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A combinatorial approach to products of Pisot substitutions

VALÉRIE BERTHÉ¹, JÉRÉMIE BOURDON², TIMO JOLIVET^{1,3}, AND ANNE SIEGEL⁴

¹*LIAFA, CNRS, Université Paris Diderot, Case 7014, 75205 Paris Cedex 13, France*

²*LINA, Université de Nantes, 2 rue de la Houssinière, 44322 Nantes Cedex, France*

³*FUNDIM, Department of Mathematics, University of Turku 20014, Finland*

⁴*INRIA, CNRS, Univ. Rennes 1, IRISA, Campus de Beaulieu, 35042 Rennes Cedex, France*

Abstract

We define a generic algorithmic framework to prove pure discrete spectrum for the substitutive symbolic dynamical systems associated with some infinite families of Pisot substitutions. We focus on the families obtained as finite products of the three-letter substitutions associated with the multidimensional continued fraction algorithms of Brun and Jacobi-Perron.

Our tools consist in a reformulation of some combinatorial criteria (coincidence conditions), in terms of properties of discrete plane generation using multidimensional (dual) substitutions. We also deduce some topological and dynamical properties of the Rauzy fractals, of the underlying symbolic dynamical systems, as well as some number-theoretical properties of the associated Pisot numbers.

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1 Introduction

Symbolic substitutions are morphisms of the free monoid which are often used to generate one-dimensional infinite words and symbolic dynamical systems. They play a prominent role among symbolic dynamical systems with zero entropy, and are intimately connected with first return maps and self-induced systems (see for instance [Fog02, Que10] for more details). Substitutions can also be defined in the framework of higher-dimensional tiling spaces [Sol97, Rob04].

According to the *Pisot substitution conjecture* (see e.g. [Fog02, BK06, BST10, ABB⁺]), the symbolic dynamical system associated with any irreducible substitution whose expansion factor is a Pisot number is conjectured to have pure discrete spectrum. This property is equivalent to being measure-theoretically isomorphic to a translation on a compact abelian group. Rauzy introduced in [Rau82] an explicit construction to establish this property by defining a bounded compact set with fractal boundary (now called the *Rauzy fractal* of the substitution) as a candidate for a fundamental domain for the underlying translation. This approach belongs to a wider scientific program which consists in studying how to code dynamical systems of an arithmetic nature as symbolic systems preserving their arithmetic properties [Sid03], and conversely, how symbolic dynamical systems can provide good simultaneous approximation algorithms in Diophantine approximation [BFZ05]. Following this line of research, the study of Rauzy fractals and Pisot substitutions has led to many developments in various domains, including symbolic dynamics and combinatorics on words, numeration dynamics, hyperbolic dynamics, fractal topology and number theory (see [Fog02, BST10] and the references therein).

It is now well established that the pure discrete spectrum property is algorithmically decidable for a given irreducible Pisot substitution in the unimodular case, that is, when the abelianization matrix of the substitution has determinant ± 1 [ST09, BST10, AL11, AL14]. Nonetheless, no general strategy exists to check whether the pure discrete spectrum property holds for all the substitutions within a given *infinite* family. A particular type of infinite families of substitutions can be obtained by fixing a finite set of substitutions S and by considering all the finite products of substitutions over S ; we call them *product families*. This is for example how Arnoux-Rauzy substitutions are constructed [AR91]. The main goal of the present paper is to pursue Rauzy's program by introducing a *generic computational framework to prove the pure discrete spectrum property* for some product families of three-letter substitutions associated with two-dimensional continued fraction algorithms defined as piecewise fractional maps (according to the formalism developed in [Bre81, Sch00]; see also [Ber11] for a discussion on the way substitutions are associated with continued fraction algorithms). We focus here on the Brun and Jacobi-Perron continued fraction algorithms, and we derive generic methods as well as dynamical, topological and number-theoretical implications.

Discrete spectrum results for infinite families have been established for two-letter unimodular irreducible Pisot substitutions [BD02, HS03], and for Arnoux-Rauzy substitutions (see also [BJS14, BJS12, BŠW13]). Our framework has also been successfully applied to a problem in discrete geometry about the critical thickness at which an arithmetic discrete plane is 2-connected [BJJP13], in connection with the ordered fully subtractive continued fraction algorithm. The modified Jacobi-Perron substitutions are studied as well in [FIY13] using a similar approach (the modified Jacobi-Perron algorithm is a two-point extension of the Brun algorithm). Recently, the pure discrete spectrum property was also announced to be proved for some infinite families of substitutions satisfying some simple combinatorial requirements on their first and last letters [Bar14], which includes β -substitutions associated with simple Parry numbers (see also [Aki00, BBK06]) and some of the families of substitutions considered below.

Combinatorial methods Our strategy relies on two main ingredients. The first one is a combinatorial definition of Rauzy fractals, introduced in [AI01], as the Hausdorff limit of planar (in the three-letter case) compact sets obtained by projecting and renormalizing some three-dimensional objects (“patterns”) consisting of finite unions of faces of unit cubes placed at

integer coordinates. These patterns are obtained by iterating a geometric, higher-dimensional substitution $\mathbf{E}_1^*(\sigma)$. According to [AI01], it can be seen as a *dual substitution* of the original (symbolic) substitution σ . A striking fact is that the resulting three-dimensional patterns obtained by iterating $\mathbf{E}_1^*(\sigma)$ on the three faces of the lower half unit cube $U := \text{◻}$ placed at the origin always lie within a *discrete plane* [AI01, Fer09]. By *discrete plane*, we mean a discretized version of a Euclidean plane; see Section 2.1 for a precise definition.

The second ingredient is the following characterization of the pure discrete spectrum property, due to [IR06], stated in terms of generation of discrete planes by dual substitutions: the symbolic dynamical system generated by a unimodular irreducible Pisot substitution σ has a pure discrete spectrum if and only if the patterns generated by iterating the dual substitution $\mathbf{E}_1^*(\sigma)$ of σ cover arbitrarily large disks in the discrete plane when starting with the initial pattern U . This geometric and combinatorial criterion, which we call the *arbitrarily large disks covering property*, is one of the numerous formulations of the so-called *coincidence conditions* used in substitutive dynamics to prove the pure discrete spectrum property. A survey is given in [BST10, ABB⁺], see also the references in [AL11]. Note that we focus here on the \mathbb{Z} -action of the shift of the substitutive symbolic dynamical system, and not on the \mathbb{R} -action of the substitutive tiling flow. However, for irreducible Pisot substitutions, pure discreteness of both actions are equivalent [CS03].

This characterization of pure discrete spectrum allows the study of finite products of substitutions over S in the following way. First, we prove a stronger result which applies not only to substitutive words, but also to S -adic words, that is, to the limits of infinite products of substitutions [BD14]: the arbitrarily large disks covering property holds when iterating any admissible *infinite* sequence $(\sigma_{i_n})_{n \geq 1}$ of substitutions over the finite set S . *Admissible* refers to the possible constraints on the allowed infinite sequences, which occur for example with the substitutions associated with some continued fraction algorithms. Then, for a given *finite* product $\sigma = \sigma_{i_1} \sigma_{i_2} \cdots \sigma_{i_k}$ over S the same result follows directly by considering the infinite periodic infinite sequence $(\sigma_{i_1} \sigma_{i_2} \cdots \sigma_{i_k})^\infty$. Examples of such sequences with the arbitrarily large disks property are shown in Figure 1 p. 10.

To prove the arbitrarily large disks covering property, we formalize and generalize the approach developed in [IO94], where discrete plane generation using dual Jacobi-Perron substitutions is studied. The intuitive idea is to prove that iterating sufficiently many substitutions eventually generates a topological annulus made of faces around the initial pattern, and that these annuli are preserved under iteration. Several new computational tools are introduced in this article, including generation graphs to prove that annuli are always eventually generated. The strategy is described in detail in Section 3.1, and more precise statements are given below and in Section 5.

Main results We focus on two families of substitutions, the substitutions associated with the Brun algorithm [Bru58], and the substitutions associated with the Jacobi-Perron algorithm [Sch73]. Note that some of the results of this paper have been announced in the extended abstract [BBS13]. Our first results concern discrete plane generation using the dual substitutions $\mathbf{E}_1^*(\sigma)$ associated with these multidimensional continued fraction algorithms. They are stated in Section 5.

- In **Theorem 5.1** we prove that there exist finite patterns V (“seeds”, not much larger than the three-face pattern U) such that iterating Brun or Jacobi-Perron substitutions from V generates an entire discrete plane (*i.e.*, arbitrarily large balls centered at the origin).
- In **Theorem 5.2** (Brun) and **Theorem 5.3** (Jacobi-Perron) we characterize the infinite sequences for which the pattern U does not suffice to generate an entire discrete plane. These characterizations are given in terms of finite state automata with few vertices, which are obtained algorithmically.
- In **Theorem 5.4** we prove that translates of arbitrarily large disks always appear in the images of U .

The above results allow us to deduce several corollaries, of topological, dynamical and number-theoretical nature, including our original motivation to prove pure discrete spectrum.

- In **Corollary 5.5** we characterize all the finite products of Brun or Jacobi-Perron substitutions for which the origin is not an inner point of the Rauzy fractal (this corresponds to the case where U is not a seed). Connectedness results for the corresponding Rauzy fractals are also obtained in **Corollary 5.6**. These topological properties are related with various number-theoretical properties, as discussed in Section 5.4.
- In **Corollary 5.7** we prove the pure discrete spectrum property for every admissible finite product of Brun or Jacobi-Perron substitutions. In **Corollary 5.8** we obtain explicit Markov partitions with connected atoms for the toral automorphisms associated with the incidence matrices of the substitutions.
- In **Theorem 5.9** we prove that for every cubic number field \mathbb{K} , there exist generators α, β with $\mathbb{K} = \mathbb{Q}(\alpha, \beta)$ and a three-letter unimodular irreducible Pisot substitution σ such that the toral translation on \mathbb{T}^2 by (α, β) is measure-theoretically conjugate to the symbolic dynamical system generated by σ . The corresponding Rauzy fractal provides a fundamental domain for the translation, a partition for a natural coding, and bounded remainder sets.

An important consequence of the present framework is that it paves the way for the study of S -adic words. In particular, our results are a key ingredient for providing symbolic representations of two-dimensional toral translations. Indeed, Brun S -adic words are proved in [BST14] (based on the results of the present paper) to provide a symbolic natural coding for *almost every* given toral translation, and not only the ones with algebraic parameters. More precisely, the pure discrete spectrum property for the Brun S -adic shift is proved for almost all of these shifts, and conversely, almost every two-dimensional toral translation admits a symbolic natural coding which is provided by the Brun continued fraction algorithm.

Organization of the paper Substitutions (and dual substitutions), discrete planes, Rauzy fractals, as well as the Brun and Jacobi-Perron continued fractions algorithms are defined in Section 2. The strategy used for generating discrete planes is then elaborated in Section 3 as generically as possible, in the sense that it can be applied to other types of continued fraction algorithms. An outline of the strategy is first given in Section 3.1. Strong coverings are introduced in Section 3.2, sufficient combinatorial criteria to establish the annulus property are given in Section 3.3, and the construction of generation graphs is described in Section 3.4. Section 4 is devoted to more technical aspects and to proofs which are specific to the Brun and Jacobi-Perron substitutions.

Our main results are then stated in Section 5. Results about discrete plane generation using Brun and Jacobi-Perron substitutions are given in Section 5.1. These results are applied in Sections 5.2, 5.3 and 5.4 to prove various properties of the dynamical substitution systems associated with finite products of Brun and Jacobi-Perron substitutions, as well as other implications of number-theoretical nature.

Computer proofs Many of the results of Section 4 have been proved using the Sage mathematics software system [Sag]. The corresponding Sage code is available on the arXiv preprint page of the current article ([arXiv:1401.0704](https://arxiv.org/abs/1401.0704)), as an attached file (“ancillary file”).

A Sage implementation of dual substitutions has been used to perform exhaustive enumerations of small patterns, in order to prove some properties of the families of substitutions under study in Sections 4.1, 4.2 and 4.3. Another type of results for which computer algebra was used is the algorithmic construction of generation graphs, in Sections 4.4 and 4.5.

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2 Preliminaries

2.1 Discrete planes and substitutions

Substitutions Let $\mathcal{A} = \{1, \dots, n\}$ be a finite set of symbols. We work here mainly with three-letter alphabets ($n = 3$). A **substitution** is a non-erasing morphism of the free monoid \mathcal{A}^* ($\sigma(a)$ is a non-empty word for every $a \in \mathcal{A}$). We denote by $\mathbf{P} : \mathcal{A}^* \rightarrow \mathbb{N}^n$ the **abelianization map** defined by $\mathbf{P}(w) = (|w|_1, \dots, |w|_n)$, where $|w|_i$ stands for the number of occurrences of i in w . The **incidence matrix** \mathbf{M}_σ of σ is the matrix of size $n \times n$ whose i th column is equal to $\mathbf{P}(\sigma(i))$ for every $i \in \mathcal{A}$. A substitution σ is **unimodular** if $\det \mathbf{M}_\sigma = \pm 1$, and it is **irreducible Pisot** if the characteristic polynomial of \mathbf{M}_σ is the minimal polynomial of a Pisot number, that is, a real algebraic integer larger than 1 whose other conjugates are smaller than 1 in modulus. It can be proved that every irreducible Pisot substitution is **primitive**, that is, there exists a positive power of its incidence matrix [CS01].

Discrete planes We denote by $(\mathbf{e}_1, \mathbf{e}_2, \mathbf{e}_3)$ the canonical basis of \mathbb{R}^3 . Before defining discrete planes we introduce (pointed) **faces** $[\mathbf{x}, i]^*$, which are defined as subsets of \mathbb{R}^3 by

$$\begin{aligned} [\mathbf{x}, 1]^* &= \{\mathbf{x} + \lambda \mathbf{e}_2 + \mu \mathbf{e}_3 : \lambda, \mu \in [0, 1]\} = \blacktriangleleft \\ [\mathbf{x}, 2]^* &= \{\mathbf{x} + \lambda \mathbf{e}_1 + \mu \mathbf{e}_3 : \lambda, \mu \in [0, 1]\} = \blacktriangleright \\ [\mathbf{x}, 3]^* &= \{\mathbf{x} + \lambda \mathbf{e}_1 + \mu \mathbf{e}_2 : \lambda, \mu \in [0, 1]\} = \blacklozenge \end{aligned}$$

where $i \in \{1, 2, 3\}$ is the **type** of $[\mathbf{x}, i]^*$, and $\mathbf{x} \in \mathbb{Z}^3$ is the **vector** of $[\mathbf{x}, i]^*$. In this paper we will refer to collections of faces as unions of faces, and by abuse of language we will often say that a face f belongs to a union of faces even if f is in fact included in it. The notation $\mathbf{x} + [\mathbf{y}, i]^*$ stands for $[\mathbf{x} + \mathbf{y}, i]^*$.

We now define discrete planes. Denote by $\langle \cdot, \cdot \rangle$ the usual scalar product. Let $\mathbf{v} \in \mathbb{R}_{>0}^3$. The **discrete plane** $\Gamma_{\mathbf{v}}$ **of normal vector** \mathbf{v} is the union of faces $[\mathbf{x}, i]^*$, with $i \in \{1, 2, 3\}$ and $\mathbf{x} \in \mathbb{Z}^3$ satisfying $0 \leq \langle \mathbf{x}, \mathbf{v} \rangle < \langle \mathbf{e}_i, \mathbf{v} \rangle$.

More intuitively, $\Gamma_{\mathbf{v}}$ can also be seen as the boundary of the union of the unit cubes with integer coordinates that intersect the lower half-space $\{\mathbf{x} \in \mathbb{R}^3 : \langle \mathbf{x}, \mathbf{v} \rangle < 0\}$. The set of its vertices in \mathbb{Z}^3 corresponds to the classic notion of a standard arithmetic discrete plane in discrete geometry [Rev91].

Observe that the lower half unit cube $U = [0, 1]^* \cup [0, 2]^* \cup [0, 3]^* = \blacklozenge$ is included in every discrete plane since the coordinates of the normal vector \mathbf{v} of a discrete plane $\Gamma_{\mathbf{v}}$ are assumed to be positive.

A **pattern** is a finite union of (pointed) faces. In order to express the fact that some patterns grow by applying (dual) substitutions, we need to define the **minimal combinatorial radius** $\text{rad}(P)$ of a pattern P containing U . It is equal to the length of the shortest sequence of faces f_1, \dots, f_n in P such that $f_1 \in U$, f_i and f_{i+1} share an edge, and f_n shares an edge with the boundary of P . Intuitively, $\text{rad}(P)$ measures the minimal distance between $\mathbf{0}$ and the boundary of P . A sequence of patterns $(P_n)_{n \geq 1}$ is said to **cover arbitrarily large disks** if the sequence of bounded sets obtained as their orthogonal projections onto the antidiagonal plane $\mathbf{x}_1 + \mathbf{x}_2 + \mathbf{x}_3 = 0$ covers arbitrarily large disks in this latter plane. In particular, if the minimal combinatorial radius of a sequence of patterns $(P_n)_{n \geq 1}$ tends to infinity, then the sequence of patterns $(P_n)_{n \geq 1}$ covers arbitrarily large disks centered at the origin.

Dual substitutions Let σ be a unimodular substitution. The *dual substitution* $\mathbf{E}_1^*(\sigma)$ is defined for any face $[\mathbf{x}, i]^*$ as

$$\mathbf{E}_1^*(\sigma)([\mathbf{x}, i]^*) = \bigcup_{(p,j,s) \in \mathcal{A}^* \times \mathcal{A} \times \mathcal{A}^* : \sigma(j)=pis} [\mathbf{M}_\sigma^{-1}(\mathbf{x} + \mathbf{P}(s)), j]^*.$$

We extend this definition to unions of faces: $\mathbf{E}_1^*(\sigma)(P \cup Q) = \mathbf{E}_1^*(\sigma)(P) \cup \mathbf{E}_1^*(\sigma)(Q)$.

