

Complexity in control-affine systems

Frédéric Jean, Dario Prandi

To cite this version:

Frédéric Jean, Dario Prandi. Complexity in control-affine systems. NETCO 2014, 2014, Tours, France. hal-01024628

HAL Id: hal-01024628 <https://inria.hal.science/hal-01024628>

Submitted on 17 Jul 2014

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Complexity in control-affine systems

Frédéric Jean¹ Dario Prandi²

¹ENSTA ParisTech (and Team GECO, INRIA Saclay)

> ²LSIS University of Toulon

NetCo 2014 – Tours, June 23th, 2014

Framework

We will consider affine-control systems, i.e., systems in the form

$$
\dot{q}(t) = f_0(q(t)) + \sum_{i=1}^m u_i(t) f_i(q(t))
$$

Here,

n the point q belongs to a smooth manifold M the f_i 's are smooth vector fields on M $u \in L^1([0, T], \mathbb{R}^m)$

This type of system appears in many applications

- Mechanical systems
- **Quantum control**
- Microswimmers (Tucsnak, Alouges)
- Neuro-geometry of vision (Mumfor, Petitot)

[Motion planning problem](#page-4-0)

[Definitions of complexity](#page-10-0)

[Asymptotic estimates in affine-control systems](#page-14-0)

1 [Motion planning problem](#page-4-0)

Problem

Given $x, y \in M$, find an admissible trajectory steering the system from x to y , possibly under some constraints.

Possible constraints:

- **1** Avoiding some obstacles
- Rendez-vous problem, i.e., being near certain places at certain times

Assumption

A metric with balls $B(q, \varepsilon)$ is fixed on M.

Different approaches are possible. We consider the following method:

Different approaches are possible. We consider the following method:

1 Find an (non-admissible) curve $\Gamma \subset M$ or a path $\gamma : [0, T] \rightarrow M$ solving the problem.

Different approaches are possible. We consider the following method:

- 1 Find an (non-admissible) curve $\Gamma \subset M$ or a path $\gamma : [0, T] \rightarrow M$ solving the problem.
- 2 Track Γ or γ with an admissible trajectory.

Different approaches are possible. We consider the following method:

- **1** Find an (non-admissible) curve $\Gamma \subset M$ or a path $\gamma : [0, T] \to M$ solving the problem. \longrightarrow global topology
- 2 Track Γ or γ with an admissible trajectory. \rightarrow local behavior of the control system

We focus on quantifying the difficulty of the second step.

2 [Definitions of complexity](#page-10-0)

Let $J : U \rightarrow [0, +\infty)$ be a cost function.

Definition (Complexity)

A measure of the cost of approximation of a given curve/path with a certain precision

In general:

1 we fix a set $Adm(\Gamma, \varepsilon)$ of admissible controls for precision ε 2 we define complexity as

$$
\sigma(\gamma,\varepsilon)=\inf_{u\in \mathsf{Adm}(\Gamma,\varepsilon)}\frac{\mathrm{cost}\; \mathrm{of}\; u}{\mathrm{cost}\; \mathrm{of}\; \mathrm{an}\; \mathrm{``}\varepsilon\; \mathrm{piece''}\; \mathrm{of}\; u}=\frac{1}{\varepsilon}\inf_{u\in \mathsf{Adm}(\Gamma,\varepsilon)}J(u,\,T).
$$

Obstacle-avoidance problem

Let
$$
\Gamma \subset M
$$
 be a curve, $\text{Tube}(\Gamma, \varepsilon) = \bigcup_{q \in \gamma} B(q, \varepsilon)$, and
\n
$$
\mathcal{A}(\Gamma, \varepsilon) = \left\{ u \in L^1([0, T], \mathbb{R}^m) \mid \begin{array}{l} T > 0, q_u(T) = y, \\ q_u(\cdot) \subset \text{Tube}(\Gamma, \varepsilon) \end{array} \right\}.
$$

