma \wedge, A^n, B^m}{\Gamma \wedge, A^n \vee B^m} \vee_{\dot{\mathcal{E}}} \quad \frac{\Gamma, A^n}{\Gamma, A} \mathbf{C}_{\dot{\mathcal{E}}}$$

Fig. 4. The cut-free systems $\text{MLL}_{\dot{\mathcal{E}}}$ and LE^{\bullet}

We first introduce the class of *entailed formulas* \mathcal{E} which are generated by a countable set $\mathcal{A} = \{a, b, \dots\}$ of *propositional variables* and the following grammar:

$$A, B ::= a \mid \bar{a} \mid A \wedge B \mid A \vee B \mid A^n \quad (9)$$

where $n > 0$. Moreover, we consider the sequents $\Gamma\{A^{n+1}\}$ and $\Gamma\{A, A^n\}$ to be equal. In other words, A^n has to be thought of as an abbreviation for the sequent A, \dots, A (n copies of A) that is allowed to occur as a subformula in a formula. We define the sequent systems $\text{MLL}_{\dot{\mathcal{E}}}$ and LE^{\bullet} on entailed formulas given by the rules in Figure 4.

Theorem 7.1 *If Γ is a sequent over \mathcal{L} then*

$$\frac{\text{LE}'}{\vdash} \Gamma \iff \frac{\text{LE}^{\bullet}}{\vdash} \Gamma$$

Proof It suffices to remark that LE^{\bullet} rules behave as LE' rules on standard NNF-formulas. \square

Let $\mathbf{C}_{\dot{\mathcal{E}}}^{\downarrow}$ be the deep inference rule $\frac{F\{A^n\}}{F\{A\}} \mathbf{C}_{\dot{\mathcal{E}}}^{\downarrow}$. Then we have a result similar to Theorem 3.6.

Theorem 7.2 *If F is a formula in \mathcal{E} then*

$$\frac{\text{LE}^{\bullet}}{\vdash} F \iff \frac{\text{MLL}_{\dot{\mathcal{E}}}}{\vdash} F' \frac{\mathbf{C}_{\dot{\mathcal{E}}}^{\downarrow}}{\vdash} \Gamma$$

Proof By rule permutations, similarly to the proof of Theorem 3.6. \square

8 Weighted cographs and fibrations

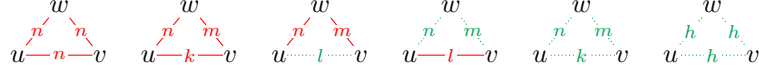
Definition 8.1 A *weighted graph* $\mathcal{G} = \langle V_{\mathcal{G}}, \overset{\mathcal{G}}{\curvearrowright}, \delta \rangle$ is a given by graph $\langle V_{\mathcal{G}}, \overset{\mathcal{G}}{\curvearrowright} \rangle$ together with a *weight function* $\delta: V_{\mathcal{G}} \times V_{\mathcal{G}} \rightarrow \mathbb{N}$ such that if $u \overset{\mathcal{G}}{\curvearrowright} v$ then $\delta(u, v) > 0$ and $\delta(u, u) = 0$.

We use the following notations: we write $u \overset{\mathcal{G}}{\curvearrowright}_k v$ iff $u \overset{\mathcal{G}}{\curvearrowright} v$ and $\delta(u, v) = k$, and we write $u \overset{\mathcal{G}}{\not\curvearrowright}_k v$ iff $u \not\overset{\mathcal{G}}{\curvearrowright} v$ and $\delta(u, v) = k$. When drawing a graph we use $v \overset{-k}{\curvearrowright} w$ for $v \overset{\mathcal{G}}{\curvearrowright}_k w$ and we use $v \overset{\dots k \dots}{\curvearrowright} w$ for $v \overset{\mathcal{G}}{\curvearrowright}_k w$. If $v \overset{0}{\curvearrowright} w$ we often draw no edges.