The term “dual” comes from the fact that $\mathbf{E}_1^*(\sigma)$ was originally introduced in [AI01] as the dual of a geometric realization of σ as a linear map in an n -dimensional vector space (where n is the size of the alphabet of σ). This is where the formula given in the above definition comes from. A more general setting is introduced in [SAI01], where the linear maps \mathbf{E}_k and \mathbf{E}_k^* are introduced for every $k \in \{0, \dots, n\}$. Intuitively, $\mathbf{E}_k(\sigma)$ acts on k -dimensional objects and its dual $\mathbf{E}_k^*(\sigma)$ acts on $(n - k)$ -dimensional objects. Some specific examples of dual substitutions are given at the end of Sections 2.2 and 2.3. Note also that the notion of a dual substitution has been successfully extended to the setting of tiling flows (\mathbb{R} -actions) associated with substitutions and under the formalism of strand spaces in [BK06].

Basic properties of dual substitutions are summarized in the proposition below. The first statement ensures that composition behaves well. The second statement can be interpreted as a form of “linearity” of \mathbf{E}_1^* , and allows us to specify a mapping \mathbf{E}_1^* simply by providing \mathbf{M}_σ and the images of $[\mathbf{0}, 1]^*$, $[\mathbf{0}, 2]^*$, $[\mathbf{0}, 3]^*$. The last two statements establish fundamental links between discrete planes and dual substitutions, which will be used throughout this paper.

Proposition 2.1 ([AI01, Fer06]). *Let σ be a unimodular substitution. We have:*

- (1) $\mathbf{E}_1^*(\sigma \circ \sigma') = \mathbf{E}_1^*(\sigma') \circ \mathbf{E}_1^*(\sigma)$ for every unimodular substitution σ' ;
- (2) $\mathbf{E}_1^*(\sigma)([\mathbf{x}, i]^*) = \mathbf{M}_\sigma^{-1}\mathbf{x} + \mathbf{E}_1^*(\sigma)([\mathbf{0}, i]^*)$ for every face $[\mathbf{x}, i]^*$;
- (3) $\mathbf{E}_1^*(\sigma)(\Gamma_{\mathbf{v}}) = \Gamma_{\mathbf{t}_{\mathbf{M}_\sigma}\mathbf{v}}$ for every discrete plane $\Gamma_{\mathbf{v}}$;
- (4) if f and g are distinct faces in a common discrete plane $\Gamma_{\mathbf{v}}$, then $\mathbf{E}_1^*(\sigma)(f) \cap \mathbf{E}_1^*(\sigma)(g)$ contains no face.

2.2 Brun substitutions

Let $\mathbf{v} \in \mathbb{R}_{\geq 0}^3$ such that $0 \leq \mathbf{v}_1 \leq \mathbf{v}_2 \leq \mathbf{v}_3$. The *Brun algorithm* [Bru58] is one of the possible natural generalizations of Euclid’s algorithm. Together with the Jacobi-Perron algorithm, it is one of the most classical multidimensional continued fraction algorithms; they both are defined as piecewise fractional maps according the formalism developed in [Bre81, Sch00]. We consider here the additive version of this algorithm, which can be defined as follows in its linear form: subtract the second largest component of \mathbf{v} from the largest and reorder the result. This yields for the algorithm in ordered form:

$$\mathbf{v} \mapsto \begin{cases} (\mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3 - \mathbf{v}_2) & \text{if } \mathbf{v}_1 \leq \mathbf{v}_2 \leq \mathbf{v}_3 - \mathbf{v}_2 \\ (\mathbf{v}_1, \mathbf{v}_3 - \mathbf{v}_2, \mathbf{v}_2) & \text{if } \mathbf{v}_1 \leq \mathbf{v}_3 - \mathbf{v}_2 < \mathbf{v}_2 \\ (\mathbf{v}_3 - \mathbf{v}_2, \mathbf{v}_1, \mathbf{v}_2) & \text{if } \mathbf{v}_3 - \mathbf{v}_2 < \mathbf{v}_1 \leq \mathbf{v}_2. \end{cases}$$

The interest of working with the additive version of the algorithm is that we will recover a finite set of associated substitutions, as described below. Iterating this map starting from $\mathbf{v}^{(0)} = \mathbf{v}$ yields an infinite sequence of vectors $(\mathbf{v}^{(n)})_{n \in \mathbb{N}}$ and the algorithm can be rewritten in matrix form:

$$\mathbf{v}^{(n)} = \mathbf{M}_{i_n}^{-1} \mathbf{v}^{(n-1)}, \tag{1}$$

where $i_n \in \{1, 2, 3\}$ and \mathbf{M}_1 , \mathbf{M}_2 and \mathbf{M}_3 are the corresponding non-negative integer matrices such that \mathbf{M}_1^{-1} is applied if $\mathbf{v}_1 \leq \mathbf{v}_2 \leq \mathbf{v}_3 - \mathbf{v}_2$, \mathbf{M}_2^{-1} is applied if $\mathbf{v}_1 \leq \mathbf{v}_3 - \mathbf{v}_2 < \mathbf{v}_2$, and \mathbf{M}_3^{-1} is applied otherwise. We thus have $\mathbf{v} = \mathbf{M}_{i_1} \dots \mathbf{M}_{i_n} \mathbf{v}^{(n)}$ for all n . We stress on the fact that the matrices \mathbf{M}_i are non-negative, which is needed to associate substitutions with them.

The **Brun expansion** of a vector $\mathbf{v} \in \mathbb{R}_{\geq 0}^3$ with $0 \leq \mathbf{v}_1 \leq \mathbf{v}_2 \leq \mathbf{v}_3$ is the infinite sequence $(i_n)_{n \geq 1}$ obtained above. For example, if $\mathbf{v} = (1, \beta, \beta^2)$, with $\beta \approx 3.21$ being the dominant Pisot cubic root of $x^3 - 3x^2 - x + 1$, then $(i_n) = 1131\ 132\ 132\ \dots$ is an eventually periodic sequence containing infinitely many 1s, 2s and 3s.

Proposition 2.2 ([Bru58]). *The Brun expansion of $\mathbf{v} \in \mathbb{R}_{\geq 0}^3$ contains infinitely many 3s if and only if \mathbf{v} is totally irrational. Moreover, for every expansion $(i_n)_{n \geq 1} \in \{1, 2, 3\}^{\mathbb{N}}$ containing infinitely many 3s, there is a unique vector \mathbf{v} whose Brun expansion is $(i_n)_{n \geq 1}$.*

The **Brun substitutions** are defined by

$$\sigma_1^{\text{Br}} : \begin{cases} 1 \mapsto 1 \\ 2 \mapsto 2 \\ 3 \mapsto 32 \end{cases} \quad \sigma_2^{\text{Br}} : \begin{cases} 1 \mapsto 1 \\ 2 \mapsto 3 \\ 3 \mapsto 23 \end{cases} \quad \sigma_3^{\text{Br}} : \begin{cases} 1 \mapsto 2 \\ 2 \mapsto 3 \\ 3 \mapsto 13 \end{cases}$$

and we denote their associated dual substitutions by $\Sigma_i^{\text{Br}} = \mathbf{E}_1^*(\sigma_i^{\text{Br}})$ for $i \in \{1, 2, 3\}$. They can be explicitly computed using the definition of dual substitutions, as shown below.

$$\Sigma_1^{\text{Br}} : \begin{cases} \blacktriangle \mapsto \blacktriangle \\ \blacktriangle \mapsto \blacklozenge \\ \blacklozenge \mapsto \blacklozenge \end{cases} \quad \Sigma_2^{\text{Br}} : \begin{cases} \blacktriangle \mapsto \blacktriangle \\ \blacktriangle \mapsto \blacklozenge \\ \blacklozenge \mapsto \blacklozenge \end{cases} \quad \Sigma_3^{\text{Br}} : \begin{cases} \blacktriangle \mapsto \blacklozenge \\ \blacktriangle \mapsto \blacktriangle \\ \blacklozenge \mapsto \blacklozenge \end{cases}$$

These substitutions and the Brun algorithm are linked by $\mathbf{M}_i = {}^t \mathbf{M}_{\sigma_i^{\text{Br}}}$ for $i \in \{1, 2, 3\}$, where the \mathbf{M}_i are the matrices associated with the Brun algorithm above. By virtue of Proposition 2.1, it follows that, if $\mathbf{v} \in \mathbb{R}_{> 0}^3$ with $0 < \mathbf{v}_1 \leq \mathbf{v}_2 \leq \mathbf{v}_3$, then $\Sigma_i^{\text{Br}}(\Gamma_{\mathbf{v}}) = \Gamma_{\mathbf{M}_i \mathbf{v}}$. An explicit link between the patterns $(\Sigma_{i_1} \cdots \Sigma_{i_n}(U))_{n \geq 1}$ and the discrete plane $\Gamma_{\mathbf{v}}$ (where the expansion of \mathbf{v} is $(i_n)_{n \geq 1}$) is given in Theorem 5.2. Concrete instances of computations using dual substitutions are found in Example 3.5, and in many of the proofs of Section 4.

In this paper we are interested in the finite products $\sigma = \sigma_{i_1}^{\text{Br}} \cdots \sigma_{i_n}^{\text{Br}}$ such that the dominant eigenvalue of \mathbf{M}_{σ} is a cubic Pisot number. Such finite products are characterized in Proposition 2.3 below. Note that such a simple characterization does not necessarily exist for general product families of substitutions.

Proposition 2.3 ([AD13]). *For every $(i_1, \dots, i_n) \in \{1, 2, 3\}^n$, the product $\sigma_{i_1}^{\text{Br}} \cdots \sigma_{i_n}^{\text{Br}}$ is irreducible Pisot if and only if $i_k = 3$ at least once.*

Proposition 2.3 above leads us to consider only the Brun expansions $(i_n)_{n \in \mathbb{N}} \in \{1, 2, 3\}^{\mathbb{N}}$ in which $i_n = 3$ infinitely often. Such sequences are called **Brun-admissible**.

Remark 2.4. The definition of Brun substitutions, and more generally of substitutions associated with a continued fraction algorithm is not canonical, since the matrices associated with the continued fraction algorithm do not impose a specific ordering of the letters in the substitutions. This specific ordering of letters is the one that yields the simplest technical proofs in Section 4, but, in the case of Brun substitutions, the same results can be obtained for any other ordering using the same techniques.

2.3 Jacobi-Perron substitutions

Let $\mathbf{v} \in \mathbb{R}_{\geq 0}^3$ be such that $\mathbf{v}_1 \leq \mathbf{v}_3$ and $\mathbf{v}_2 \leq \mathbf{v}_3$. The **Jacobi-Perron algorithm** [Jac68, Per07] consists in iterating the map

$$\mathbf{v} \mapsto (\mathbf{v}_2 - a\mathbf{v}_1, \mathbf{v}_3 - b\mathbf{v}_1, \mathbf{v}_1), \quad \text{where } a = \lfloor \mathbf{v}_2/\mathbf{v}_1 \rfloor, \quad \text{and } b = \lfloor \mathbf{v}_3/\mathbf{v}_1 \rfloor.$$

Like with the Brun algorithm, we obtain an infinite sequence of vectors $\mathbf{v}^{(0)} = \mathbf{v}, \mathbf{v}^{(1)}, \mathbf{v}^{(2)}, \dots$ such that $\mathbf{v}^{(n)} = \mathbf{M}_{a_n, b_n}^{-1} \mathbf{v}^{(n-1)}$ where $\mathbf{M}_{a, b}$ is a matrix with non-negative integer entries. The

Jacobi-Perron expansion of \mathbf{v} is the infinite sequence $(a_n, b_n)_{n \geq 1}$. It can be proved that $(a_n, b_n)_{n \geq 1}$ is the Jacobi-Perron expansion of some vector \mathbf{v} if and only if for every $n \geq 1$, we have $0 \leq a_n \leq b_n$, $b_n \neq 0$, and $a_n = b_n$ implies $a_{n+1} \neq 0$ (see [Sch73, Sch00, Bre81]).

A sequence $(a_n, b_n)_{n \geq 1}$ is **Jacobi-Perron-admissible** if, for every $n \geq 1$, we have $0 \leq a_n \leq b_n$, $b_n \neq 0$, and $a_n = b_n$ implies $a_{n+1} \neq 0$. Here again, we can state a convenient characterization of expansions of totally irrational vectors.

Proposition 2.5 ([Per07]). *A sequence $(a_n, b_n)_{n \geq 1}$ is a Jacobi-Perron-admissible sequence if and only if it is the Jacobi-Perron expansion of totally irrational vector $\mathbf{v} \in \mathbb{R}_{\geq 0}^3$. Moreover, for every such expansion $(a_n, b_n)_{n \geq 1}$, there is a unique vector \mathbf{v} whose expansion is $(a_n, b_n)_{n \geq 1}$.*

Let the **Jacobi-Perron substitutions** be defined by $\sigma_{a,b}^{\text{JP}} : 1 \mapsto 3, 2 \mapsto 13^a, 3 \mapsto 23^b$ for all $a, b \geq 0$. Their associated duals are defined by $\Sigma_{a,b}^{\text{JP}} = \mathbf{E}_1^*(\sigma_{a,b}^{\text{JP}})$ and are given by

$$\Sigma_{a,b}^{\text{JP}} : \begin{cases} [0, 1]^* & \mapsto [a\mathbf{e}_1, 2]^* \\ [0, 2]^* & \mapsto [b\mathbf{e}_1, 3]^* \\ [0, 3]^* & \mapsto [0, 1]^* \cup \bigcup_{k=0}^{a-1} [k\mathbf{e}_1, 2]^* \cup \bigcup_{k=0}^{b-1} [k\mathbf{e}_1, 3]^*. \end{cases}$$

These substitutions are chosen in such a way that $\mathbf{M}_{a,b} = {}^t\mathbf{M}_{\sigma_{a,b}^{\text{JP}}}$ for every $a, b \geq 0$, so the correspondence with dual substitutions and discrete planes is analogous to the one described above for the Brun substitutions. We also have:

Proposition 2.6 ([DFPLR04]). *Every product $\sigma_{a_1, b_1}^{\text{JP}} \cdots \sigma_{a_n, b_n}^{\text{JP}}$ is irreducible Pisot if $0 \leq a_n \leq b_n$ and $b_n \neq 0$ for all $n \geq 1$.*

Note that, unlike Brun substitutions, Jacobi-Perron substitutions do not consist of a finite number of substitutions, but of infinitely many substitutions parametrized by $a, b \in \mathbb{N}$. In essence this reflects the fact that the associated continued fraction algorithm is *multiplicative*, in opposition with the *additive* version of the Brun algorithm given in Section 2.2. Such notions are defined formally in [Sch00].

Additive version of Jacobi-Perron substitutions When dealing with Jacobi-Perron substitutions in Section 3, we will work directly on the substitutions $\Sigma_{a,b}^{\text{JP}}$, except in Section 4.5, where the generation graphs that we will construct apply only to products of *finitely many* substitutions (but there are infinitely many substitutions $\sigma_{a,b}^{\text{JP}}$). Hence we will need to decompose $\sigma_{a,b}^{\text{JP}}$ in products using finitely many different substitutions. A first natural additive substitutive realization of Jacobi-Perron algorithm is given by $\tau_1 : 1 \mapsto 1, 2 \mapsto 21, 3 \mapsto 3$, $\tau_2 : 1 \mapsto 1, 2 \mapsto 2, 3 \mapsto 31$, $\tau_3 : 1 \mapsto 3, 2 \mapsto 1, 3 \mapsto 2$, with $\sigma_{a,b}^{\text{JP}} = \tau_3 \tau_1^a \tau_2^b$. (Note that $\tau_1 \tau_2 = \tau_2 \tau_1$.) The problem with this decomposition is that it is not restrictive enough, which makes the generation graphs associated with this family too difficult to handle.

So we choose another additive decomposition, which enforces the constraint $0 \leq a \leq b$ and $b \geq 1$ to be satisfied in the additive products. By setting $\theta_1 = \tau_2$, $\theta_2 = \tau_1 \tau_2$, $\theta_3 = \tau_3 \tau_2$ and $\theta_4 = \tau_3 \tau_1 \tau_2$, we have

$$\sigma_{a,b}^{\text{JP}} = \begin{cases} \theta_3 \theta_1^{b-1} & \text{if } a = 0 \\ \theta_3 \theta_1^{b-a-1} \theta_2^a = \theta_4 \theta_1^{b-a} \theta_2^{a-1} & \text{if } 0 < a < b \\ \theta_4 \theta_2^{a-1} & \text{if } a = b. \end{cases}$$

Finally, we define $\Theta_i = \mathbf{E}_1^*(\theta_i)$ for $i \in \{1, 2, 3, 4\}$. These are the substitutions that will be used in Section 3.4 to construct generation graphs for the Jacobi-Perron substitutions.

Remark 2.7. Note that the ‘‘rhythm’’ is provided by θ_3 and θ_4 , in the sense that every product $\theta_{i_1} \cdots \theta_{i_n}$ starting with θ_3 or θ_4 can be uniquely decomposed into a product of $\sigma_{a_k, b_k}^{\text{JP}}$. Consequently, if an infinite sequence $(i_n)_{n \geq 1} \in \{1, 2, 3, 4\}^{\mathbb{N}}$ is the additive expansion of an admissible Jacobi-Perron expansion, then it must contain infinitely many occurrences of 3 and 4.

2.4 Substitutive dynamics and symbolic codings

Pure discrete spectrum We endow $\mathcal{A}^{\mathbb{Z}}$ with the product of discrete topologies. Let S stand for the (two-sided) shift map over $\mathcal{A}^{\mathbb{Z}}$, that is, $S((u_n)_{n \in \mathbb{Z}}) = (u_{n+1})_{n \in \mathbb{Z}}$. Let σ be a primitive substitution over \mathcal{A} and let $u \in \mathcal{A}^{\mathbb{Z}}$ be such that $\sigma^k(u) = u$ for some $k \geq 1$ and such that $u_{-1}u_0$ is a factor of some $\sigma^\ell(i)$, $i \in \mathcal{A}$ (such a word u exists by primitivity of σ). Let $\overline{\mathcal{O}(u)}$ be the orbit closure of the two-sided word u under the action of the shift S . The substitutive symbolic dynamical system (X_σ, S) generated by σ is defined as $X_\sigma = \overline{\mathcal{O}(u)}$. One easily checks by primitivity that (X_σ, S) does not depend on the choice of the two-sided word u fixed by some power of σ . For more details, see [Que10, Fog02].

