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Siegfried M. Rump, Florian Bünger, Claude-Pierre Jeannerod

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# Improved Error Bounds for Floating-Point Products and Horner's Scheme

**Siegfried M. Rump · Florian Bünger · Claude-Pierre Jeannerod**

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**Abstract** Let $\mathbf{u}$ denote the relative rounding error of some floating-point format. Recently it has been shown that for a number of standard Wilkinson-type bounds the typical factors $\gamma_k := k\mathbf{u}/(1-k\mathbf{u})$ can be improved into $k\mathbf{u}$, and that the bounds are valid without restriction on $k$. Problems include summation, dot products and thus matrix multiplication, residual bounds for $LU$- and Cholesky-decomposition, and triangular system solving by substitution.

In this note we show a similar result for the product $\prod_{i=0}^{k} x_i$ of real and/or floating-point numbers $x_i$, for computation in any order, and for any base $\beta \geqslant 2$. The derived error bounds are valid under a mandatory restriction of $k$. Moreover, we prove a similar bound for Horner's polynomial evaluation scheme.

**Keywords** floating-point product · IEEE 754 standard · Wilkinson type error estimates · Horner scheme

**CR Subject Classification** 65G50 · 65F05

Siegfried M. Rump
Institute for Reliable Computing, Hamburg University of Technology, Schwarzenbergstraße 95, Hamburg 21071, Germany, and Faculty of Science and Engineering, Waseda University, 3-4-1 Okubo, Shinjuku-ku, Tokyo 169-8555, Japan (rump@tuhh.de).

Florian Bünger
Institute for Reliable Computing, Hamburg University of Technology, Schwarzenbergstraße 95, Hamburg 21071, Germany (florian.buenger@tuhh.de).

Claude-Pierre Jeannerod
Inria, Laboratoire LIP (CNRS, ENS de Lyon, Inria, UCBL), Université de Lyon, 46 allée d'Italie 69364 Lyon cedex 07, France (claude-pierre.jeannerod@ens-lyon.fr).

## 1 Introduction and notation

Denote by $\mathbb{F}$ a set of floating-point numbers with $p$ digits precision in base $\beta$, and with operations according to IEEE 754 standard [3] in rounding to nearest with any tie breaking rule. Then, $\mathbf{u} := \frac{1}{2}\beta^{1-p}$ denotes the relative rounding error unit. Throughout the paper we assume that $\beta \geqslant 2$ and $p \geqslant 1$, and that neither overflow nor underflow occurs.

As usual, for $\circ \in \{+, -, \cdot, /\}$ and $a, b \in \mathbb{F}$, the floating-point result of an operation $a \circ b$ is defined to be $\mathrm{fl}(a \circ b)$ for a rounding to nearest $\mathrm{fl}\colon \mathbb{R} \to \mathbb{F}$. It follows [2, p. 38] that $|\mathrm{fl}(x) - x| \leqslant \mathbf{u}|x|$ for $x \in \mathbb{R}$, and in particular

$$|\mathrm{fl}(a \circ b) - (a \circ b)| \leqslant \mathbf{u}|a \circ b|. \tag{1.1}$$

For matrices $A \in \mathbb{F}^{m\times k}$ and $B \in \mathbb{F}^{k\times n}$, denote by $\widehat{C}$ the floating-point result of the exact product $C := AB$ computed using (blocked versions of) the classical algorithm, with any ordering for the inner products. A rounding error analysis *à la* Wilkinson then leads typically to $|\widehat{C} - C| \leqslant \gamma_k|A||B|$ with $\gamma_k := \frac{k\mathbf{u}}{1-k\mathbf{u}} = k\mathbf{u} + O(\mathbf{u}^2)$; see for example [2, p. 71]. This standard estimate has been improved in [4] into

$$|\widehat{C} - C| \leqslant k\mathbf{u}|A||B| \tag{1.2}$$

without restriction on the integer $k$ and, in [9], similar improvements have been obtained for the residuals of the computed LU and Cholesky factors as well as for triangular system solutions.

A similar result was recently shown by Graillat, Lefèvre, and Muller [1] for binary arithmetic:

**Theorem 1.1** *Assume $\beta = 2$ and let $x \in \mathbb{F}$ and $k \in \mathbb{N}$ be given. If the power $x^{k+1}$ is computed by successive multiplications by $x$, then, in absence of underflow and overflow, the computed approximation $\widehat{r}$ satisfies*

$$\left|\widehat{r} - x^{k+1}\right| \leqslant k\mathbf{u}\left|x^{k+1}\right| \qquad \text{if} \quad k+1 \leqslant \sqrt{2^{1/3}-1}\cdot\mathbf{u}^{-1/2}. \tag{1.3}$$

This improves the classical Wilkinson-type estimate $|\widehat{r} - x^{k+1}| \leqslant \gamma_k|x^{k+1}|$. They also note that for $k \approx \mathbf{u}^{-1}$ the relative error on $\widehat{r}$ can indeed be larger than $k\mathbf{u}$, thus suggesting that in the case of integer powers, the price to be paid for the refined constant $k\mathbf{u}$ is a necessary restriction on the range of $k$. This is in contrast with bounds like (1.2) and the results in [4, 9], where restrictions on $k$ can be avoided.

As Muller [8] mentioned, repeated multiplication may not be the method of choice to evaluate $x^{k+1}$. However, for better methods like binary exponentiation no improvement on the classical constant $\gamma_k$ seems to be known.

In this note we generalize Theorem 1.1 to products of real and/or floating-point numbers, to any base, and to any evaluation scheme using $k$ multiplications. Our restriction on $k$ is weaker than the one in (1.3), though of the same order, and we show that it is essentially sharp.

