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Vertex-Coloring Edge-Weighting of Bipartite Graphs with Two Edge Weights

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Let \( G \) be a graph and \( S \) be a subset of \( \mathbb{Z} \). A vertex-coloring \( S \)-edge-weighting of \( G \) is an assignment of weights by the elements of \( S \) to each edge of \( G \) so that adjacent vertices have different sums of incident edges weights.

It was proved that every 3-connected bipartite graph admits a vertex-coloring \( S \)-edge-weighting for \( S = \{1, 2\} \) (H. Lu, Q. Yu and C. Zhang, Vertex-coloring 2-edge-weighting of graphs, European J. Combin., 32 (2011), 22-27). In this paper, we show that every 2-connected and 3-edge-connected bipartite graph admits a vertex-coloring \( S \)-edge-weighting for \( S \in \{\{0, 1\}, \{1, 2\}\} \). These bounds we obtain are tight, since there exists a family of infinite bipartite graphs which are 2-connected and do not admit vertex-coloring \( S \)-edge-weightings for \( S \in \{\{0, 1\}, \{1, 2\}\} \).

Keywords: edge-weighting, vertex-coloring, 2-connected, bipartite graph

1 Introduction

In this paper, we consider only finite, undirected and simple connected graphs. For a vertex \( v \) of graph \( G = (V, E) \), \( N_G(v) \) denotes the set of vertices which are adjacent to \( v \) and \( d_G(v) = |N_G(v)| \) is called the degree of vertex \( v \). Let \( \delta(G) \) and \( \Delta(G) \) denote the minimum degree and maximum degree of graph \( G \), respectively. For \( v \in V(G) \) and \( r \in \mathbb{Z}^+ \), let \( N^r_G(v) = \{u \in N(v) \mid d_G(u) = r\} \). If \( v \in V(G) \) and \( e \in E(G) \), we use \( v \sim e \) to denote that \( v \) is an end-vertex of \( e \). For two disjoint subsets \( S, T \) of \( V(G) \), let \( E_G(S, T) \) denote the subset of edges of \( E(G) \) with one end in \( S \) and other end in \( T \) and let \( e_G(S, T) = |E_G(S, T)| \). Let \( G = (U, W; E) \) denote a bipartite graph with bipartition \((U, W)\) and edge set \( E \).

Let \( S \) be a subset of \( \mathbb{Z} \). An \( S \)-edge-weighting of a graph \( G \) is an assignment \( w : E(G) \to S \). An \( S \)-edge-weighting \( w \) of a graph \( G \) induces a coloring of the vertices of \( G \), where the color of vertex \( v \), denoted by \( c(v) \), is \( \sum_{e \sim v} w(e) \). An \( S \)-edge-weighting of a graph \( G \) is a vertex-coloring if for every edge \( e = uv \), \( c(u) \neq c(v) \) and we say that \( G \) admits a vertex-coloring \( S \)-edge-weighting. If \( S = \{1, 2, \ldots, k\} \), then a vertex-coloring \( S \)-edge-weighting of a graph \( G \) is usually called a vertex-coloring \( k \)-edge-weighting.

For vertex-coloring edge-weighting, Karoński et al. (2004) posed the following conjecture:

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Conjecture 1.1 Every graph without isolated edges admits a vertex-coloring 3-edge-weighting.

This conjecture is still wide open. [Karoński et al. (2004)] showed that Conjecture 1.1 is true for 3-colorable graphs. Recently, [Kalkowski et al. (2010)] showed that every graph without isolated edges admits a vertex-coloring 5-edge-weighting. This result is an improvement on the previous bounds on k established by [Addario-Berry et al. (2007)], [Addario-Berry et al. (2008)], and [Wang and Yu (2008)], who obtained the bounds \( k = 30, k = 16, \) and \( k = 13, \) respectively.

Many graphs actually admit a vertex-coloring 2-edge-weighting (in fact, experiments suggest (see [Addario-Berry et al. (2008)]) that almost all graphs admit a vertex-coloring 2-edge-weighting), however it is not known which ones do not. [Khatirinejad et al. (2012)] explored the problem of classifying those graphs which admit a vertex-coloring 2-edge-weighting. [Chang et al. (2011)] had made some progress in determining which classes of graphs admit vertex-coloring 2-edge-weightings, and proved that there exists a family of infinite bipartite graphs (e.g., the generalized \( \theta \)-graphs) which are 2-connected and admit a vertex-coloring 3-edge-weighting but not vertex-coloring 2-edge-weightings. [Lu et al. (2011)] showed that every 3-connected bipartite graph admits a vertex-coloring 2-edge-weighting.