With this set we define the tubular approximation complexity

$$
\Sigma_a(\Gamma,\varepsilon)=\frac{1}{\varepsilon}\inf_{u\in\mathcal{A}(\Gamma,\varepsilon)}J(u,\,\mathcal{T}).
$$

Rendez-vous problem

Let
$$
\gamma : [0, T] \to M
$$
 be a path and
\n
$$
\mathcal{N}(\gamma, \varepsilon) = \left\{ u \in L^{1}([0, T], \mathbb{R}^{m})) \mid \begin{array}{l} q_{u}(T) = y \text{ and } q_{u}(t) \in B(\gamma(t), \varepsilon) \\ \text{for any } t \in [0, T] \end{array} \right\}.
$$

This set defines the neighborhood approximation complexity

$$
\sigma_n(\gamma,\varepsilon)=\frac{1}{\varepsilon}\inf_{u\in\mathcal{N}(\gamma,\varepsilon)}J(u,\,\mathcal{T}).
$$

3 [Asymptotic estimates in affine-control systems](#page-14-0)

Particular case: nonholonomic control systems

Nonholonomic control system $=$ control-affine system without drift

$$
\dot{q}(t)=\sum_{i=1}^m u_i(t)f_i(q(t)),
$$

that satisfies the Hörmander condition, i.e., such that

$$
\mathsf{Lie}_q\{f_1,\ldots,f_m\} = \mathcal{T}_qM, \quad \text{for any } q \in M.
$$

- \blacksquare The value function associated to this system w.r.t. the L^1 cost is a distance, called sub-Riemannian distance.
- 2 Due to the linearity of the system, we can always reparametrize trajectories without changing their L^1 cost. Hence,

Tubular approximation complexity Neighborhood approximation complexity

Introduced by Gromov (1996) in a different context.

Weak equivalence:

$$
\sigma(\Gamma, \varepsilon) \asymp g(\varepsilon) \iff \mathcal{C}_1 \leq \frac{\sigma(\Gamma, \varepsilon)}{g(\varepsilon)} \leq \mathcal{C}_2 \quad \text{ for } \varepsilon \downarrow 0.
$$

Complete results (Jean 2003).

Strong equivalence:

$$
\sigma(\Gamma,\varepsilon) \simeq g(\varepsilon) \iff \lim_{\varepsilon \downarrow 0} \frac{\sigma(\Gamma,\varepsilon)}{g(\varepsilon)} = 1.
$$

Results in particular cases (Gauthier, Zakalyukin, et al., 2004-2013)

General case

Recall the general form of a control-affine system

$$
\dot{q}(t) = f_0(q(t)) + \sum_{i=1}^m u_i(t) f_i(q(t)).
$$

We will consider:

- **strong Hörmander condition:** Lie_g $\{f_1, \ldots, f_m\} = T_qM$ for any $q \in M$.
- The set of controls is

$$
\mathcal{U} = \bigcup_{\mathcal{T} \in (0,\mathcal{T}]} \mathsf{L}^1([0,\mathcal{T}],\mathbb{R}^m).
$$

The cost J is the L^1 -norm of u.

Consequences:

- **1** Small time local controllability.
- The associated driftless system ($f_0 = 0$) is a nonholonomic system.

Complexities for control-affine systems

- We will use the sub-Riemannian metric to define the complexities.
- Since the system is not linear, we cannot reparametrize the trajectories, and hence

Tubular approximation \limsup approximation $\lim_{x\to\infty}$ Neighborhood approximation complexity complexity

For any $q \in M$, $s \in \mathbb{N}$, let

$$
\Delta^s(q) = \text{span}\{[f_{i_1}, [f_{i_2}, [\ldots, f_{i_k}] \ldots]](q) \mid 1 \leq k \leq s, 1 \leq i_j \leq m\}.
$$

$$
\Delta^1(q) \subset \Delta^2(q) \subset \ldots \subset \Delta^r(q) = T_q M
$$

Hypothesis

Equiregularity: for any $s \in \mathbb{N}$, dim Δ^s does not depend on the point $q \in M$.

