

Hybrid information flow analysis against web tracking (invited talk)

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Joint work with:





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Hybrid information flow analysis against web tracking

Thomas Jensen, Inria

Keynote, CRISIS'2017

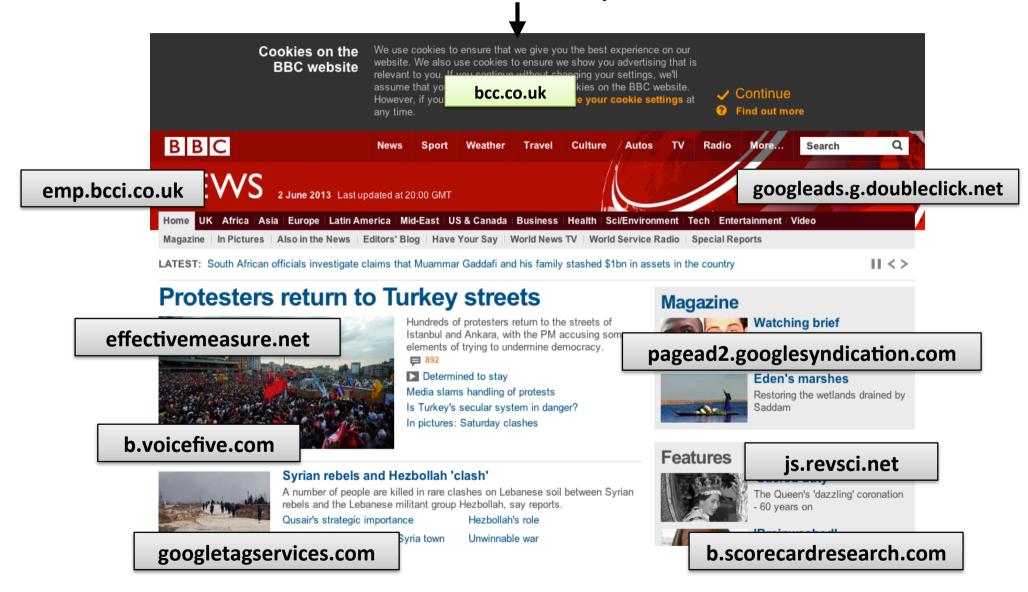


Overview

- 1. Web tracking.
- 2. Knowledge-based information flow analysis.
- 3. Browser randomisation.



Thanks to EU ePrivacy Directive



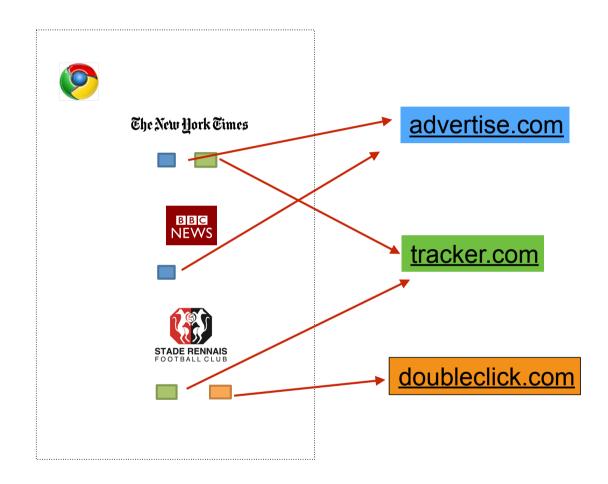


Web tracking



Bigger browsing profiles

- = increased value for trackers
- = reduced privacy for users



(Hypothetical tracking relationships only.)



Browser extensions to the rescue?