We write

\[ G_{12} = \{ G | G \text{ admits a vertex-coloring } \{1, 2\}\text{-edge-weighting} \}; \]
\[ G_{01} = \{ G | G \text{ admits a vertex-coloring } \{0, 1\}\text{-edge-weighting} \}; \]
\[ G^*_{12} = \{ G | G \text{ is bipartite and admits a vertex-coloring } \{1, 2\}\text{-edge-weighting} \}; \]
\[ G^*_{01} = \{ G | G \text{ is bipartite and admits a vertex-coloring } \{0, 1\}\text{-edge-weighting} \}. \]

[Dudek and Wajc (2011)] showed that determining whether a graph belongs to \( G_{12} \) or \( G_{01} \) is NP-complete. Moreover, they showed that \( G_{12} \neq G_{01} \). The counterexamples constructed by [Dudek and Wajc (2011)] are non-bipartite.

Now we construct a bipartite graph, which admits a vertex-coloring 2-edge-weighting but not vertex-coloring \( \{0, 1\}\)-edge-weightings. Let \( C_6 \) be a cycle of length six and \( \Gamma \) be a graph obtained by connecting an isolated vertex to one of the vertices of \( C_6 \). Take two disjoint copies of \( \Gamma \). Connect two vertices of degree one of the two copies and this gives a connected bipartite graph \( G \). It is easy to prove that \( G \) admits a vertex-coloring 2-edge-weighting but not vertex-coloring \( \{0, 1\}\)-edge-weighting. Hence \( G^*_{01} \neq G^*_{12} \).

Next we would like to propose the following problem.

Problem 1 Determining whether a graph \( G \in G^*_{12} \) or \( G \in G^*_{01} \) is polynomial?

In this paper, we characterize bipartite graphs which admit a vertex-coloring \( S \)-edge-weighting for \( S \in \{\{0, 1\}, \{1, 2\}\} \), and obtain the following result.

Theorem 1.2 Let \( G \) be a 3-edge-connected bipartite graph \( G = (U, W, E) \) with minimum degree \( \delta(G) \). If \( G \) contains a vertex \( u \) of degree \( \delta(G) \) such that \( G - u \) is connected, then \( G \) admits a vertex-coloring \( S \)-edge-weighting for \( S \in \{\{0, 1\}, \{1, 2\}\} \).

By Theorem 1.2, it is easy to obtain the following result, which improves and extends the result obtained by [Lu et al. (2011)].

Theorem 1.3 Every 2-connected and 3-edge-connected bipartite graph \( G = (U, W, E) \) admits a vertex-coloring \( S \)-edge-weighting for \( S \in \{\{0, 1\}, \{1, 2\}\} \).
So far, all known counterexamples of bipartite graphs, which do not have vertex-coloring \{0, 1\}-edge-weightings or vertex-coloring \{1, 2\}-edge-weightings are graphs with minimum degree 2. So we would like to propose the following problem.

**Problem 2** Does every bipartite graph with \(\delta(G) \geq 3\) admit a vertex-coloring \(S\)-edge-weighting, where \(S \in \{\{0,1\}, \{1,2\}\}\)?

A factor of a graph \(G\) is a spanning subgraph. For a graph \(G\), there is a close relationship between 2-edge-weighting and graph factors. Namely, a 2-edge-weighting problem is equivalent to finding special factors of graphs (see Addario-Berry et al. (2007) and Addario-Berry et al. (2008)). So to find factors with pre-specified degree is an important part of edge-weighting.

Let \(g, f : V(G) \rightarrow \mathbb{Z}\) be two integer-valued functions such that \(g(v) \leq f(v)\) and \(g(v) \equiv f(v) \mod 2\) for all \(v \in V(G)\). A factor \(F\) of \(G\) is called \((g, f)\)-parity factor if \(g(v) \leq d_F(v) \leq f(v)\) and \(d_F(v) \equiv f(v) \mod 2\) for all \(v \in V(G)\). For \(X \subseteq V(G)\), we write \(g(X) = \sum_{x \in X} g(x)\) and \(f(X)\) is defined similarly. For \((g, f)\)-parity factors, Lovász obtained a sufficient and necessary condition.

