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Patrice Raveneau, Hervé Rivano

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Patrice Raveneau , Hervé Rivano

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Abstract: This document presents the first part of the DTN (Delay Tolerant Network) study within IOT (Internet Of Things) context. The motivation for using generic protocols able to handle the constraints due to the IOT is highlighted with the choice of Bundle Protocol. A study of existing implementations of this protocol is realised within a sensor context. We justify the choices made for our implementation, then we define the mechanisms which we will test with the IOT-Lab platform by following the protocol of tests we have developed. The performance analysis are also presented.

Key-words: DTN, IOT, IOT-Lab

**RESEARCH CENTRE
GRENOBLE – RHÔNE-ALPES**

Inovallée
655 avenue de l'Europe Montbonnot
38334 Saint Ismier Cedex

Scénario de tests pour les DTN pour l'IOT

Collaboration III Urbanet

Résumé : Ce document présente la première partie de l'étude des DTN (Delay Tolerant Network) dans un contexte d'IOT (Internet Of Things). La motivation de l'utilisation de protocoles génériques capables de supporter les contraintes inhérentes au contexte IOT est mise en avant avec le choix du Bundle Protocol. Une étude des implantations existantes de ce protocole est faite dans un contexte capteur. Nous proposons notre implantation en justifiant les choix réalisés, puis nous définissons des mécanismes que nous testerons avec la plate-forme IOT-Lab en suivant le protocole de tests que nous avons élaboré. Les performances à analyser sont également présentées.

Mots-clés : DTN, IOT, IOT-Lab

1 Introduction

The Internet Of Things (IOT) refers to one part of the smart world developing recently where each device, each person and each object is connected through one or several communicating devices and interact with each other [1].

The objective of this document is to define protocol architectures to respond to the heterogeneity that may arise in a smart context with several devices with several mobility behaviours. We also define a test scenario to analyse the performance of a network using our proposal in an IOT context. For an IOT scenario, there is a lot of data communications but generated by small data volumes such as a measurement of temperature, air quality or parking occupancy.

We describe the context of using sensors and raise the difficulties of this context. Subsequently, we propose a protocol architecture fitting the requirements of several IOT scenarios before considering the implementations constraints on small devices. Once this framework defined, we propose mechanisms able to use mobility of devices to reduce the number of useless communications in order to better use network capacity. Eventually, we define the tests scenario to validate our proposals and analyse the performance of the network.

2 Context

The sensors will be considered static or mobile as required by the intended IOT application. In the case of static nodes, their distribution may be regular or random depending on the conditions of the deployment. However, the literature is abundant on static sensor networks and addresses several problems. The major problems are related to energy management [2]. The solutions proposed to solve this problem based on routing protocols and algorithms using the geographical distribution of nodes [3]. Load distribution techniques in the network [4] and selecting routing trees minimizing the period of data collection [5] allow also to reduce the energy consumption. Energy can also be saved by using an optimal transmission range [6] and cutting the network into a number of optimal groups [7].

Sensor networks deployed within a city might count thousands of nodes. Nevertheless, for our tests, we will select key elements to represent the several nodes behaviours that might occur in such a network.

For our study, we consider static and mobile sensors communicating to a single sink through one or several mobile gateways. We consider static sensors to monitor environment and mobile sensors which are able to collect data from the static ones. We use mobile sensors as relays between two unconnected sensors segments of a sensor network.

The information exchanged over the network are control messages of data. Included in the control signaling for determining routes between nodes and data acknowledgments.

3 Architecture protocol

We now present the protocols used within the network.

We highlight the main constraints of our network, the heterogeneous technologies from urban environment, the lack of connectivity, the dynamic mobility and long links interruptions. Protocols like Internet Protocol (IP) are not able to manage these constraints. If a link is unavailable, then the datagram is lost. A reliable protocol such as Transmission Control Protocol (TCP) will handle retransmissions. But the duration of such interruptions are too large to allow TCP to be effective. That's why we decided to focus on Delay/Disruption Tolerant Networks (DTN).