Let μ be a shift-invariant Borel probability measure for (X_σ, S) . An eigenfunction for (X_σ, S, μ) is an $L^2(X_\sigma, \mu)$ function f for which there is an associated eigenvalue $\alpha \in \mathbb{C}$ (with modulus 1) such that $f \circ S = \alpha f$. The symbolic dynamical system (X_σ, S) is said to have **pure discrete spectrum** if the linear span of the eigenfunctions is dense in $L^2(X_\sigma, \mu)$. Note that measure-theoretic discrete spectrum and topological discrete spectrum are proved to be equivalent for primitive substitutive dynamical systems [Hos86].

Rauzy fractals Let σ be a unimodular irreducible Pisot substitution with incidence matrix \mathbf{M}_σ . We denote by π_c the projection onto the contracting eigenhyperplane of \mathbf{M}_σ along the expanding eigendirection of \mathbf{M}_σ . The **Rauzy fractal** \mathcal{T}_σ associated with σ is the Hausdorff limit of the sequence $(\mathcal{D}_n)_{n \geq 0}$ with $\mathcal{D}_n = \mathbf{M}_\sigma^n \circ \pi_c \circ \mathbf{E}_1^*(\sigma)^n(U)$ for all n , with $U = [\mathbf{0}, 1]^* \cup [\mathbf{0}, 2]^* \cup [\mathbf{0}, 3]^* = \diamond$. It is divided into the **subtiles** $\mathcal{T}_\sigma(i)$, for $i \in \mathcal{A}$, provided by the Hausdorff limit of the sequence of compact sets $\mathcal{D}_n(i) = \mathbf{M}_\sigma^n \circ \pi_c \circ \mathbf{E}_1^*(\sigma)^n([\mathbf{0}, i]^*)$ for $n \geq 0$. Examples of Rauzy fractals are given in Figure 4 p. 26.

A Rauzy fractal provides a geometrical realization for the substitutive dynamics of a unimodular irreducible Pisot substitution σ . First, a domain exchange transformation (\mathcal{T}_σ, E) can be defined acting on the three subtiles of the Rauzy fractal; it is defined by translating each subtile $\mathcal{T}_\sigma(i)$ by the vector $\pi_c(\mathbf{e}_i)$. A combinatorial condition called the strong coincidence condition guarantees that the subtiles $\mathcal{T}_\sigma(i)$ are disjoint in measure [AI01]. The domain exchange transformation (\mathcal{T}_σ, E) is proved to be measure-theoretically conjugate to (X_σ, S) if σ satisfies the strong coincidence condition. Furthermore this domain exchange transformation is conjectured to factor onto a toral translation, according to the Pisot conjecture. The underlying lattice for the factorization is given by the vectors $\pi_c(\mathbf{e}_i - \mathbf{e}_j)$ for $i \neq j$. This can be reformulated in terms of tilings: (\mathcal{T}_σ, E) is measure-theoretically conjugate to a translation (via this factorization) if and only if it translates by the vectors of the lattice $\pi_c(\mathbb{Z}^2)$ of the subtiles of the Rauzy fractal tile the contracting eigenhyperplane of \mathbf{M}_σ . For more details, see [AI01, CS01]. Second, when pure discrete spectrum holds, the Rauzy fractal and its subtiles can be used to construct a Markov partition of the toral automorphism provided by the incidence matrix \mathbf{M}_σ of σ (see [IO93, KV98, Pra99, IR06, Sie00]). This will be developed in Section 5.3.

Symbolic natural codings and bounded remainder sets Given an invertible transformation T on a set X and a partition $\mathcal{P} = \{P_0, \dots, P_{k-1}\}$ of X , a **coding** of $x \in X$ is a sequence $u \in \{0, \dots, k-1\}^{\mathbb{Z}}$ such that $u_n = i$ if $T^n x \in P_i$, for all $n \in \mathbb{Z}$. A symbolic dynamical system (Ω, S) is a **symbolic coding** of (X, T) if there exists a finite partition of X such that every element of Ω is a coding of the orbit of some point of X , and if, furthermore, (Ω, S) and (X, T) are measure-theoretically conjugate. A **toral translation** on \mathbb{T}^2 is a map $R_\alpha : \mathbb{T}^2/L \rightarrow \mathbb{T}^2/L$, $x \mapsto x + \alpha \pmod{L}$, where $\alpha \in \mathbb{R}^2$, and where L is a full rank lattice in \mathbb{R}^2 . A symbolic dynamical system (Ω, S) is a **symbolic natural coding** of $(\mathbb{T}^2/L, R_\alpha)$ if it is a symbolic coding defined with respect to a fundamental domain for the lattice L in \mathbb{R}^2 together with a finite partition of this fundamental domain such that the restriction of the map R_α is provided by the translation by some vector on each element of the partition. Lastly, a set $Y \subset X$ is called a **bounded remainder set** for the minimal dynamical system (X, T, μ) if its indicator function

yields bounded Birkhoff sums, that is, if there exists $C > 0$ such that for a.e. $x \in X$, we have $|\#\{n \in \mathbb{Z} : |n| \leq N, T^n(x) \in Y\} - (2N + 1)\mu(Y)| \leq C$. In case of pure discrete spectrum of (X_σ, S) , the Rauzy subtiles $\mathcal{T}(i)$ provide bounded remainder sets for the associated toral translation (for more details, see the proof of Corollary 5.7).

3 Generating discrete planes with substitutions

In this section we establish a generic strategy to generate discrete planes with substitutions associated with multidimensional continued fraction algorithms, in the sense that this strategy can be applied to other product families of substitutions and continued fraction algorithms. We have in mind fibred algorithms in the sense of [Sch95], which are Markovian (“without memory”) and which are defined by piecewise fractional maps such as considered in [Bre81, Sch00]. The underlying matrices are thus nonnegative unimodular matrices, so we can use the formalism of dual substitutions.

A detailed outline of the strategy is first given in Section 3.1. In Section 3.2 we introduce *strong coverings*, which serve as a combinatorial restriction on the patches of discrete planes, and which are crucial in order to prove the results of Section 3.3, where we give sufficient combinatorial criteria to establish the *annulus property*. In Section 3.4 we construct *generation graphs* to characterize the sequences of substitutions that fail to generate an entire discrete plane when starting from a finite pattern.

3.1 Outline of the strategy

We recall that a *pattern* is a finite union of faces and that a *seed* is a pattern V such that iterating any sequence of respectively Brun or Jacobi-Perron dual substitutions from V yields patterns with arbitrarily large minimal combinatorial radii centered at the origin. We will provide adequate tools in order to prove Theorem 5.1 (existence of seeds), Theorems 5.2 and 5.3 (characterization of sequences of which U is not a seed), and Theorem 5.4 (iterating from U generates translates of arbitrarily large disks).

This is illustrated by the following example for Brun substitutions. Let $(i_n) = 232\ 232\ \dots$ and $(j_n) = 2311\ 2311\ \dots$ be two infinite (periodic) sequences. Figure 1 (left) suggests that $\Sigma_{i_1}^{\text{Br}} \dots \Sigma_{i_n}^{\text{Br}}(U)$ covers arbitrarily large disks centered at the origin as $n \rightarrow \infty$, but Figure 1 (right) suggests that $\Sigma_{j_1}^{\text{Br}} \dots \Sigma_{j_n}^{\text{Br}}(U)$ does **not** cover arbitrarily large disks centered at the origin as $n \rightarrow \infty$. This can be proved rigorously thanks to Theorem 5.3, and Theorem 5.4 guarantees that *translates* of arbitrarily large disks do occur.

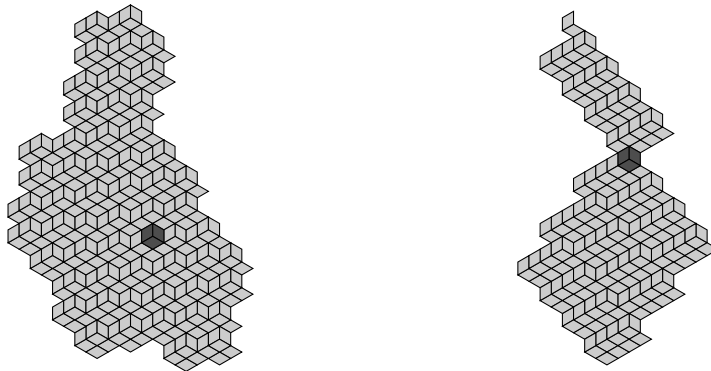


Figure 1: On the left, the pattern $(\Sigma_2^{\text{Br}} \Sigma_3^{\text{Br}} \Sigma_2^{\text{Br}})^4(U)$. On the right, the pattern $(\Sigma_2^{\text{Br}} \Sigma_3^{\text{Br}} \Sigma_1^{\text{Br}} \Sigma_1^{\text{Br}})^3(U)$. The pattern U at the origin is shown in dark gray. The Rauzy fractals associated with these two products of substitutions are plotted in Figure 4 p. 26. We will see that in both cases, the arbitrarily large disks covering property is satisfied.

Annuli and covering properties The patterns generated by iterating dual substitutions can have complicated shapes, so there is no obvious way of proving that such a sequence covers arbitrarily large disks. One approach to prove this property in the context of dual substitutions has been initiated by Ito and Ohtsuki [IO94]. The idea is to make sure that the generated patterns contain an increasing number of nested concentric annuli of positive width, and hence cover arbitrarily large disks. This requires that the *annulus property* holds, *i.e.*, that the image of an annulus by a dual substitution remains an annulus.

The action of a dual substitution is governed mainly by the action of the inverse of the incidence matrix of the substitution; roughly speaking, a dual substitution can be seen as a discrete analogue of an affine transformation. Affine transformations preserve the topological property of being an annulus. However the discretization step makes things more complicated, and there exist annuli not preserved under the action of a dual substitution: the annulus property is wrong if no additional assumptions are made on the annuli, as shown in Example 3.5.

One way around this is the notion of \mathcal{L} -covering (introduced in [IO94]), which intuitively means that a pattern is path-connected by paths of contiguous connected patterns (the elements of \mathcal{L}), which provides some combinatorial restrictions on the annuli. Contrary to as stated in [IO94], \mathcal{L} -coverings alone are not sufficient to ensure that the image of an annulus remains an annulus. We thus have extended this notion by considering patterns of more than two faces to be able to find suitable sets \mathcal{L} for the substitutions we study. (In particular, the set of 7 patterns given in [IO94] for Jacobi-Perron substitutions is not sufficient.)

We introduce and discuss covering properties in Section 3.2, as well as the stronger condition (strong \mathcal{L} -covering) under which the desired annulus property holds. These restrictions allow us to prove the annulus property in Section 3.3. The proofs require many verifications which are specific to the set of substitutions under study. Hence the results given in Sections 3.2 and 3.3 are stated in a generic way (they do not depend on the choice of substitutions), and proving that they hold for our families of substitutions is reduced to technical verifications, which are carried out in Section 4.

Generation graphs The annulus property alone is not sufficient: we must also make sure that, starting from an initial pattern expected to be a candidate seed, at least one annulus is eventually generated by iterating sufficiently many substitutions.

A first natural strategy to check this is to track all the possible images of the initial pattern U under iteration of admissible products, and prove that an annulus centered at the origin is always eventually generated around U . This is the strategy which was originally taken in [IO94] with the Jacobi-Perron substitutions, by constructing graphs whose vertices are patterns. The problem with this approach is that the obtained graphs are too large and difficult to analyze, and require the proof of numerous lemmas. In some cases, like with the Brun substitutions, this approach is impracticable; this is mostly due to the fact that for some admissible sequences of Brun (or Jacobi-Perron), the pattern U is not sufficient to generate an entire discrete plane, as illustrated in Figure 1 (right).

Hence we introduce in Section 3.4 a new tool, *generation graphs*, which are much more manageable (the vertices are single faces and not patterns), which can be constructed algorithmically, and which contain the information needed about the annulus generation properties around seeds. This prevents us from having to deal with numerous cases by hand (in contrast with the graphs described in [IO94] for example). The aim of generation graphs is to track all the possible sequences of preimages of each *face* of the annuli that we want to generate, instead of tracking all the possible patterns around a given initial pattern. The construction given in Section 3.4 is generic, and specific generation graphs are computed in Sections 4.4 and 4.5 for the Brun and Jacobi-Perron substitutions.

Remark 3.1. Given a *single* unimodular irreducible Pisot substitution σ , it is not difficult to see that there exists a finite seed V such that iterating $\mathbf{E}_1^*(\sigma)$ from V generates an entire

discrete plane. This follows from the fact that \mathbf{M}_σ is contracting on the contracting plane of \mathbf{M}_σ thanks to the Pisot condition, so the discrete affine map $\mathbf{E}_1^*(\sigma)$ is also contracting on the associated (invariant) discrete plane (for a proof, see [BST10]). In the case of a product family of substitutions, the existence of such finite seeds V is not guaranteed, especially if some of the substitutions are not Pisot (for example σ_3^{Br} is Pisot, but σ_1^{Br} and σ_2^{Br} are not). Constructing generation graphs will allow us to prove the existence of seeds even in such cases.

3.2 Covering properties

We now introduce \mathcal{L} -coverings and strong \mathcal{L} -coverings, which are the combinatorial tools that will be used in order to prove the annulus property in Section 3.3.

A pattern is said to be **edge-connected** if any two faces are connected by a path of faces f_1, \dots, f_n such that f_k and f_{k+1} share an edge, for all $k \in \{1, \dots, n-1\}$. In the definitions below, \mathcal{L} will always denote a set of patterns which is closed by translation of \mathbb{Z}^3 , so we will define such sets by giving only one element of each translation class. Let \mathcal{L} be a set of patterns. A pattern P is **\mathcal{L} -covered** if for all faces $e, f \in P$, there exist patterns $Q_1, \dots, Q_n \in \mathcal{L}$ such that

- (1) $e \in Q_1$ and $f \in Q_n$;
- (2) $Q_k \cap Q_{k+1}$ contains at least one face, for all $k \in \{1, \dots, n-1\}$;
- (3) $Q_k \subseteq P$ for all $k \in \{1, \dots, n\}$.

The next proposition, due to [IO94], gives a sufficient combinatorial criterion to prove that a dual substitution preserves \mathcal{L} -covering.

Proposition 3.2. *Let \mathcal{L} be a set of patterns, P be an \mathcal{L} -covered pattern and Σ be a dual substitution. If $\Sigma(Q)$ is \mathcal{L} -covered for every $Q \in \mathcal{L}$, then $\Sigma(P)$ is \mathcal{L} -covered.*

We will see in Example 3.5 that \mathcal{L} -coverings are not sufficient for our purposes, so we introduce **strong \mathcal{L} -coverings**. A pattern P is **strongly \mathcal{L} -covered** if P is \mathcal{L} -covered and if for every pattern $X \subseteq P$ that is edge-connected and consists of two faces, there exists a pattern $Y \in \mathcal{L}$ such that $X \subseteq Y \subseteq P$.

Example 3.3. Let $\mathcal{L}^{\text{Br}} = \{\text{[diagram 1]}, \text{[diagram 2]}, \text{[diagram 3]}, \text{[diagram 4]}, \text{[diagram 5]}, \text{[diagram 6]}, \text{[diagram 7]}\}$ (a set of patterns which will later be used with the Brun substitutions in Section 4), and let $P_1 = \text{[diagram 8]}$, $P_2 = \text{[diagram 9]}$, $P_3 = \text{[diagram 10]}$, $P_4 = \text{[diagram 11]}$. Then P_1 is neither \mathcal{L}^{Br} - nor \mathcal{L}^{JP} -covered; P_2 is not \mathcal{L}^{Br} -covered, but it is strongly \mathcal{L}^{JP} -covered; P_3 is \mathcal{L}^{Br} - and \mathcal{L}^{JP} -covered, but not strongly covered.; and P_4 is strongly \mathcal{L}^{Br} - and \mathcal{L}^{JP} -covered.

3.3 The annulus property

In this section we define \mathcal{L} -annuli and we introduce Property A, which is a sufficient condition under which we can prove that a dual substitution Σ verifies the **annulus property**: if A is an \mathcal{L} -annulus of P , then $\Sigma(A)$ is an \mathcal{L} -annulus of $\Sigma(P)$ (Proposition 3.7).

The **boundary** ∂P of a pattern P is the union of the edges e of the faces of f such that e is contained in one face only. An **\mathcal{L} -annulus** of a simply connected pattern P is a pattern A such that A is strongly \mathcal{L} -covered, A and P have no face in common, and $P \cap \partial(P \cup A) = \emptyset$. Note that the condition $P \cap \partial(P \cup A) = \emptyset$ in the above definition is a concise way to express the fact that the \mathcal{L} -covered set A is a “good surrounding” of U .

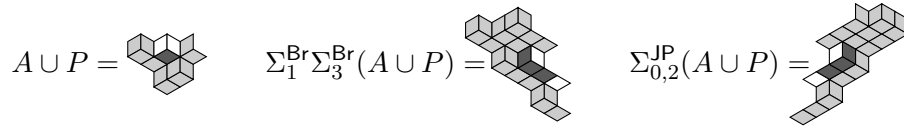
Example 3.4. Let A_1, A_2, A_3 and A_4 be defined by

$$U \cup A_1 = \text{[diagram 12]} \quad U \cup A_2 = \text{[diagram 13]} \quad U \cup A_3 = \text{[diagram 14]} \quad U \cup A_4 = \text{[diagram 15]},$$

where U is depicted in dark gray. Then, A_1 is not an annulus of U because $U \cap \partial(U \cup A_1)$ is non-empty (it contains an edge); A_2 is not an annulus of U because $U \cap \partial(U \cup A_2)$ is non-empty

(it contains a point); A_3 is not an \mathcal{L}^{Br} -annulus of U because it is not strongly \mathcal{L}^{Br} -covered; indeed, if $X = \text{⬠} \subseteq A_3$ is the pattern depicted in white in the picture, there does not exist a pattern $Y \in \mathcal{L}^{\text{Br}}$ such that $X \subseteq Y \subseteq A_3$; and A_4 is an \mathcal{L}^{Br} -annulus of U .