**Theorem 1.2** *Let $x_0, x_1, \dots, x_k \in \mathbb{R}$ be given and suppose that $\ell$ of them are in $\mathbb{F}$. Let also*

$$K := 2k + 1 - \ell \qquad \text{and} \qquad \omega := \begin{cases} 1 & \text{if } \beta \text{ is odd,} \\ 2 & \text{if } \beta \text{ is even.} \end{cases} \tag{1.4}$$

*Then, any order of evaluation of the product of $\prod_{i=0}^{k} \mathrm{fl}(x_i)$ produces an approximation $\widehat{r}$ such that, in absence of underflow and overflow,*

$$\Big|\widehat{r} - \prod_{i=0}^{k} x_i\Big| \leqslant K\mathbf{u} \Big|\prod_{i=0}^{k} x_i\Big| \qquad \text{if} \quad K < \sqrt{\frac{\omega}{\beta}}\, \mathbf{u}^{-1/2}. \tag{1.5}$$

*In particular, if $\beta = 2$ and all the $x_i$ are in $\mathbb{F}$, then $(K, \omega) = (k, \beta)$ and (1.5) becomes*

$$\Big|\widehat{r} - \prod_{i=0}^{k} x_i\Big| \leqslant k\mathbf{u} \Big|\prod_{i=0}^{k} x_i\Big| \qquad \text{if} \quad k < \mathbf{u}^{-1/2}. \tag{1.6}$$

*For $\beta = 2$ and $p \geqslant 4$, the constraint in (1.6) cannot be replaced by $k < 12\mathbf{u}^{-1/2}$.*

Remark. Note that for $\beta = 2$ and all the $x_i$ in $\mathbb{F}$ the restriction $k < \mathbf{u}^{-1/2}$ improves on the restriction $k + 1 \leqslant \sqrt{2^{1/3} - 1} \cdot \mathbf{u}^{-1/2} = 0.509... \cdot \mathbf{u}^{-1/2}$ in (1.3).

The techniques to prove Theorem 1.2 can be used to obtain similar results for other evaluation schemes. As an example we show how to improve the classical factor $\gamma_{2n}$ for Horner's scheme [2, p. 95].

**Theorem 1.3** *Let $x, a_0, a_1, \dots, a_n \in \mathbb{F}$ be given and let $\widehat{r}$ be the approximation to $\sum_{i=0}^{n} a_i x^i$ produced by Horner's scheme. Then, in absence of underflow and overflow,*

$$\Big|\widehat{r} - \sum_{i=0}^{n} a_i x^i\Big| \leqslant 2n\mathbf{u} \sum_{i=0}^{n} \big|a_i x^i\big| \qquad \text{if} \quad n < \frac{1}{2}\left(\sqrt{\frac{\omega}{\beta}}\, \mathbf{u}^{-1/2} - 1\right)$$

*using $\omega$ defined in* (1.4)*.*

## 2 Products

We need some preliminaries to prove Theorem 1.2. If some $x_i$ is zero, then $\widehat{r} = 0$ because no overflow occurs, and the results in Theorem 1.2 are trivial. If all the $x_i$ are nonzero, then $\widehat{r} \neq 0$ because, by assumption, no underflow occurs. Furthermore, using $\mathbb{F} = -\mathbb{F}$ and $\mathrm{fl}(-x) = -\mathrm{fl}(x)$, we may henceforth assume without loss of generality that all the $x_i$ are positive, so that all the $\widehat{r_i}$ are positive as well.

The standard estimate (1.1) can be improved in two ways. First, it is known that

$$x \in \mathbb{R}: \qquad |\mathrm{fl}(x) - x| \leqslant \frac{\mathbf{u}}{1 + \mathbf{u}}|x| \tag{2.1}$$

and that this bound is sharp; see for example [6, p. 232] and [5]. Second, we use the *unit in the first place* (ufp): a real number $x$ being given, we set $\mathrm{ufp}(0) = 0$ and, if

$x \neq 0$, $\operatorname{ufp}(x) := \beta^{\lfloor \log_\beta |x| \rfloor}$. Thus, $\operatorname{ufp}(x)$ can be thought of as the weight of the first nonzero digit of $x$ in its base-$\beta$ representation. Then,

$$x \in \mathbb{R}: \qquad |\mathrm{fl}(x) - x| \leqslant \mathbf{u} \operatorname{ufp}(x). \tag{2.2}$$

This estimate is sharp as well; for more details, see [10]. Combining (2.1) and (2.2) yields the improved estimate

$$x \in \mathbb{R}\backslash\{0\}: \qquad \mathrm{fl}(x) = x(1+\varepsilon) \quad \text{with} \quad |\varepsilon| \leqslant \min\left[\frac{\mathbf{u}}{1+\mathbf{u}}, \mathbf{u}\frac{\operatorname{ufp}(x)}{|x|}\right]. \tag{2.3}$$

In the following we will use

$$x \in \mathbb{R}\backslash\{0\}: \qquad \operatorname{ufp}(x) \leqslant |x| < \beta \operatorname{ufp}(x), \tag{2.4}$$

as well as

$$\begin{aligned} f \in \mathbb{F} \cap [1, \beta] \qquad &\Rightarrow \qquad f = 1 + 2n\mathbf{u} \quad \text{with } n \in \mathbb{N}_0, \\ f \in \mathbb{F} \cap [\beta^{-1}, 1] \qquad &\Rightarrow \qquad f = 1 - \tfrac{2n}{\beta}\mathbf{u} \quad \text{with } n \in \mathbb{N}_0\,. \end{aligned} \tag{2.5}$$

Some notation is necessary to formalize the computation of the floating-point approximation $\widehat{r}$ in (1.5). The evaluation of $\prod_{i=0}^{k} \mathrm{fl}(x_i)$ in any given order by means of $k$ floating-point multiplications is represented by a binary tree $B$ whose $k+1$ leafs correspond to the $\mathrm{fl}(x_i)$ and whose $k$ inner nodes correspond to the multiplications. Thus, $B$ has $2k+1$ nodes $N_i$ in total.

Since the order of evaluation is arbitrary, we may assume without loss of generality that $x_0, \ldots, x_L \in \mathbb{F}$ with $L := \ell - 1$. The numbering of the nodes shall be such that $N_i$ corresponds to $x_{i+L}$ for $i = -L, \ldots, k-L$, and $N_{k-L+1}, \ldots, N_K$ are the inner nodes. Moreover, $N_K$ shall be the root of $B$.

Each node $N_i$ is the root of a tree $B_i$ and is identified with the floating-point value $\widehat{r_i} = \mathrm{fl}(r_i)$ computed by $B_i$. It follows in particular that $\widehat{r} = \widehat{r_K}$. More precisely, define $r_i := x_{i+L}$ for $i = -L, \ldots, k-L$ and, by means of a recursive definition, if an inner node $N_i$, $i \in \{k-L+1, \ldots, K\}$, has children $N_{i_\nu}$, $1 \leqslant \nu \leqslant 2$, for which $\widehat{r_{i_1}}, \widehat{r_{i_2}}$ are already known, define $r_i := \widehat{r_{i_1}} \cdot \widehat{r_{i_2}}$. Since the $x_i$ and $\widehat{r_i}$ have been assumed to be positive, the same holds for the $r_i$.