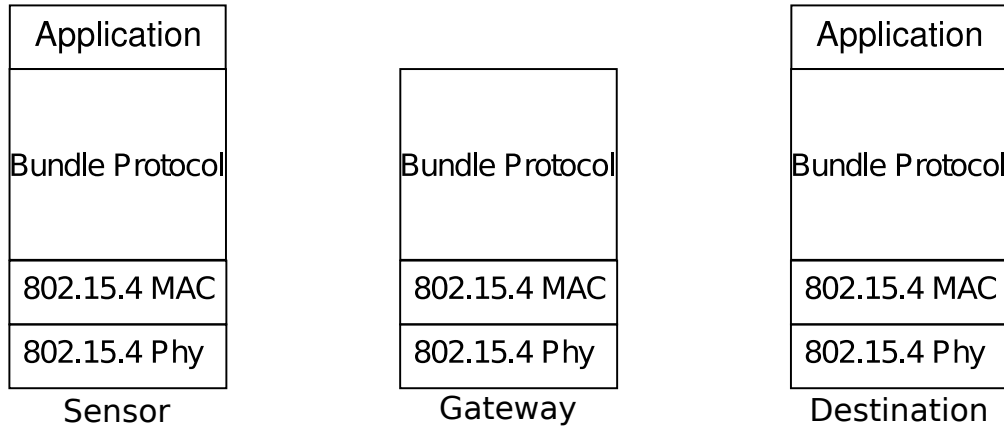


Figure 1: Architecture with Bundle Protocol

The Bundle Protocol which meets the requirements of DTN architecture is able to manage the key constraints. In addition, it allows to interconnect networks with different protocol architectures.

3.1 A DTN architecture

For sensors in an IOT purpose, it is necessary to limit the number of protocols layers to minimize the use of memory resources. Thus we arrived at the definition of the protocol architecture in figure 1.

On the protocol architecture we defined, the Bundle Protocol is implemented on each system, even on sensors. Sensors are able to store data they do not have themselves created. We clarified that, given the small memory sensors, we had to limit the number of protocols used on that equipment. That's why we do not use network layer as IPv6 Low power Wireless Personal Area Networks (6LoWPAN). Indeed, the protocol 6LoWPAN [8] brings nothing more to our architecture and can not manage long discontinuities. The Bundle Protocol actually allows naming all entities in a network. In addition, the fields indicating the source and destination are Self-Delimiting Numeric Value (SDNV). The goal is to reduce the overhead on sensors communications with small headers.

The main drawback of this architecture is that it consumes a part of the sensors memory. Adding a protocol level using several mechanisms such as custody management, storage or reports consumes memory resources. However, in a context where nodes are mobile, it is necessary that the terminal can store data from other sources in order to increase the probability of delivery.

It will be useful to be able to select the best data to share and store. Such a solution makes better use of resources. Similarly, the use of acknowledgment mechanisms increases the memory usage; however it reduces the transmissions of Bundles that have already been delivered to a gateway.

We do not consider the Bundle Protocol as an overlay protocol, linking very different networks, but as a way to overcome the constraints caused by long interruptions and dynamic mobility scenarios.

The use of the Bundle Protocol on the entire network allows to consider gateways as a relay and not as the destination data from the sensors network. Nevertheless, this relay plays a role of custodian and reduces the load in the sensor network indicating that Bundles left this part

Table 1: Advantages and drawbacks of the proposed architecture

Advantages	Drawbacks
Full DTN architecture	Overhead in the network
All equipments interacting	Memory limitation on sensors

of the network. The last node to have transmitted Bundles to the gateway removes it from its memory.

The architecture we have proposed is intended generic. We do not wish to provide architectures that are tailored to a single specific scenario. Our architecture can adapt to other scenarios. The mobile gateway could be replaced by a static one, the proportion of static and mobile nodes could change and the architecture would remain the same.

3.2 Bundle Protocol implementation on sensors

We want to prove that the Bundle Protocol could be implemented on sensors. For this, we studied the existing codes before achieving ours.

