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# Graph Operators for Coupling-aware Graph Partitioning Algorithms

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CIMI15 July 2015

### Outline

- Introduction
- 2 Algorithms
- 3 Preliminary Experimental Results
- 4 Conclusion & Prospects

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#### Context

Efficiency of numerical simulations requires a distribution of the computational mesh to several processors

## Load Balancing Problem

Distribute equally the load is crucial for the performance of simulations

A common approach is based on graph representation:

- $\bullet \ \ \text{vertex} \to \text{process}$
- edge → dependency

#### Graph Partitioning Problem

Divide the graph in equal parts and assign them to different processors



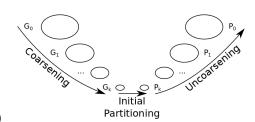
NP-hard problem  $\rightarrow$  heuristic algorithms

## Multilevel Graph Partitioning

Most common graph partitioning techiques use a multilevel framework to simplify the partitioning

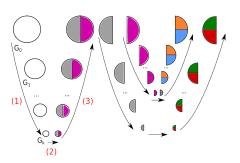
#### 3 phases of a V-cycle:

- Coarsening Phase vertex contraction: HEM
- Initial Partitioning Phase perform partitioning with any known algorithm: bisection, spectral, etc.
- Uncoarsening Phase Projection of the initial partitioning back to the initial graph (refinement)



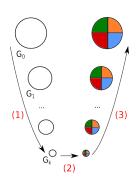
## Multilevel Graph Partitioning

## Multilevel Recursive Bisection (MLRB)



k-1 V-cycles for a k-way partitioning

## Multilevel k-way Partitioning (MLKW)

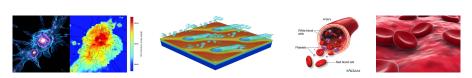


1 V-cycle with direct k-way partitioning

## Challenges

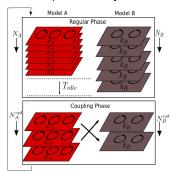
- emerging complex simulations where traditional load balancing strategies are not adequate
- multi-scale, multi-physics, coupled simulations

Domain	Simulation	Туре	Components
astrophysics	structure and	multiplysics	Newtonian gravitational interactions model
	evolution of stars	multiplysics	star aging model
biology	modeling thrombus	multiscale	macroscopic dynamics of bload flow
	development	multiscale	microscopic interactions between cells
aerospace	combustion chamber	multiphysics,	Navier-Stokes equations for fluid flow
engineering	combustion chamber	coupled	heat tranfer solver for solid
environemental science	climate modeling		atmospheric model
		multiphysics,	ocean model
		coupled	sea-ice model
			land model



## Coupled Simulation Model

A coupled simulation consists of a number of component simulations that interact periodically



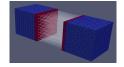
Coupled simulation model between components A and B

#### Regular Phase

- 1. Components solve individual systems
- 2. Parallel execution of components
- 3. Synchronization step  $(T_{idle})$

## Coupling Phase

- 1. Domains intersect coupling interface
- 2. Data exchange between components



#### Related Work

We identify two sub-problems of coupled simulations:

- Resource Distribution Problem: find nb of processors for the components on both phases
- <u>Data Distribution Problem</u>: find a good data distribution for both phases of the coupled simulation

#### Related Work

We identify two sub-problems of coupled simulations:

- Resource Distribution Problem: find nb of processors for the components on both phases
  - empiric approach as AVBP-AVTP [F. Duchaine, et al. (2012)]
  - dynamic adaptation of processors as CESM [Graig, A.P., et al. (2012)]
- Data Distribution Problem:

find a good data distribution for both phases of the coupled simulation

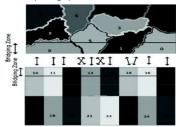
- NAIVE approach: partition each component independently
  - $\rightarrow$  imbalance in coupling phase
- multi-constraint graph partitioning [Karypis, G. (2003)]
  - $\rightarrow$  no control on the nb of processors in different phases



#### Motivation

With a classic graph partitioning, during a coupling phase:

- Load is not well balanced
- Bad inter-components communications



multi-scale simulation in material physics

## Coupling-aware partitioning (co-partitioning)

Partitioning that takes into account the coupling phase explicitly ⇒ as a trade-off to the partitioning in regular phase

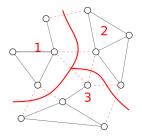
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## Classic Graph Partitioning

k-way graph partitioning  $\rightarrow$  divide graph in k parts. Assign it to k processors