By assumption, $\widehat{r_i} = \mathrm{fl}(r_i) = x_{i+L}$ for $i = -L, \ldots, 0$. Moreover, for $i = 1, \ldots, K$ we have

$$\widehat{r_i} = \mathrm{fl}(r_i) =: (1+\varepsilon_i) r_i \qquad \text{with} \quad |\varepsilon_i| \leqslant \min\left[\frac{\mathbf{u}}{1+\mathbf{u}}, \mathbf{u}\frac{\operatorname{ufp}(r_i)}{r_i}\right] < \mathbf{u}. \tag{2.6}$$

For $i \in \{1, \ldots, k-L\}$, the relative errors $\varepsilon_i$ correspond to the rounding of $x_{i+L}$ into $\mathrm{fl}(x_{i+L})$, while for the remaining indices $i \in \{k-L+1, \ldots, K\}$ they correspond to the $k$ multiplications. This implies $\prod_{i=0}^{k} \mathrm{fl}(x_i) = \prod_{i=1}^{k-L} (1+\varepsilon_i) \cdot \prod_{i=0}^{k} x_i$, and therefore

$$\widehat{r_K} - \prod_{i=0}^{k} x_i = \Big( \prod_{i=1}^{K} (1+\varepsilon_i) - 1 \Big) \cdot \prod_{i=0}^{k} x_i. \tag{2.7}$$

Since all factors $x_i$ are positive, (1.5) is equivalent to $\left|\prod_{i=1}^{K}(1+\varepsilon_i)-1\right| \leqslant K\mathbf{u}$, and because $\prod_{i=1}^{K}(1+\varepsilon_i) \geqslant (1-\mathbf{u})^K \geqslant 1-K\mathbf{u}$ it suffices to prove

$$\prod_{i=1}^{K}(1+\varepsilon_i) \leqslant 1+K\mathbf{u}. \tag{2.8}$$

Hence, we need only upper bounds on the $\varepsilon_i$ for the proof of Theorem 1.2.

Furthermore, the lemma below shows that, under weaker assumptions on the maximum $K$, the estimate (1.5) in Theorem 1.2 is true if a single $\varepsilon_i$ is not positive, that is, if any of the $k-L$ real $x_i$ or any single intermediate product is not rounded upwards. A similar observation was already made in [1, Lemma 3].

**Lemma 2.1** *With the notation above, in particular (2.6), assume* $K \leqslant \sqrt{2}\,\mathbf{u}^{-1/2}$. *If there exists an index* $i \in \{1,\dots,K\}$ *with* $\varepsilon_i \leqslant 0$*, then (1.5) holds true.*

*Proof.* By (2.6) and (2.8), it suffices to show $Z := (1+\mathbf{u})^{K-1} \leqslant 1+K\mathbf{u}$. Using $K^2\mathbf{u} \leqslant 2$ gives

$$\ln(Z) = (K-1)\ln(1+\mathbf{u}) \leqslant (K-1)\mathbf{u} \leqslant K\mathbf{u} - \tfrac{1}{2}K^2\mathbf{u}^2 \leqslant \ln(1+K\mathbf{u}). \qquad \square$$

**Proof of Theorem 1.2.** With the notation above, in particular using (2.6), we have to prove (2.8). For $K \in \{0,1\}$ the assertion is trivial so that henceforth we assume $K \geqslant 2$. By Lemma 2.1 we can also assume that

$$\varepsilon_i > 0 \quad \text{for all } i \in \{1,\dots,K\}. \tag{2.9}$$

Let $\varphi \in \mathbb{N}$ be the largest integer satisfying

$$\varphi < \sqrt{\frac{\omega}{\beta}}\,\mathbf{u}^{-1/2}. \tag{2.10}$$

Note that $\varphi \geqslant 2$ because $2 \leqslant K < \sqrt{\omega/\beta}\,\mathbf{u}^{-1/2} \leqslant \varphi+1$. Define $I \subseteq \{1,\dots,K\}$ to be the index set with

$$i \in I \qquad :\Leftrightarrow \qquad \varepsilon_i > \frac{\mathbf{u}}{1+\varphi\mathbf{u}}. \tag{2.11}$$

The following two properties will be proved for distinct $i, j \in I$:

a) The nodes $N_i$ and $N_j$ are not adjacent in the tree $B$. (2.12)

b) The nodes $N_i$ and $N_j$ do not have the same parent node in $B$. (2.13)

*Proof of* (2.12). In order to derive a contradiction suppose that $N_i$ is a child of $N_j$. It follows that $r_j = \widehat{r_i}\,\widehat{q}$, where $\widehat{q} \in \mathbb{F}$ is a (rounded) $x_i$ or some intermediate result. If $\operatorname{ufp}(\widehat{r_i}) = \widehat{r_i}$, then $\widehat{r_i}$ is a power of $\beta$ and $\varepsilon_j = 0$ contradicting (2.9), so that (2.6) and $i \in I$ imply

$$\operatorname{ufp}(r_i) = \operatorname{ufp}(\widehat{r_i}) < r_i < (1+\varphi\mathbf{u})\operatorname{ufp}(r_i) \qquad \text{for} \quad i \in I. \tag{2.14}$$

Since the second inequality is strict and $1+\varphi\mathbf{u} < 1+\sqrt{\mathbf{u}} < \beta$, it follows by (2.5), no matter whether $\varphi$ is odd or even, that

$$\operatorname{ufp}(r_i) = \operatorname{ufp}(\widehat{r_i}) < \widehat{r_i} \leqslant (1+\varphi\mathbf{u})\operatorname{ufp}(r_i) \qquad \text{for} \quad i \in I. \tag{2.15}$$

By (2.15) and (2.5) we have

$$\widehat{r_i} = \mathrm{ufp}(\widehat{r_i})(1 + m\mathbf{u}) \quad \text{for even } m \in \mathbb{N} \text{ with } 2 \leqslant m \leqslant \varphi. \tag{2.16}$$

Hence, $r_j = \widehat{r_i}\,\widehat{q}$, (2.4), (2.16), $j \in I$, and (2.14) imply

$$\frac{R}{1 + m\mathbf{u}} \leqslant \frac{r_j}{\widehat{r_i}} = \widehat{q} \leqslant (1 + \varphi\mathbf{u})R \qquad \text{abbreviating } R := \frac{\mathrm{ufp}(r_j)}{\mathrm{ufp}(\widehat{r_i})}. \tag{2.17}$$