Example 3.5. Let P be a pattern that consists of the single face $[0, 3]^*$ (shown in dark gray below) and let A be the set of faces surrounding P defined in the picture below.



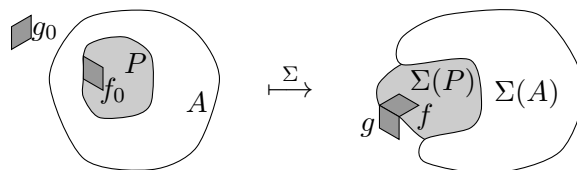
The pattern A is both \mathcal{L}^{Br} - and \mathcal{L}^{JP} -covered, but *not* strongly (so it is not a valid annulus). Its images by Brun or Jacobi-Perron substitutions fail to be topological annuli (as shown above), which illustrates the necessity of *strong* coverings if we want the image of an annulus to remain an annulus. The faulty two-face pattern and its images are shown in white.

Definition 3.6. Let Σ be a dual substitution and let \mathcal{L} a set of edge-connected patterns. We say that **Property A** holds for Σ with respect to \mathcal{L} when restricted to a family of discrete planes if for every connected two-face pattern $f \cup g$ of this family of discrete planes and for every disconnected two-face pattern $f_0 \cup g_0$ such that $f \in \Sigma(f_0)$ and $g \in \Sigma(g_0)$, the following holds: there do not exist a pattern P and an \mathcal{L} -annulus A of P (both included in a common discrete plane Γ of the given family of discrete planes) such that $f_0 \in P$ and $g_0 \in \Gamma \setminus (A \cup P)$.

The interest of the above technical definition of Property A becomes apparent in the proof of Proposition 3.7 below, thanks to which Property A (and some \mathcal{L} -covering assumptions) provide sufficient conditions for the annulus property. The main interest of Property A is that it can be checked by handling finitely many cases, namely by enumerating all the two-face connected patterns $f \cup g$ that admit a disconnected preimage. This is done for the Brun and Jacobi-Perron substitutions in Proposition 4.4 and 4.5.

Proposition 3.7. *Let Σ be a dual substitution and \mathcal{L} be a set of edge-connected patterns such that Property A holds for Σ with respect to \mathcal{L} and to a given family of discrete planes. Assume that the image by Σ of every strongly \mathcal{L} -covered pattern is strongly \mathcal{L} -covered. Let P be a pattern and A be an \mathcal{L} -annulus of P , both included in a common discrete plane of the given family. Then $\Sigma(A)$ is an \mathcal{L} -annulus of $\Sigma(P)$.*

Proof. The pattern A is strongly \mathcal{L} -covered because it is an \mathcal{L} -annulus, so $\Sigma(A)$ is also strongly \mathcal{L} -covered, by assumption. It remains to show that $\Sigma(P) \cap \partial(\Sigma(P) \cup \Sigma(A)) = \emptyset$. Suppose the contrary. This means that there exist faces f, g, f_0, g_0 such that $f \in \Sigma(f_0)$, $g \in \Sigma(g_0)$, $f \cup g$ is connected, and $f_0 \cup g_0$ is disconnected, as shown below.



These are precisely the conditions stated in Property A, so such a situation cannot occur and the proposition holds. \square

The annulus property will be crucially used in the proof of Theorem 5.1: it ensures that the minimal combinatorial radius is strictly increasing under iteration of dual substitutions.

Remark 3.8. In the definition of Property A, we consider a restriction to a given family of planes, because it is a natural assumption for the sets of substitutions that we are studying. Indeed, for Brun substitutions, vectors \mathbf{v} that are expanded satisfy $0 < \mathbf{v}_1 < \mathbf{v}_2 < \mathbf{v}_3$, and for Jacobi-Perron substitutions, vectors \mathbf{v} that are expanded satisfy $0 < \mathbf{v}_1 < \mathbf{v}_3$ and $0 < \mathbf{v}_2 < \mathbf{v}_3$. These restrictions allow us to significantly simplify the proofs of Section 4.

3.4 Generation graphs

Let $\Sigma_1, \dots, \Sigma_\ell$ be dual substitutions, let \mathcal{V} be a finite family of candidate seed patterns, and let \mathcal{W} be a finite family of patterns that we want to “reach” in the following sense: given a sequence $(i_1, \dots, i_n) \in \{1, \dots, \ell\}^n$ and given $V \in \mathcal{V}$, we want to characterize when $\Sigma_{i_1} \cdots \Sigma_{i_n}(V)$ contains a pattern $W \in \mathcal{W}$ for every sufficiently large n . (Note that W might depend on n .) Typically, \mathcal{V} will be a set of patterns containing \diamond , and \mathcal{W} will be the set of all the possible minimal \mathcal{L} -annuli of the $V \in \mathcal{V}$. The goal is then to prove that iterating substitutions starting from some $V \in \mathcal{V}$ eventually generates a pattern containing an annulus W around a pattern in \mathcal{V} , in view of “initializing” the strategy described in Section 3.1.

This is achieved thanks to the algorithmic construction of generation graphs below, where we recursively backtrack all the possible preimages of the faces in the patterns of \mathcal{W} , and we check that they eventually come back to a pattern of \mathcal{V} . This is formalized in Proposition 3.11 below, and used effectively in Sections 4.4 and 4.5.

For the above method to work in practice, we will need to filter out some useless faces when backtracking preimages of faces. We will then use a **filter set** \mathcal{F} , by restricting the allowed preimages to \mathcal{F} only. For example, the set \mathcal{F}^{Br} that will be used with the Brun substitutions is the set of all the faces that belong to a discrete plane $\Gamma_{\mathbf{v}}$ with $0 < \mathbf{v}_1 < \mathbf{v}_2 < \mathbf{v}_3$.

Definition 3.9. Let $\Sigma_1, \dots, \Sigma_\ell$ be dual substitutions, let \mathcal{F} be a family of faces (the filter set) and let \mathcal{X} be a finite set of faces (the initial set). The **generation graph** associated with $\Sigma_1, \dots, \Sigma_\ell$, \mathcal{F} and \mathcal{X} is defined as the graph $\mathcal{G} = \bigcup_{n \geq 0} \mathcal{G}_n$ where $(\mathcal{G}_n)_{n \geq 0}$ is the sequence of directed graphs (whose vertices are faces), defined by induction as follows.

- (1) *Initialization.* \mathcal{G}_0 has no edges and its set of vertices is \mathcal{X} .
- (2) *Iteration.* Suppose that \mathcal{G}_n is constructed for some $n \geq 0$. Start with \mathcal{G}_{n+1} equal to \mathcal{G}_n . Then, for each vertex f of \mathcal{G}_n , for each $i \in \{1, \dots, \ell\}$, and for each $g \in \mathcal{F}$ such that $f \in \Sigma_i(g)$, add the vertex g and the edge $g \xrightarrow{i} f$ to \mathcal{G}_{n+1} .

Remark 3.10. The non-decreasing union $\mathcal{G} = \bigcup_{n \geq 1} \mathcal{G}_n$ considered in Definition 3.9 is not necessarily finite. It is finite if and only if $\mathcal{G}_n = \mathcal{G}_{n+1}$ for some $n \geq 1$. This is the case for Brun substitutions (see Section 4.4). However, even if $\mathcal{G}_n = \mathcal{G}_{n+1}$ never occurs, the infinite graph \mathcal{G} can still be successfully exploited, as will be done for Jacobi-Perron substitutions (see Section 4.5).

The orientation of edges in generation graphs agrees with the usual notation of function composition. In particular, for every path $f_n \xrightarrow{i_n} \cdots \xrightarrow{i_2} f_1 \xrightarrow{i_1} f_0$, we have $f_0 \in \Sigma_{i_1} \cdots \Sigma_{i_n}(f_n)$.

Proposition 3.11. *Let*

- $\Sigma_1, \dots, \Sigma_\ell$ be dual substitutions;
- \mathcal{F} be a set of faces such that if $\Gamma_{\mathbf{v}} \subseteq \mathcal{F}$, then $\Sigma_i(\Gamma_{\mathbf{v}}) \subseteq \mathcal{F}$ for all $i \in \{1, \dots, \ell\}$;
- \mathcal{W} be a set of patterns such that for every $\Gamma_{\mathbf{v}} \subseteq \mathcal{F}$ we have $W \subseteq \Gamma_{\mathbf{v}}$ for some $W \in \mathcal{W}$;
- \mathcal{G} be the generation graph constructed with substitutions $\Sigma_1, \dots, \Sigma_\ell$, filter set \mathcal{F} and initial set $\mathcal{X} = \bigcup_{W \in \mathcal{W}} W$;
- \mathcal{V} be a set of patterns such that
 - for every $\Gamma_{\mathbf{v}} \subseteq \mathcal{F}$ we have $V \subseteq \Gamma_{\mathbf{v}}$ for some $V \in \mathcal{V}$,
 - for every $\Gamma_{\mathbf{v}} \subseteq \mathcal{F}$ there exists $V \in \mathcal{V}$ such that $\Gamma_{\mathbf{v}} \cap \bigcup_{V \in \mathcal{V}} V = V$,
 - for every $V \in \mathcal{V}$ there exists $\Gamma_{\mathbf{v}} \subseteq \mathcal{F}$ such that $V \subseteq \Gamma_{\mathbf{v}} \subseteq \mathcal{F}$.
- $(i_1, \dots, i_n) \in \{1, \dots, \ell\}^n$.

We have:

- (1) Suppose that for every path $f_n \xrightarrow{i_n} \cdots \xrightarrow{i_1} f_0$ in \mathcal{G} we have $f_n \in \bigcup_{V \in \mathcal{V}} V$. Then for every $V \in \mathcal{V}$ there exists $W \in \mathcal{W}$ such that $\Sigma_{i_1} \cdots \Sigma_{i_n}(V)$ contains W .
- (2) For every path $f_n \xrightarrow{i_n} \cdots \xrightarrow{i_1} f_0$ in \mathcal{G} we have $f_0 \notin \Sigma_{i_1} \cdots \Sigma_{i_n}(U)$ if $f_n \notin U$.

Proof. We first prove Statement (1). Let $V \in \mathcal{V}$ and let $\Gamma_{\mathbf{v}} \subseteq \mathcal{F}$ be a discrete plane containing V . Let $W \in \mathcal{W}$ be such that $W \subseteq \Sigma_{i_1} \cdots \Sigma_{i_n}(\Gamma_{\mathbf{v}})$ (we use the second and third assumptions of the proposition). We will prove that $\Sigma_{i_1} \cdots \Sigma_{i_n}(V)$ contains W . Let $f \in W$ and set $f_0 = f$. By Proposition 2.1, there exists f_1 in the discrete plane $\Sigma_{i_1}(\Gamma_{\mathbf{v}})$ such that $f_0 \in \Sigma_{i_1}(f_1)$. One has $f_1 \in \mathcal{F}$ (we again use the second assumption). By repeatedly applying Proposition 2.1 and by definition of generation graphs, there exists a path $f_n \xrightarrow{i_n} \cdots f_1 \xrightarrow{i_1} f_0$ in \mathcal{G} . Hence, $f_n \in \bigcup_{V \in \mathcal{V}} V$ according to the assumption on i_1, \dots, i_n . Furthermore $f_n \in \Sigma_{i_1} \cdots \Sigma_{i_n}(\Gamma_{\mathbf{v}})$. Hence, by assumption on \mathcal{V} , we have $f_n \in V$, so $f \in \Sigma_{i_1} \cdots \Sigma_{i_n}(V)$. This holds for all $f \in W$ so the statement is proved.

We now prove Statement (2). Let $f_n \xrightarrow{i_n} \cdots f_1 \xrightarrow{i_1} f_0$ be a path in \mathcal{G} such that $f_n \notin U$. The patterns U and f_n are disjoint patterns included in a common discrete plane (because U is included in every discrete plane). Hence, Proposition 2.1 implies that $\Sigma_{i_1} \cdots \Sigma_{i_n}(U)$ and $\Sigma_{i_1} \cdots \Sigma_{i_n}(f_n)$ do not have any face in common for all $n \geq 1$, which implies that $f_0 \notin \Sigma_{i_1} \cdots \Sigma_{i_n}(U)$ since $f_0 \in \Sigma_{i_1} \cdots \Sigma_{i_n}(f_n)$. \square

Remark 3.12. The technical assumptions on \mathcal{V} and \mathcal{W} in the previous proposition cannot be avoided (we rely on the properties of discrete planes and minimal annuli). Also, the assumption on \mathcal{F} always holds for the set of faces of discrete planes and for any complete multidimensional continued fraction algorithm, which includes the additive Brun and Jacobi-Perron algorithms considered here. Complete means that the continued fraction transformation is onto.

Statement (1) of Proposition 3.11 will be used to obtain “positive” results, such as the fact that a given initial pattern V always generates an entire discrete plane (see Theorem 5.1). The technical assumptions on \mathcal{F} , \mathcal{V} and \mathcal{W} are unavoidable but are natural and easy to check for the substitutions defined by multidimensional continued fraction algorithms such as the Brun and Jacobi-Perron substitutions in Section 4.1.

Conversely, Statement (2) will be used in Theorems 5.2 and 5.3 to characterize the sequences $(i_n)_{n \geq 1}$ for which iterating the corresponding sequence of dual substitutions *fails* to generate an entire discrete plane when starting from U .

4 Technical proofs

We have gathered in this section all the proofs which make specific use of Brun and Jacobi-Perron substitutions. Proofs devoted to minimal annuli and seeds are handled in Section 4.1, to covering properties in Section 4.2, to Property A in Section 4.3, to generation graphs for the Brun substitutions in Section 4.4, to generation graphs for the Jacobi-Perron substitutions in Section 4.5, and to generating translates of seeds in Section 4.6.

The following two sets of patterns will be used throughout this section:

$$\begin{aligned} \mathcal{L}^{\text{Br}} &= \{ \text{[diagram 1]}, \text{[diagram 2]}, \text{[diagram 3]}, \text{[diagram 4]}, \text{[diagram 5]}, \text{[diagram 6]}, \text{[diagram 7]} \}, \\ \mathcal{L}^{\text{JP}} &= \{ \text{[diagram 1]}, \text{[diagram 2]}, \text{[diagram 3]}, \text{[diagram 4]}, \text{[diagram 5]}, \text{[diagram 6]}, \text{[diagram 7]}, \text{[diagram 8]}, \text{[diagram 9]} \}. \end{aligned}$$

Remark 4.1. We will often use the arithmetic restrictions in the definition of discrete planes in order to simplify the combinatorics of the patterns that occur. For example, if $\mathbf{v}_1 \leq \mathbf{v}_3$ and $\mathbf{v}_2 \leq \mathbf{v}_3$, then $\Gamma_{\mathbf{v}}$ cannot contain any translate of the two-face pattern $[0, 1]^* \cup [(0, 1, 0), 1]^* = \text{[diagram]}$ or $[0, 2]^* \cup [(0, 0, 1), 2]^* = \text{[diagram]}$. If moreover $\mathbf{v}_1 \leq \mathbf{v}_2$, then the pattern $[0, 1]^* \cup [(0, 1, 0), 1]^* = \text{[diagram]}$ also never occurs.

Preimages by dual substitutions In some of the following proofs we will need to compute preimages of faces by Σ_i^{Br} or Σ_i^{JP} . By abuse of notation, given a dual substitution Σ , we will write $\Sigma^{-1}([\mathbf{x}, i]^*)$ for the union of faces $[\mathbf{y}, j]^*$ such that $\Sigma([\mathbf{y}, j]^*)$ contains the face $[\mathbf{x}, i]^*$. Computing $\Sigma^{-1}([\mathbf{x}, i]^*)$ can be done directly from the definition of dual substitutions, so such computations will be omitted in the following.

4.1 Minimal annuli and seeds

The following two propositions list all the possible \mathcal{L}^{Br} - and \mathcal{L}^{JP} -annuli in an admissible discrete plane. Proposition 4.3 can be proved in the same way as Proposition 4.2.

Proposition 4.2 (Brun minimal annuli). *Let P be a pattern contained in a discrete plane $\Gamma_{\mathbf{v}}$ with $0 < \mathbf{v}_1 < \mathbf{v}_2 < \mathbf{v}_3$. If P contains U and an \mathcal{L}^{Br} -annulus of U , then P must contain one of the following two patterns, each of which contains an \mathcal{L}^{Br} -annulus (shown in light gray) of U (shown in dark gray):*

$$V_1^{\text{Br}} = \text{[diagram]} \quad V_2^{\text{Br}} = \text{[diagram]} .$$

Proof. According to Remark 4.1, the algebraic restriction $0 < \mathbf{v}_1 < \mathbf{v}_2 < \mathbf{v}_3$ enables us to enumerate all the surroundings of the pattern U in an admissible discrete plane. It is then easy to check that such a surrounding either contains V_1^{Br} or V_2^{Br} , or contains [diagram] , [diagram] or [diagram] , which is forbidden by Remark 4.1 because $\mathbf{v}_1 < \mathbf{v}_2 < \mathbf{v}_3$. \square

Proposition 4.3 (Jacobi-Perron minimal annuli). *Let P be a pattern contained in a discrete plane $\Gamma_{\mathbf{v}}$ with $0 < \mathbf{v}_1 < \mathbf{v}_3$ and $0 < \mathbf{v}_2 < \mathbf{v}_3$. If P contains U and an \mathcal{L}^{JP} -annulus of U , then P must contain one of the following four patterns, each of which contains an \mathcal{L}^{JP} -annuli (shown in light gray) of U (shown in dark gray):*

$$V_1^{\text{JP}} = \text{[diagram]} \quad V_2^{\text{JP}} = \text{[diagram]} \quad V_3^{\text{JP}} = \text{[diagram]} \quad V_4^{\text{JP}} = \text{[diagram]} .$$

4.2 Covering properties

Proposition 4.4 (Strong \mathcal{L}^{Br} -covering). *Let P be an \mathcal{L}^{Br} -covered pattern such that the patterns [diagram] , [diagram] and [diagram] do not occur in P . Then $\Sigma_i^{\text{Br}}(P)$ is strongly \mathcal{L}^{Br} -covered for each $i \in \{1, 2, 3\}$.*

Observe that the patterns given in the statement of Proposition 4.4 cannot occur in a discrete plane with normal vector \mathbf{v} for which $0 < \mathbf{v}_1 < \mathbf{v}_2$ according to Remark 4.1, which is the natural domain of definition for Brun algorithm. The same holds for Proposition 4.5 with Jacobi-Perron admissible discrete planes.