AdBlockPlus: blocks scripts/requests only from known advertisement companies



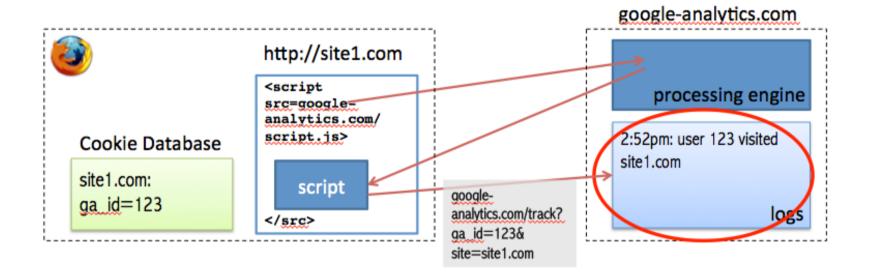
Ghostery: blocks scripts/requests only from known tracking companies

- No protection from tracking by other companies
- No protection from tracking by the main (first-party) web site



Tracking by storing identity

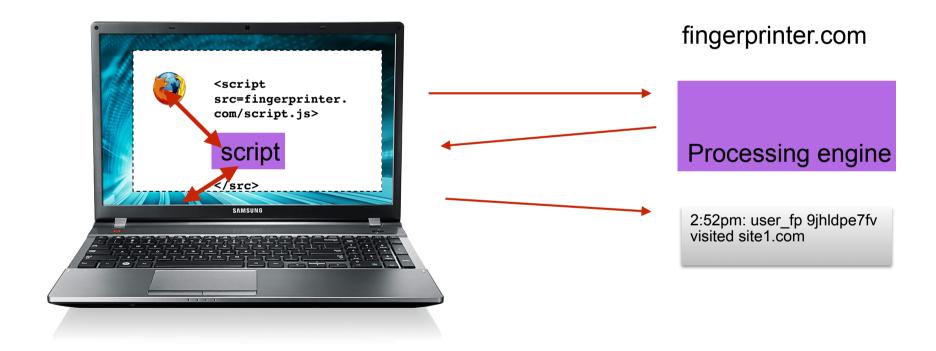
Cookies are used to track repeated visits to a site.





Tracking by finger-printing

Browser and operating system properties are used to track repeated visits to a site.





Panopticlick: an early study on fingerprinting





Your browser fingerprint appears to be unique among the 2,419,678 tested so far.

Currently, we estimate that your browser has a fingerprint that conveys at least 21.21 bits of identifying information.

- A study from 2010 showed that 83,6 % of browser fingerprints were unique among the 500 000 browsers tested.
- Fingerprints: HTTP headers, browser and OS features,:
 language, plug-ins, fonts, ...



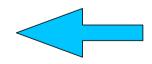
Protection mechanisms

Stateful tracking: well-known and being addressed at different levels:

- Third-party cookie blocking
- EU e-Privacy directive
- Non-interference analysis for JavaScript

Stateless tracking: less clear

- IP address tracking
- Web browser finger-printing





2. Knowledge-based information flow analysis



Use information flow to spot web tracking

Scripts can access information about configurations

Goal: Estimate the information that a script learns about configurations.

But some information leakage is reasonable!

Scripts may need information about browser and execution environment for optimal user experience.



Information flow control

Non-interference

Secret input does not flow into public output

Yes/no?

Attacker knowledge

What information about the secret flows to output in an execution.

Secret > 42

Quantitative information flow

How much information about secrets is leaked?

Leaks 17 bits



Attacker knowledge

```
var x = 0;
if (name == "Firefox") {
    x = 1;
}
else {
    if (fonts == fontsSet1) {
        x = 2;
    }
}
output x;
=> name = "Firefox" &&
    fonts = fontsSet1

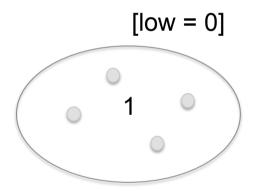
=> name = "Firefox" &&
    fonts = fontsSet1
```

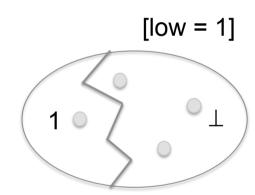
Depending on the browser, different executions leak different amount of information.