3.2.1 ION

The first implementation we present is Interplanetary Overlay Network (ION) developed by Jet Propulsion Laboratory (JPL). This implementation was performed with the goal of being deployed in a context of interplanetary communications. To ensure the robustness of their systems, the missions of the JPL do not use dynamic memory allocation. In addition, for interplanetary missions, it is essential to recover data, even partial, so overhead induced by the Bundle Protocol should be minimal [9]. ION is a complete implementation as it offers adaptation layers with TCP, User Datagram Protocol (UDP) and Licklider Transmission Protocol (LTP). The implementation of LTP is also provided.

While this implementation contains all features and limits the use of dynamic memory mechanisms, it is not appropriate to our sensor context. The use of pre-allocation memory mechanisms is perfect for our purpose. Nevertheless, the existence of various libraries increases the code size and does not meet one of the major constraints implementing on sensor, low memory.

We focus now on the Bundle Protocol reference implementation.

3.2.2 DTN2

DTN2 is the implementation of the Research Group DTN and meets requirements of [10]. This implementation provides several layers including IP and Ethernet convergence layers. Convergence layers to the level 2 and 3 protocols show that it is possible to deploy the Bundle Protocol on low-level protocols and then limit the overhead, which is an important aspect in the IOT context.

This implementation, however, has significant disadvantages for implementation on sensors. Unlike ION, the usage of dynamic memory is not restricted. In addition, the programming language used is C++ which is too resource hungry for deployment on sensors.

We turn our study to implementations thought for embedded systems.

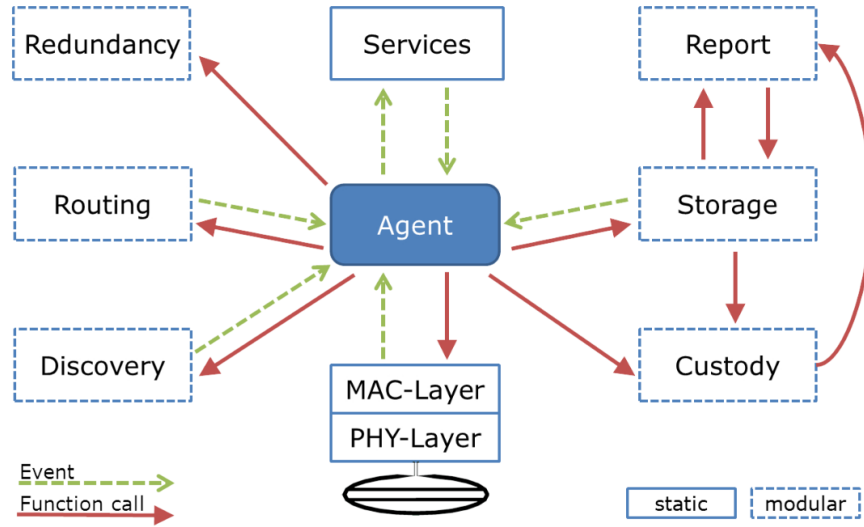


Figure 2: Architecture of μ DTN taken from [13]

3.2.3 IBR-DTN

IBR-DTN is an implementation for embedded systems [11]. Unlike other solutions designed for small systems IBR-DTN can work with equipment using other implementations of the Bundle Protocol [12]. In addition to convergence layers on UDP and TCP, this implementation also includes a layer of convergence to the IEEE 802.15.4 protocol. This implementation is adapted to an IOT context.

The problem is that the code size of this implementation is still larger than the 48K that we have at our disposal. We continue our search for implementations that would suit our equipment.

3.2.4 μ DTN

μ DTN is designed to be suitable for sensor networks [13]. Unlike DTNLite [14] that uses the DTN concept without implementing the Bundle Protocol and ContikiDTN [15] using the Bundle Protocol over a convergence layer on TCP, μ DTN works directly above a convergence layer to the IEEE 802.15.4 protocol. We present in figure 2, the architecture implementing μ DTN.

However, while μ DTN reduces to an absolute minimum the protocol stack, the code is still too large to be deployed on the sensors of IOT-Lab. We have therefore chosen to offer a lighter version of μ DTN, we name nanoDTN.

