

# Data Discovery in RDF Graphs

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# Data Discovery in RDF Graphs

Ioana Manolescu

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DEXA Conference, Aug 29, 2017

### Outline

- Background: semantic RDF graphs
- Summarizing semantic-rich RDF graphs [ČGM15a, ČGM15b, ČGM17a]
  - Joint work with Šejla Čebirić (Inria) and François Goasdoué (U. Rennes 1 and Inria)
- Finding insights in RDF graphs [DMS17]
  - Joint work with Yanlei Diao and Shu Shang (Ecole Polytechnique and Inria)

Part I

# Background: RDF graphs

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#### Big Data needs semantics

AI Magazine, Spring 2015





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Discovering RDF Graphs

#### Do we really need the semantics?

#### Yes. All the time.

#### Application knowledge / constraints:

- Every Senator is an ElectedOfficial which is a Person
- (On Wikipedia) being BornInAPlace means being a Person

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#### Without the semantics, we may miss query answers

Data	Constraints	Query
John is a <u>Senator</u>	Every Senator is a Person	Who is a <u>Person</u> ?

#### Do we really need the semantics?

#### Yes. All the time.

#### Application knowledge / constraints:

- Every Senator is an ElectedOfficial which is a Person
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#### Semantic contraints are a **compact** way of encoding information

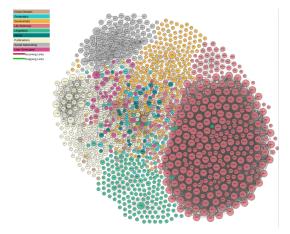
"Every  $\underline{\text{ElectedOfficial}}$  is a  $\underline{\text{Person}}$ " stated only once even if thousands of  $\underline{\text{ElectedOfficials}}$ .

#### Semantics for Web data

Data and metadata on the Web is often structured in graphs, e.g., **RDF** (W3C's Resource Description Framework)

RDF

• Famous application: the Linked Open Data cloud (2017)



## The Resource Description Framework (RDF)

#### RDF graph: set of triples

Assertion	Triple	Relational notation	Intuition
Class	s rdf:type o	o(s)	"s is an o"
Property	spo	p(s,o)	"The p of s is o"

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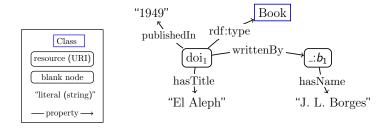


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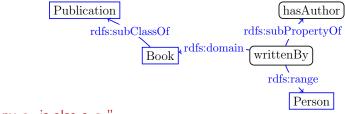
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# RDF Schema (RDFS)

#### Declare deductive constraints between classes and properties

Constraint	Triple	OWA interpretation
Subclass	$c_1$ rdfs:subClassOf $c_2$	$c_1\subseteqc_2$
Subproperty	$p_1$ rdfs:subPropertyOf $p_2$	$\mathtt{p}_1\subseteq \mathtt{p}_2$
Domain typing	p rdfs:domain c	${\sf \Pi}_{ m domain}({\tt p})\subseteq{\tt c}$
Range typing	p rdfs:range c	$\Pi_{\mathrm{range}}(\mathtt{p})\subseteq \mathtt{c}$

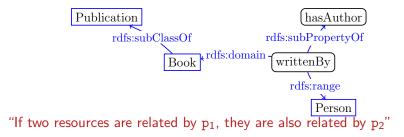


RDF RDFS

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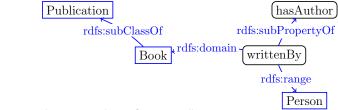
#### "Anyone having p is a c"

RDF RDFS

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"Anyone who is a value of p is a c"

### Open-world assumption and RDF entailment

RDF data model based on the open-world assumption.

Deductive constraints lead to **implicit triples**: part of the graph even though not explicitly present

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# Open-world assumption and RDF entailment

RDF data model based on the open-world assumption.