**Theorem 1.4 (Lovász (1972))** A graph \(G\) contains a \((g, f)\)-parity factor if and only if for any two disjoint subsets \(S\) and \(T\) of \(V(G)\), it follows that

\[
\eta(S, T) = f(S) - g(T) + \sum_{x \in T} d_{G-S}(x) - \tau(S, T) \geq 0,
\]

where \(\tau(S, T)\) denotes the number of components \(C\), called \(g\)-odd components of \(G - S - T\) such that \(g(V(C)) + c_G(V(C), T) \equiv 1 \mod 2\).

In the proof of the main theorems, we also need the following three lemmas.

**Theorem 1.5 (Chang et al. (2011))** Every non-trivial connected bipartite graph \(G = (A, B, E)\) with \(|A|\) even, admits a vertex-coloring 2-edge-weighting \(w\) such that \(c(u)\) is odd for \(u \in A\) and \(c(v)\) is even for \(v \in B\).

**Theorem 1.6 (Chang et al. (2011))** Let \(r \geq 3\) be an integer. Every \(r\)-regular bipartite graph \(G\) admits a vertex-coloring 2-edge-weighting.

**Theorem 1.7 (Khatirinejad et al. (2012))** Every \(r\)-regular graph \(G\) admits a vertex-coloring 2-edge-weighting if and only if it admits a vertex-coloring \{0, 1\}-edge-weighting.

## 2 Proof of Theorem 1.2

**Corollary 2.1** Every non-trivial connected bipartite graph \(G = (A, B, E)\) with \(|A|\) even admits a vertex-coloring \{0, 1\}-edge-weighting.

**Proof:** By Theorem 1.5 \(G\) admits a vertex-coloring 2-edge-weighting \(w\) such that \(c(u)\) is odd for \(u \in A\) and \(c(v)\) is even for \(v \in B\). Let \(w'(e) = 0\) if \(w(e) = 2\) and \(w'(e) = 1\) if \(w(e) = 1\). Then \(w'\) is a vertex-coloring \{0, 1\}-edge-weighting of graph \(G\).

For completing the proof of Theorem 1.2 we need the following two technical lemmas.
Lemma 2.2 Let $G$ be a bipartite graph with bipartition $(U, W)$, where $|U| \equiv |W| \equiv 1 \pmod{2}$. Let $\delta(G) = \delta$ and $u \in U$ such that $d_G(u) = \delta$. If one of the following two conditions holds, then $G$ contains a factor $F$ such that $d_F(u) = \delta$, $d_F(x) \equiv \delta + 1 \pmod{2}$ for all $x \in U - u$, $d_F(y) \equiv \delta \pmod{2}$ for all $y \in W$ and $d_F(y) \leq \delta - 2$ for all $y \in N^\delta_G(u)$.

(i) $\delta(G) \geq 4$, $G$ is 3-edge-connected and $G - u$ is connected.
(ii) $\delta(G) = 3$, $G$ is 3-edge-connected and $|N^\delta_G(u)| \leq 2$.

Proof: Let $M$ be an integer such that $M \geq \triangle(G)$ and $M \equiv \delta \pmod{2}$. Let $m \in \{0, -1\}$ such that $m \equiv \delta \pmod{2}$. Let $g, f : V(G) \to \mathbb{Z}$ such that

$$g(x) = \begin{cases} \delta & \text{if } x = u, \\ m - 1 & \text{if } x \in U - u, \\ m & \text{if } x \in W, \end{cases}$$

and

$$f(x) = \begin{cases} M + 1 & \text{if } x \in U - u, \\ M & \text{if } x \in (W \cup \{u\}) - N^\delta_G(u), \\ \delta - 2 & \text{if } x \in N^\delta_G(u). \end{cases}$$

