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► **To cite this version:**

Manfred Kufleitner, Alexander Lauser. The Join of the Varieties of R-trivial and L-trivial Monoids via Combinatorics on Words. *Discrete Mathematics and Theoretical Computer Science*, 2012, Vol. 14 no. 1 (1), pp.141-146. 10.46298/dmtcs.564 . hal-00992870

HAL Id: hal-00992870

<https://inria.hal.science/hal-00992870>

Submitted on 19 May 2014

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The Join of the Varieties of \mathcal{R} -trivial and \mathcal{L} -trivial Monoids via Combinatorics on Words

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received 10th January 2012, revised 25th May 2012, accepted 26th May 2012.

The join of two varieties is the smallest variety containing both. In finite semigroup theory, the varieties of \mathcal{R} -trivial and \mathcal{L} -trivial monoids are two of the most prominent classes of finite monoids. Their join is known to be decidable due to a result of Almeida and Azevedo. In this paper, we give a new proof for Almeida and Azevedo's effective characterization of the join of \mathcal{R} -trivial and \mathcal{L} -trivial monoids. This characterization is a single identity of ω -terms using three variables.

Keywords: finite semigroup theory, join of pseudovarieties, Green's relations, combinatorics on words

1 Introduction

Green's relations \mathcal{R} and \mathcal{L} are a standard tool in the study of semigroups [5]. In the context of finite monoids, among other results, they have been used to give effective characterizations of language classes such as star-free languages [3, 11] and piecewise testable languages [6, 12]. A deterministic extension of piecewise testable languages yields the class of languages corresponding to \mathcal{R} -trivial monoids, and a codeterministic extension corresponds to \mathcal{L} -trivial monoids [4, 9].

Almeida and Azevedo gave an effective characterization for the least variety of finite monoids containing all \mathcal{R} -trivial and all \mathcal{L} -trivial monoids [2], *i.e.*, for the *join* of the two varieties. Their proof is based on sophisticated algebraic techniques, on Reiterman's Theorem [10], and on a combinatorial result of König [7]. In this paper, we give a new proof of Almeida and Azevedo's Theorem. The current proof was inspired by another proof of the authors [8], which in turn uses ideas of Klíma [6]. The main ingredient is a system of congruences which relies on simple combinatorics on words.

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2 Preliminaries

Let A be a finite alphabet. The set of finite words over A is denoted by A^* . It is the free monoid over A . The *empty word* is 1. The *content* of a word $u = a_1 \cdots a_n$ with $a_i \in A$ is $\alpha(u) = \{a_1, \dots, a_n\}$, and its *length* is $|u| = n$. The length of the empty word is 0. A word u is a *prefix* (respectively *suffix*) of v if there exists $x \in A^*$ such that $ux = v$ (respectively $xu = v$); if $x \neq 1$, then u is a *proper prefix*.

For more details concerning the algebraic concepts introduced in the remainder of this section, we refer the reader to textbooks such as [1, 4, 9]. *Green's relations* \mathcal{R} and \mathcal{L} are important tools in the study of finite monoids. Let M be a finite monoid. We set $u \mathcal{R} v$ for $u, v \in M$ if $uM = vM$, and the latter condition is equivalent to the existence of $x, y \in M$ with $u = vx$ and $v = uy$. Symmetrically, $u \mathcal{L} v$ if $Mu = Mv$. The monoid M is \mathcal{R} -trivial (respectively \mathcal{L} -trivial) if \mathcal{R} (respectively \mathcal{L}) is the identity relation on M . We write $u <_{\mathcal{R}} v$ if $uM \subsetneq vM$, and we write $u <_{\mathcal{L}} v$ if $Mu \subsetneq Mv$.