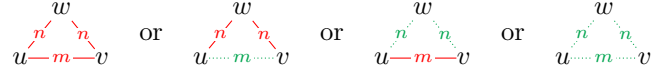
Definition 8.2 A *weighted cograph* is a weighted graph such that:

1. the graphs $\langle V_{\mathcal{G}}, \overset{\mathcal{G}}{\curvearrowright}_i \rangle$ and $\langle V_{\mathcal{G}}, \overset{\mathcal{G}}{\curvearrowright}_i \rangle$ are Z-free for all $i \neq 0$;

2. for all $u, v, w \in V_G$, and any $n, m, k, l, h \in \mathbb{N}$, with n, m, k being pairwise distinct and $h > 0$, the following configurations are forbidden:



3. for all $u, v, w \in V_G$ with



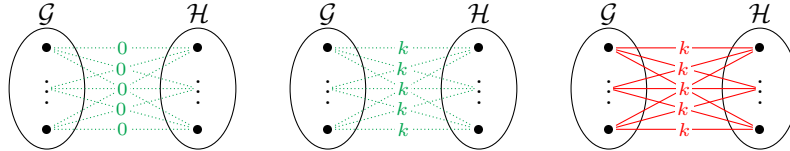
either $n = 0$ or $m = 0$ or $n > m$.

Remark 8.3 A cograph is a weighted cograph \mathcal{G} with $\delta : V_G \times V_G \rightarrow \{0, 1\}$.

We define the *juxtaposition*, *graded union* and *graded join* operations:

$$\begin{aligned} \mathcal{G} \star \mathcal{H} &= \langle V_G \cup V_{\mathcal{H}}, \overline{\mathcal{G}} \cup \overline{\mathcal{H}}, \delta_{\mathcal{G}} \cup \delta_{\mathcal{H}} \cup \delta^0 \rangle \\ \mathcal{G} \vee \mathcal{H} &= \langle V_G \cup V_{\mathcal{H}}, \overline{\mathcal{G}} \cup \overline{\mathcal{H}}, \delta_{\mathcal{G}} \cup \delta_{\mathcal{H}} \cup \delta^{\smile} \rangle \\ \mathcal{G} \wedge \mathcal{H} &= \langle V_G \cup V_{\mathcal{H}}, \overline{\mathcal{G}} \cup \overline{\mathcal{H}} \cup \{\{u, v\} \mid u \in \mathcal{G}, v \in \mathcal{H}\}, \delta_{\mathcal{G}} \cup \delta_{\mathcal{H}} \cup \delta^{\frown} \rangle \end{aligned}$$

where δ^0 is the weight function which assigns to each $(u, v) \in V_G \times V_{\mathcal{H}} \cup V_{\mathcal{H}} \times V_G$ the weight 0, while δ^{\frown} (resp. δ^{\smile}) is the weight function which assigns to each $(u, v) \in V_G \times V_{\mathcal{H}} \cup V_{\mathcal{H}} \times V_G$ the weight $k = 1 + \max\{\delta(w, z) \mid w \overline{\mathcal{G}} z \text{ or } w \overline{\mathcal{H}} z\}$ (respectively $k = 1 + \max\{\delta(w, z) \mid w \overline{\mathcal{G}} z \text{ or } w \overline{\mathcal{H}} z\}$). We represent these operations as follows:



We associate to each entailed formula F (sequent Γ) a graded relation web:

$$\begin{aligned} \llbracket a \rrbracket &= \bullet_a, & \llbracket \bar{a} \rrbracket &= \bullet_{\bar{a}}, & \llbracket A \vee B \rrbracket &= \llbracket A \rrbracket \vee \llbracket B \rrbracket, \\ \llbracket A \wedge B \rrbracket &= \llbracket A \rrbracket \wedge \llbracket B \rrbracket, & \llbracket A, B \rrbracket &= \llbracket A \rrbracket \star \llbracket B \rrbracket \end{aligned}$$