Non-interference

A program is termination-insensitive non-interferent if it produces the same output for all low-equivalent input.





Example with two binary secrets h1, h2 \in {0,1}

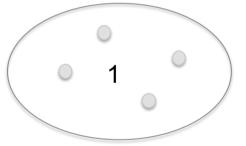
Secure and insecure executions

The following program is not secure - but it has secure executions.

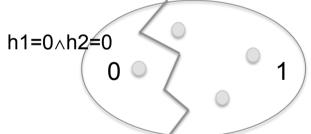
```
if h1 = 1 then x = 1
else skip;
if h2 = 1 then low = 1
else low = x;
output low
```

Initial values:

$$[low = 1, x = 1]$$



[low = 1, x = 0]



Secure

Leaks information about secret

Monitoring information flow

Dynamic information flow control

- Label values with security labels.
- Propagate labels during execution.
- Halt execution before outputting secret values.

Non-interference needs "no-sensitive- upgrade" principle: no assignment to public variables under secret control

```
[secret = 1]

if secret = 1 then

output = 1

else output = public;
return output
```



Hybrid information flow analysis

Combine dynamic IF control with a static analysis

- statically analyse the non-executed branches
- admits more executions but limited by the precision of the static analysis

```
[secret = 1, public = 1]

if secret = 1 then
  output = 1
else output = public;
return output
Static analysis says
public : [1]
```



Hybrid monitor of attacker knowledge

Define a hybrid monitor that tracks the knowledge of an attacker.

Use knowledge to

- enforce non-interference,
- accept more secure executions.

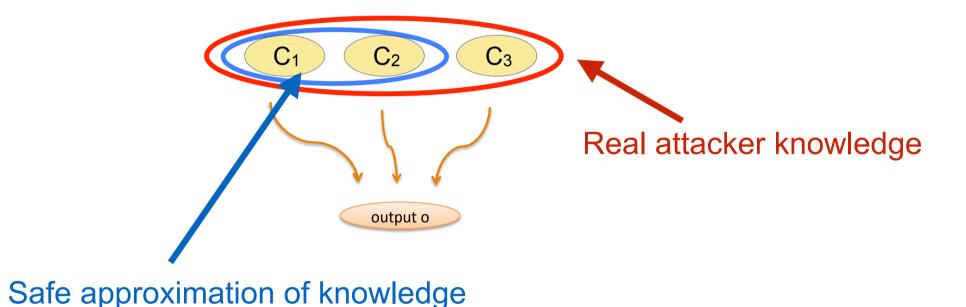
Solution:

- A domain for representing attacker knowledge.
- Monitor formalised as an operational semantics.



Attacker knowledge

Knowledge: What does an attacker learn about the input state by seeing a particular output.





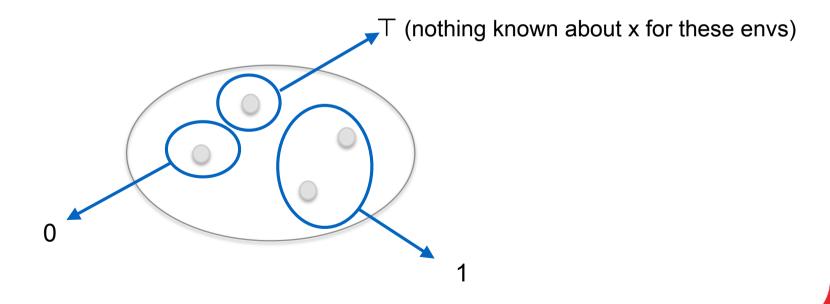
Domain of knowledge

$$\kappa$$
: Var \rightarrow K

 $\mathbf{K} \triangleq \mathrm{Env} \rightarrow \mathrm{Value} \cup \{\top, \bot\}$

The knowledge in a variable $\kappa(x)$:

which initial environments leads to a particular value of x?