By definition, we have $g(v) \equiv f(v) \pmod{2}$ for all $v \in V(G)$. It is sufficient for us to show that $G$ contains a $(g, f)$-parity factor. Indirectly, suppose that $G$ contains no $(g, f)$-parity factors. By Theorem [1.4] there exist two disjoint subsets $S$ and $T$ such that

$$\eta(S, T) = f(S) - g(T) + \sum_{x \in T} d_{G-S}(x) - \tau(S, T) < 0,$$

where $\tau(S, T)$ denotes the number of $g$-odd components of $G - S - T$. Since $f(V(G))$ is even, by parity, we have

$$\eta(S, T) = f(S) - g(T) + \sum_{x \in T} d_{G-S}(x) - \tau(S, T) \leq -2. \quad (1)$$

Hence $S \cup T \neq \emptyset$. We choose $S$ and $T$ such that $S \cup T$ is minimal. Let $A = V(G) - S - T$.

Claim 1. $T \subseteq \{u\}$. 


Otherwise, let \( v \in T - u \) and \( T' = T - v \). We have
\[
\eta(S, T') = f(S) - g(T') + \sum_{x \in T'} d_{G-S}(x) - \tau(S, T')
\]
\[
= f(S) - (g(T) - g(v)) + \left( \sum_{x \in T} d_{G-S}(x) - d_{G-S}(v) \right) - \tau(S, T')
\]
\[
\leq f(S) - g(T) + \left( \sum_{x \in T} d_{G-S}(x) - d_{G-S}(v) \right) - (\tau(S, T) - e_G(v, A)) + g(v)
\]
\[
\leq f(S) - g(T) + \sum_{x \in T} d_{G-S}(x) - \tau(S, T) + (g(v) - d_{G-S}(v) + e_G(v, A))
\]
\[
= \eta(S, T) + (g(v) - d_{G-S}(v) + e_G(v, A))
\]
\[
\leq \eta(S, T) - (d_{G-S}(v) - e_G(v, A))
\]
\[
\leq \eta(S, T) \leq -2,
\]
contradicting the choice of \( S \) and \( T \).

**Claim 2.** \( S \subseteq N_G^G(u) \).

Otherwise, suppose that \( S - N_G^G(u) \neq \emptyset \) and let \( v \in S - N_G^G(u) \). Let \( S' = S - v \). We have
\[
\eta(S', T) = f(S') - g(T) + \sum_{x \in T} d_{G-S'}(x) - \tau(S', T)
\]
\[
= (f(S) - f(v)) - g(T) + \left( \sum_{x \in T} d_{G-S}(x) + e_G(v, T) \right) - \tau(S', T)
\]
\[
\leq f(S) - g(T) + \sum_{x \in T} d_{G-S}(x) - (\tau(S, T) - e_G(v, A)) - f(v) + e_G(v, T)
\]
\[
= f(S) - g(T) + \sum_{x \in T} d_{G-S}(x) - \tau(S, T) + (e_G(v, T) + e_G(v, A) - f(v))
\]
\[
\leq \eta(S, T) + (d_G(v) - f(v))
\]
\[
\leq \eta(S, T) \leq -2,
\]
contradicting the choice of \( S \) and \( T \) again.

We write \( \tau(S, T) = \tau \). By Claims 1 and 2, we have
\[
\eta(S, T) = f(S) - g(T) + \sum_{x \in T} d_{G-S}(x) - \tau
\]
\[
= (\delta - 2)|S| - |T| + |T|(|S| - |S|) - \tau \quad \text{(by Claims 1 and 2)}
\]
\[
= (\delta - 2 - |T||S|) - \tau \quad \text{(by Claim 1)}
\]
\[
\leq -2,
\]
i.e.,
\[(\delta - 2 - |T||S| + 2 \leq \tau, \tag{2}\]
which implies \(\tau \geq 2\) since \(|T| \leq 1\).