Two weighted graphs \mathcal{G} and \mathcal{H} are *isomorphic* (denoted $\mathcal{G} \simeq \mathcal{H}$) if there is a bijection ϕ between V_G and $V_{\mathcal{H}}$ which preserves edges and weights order, that is $u \overline{\mathcal{G}} v$ iff $\phi(u) \overline{\mathcal{H}} \phi(v)$, and $\delta(u, v) > \delta(u', v')$ iff $\delta(\phi(u), \phi(v)) > \delta(\phi(u'), \phi(v'))$. Then Theorem 4.1 can be extended to the following:

Theorem 8.4 A \mathcal{A} -labeled weighted graph \mathcal{G} is a weighted cograph iff there is a sequent Γ of entailed formulas such that $\mathcal{G} \simeq \llbracket \Gamma \rrbracket$.

Proof The proof is similar to the one of Theorem 4.1. However, the condition $\mathcal{G} \simeq \llbracket \Gamma \rrbracket$ (instead of $\mathcal{G} = \llbracket \Gamma \rrbracket$) is due to the existence of weighted cographs not of the form $\llbracket \Gamma \rrbracket$. By means of example take $a-2-b \simeq \llbracket a \wedge b \rrbracket = a-1-b$. \square

Definition 8.5 A *weighted skew fibration* $f: \mathcal{G} \rightarrow \mathcal{H}$ is a skew fibration between weighted graphs that preserves the weights.

Note that this means in particular that $f(u) = f(v)$ implies that $\delta(u, v) = 0$.

Theorem 8.6 Let Γ and Γ' be sequents over \mathcal{E} . Then $\Gamma' \xrightarrow{\text{C}_E^\downarrow} \Gamma$ iff there is a weighted relevant skew fibration $f: \llbracket \Gamma' \rrbracket \rightarrow \llbracket \Gamma \rrbracket$.

Proof The proof is similar to the one for (non-weighted) relevant skew fibrations. First, let $\Gamma' = \Gamma_0, \Gamma_1, \dots, \Gamma_n = \Gamma$ be a sequence of sequents such that

$$\frac{\Gamma_i \{A_i, A_i\}}{\Gamma_{i+1} = \Gamma_i \{A_i\}} \text{C}_E^\downarrow .$$

By definition of juxtaposition, join and union cograph operations we have that $f_i: \llbracket \Gamma_i \rrbracket \rightarrow \llbracket \Gamma_{i+1} \rrbracket$ is a relevant skew fibration and preserves \curvearrowright and weights. Then also $f = f_{n-1} \circ \dots \circ f_0$ is a weighted relevant skew fibration.

The converse follows by remarking that $f(u) = f(v)$ iff $u \curvearrowright_0 v$. \square

9 Weighted RB-cographs

Definition 9.1 A *weighted RB-cograph* is a tuple $\mathcal{G} = \langle V_{\mathcal{G}}, \overset{\mathcal{G}}{\curvearrowright}, \delta_{\mathcal{G}}, \underset{\mathcal{G}}{\curvearrowleft} \rangle$ where:

- $\mathcal{G}^{R\delta} = \langle V_{\mathcal{G}}, \overset{\mathcal{G}}{\curvearrowright}, \delta_{\mathcal{G}} \rangle$ is a weighted cograph;
- $\underset{\mathcal{G}}{\curvearrowleft}$ is a perfect matching on $V_{\mathcal{G}}$;

A weighted RB-cograph $\mathcal{G} = \langle V_{\mathcal{G}}, \overset{\mathcal{G}}{\curvearrowright}, \delta_{\mathcal{G}}, \underset{\mathcal{G}}{\curvearrowleft} \rangle$ is *\mathfrak{x} -connected* (*\mathfrak{x} -acyclic*) if the RB-cograph $\mathcal{G}^{RB} = \langle V_{\mathcal{G}}, \overset{\mathcal{G}}{\curvearrowright}, \underset{\mathcal{G}}{\curvearrowleft} \rangle$ is an \mathfrak{x} -connected (\mathfrak{x} -acyclic) RB-cograph. A weighted RB-cograph is *entailed* if it is \mathfrak{x} -connected, \mathfrak{x} -acyclic, and satisfies the following condition:

- if $a, b, c \in V$ such that $a \curvearrowright_m b$ for $m > 0$, and $c \curvearrowright_n a$ and $c \curvearrowright_n b$, with $n > m$ or $n = 0$, then there is $d \in V$ such that



Theorem 9.2 If Γ is a sequent of entailed formulas then

$$\xrightarrow{\text{MLL}_E^\bullet} \Gamma \iff \text{there is an entailed weighted RB-cograph } \mathcal{G} \text{ with } \mathcal{G}^{R\delta} = \llbracket \Gamma \rrbracket$$

Proof The proof piggybacks on Retoré's sequentialization proof [15]. Each proof in MLL_E induces the construction of an entailed weighted cograph \mathcal{G} by the

$$\frac{\langle \llbracket a \rrbracket \star \llbracket \bar{a} \rrbracket \mid \vee = \{\{a, \bar{a}\}\} \rangle \llbracket \text{ax} \rrbracket}{\langle \llbracket \Gamma \wedge \rrbracket \star \llbracket A \rrbracket \star \llbracket B \rrbracket \mid \vee \rangle} \llbracket \vee \dot{\varepsilon} \rrbracket} \frac{\langle \llbracket \Gamma \wedge \rrbracket \star (\llbracket A \rrbracket \vee \llbracket B \rrbracket) \mid \vee \rangle}{\langle \llbracket \Gamma \rrbracket \star \llbracket A \rrbracket \mid \overset{\Gamma, A}{\vee} \rangle \quad \langle \llbracket B \rrbracket \star \llbracket \Delta \rrbracket \mid \overset{B, \Delta}{\vee} \rangle} \llbracket \wedge \dot{\varepsilon} \rrbracket}$$

Fig. 5. Construction rules for entailed weighted RB-cographs.

operations shown in Figure 5. In fact, each of these operations preserves ε -connectedness, ε -acyclicity and entailment conditions.

Conversely, let Γ be the sequent such that $\llbracket \Gamma \rrbracket = \mathcal{G}^{R\delta}$ and let F_Γ be the formula in \mathcal{L} obtained by substituting each comma occurring in Γ by a \vee . By Theorem 6.3 we have derivation π_{MLL} of F_Γ in MLL. We construct a derivation π_{LE}^\bullet of Γ in $\text{MLL}_\varepsilon^\bullet$ by induction over the rules in π_{MLL} :

- If the last rule in $\pi_{\mathcal{L}}$ is an ax -rule, then the last rule in π_{LE}^\bullet is a ax -rule;
- If the last rule in $\pi_{\mathcal{L}}$ is an \vee -rule of the form $\frac{\Gamma, A, B}{\Gamma, A \vee B} \vee$, then $\delta(a, b) = \delta(a', b')$ for all $a, a' \in V_{\llbracket A \rrbracket}$ and $b, b' \in V_{\llbracket B \rrbracket}$. If $\delta(u, v) = 0$ we skip this rule inference in the construction of π_{LE}^\bullet (the \vee introduced by this rule in F_Γ is a comma in Γ). Otherwise, the last rule in π_{LE}^\bullet is a $\vee \dot{\varepsilon}$ -rule. In fact, for each $c \in V_{\llbracket \Gamma \rrbracket}$ by entailment condition there are $d \in V_{\llbracket \Gamma \rrbracket}$ such that $c \frown d$; that is $\Gamma = \Gamma \wedge$.
- If the last rule in $\pi_{\mathcal{L}}$ is an \wedge -rule, then the last rule in π_{LE}^\bullet is a $\wedge \dot{\varepsilon}$ -rule. \square