Proof. First, $\Sigma_i^{\text{Br}}(P)$ is \mathcal{L}^{Br} -covered thanks to Proposition 3.2, because $\Sigma_i^{\text{Br}}(Q)$ is \mathcal{L}^{Br} -covered for every $Q \in \mathcal{L}^{\text{Br}}$, as can be verified from the equalities below:

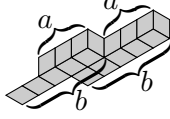
$$\begin{aligned} \Sigma_1^{\text{Br}}(\mathcal{L}^{\text{Br}}) &= \{ \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]} \} \\ \Sigma_2^{\text{Br}}(\mathcal{L}^{\text{Br}}) &= \{ \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]} \} \\ \Sigma_3^{\text{Br}}(\mathcal{L}^{\text{Br}}) &= \{ \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]}, \text{[diagram]} \}. \end{aligned}$$

It remains to prove that $\Sigma_i^{\text{Br}}(P)$ is *strongly* \mathcal{L}^{Br} -covered. Let $X \subseteq \Sigma_i^{\text{Br}}(P)$ be an edge-connected two-face pattern. If X is a translate of one of the first 6 patterns of \mathcal{L}^{Br} , then the result trivially holds (take $Y = X$ in the definition of strong coverings).

If X is a translate of $f \cup g = \text{[diagram]}$, then there must exist $f_0, g_0 \in P$ with $f_0 \neq g_0$ such that $f \in \Sigma_i^{\text{Br}}(f_0)$ and $g \in \Sigma_i^{\text{Br}}(g_0)$. Indeed, we can check that X cannot be contained in the image of a single face, by definition of the Σ_i^{Br} . Enumerating all the possible such preimages f_0, g_0 gives $f_0 \cup g_0 = \text{[diagram]}$ (for Σ_1^{Br}), $f_0 \cup g_0 = \text{[diagram]}$ (for Σ_2^{Br}), or $f_0 \cup g_0 = \text{[diagram]}$ (for Σ_3^{Br}). For $i \in \{1, 2, 3\}$ we have $\Sigma_i^{\text{Br}}(f_0 \cup g_0) = \text{[diagram]}$ so $X \subseteq \text{[diagram]} \subseteq \Sigma_i^{\text{Br}}(P)$, which enables us to take $Y = \text{[diagram]} \in \mathcal{L}^{\text{Br}}$. The cases $X = \text{[diagram]}$ or [diagram] can be settled exactly in the same way. We have treated 9 cases over 12 possible cases for X , so the proof is complete because we have ruled out the 3 remaining patterns in the statement of the proposition. \square

Proposition 4.5 (Strong \mathcal{L}^{JP} -covering). *Let P be an \mathcal{L}^{JP} -covered pattern avoiding the patterns $\begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix}$ and $\begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix}$. Then $\Sigma_{a,b}^{\text{JP}}(P)$ is strongly \mathcal{L}^{JP} -covered for all $0 \leq a \leq b$, with $b \neq 0$.*

Proof. First we must prove that $\Sigma_{a,b}^{\text{JP}}(P)$ is \mathcal{L}^{JP} -covered. Thanks to Proposition 3.2 it is enough to prove that $\Sigma_{a,b}^{\text{JP}}(Q)$ is \mathcal{L}^{JP} -covered for every $Q \in \mathcal{L}^{\text{JP}}$. Suppose that $Q = \begin{smallmatrix} \diamond \\ \diamond \end{smallmatrix}$. The pattern $\Sigma_{a,b}^{\text{JP}}(Q)$ is of the form



(in this example $a = 3$ and $b = 5$). It is then easy to check that if $a = 0$, $\Sigma_{a,b}^{\text{JP}}(Q)$ can be \mathcal{L}^{JP} -covered using the patterns $\begin{smallmatrix} \diamond \\ \diamond \end{smallmatrix}$, $\begin{smallmatrix} \diamond \\ \diamond \end{smallmatrix}$ and $\begin{smallmatrix} \diamond \\ \diamond \end{smallmatrix}$, and if $a \neq 0$, $\Sigma_{a,b}^{\text{JP}}(Q)$ can be \mathcal{L}^{JP} -covered using the patterns $\begin{smallmatrix} \diamond \\ \diamond \end{smallmatrix}$, $\begin{smallmatrix} \diamond \\ \diamond \end{smallmatrix}$, $\begin{smallmatrix} \diamond \\ \diamond \end{smallmatrix}$, $\begin{smallmatrix} \diamond \\ \diamond \end{smallmatrix}$ and $\begin{smallmatrix} \diamond \\ \diamond \end{smallmatrix}$. A similar reasoning can be carried out for each of the other 9 patterns $Q \in \mathcal{L}^{\text{JP}}$, which proves that P is \mathcal{L}^{JP} -covered.

It remains to prove the *strong* \mathcal{L}^{JP} -covering. Let $X \subseteq \Sigma_{a,b}^{\text{JP}}(P)$ be an edge-connected two-face pattern. If X is a translate of one of the first 7 patterns of \mathcal{L}^{JP} , then the result is obvious, because X is itself in \mathcal{L}^{JP} . If X is a translate of $\begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix}$, then we must have $X = \begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix} \subseteq \begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix} \subseteq \Sigma_{a,b}^{\text{JP}}(P)$, because in an image by $\Sigma_{a,b}^{\text{JP}}$, a face of type 1 must come from a face of type 3, so there must also be a face of type 3 at the same position (because $b \neq 0$). If X is a translate of $\begin{smallmatrix} \blacktriangleright \\ \blacktriangleright \end{smallmatrix}$, a similar reasoning yields $X = \begin{smallmatrix} \blacktriangleright \\ \blacktriangleright \end{smallmatrix} \subseteq \begin{smallmatrix} \blacktriangleright \\ \blacktriangleright \end{smallmatrix} \subseteq \Sigma_{a,b}^{\text{JP}}(P)$. If X is a translate of $\begin{smallmatrix} \blacktriangleleft \\ \blacktriangleleft \end{smallmatrix}$, then X is in the image of a face of type 3, or has one face in the image of a face of type 2, and the other face in the image of a face of type 3. In both cases, we must have $X = \begin{smallmatrix} \blacktriangleleft \\ \blacktriangleleft \end{smallmatrix} \subseteq \begin{smallmatrix} \blacktriangleleft \\ \blacktriangleleft \end{smallmatrix} \subseteq \Sigma_{a,b}^{\text{JP}}(P)$ (since $b \geq a$).

We have handled 10 cases over 12 possible cases for X , and the two remaining patterns have been ruled out in the statement of the proposition. \square

4.3 Property A

Proposition 4.6 (Property A for Brun). *Property A holds for Brun substitutions with \mathcal{L}^{Br} , when restricted to planes $\Gamma_{\mathbf{v}}$ with $0 < \mathbf{v}_1 < \mathbf{v}_2 < \mathbf{v}_3$.*

Proof. There are finitely many two-face connected patterns $f \cup g$, so we can enumerate all the faces f, g, f_0, g_0 that satisfy $f \in \Sigma(f_0)$, $g \in \Sigma(g_0)$, $f \cup g$ is connected and $f_0 \cup g_0$ is disconnected (see Definition 3.6), for $\Sigma = \Sigma_1^{\text{Br}}, \Sigma_2^{\text{Br}}$ and Σ_3^{Br} . It turns out that there are 9 such possibilities, where the corresponding values for $f_0 \cup g_0$ are shown in the table below.

Σ_1^{Br}	Σ_2^{Br}	Σ_3^{Br}
$[\mathbf{0}, 2]^* \cup [(0, 1, 0), 1]^*$	$[\mathbf{0}, 3]^* \cup [(1, 0, -1), 3]^*$	$[\mathbf{0}, 3]^* \cup [(0, 1, -1), 3]^*$
$[\mathbf{0}, 2]^* \cup [(1, -1, 0), 2]^*$	$[\mathbf{0}, 3]^* \cup [(0, 1, 1), 1]^*$	$[\mathbf{0}, 3]^* \cup [(0, 0, 1), 2]^*$
$[\mathbf{0}, 2]^* \cup [(0, 1, 1), 1]^*$	$[\mathbf{0}, 3]^* \cup [(0, 0, 1), 1]^*$	$[\mathbf{0}, 3]^* \cup [(1, 0, 1), 2]^*$

Let us treat the case $f_0 \cup g_0 = [\mathbf{0}, 2]^* \cup [(1, -1, 0), 2]^*$. Suppose that there exist a pattern P and an \mathcal{L}^{Br} -annulus A of P that is included in a discrete plane such that $f_0 \in P$ and $g_0 \notin A$. Because A is an annulus of P , every extension of $f_0 \cup g_0$ within a discrete plane must be of the form $\begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix}$ or $\begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix}$, where $f_0 \cup g_0$ is shown in light gray and the dark gray faces are in A .

Now, an occurrence of $\begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix}$ is forbidden, since we are restricted to discrete planes $\Gamma_{\mathbf{v}}$ with $0 < \mathbf{v}_1 < \mathbf{v}_2 < \mathbf{v}_3$ (see Remark 4.1). An occurrence of $\begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix}$ also cannot happen, because A is strongly \mathcal{L}^{Br} -covered. Indeed, $\begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix} \subseteq A$, so there must exist a translate of a pattern of \mathcal{L}^{Br} that is included in A and that contains $\begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix}$. The only such pattern in \mathcal{L}^{Br} is $\begin{smallmatrix} \blacktriangleright \\ \blacktriangleright \end{smallmatrix}$ (note that $\begin{smallmatrix} \blacksquare \\ \blacksquare \end{smallmatrix} \notin \mathcal{L}^{\text{Br}}$). This is impossible: otherwise $\begin{smallmatrix} \blacktriangleright \\ \blacktriangleright \end{smallmatrix}$ and $f_0 \cup g_0$ overlap, which is a contradiction because $f_0, g_0 \notin A$ and $\begin{smallmatrix} \blacktriangleright \\ \blacktriangleright \end{smallmatrix} \subseteq A$. The same reasoning applies to the eight other cases. \square

In order to prove Property A for Jacobi-Perron substitutions (Proposition 4.8 below), we first need to prove Lemma 4.7 below, which describes all the possible disconnected preimages by

$\Sigma_{a,b}^{\text{JP}}$ of a two-face connected pattern. A striking fact is that there are only three possible such preimages, despite the fact that a and b can take infinitely many values.

Lemma 4.7. *Let f and g be two faces such that*

- (1) $f \cup g$ is included in a discrete plane $\Gamma_{\mathbf{v}}$ with $0 < \mathbf{v}_1 < \mathbf{v}_3$ and $0 < \mathbf{v}_2 < \mathbf{v}_3$;
- (2) $f \cup g$ is disconnected;
- (3) $\Sigma_{a,b}^{\text{JP}}(f \cup g)$ is connected for some $0 \leq a \leq b, b \neq 0$.

Then $f \cup g$ is a translate of one of the three patterns $P_1 = [\mathbf{0}, 3]^* \cup [(1, 0, -1), 3]^*$, $P_2 = [\mathbf{0}, 3]^* \cup [(1, -1, -1), 3]^*$, or $P_3 = [\mathbf{0}, 3]^* \cup [(-1, 1, -1), 3]^*$.

Proof. We first need the following easy preliminary fact (Equation (2) below), which can be checked by inspection of finitely many cases. A two-face pattern $P = [\mathbf{x}, i]^* \cup [\mathbf{y}, j]^*$ is edge-connected if and only $\mathbf{y} - \mathbf{x} \in V_{ij}$, where

$$\begin{aligned}
V_{11} &= \{\pm(0, 1, 0), \pm(0, 0, 1), \pm(0, 1, -1), \pm(0, 1, 1)\}, \\
V_{22} &= \{\pm(1, 0, 0), \pm(0, 0, 1), \pm(1, 0, -1), \pm(1, 0, 1)\}, \\
V_{33} &= \{\pm(1, 0, 0), \pm(0, 1, 0), \pm(1, -1, 0), \pm(1, 1, 0)\}, \\
V_{12} &= \{(0, 0, 0), (-1, 1, 0), \pm(0, 0, 1), (0, 1, -1), (-1, 1, -1), (-1, 1, 1), (-1, 0, 1)\}, \\
V_{13} &= \{(0, 0, 0), (-1, 0, 1), \pm(0, 1, 0), (-1, 1, 0), (-1, 1, 1), (-1, -1, 1), (0, -1, 1)\}, \\
V_{23} &= \{(0, 0, 0), (0, -1, 1), \pm(1, 0, 0), (-1, 0, 1), (-1, -1, 1), (1, -1, 1), (1, -1, 0)\}, \\
V_{21} &= -V_{12}, \quad V_{31} = -V_{13}, \quad V_{32} = -V_{22}.
\end{aligned} \tag{2}$$

In contrast with the Brun substitutions, we cannot prove Lemma 4.7 by enumerating all the possibilities because a and b can take infinitely many values. However, the problem can be reduced to solving a simple family of linear equations, which then allows a systematic study of the disconnected preimages of $f \cup g$.

We will show that if Condition (3) holds, then Condition (1) or (2) fails, unless $f \cup g$ is a translate of P_1 , P_2 or P_3 . We will only treat the case where the types of f and g are both equal to 3. This corresponds to the most complicated case; the remaining cases can be handled analogously. Recall that $\mathbf{M}_{a,b} = {}^t\mathbf{M}_{\sigma_{a,b}^{\text{JP}}}$. Let $f = [\mathbf{x}, 3]^*$ and $g = [\mathbf{y}, 3]^*$ be two faces contained in a same discrete plane, and assume that Condition (3) holds. According to Equation (2), the fact that $\Sigma_{a,b}^{\text{JP}}(f \cup g)$ is connected implies one of the following six relations:

$$\begin{aligned}
{}^t\mathbf{M}_{a,b}^{-1}(\mathbf{y} - \mathbf{x}) &\in V_{11}, \\
{}^t\mathbf{M}_{a,b}^{-1}(\mathbf{y} - \mathbf{x}) + (k, 0, 0) &\in V_{12} \quad (0 \leq k \leq a - 1), \\
{}^t\mathbf{M}_{a,b}^{-1}(\mathbf{y} - \mathbf{x}) + (\ell, 0, 0) &\in V_{13} \quad (0 \leq \ell \leq b - 1), \\
{}^t\mathbf{M}_{a,b}^{-1}(\mathbf{y} - \mathbf{x}) + (k - k', 0, 0) &\in V_{22} \quad (0 \leq k, k' \leq a - 1), \\
{}^t\mathbf{M}_{a,b}^{-1}(\mathbf{y} - \mathbf{x}) + (\ell - k, 0, 0) &\in V_{23} \quad (0 \leq k \leq a - 1, 0 \leq \ell \leq b - 1), \\
{}^t\mathbf{M}_{a,b}^{-1}(\mathbf{y} - \mathbf{x}) + (\ell - \ell', 0, 0) &\in V_{33} \quad (0 \leq \ell, \ell' \leq b - 1),
\end{aligned}$$

for some $0 \leq a \leq b$ with $b \neq 0$. This is equivalent to

$$\mathbf{y} - \mathbf{x} \in \left\{ \begin{pmatrix} 0 \\ 0 \\ s \end{pmatrix}, \pm \begin{pmatrix} 1 \\ 0 \\ s \end{pmatrix}, \begin{pmatrix} -1 \\ 1 \\ s \end{pmatrix}, \pm \begin{pmatrix} 0 \\ 1 \\ t \end{pmatrix}, \pm \begin{pmatrix} 1 \\ 1 \\ t \end{pmatrix}, \begin{pmatrix} 1 \\ -1 \\ -t \end{pmatrix} : s \in \mathbb{Z}, t \geq 0 \right\}.$$

This either contradicts Condition (1) or (2), or else implies that $f \cup g$ is equal to P_1 , P_2 or P_3 . Indeed, let $\Gamma_{\mathbf{v}}$ be a discrete plane that contains both f and g . Thanks to Remark 4.1, we can restrict the possible values of $\mathbf{x} - \mathbf{y}$ even further. The remaining valid solutions are shown in the following table, where $s \in \mathbb{Z}$, $t \geq 0$, $0 < \mathbf{v}_1 < \mathbf{v}_3$ and $0 < \mathbf{v}_2 < \mathbf{v}_3$. We then observe that either


$\mathbf{y} - \mathbf{x} \in V_{33}$ (so Condition (2) is violated because $f \cup g$ is connected thanks to Equation (2)), or that $f \cup g$ is equal to P_1 , P_2 or P_3 .

