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Sparse supernodal solver with low-rank compression for solving the frequency-domain Maxwell equations discretized by a high order HDG method

November 20th, 2017 - JJC

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Introduction

PaStiX: Sparse Direct Solver

- Originally designed for distributed heterogeneous architectures
- Recent works introduced low-rank compression to reduce the burden on time and memory complexities

HORSE: HDG solver

- Solve a reduced system on faces instead of elements
- Lead to a sparse system smaller than the one using CG or DG
- · Domain-Decomposition is then used
- Use a sparse direct solver on each subdomain

Objective for this talk: evaluate the impact of using low-rank compression to enhance subdomains' factorization in HORSE.



Outline

- 1. Background on Sparse Direct Solvers
- 2. Block Low-Rank Operations
- 3. Numerical Experiments
- 4. Case study: PASTIX inside a HDG method



Background on Sparse Direct Solvers



PaStiX: a sparse direct solver

Current sparse direct solvers for 3D problems

- $\Theta(n^2)$ time complexity
- $\Theta(n^{\frac{4}{3}})$ memory complexity
- BLAS 3 operations: efficient operations

Block Low-Rank solver

- In many applications, the machine precision is not required
- Minimal Memory strategy allows to save memory
- Just-In-Time strategy allows to reduce time-to-solution

Objective: build an algebraic low-rank solver following the supernodal approach of PASTIX, with block-data structures



Sparse Direct Solvers approach

Problem: solve Ax = b

- Cholesky: factorize $A = LL^t$ (symmetric pattern $(A + A^t)$ for LU)
- Solve Ly = b
- Solve $L^t x = y$

PaStiX steps

- 1. Order unknowns to minimize the fill-in
- 2. Compute a symbolic factorization to build L structure
- **3.** Factorize the matrix in place on L structure
- 4. Solve the system with forward and backward triangular solves



Ordering with Nested Dissection

Entry graph $G = (V, E, \sigma_p)$

V: vertices, E: edges, σ_p : unknowns permutation

Algorithm to compute σ_p

- **1.** Partition $V = A \cup B \cup C$
- **2.** Order C with larger numbers: $V_A < V_C$ and $V_B < V_C$
- **3.** Apply the process recursively on A and B



Three-levels of nested dissection on a regular cube.

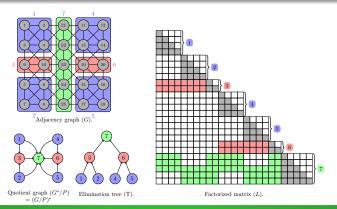
Metis and Scotch are mostly used to perform this partitioning.



Symbolic Factorization

General approach

- 1. Build a partition with the nested dissection process
- 2. Compress information on data blocks
- 3. Compute the block elimination tree using the block quotient graph

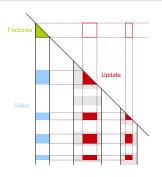




Numerical Factorization

Algorithm to eliminate the column block *k*

- 1. Factorize the diagonal block (POTRF/GETRF)
- 2. Solve off-diagonal blocks in the current column (TRSM)
- 3. Update the trailing matrix with the column's contribution (GEMM)



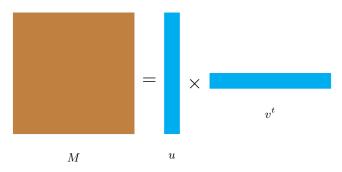


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Block Low-Rank Operations



Low-Rank Compression

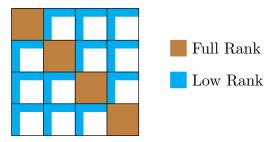


$$M \in \mathbb{R}^{n \times n}$$
; $u, v \in \mathbb{R}^{n \times r}$

Memory Consumption reduced: 2nr instead of n^2 Several methods can be used: SVD, RRQR, BDLR, ACA... The objective is to compute \hat{A} such that $||A - \hat{A}||_2 < \tau$