Since $\widehat{q} \in \mathbb{F}$, $R$ is a power of $\beta$, and $R/(1 + m\mathbf{u}) > R(1 - m\mathbf{u}) \in \mathbb{F}$, (2.5) implies that there exists $\nu \in \mathbb{Q}$ such that

$$\widehat{q} = R(1 + \nu\mathbf{u}) \qquad \text{and} \qquad -m < \nu \leqslant \varphi. \tag{2.18}$$

Moreover, if $\nu$ is non-negative, then $\nu$ is a non-negative even integer by (2.5). From (2.18) and (2.16) we get $|\nu| \leqslant \varphi$. Now $r_j = \widehat{r_i}\,\widehat{q}$, (2.18), and (2.16) give

$$\mathrm{ufp}(r_j) \leqslant r_j = \mathrm{ufp}(r_j)(1 + (m + \nu)\mathbf{u} + m\nu\mathbf{u}^2), \tag{2.19}$$

and (2.14) together with $j \in I$ yields

$$0 \leqslant (m + \nu)\mathbf{u} + m\nu\mathbf{u}^2 \leqslant \varphi\mathbf{u}. \tag{2.20}$$

First, assume that $\nu$ is an even integer. Then, $m + \nu > 0$ is also even by (2.16), so that $1 + (m + \nu)\mathbf{u} \in \mathbb{F}$ and $|m\nu\mathbf{u}^2| \leqslant \varphi^2\mathbf{u}^2 < \mathbf{u}$ imply $\widehat{r_j} = \mathrm{ufp}(r_j)(1 + (m + \nu)\mathbf{u})$ and

$$\varepsilon_j = \frac{\widehat{r_j} - r_j}{r_j} = -\frac{\mathrm{ufp}(r_j)m\nu\mathbf{u}^2}{r_j} \leqslant \varphi|\nu|\mathbf{u}^2. \tag{2.21}$$

If $\nu \geqslant 0$, then $\varepsilon_j \leqslant 0$, a contradiction. Otherwise, (2.18) and $-\nu \in \mathbb{N}$ give $|\nu| = -\nu \leqslant m - 1 \leqslant \varphi - 1$, so that $\varphi < \mathbf{u}^{-1/2}$ implies

$$\varphi|\nu|\mathbf{u}^2(1 + \varphi\mathbf{u}) < \frac{1}{\sqrt{\mathbf{u}}}\left(\frac{1}{\sqrt{\mathbf{u}}} - 1\right)\mathbf{u}^2(1 + \sqrt{\mathbf{u}}) = (1 - \sqrt{\mathbf{u}})\mathbf{u}(1 + \sqrt{\mathbf{u}}) \leqslant \mathbf{u}.$$

Hence, $\varepsilon_j < \frac{\mathbf{u}}{1+\varphi\mathbf{u}}$ by (2.21), again a contradiction to $j \in I$ by (2.11).

Second, assume that $\nu$ is not an even integer. Then, (2.18) and (2.5) give $\nu < 0$. Write $\nu = 2n/\beta =: s + r/\beta$ with $n, s, r \in \mathbb{Z}_{\leqslant 0}$ with $|r| := (2|n|) \bmod \beta$. Since $2n$ is even, necessarily

$$|r| \leqslant \begin{cases} \beta - 2 \text{ if } \beta \text{ is even,} \\ \beta - 1 \text{ if } \beta \text{ is odd,} \end{cases} \qquad \Rightarrow \qquad \left|\frac{r}{\beta}\right| \leqslant 1 - \frac{\omega}{\beta} \tag{2.22}$$

using $\omega$ as in (1.4). In particular, for $\beta = 2$ this means $r = 0$. Now, (2.19) becomes

$$r_j = \mathrm{ufp}(r_j)(1 + (m + s + \delta)\mathbf{u}) \quad \text{with} \quad \delta := \frac{r}{\beta} + m\nu\mathbf{u} < 0 \tag{2.23}$$

because $r \leqslant 0$ and $-m \leqslant \nu < 0$. Using (2.22) and (2.10) we obtain

$$|\delta|\mathbf{u} \leqslant \left(1 - \frac{\omega}{\beta} + \varphi^2\mathbf{u}\right)\mathbf{u} < \mathbf{u}. \tag{2.24}$$

If $s$ is odd, then $\delta < 0$ and (2.24) yield $\widehat{r_j} = \mathrm{ufp}(r_j)(1 + (m + s - 1)u)$ and $\varepsilon_j < 0$, a contradiction. If $s$ is even, then $\widehat{r_j} = \mathrm{ufp}(r_j)(1 + (m + s)\mathbf{u})$ and

$$\varepsilon_j = -\delta\mathbf{u}\frac{\mathrm{ufp}(r_j)}{r_j} \leqslant |\delta|\mathbf{u}. \tag{2.25}$$

Note that $s$ even implies $r \neq 0$ as $\nu$ is not an even integer.[1] By (2.18) we have $-m < \nu = s + r/\beta$. Since $m, s$ are even integers and $r/\beta < 0$, it follows $-m + 2 \leqslant s = \nu - r/\beta$, so that (2.22) yields

$$|\nu| = -\nu \leqslant m - 2 - \frac{r}{\beta} \leqslant \varphi - 1 - \frac{\omega}{\beta}. \tag{2.26}$$

From (2.25), (2.23), (2.22), (2.16), (2.26), and (2.10) we deduce the final contradiction to $j \in I$ and (2.11):

$$\frac{\varepsilon_j}{\mathbf{u}} \leqslant |\delta| \leqslant 1 - \frac{\omega}{\beta} + \varphi\left(\varphi - 1 - \frac{\omega}{\beta}\right)\mathbf{u} < 1 - \varphi\mathbf{u} - \frac{\omega}{\beta} + \varphi^2\mathbf{u} < 1 - \varphi\mathbf{u} < \frac{1}{1 + \varphi\mathbf{u}}.$$

This finishes the proof of (2.12).