Inequalities	Solutions
$-\mathbf{v}_3 < s\mathbf{v}_3 < \mathbf{v}_3$	$s = 0: \mathbf{y} - \mathbf{x} \in V_{33}$
$-\mathbf{v}_3 < \mathbf{v}_1 + s\mathbf{v}_3 < \mathbf{v}_3$	$s = 0: \mathbf{y} - \mathbf{x} \in V_{33}; s = 1: P_1$
$-\mathbf{v}_3 < -\mathbf{v}_1 + \mathbf{v}_2 + s\mathbf{v}_3 < \mathbf{v}_3$	$s = 0: \mathbf{y} - \mathbf{x} \in V_{33}; s = -1: P_3; s = 1: P_2$
$-\mathbf{v}_3 < \mathbf{v}_2 + t\mathbf{v}_3 < \mathbf{v}_3$	$t = 0: \mathbf{y} - \mathbf{x} \in V_{33}$
$-\mathbf{v}_3 < \mathbf{v}_1 + \mathbf{v}_2 + t\mathbf{v}_3 < \mathbf{v}_3$	$t = 0: \mathbf{y} - \mathbf{x} \in V_{33}$
$-\mathbf{v}_3 < \mathbf{v}_1 - \mathbf{v}_2 - t\mathbf{v}_3 < \mathbf{v}_3$	$t = 0: \mathbf{y} - \mathbf{x} \in V_{33}; t = 1: P_3$

□

Proposition 4.8 (Property A for Jacobi-Perron). *Property A holds for Jacobi-Perron substitutions with \mathcal{L}^{JP} , when restricted to discrete planes $\Gamma_{\mathbf{v}}$ with $0 < \mathbf{v}_1 < \mathbf{v}_3$ and $0 < \mathbf{v}_2 < \mathbf{v}_3$.*

Proof. Thanks to Lemma 4.7, if f, g, f_0, g_0 satisfy the conditions required in Definition 3.6 for some Jacobi-Perron substitution, then $f_0 \cup g_0$ must be equal to a translate of one of the three patterns P_1, P_2, P_3 of the statement of Lemma 4.7.

Suppose that there exists a pattern P and an \mathcal{L}^{JP} -annulus A of P that is included in a discrete plane such that $f_0 \in P$ and $g_0 \notin A \cup P$. Because A is an annulus of P , every extension of $f_0 \cup g_0$ within a discrete plane must be of the form , where $f_0 \cup g_0$ is shown in light gray and the dark gray faces are in A . Now, similarly as in the proof of Proposition 4.6, a contradiction must occur in each case, thanks to the strong \mathcal{L}^{JP} -covering of A and the precise choice of patterns in \mathcal{L}^{JP} . □

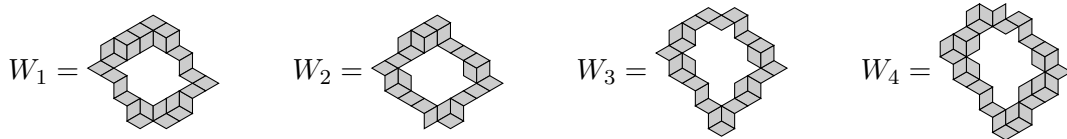
4.4 Generation graphs for the Brun substitutions

We now construct generation graphs for the substitutions $\Sigma_1^{\text{Br}}, \Sigma_2^{\text{Br}}, \Sigma_3^{\text{Br}}$. To construct the graphs below we use the filter \mathcal{F}^{Br} equal to the set of all the faces f that belong to a discrete plane $\Gamma_{\mathbf{v}}$ with $0 < \mathbf{v}_1 < \mathbf{v}_2 < \mathbf{v}_3$.

Graph \mathcal{G}^{Br} for the seed U Let \mathcal{G}^{Br} be the generation graph (see Definition 3.9) associated with $\Sigma_1^{\text{Br}}, \Sigma_2^{\text{Br}}, \Sigma_3^{\text{Br}}$, the filter \mathcal{F}^{Br} and the initial set of faces $\mathcal{X} = V_1^{\text{Br}} \cup V_2^{\text{Br}}$ (the union of the minimal annuli given in Proposition 4.2). Its computation stops after two iterations of the algorithm. It has 19 vertices and 47 edges.

Now we also define $\tilde{\mathcal{G}}^{\text{Br}}$ as the subgraph of \mathcal{G}^{Br} containing only the vertices which are not contained in U and which are at the end of an infinite backward path containing infinitely many 3s. The graph $\tilde{\mathcal{G}}^{\text{Br}}$ is plotted in Figure 2. It will be used in the proofs of Lemma 4.12 and Theorem 5.2 to characterize the sequences $(i_n)_{n \geq 1}$ for which the patterns $\Sigma_{i_1}^{\text{Br}} \dots \Sigma_{i_n}^{\text{Br}}(U)$ generate an entire discrete plane.

Graph \mathcal{H}^{Br} for the seeds V_i^{Br} We now construct another generation graph for the Brun substitutions, with the seeds V_1^{Br} and V_2^{Br} instead of U . This time, every Brun-admissible sequence will be proved to generate an \mathcal{L}^{Br} -annulus around V_1^{Br} or V_2^{Br} . Before computing \mathcal{H}^{Br} we first enumerate all the possible minimal \mathcal{L}^{Br} -annuli of V_1^{Br} or V_2^{Br} (minimal in the sense that any \mathcal{L}^{Br} -annulus of V_i^{Br} must contain one of these patterns). This can be done similarly as in the proof of Proposition 4.2, and yields the following four possible minimal annuli.



Let \mathcal{H}^{Br} be the generation graph (according to Definition 3.9) constructed with substitutions $\Sigma_1^{\text{Br}}, \Sigma_2^{\text{Br}}, \Sigma_3^{\text{Br}}$, filter \mathcal{F}^{Br} and initial set $\mathcal{X} = W_1 \cup W_2 \cup W_3 \cup W_4$ (a total of 60 faces). Its

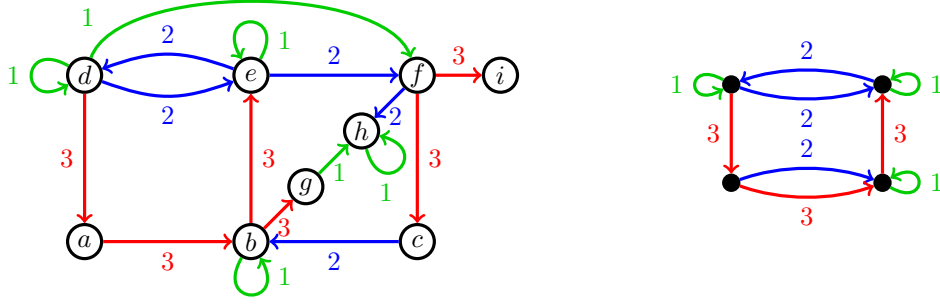


Figure 2: The graph $\tilde{\mathcal{G}}^{\text{Br}}$ (left). The faces corresponding to the vertices of the graph are $a = [(1, 1, -1), 1]^*$, $b = [(1, -1, 1), 3]^*$, $c = [(1, 1, -1), 2]^*$, $d = [(-1, 1, 0), 2]^*$, $e = [(-1, 0, 1), 3]^*$, $f = [(-1, 1, 0), 3]^*$, $g = [(-1, 0, 1), 2]^*$, $h = [(-1, -1, 1), 3]^*$, $i = [(1, 1, -1), 3]^*$. The graph on the right is obtained from $\tilde{\mathcal{G}}^{\text{Br}}$ by merging some vertices. It is a reduced version of $\tilde{\mathcal{G}}^{\text{Br}}$ in the sense that the language of the edge labellings of its paths is equal to that of $\tilde{\mathcal{G}}^{\text{Br}}$.

computation stops after six iterations of the algorithm. It has 101 vertices and 240 edges. We do not plot the graph below, but it can be easily recovered from the Sage source code provided with this article. The aim of computing \mathcal{H}^{Br} is to prove the following lemma, which will be used in the proof of Theorem 5.1.

Lemma 4.9. *Let $V \in \{V_1^{\text{Br}}, V_2^{\text{Br}}\}$. There exists N such that for every $(i_1, \dots, i_n) \in \{1, 2, 3\}^n$ containing more than N occurrences of 3, the pattern $\Sigma_{i_1}^{\text{Br}} \dots \Sigma_{i_n}^{\text{Br}}(V)$ contains an \mathcal{L}^{Br} -annulus of V_1^{Br} or V_2^{Br} .*

Proof. Let $N \geq 1$ be such that for every path $f_n \xrightarrow{i_n} f_{n-1} \xrightarrow{i_{n-1}} \dots \xrightarrow{i_1} f_0$ in \mathcal{H}^{Br} with $i_n = 3$ more than N times, we have $f_n \in V_1^{\text{Br}} \cup V_2^{\text{Br}}$. The existence of such an N can be proved computationally in the following way: one checks that every strongly connected component containing an edge labelled by 3 consists only of vertices in V_1^{Br} or V_2^{Br} . Any path containing sufficiently many 3s must originate from such a component, and hence from a face with vertex in V_1^{Br} or V_2^{Br} .

The proposition now follows directly from Statement (1) of Proposition 3.11, applied to a sequence (i_1, \dots, i_n) containing at least N occurrences of 3. The hypotheses on \mathcal{F}^{Br} , $\mathcal{V} = \{V_1^{\text{Br}}, V_2^{\text{Br}}\}$, $\mathcal{W} = \{W_1, W_2, W_3, W_4\}$ needed in Proposition 3.11 can be checked easily. In particular, the fact that $\Gamma_{\mathbf{v}} \subseteq \mathcal{F}^{\text{Br}}$ implies $\Sigma_i^{\text{Br}}(\Gamma_{\mathbf{v}}) \subseteq \mathcal{F}^{\text{Br}}$ for all $i \in \{1, 2, 3\}$ follows directly from the definition of the Brun algorithm and from Proposition 2.1. \square

4.5 Generation graphs for the Jacobi-Perron substitutions

In this section we take the filter \mathcal{F}^{JP} to be the set of faces f that belong to a discrete plane $\Gamma_{\mathbf{v}}$ such that $0 < \mathbf{v}_1 < \mathbf{v}_3$ and $0 < \mathbf{v}_2 < \mathbf{v}_3$.

The graph \mathcal{G}^{JP} for the additive Jacobi-Perron substitutions Let \mathcal{G}^{JP} be the generation graph (see Definition 3.9) associated with the additive Jacobi-Perron substitutions $\Theta_1, \Theta_2, \Theta_3, \Theta_4$, and the starting set $\mathcal{X} = V_1^{\text{JP}} \cup V_2^{\text{JP}} \cup V_3^{\text{JP}} \cup V_4^{\text{JP}}$ (defined in Proposition 4.3). Unlike the graph \mathcal{G}^{Br} , the graph \mathcal{G}^{JP} is infinite (we never have $\mathcal{G}_n^{\text{JP}} = \mathcal{G}_{n+1}^{\text{JP}}$). However, the structure of \mathcal{G}^{JP} is simple enough, so we will be able to describe it precisely below.

Let $\mathcal{G}_3^{\text{JP}}$ be the graph obtained at the third step in the computation of \mathcal{G}^{JP} , defined in Definition 3.9. This graph has 33 vertices and 93 edges. In order to use $\mathcal{G}_3^{\text{JP}}$ to study \mathcal{G}^{JP} we need the following lemma, which can be proved easily by studying the preimages of the additive Jacobi-Perron substitutions $\Theta_1, \Theta_2, \Theta_3$ and Θ_4 .

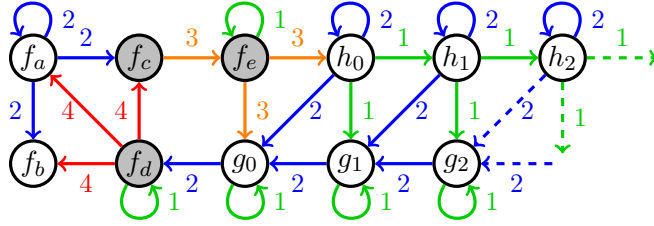
Lemma 4.10. *Define $e_n = [(-n, 1, 0), 3]^*$, $f_n = [(-n, 1, 0), 1]^*$, $g_n = [(-n, -1, 1), 3]^*$, $h_n = [(-n, -1, 1), 1]^*$ for $n \geq 0$. For every integer $n \geq 1$, the only preimages of e_n, f_n, g_n, h_n by one*

of the additive Jacobi-Perron substitutions $\Theta_1, \Theta_2, \Theta_3$ or Θ_4 are given by:

$$\begin{array}{llll} e_n \in \Theta_1(e_{n+1}) & f_n \in \Theta_1(f_n) & g_n \in \Theta_1(g_n) & h_n \in \Theta_1(h_{n-1}) \\ e_n \in \Theta_1(f_n) & f_n \in \Theta_2(h_{n-1}) & g_n \in \Theta_1(h_{n-1}) & h_n \in \Theta_2(h_n). \\ e_n \in \Theta_2(e_n) & & g_n \in \Theta_2(h_n) & \\ e_n \in \Theta_2(f_{n-1}) & & g_n \in \Theta_2(h_{n+1}) & \end{array}$$

We deduce from Lemma 4.10 that \mathcal{G}^{JP} has two infinite branches, the first one consisting of vertices e_n, f_n , and the second one consisting of vertices g_n, h_n , with the only edge labels occurring in these infinite branches being 1 and 2. Indeed, the only faces of $\mathcal{G}_3^{\text{JP}}$ with preimages allowed by the filter set \mathcal{F}^{JP} which are not vertices of the graph $\mathcal{G}_3^{\text{JP}}$ are e_4, f_3, g_4 and h_3 .

Since we are interested only in the additive Jacobi-Perron-admissible sequences $(i_n)_{n \geq 1}$ (that is, containing infinitely many edges labelled by 3 or 4 by Remark 2.7), we remove from \mathcal{G}^{JP} all the vertices which are not at the beginning of a backward infinite path labelled by such a sequence $(i_n)_{n \geq 1}$. This yields the following subgraph of \mathcal{G}^{JP} (which now contains only one infinite branch):



where the vertices are given by $f_a = [(1, 1, -1), 1]^*$, $f_b = [(1, 1, -1), 3]^*$, $f_c = [(1, 1, -1), 2]^*$, $f_d = [(1, -1, 1), 3]^*$, $f_e = [(-1, 1, 1), 3]^*$, and the edges are the ones shown above, plus, for every $n \geq 0$, the edges $g_n \xrightarrow{1} g_n$, $h_n \xrightarrow{2} h_n$, $h_n \xrightarrow{2} g_n$, $h_n \xrightarrow{1} g_{n+1}$, $h_n \xrightarrow{1} h_{n+1}$ and $g_{n+1} \xrightarrow{2} g_n$. Moreover we have deleted the edge $f_e \xrightarrow{4} f_d$ in the above graph, because every infinite path containing it is incompatible with the condition $a_n = b_n \Rightarrow a_{n+1} \neq 0$ on Jacobi-Perron expansions $(a_n, b_n)_{n \geq 1}$ (see Section 2.3). Indeed, if the edge $f_e \xrightarrow{4} f_d$ were allowed, then the forbidden product $\Theta_4 \Theta_1^k \Theta_3 = \Sigma_{1,1}^{\text{JP}} \Sigma_{0,k+1}^{\text{JP}}$ would be allowed for some $k \geq 0$, a contradiction.

The graph \mathcal{H}^{JP} We now construct a generation graph \mathcal{H}^{JP} for the seeds $V_1^{\text{JP}}, V_2^{\text{JP}}, V_3^{\text{JP}}$ and V_4^{JP} . Before computing \mathcal{H}^{JP} we must enumerate all the possible minimal \mathcal{L}^{JP} -annuli of the V_i^{JP} . This can be done similarly as for the proof of Proposition 4.2, and yields eight possible minimal annuli W_1, \dots, W_8 , which are similar to the annuli W_1, \dots, W_4 computed for \mathcal{H}^{Br} .

Let \mathcal{H}^{JP} be the generation graph (according to Definition 3.9) associated with the additive Jacobi-Perron substitutions $\Theta_1, \Theta_2, \Theta_3, \Theta_4$, the initial set $\mathcal{X} = W_1 \cup \dots \cup W_8$ and the filter \mathcal{F}^{JP} . Thanks to the structure of \mathcal{H}^{JP} we have the following lemma, which can be proved in exactly the same way as the corresponding Lemma 4.9 for the Brun substitutions.

Lemma 4.11. *Let $V \in \{V_1^{\text{JP}}, V_2^{\text{JP}}, V_3^{\text{JP}}, V_4^{\text{JP}}\}$. There exists N such that for every $(i_1, \dots, i_n) \in \{1, 2, 3, 4\}^n$ containing more than N occurrences of 3 or 4, the pattern $\Theta_{i_1} \dots \Theta_{i_n}(V)$ contains an \mathcal{L}^{JP} -annulus of $V_1^{\text{JP}}, V_2^{\text{JP}}, V_3^{\text{JP}}$ or V_4^{JP} .*

4.6 Generating translates of seeds

Lemma 4.12. *There exists N such that if a finite sequence $(i_1, \dots, i_n) \in \{1, 2, 3\}^n$ contains more than N occurrences of 3, then $\Sigma_{i_1}^{\text{Br}} \dots \Sigma_{i_n}^{\text{Br}}(U)$ contains a translate of V_1^{Br} or V_2^{Br} .*

Proof. We can assume that the sequences (i_1, \dots, i_n) we consider are the edge labellings of a backward path $\bullet \xrightarrow{i_n} \dots \bullet \xrightarrow{i_1} \bullet$ in the four-vertex graph given in Section 4.4, Figure 2. Indeed, the language of the labellings of the four-vertex graph is the same as the one of the graph $\tilde{\mathcal{G}}^{\text{Br}}$ (given

in the same figure), and every path containing sufficiently many 3s which is not contained in $\tilde{\mathcal{G}}^{\text{Br}}$ must originate from U (by construction of $\tilde{\mathcal{G}}^{\text{Br}}$).

Using this restriction, we can assume that (i_n, \dots, i_1) necessarily starts with a word in one of the following languages, for n large enough: $21^*321^*31^*2$, $21^*331^*31^*21^*2$, $21^*331^*31^*21^*33$, $21^*331^*31^*21^*321^*31^*2(1 \cup 2 \cup 3)$. This enumeration can be carried out easily by enumerating the possible path labellings in the four-vertex graph (we must turn counter-clockwise when sufficiently many 3s are read). Also note that each of these languages starts with 21^*32 or 21^*33 , which corresponds to starting from the top-left vertex in the four-vertex graph. This is not a restriction: indeed, if (i_n, \dots, i_1) does not start in this way, then we can consider longer sequences (i_{n+k}, \dots, i_1) which agree with the above languages.