3 parts of weight= 4 and edgecut= 8

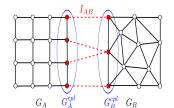
#### Objectives:

- balance weight between parts
- minimize number of edge cut

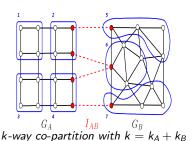
#### Coupled simulations

enrich the classic graph model to take into account multiple components and the coupling phase

## Definition of Co-partitioning



Coupling interface with interedges



#### Graph Model

- $G_X$ : component in regular phase
- $G_X^{cpl}$ : component in coupling phase
- *I<sub>XY</sub>*: inter-component communication

## *k*-way co-partitioning Objectives

- ullet For each X component:
  - 1.  $k_X$  partitioning in regular phase
  - 2.  $k_X^{cpl}$  partitioning in coupling interface
- For each coupling phase XY:
  - 3. minimize inter-component communication

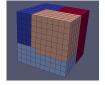
## Algorithms

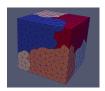
State-of-the-art approach: each component is partitioned independently (NAIVE)





coupling interfaces of meshA and meshB





naive partitions of meshA and meshB

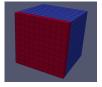
### Problem:

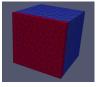
Components are not aware of their coupling interfaces  $\Rightarrow$ 

Imbalance during coupling phase!

## Algorithms

State-of-the-art approach: each component is partitioned independently (NAIVE)

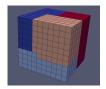


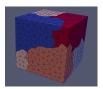


Problem:

Components are not aware of their coupling interfaces ⇒ Imbalance during coupling phase!

coupling interfaces of meshA and meshB





naive partitions of meshA and meshB

We propose two k-way co-partitioning algorithms:

- AWARE
- PROJREPART

Algorithmic description as a sequence of graph operators:

Partition, Restriction, Projection, Repartition, Extension [Predari, M., Esnard, A. (14)]

#### **Partition**

computes a partition of a graph with any algorithm



#### Restriction

finds the subgraph in coupling phase

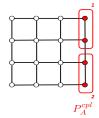


#### Projection

computes a partition similar to another one

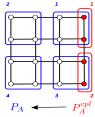


Extension: extends an initially smaller partition to the rest of a graph



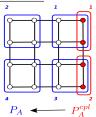
- Input: initial partition  $P_A^{clp}$  on a subgraph
- Constraint: vertices in  $P_A^{clp}$  should remain fixed

Extension: extends an initially smaller partition to the rest of a graph



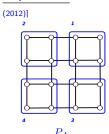
- Input: initial partition  $P_A^{clp}$  on a subgraph
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- Output: partition  $P_A$  of the whole graph

Extension: extends an initially smaller partition to the rest of a graph



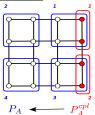
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Repartition: changes the numbers of part in a partition [Vuchener, C., Esnard, A.



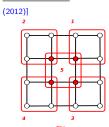
- Input: initial partition  $P_A$  on the whole graph
- Objective: good migration matrix (greedy strategy)

Extension: extends an initially smaller partition to the rest of a graph



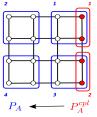
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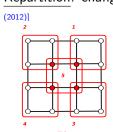
- Input: initial partition  $P_A$  on the whole graph
- Objective: good migration matrix (greedy strategy)
- ullet Output: balanced partition  $P_A^{'}$  on the whole graph

Extension: extends an initially smaller partition to the rest of a graph



- Input: initial partition  $P_A^{clp}$  on a subgraph
- Constraint: vertices in  $P_A^{clp}$  should remain fixed
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- Input: initial partition  $P_A$  on the whole graph
- Objective: good migration matrix (greedy strategy)
- ullet Output: balanced partition  $P_A^{'}$  on the whole graph

Partitioning implementation should handle fixed vertices!

## Graph Partitiong Tools for Initial Fixed Vertices

List of graph partitioning tools that handle fixed vertices.

