Divide and Conquer Symmetric Tridiagonal Eigensolver for Multicore Architectures

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Conclusion

Scientific Context

PLASMA

- Dense linear algebra library
- Designed for shared-memory systems

Symmetric Eigenproblem

- Solve $AX = \Lambda X$, where A is a symmetric square matrix
- A are the eigenvalues, X the eigenvectors

Eigenproblems in PLASMA

Three Stages Symmetric Eigensolvers

- $A = QTQ^T$: reduction to tridiagonal form: $\Theta(n^3)$
- 2 $T = V \wedge V^T$: tridiagonal eigensolver
- $A = (QV) \wedge (QV)^T$: back-transformation: $\Theta(n^3)$

Q, V orthogonal; T tridiagonal

Objectives

- Introduce a task-based tridiagonal eigensolver
- Towards a fully task-based eigensolver in PLASMA

Outline

Overview in LAPACK

- **2** Algorithm
- Implementation
- 4 Results

Tridiagonal Eigenproblems in LAPACK

Algorithm	Operational Cost	Workspace	Accuracy
QR	$\Theta(n^3)$	$\Theta(n)$	$\Theta(\sqrt{n}\epsilon)$
BI	$\Theta(n^3)$	$\Theta(n)$	$\Theta(n\epsilon)$
DC	$\Theta(n^2)$ to $\Theta(n^3)$	$\Theta(n^2)$	$\Theta(\sqrt{n}\epsilon)$
MRRR	$\Theta(n^2)$	$\Theta(n)$	$\Theta(n\epsilon)$

Computing all eigenpairs

Particular Cases

- Computing all eigenvalues: $\Theta(n^2)$ using QR-PWK
- Computing k eigenpairs: Θ(nk) using MRRR

Computing All Eigenpairs

MRRR vs D&C

- Depend on the matrix properties: "clustered" eigenvalues
- The average performance on practical matrices for a sequential run is comparable (*Performance and Accuracy of LAPACK's Symmetric Tridiagonal Eigensolvers* by Demmel et al.)
 - D&C: Θ(n^{2.5})
 - MRRR: $\Theta(n^{2.2})$

Choice

- D&C was not implemented in a task-based fashion
- Depend on the application: Cost vs Accuracy

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Divide & Conquer – Divide

Strategy

- Divide the problem into two halves subproblems
- Recursively solve each subproblem
- Merge those subproblems according to the approximation

Divide with a rank-1 approximation:

$$\begin{pmatrix} \ast & \ast & \varnothing & \emptyset & \emptyset \\ \ast & \ast & & \beta & \emptyset \\ \emptyset & \beta & & \ast & \ast \\ \emptyset & \emptyset & & \ast & \star \end{pmatrix} = \begin{pmatrix} \ast & \ast & & \emptyset & \emptyset & \emptyset \\ \ast & \ast^{*} & & \emptyset & \emptyset & \emptyset \\ \emptyset & \emptyset & & \ast^{*} & \star & \star \end{pmatrix} + \begin{pmatrix} \emptyset & \emptyset & \emptyset & \emptyset & \emptyset \\ \emptyset & \beta & \beta & \emptyset \\ \emptyset & \emptyset & \beta & \beta & \emptyset \\ \emptyset & \emptyset & \emptyset & \emptyset & \emptyset \end{pmatrix}$$

Divide & Conquer – Conquer



Since *T* has been decomposed as:

$$T = \begin{pmatrix} V_1 D_1 V_1^T & \emptyset \\ \emptyset & V_2 D_2 V_2^T \end{pmatrix} + \beta u u^T$$

We have:

$$T = \tilde{V}(D + \beta z z^T) \tilde{V}^T$$

The merging equation is:

$$(D + \beta z z^T) = X \wedge X^T$$

Final *T* eigenvectors are:

$$V = \tilde{V} \times X$$

Relations

$$D = diag(D_1, D_2)$$
$$V = diag(V_1, V_2)$$
$$z = \tilde{V}^T u$$

Deflation Process

The deflation occurs in $(D + \beta z z^T)$ when

- An eigenvalue has a multiplicity greater than one
- An entry in z is null

For those eigenvalues, the solution provided by the subproblems is also the solution of the current problem