It now remains to prove that if (i_n, \dots, i_1) is in one of the following languages then $\Sigma_{i_1}^{\text{Br}} \dots \Sigma_{i_n}^{\text{Br}}(U)$ contains a translate of V_1^{Br} or V_2^{Br} . To do so we track all the possible sequences of iterations starting from U , and we explicitly check that a translate of one of the seeds V_1 or V_2 is eventually generated. This is done thanks to the graph shown in Figure 3. The vertices are (translation classes of) patterns, and for every edge $P \xrightarrow{i} Q$ in the graph, $\Sigma_i^{\text{Br}}(P)$ contains a translated copy of Q . The pattern $P_0 = [(0, -1, 1), 3]^*$, $[(0, 0, 0), 2]^*$, $[(1, 0, 0), 2]^*$, $[(1, 0, 0), 3]^*$ (the leftmost vertex of the graph) can be easily proved to belong to the image of U after iterating any sequence of substitutions containing Σ_3^{Br} at least 3 times.

Note that unlike the generation graphs constructed in Sections 4.4 and 4.5, the graph of Figure 3 used in the present proof has not been constructed algorithmically. Constructing such a graph is tedious, but checking its validity is easy (by using computer algebra). \square

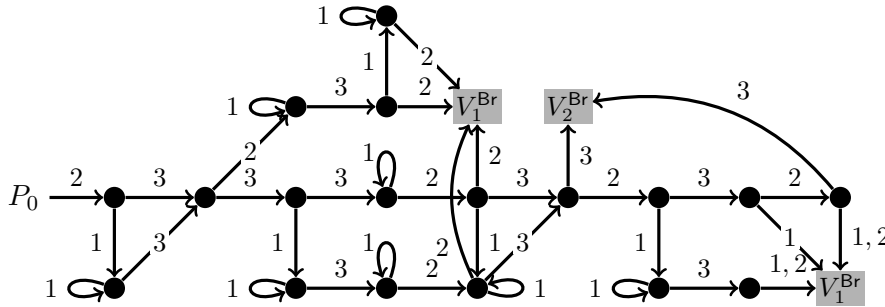
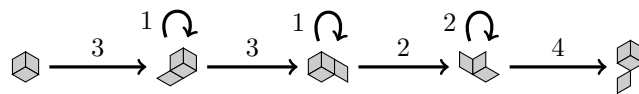


Figure 3: The graph used in the proof of Lemma 4.12. The vertices are patterns (up to translation). There is an edge from P to Q labelled by i if $\Sigma_i^{\text{Br}}(P)$ contains a *translated copy* of Q .

We now prove an analogue of the previous lemma for Jacobi-Perron substitutions. The proof is easier (we need a smaller graph); indeed the set of “bad” Jacobi-Perron expansions is more constrained (see Theorem 5.3).

Lemma 4.13. *Let $(a_n, b_n)_{n \geq 1}$ be an admissible Jacobi-Perron expansion. Then there exists $N \geq 0$ such that $\Sigma_{a_1, b_1}^{\text{JP}} \dots \Sigma_{a_N, b_N}^{\text{JP}}(U)$ contains a translate of V_1^{JP} , V_2^{JP} , V_3^{JP} or V_4^{JP} .*

Proof. Similarly as in the proof of Lemma 4.12, we can apply Theorem 5.3, so that the result only needs to be proved for products of the form $\Theta_{i_1} \dots \Theta_{i_n}$ with $i_1 \dots i_n \in 31^*31^*22^*4$, which is a translation of the condition given in Theorem 5.3 in terms of additive Jacobi-Perron substitutions. This is settled thanks to the following graph, in which every edge $P \xrightarrow{i} Q$ means that $\Theta_i(P)$ contains a translated copy of Q .



Finally, it can be checked that the last pattern $\begin{array}{c} \blacksquare \\ \blacksquare \\ \blacksquare \end{array}$ is sufficient to generate a seed V_i^{JP} , thanks to Theorem 5.3 below (the proof of which does not make use of Lemma 4.13). \square

5 Main results

We first state discrete plane generation results in Section 5.1. We then state corollaries of topological, dynamical and number-theoretical nature in Sections 5.2, 5.3 and 5.4. We recall that a product $\sigma_{i_1}^{\text{Br}} \cdots \sigma_{i_n}^{\text{Br}}$ is **Brun-admissible** if $i_k = 3$ for some $n \in \{1, \dots, n\}$, and that a product $\sigma_{a_1, b_1}^{\text{JP}} \cdots \sigma_{a_n, b_n}^{\text{JP}}$ is **Jacobi-Perron-admissible** if for every $1 \leq k \leq n$, we have (1) $0 \leq a_k \leq b_k$, (2) $b_k \neq 0$, and (3) $a_k = b_k$ implies $a_{k+1} \neq 0$. We also recall that the **combinatorial radius** $\text{rad}(P)$ of a pattern P containing the origin measures the minimal distance between $\mathbf{0}$ and the boundary of P (see Section 2.1 for a precise definition).

5.1 Discrete plane generation

Existence of finite seeds We start by stating the existence of finite seeds such that iterating *any* (Brun or Jacobi-Perron) admissible sequence of substitutions on them generates patterns with arbitrarily large minimal combinatorial radius. These finite seeds $V_1^{\text{Br}}, V_2^{\text{Br}}$ and $V_1^{\text{JP}}, V_2^{\text{JP}}, V_3^{\text{JP}}, V_4^{\text{JP}}$ are given explicitly in Propositions 4.2 and 4.3, and are proved to be minimal in the sense that every admissible discrete plane must contain one of them.

Theorem 5.1. *For every $R \geq 0$ there exists $N \geq 0$ such that*

- (1) $\text{rad}(\Sigma_{i_1}^{\text{Br}} \cdots \Sigma_{i_n}^{\text{Br}}(V)) \geq R$ for every $(i_1 \cdots i_n) \in \{1, 2, 3\}^n$ that contains more than N occurrences of 3, and for every $V \in \{V_1^{\text{Br}}, V_2^{\text{Br}}\}$;
- (2) $\text{rad}(\Sigma_{a_1, b_1}^{\text{JP}} \cdots \Sigma_{a_n, b_n}^{\text{JP}}(V)) \geq R$ for every admissible Jacobi-Perron expansion $(a_1, b_1), \dots, (a_n, b_n)$ such that $n \geq N$, and for every $V \in \{V_1^{\text{JP}}, V_2^{\text{JP}}, V_3^{\text{JP}}, V_4^{\text{JP}}\}$.

Proof of Theorem 5.1. We fix $V \in \{V_1^{\text{Br}}, V_2^{\text{Br}}\}$. The idea of the proof is that given n concentric nested annuli around the initial pattern V , applying dual substitutions preserves these n concentric nested annuli (thanks to the annulus property), and moreover eventually generates a new annulus around V (thanks to generation graphs). This gives $n + 1$ concentric nested annuli, and repeating this reasoning yields the result.

We now formalize the above reasoning to prove Statement (1) of Theorem 5.1 about Brun substitutions. By virtue of Lemma 4.9, there exists $M \geq 1$ such that for every $(i_1, \dots, i_n) \in \{1, 2, 3\}^n$ containing more than M occurrences of 3, the pattern $\Sigma_{i_1}^{\text{Br}} \cdots \Sigma_{i_n}^{\text{Br}}(V)$ contains an \mathcal{L}^{Br} -annulus of V_1^{Br} or V_2^{Br} .

Let $R \geq 0$, let $N = R \times M$, and let $(i_1, \dots, i_n) \in \{1, 2, 3\}^n$ containing more than N occurrences of 3. Write $\Sigma_{i_1}^{\text{Br}} \cdots \Sigma_{i_n}^{\text{Br}} = \Sigma_1 \cdots \Sigma_R$, where each Σ_k is a product of Σ_i^{Br} containing Σ_3^{Br} at least M times. To prove that $\text{rad}(\Sigma_1 \cdots \Sigma_R(V)) \geq R$, it is enough to prove that $\Sigma_1 \cdots \Sigma_R(V)$ contains R nested annuli A_1, \dots, A_R , that is, annuli such that A_1 is an \mathcal{L}^{Br} -annulus of V_j^{Br} for some $j \in \{1, 2\}$, and for all $k \in \{1, \dots, R-1\}$, A_k is an \mathcal{L}^{Br} -annulus of $A_{k-1} \cup \dots \cup A_1 \cup V_j^{\text{Br}}$. Observe that V is not necessarily equal to V_j^{Br} .

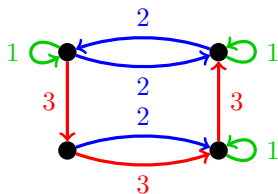
The reasoning goes by induction. First, $\Sigma_R(V)$ contains an \mathcal{L}^{Br} -annulus of V_j^{Br} for some $j \in \{1, 2\}$, since Σ_R contains at least M occurrences of Σ_3^{Br} , according to Lemma 4.9. Now suppose that for some $k \in \{1, \dots, R-1\}$, the pattern $\Sigma_{R-k+1} \cdots \Sigma_R(V)$ contains k concentric nested annuli A_1, \dots, A_k around V_j^{Br} , with $j \in \{1, 2\}$. We use the annulus property for Brun substitutions: thanks to Proposition 3.7 (the annulus property), Proposition 4.4 (strong covering conditions) and Proposition 4.6 (Property A), the patterns $\Sigma_{R-k}(A_1), \dots, \Sigma_{R-k}(A_k)$ are k concentric \mathcal{L}^{Br} -annuli. Moreover, $\Sigma_{R-k}(V_j^{\text{Br}})$ contains an \mathcal{L}^{Br} -annulus of V_1^{Br} or V_2^{Br} . In total this makes $k + 1$ concentric annuli contained in $\Sigma_{R-k} \Sigma_{R-k+1} \cdots \Sigma_R(V)$. They have no face in common thanks to Proposition 2.1, so the induction step holds and Statement (1) is proved.

Statement (2) for Jacobi-Perron substitutions can be proved by the same reasoning. We apply Lemma 4.9 (to generate an annulus around the pattern), Proposition 3.7 (the annulus property), Proposition 4.5 (strong covering conditions), and Proposition 4.8 (Property A).

Note that in the above proofs, we have implicitly used the fact that if a pattern contains U and an \mathcal{L}^{Br} - or \mathcal{L}^{JP} -annulus of U , then it must contain one of the patterns V_i^{Br} or V_i^{JP} . This is proved in Proposition 4.2 for Brun substitutions and in Proposition 4.3 for Jacobi-Perron substitutions. \square

Characterizing sequences for which U is not a seed We now give characterizations of the sequences of Brun and Jacobi-Perron substitutions such that the pattern $U = \diamond$ is not a seed, that is, such that the patterns generated by iterating the substitutions from U do not generate an entire discrete plane.

Theorem 5.2. *Let $(i_n)_{n \geq 1} \in \{1, 2, 3\}^{\mathbb{N}}$ be a Brun-admissible sequence and let $P_n = \Sigma_{i_1}^{\text{Br}} \dots \Sigma_{i_n}^{\text{Br}}(U)$ for $n \geq 1$. Then $(P_n)_{n \geq 0}$ is an increasing sequence of patterns which are contained in the discrete plane $\Gamma_{\mathbf{v}}$, where $\mathbf{v} \in \mathbb{R}_{>0}^3$ is the vector whose Brun expansion is $(i_n)_{n \geq 1}$. Moreover, we have $\bigcup_{n \geq 0} P_n \subsetneq \Gamma_{\mathbf{v}}$ if and only if there exist k such that $(i_n)_{n \geq k}$ is the labelling of an infinite backward path $\dots \xrightarrow{i_{k+1}} \bullet \xrightarrow{i_k} \bullet$ in the following graph:*



We remark that for a finite product of Brun substitutions $\sigma = \sigma_{i_1}^{\text{Br}} \dots \sigma_{i_m}^{\text{Br}}$, Theorem 5.2 can be interpreted in the following way: the sequence of patterns $(\mathbf{E}_1^*(\sigma)^n(U))_{n \geq 1}$ does not generate an entire discrete plane if and only if there is a loop labelled by (some power of) the word $i_m \dots i_1$ in the graph. This is the case for example with $i_4 i_3 i_2 i_1 = 3332$, but not with $i_3 i_2 i_1 = 332$ (it is easy to check that there is no loop labelled by a power of 332). Note that in some cases it is necessary to take a power of $i_m \dots i_1$ for such a loop to exist, for example with $i_2 i_1 = 32$: indeed, there is a loop labelled by 3232 but no loop labelled by 32.

Proof. The fact that $(P_n)_{n \geq 0}$ is an increasing sequence contained in $\Gamma_{\mathbf{v}}$ easily follows from Proposition 2.1 and the fact that U is contained in every discrete plane. We now prove the second part of the proposition. First, recall that the four-vertex graph given in the statement of Theorem 5.2 and the graph $\tilde{\mathcal{G}}^{\text{Br}}$ given in Section 4.4, Figure 2, admit the same sets of infinite backward path labellings with infinitely many 3s. (The four-vertex graph is simply a “reduced” version of $\tilde{\mathcal{G}}^{\text{Br}}$ with some identified vertices.)

Assume now that there does not exist k such that $(i_n)_{n \geq k}$ is the labelling of an infinite backward path in the four-vertex graph. Then, for every $k \geq 0$, there exists $M \geq 0$ such that for every path $f_M \xrightarrow{i_{k+M}} f_{M-1} \dots f_1 \xrightarrow{i_k} f_0$ in \mathcal{G}^{Br} , we have $f_M \in U$. Indeed, the only vertices in $\mathcal{G}^{\text{Br}} \setminus \tilde{\mathcal{G}}^{\text{Br}}$ which are at the end of a backward infinite path with infinitely many 3s are the vertices of U . It then follows from Statement (1) of Proposition 3.11 that $\Sigma_{i_k}^{\text{Br}} \dots \Sigma_{i_{k+M}}^{\text{Br}}(U)$ contains an \mathcal{L}^{Br} -annulus of U . By considering k large enough in the above reasoning, we can now apply Theorem 5.1 with $\mathcal{V} = \{U\}$ and $\mathcal{W} = \{V_1^{\text{Br}}, V_2^{\text{Br}}\}$ to prove that the patterns $P_n = \Sigma_{i_1}^{\text{Br}} \dots \Sigma_{i_n}^{\text{Br}}(U)$ have arbitrarily large combinatorial radius, which implies that $\bigcup_{n \geq 0} P_n = \Gamma_{\mathbf{v}}$.

Conversely, if there exists k such that $(i_n)_{n \geq k}$ is the labelling of an infinite backward path in the four-vertex graph, then Proposition 3.11 Statement (2) implies that the patterns $(\Sigma_{i_k}^{\text{Br}} \dots \Sigma_{i_{k+n}}^{\text{Br}}(U))_{n \geq 0}$ contain no \mathcal{L}^{Br} -annulus of U for any $n \geq 0$, so the discrete plane $\Gamma_{\mathbf{v}}$ is not generated by the patterns $(\Sigma_{i_1}^{\text{Br}} \dots \Sigma_{i_{k+n}}^{\text{Br}}(U))_{n \geq 1} = (\Sigma_{i_1}^{\text{Br}} \dots \Sigma_{i_{k-1}}^{\text{Br}} \Sigma_{i_k}^{\text{Br}} \dots \Sigma_{i_{k+n}}^{\text{Br}}(U))_{n \geq 0}$. \square

Theorem 5.3 below recovers the same characterization as the one stated in [IO94], by using a different notion of generation graphs, and which thus does not require the same succession of lemmas as found in [IO94].

Theorem 5.3. *Let $(a_n, b_n)_{n \geq 1}$ be an admissible Jacobi-Perron sequence and for $n \geq 1$ define $P_n = \Sigma_{a_1, b_1}^{\text{JP}} \cdots \Sigma_{a_n, b_n}^{\text{JP}}(U)$. Then $(P_n)_{n \geq 0}$ is an increasing sequence of patterns which are contained in the discrete plane $\Gamma_{\mathbf{v}}$, where $\mathbf{v} \in \mathbb{R}_{>0}^3$ is the vector whose Jacobi-Perron expansion is $(a_n, b_n)_{n \geq 1}$. Moreover, we have $\bigcup_{n \geq 0} P_n \subsetneq \Gamma_{\mathbf{v}}$ if and only if there exists $\ell \geq 1$ such that for every $k \geq 0$, we have (1) $a_{\ell+3k} = 0$, (2) $a_{\ell+3k+1} = b_{\ell+3k+1}$ and (3) $0 < a_{\ell+3k+2} < b_{\ell+3k+2}$.*

Proof. Let $(i_n)_{n \geq 1}$ be the additive Jacobi-Perron expansion corresponding to $(a_n, b_n)_{n \geq 1}$ (defined in Section 2.3). Since this additive expansion is admissible, it must contain infinitely many 3s or 4s (see Remark 2.7). Similarly as in the proof of Theorem 5.2, the infinite edge labellings of the graph \mathcal{G}^{JP} defined in Section 4.5 give us the desired characterization, thanks to Proposition 3.11 and Theorem 5.1.

Indeed, in the graph \mathcal{G}^{JP} , following an infinite path containing infinitely many 3s or 4s forces us to turn clockwise, visiting f_c, f_e and f_d (and possibly other states between f_d and f_e) cyclically and in this order (see Section 4.5). The result then follows from the definition of the additive decomposition of $\Sigma_{a,b}^{\text{JP}}$ by $\Theta_1, \Theta_2, \Theta_3, \Theta_4$, and by the following facts, which can easily be checked: any path from f_c to f_e in \mathcal{G}^{JP} corresponds to $a_n = 0$; any path from f_d to f_c in \mathcal{G}^{JP} corresponds to $a_n = b_n$; any path from f_e to f_d in \mathcal{G}^{JP} corresponds to $0 < a_n < b_n$. \square

Translates of seeds always occur The next theorem states that the initial pattern U is always sufficient to generate *translates* of patterns with arbitrarily large radius (even though U is not a seed). Figure 1 p. 10 illustrates an example of such a situation for Brun substitutions.