*Proof of* (2.13). Again, in order to derive a contradiction, assume that $N_i$ and $N_j$ are the left and right children of an inner node $N_a$, $a \in \{k - L + 1, \ldots, K\}$, that is, $r_a = \widehat{r_i}\widehat{r_j}$ and $\widehat{r_a} = \mathrm{fl}(r_a)$. Then, like in the proof of (2.12), $i, j \in I$ implies

$$\begin{aligned}\mathrm{ufp}(r_i) = \mathrm{ufp}(\widehat{r_i}) < r_i < \widehat{r_i} = (1 + m\mathbf{u})\mathrm{ufp}(r_i) \leqslant (1 + \varphi\mathbf{u})\mathrm{ufp}(r_i),\\ \mathrm{ufp}(r_j) = \mathrm{ufp}(\widehat{r_j}) < r_j < \widehat{r_j} = (1 + n\mathbf{u})\mathrm{ufp}(r_j) \leqslant (1 + \varphi\mathbf{u})\mathrm{ufp}(r_j)\end{aligned}$$

with even $m, n \in \mathbb{N}_{\leqslant\varphi}$. Thus,

$$r_a = (1 + (m + n)\mathbf{u} + mn\mathbf{u}^2)\mathrm{ufp}(r_i)\mathrm{ufp}(r_j), \tag{2.27}$$

and $(m + n)\mathbf{u} \leqslant 2\varphi\mathbf{u} < 2\sqrt{\omega/\beta}\,\mathbf{u}^{1/2} \leqslant \frac{2\omega}{K\beta} \leqslant \frac{2}{K} \leqslant 1$ because $K \geqslant 2$. Moreover, $m + n$ is even and $mn\mathbf{u}^2 \leqslant \varphi^2\mathbf{u}^2 < \mathbf{u}$. Thus (2.27) yields $\mathrm{ufp}(r_a) = \mathrm{ufp}(\widehat{r_a}) = \mathrm{ufp}(r_i)\mathrm{ufp}(r_j)$, $\widehat{r_a} = (1 + (m + n)\mathbf{u})\mathrm{ufp}(\widehat{r_a})$, and $\varepsilon_a = -mn\mathbf{u}^2\mathrm{ufp}(\widehat{r_a})/r_a < 0$ contradicting (2.9). This finishes the proof of (2.13).

For $I$ consisting of $k'$ indices, (2.6) and (2.11) give

$$\prod_{i=1}^{K}(1 + \varepsilon_i) \leqslant \left(1 + \frac{\mathbf{u}}{1 + \mathbf{u}}\right)^{k'}\left(1 + \frac{\mathbf{u}}{1 + \varphi\mathbf{u}}\right)^{K-k'}. \tag{2.28}$$

Using (2.12) and (2.13) we will show by Lemma 2.2 in Subsection 2.1 that $k' \leqslant \left\lfloor \frac{K+1}{2} \right\rfloor$. This implies

$$\prod_{i=1}^{K}(1 + \varepsilon_i) \leqslant \left(1 + \frac{\mathbf{u}}{1 + \mathbf{u}}\right)^{\left\lfloor \frac{K+1}{2} \right\rfloor}\left(1 + \frac{\mathbf{u}}{1 + \varphi\mathbf{u}}\right)^{\left\lceil \frac{K-1}{2} \right\rceil}. \tag{2.29}$$

[1] Thus, for the classical case $\beta = 2$ a contradiction to $\{i, j\} \subseteq I$ is already obtained.

Hence, according to (2.8) and using $\frac{\mathbf{u}}{1+\mathbf{u}} \geqslant \frac{\mathbf{u}}{1+\varphi\mathbf{u}}$, the proof is finished if we show

$$F(K) := \left(1 + \frac{\mathbf{u}}{1+\mathbf{u}}\right)^{\frac{K+1}{2}} \left(1 + \frac{\mathbf{u}}{1+\varphi\mathbf{u}}\right)^{\frac{K-1}{2}} \leqslant 1 + K\mathbf{u}. \tag{2.30}$$

For later use, we do this by proving for real $\psi$ the following stronger statement

$$G(\psi) := \left(1 + \frac{\mathbf{u}}{1+\mathbf{u}}\right)^{\frac{\psi+1}{2}} \left(1 + \frac{\mathbf{u}}{1+\psi\mathbf{u}}\right)^{\frac{\psi-3}{2}} \leqslant 1 + (\psi-1)\mathbf{u} \tag{2.31}$$

provided that $1 \leqslant \psi \leqslant \sqrt{\frac{\omega}{\beta}}\,\mathbf{u}^{-1/2}$. If this is true, then for $1 \leqslant K \leqslant \varphi$ we obtain

$$F(K) \leqslant G(K)\left(1 + \frac{\mathbf{u}}{1+\varphi\mathbf{u}}\right) \leqslant (1 + (K-1)\mathbf{u})\left(1 + \frac{\mathbf{u}}{1+(K-1)\mathbf{u}}\right) = 1 + K\mathbf{u}$$

which is (2.30). A computation yields the Taylor expansion

$$G(\psi) = 1 + (\psi-1)\mathbf{u} + \frac{1}{2}G''(\xi)\xi^2 \qquad \text{with} \quad G''(\xi) =: \alpha N(\xi)$$

for some $0 < \xi < \mathbf{u}$ and

$$\alpha := -\frac{(\psi-1)\left(\frac{1+(\psi+1)\xi}{1+\psi\xi}\right)^{\frac{\psi-1}{2}}\left(1+\frac{\xi}{1+\xi}\right)^{\frac{\psi-1}{2}}}{4\,(1+2\xi)\,(1+\xi)^3\,(1+(\psi+1)\xi)^3\,(1+\psi\xi)} < 0.$$

It suffices to show $N(\xi) \geqslant 0$ for $0 < \xi < \mathbf{u}$. Now $N(\xi) = \sum_{\nu=0}^{5} c_\nu$ and $\psi^2\mathbf{u} \leqslant 1$ with

$$\begin{aligned}
c_0 &= 60\,\xi^4 + 160\,\xi^3 + 144\,\xi^2 + 48\,\xi + 4 &&> 4 + 48\,\xi \\
c_1 &= \left(48\,\xi^5 + 192\,\xi^4 + 248\,\xi^3 + 124\,\xi^2 + 20\,\xi\right)\psi &&> 20\,\psi\,\xi + 124\,\psi\,\xi^2 \\
c_2 &= \left(72\,\xi^5 + 187\,\xi^4 + 140\,\xi^3 + 24\,\xi^2 - 4\xi\right)\psi^2 &&> -4 \\
c_3 &= \left(32\,\xi^5 + 41\,\xi^4 - 8\,\xi^2\right)\psi^3 &&> -8\,\psi\,\xi \\
c_4 &= \left(8\,\xi^5 + \xi^4 - 4\,\xi^3\right)\psi^4 &&> -4\,\xi \\
c_5 &= -\,\xi^4\,\psi^5 &&> -\psi\,\xi^2.
\end{aligned}$$

The series expansions were computed by the Symbolic Math Toolbox of MATLAB [7]. It follows $N(\xi) > 0$ for $0 < \xi < \mathbf{u}$, and this proves (1.5) and (1.6).