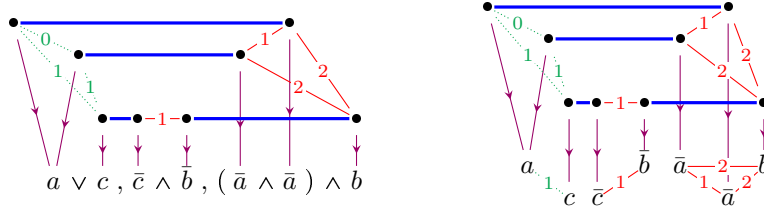
Definition 9.3 A combinatorial LE' -proof of a sequent Γ in \mathcal{L} is given by an axiom-preserving weighted relevant skew fibration $f: \mathcal{C}^{R\delta} \rightarrow \llbracket \Gamma \rrbracket$ where \mathcal{C} is an entailed weighted RB-cograph.

Theorem 9.4 Let Γ be a sequent in \mathcal{L} then

$$\frac{\text{LE}'}{\vdash \Gamma} \iff \text{there is a combinatorial LE}'\text{-proof } f: \mathcal{C}^{R\delta} \rightarrow \llbracket \Gamma \rrbracket$$

Proof This follows from Theorems 7.1, 7.2, 8.6 and 9.2. \square

Below is an example of a combinatorial LE' -proof. On the left the conclusion is shown as sequent, and on the right as weighted cograph.



Theorem 9.5 Let Γ be a sequent and \mathcal{G} a graph together with a perfect matching and a weight function on its edges, and let f be a map from \mathcal{G} to $\llbracket \Gamma \rrbracket$. It can be decided in polynomial time in $|V_{\mathcal{G}}| + |\Gamma|$ whether $f: \mathcal{G} \rightarrow \llbracket \Gamma \rrbracket$ is a combinatorial LE' -proof.

Proof All necessary properties (forbidden edges configurations for \mathcal{G} being a weighted cograph, \ae -connectedness and \ae -acyclicity, and f being a weighted relevant skew fibration) can be checked in polynomial time. \square

10 Conclusion

In this paper we presented combinatorial proofs for entailment logic $\mathbf{E}_\text{\ae}$, classical relevant logics $\mathbf{R}_\text{\ae}$ and classical relevant logic with mingle $\mathbf{RM}_\text{\ae}$. In some sense, combinatorial proof for entailment logic can be considered as a case study for logics with commutative but not associative connectives.

In fact, this paper can be seen as a small step in a larger research project showing that combinatorial proofs are a uniform, modular and bureaucratic-free way of representing proofs for a large class of logics. Apart from the logics studied in this paper, this goal has been achieved for multiplicative linear logic with and without mix in [15], for classical propositional logic in [11,12,17], and for intuitionistic propositional logic in [?]. For first-order logic, modal logics, and larger fragments of linear logic, this is work in progress.

A necessary condition for a logic to have combinatorial proofs seems to be the ability to separate the multiplicative (linear) fragment from the additive (contraction+weakening) fragment. This can happen inside some form of deep inference proof system [6,10], and is realized in this paper in Theorems 3.6 and 7.2.

A crucial condition that combinatorial proofs should obey, in order to be called *combinatorial proofs* for a chosen logic, is that all combinatorial properties needed for correctness of a given proof object can be checked in polynomial time with respect to its size. Then combinatorial proofs form a proof system (in the sense of Cook and Reckhow [7]) for the chosen logic. The combinatorial proofs we give in this paper have this property.

Thanks to their combinatorial (or bureaucracy-free) nature, combinatorial proofs allow us to capture a less coarser notion of proof identity with respect to the one given by syntactic formalisms like sequent calculus and analytic tableaux. Following the work in [12,1,?] we put forward the following notion of proof identity:

Two proofs are the same iff they have the same combinatorial proof.

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