The hybrid monitor

Defined as an operational reduction semantics:

$$(P, \rho, \kappa) \Downarrow (\rho', \kappa')$$

- ρ , ρ ': Env $\cup \{\cdot\}$,
- κ , κ ': Var \rightarrow K labeling of variables with knowledge,
- Env: Var → Val for dynamic analysis,
- · · "empty" environment for static analysis.

Semantic rules

Rules come in pairs: one dynamic and one static

Dynamic analysis computes real values and knowledge

Static analysis called on-the-fly and propagates only knowledge

$$\text{IFDYN } \frac{C[\![e]\!]_{\rho} = \alpha \qquad (c_{\alpha}, (\rho, \kappa)) \Downarrow (\rho', \kappa_{\alpha}) \qquad (c_{\bar{\alpha}}, (\cdot, \kappa)) \Downarrow (\cdot, \kappa_{\bar{\alpha}}) }{(\text{if } e \text{ then } c_{tt} \text{ else } c_{f\!f}, (\rho, \kappa)) \Downarrow (\rho', \mathbb{IF}(\{\![e]\!]_{\kappa}, \kappa_{tt}, \kappa_{f\!f}))}$$

IFSTAT
$$\frac{(c_{tt}, (\cdot, \kappa)) \Downarrow (\cdot, \kappa_{tt}) \quad (c_{ff}, (\cdot, \kappa)) \Downarrow (\cdot, \kappa_{ff})}{(\text{if } e \text{ then } c_{tt} \text{ else } c_{ff}, (\cdot, \kappa)) \Downarrow (\cdot, \mathbb{IF}(\{e\}_{\kappa}, \kappa_{tt}, \kappa_{ff}))}$$



Operator combining knowledge obtained in branches with knowledge from the condition

Enforcing non-interference

Before any output, the monitor checks non-interference:

if non-interference then output is not modified.

OUTNI
$$\frac{(c, init(\rho)) \Downarrow (\rho', \kappa) \quad \rho'(x) = v \quad NI(\rho, \kappa(x), v)}{(c; \text{output } x, \rho) \Downarrow v}$$

The non-interference check is defined as:

$$NI(\rho,K,v) \stackrel{\triangle}{=} (\rho) \subseteq K^{\text{-}1}(v) \cup K^{\text{-}1}(\bot).$$

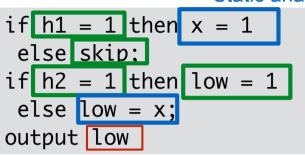




Example analysis

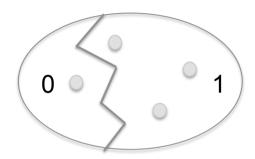
[h1 = 0, h2 = 1]

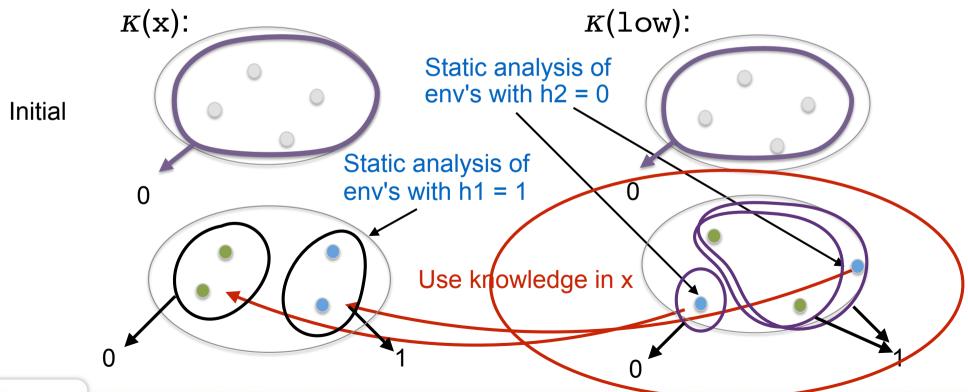
Static analysis



Knowledge

[low = 0, x = 0]







Stock-taking

Knowledge-based information flow analysis

- computes a correct approximation of attacker knowledge,
- enforces non-interference,
- more permissive than existing hybrid monitors,
- can be combined with other monitors to enhance precision.