The assertion on possible constraints of $k$ is deferred to the appendix. This finishes the proof of Theorem 1.2. □

In the proof of Theorem 1.2 we defined $\varphi$ to be the largest integer less than $\sqrt{\omega/\beta}\,\mathbf{u}^{-1/2}$, which reduces to $\varphi < \mathbf{u}^{-1/2}$ for binary arithmetic. Switching from binary arithmetic to another basis requires indeed an adapted definition of $\varphi$. Consider $p := 5$ decimal digits, that is, $\mathbf{u} = 0.5 \cdot 10^{-4}$. Then, $\widehat{r_i} := \mathrm{fl}(1.3033 \cdot 0.7697) = 1.0032$ and $\widehat{q} := 0.99696$ yield $\widehat{r_j} = 1.0002$. Moreover, $\varphi = 63$ whilst the largest integer less than $\mathbf{u}^{-1/2}$ is $\varphi' = 141$. However, both $\varepsilon_i$ and $\varepsilon_j$ would satisfy (2.11) if $\varphi$ was replaced by $\varphi'$, and indices of adjacent nodes would belong to $I$.

## 2.1 A result on colored trees

In (2.29) in the proof of Theorem 1.2 we used the upper bound $\lfloor \frac{K+1}{2} \rfloor$ for the number $k'$ of nodes in the index set $I$. This bound is a consequence of the following lemma.

**Lemma 2.2** *Let $T$ be a tree with $M$ nodes, each having at most two children. Assume that $C$ nodes of $T$ are colored according to the following rules:*

(i) *colored nodes are not adjacent;*
(ii) *each node has at most one colored child.*

*Then,*

$$C \leqslant \begin{cases} \left\lfloor \frac{M+1}{2} \right\rfloor & \text{if the root of } T \text{ is colored,} \\ \left\lfloor \frac{M}{2} \right\rfloor & \text{otherwise.} \end{cases}$$

*Furthermore, these inequalities are sharp for all $M$.*

*Proof* The result is trivial for $M = 1$, so assume $M \geqslant 2$ and that the result is true up to $M - 1$. The root $R$ of $T$ is then connected to a tree $T_1$ and, possibly, also to another tree $T_2$ disjoint from $T_1$. Let $T_1$ have $M_1$ nodes, $C_1$ of which being colored. Define $M_2$ and $C_2$ similarly if $T_2$ exists, and let $C_2 = M_2 = 0$ otherwise. Clearly, $M = M_1 + M_2 + 1$ and $0 \leqslant C_i \leqslant M_i \leqslant M - 1$ for $i = 1, 2$.

If $R$ is colored, then $C = C_1 + C_2 + 1$ and (i) implies the root of $T_1$ is not colored. Hence, by induction, $C_1 \leqslant \lfloor M_1/2 \rfloor \leqslant M_1/2$. Similarly, $C_2 \leqslant M_2/2$, so that

$$C \leqslant \frac{M_1}{2} + \frac{M_2}{2} + 1 = \frac{M+1}{2}.$$

If $R$ is not colored, then $C = C_1 + C_2$ and (ii) implies that $R$ has at most one colored child. Hence, for $M_2$ either zero or nonzero,

$$C \leqslant \frac{M_1}{2} + \frac{M_2}{2} + \frac{1}{2} = \frac{M}{2}.$$

Since $C$ is an integer, the claimed bounds follow for $M \geqslant 2$. Finally, trees with all internal nodes having exactly one child ("linked lists") and whose colored and uncolored nodes alternate show that the bound is attained for any $M$. □

Now, the upper bound for $k'$ in the proof of Theorem 1.2 is obtained as follows. First, we construct a tree $T$ by removing from the binary tree $B$ the leafs $N_{-L}, \ldots, N_0$ associated with the $\ell$ operands $x_i$ already in $\mathbb{F}$. The nodes of $T$ are the nodes $N_1, \ldots, N_K$ of $B$, and the nodes $N_i$ with $i \in I$ are considered as colored. Then, (2.12) and (2.13) imply that $T$ follows the rules (i) and (ii) of Lemma 2.2, so that $|I| = k' \leqslant \left\lfloor \frac{K+1}{2} \right\rfloor$.

Optimality of the bounds in Lemma 2.2 is established by linked lists which represent recursive multiplication of floating-point numbers. We note that optimal bounds are attained for other evaluation schemes as well. Examples for all $M$ for trees with colored root are sketched in Figure 2.1; examples with uncolored root follow similarly.

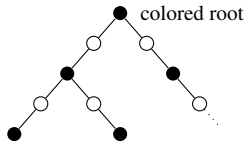


**Fig. 2.1** Trees attaining the bound $C = \left\lfloor \frac{M+1}{2} \right\rfloor$ for colored root.

## 3 Horner scheme

Using the techniques of the previous section we prove Theorem 1.3. For $n = 0$ the assertion is trivial so that we may assume $n \geqslant 1$. The Horner scheme computes

$$\widehat{r_0} := \mathrm{fl}(a_n x); \qquad \widehat{r_i} := \mathrm{fl}(\mathrm{fl}(\widehat{r_{i-1}} + a_{n-i})x), \quad i = 1, \ldots, n-1; \qquad \widehat{r} = \widehat{r_n} := \mathrm{fl}(\widehat{r_{n-1}} + a_0).$$

For $i = 1, \ldots, n$, let the relative error of the $i$-th addition and multiplication be denoted by $\varepsilon_i$ and $\varepsilon'_{i-1}$, respectively. Then,

$$\begin{aligned} \widehat{r_0} &= a_n x(1 + \varepsilon'_0), \\ \widehat{r_i} &= (\widehat{r_{i-1}} + a_{n-i})x(1 + \varepsilon_i)(1 + \varepsilon'_i), \quad i = 1, \ldots, n-1, \\ \widehat{r} &= (\widehat{r_{n-1}} + a_0)(1 + \varepsilon_n). \end{aligned} \tag{3.1}$$