Туре	Fixed	Parallel	Multilevel Scheme	Initial Part.	Available
graph	yes	no	MLKW	RB	source
graph	only k	no	MLKW	greedy	no
hypergraph	yes	no	MLRB	-	binary
hypergraph	yes	no	MLRB	-	binary
hypergraph	yes	no	MLKW	RB*	no
hypergraph	yes	yes	MLRB	-	source
	graph graph hypergraph hypergraph hypergraph	graph yes graph only k hypergraph yes hypergraph yes hypergraph yes	graph yes no graph only k no hypergraph yes no hypergraph yes no hypergraph yes no	graph         yes         no         MLKW           graph         only k         no         MLKW           hypergraph         yes         no         MLRB           hypergraph         yes         no         MLRB           hypergraph         yes         no         MLKW	graph         yes         no         MLKW         RB           graph         only k         no         MLKW         greedy           hypergraph         yes         no         MLRB         -           hypergraph         yes         no         MLRB         -           hypergraph         yes         no         MLKW         RB*

Limitations of Recursive Bisection (RB) based Algorithms Initial fixed vertices scheme



Partitioning quality of RB





before refinements after refinements

RB fails to partition a simple graph in 4 parts with initial fixed vertices ⇒ inherent numbering constraint VS fixed vertex constraint

## k-way Greedy Graph Growing Partitioning (KGGGP)

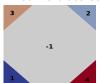
KGGGP: initial partitioning algorithm inside SCOTCH's framework:

- extension of greedy bisection to k parts
- at each step, selection of the best global displacement (v, k)
- selection based on an edgecut minimization criterion
- use of FM-like structures

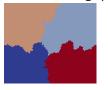
Time Complexity: O(k|E|)

Optimization: local selection of best displacement ([Karypis, G., Kumar, V. (1998)])

Initial fixed vertices scheme



Partitioning quality of KGGGP





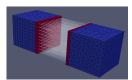
before refinements

after refinements

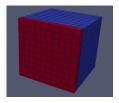
## AWARE Algorithm

#### Each component is aware of it's coupling interface

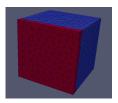
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coupling overview with interedges



mesh A



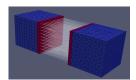
mesh B



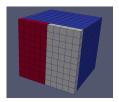
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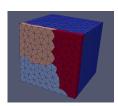
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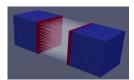
mesh B



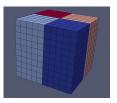
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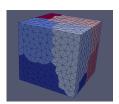
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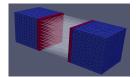


mesh B

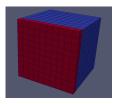


#### Components are also aware of other component's coupling interface

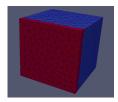
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coupling overview with interedges



mesh A

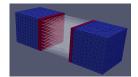


mesh B

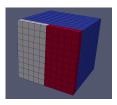


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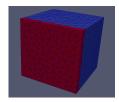
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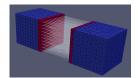


mesh B

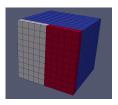


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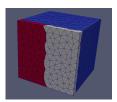
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mesh A

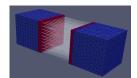


mesh B

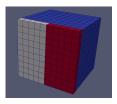


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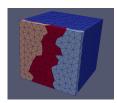
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coupling overview with interedges



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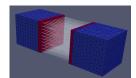


mesh B

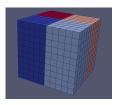


#### Components are also aware of other component's coupling interface

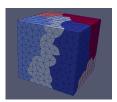
- 1.  $Rest(G_A) \to G_A^{cpl}$ 2.  $Rest(G_B) \to G_B^{cpl}$ 3.  $Part(G_A^{cpl}) \to P_B^{cpl}$ 4.  $Proj(G_A^{cpl}, G_B^{cpl}, P_A^{cpl}) \to \tilde{P}_B^{cpl}$ 5.  $Repart(G_B^{cpl}, \tilde{P}_B^{cpl}) \to P_B^{cpl}$ 6.  $Ext(G_A, G_A^{cpl}, P_A^{cpl}) \to P_A$ 7.  $Ext(G_B, G_B^{cpl}, P_B^{cpl}) \to P_B$



coupling overview with interedges



mesh A



mesh B



### Outline

- Introduction
- 2 Algorithms
- 3 Preliminary Experimental Results
- 4 Conclusion & Prospects

## Experimental Description for KGGGP algorithm

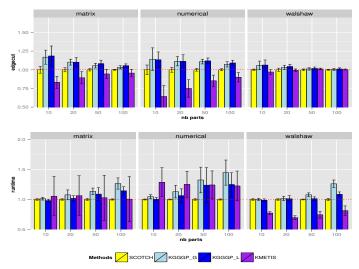
collection	graph	# vtx	# edges	avg d°	min d°	max d°
walshaw	fe_rotor	99 617	662 431	13.30	5	125
walshaw	144	144 649	1 074 393	14.86	4	26
walshaw	wave	156 317	1 059 331	13.55	3	44
walshaw	m14b	214 765	1 679 018	15.64	4	40
matrix	audikw1	943 695	38 354 076	81.28	20	344
matrix	ecology1	1 000 000	1 998 000	4.00	2	4
matrix	thermal2	1 227 087	3 676 134	5.99	2	10
matrix	af_shell10	1 508 065	25 582 130	33.93	14	34
numerical	NACA0015	1 039 183	3 114 818	5.99	3	10
numerical	333SP	3712815	11 108 633	5.98	2	28
numerical	nlr	4 163 763	12 487 976	6.00	3	20
numerical	channel	4802000	42 681 372	17.78	6	18
numerical	adaptive	6815744	13 624 320	4.00	2	4