Block Low-Rank Compression for a dense matrix

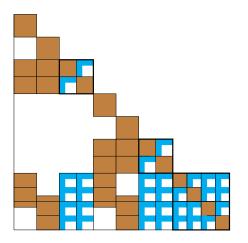


Some related works for dense matrices

- LSTC
- · Airbus (hmat-oss)
- CEA



Block Low-Rank Compression – Symbolic Factorization



Large off-diagonal are low-rank, in the form uv^t



Some Related Works

Sparse Block Low-Rank solver: BLR-MUMPS

- They introduced an approach to reduce the time-to-solution
- The main difference is the difference between multifrontal and supernodal

Some other approaches for sparse

- Strumpack HSS by Ghysels et al. uses randomized sampling to perform an efficient extend-add
- HODLR by Darve et al. uses pre-selection of rows and columns (BDLR)
- \bullet $\,\mathcal{H}\text{-LU}$ by Hackbusch et al. does not fully exploit the symbolic factorization
- HODLR/HSS like method by Chadwick et al. uses estimation of ranks and the geometry of the problem





Block Low-Rank Algorithm

Approach

- Large supernodes are partitioned into a set of smaller supernodes
- · Large off-diagonal blocks are represented as low-rank blocks

Operations

- Diagonal blocks are dense
- TRSM are performed on low-rank off-diagonal blocks
- GEMM are performed between low-rank off-diagonal blocks. It creates contributions to dense or low-rank blocks: this is the extend-add problem

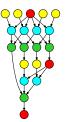
Compression techniques

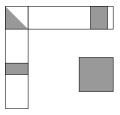
- SVD, RRQR for now
- Possible extension to any algebraic method: ACA, randomized techniques...



- 1. Eliminate each column block
 - 1.1 Factorize the dense diagonal block Compress off-diagonal blocks belonging to the supernode
 - 1.2 Apply a TRSM on LR blocks (cheaper)
 - 1.3 LR update on dense matrices (LR2GE extend-add)
- 2. Solve triangular systems with low-rank blocks



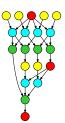


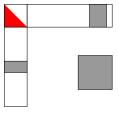




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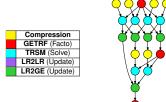


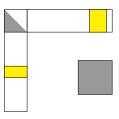






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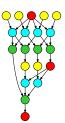


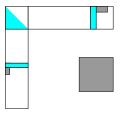




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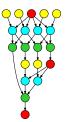


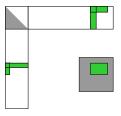




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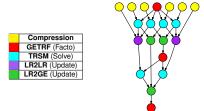


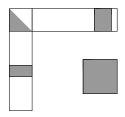






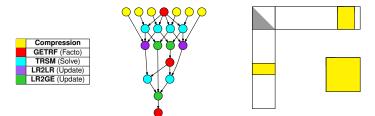
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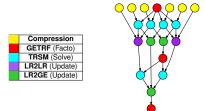


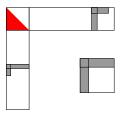
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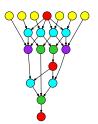


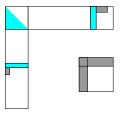




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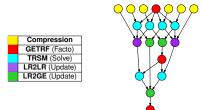


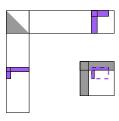






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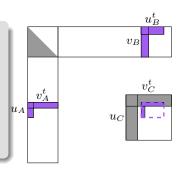






Focus on the LR2LR kernel: build the $u_{AB}v_{AB}^t$ contribution

- Update of C with contribution from blocks A and B
- $\hbox{ The low-rank matrix } u_{AB}v_{AB}^t \\ \hbox{ is added to } u_Cv_C^t \\$
- $u_{AB}v_{AB}^t = (u_A(v_A^t v_B))u_B^t$ or $u_{AB}v_{AB}^t = u_A((v_A^t v_B)u_B^t)$
- Eventually, recompression of $(v_A^t v_B)$





Focus on the LR2LR kernel: add the contribution into $u_C v_C^t$

A low-rank structure $u_C v_C^t$ receives a low-rank contribution $u_{AB} v_{AB}^t$.