For each $i \in \{1, \ldots, n-1\}$ we apply Theorem 1.2 to the product $x_0 x_1$ with $x_0 := \widehat{r_{i-1}} + a_{n-i} \in \mathbb{R}$ and $x_1 := x \in \mathbb{F}$. Then, $k = 1$, $\ell = 1$ and therefore $K = 2$, so that (2.29) with the constant $\varphi$ defined in (2.10) yields[2]

$$(1 + \varepsilon_i)(1 + \varepsilon'_i) \leqslant \left(1 + \frac{\mathbf{u}}{1 + \mathbf{u}}\right)\left(1 + \frac{\mathbf{u}}{1 + \varphi\mathbf{u}}\right), \quad i = 1, \ldots, n-1. \tag{3.2}$$

Furthermore, (2.6) gives

$$(1 + \varepsilon'_0)(1 + \varepsilon_n) \leqslant \left(1 + \frac{\mathbf{u}}{1 + \mathbf{u}}\right)^2. \tag{3.3}$$

From the equalities in (3.1) we deduce that $\widehat{r} = \sum_{i=0}^{n} a_i(1 + \alpha_i)x^i$, where

$$\begin{aligned} 1 + \alpha_n &= (1 + \varepsilon'_0) \cdot \textstyle\prod_{j=1}^{n-1}(1 + \varepsilon_j)(1 + \varepsilon'_j) \cdot (1 + \varepsilon_n), \\ 1 + \alpha_i &= \textstyle\prod_{j=n-i}^{n-1}(1 + \varepsilon_j)(1 + \varepsilon'_j) \cdot (1 + \varepsilon_n), \quad i = 1, \ldots, n-1, \\ 1 + \alpha_0 &= 1 + \varepsilon_n. \end{aligned}$$

Hence, (1.1), (3.2) and (3.3) imply

$$(1 - \mathbf{u})^{2n} \leqslant 1 + \alpha_n \leqslant \left(1 + \frac{\mathbf{u}}{1 + \mathbf{u}}\right)^{n+1}\left(1 + \frac{\mathbf{u}}{1 + \varphi\mathbf{u}}\right)^{n-1} =: H_n$$

[2] In fact, (2.29) is applied to $|x_0|, |x_1|$ because the proof of Theorem 1.2 assumes positive factors.

and, for $i = 0, \dots, n-1$,

$$(1-\mathbf{u})^{2i+1} \leqslant 1+\alpha_i \leqslant \left(1+\frac{\mathbf{u}}{1+\mathbf{u}}\right)^{i+1}\left(1+\frac{\mathbf{u}}{1+\varphi\mathbf{u}}\right)^{i}.$$

Then, using $1-2n\mathbf{u} < (1-\mathbf{u})^{2n}$, we see that $1-2n\mathbf{u} \leqslant 1+\alpha_i \leqslant H_n$ for all $i = 0, 1, \dots, n$. The assumption $n < \frac{1}{2}\left(\sqrt{\frac{\omega}{\beta}}\,\mathbf{u}^{-1/2} - 1\right)$ implies $2n+1 \leqslant \varphi$. Thus, (2.31) proves $H_n \leqslant G(2n+1) \leqslant 1+2n\mathbf{u}$. □

We close this note with an application of Theorem 1.2.

**Corollary 3.1 (Evaluation of a polynomial given by its roots)**
*Given $z, z_1, \dots, z_n, a_n \in \mathbb{F}$, let $\widehat{r} \in \mathbb{F}$ be a floating-point approximation to*

$$r = a_n \prod_{i=1}^{n} (z - z_i)$$

*obtained by first evaluating the n differences and then, in any order, a product of $n+1$ terms. If $n < \frac{1}{2}\sqrt{\frac{\omega}{\beta}}\,\mathbf{u}^{-1/2}$ then, in absence of underflow and overflow,*

$$|\widehat{r} - r| \leqslant 2n\mathbf{u}|r|.$$

*Proof* Define $x_0 := a_n \in \mathbb{F}$ and $x_i := z - z_i \in \mathbb{R}$ for $i = 1, \dots, n$. Then, Theorem 1.2 with $k = n$, $\ell = 1$, $K = 2k + 1 - \ell = 2n < \sqrt{\frac{\omega}{\beta}}\,\mathbf{u}^{-1/2}$ yields the assertion. □

## 4 Appendix

The goal of this appendix is to prove that for $\beta = 2$ and $p \geqslant 4$ the constraint $k < \mathbf{u}^{-1/2}$ in Theorem 1.2 cannot be replaced by $k < 12\mathbf{u}^{-1/2}$. To do that[3] we construct $x_0, x_1, x_2 \in \mathbb{F}$ for given precision $p$ such that $x_1 x_2 < 1$ and $\mathrm{fl}(\mathrm{fl}(x_0 x_1)x_2) = x_0$. Subsequent multiplications by $x_1 x_2$ produce an exponential growth of the rounding error, eventually exceeding $k\mathbf{u}$.

Define $s := \lfloor \mathbf{u}^{-1/2} \rfloor \in \mathbb{N}$, so that $s = \mathbf{u}^{-1/2} - \delta$ with $0 \leqslant \delta < 1$. We henceforth assume $p \geqslant 15$ and treat the case $p \leqslant 14$ later. Note that $\beta = 2$ and $p \geqslant 15$ imply $s \geqslant 181$. We distinguish two cases.

First, assume $s$ is odd. Set

$$x_0 := 1 + (2s+8)\mathbf{u}, \quad x_1 := 1 - (s-4)\mathbf{u}, \quad \text{and} \quad x_2 := 1 + (s-5)\mathbf{u},$$

so that $x_i \in \mathbb{F}$. Then, $x_0 x_1 = 1 + (s+10)\mathbf{u} + \mu_1\mathbf{u}$ with $\mu_1 := 4\delta\sqrt{\mathbf{u}} + (32 - 2\delta^2)\mathbf{u}$, so that $0 < \mu_1 < 1$ and $s$ odd imply $\mathrm{fl}(x_0 x_1) = 1 + (s+11)\mathbf{u}$. Moreover, $\mathrm{fl}(x_0 x_1)x_2 = 1 + (2s+7)\mathbf{u} + \mu_2\mathbf{u}$ with

$$\mu_2 := \sqrt{\mathbf{u}}\left(6 - 55\sqrt{\mathbf{u}} + \Phi\delta\right) \quad \text{with} \quad \Phi := (\delta - 6)\sqrt{\mathbf{u}} - 2.$$

[3] In [1] long sequences $x_i \in \mathbb{F}$ with $\mathrm{fl}((\dots(\mathrm{fl}(x_0 x_1)x_2)\dots)x_k) = x_0$ are constructed for some precisions.