Deductive constraints lead to **implicit triples**: part of the graph even though not explicitly present

explicit triples		
+	$\rightarrow$	implicit triples
entailment rules		

Exhaustive application of entailment leads to saturation (closure)

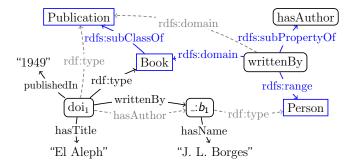
#### The semantics of an RDF graph G is its saturation $G^{\infty}$

Sample instance entailment rules from schema and instance triples

 $\mathtt{c_1} \ \mathrm{rdfs:subClassOf} \ \mathtt{c_2} \ \land \ \mathtt{s} \ \mathrm{rdf:type} \ \mathtt{c_1} \vdash_{\mathtt{RDF}} \mathtt{s} \ \mathrm{rdf:type} \ \mathtt{c_2}$ 

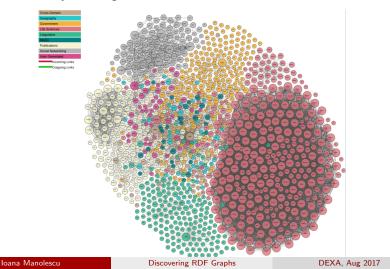
- $p_1 \; \mathrm{rdfs:subPropertyOf} \; p_2 \; \land \; \texttt{s} \; p_1 \; \texttt{o} \vdash_{\mathrm{RDF}} \texttt{s} \; p_2 \; \texttt{o}$
- p rdfs:domain c  $\land$  s p o  $\vdash_{RDF}$  s rdf:type c

p rdfs:range c  $\,\wedge\,$  s p o  $\vdash_{\rm RDF}$  o rdf:type c



### RDF graph discovery

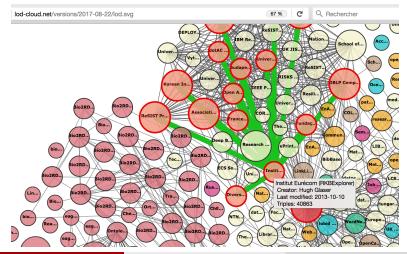
An RDF graph can be large and complex, lack a fixed schema, include many heterogeneous values...



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### RDF graph discovery

An RDF graph can be large and complex, lack a fixed schema, include many heterogeneous values...



# RDF graph discovery

Two approaches:

- RDF summarization: compactly representing the explicit and implicit structure of a graph
- Insight discovery in RDF graphs: automatically identify aggregation queries with interesting results

# Part II

# **RDF** summarization

#### **RDF** summaries

#### Problem

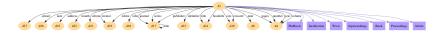
RDF graph G is large, heterogeneous, partially implicit. How to compactly represent all its structure?

#### Existing solutions

Partial representation (frequent patterns, statistics etc.)
e.g., [NM11, LYL13]
Potentially not compact e.g., [GW97, CFKP15]
Only for explicit data, e.g., [CDT13, ZDYZ14]

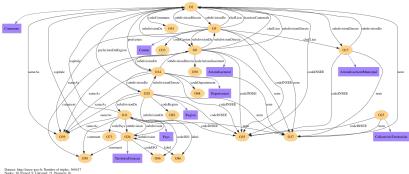
# A summary of DBLP data

#### 150M triples



#### A summary of geographic data

French territory division in regions, departments, urban areas, cities, districts etc. 368K triples



Dataset: http://insee-geo.fr. Number of triples: 368417 Nodes: 30 (Typed: 9, Untyped: 21, Property: 0) Edges: 69 (Data edges: 69, Schema edges: 0)

### **RDF** summaries

#### We define

- ORDF node equivalence relation ≡: equivalence relation such that class and property nodes are only equivalent to themselves
- ② RDF summary G<sub>/=</sub> of an RDF graph G: the quotient of G through =

Recall: quotient of a directed graph G by  $\equiv$ 

- G = (V, E),  $\equiv$  equivalence relation on V
  - $G_{/\equiv}$  nodes: one for  $\equiv$  equivalence class of V
  - $G_{/\equiv}$  edges:  $n_{\equiv}^1 \xrightarrow{a} n_{\equiv}^2$  iff  $\exists n_1 \xrightarrow{a} n_2 \in G$  such that  $n_1$  represented by  $n_{/\equiv}^1$ ,  $n_2$  represented by  $n_{/\equiv}^2$

Why not use any node equivalence? E.g., forward and backward bisimilarity  $\sim_{\tt fb} [\rm HHK95]$ 

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Sample graph G and its quotient through  $\sim_{\mathtt{fb}}$ 

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Sample graph G and its quotient through  $\sim_{\mathtt{fb}}$ 

Loss of class and (some) property names

Why not use any graph node equivalence? E.g., forward and backward bisimilarity  $\sim_{\tt fb}$ 

Sample graph G and its quotient through  $\sim_{\mathtt{fb}}$ 

#### Loss of schema triples

Why not use any graph node equivalence? E.g., forward and backward bisimilarity  $\sim_{\tt fb}$ 