A *variety* of finite monoids is a class of monoids closed under finite direct products, submonoids, and quotients. A variety of finite monoids is often called a *pseudovariety* in order to distinguish from varieties in Birkhoff's sense. Since we do not need this distinction in the current paper, whenever we use the term *variety* we mean a variety of finite monoids. The *join* $\mathbf{V}_1 \vee \mathbf{V}_2$ of two varieties \mathbf{V}_1 and \mathbf{V}_2 is the smallest variety containing $\mathbf{V}_1 \cup \mathbf{V}_2$. A monoid M is in $\mathbf{V}_1 \vee \mathbf{V}_2$ if and only if there exist $M_1 \in \mathbf{V}_1$ and $M_2 \in \mathbf{V}_2$ such that M is a quotient of a submonoid of $M_1 \times M_2$. If M is a finite monoid, then there exists an integer $\omega_M \geq 1$ such that, for all $u \in M$, the element u^{ω_M} is idempotent. Moreover, the element u^{ω_M} is the unique idempotent generated by u . Usually, the monoid M is clear from the context and thus, we simply write ω instead of ω_M . This leads to the following definition. An ω -term over a finite alphabet X is either a word in X^* , or of the form t^ω for some ω -term t , or the concatenation $t_1 t_2$ of two ω -terms t_1, t_2 . A homomorphism $\varphi : X^* \rightarrow M$ to a finite monoid M uniquely extends to ω -terms over X by setting $\varphi(t^\omega) = \varphi(t)^{\omega_M}$. Let u, v be two ω -terms over X . A finite monoid M *satisfies* the identity $u = v$ if $\varphi(u) = \varphi(v)$ for all homomorphisms $\varphi : X^* \rightarrow M$. The class of finite monoids satisfying the identity $u = v$ is denoted by $\llbracket u = v \rrbracket$. For all ω -terms u, v , the class $\llbracket u = v \rrbracket$ forms a variety. We need the following three varieties in this paper:

$$\begin{aligned} \mathbf{R} &= \llbracket (xy)^\omega x = (xy)^\omega \rrbracket, \\ \mathbf{L} &= \llbracket x(zx)^\omega = (zx)^\omega \rrbracket, \\ \mathbf{W} &= \llbracket (xy)^\omega x(zx)^\omega = (xy)^\omega (zx)^\omega \rrbracket. \end{aligned}$$

A monoid is in \mathbf{R} if and only if it is \mathcal{R} -trivial. Symmetrically, a monoid is in \mathbf{L} if and only if it is \mathcal{L} -trivial. The aim of this paper is to give a new proof of Almeida and Azevedo's result $\mathbf{R} \vee \mathbf{L} = \mathbf{W}$. The inclusion $\mathbf{R} \vee \mathbf{L} \subseteq \mathbf{W}$ is trivial since $\mathbf{R} \cup \mathbf{L} \subseteq \mathbf{W}$ and \mathbf{W} is a variety.

3 Congruences

In this section, we introduce the main combinatorial tool for our proof. It is a family of congruences \equiv_n on A^* for some finite alphabet A such that $A^*/\equiv_n \in \mathbf{R} \vee \mathbf{L}$ for all integers $n \geq 0$, see Lemma 2 below. As a first step towards the definition of \equiv_n we need to introduce an asymmetric, weaker congruence $\equiv_n^{\mathcal{R}}$. Let $u, v \in A^*$. We let $u \equiv_0^{\mathcal{R}} v$ if $\alpha(u) = \alpha(v)$. For $n \geq 0$, we let $u \equiv_{n+1}^{\mathcal{R}} v$ if the following conditions hold:

1. $\alpha(u) = \alpha(v)$,
2. for all factorizations $u = u_1 a u_2$ and $v = v_1 a v_2$ with $a \in A \setminus (\alpha(u_1) \cup \alpha(v_1))$ we have $u_1 \equiv_n^{\mathcal{R}} v_1$ and $u_2 \equiv_n^{\mathcal{R}} v_2$, and

3. for all factorizations $u = u_1 a u_2$ and $v = v_1 a v_2$ with $a \in A \setminus (\alpha(u_2) \cup \alpha(v_2))$ we have $u_1 \equiv_n^{\mathcal{R}} v_1$.

By a straightforward verification we see that $\equiv_n^{\mathcal{R}}$ is an equivalence relation. The factorization $u_1 a u_2$ with $a \in A \setminus \alpha(u_1)$ is unique. Therefore, induction on n shows that the index of $\equiv_n^{\mathcal{R}}$ is finite. If $u \equiv_{n+1}^{\mathcal{R}} v$, then $u \equiv_n^{\mathcal{R}} v$. Moreover, if $u \equiv_n^{\mathcal{R}} v$ and $a \in A$, then $au \equiv_n^{\mathcal{R}} av$ and $ua \equiv_n^{\mathcal{R}} va$. Therefore, the relation $\equiv_n^{\mathcal{R}}$ is a finite index congruence on A^* .