Since \(G - u\) is connected, we may see that \(S \neq \emptyset\). Note that \(G\) is 3-edge-connected, by Claims 1 and 2, we have
\[3\tau \leq e_G(A, S \cup T)
= e_G(A, S) + e_G(A, T)
\leq (\delta - |T||S| + |T|)(\delta - |S|) \quad \text{(by Claims 1 and 2)}
= (\delta - 2|T||S| + |T||S| + |T|, \delta),\]
i.e.,
\[3\tau \leq (\delta - 2|T||S| + |T||S|, \delta). \tag{3}\]
Combining (2) and (3), we may see that
\[\delta|T| \geq (2\delta - |T| - 6)|S| + 6. \tag{4}\]

If \(\delta \geq 4\), then we have
\[\delta \geq \delta|T| \quad \text{(since } |T| \leq 1)\]
\[\geq (2\delta - |T| - 6)|S| + 6 \quad \text{(since } |S| \geq 1)\]
\[\geq 2\delta - |T|\]
\[\geq 2\delta - 1,\]
a contradiction. So we may assume that \(\delta = 3\). Note that \(|S| \leq |N^6_G(u)| \leq 2\). By (4), we have
\[3 \geq \delta|T| \geq -|T||S| + 6 \geq 4, \tag{5}\]
a contradiction again.
This completes the proof. \(\square\)

**Lemma 2.3** Let \(G\) be a bipartite graph with bipartition \((U, W)\), where \(|U| \equiv |W| \equiv 1 \pmod{2}\). Let \(\delta(G) = \delta\) and \(u \in U\) such that \(d_G(u) = \delta\). If one of the following two conditions holds, then \(G\) contains a factor \(F\) such that \(d_F(u) = 0\), \(d_F(x) \equiv 1 \pmod{2}\) for all \(x \in U - u\), \(d_F(y) \equiv 0 \pmod{2}\) for all \(x \in W\) and \(d_F(y) \geq 2\) for all \(y \in N_G(u)\).

(i) \(\delta(G) \geq 4\), \(G\) is 3-edge-connected and \(G - u\) is connected.

(ii) \(\delta(G) = 3\), \(G\) is 3-edge-connected and there exists a vertex \(v \in N_G(u)\) such that \(d_G(v) > 3\).

**Proof:** Let \(M\) be an even integer such that \(M \geq \Delta(G)\). Let \(g, f : V(G) \to Z\) such that
\[g(x) = \begin{cases} 0 & \text{if } x \in (\{u\} \cup W) - N_G(u), \\ 2 & \text{if } x \in N_G(u), \\ -1 & \text{if } x \in U - u, \end{cases}\]
and

\[ f(x) = \begin{cases} 
   M & \text{if } x \in W \\
   0 & \text{if } x = u, \\
   M + 1 & \text{if } x \in U - u.
\end{cases} \]

Clearly, \( g(v) \equiv f(v) \pmod{2} \) for all \( v \in V(G) \) and \( g(V(G)) \) is even. It is also sufficient for us to show that \( G \) contains a \((g, f)\)-parity factor.

Indirectly, suppose that \( G \) contains no \((g, f)\)-parity factors. By Theorem 1.4, there exist two disjoint subsets \( S \) and \( T \) such that

\[ \eta(S, T) = f(S) - g(T) + \sum_{x \in T} d_{G - S}(x) - \tau(S, T) \leq -2, \tag{6} \]

where \( \tau(S, T) \) denotes the number of \( g \)-odd components of \( G - S - T \). We choose \( S \) and \( T \) such that \( S \cup T \) is minimal. Let \( A = V(G) - S - T \).

**Claim 1.** \( S \subseteq \{u\} \).

Otherwise, suppose that there exists a vertex \( v \in S - u \). Let \( S' = S - v \). Then we have

\[
\eta(S', T) = f(S') - g(T) + \sum_{y \in T} d_{G - S'}(y) - \tau(S', T)
\]

\[
= (f(S) - f(v)) - g(T) + \left( \sum_{y \in T} d_{G - S}(y) + e_G(v, T) \right) - \tau(S', T)
\]

\[
\leq f(S) - g(T) + \sum_{y \in T} d_{G - S}(y) - f(v) + e_G(v, T) - (\tau(S, T) - e_G(v, A))
\]

\[
= \eta(S, T) - (f(v) - e_G(v, T) - e_G(v, A))
\]

\[
\leq \eta(S, T) - (f(v) - d_G(v))
\]

\[
\leq \eta(S, T) \leq -2,
\]

contradicting the the choice of \( S \cup T \).