Theorem

Let $f_0 \subset \Delta^s \setminus \Delta^{s-1}$.

Let $\Gamma \subset M$ be a smooth curve. Let k such that $T\Gamma \subset \Delta^k$ and TΓ $\not\subset \Delta^{k-1}$. Then, if T is sufficiently small, we have

$$
\Sigma_a(\Gamma,\varepsilon)\asymp \frac{1}{\varepsilon^k}
$$

Let $\gamma:[0,\, \mathcal T]\to M$ be a path and k such that $\dot\gamma\in \Delta^k$ and $\dot\gamma\notin\Delta^{k-1}.$ If, moreover, $\mathsf{s}=k$, we assume that $\dot\gamma\neq\mathit{f}_0(\gamma)$ mod $\Delta^{s-1}(\gamma)$. Then

$$
\sigma_{n}(\gamma,\varepsilon)\asymp \frac{1}{\varepsilon^{\max\{s,k\}}}
$$

- \blacksquare The complexity of curves is not sensible to the drift.
- \blacksquare The complexity of **paths** depends on the drift. In particular, when $f_0 \subset \Delta^r \setminus \Delta^{r-1}$ where r is such that $\Delta^r = T_a M$, the complexity is always maximal, i.e., $\sigma_n(\gamma, \varepsilon) \asymp \varepsilon^{-r}$.

Techniques and Remarks

Neak estimates of the value function near a point (generalization of the sub-Riemannian Ball-Box theorem).

Example

- f_1 and f_2 control vector fields on \mathbb{R}^3 satisfying the Hörmander condition,
- Drift s.t. $f_0 \not\subset \Delta^1 = \text{span}\{f_1, f_2\}.$

Nonholonomic system.

Control-affine system.

Estimates obtained by reducing the control system with drift to a driftless but time-dependent system.

$$
\dot{q} = f_0(q) + \sum_{i=1}^m u_i f_i(q) \longrightarrow \dot{q} = \sum_{i=1}^m u_i (e^{-tf_0})_* f_i(q).
$$

For this system we can define a generalization of the nilpotent m. approximation, that yields the estimates.

Final remarks

■ We studied also two other notions of complexity, where we track the curve/path by interpolation, and no metric is assumed.

We studied also another cost

$$
\mathcal{I}(u,\mathcal{T})=\int_0^{\mathcal{T}}\sqrt{1+\sum_{i=1}^m u_i(t)^2}\,dt.
$$

Thank you for your attention.

Ball-Box

Let $\{\partial_{z_i}\}_{i=1}^n$ be the canonical basis of \mathbb{R}^n and $\mathcal{R}_{f_0}(q,\varepsilon)$ the reachable set from q with cost $\leq \varepsilon$. We define

$$
\Xi(\eta) = \bigcup_{0 \leq \xi \leq \mathcal{T}} (\xi \partial_{z_{\ell}} + \text{Box}(\eta))
$$

$$
\Pi(\eta) = \bigcup_{0 \leq \xi \leq \mathcal{T}} \{z \in \mathbb{R}^n : |z_{\ell} - \xi| \leq \eta^s, |z_i| \leq \eta^{w_i} + \eta \xi^{\frac{w_i}{s}} \text{ pour } w_i \leq s, i \neq k,
$$

$$
\text{ et } |z_i| \leq \eta(\eta + \xi^{\frac{1}{s}})^{w_i - 1} \text{ pour } w_i > s \},
$$

Theorem

Let $z = (z_1, \ldots, z_n)$ a privileged coordinate system at q for $\{f_1, \ldots, f_m\}$, rectifying f_0 as the k-th coordinate vector field $\partial_{\mathsf{z}_\ell}$, for some $1\leq \ell\leq n$. Them, there exist C, ε_0, T_0 s.t., if $T < T_0$, it holds

$$
\Xi\left(\frac{1}{C}\varepsilon\right)\subset\mathcal{R}_{f_0}(q,\varepsilon)\subset\Pi(C\varepsilon),\qquad\text{for }\varepsilon<\varepsilon_0.
$$