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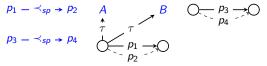
Sample graph G and its quotient through  $\sim_{\mathtt{fb}}$ 

#### Loss of implicit triples

Sample graph G and its quotient through  $\sim_{\mathtt{fb}}$ 

Quotient of the same graph through

the RDF node equivalence  $\equiv_{fb}$ 



## Formal summary properties

For any RDF equivalence relation  $\equiv$ :

Size limit	The summary is at most as large as the graph.
Schema	The schema of $G_{/\equiv}$ is the schema of G.
preservation	
Representativeness	Any conjunctive query $q$ with answers on G also
	has answers on its summary:
	$q(G^{\infty}) \neq \emptyset \Rightarrow q((G_{/\equiv})^{\infty}) \neq \emptyset$ This enables <b>query pruning (for query an</b> -
	This enables query pruning (for query an-
	swering) without saturating G

### Which equivalence relations to use?

#### Equivalence notions previously studied

- Forward / backward / forward and backward simulation
- Forward / backward / forward and backward bisimulation

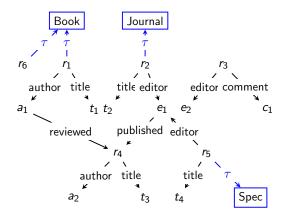
Adapted to semantic RDF graphs

Novel equivalence notions we introduce (see next)

- Flexible similarity suited to heterogeneous graphs
- Based on property cliques and possibly on RDF types

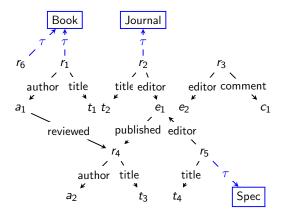
### RDF node equivalence based on property cliques

Intuition:  $a_1, a_2$  are similar;  $r_1, r_2, r_3, r_4, r_5$  are similar



### RDF node equivalence based on property cliques

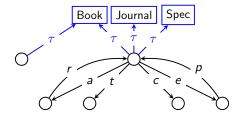
**Output property cliques:**  $\{a, t, e, c\}$ ;  $\{r\}$ ;  $\{p\}$ ;  $\emptyset$ **Input property cliques:**  $\{a\}$ ;  $\{t\}$ ;  $\{e\}$ ;  $\{c\}$ ;  $\{r, p\}$ ;  $\emptyset$ 



## Weak clique-based summaries

Two nodes are weakly equivalent ( $\equiv_W$ ) iff they have the same input clique or the same output clique.

Weak summary  $G_{\equiv W}$  of the sample RDF graph G:

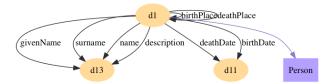


Property: In  $G_{\equiv W}$ , each data property appears exactly once  $\Rightarrow$  its nodes are "source of p, target of p" for each p [ČGM15b].

## Weak clique-based summaries

Property:  $G_{\equiv W}$  nodes are "source of p, target of p" for each p.

**Detecting errors in the data:** why do the birthplace and deathplace loop?



#### Looking in the data, we find:

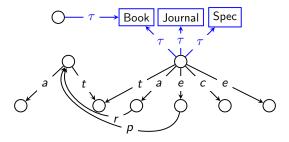
 $\label{eq:linear} $$ $ \frac{1999/02/22-rdf-syntax-ns\#type}{http://xmlns.com/foaf/0.1/Person} .$ 

(http://dbpedia.org/resource/Kunitomo\_lkkansai) (http://dbpedia.org/ontology/birthPlace)
(http://dbpedia.org/resource/Kunitomo\_lkkansai) .

## Strong clique-based summaries

Two nodes are strongly equivalent  $(\equiv_S)$  iff they have the same input clique and the same output clique.

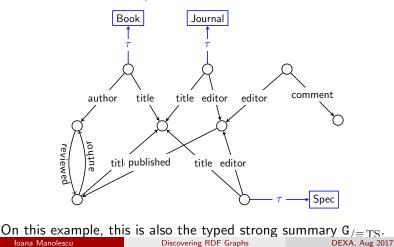
Strong summary  ${\tt G}_{/\equiv\,{\rm S}}$  of the sample RDF graph G:



## Using types for summarization

Group nodes **first by their types**; then group untyped nodes by their property cliques.