Algorithm

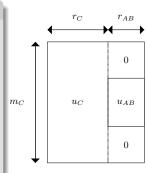
$$A = u_C v_C^t + u_{AB} v_{AB}^t = \left(\left[u_C, u_{AB} \right] \right) \times \left(\left[v_C, v_{AB} \right] \right)^t$$

• QR:
$$[u_C, u_{AB}] = Q_1 R_1 \Theta(m(r_C + r_{AB})^2)$$

• QR:
$$[v_C, v_{AB}] = Q_2 R_2 \quad \Theta(n(r_C + r_{AB})^2)$$

• SVD:
$$R_1 R_2^t = u \sigma v^t \quad \Theta((r_C + r_{AB})^3)$$

$$A = (Q_1 u \sigma) \times (Q_2 v)^t$$



Optimizations using RRQR to reduce complexity to $\Theta(n(r_C+r_{AB})r_C^*)$



Comparison of both strategies

Memory consumption

- Minimal Memory strategy really saves memory
- Just-In-Time strategy reduces the size of L' factors, but supernodes are allocated dense at the beginning: no gain in pure right-looking

Low-Rank extend-add

- Minimal Memory strategy requires expensive extend-add algorithms to update (recompress) low-rank structures with the LR2LR kernel
- Just-In-Time strategy continues to apply dense update at a smaller cost through the LR2GE kernel





3

Numerical Experiments



Experimental setup

Machine: 2 INTEL Xeon E5-2680v3 at 2.50 GHz

- 128 GB
- 24 cores

3D Matrices from The SuiteSparse Matrix Collection

- Audi: structural problem (943 695 DoFs)
- Atmosmodj: atmospheric model (1 270 432 DoFs)
- Geo1438: geomechanical model of earth (1 437 960 DoFs)
- Hook: model of a steel hook (1 498 023 DoFs)
- Serena: gas reservoir simulation (1 391 349 DoFs)
- + Laplacian: Poisson problem (7-points stencil)

Parallelism is obtained with PASTIX static scheduling for multi-threaded architectures



Parameters

Entry parameters

- Tolerance τ : absolute parameter (normalized for each block)
- Compression method: SVD or RRQR
- Compression strategy: Minimal Memory or Just-In-Time
- Blocking sizes: between 128 and 256 in following experiments
- Compression sizes: 128 for width and 20 for height

Strategy Minimal Memory

- Blocks are compressed at the beginning
- Each contribution implies a recompression

Strategy Just-In-Time

- · Blocks are compressed just before a supernode is eliminated
- Those blocks are never uncompressed



	Complete	Just-In-Time RRQR	Minimal Memory RRQR
Factorization time (s)			
Compression	-	3.4e+01	5.6+00
Block factorization (GETRF)	7.2e-01	7.3e-01	7.6e-01
Panel solve (TRSM)	1.7e+01	7.4e+00	7.9e+00
Update			
Formation of contribution	-	-	4.2e+01
Addition of contribution	-	-	7.3e+02
Dense udpate (GEMM)	4.6e+02	9.7e+01	2.4e+01
Total	4.7e+02	1.4e+02	8.1e+02
Solve time (s)	6.3e+00	3.0e+00	3.2e+00
Factors final size (GB)	16.3	7.49	7.31
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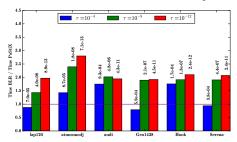
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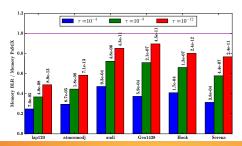
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Behaviour of RRQR/Minimal Memory