Merging Phase

- Find the deflated eigenvalues
- Solve the merging equation for non-deflated eigenvalues
- Output the previous system for non-deflated eigenvalues

Merge Step - Deflate



- 1: Deflation due to equal eigenvalues from D_1 and D_2
- 2: Deflation due to zero entries in z

Merge Step - Compute Eigenvectors

Operations on Non-Deflated Eigenvectors

- Solve the secular equation (rely on LAPACK)
- Reduction to compute the stabilization vector W
- Stabilize eigenvectors according to $W \rightarrow V_{new}$



Merge Step - Update Eigenvectors



Performance

- *V_{up}*: computational bound
- *V_{def}*: memory bound

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Related Work

MKL LAPACK

- LAPACK algorithm linked with optimised BLAS 3 kernels
- Only GEMMs in parallel
- Interesting speedup: GEMMs represent 80% of the overall time

MKL ScaLAPACK

- Independent subproblems can be solved in parallel
- GEMMs in parallel, as well as solving secular equations
- Permutations are different with LAPACK version
- Main limit: static scheduling

PLASMA Programming Model

Using a Direct Acyclic Graph (DAG)

- Algorithm expressed as a flow of sequential tasks
- Rely on the QUARK dynamic runtime
- Operations on vectors panels (1D)

Dependencies

- INPUT: read
- INOUT: read/write
- OUTPUT: write

Direct Acyclic Graph (1)



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Direct Acyclic Graph (2)



Properties

- Generic DAG
- Constant number of dependencies

Without Deflation - Parallel GEMMs

Traces on a *n*-by-*n* matrix of type 4 with $n = 10\ 000$ on a 16 cores machine



Equivalent to MKL LAPACK (4.3s)

UpdateVect	ComputeVect
LAED4	ComputeLocalW



Operations within a merge in parallel (1.8s)



Independent subproblems in parallel (1.6s)



Operations within a merge in parallel (1.8s)



Independent subproblems in parallel (1.6s)



Operations within a merge in parallel (1.8s)



Independent subproblems in parallel (1.6s)



Operations within a merge in parallel (1.8s)

THEFT A



Independent subproblems in parallel (1.6s)

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With Deflation

Trace on a *n*-by-*n* matrix of type 2 with $n = 10\ 000$ on a 16 cores machine (0.4s)



Permute V	CopyBackDeflated
LASET	SortEigenvectors

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- **B** Implementation



Validation on a Large Set of Matrices

Туре	Description
1	$\lambda_1 = 1, \lambda_{i,i\in[2,n]} = \frac{1}{k}$
2	$\lambda_{i,i\in[1,n-1]}=1,\lambda_n=\frac{1}{k}$
3	$\lambda_{i,i\in[1,n]}=k^{-\frac{i-1}{n-1}}$
4	$\lambda_{i,i\in[1,n]} = 1 - (\frac{i-1}{n-1})(1 - \frac{1}{k})$
5	n random numbers with $log(i)$ uniformly distributed
6	n random numbers
7	$\lambda_{i,i\in[1,n-1]} = ulp imes i, \lambda_n = 1$
8	$\lambda_1 = \textit{ulp}, \lambda_{i,i \in [2,n-1]} = 1 + i imes \sqrt{\textit{ulp}}, \lambda_n = 2$
9	$\lambda_1 = 1, \lambda_{i,i \in [2,n]} = \lambda_{i-1} + 100 imes ulp$
10	(1,2,1) tridiagonal matrix
11	Wilkinson matrix
12	Clement matrix
13	Legendre matrix
14	Laguerre matrix
15	Hermite matrix

Results

Conclusior

Scalability with 16 cores and $n = 10\ 000$



Different levels of deflation

Relative speedup vs Existing Solutions

With 16 cores



Relative speedup vs MR3-SMP

With 16 cores





Results

Accuracy vs MRRR

Stands for classical LAPACK implementations



Conclusion

Results

- Outperforms existing D&C implementations
- Competitive with respect to MR³-SMP
- Demonstrates the asset of a task-based approach

Future Work

- Solve the SVD problem using the same approach
- Expand to distributed/heterogeneous machines
- Gather with the reduction to tridiagonal form

Thank you!

Available in PLASMA 2.7.1 http://icl.cs.utk.edu/plasma/