We synthesize in table 2 results analysis of existing implementations compared to metrics that we are interested in for our deployment.

3.2.5 nanoDTN

We named our architecture, nanoDTN, since it is based on μ DTN architecture, which we removed features while retaining basic interoperability provided by the μ DTN. This implementation being too large to be carried on WSN430, we created ours, lighter, for deployment. We kept the operating system Contiki [17] using the C language and therefore allows the use of multiple standard libraries.

Table 2: Comparison of different implementations for deployment on WSN430

Implantation	Language Appropriate	Memory pre-allocated	Size code	Bundle Protocol interoperable	Layer low convergence
ION [9]	✓	✓	✗	✓	✗
DTN2 [16]	✗	✗	✗	✓	✗
IBR-DTN [11]	✓	✓	✗	✓	✗
μ DTN [13]	✓	✓	✗	✓	✓
DTNLite [14]	✓	✓	✓	✗	✓
ContikiDTN [15]	✓	✓	✓	✓	✗

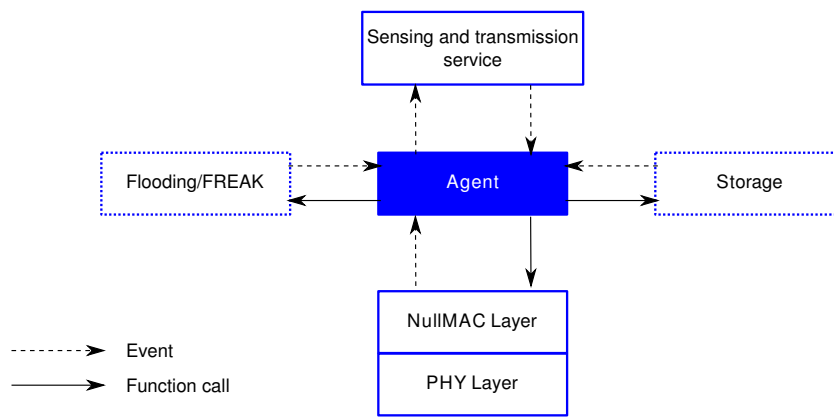


Figure 3: nanoDTN architecture

We started by removing the functions of reports. These features are useful for custody. From there, we have also chosen to remove the management of custody that could no longer exist without the reports. We note that this is not because we delete the functionality of custody that nodes can no more carry Bundles from other sources. The difference is that now a node carrying a bundle will no longer be able to withdraw it from its memory after custody taken for this Bundle by another node. A Bundle will be removed because delivered to the destination, acknowledgment reception for this Bundle or expiration of its lifetime.

In order to reduce the memory footprint, we have also chosen to delete the redundancy function that allowed not to convey a new bundle that could have been delivered. We follow on functionalities removal by deleting the Neighbor Discovery. So we chose periodic transmissions.

Our field of application does not require reliability. We decided to use only one MAC layer, which does not support carrier sense and provides no reliability. This MAC transmits data to the upper level and the lower level. We have also considered that these sensors could not have other features than collection and transmission of data. So we have restricted to a service composed of sensing, gathering and periodic transmissions. Finally, we restricted the routing to the routing functionality that is the simplest, a routing by flooding. This is a simple routing, robust communications intensive but very few resources, so adapted to our context.

Figure 3 shows the architecture of the nanoDTN implementation.

We have never changed functionality inherent to the Bundles headers to maintain interoperability with other implementations of the Bundle Protocol.

We will modify the routing function to test the FREAK scheme that we define next.

4 Using meetings with gateway = FREAK

In [18], the authors outline an assumption by writing that *the future rate nodes meet can be roughly predicted by the previous rate*. We consider this hypothesis and suggest a mechanism we call FREAK.

4.1 The FREAK algorithm

FREAK (Frequency Routing, Encounters And keenness) is so named because it offers a frequency-based routing meeting and is optimistic. Our proposal is optimistic because unlike solutions like Encounter-Based Routing (EBR) [18], MaxProp [19] or Probabilistic Routing Protocol using History of Encounters and Transitivity (PRoPHET) [20] that use information meeting on the set of nodes of the network, our algorithm uses only the frequency of meeting with all destination nodes. This restriction allows to respect the constraint on the memory of nodes we have analysed earlier.