Now $\Phi < 0$ for any value of $\delta$, so that $0 < 4\sqrt{\mathbf{u}} - 60\mathbf{u} \leqslant \mu_2 \leqslant 6\sqrt{\mathbf{u}} - 55\mathbf{u} < 1$. Thus,

$$\mathrm{fl}(\mathrm{fl}(x_0 x_1)x_2) = x_0. \tag{4.1}$$

Define a vector $X := [x_0 \ x \ x \dots x] \in \mathbb{F}^{2m+1}$ with $m$ times repeating the row vector $x = [x_1 \ x_2] \in \mathbb{F}^2$. Denoting $\widehat{r}_0 := x_0$ and $\widehat{r}_i := \mathrm{fl}(\widehat{r}_{i-1} X_i)$ for $i \geqslant 1$ yields $\widehat{r}_2 = v_0$. Then, abbreviating $\pi := x_1 x_2$ and using $\widehat{r}_{2m} = \widehat{r}_2 = x_0$ gives

$$\widehat{r}_{2m} - \prod_{i=0}^{2m} X_i = x_0 - x_0 \pi^m = (\pi^{-m} - 1) \prod_{i=0}^{2m} X_i \qquad \text{for} \quad 1 \leqslant m \in \mathbb{N}. \tag{4.2}$$

Now,

$$\pi = 1 - \left(2 - (9 + 2\delta)\sqrt{\mathbf{u}}\right)\mathbf{u} - (20 + 9\delta + \delta^2)\mathbf{u}^2 < 1 - \left(2 - 11\sqrt{\mathbf{u}}\right)\mathbf{u} =: 1 - \gamma\mathbf{u},$$

and for $m \in \mathbb{N}$,

$$\pi^{-m} > 1 + m\gamma\mathbf{u} + \frac{m(m-1)}{2}\gamma^2\mathbf{u}^2 = 1 + 2m\mathbf{u} + \frac{m\mathbf{u}\sqrt{\mathbf{u}}}{2}\left[(m-1)\gamma^2\sqrt{\mathbf{u}} - 22\right].$$

The assumption $p \geqslant 15$ implies

$$(6 - 2\sqrt{\mathbf{u}})\gamma^2 - 22 = 2 - 272\sqrt{\mathbf{u}} + (814 - 242\sqrt{\mathbf{u}})\mathbf{u} > 2 - 272\sqrt{\mathbf{u}} > 0,$$

and therefore

$$m \geqslant 6\mathbf{u}^{-1/2} - 1 \quad \Rightarrow \quad \pi^{-m} > 1 + 2m\mathbf{u}. \tag{4.3}$$

Combining this with (4.2) shows that the error bound in (1.6) is not satisfied for $k = 2\left\lceil 6\mathbf{u}^{-1/2} - 1 \right\rceil < 12\mathbf{u}^{-1/2}$, and that finishes the first part.

Second, assume $s$ is even and define as before

$$y_0 := 1 + (2s + 6)\mathbf{u}, \quad y_1 := 1 - (s - 3)\mathbf{u}, \quad \text{and} \quad y_2 := 1 + (s - 4)\mathbf{u}. \tag{4.4}$$

Then, $y_i \in \mathbb{F}$. Furthermore, $y_0 y_1 = 1 + (s + 7)\mathbf{u} + \mu_1\mathbf{u}$ with $\mu_1 := 4\delta\sqrt{\mathbf{u}} + (18 - 2\delta^2)\mathbf{u}$, so that $0 < \mu_1 < 1$ and $s$ even imply $\mathrm{fl}(y_0 y_1) = 1 + (s + 8)\mathbf{u}$. Moreover, $\mathrm{fl}(y_0 y_1)y_2 = 1 + (2s + 5)\mathbf{u} + \mu_2\mathbf{u}$ with

$$\mu_2 := \sqrt{\mathbf{u}}\left(4 - 32\sqrt{\mathbf{u}} + \Phi\delta\right) \quad \text{with} \quad \Phi := (\delta - 4)\sqrt{\mathbf{u}} - 2.$$

As before, $\Phi < 0$ for any value of $\delta$. Thus, $0 < 2\sqrt{\mathbf{u}} - 35\mathbf{u} \leqslant \mu_2 \leqslant 4\sqrt{\mathbf{u}} - 32\mathbf{u} < 1$. Hence, similar to (4.1), $\mathrm{fl}(\mathrm{fl}(y_0 y_1)y_2) = y_0$ is again true. Now for the values $y_1, y_2$ in (4.4) we obtain

$$y_1 y_2 = (1 - (s - 3)\mathbf{u})(1 + (s - 4)\mathbf{u}) < x_1 x_2,$$

and the result follows as before. Finally, for the cases $4 \leqslant p \leqslant 14$, consider

| p | $m_0$ | $m_1$ | $m_2$ | F |
|---|---|---|---|---|
| 4 | 2 | -4 | 4 | 9.6 |
| 5 | 20 | -3 | 2 | 8.9 |
| 6 | 32 | -14 | 16 | 5.8 |
| 7 | 28 | -9 | 8 | 6.8 |
| 8 | 52 | -39 | 44 | 5.8 |
| 9 | 48 | -21 | 20 | 4.6 |
| 10 | 140 | -117 | 130 | 5.2 |
| 11 | 94 | -43 | 42 | 5.8 |
| 12 | 186 | -154 | 158 | 4.0 |
| 13 | 184 | -89 | 88 | 4.1 |
| 14 | 262 | -125 | 124 | 7.2 |

For precision $p$ define $x_i := 1 + m_i\mathbf{u}$. Then, (4.1) is satisfied, and the error bound in (1.6) is not true for $k < F\mathbf{u}^{-1/2}$. This finishes the proof. □

We finally mention that it is easy to see that, if $1 \leqslant p \leqslant 2$, then the error bound in (1.6) is satisfied for all $k \in \mathbb{N}$, and if $p = 3$, then the minimum value of $k$ for which it is not satisfied is $k = 72 \approx 25\mathbf{u}^{-1/2}$.

## 5 Summary

In previous papers, the factor $\gamma_k$ has been replaced by $k\mathbf{u}$ in a number of classical error estimates in numerical analysis together with removing the restriction on $k$. We proved that $k\mathbf{u}$ can be used for general products and for the Horner scheme, however, with a mandatory restriction on $k$. So, as by Theorem 1.2, a general principle to replace $\gamma_k$ by $k\mathbf{u}$ is necessarily restricted to $k \lesssim \mathbf{u}^{-1/2}$.

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