#### Metrics: partitioning quality, time performance

- our implementations : KGGGP\_G, KGGGP\_L
- tools: SCOTCH 6.0.3, kMETIS 5.1.0, PaToH 3.0, Zoltan 3.81
- imbalance tolerance = 5%
- results from 3 experiments on average values of each graph

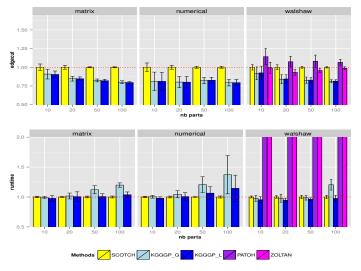
## Exp1: Evaluation of KGGGP without fixed vertices

#### Relative results to SCOTCH



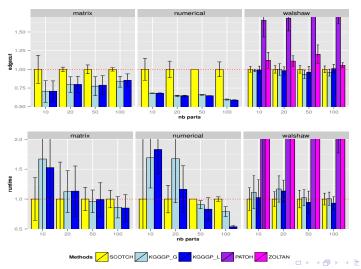
## Exp2: Evaluation of KGGGP with fixed vertices

vertices are fixed based on a bubble scheme (seeds & levelset)



## Exp3: Evaluation of KGGGP with fixed vertices

 vertices are fixed based on a repartition scheme (initial imbalance partition)

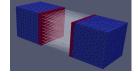


## Experimental Description for co-partitioning

- synthetically generated mesh structures
- use of KGGGP implementation

Description of the experiments

Bescription of the experiments						
	Exp.	Graph A	Graph B	]		
	exp1	cube-hexa-25x25x25	cube-hexa-100x100x100	]		
	exp2	cube-hexa-25x25x25	cube-hexa-70x70x70	l		
	exp3	cube-tetra-40630	cube-hexa-100x100x100	l		
	exp4	cube-tetra-40630	cube-tetra-486719	l		



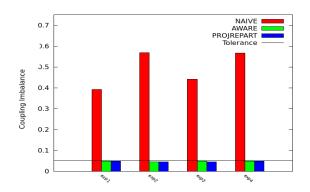
#### Metrics:

- load balancing in coupling phase
- global edgecut
- number of messages exchanged

Results shown for  $k_A = 16$  and  $k_B = 48$ 



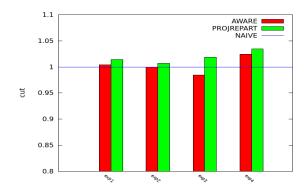
## Load Balancing in Coupling Phase



Load balance during the coupling phase of mesh B

- AWARE and PROJREPART respect the imbalance tolerance (typically 5% of the ideal weight)
- NAIVE highly imbalanced during coupling phase

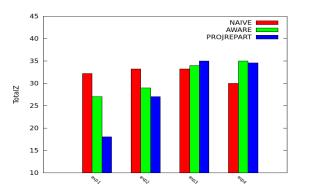
## Edgecut



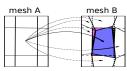
Total edgecut of mesh B

- NAIVE is used as the reference
- AWARE and PROJREPART do not severely degrade the total egdecut

## Number of Messages



mesh A mesh B well-aligned



Total number of messages exchanged during coupling interface

bad-aligned

- exp1 exp2: hexaedric elements are well aligned
- exp3 exp4: elements of different discretization misaligned

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#### Conclusion and Future Work

#### Conclusion

- 1. A new algorithm for graph partitioning with fixed vertices KGGGP
- 2. Two new algorithms for graph partitioning of coupled simulations: AWARE and PROJREPART

#### Experimental results:

- 2. Good load balancing during the coupling phase for both methods at a slight increase of edge-cut
- In simple experiments the number of messages exchanged between the two graph is minimized

#### Prospects for co-partitioning

- More experiments with real-life and larger graphs
- Parallel version



## End

Thank you for you attention. Any questions?