One might expect that replication is not appropriate in a context where memory limitations are strong. However, by depriving themselves of this mechanism, we decrease the number of contacts per message. With the replication, a message has as many opportunities as a carrier node has encounters. In the case of transmissions without replication, there is only one message bearer. The transmission reduces the load but it also reduces the number of useful contacts per message and provides worse than using replication mechanisms, more specifically in an opportunistic environment.

Our idea is to assume that some terminals will encounter more often the destination than others. Rather than determining the number of meetings with each node [18], we calculate the average frequency of encounters with the gateway. The best relays are the nodes meeting more often destination. In conclusion, the higher the frequency of meeting with the destination, the better this relay is. The FREAK operation is simple. When two nodes meet the node which sees less often the base station sends copies of its Bundles to another. The algorithm 1 summarizes this operation.

The metric is updated only when the node meets the gateway and carries Bundles to deliver to it.

4.2 The communication protocol for FREAK

FREAK is the algorithm which, based on a metric of encounters of nodes with the destination, decides whether a node has to transmit Bundles or not. We need to use a protocol to exchange this metric when two nodes are in contact.

We define our discovery protocol such as each node is periodically sending a control message containing its address and the value of its metric. When a node receives such a message, it creates a new message containing its own value of the metric and sends this answer to the former node.

Depending on the values of the metric, the FREAK algorithm computes on each node whether it has to send or receives Bundles. There is no reliability provided, but the goal of such a scheme is to decrease the number of useless transmissions when using a routing scheme such as flooding and also to use low memory and computation skills to fit on a sensor.

Algorithm 1 The FREAK mechanism

```

Let A be the local node
nbrContacts = 0
freq = 0
for all met node (called B) do
  if B is the destination then
    nbrContacts ++ # If this is the destination the metric is updated
    freq =  $\frac{nbrContacts}{CurrentTime}$ 
    Send all Bundles # and all Bundles are transmitted to B
    Delete delivered Bundles
  else
    if  $contact_{freq}(A) < contact_{freq}(B)$  then # If met node sees the destination more
    frequently
      send Bundles to B # then all Bundles are transmitted
    else
      wait Bundles from B # otherwise, expecting to receive those from B
    end if
  end if
end for

```

Table 3: table summarizing the characteristics of the scenario

Traffic renewal	Periodical
Lifetime	Few hours
Sensor Types	Static and mobile
Volume of data	10 bytes

5 Selected scenario

We select a scenario of tests compliant with an urban environment. We consider that a portion of the nodes would remain static and that other nodes which could be carried by citizens, vehicles, public transports would be mobile. Then, the mobile nodes would serve as gateways or as Mobile Ubiquitous LAN Extension (mule) to gather data from static nodes. Data collected by such a system might represent a huge volume, but from a lot of sources providing low volume information, such as measurements of temperature, noise or atmospheric pollution.

5.1 Traffic

Measurements are performed periodically. The extent of the monitored areas is large. It will be necessary that the network be able to handle data from different sources.

We assume that the lifetime of the collected data should be within a few hours. While gathering data from an urban environment, a part of the collected data might represent an interest for near real time applications while most of the collected data is transmitted, then stored to be analysed offline. Applications such as smart parking or sharing transport systems require near real time data.

We synthesize the choices made in the table 3 .