Typed weak summary  ${\tt G}_{/\equiv\,\rm TW}$  of the sample RDF graph G:



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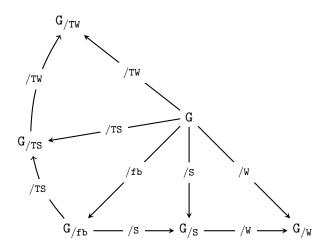
## RDF summaries outline

Summary	Weak?	Strong?	Types first?
$G_{\equiv W}$	$\checkmark$		
G <sub>/=S</sub>		$\checkmark$	
$G_{\equiv TW}$	$\checkmark$		$\checkmark$
$G_{\equiv TS}$		$\checkmark$	$\checkmark$

## RDF summaries outline

Summary	Weak?	Strong?	FW bisim?	BW bisim?	Types first?
$G_{\equiv W}$	$\checkmark$				
G <sub>/=S</sub>		$\checkmark$			
$G_{\equiv TW}$	$\checkmark$				$\checkmark$
$G_{\equiv TS}$		$\checkmark$			$\checkmark$
$G_{/\equiv  fw}$			$\checkmark$		
$G_{\equiv bw}$				$\checkmark$	
$G_{\equiv fb}$			$\checkmark$	$\checkmark$	
$G_{\equiv fw,T}$			$\checkmark$		$\checkmark$
$G_{\equiv bw,T}$				$\checkmark$	$\checkmark$
$G_{/\equiv  {\rm fb}, { m T}}$			$\checkmark$	$\checkmark$	$\checkmark$

# Relations between RDF summaries [ČGM17b]



# Summary size comparison (more in [ČGM17b])

Graph G	G	Summary $G_{/\equiv}$	$ G_{/\equiv} $	$cf_{\equiv}$
DBLP	150,787,464	$G_{ m W}$	71	2,123,767
DBLP	150,787,464	G <sub>/S</sub>	206	731,978
DBLP	150,787,464	$G_{ m fw}$	262,695	574
LUBM1M	1,227,868	$G_{ m /W}$	161	7,579
LUBM1M	1,227,868	G <sub>/S</sub>	207	5,903
LUBM1M	1,227,868	G <sub>/fw</sub>	1982	617
LUBM10M	11,990,183	$G_{ m /W}$	162	74,013
LUBM10M	11,990,183	G <sub>/S</sub>	206	58,204
LUBM10M	11,990,183	G <sub>/fw</sub>	24,958	480
LUBM10M	11,990,183	G/bw	6,162	1,944
LUBM10M	11,990,183	G/fb	11,990,076	1

## Summarizing $\mathtt{G}^\infty$

Recall: With an RDF Schema, the semantics of G is  $G^\infty \Rightarrow$  We really need  $(G^\infty)_{/\equiv}!$ 

- Saturate G, then summarize
- 2 Can we avoid saturating G?...

## Summarizing $\mathtt{G}^\infty$

Recall: With an RDF Schema, the semantics of G is  $G^{\infty} \Rightarrow$ We really need  $(G^{\infty})_{/\equiv}!$ 

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- 2 Can we avoid saturating G?...

#### Shortcut theorem [ČGM17a]

For the summaries  ${\tt G}_{\rm /W},~{\tt G}_{\rm /S},~{\tt G}_{\rm /fw},~{\tt G}_{\rm /bw},~{\tt G}_{\rm /fb}:$ 

 $({\tt G}^\infty)_{/\equiv}$  is the same as  $(({\tt G}_{/\equiv})^\infty)_{/\equiv}$ 

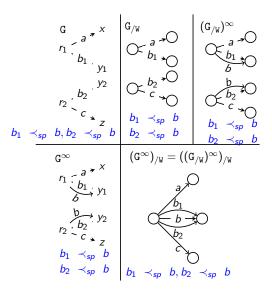
Also: sufficient condition for any  $\equiv$  to admit the shortcut.

### Shortcut toward the summary of ${\tt G}^\infty$

 $\begin{array}{lll} \text{Direct} & {\tt G} \to {\tt sat.} \to {\tt G}^\infty \to {\tt summ.} \to ({\tt G}^\infty)_\equiv \\ \text{Shortcut} & {\tt G} \to {\tt summ.} \to {\tt G}_\equiv \to {\tt sat.} \to ({\tt G}_\equiv)^\infty \to {\tt summ.} \to (({\tt G}_\equiv)^\infty)_\equiv \end{array}$ 

If  $G_{\equiv}$  is much smaller than G, **the shortcut may be faster!** Up to 20 times in our experiments [ČGM17b]