**Claim 2.** \( T \subseteq N_G(u) \).
Otherwise, suppose that \( T - N_G(u) \neq \emptyset \). Let \( x \in T - N_G(u) \) and let \( T' = T - x \). Then we have

\[
\eta(S, T') = f(S) - g(T') + \sum_{y \in T'} d_{G-S}(y) - \tau(S, T')
\]

\[
= f(S) - (g(T) - g(x)) + \left( \sum_{y \in T} d_{G-S}(y) - d_{G-S}(x) \right) - \tau(S, T')
\]

\[
\leq f(S) - (g(T) - g(x)) + \left( \sum_{y \in T} d_{G-S}(y) - d_{G-S}(x) \right) - (\tau(S, T) - e_G(x, A))
\]

\[
= f(S) - g(T) + \sum_{y \in T} d_{G-S}(y) - \tau(S, T) - (d_{G-S}(x) - e_G(x, A)) + g(x)
\]

\[
= \eta(S, T) - (d_{G-S}(x) - e_G(x, A)) + g(x)
\]

\[
\leq \eta(S, T) \leq -2,
\]

contradicting the choice of \( S \cup T \).

By Claims 1 and 2, we may see that \( f(S) = 0 \) and \( g(T) = 2 |T| \). For simplicity, we write \( \tau(S, T) = \tau \).

By (6), we see that

\[
\tau \geq \sum_{x \in T} (d_G(x) - |S|) - 2 |T| + 2,
\]

which implies

\[
\tau \geq \sum_{x \in T} (\delta - 1) - 2 |T| + 2 \geq 2.
\]

(7)

(8)

Note that \( G - u \) is connected, so we have \( |T| \geq 1 \). Since \( G \) is 3-edge-connected, we have

\[
3 \tau \leq \sum_{x \in T} (d_G(x) - |S|) + (\delta - |T|)|S|
\]

\[
= \sum_{x \in T} d_G(x) + (\delta - 2 |T|)|S|,
\]

i.e.,

\[
3 \tau \leq \sum_{x \in T} d_G(x) + (\delta - 2 |T|)|S|.
\]

(9)

Inequalities (7) and (9) implies

\[
2 \sum_{x \in T} d_G(x) + 6 \leq |S||T| + 6 |T| + \delta |S| \quad \text{(since } |S| \leq 1 \text{ and } |T| \geq 1 \text{)}
\]

\[
\leq 7 |T| + \delta,
\]
Vertex-Coloring Edge-Weighting of Bipartite Graphs with Two Edge Weights

i.e.,

\[ 7|T| \geq 2 \sum_{x \in T} d_G(x) + 6 - \delta. \] (10)

If \( \delta \geq 4 \), by (10), it follows

\[ 7|T| \geq 6 + \delta(2|T| - 1) \geq 8|T| + 2, \]

a contradiction. So we may assume that \( \delta = 3 \). By condition (ii), \( \sum_{x \in T} d_G(x) \geq 3|T| + 1 \). Combining (10),

\[ 7|T| \geq 2 \sum_{x \in T} d_G(x) + 6 \]

\[ \geq 2(3|T| + 1) + 3, \]

which implies \( |T| \geq 5 \), a contradiction since \( |T| \leq |N_G(u)| \leq 3 \).

This completes the proof. \( \square \)

**Proof of Theorem 1.2:** By Theorem 1.5 and Corollary 2.1 we can assume that both \( |A| \) and \( |B| \) are odd.

Firstly, we consider \( S = \{0, 1\} \). If \( G \) is 3-regular, by Theorem 1.6, then \( G \) admits a vertex-coloring 2-edge-weighting. By Theorem 1.7, \( G \) also admits a vertex-coloring \( \{0, 1\} \)-edge-weighting. So we can assume that \( \delta(G) \geq 3 \) and \( G \) is not 3-regular. If \( \delta(G) = 3 \), since \( G \) is 3-edge-connected, then \( G - x \) is connected for every vertex \( x \) of \( G \) with degree three. Hence there exists a vertex \( v \) with degree three such that \( N_G(v) \) contains a vertex with degree at least four. Let

\[ u^* = \begin{cases} u & \text{if } \delta \geq 4, \\ v & \text{if } \delta = 3. \end{cases} \]

Without loss generality, we may assume that \( u^* \in U \) and so it is a vertex satisfying the conditions of Lemma 2.3. Hence by Lemma 2.3 \( G \) contains a factor \( F \), which satisfies the following three conditions.