3. Browser randomisation against web tracking



How to enforce privacy?

The hybrid monitor monitor evaluates how much a tracker learns about an individual user.

How to react when a script tracks too much?

```
p(name) =
Firefox 0.45
Chrome 0.45
Opera 0.10
```

- With 10 users, our Opera user will be uniquely identified.
- Halting the program or suppress output might still make you identifiable.
- Need collaboration between users.



A solution

- Users switch between configurations
- How many configurations are needed to guarantee privacy?
- Usability issue: users want to switch as little as possible





Challenge: find a distribution on configurations for all users such that

- Privacy is guaranteed
- Usability is maximised



Privacy: probability of guessing

```
if (name == "Opera") x = A;
else x = B;
output x;
```

A priori

```
p(name) =
Firefox 0.45
Chrome 0.45
Opera 0.10
```

$$x = A$$

A posteriori

```
p(name) =
Firefox 0
Chrome 0
Opera 1
```

Probability Pr^G: guessing the secret in case of worst observation.

Stronger that definition based on Bayesian risk: $p(A) \max_i p(i|A) + p(B) \max_i p(i|B) = 0.55$



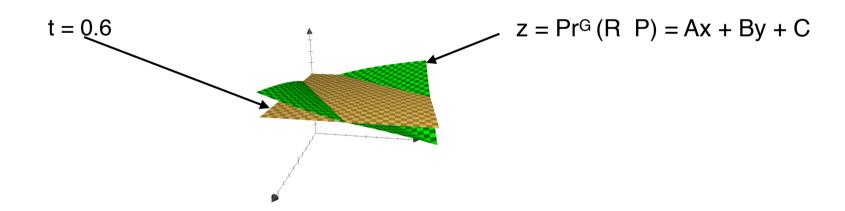
Achieving privacy

Find a randomisation R of user state such that when composed with tracking script P

 $Pr^{G}(R; P) < privacy threshold.$

R	"Firefox"	"Opera"
id_1	X	1-x
id_2	1-y	y

P	01	02	о3
"Firefox"	1/2	1/3	1/6
"Opera"	1/6	1/2	1/3

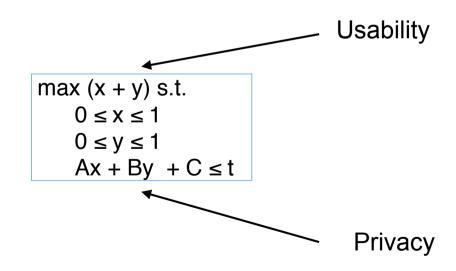


Adding usability

Users don't want to switch configuration:

- choose randomisation so that they get their original configuration as much as possible
- linear constraint on the randomisation matrix.
- problem reduces to linear programming.

R	"Firefox"	"Opera"
id_1	X	1-x
id ₂	1-y	y





Wrapping up

Web tracking is done by different techniques

cookies, other browser storage, fingerprinting

Information flow analysis against fingerprinting

- Hybrid information flow monitoring with static analysis
- Computes tracker knowledge.

Enforcing browser anonymity

- Randomisation of configurations
- Combining privacy and usability



Based on a true story...

- F. Besson, N. Bielova, T. Jensen: Hybrid Information Flow monitoring against web tracking, CSF 2013.
- F. Besson, N. Bielova, T. Jensen: Browser randomization against fingerprinting: A Quantitative information flow approach, NordSec 2014.
- F. Besson, N. Bielova, T. Jensen: Hybrid monitoring of attacker knowledge, CSF 2016.