Theorem 5.4. *For every $R \geq 0$ there exists $N \geq 0$ such that*

- (1) $\Sigma_{i_1}^{\text{Br}} \cdots \Sigma_{i_n}^{\text{Br}}(U)$ contains a translate of a pattern P with $\text{rad}(P) \geq R$ for every $(i_1 \cdots i_n) \in \{1, 2, 3\}^n$ that contains more than N occurrences of 3;
- (2) $\Sigma_{a_1, b_1}^{\text{JP}} \cdots \Sigma_{a_n, b_n}^{\text{JP}}(U)$ contains a translate of a pattern P with $\text{rad}(P) \geq R$ for every admissible Jacobi-Perron expansion $(a_1, b_1), \dots, (a_n, b_n)$ such that $n \geq N$.

Proof. Since we are dealing with *translates* of seeds, we cannot use generation graphs like in the previous theorems, we must track “by hand” all the possible sequences of iterations starting from U , and explicitly check that a translate of one of the seeds of Theorem 5.1 is eventually generated. This was done in detail in Lemma 4.12 for Statement (1) and in Lemma 4.13 for Statement (2). The result then follows directly from Theorem 5.1: once a translate of a seed occurs, its images are translates of patterns with arbitrarily large minimal combinatorial radius. (Note that by “linearity” of \mathbf{E}_1^* dual substitutions (Proposition 2.1 (2)), the image of a translate of a pattern is always a translate of its image.) \square

5.2 Topological properties of Rauzy fractals for products of substitutions

Before stating the results of this section, we stress the fact that there are many connections between topological properties of Rauzy fractals (such as being connected or having zero interior point) and number-theoretical properties of the associated dominant Pisot eigenvalue. This is described in more detail in Section 5.4.

Zero interior point We consider a further interpretation of a discrete plane generation property in terms of Rauzy fractals. More precisely, the origin is an interior point of the Rauzy fractal of σ if and only if the patterns $\mathbf{E}_1^*(\sigma)^n(U)$ generate patterns with arbitrarily large minimal combinatorial radius centered at the origin (see [BS05] or [ST09]). An immediate consequence is the following theorem for Brun and Jacobi-Perron substitutions (obtained thanks to Theorem 5.2 and Theorem 5.3), which is illustrated on some finite products of Brun substitutions in Figure 4 p. 26.

Corollary 5.5. *We have:*

- The origin is **not** an interior point of the Rauzy fractal of an admissible product $\sigma_{i_1}^{\text{Br}} \cdots \sigma_{i_n}^{\text{Br}}$ if and only if there exists a cycle labelled by i_1, \dots, i_n in the directed graph given in Theorem 5.2.
- The origin is **not** an interior point of the Rauzy fractal of an admissible product $\sigma_{a_1, b_1}^{\text{JP}} \cdots \sigma_{a_n, b_n}^{\text{JP}}$ if and only if the infinite sequence $((a_1, b_1), \dots, (a_n, b_n))^\infty$ satisfies the condition given in Theorem 5.3.

Note that in the case of Arnoux-Rauzy substitutions, the seed U always generates an entire discrete plane when iterating an admissible sequence, so for Arnoux-Rauzy substitutions the origin is always an interior point of the Rauzy fractal [BJS12]. Also note that the fact that the origin is an interior point of the Rauzy fractal implies pure discrete spectrum of the associated dynamical system [ST09, BST10]. It also has interpretations in number-theoretical terms, which will be discussed in Sections 5.3 and 5.4.

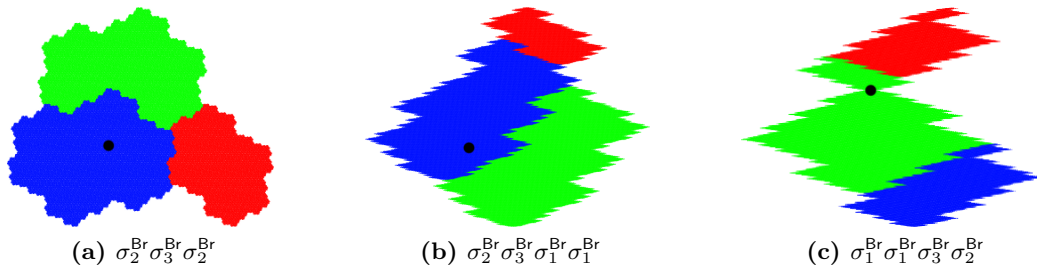


Figure 4: Some Rauzy fractals of products of Brun substitutions. Only (c) does not have the origin as an interior point. An example of a product of substitutions whose Rauzy fractal has zero as an interior point but whose reverse product does not is provided by (b) (and its reverse product (c)). The fractals of (a) and (c) correspond to the two examples of substitutions shown in Figure 1 p. 10.

Connectedness The following result states that the Rauzy fractals of the substitutions in the families we consider are all connected. It has been proved in the case of Arnoux-Rauzy substitutions using similar tools [BJS14]. One interest of such a topological property comes from the fact that the Rauzy fractal provides a geometric realization for the two dynamics that act on (X_σ, S) , namely the dynamics of the shift (the toral translation in the case of pure discrete spectrum), but also the dynamics provided by the (positive entropy) action of the substitution. Indeed, the subtiles of the Rauzy fractals form a Markov partition for the toral automorphism given by the incidence matrix of the substitution (in case of pure discrete spectrum). See Corollary 5.7 and 5.8 below for more detail.

Corollary 5.6. *Let σ be an admissible finite product of Brun or Jacobi-Perron substitutions. The Rauzy fractal \mathcal{T}_σ and its subtiles $\mathcal{T}_\sigma(1)$, $\mathcal{T}_\sigma(2)$, $\mathcal{T}_\sigma(3)$ are all connected.*

Note that such fractals are not necessarily simply connected. This raises the question of characterizing the admissible finite product of Brun or Jacobi-Perron substitutions that yield a simply connected fractal, which is an interesting and seemingly difficult question.

Proof. Let σ be an admissible product of Brun or a Jacobi-Perron substitutions. Thanks to Propositions 4.4 and Proposition 4.5, the patterns $\mathbf{E}_1^*(\sigma)^n(U)$ are all \mathcal{L}^{Br} -connected (or respectively \mathcal{L}^{JP} -connected). Let \mathcal{L} stand for \mathcal{L}^{Br} or for \mathcal{L}^{JP} according to the fact that σ is an admissible product of Brun or a Jacobi-Perron substitutions. Since \mathcal{L} consists of connected patterns only, the definition of \mathcal{L} -coverings implies that every \mathcal{L} -covered pattern is path-connected.

It follows that the patterns $\mathbf{E}_1^*(\sigma)^n(U)$ are connected for every $n \geq 0$. Consequently, since the Rauzy fractal \mathcal{T}_σ of σ is a Hausdorff limit of connected sets, it is also connected. The same reasoning yields the connectedness of $\mathcal{T}_\sigma(1)$, $\mathcal{T}_\sigma(2)$, $\mathcal{T}_\sigma(3)$, since it can easily be checked that the images of $[0, 1]^*$, $[0, 2]^*$, $[0, 3]^*$ are all eventually \mathcal{L} -connected. \square

5.3 Dynamical properties of products of substitutions

The tools developed in Section 3 enable us to establish dynamical properties of the infinite product families consisting of arbitrary finite products of Brun or Jacobi-Perron substitutions. This will be done in particular in Corollary 5.7 below. Even though our tools are not fully algorithmic, they can be seen as an extension of the algorithms used in the study of a *single* substitution, where many algorithms have been developed based on the constructions of graphs which are proved to be finite thanks to the Pisot property (see [ST09, BST10, AL11]).

Corollary 5.7. *Let σ be an admissible finite product σ of Brun or Jacobi-Perron substitutions. Then the system (X_σ, S) has pure discrete spectrum. Furthermore, the Rauzy fractal of σ (and its subtiles) provide the domain and partition of a symbolic natural coding of this translation. These subtiles are moreover connected bounded remainder sets for this translation.*

Proof. This theorem follows from Theorem 5.4 together with [IR06], where it is proved that (X_σ, S) is measure-theoretically conjugate to a toral translation if the sequence $(\mathbf{E}_1^*(\sigma)^n(U))_n$ satisfies the arbitrarily large disks covering property, that is, if it contains translates of patterns with arbitrarily large minimal combinatorial radius. One then deduces that (X_σ, S) is a symbolic natural coding of this translation, with the Rauzy fractal \mathcal{T}_σ together with its subtiles providing the fundamental domain and its partition (the translation vectors $\pi_c(\mathbf{e}_i)$ of the exchange of pieces (\mathcal{T}_σ, E) described in Section 2.4 provide the restriction of the translation to the fundamental domain \mathcal{T}_σ). Connectedness is a consequence of Corollary 5.6.

It remains to check that the subtiles of the Rauzy fractal provide bounded remainder sets. A bi-infinite word $u \in \mathcal{A}^{\mathbb{Z}}$ is said to have bounded discrepancy if there exist f_i , $i \in \mathcal{A}$, and $C > 0$ such that $\limsup_{i \in \mathcal{A}, N \in \mathbb{N}} ||u_{-N} \cdots u_0 u_1 \cdots u_N|_i - (2N + 1)f_i| \leq C$. If a coding u of a minimal translation is such that its elements have bounded symbolic discrepancy, then the elements of the associated partition are bounded remainder sets. We conclude by noticing that bi-infinite words in (X_σ, S) have bounded symbolic discrepancy since σ is an irreducible Pisot substitution (see e.g. [Ada04b]). \square

Markov partition for toral automorphisms Markov partitions for the toral automorphisms $(\mathbb{T}^3, \mathbf{M})$ are known to exist in the case where \mathbf{M} is hyperbolic (no eigenvalues lie on the unit circle) [Bow08, Man02]. There is a general way to construct them explicitly in dimension two with rectangles [AW70, Adl98], but no such constructions are known for dimensions $d \geq 3$. Indeed a result of Bowen [Bow78] states that such partitions must have fractal boundary if $d \geq 3$, which reduces the hopes of finding explicit Markov partitions with a simple shape.

In the particular case where \mathbf{M} is the incidence matrix of a Pisot unimodular substitution (or the companion matrix of a unit Pisot number β in the framework of beta-numeration), Rauzy fractals provide a way to construct explicit Markov partitions for $(\mathbb{T}^n, \mathbf{M})$. Such an approach was initiated in [IO93, KV98, Pra99] and was proved to be successful in the Pisot case [IR06, Sie00], provided that the associated Rauzy fractal satisfies some tiling properties (which are equivalent to pure discrete spectrum of the associated symbolic dynamical system).

Consequently, another dynamical application of Theorems 5.1 and 5.4 is that for every matrix \mathbf{M} which is the incidence matrix of an admissible product of Brun or Jacobi-Perron substitutions, an explicit Markov partition can be constructed for the toral automorphisms $(\mathbb{T}^3, \mathbf{M})$.

Corollary 5.8. *Let σ be an admissible finite product σ of Brun or Jacobi-Perron substitutions. Then, there exists a three-set partition of the torus \mathbb{T}^3 with connected elements which is a Markov partition for the toral automorphism $(\mathbb{T}^3, \mathbf{M}_\sigma)$ provided by the incidence matrix \mathbf{M}_σ of σ .*

This result is a consequence of the results from [IO93, Pra99, Sie00] (which establish links between the Rauzy fractal of σ and Markov partitions of \mathbf{M}_σ), together with Corollary 5.7. A similar result for Arnoux-Rauzy substitutions is stated in [BJS14]. Also note that the connectedness of the domains of the Markov partition in question is due to the connectedness of Rauzy fractals (Corollary 5.6 in Section 5.2).

5.4 Number-theoretical implications

Fractal tiles for cubic number fields By a result of Paysant-Le-Roux and Dubois [PD84] together with Corollary 5.7, we are able to associate natural codings with respect to a domain and a partition provided by Rauzy fractals with every cubic field.

Theorem 5.9. *For every cubic extension \mathbb{K} of \mathbb{Q} , there exist $\alpha, \beta \in \mathbb{K}$ and a three-letter unimodular irreducible Pisot substitution σ such that $\mathbb{K} = \mathbb{Q}(\alpha, \beta)$ and such that the toral translation $(\mathbb{T}^2, x \mapsto x + (\frac{\alpha}{\beta}))$ is measure-theoretically conjugate to (X_σ, S) . Furthermore, it admits a symbolic natural coding with respect to a partition made of connected bounded remainder sets.*

Proof. According to [PD84, Proposition IV], any cubic field is generated by some Pisot number β with minimal polynomial $P(X) = X^3 - c_2X^2 + c_1X - 1 = 0$, where $c_2 \geq 2c_1 - 2$ and $c_1 \geq 3$. If we set $\sigma = \sigma_{0,1}^{\text{JP}} \sigma_{0,1}^{\text{JP}} \sigma_{c_1-3, c_2-c_1}^{\text{JP}}$ (an admissible product because $c_2 - c_1 > c_1 - 2$), then \mathbf{M}_σ has characteristic polynomial $P(X)$. Since \mathbf{M}_σ is primitive, it admits a positive eigenvector \mathbf{v} with $\mathbf{v}_1 + \mathbf{v}_2 + \mathbf{v}_3 = 1$ and we have $\mathbb{Q}(\mathbf{v}_2, \mathbf{v}_3) = \mathbb{K}$. The Jacobi-Perron expansion of \mathbf{v} has period $(c_1 - 3, c_2 - c_1)(0, 1), (0, 1)$ (this is an admissible expansion). Now consider the projection $\pi_{\mathbf{v}, \mathbf{1}^\perp}$ along \mathbf{v} onto the antidiagonal plane with normal vector $(1, 1, 1)$. This projection expresses as

$$\pi_{\mathbf{v}, \mathbf{1}^\perp}(x, y, z) = (\mathbf{v}_2(x + y + z) - x)(\mathbf{e}_3 - \mathbf{e}_1) + (\mathbf{v}_3(x + y + z) - y)(\mathbf{e}_3 - \mathbf{e}_2).$$

In particular, the vector $\pi_{\mathbf{v}, \mathbf{1}^\perp}(\mathbf{e}_3)$ expressed in the basis $(\mathbf{e}_3 - \mathbf{e}_1, \mathbf{e}_3 - \mathbf{e}_2)$ has coordinates $(\mathbf{v}_2, \mathbf{v}_3)$. Consider the translation by the vector $(\pi_{\mathbf{v}, \mathbf{1}^\perp} \mathbf{e}_3)$ modulo the lattice generated by $\pi_{\mathbf{v}, \mathbf{1}^\perp}(\mathbf{e}_3 - \mathbf{e}_1)$ and $\pi_{\mathbf{v}, \mathbf{1}^\perp}(\mathbf{e}_3 - \mathbf{e}_2)$. It is measure-theoretically conjugate to the translation by $(\mathbf{v}_2, \mathbf{v}_3)$ modulo \mathbb{Z}^2 . The projection $\pi_{\mathbf{v}, \mathbf{1}^\perp}$ can be used to define a Rauzy fractal for σ instead of the projection π_σ used in the definition of Rauzy fractals in Section 2.4. This means that we are working with a domain exchange whose translation vectors are given by π instead of π_c . According to Corollary 5.7, the domain exchange $(\mathcal{T}_\sigma(i), E)$ (obtained with π_c) factors onto a translation, so the domain exchange provided by π also factors onto a translation [IR06]. This yields the desired result by taking $\alpha = \mathbf{v}_2, \beta = \mathbf{v}_3$. \square

Observe that the existence of a symbolic natural coding provides equidistribution properties for the associated translation. See [Ada04a] for an illustration in the case of one-dimensional translations.

Number-theoretical applications of topological properties of Rauzy fractals Several of the topological properties of Rauzy fractals have direct number-theoretical consequences. Consider first the framework of (Pisot) β -numeration. The connectedness of the Rauzy fractal (called central tile in this context) is conjectured to guarantee explicit relations between the norm of β and the β -expansion of 1 [AG05]. Furthermore, the properties of rational numbers with purely periodic β -expansions are now known to be closely related to the shape of the boundary of the Rauzy fractal [ABBS08, AFSS10]. More generally, at the interplay between substitutions and Diophantine approximation for cubic number fields, explicit computation of the size of the largest ball contained in the Rauzy fractal is used in [HM06] to determine the sequence of best approximations with respect to a specific norm, for some two-dimensional vectors provided by non-totally real cubic Pisot units. See also [IFHY03, IY07] for a study of the limit set of

the points $\sqrt{q}(\|q\alpha\|, \|q\beta\|)$, with q being a positive integer, and with (α, β) such that $(1, \alpha, \beta)$ forms a basis of a non-totally real cubic number field. This limit set is known to be a union of homothetic ellipses centered at the origin. The question of when the closest of these ellipses is provided by Brun algorithm is investigated in [IFHY03, IY07].

The fact that the origin is an interior point of the Rauzy fractal is also a particularly interesting property. It is closely related to the so-called finiteness properties in numeration. Indeed, the *(F) property* (introduced and called finiteness property in [FS92]), which expresses some finiteness properties of digital expansions in a Pisot base β , is equivalent to the fact that the origin is an interior point of the central tile associated with β [Aki99, Aki02]. Several variants of the (F) property have then been proposed, one of them being the *extended (F) property* introduced in [BS05, FT06] to extend the classical (F) property to the numeration systems associated with substitutions (called the Dumont-Thomas numeration). The extended (F) property can again be stated in topological terms: it holds if and only if the origin is an interior point of the Rauzy fractal associated with σ (see [ST09, BST10]). Consequently, Corollary 5.5 above, which characterizes the products of Arnoux-Rauzy, Brun and Jacobi-Perron substitutions for which the origin is an interior point of the Rauzy fractal, also provides a characterization of when such a finite product of substitutions satisfies the extended (F) property. Furthermore, this is a sufficient condition for pure discrete spectrum which can also be interpreted as a coincidence type condition [BST10].

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