(i) \( d_F(u^*) = 0 \);

(ii) \( d_F(x) \equiv 1 \pmod{2} \) for all \( x \in U - u^* \);

(iii) \( d_F(y) \equiv 0 \pmod{2} \) for all \( x \in W \) and \( d_F(y) \geq 2 \) for all \( y \in N_G(u^*) \).

Clearly, \( d_F(x) \neq d_F(y) \) for all \( xy \in E(G) \). We assign weight 1 for each edge of \( E(F) \) and weight 0 for each edge of \( E(G) - E(F) \). Then we obtain a vertex-coloring \( \{0, 1\} \)-edge-weighting of graph \( G \).

Secondly, we show that \( G \) admits a vertex-coloring 2-edge-weighting. By Theorem 1.6 we may assume that \( G \) is not 3-regular. If \( \delta = 3 \), since \( G \) is 3-edge-connected, then \( G \) contains a vertex \( v' \) such that \( d_G(v') = 3 \), \( G - v' \) is connected and \( |N_G^\delta(v')| \leq 2 \). Let

\[ u^* = \begin{cases} u & \text{if } \delta \geq 4, \\ v' & \text{if } \delta = 3. \end{cases} \]

Then \( u^* \) is a vertex satisfying the conditions of Lemma 2.2. Hence by Lemma 2.2 \( G \) contains a factor \( F \) such that
(i) $d_F(u^*) = \delta$;
(ii) $d_F(x) \equiv \delta \pmod{2}$ for all $x \in W$ and $d_F(x) \leq \delta - 2$ for all $x \in N^\delta_G(u^*)$;
(iii) $d_F(y) \equiv \delta + 1 \pmod{2}$ for all $y \in U - u^*$.

Let $w : E(G) \to \{1, 2\}$ be a 2-edge-weighting such that $w(e) = 1$ for each $e \in E(F)$ and $w(e') = 2$ for each $e' \in E(G) - E(F)$. Clearly, $c(u^*) = \delta$. If $y \in N^\delta_G(u^*)$, since there exists an edge $e \sim y$ such that $e \notin E(F)$, then $c(y) = \sum_{e \sim y} w(e) > \delta$. If $y \in N_G(u^*) - N^\delta_G(u^*)$, then $c(y) \geq d_G(y) > \delta$.

Hence $c(y) \neq c(u^*)$ for all $y \in N_G(u^*)$. For each $xy \in E(G)$, where $x \in U - u^*$ and $y \in W$, by the choice of $F$, we have $c(x) \equiv \delta + 1 \pmod{2}$ and $c(y) \equiv \delta \pmod{2}$. Hence $w$ is a vertex-coloring \{1, 2\}-edge-weighting of the graph $G$.

This completes the proof. \hfill $\Box$

**Corollary 2.4** Let $G$ be a 3-edge-connected bipartite graph. If $3 \leq \delta(G) \leq 5$, then $G$ admits a vertex-coloring $S$-edge-weighting for $S \in \{\{0, 1\}, \{1, 2\}\}$.

**Proof:** Since $3 \leq \delta \leq 5$ and $G$ is 3-edge-connected, then for every vertex $v$ of degree $\delta$, $G - v$ is connected. By Lemma 2.2 and Theorem 1.2 with the same proof, $G$ admits a vertex-coloring $S$-edge-weighting for $S \in \{\{0, 1\}, \{1, 2\}\}$. \hfill $\Box$

3 Conclusions

In this paper, we prove that every 2-connected and 3-edge-connected bipartite graph admits a vertex-coloring $S$-edge-weighting for $S \in \{\{0, 1\}, \{1, 2\}\}$. The generalized $\delta$-graphs is 2-connected and has a vertex-coloring 3-edge-weighting but not vertex-coloring $\{0, 1\}$-edge-weighting or vertex-coloring 2-edge-weighting. So it is an interesting problem to classify all 2-connected bipartite graphs admitting a vertex-coloring $S$-edge-weighting. Since the parity-factor problem is polynomial, then there exists a polynomial algorithm to find a vertex-coloring $S$-edge-weighting of bipartite graphs satisfying the conditions of Theorem 1.2.

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