#### Shortcut example: G<sub>W</sub>



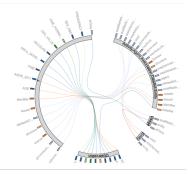
## Shortcut counter-example: $G_{TW}$

### Summary-enabled LOD cloud exploration

#### ILDA Inria team (E. Pietriga, H. Ozaygen) Use summary to derive visualisation instead of the original graph (smaller, faster)

abs-linked-data : Australian Bureau of Statistics (ABS) Linked Data 🛛 🗾

	· · · · · · · · · · · · · · · · · · ·
HEB O FD	AGE
	AGE ASGS 2011
	FREQUENCY
esources	MFASURF
la l	PROJ SERIES
ABS Datasets	SEX ABS
ADD Datasets	TIME
ABORIGINAL POP PROJ.rd	creator
ABORIGINAL_POP_PROUT	identifier
	issued
ABORIGINAL_POP_PROJ_R	license
(80,826 triples)	apent
	agent atTime
ABS ABORIGINAL POPPRC	entity
(879,162 triples)	
()	generated ceneratedAtTime
ABS ANNUAL ERP ASGS /	badBole
ADS_ANNOAL_ENT_ASSSS	gualifiedAssociation
ABS ANNUAL ERP LGA.rdf	qualifiedUsage
ABS_ANNOAL_ERP_LOA.IO	startedAtTime
	used
ABS_CENSUS2011_B01_LG	wasAssociatedWith
triples)	wasAttributedTo
	wasDerivedFrom
ABS_CENSUS2011_B01.rdf (	wasGeneratedBy
	wasinformedBy
ABS_CENSUS2011_B01_SA	wasStartedBy
(4,860,150 triples)	TIME FORMAT
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ABS CENSUS2011 B02 LG.	dataSet
triples)	structure
uipies)	



# Part III

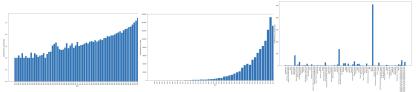
# Finding insights in RDF graphs

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## Insight in an RDF graph

We consider an insight to be the result of an aggregation query over the RDF graph

We focus one-dimensional aggregates  $\Rightarrow$  2D layout



An insight is interesting if a certain measure (e.g., variance) on its set of aggregation values is high

Problem

**Problem:** given a graph G, find the top-k insights

**Dagger**: Digging for Interesting Aggregates in RDF Graphs [DMS17] (ongoing)

1. Candidate facts Resources from G: of a certain type, or having certain property sets

2. Candidate dimension Properties of the candidate facts, with strong support and relatively few distinct values.

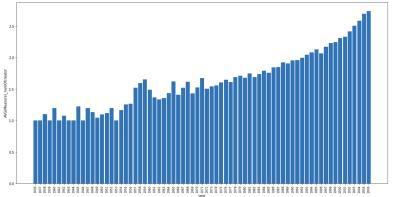
Also: derived properties, e.g., authors count;

3. Candidate measure Another property of the candidate facts Also: automatic value typing

4. Candidate aggregation function Chosen depending on the measure type

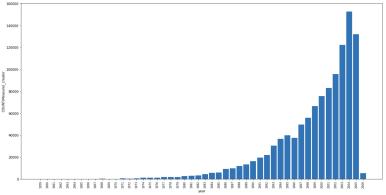
## Dagger-selected aggregate in DBLP data

Average number of authors of journal articles, per publication year



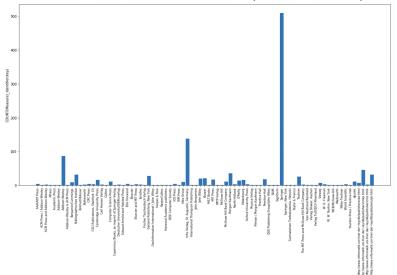
## Dagger-selected aggregate in DBLP data

#### Number of book authors, per book publication year



## Dagger-selected aggregate in DBLP data

The number of books by each publisher (highest: Springer)



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## Part IV

## Conclusion

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## The need for RDF graph discovery tools

- RDF graphs can be **large and complex**, they lack a prescriptive schema
- Semantic rules lead to implicit data
- Toward helping users to discover RDF graphs:
  - Structural quotient summaries representing the complete graph structure; compact clique-based summaries; available at:

https://team.inria.fr/cedar/projects/rdfsummary/

Insight discovery: interesting aggregate queries; project Web page:

https://team.inria.fr/cedar/projects/dagger/

 Many follow-up directions: parallelization, more interestingness measures, extensions to ML.

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