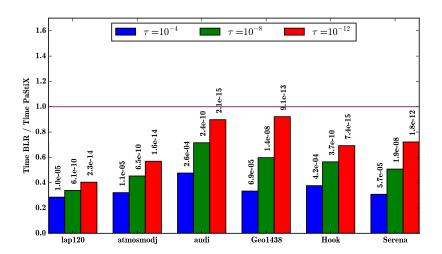
Performance



Memory footprint

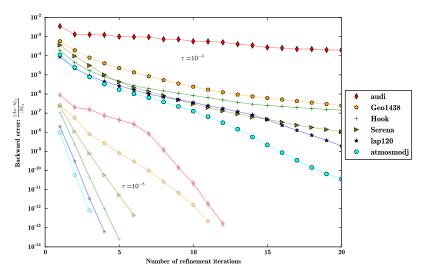


Performance of RRQR/Just-In-Time



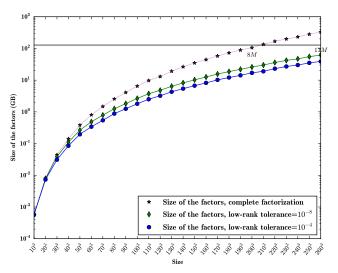


Convergence of RRQR/Minimal Memory





Scaling on Laplacians: size of the factors using RRQR/Minimal Memory







4

Case study: PASTIX inside a HDG method



HORSE

Method

- Introduce an hybrid variable to solve a system smaller than the one obtained through DG
- · Perform Domain-Decomposition to solve this system
- Use a direct solver for subdomains

Solve 3D frequency-domain Maxwell's equations

- $\mathbf{i}\omega\varepsilon_r\mathbf{E} \mathbf{curl}\,\mathbf{H} = -\mathbf{J}$, in Ω
- $\mathbf{i}\omega\mu_r\mathbf{H} + \mathbf{curl}\,\mathbf{E} = 0$, in Ω
- $\mathbf{n} \times \mathbf{E} = 0$, on Γ_m
- $\mathbf{n} \times \mathbf{E} + \mathbf{n} \times (\mathbf{n} \times \mathbf{H}) = \mathbf{n} \times \mathbf{E}^{\mathsf{inc}} + \mathbf{n} \times (\mathbf{n} \times \mathbf{H}^{\mathsf{inc}}), \text{ on } \Gamma_a$





The challenge with DG for frequency-domain problems

Classical DG formulation

- Naturally adapted to heterogeneous media and discontinuous solutions
- Can easily deal with unstructured, possibly non-conforming meshes (h-adaptivity)
- High order with compact stencils and non-conforming approximations (p-adaptivity)
- Usually rely on polynomial interpolation but can also accommodate alternative functions (e.g. plane waves)
- Well-suited for efficient parallelization
- But leads to larger problems compared to continuous finite element methods



Hybridizable DG method in 3D

Principles of a HDG formulation

- Keep the advantages of classical DG methods
- Introduce an hybrid variable to decouple local problems defined at the element level
- Solve a reduced linear system for the hybrid variable unknowns only

Complexity: number of globally coupled degrees of freedom

• Classical DG method with \mathbb{P}_p interpolation

$$(p+1)(p+2)(p+3)N_e$$
 , N_e is the # of elements

• HDG method with \mathbb{P}_p interpolation

$$(p+1)(p+2)N_f$$
 , N_f is the # of faces

• Continuous finite element formulation based on Nedelec's first family of face/edge elements in a simplex (tetrahedron): $\frac{p(p+2)(p+3)}{2}N_e$



Hybridizable DG method in 3D

ANR TECSER project (May 2014 - April 2017)
Funded by DGA (French armament procurement agency)
http://www-sop.inria.fr/nachos/projects/tecser

- Implementation of HDG for arbitrary high order interpolation
- Local definition (element-wise and face-wise) of the interpolation degree
- Extension of the formulation to a non-conforming hybrid hexahedral/tetrahedral mesh
- Scalability improvement
 PDE-based Schwartz domain decomposition algorithm
 PaStiX as a local (subdomain) solver
- Coupling with a BEM for an accurate treatment of far field radiation
- HORSE software
 High Order solver for Radar cross Section Evaluation