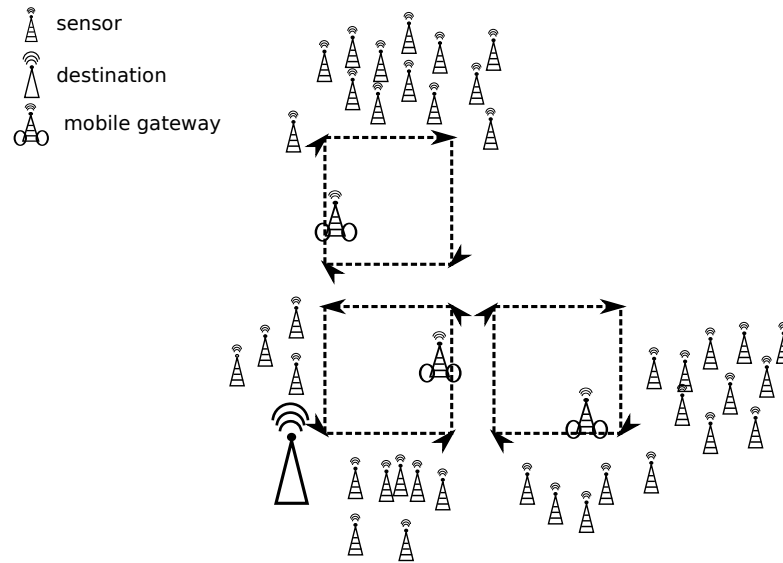


Figure 4: Topology Scenario

5.2 Hardware constraints

In the network of sensors, nodes are static or mobile. In the case where the sensors are mobile, this would mean that they would be disposed on human beings or vehicles. Having no knowledge of the frequency of meetings of these elements, we cannot program adequate awakenings. In order to reduce the number of missed contacts with other terminals, we choose to let radio equipment on continuously.

5.3 Topology

As we mentioned earlier, the network will deal with static and mobile nodes. In order to take benefit from DTN and the Bundle Protocol, some parts of the network will remain disconnected. Only some mobile nodes will handle a partial connectivity between two parts of the network.

The topology we thought about to test our mechanism contains three network segments with static nodes. These nodes can only communicate with nodes from another segment thanks to one of the mobile nodes. We also consider three mobile nodes which will have a pattern whose parameters such as speed and period will change from one test to another.

This topology is shown on figure 4.

We note that each mobile node visits only one static area. We will be able to analyse the performance of a network depending on the use of a basic routing scheme such as flooding or on the use of a mechanism which tries to infer the future based on measurements of the past encounters.

5.4 Technological choices

The sensors that we will use will be the WSN430 type. These sensors are the equipments of the IOT-Lab platform. These elements have an additional external flash memory of 1 Megabytes. With our data on 10 bytes, these terminals can store several tens of thousands of data.

Table 4: Selected settings

Sensors memory	48 kB
Number of sensors	20
Distribution of mobile/static nodes	50%/50% – 25%/75%

These devices have limited memory capacity, with 48 kilobytes of ROM and 10 kilobytes of RAM. Depending on the version of the WSN430, transmissions are either made with a frequency of 868 MHz or 2400 MHz with a 802.15.4-compliant chipset [21]. These nodes are the IOT-Lab nodes that we will use to achieve the implementations and experiments.

5.5 Performance to evaluate

We define here the performance that we will evaluate with our tests.

- The delivery ratio of Bundles which reached the destination.
- The overhead of communications is the ratio of transmitted Bundles within the network over the number of delivered Bundles.
- The delivery delay between the source and the destination.
- The energy consumption of each sensor, in order to analyse how communications impact the energy drain of the sensor node.

6 Conclusion

We defined in this report, an architecture compliant with several constraints of IOT purpose, such as mobility management and low memory devices, and fitting several applications it is not designed for one particular application. The implementation of such an architecture was analysed and lead to the definition of a new one, nanoDTN. Based on this architecture, we propose a DTN routing algorithm to determine whether or not a node shall send its Bundles during a contact. Finally, we design the tests to analyse our architecture, algorithm and control protocol.

The next step is to run the tests and draw conclusions about the use of DTN and the Bundle Protocol in an IOT environment. The interoperability gain is obvious, since we use a standard protocol, there is no need of applicative gateways, which are often used in a sensor context. Nevertheless, the question that we will handle is, is the network performance good enough for IOT applications.

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**RESEARCH CENTRE
GRENOBLE – RHÔNE-ALPES**

Inovallée
655 avenue de l'Europe Montbonnot
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