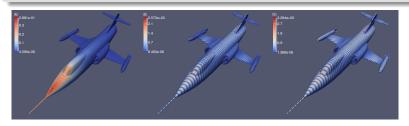


Hybridizable DG method in 3D

Scattering of a plane wave by a jet

- Unstructured tetrahedral mesh: 1,645,874 elements / 3,521,251 faces
- Frequency: 600 MHz, Wavelength: $\lambda \simeq 0.5$ m, Penalty parameter: $\tau = 1$

HDG method	# DoF Λ field	# DoF EM field
$HDG ext{-}\mathbb{P}_1$	21,127,506	39,500,976
$HDG ext{-}\mathbb{P}_2$	42,255,012	98,752,440
$HDG ext{-}\mathbb{P}_3$	70,425,020	197,504,880



Contour line of $|\mathbf{E}|$ - HDG $-\mathbb{P}_1$ to HDG $-\mathbb{P}_3$



Practical example

Study case

- Jet with \mathbb{P}_2 interpolation
- $42\ 255\ 012$ unknowns for Λ
- A single factorization is performed on each subdomain
- There may be several solves for iterative refinement between subdomains

Machine

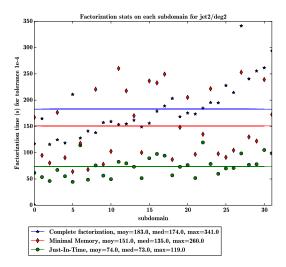
- Occigen: 24-cores nodes with 128 GB
- Each MPI process is hold on a socket with 12 cores
- In next experiments, we use 32, 48 or 64 subdomains
- Subdomains are ordered accordingly to the number of operations for the complete factorization





Factorization on each subdomain – 32 processes with

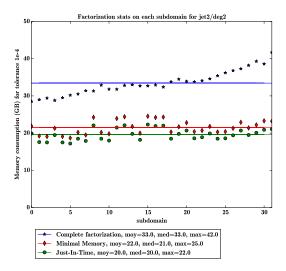
 $\tau = 1e - 4$







Memory Consumption on each subdomain – 32 processes with $\tau=1e-4$







Impact on HORSE

Subs	τ	Method	Fact(s)	Nb. of iters	Refinement (s)	HDGM (s)	Memory (GB)
32	-	Complete-Factorization	526.0	9	254.5	814.0	41.7
	1e-4	Just-In-Time	209.3	9	164.8	405.8	41.7 (22.3)
		Minimal-Memory	516.7	15	273.2	822.0	24.5
	1e-8	Just-In-Time	325.2	9	192.9	550.4	41.7 (29.4)
		Minimal-Memory	600.1	9	193.2	825.8	30.5
	-	Complete-Factorization	237.3	8	118.6	374.6	24.9
48	1e-4	Just-In-Time	112.2	9	106.1	237.1	24.9 (14.1)
		Minimal-Memory	229.3	13	159.7	407.5	15.3
	1e-8	Just-In-Time	171.5	8	105.8	296.1	24.9 (18.1)
		Minimal-Memory	319.4	8	109.8	447.9	18.8
	-	Complete-Factorization	179.8	9	104.7	298.3	17.4
64	1e-4	Just-In-Time	79.7	10	91.1	184.1	17.4 (10.0)
		Minimal-Memory	138.1	13	120.7	272.2	11.0
	1e-8	Just-In-Time	124.6	9	90.0	228.2	17.4 (12.9)
		Minimal-Memory	239.2	9	91.5	344.5	13.7



Conclusion

Block Low-Rank solver

- · Exploits the symbolic structure
- Demonstrate the extend-add issues in a supernodal context
- The version presented follows the parallelism of PASTIX
- On-going work to implement efficiently this approach over runtime systems

PaStiX in HORSE

- Low-rank compression is useful to reduce memory footprint and/or time-to-solution
- · Future work: study load imbalance between subdomains
- Future work: impact with $\mathbb{P}_3/\mathbb{P}_4$ interpolation





PASTIX 6.0.0alpha is now available!

http://gitlab.inria.fr/solverstack/pastix

- Support shared memory with different schedulers:
 - sequential
 - ► static scheduler
 - ► PaRSEC runtime system
- Low-rank support
- Cholesky and LU factorizations

Thank you.