Lemma 1 For every finite alphabet A and every integer $n \geq 0$ we have $A^* / \equiv_n^{\mathcal{R}} \in \mathbf{R}$.

Proof: It suffices to show $(xy)^{n+1} x \equiv_n^{\mathcal{R}} (xy)^{n+1}$ for all words $x, y \in A^*$. We note that for $y = 1$ this yields $x^{n+2} \equiv_n^{\mathcal{R}} x^{n+1}$. The proof is by induction on n . For $n = 0$, the claim is true since $\alpha(xy) = \alpha(x)$. Let now $n > 0$. As before, $\alpha((xy)^{n+1} x) = \alpha((xy)^{n+1})$. Suppose $(xy)^{n+1} x = u_1 a u_2$ and $(xy)^{n+1} = v_1 a v_2$ for $a \in A \setminus (\alpha(u_1) \cup \alpha(v_1))$. Then $u_1 = v_1$ and both are proper prefixes of xy . Thus $u_2 = p(xy)^n x$ and $v_2 = p(xy)^n$ for some $p \in A^*$. By induction $(xy)^n x \equiv_{n-1}^{\mathcal{R}} (xy)^n$ and hence, $u_2 \equiv_n^{\mathcal{R}} v_2$.

Suppose now $(xy)^{n+1} x = u_1 a u_2$ and $(xy)^{n+1} = v_1 a v_2$ for $a \in A \setminus (\alpha(u_2) \cup \alpha(v_2))$. Then av_2 is a suffix of xy and au_2 is a suffix of yx . We can therefore write $v_1 = (xy)^n p'$ for some prefix p' of xy . Similarly, $u_1 = (xy)^k p$ for some $k \in \{n, n+1\}$ and some prefix p of xy , i.e., we have $pq = xy$ for some $q \in A^*$. By induction, we have $(xy)^{n+1} \equiv_{n-1}^{\mathcal{R}} (xy)^n$ and thus $(xy)^{n+1} p \equiv_{n-1}^{\mathcal{R}} (xy)^n p$. We can therefore assume $k = n$. Without loss of generality, let $|p| \leq |p'|$, i.e., $p' = ps$ for some $s \in A^*$. It follows

$$u_1 = (pq)^n p \quad \text{and} \quad v_1 = (pq)^n ps.$$

Since $p' = ps$ is a prefix of $xy = pq$, the word s is a prefix of q . In particular, there exists $t \in A^*$ such that $qp = st$. This yields

$$u_1 = p(st)^n \quad \text{and} \quad v_1 = p(st)^n s.$$

By induction, $(st)^n \equiv_{n-1}^{\mathcal{R}} (st)^n s$ and thus $u_1 \equiv_{n-1}^{\mathcal{R}} v_1$. This shows $(xy)^{n+1} x \equiv_n^{\mathcal{R}} (xy)^{n+1}$ which concludes the proof. \square

There is a left-right symmetric congruence $\equiv_n^{\mathcal{L}}$ on A^* . It can be defined by setting $u \equiv_n^{\mathcal{L}} v$ if and only if $u^\rho \equiv_n^{\mathcal{R}} v^\rho$. Here, $u^\rho = a_n \cdots a_1$ is the reversal of the word $u = a_1 \cdots a_n$ with $a_i \in A$. It satisfies $A^* / \equiv_n^{\mathcal{L}} \in \mathbf{L}$ for every $n \geq 0$. We define $u \equiv_n v$ if and only if both $u \equiv_n^{\mathcal{R}} v$ and $u \equiv_n^{\mathcal{L}} v$. The following lemma puts together some properties of the finite index congruence \equiv_n .

Lemma 2 For every finite alphabet A and every integer $n \geq 0$ the following properties hold:

1. $A^* / \equiv_n \in \mathbf{R} \vee \mathbf{L}$.
2. If $u_1 a u_2 \equiv_{n+1} v_1 a v_2$ for $a \in A \setminus (\alpha(u_1) \cup \alpha(v_1))$, then $u_1 \equiv_n^{\mathcal{R}} v_1$ and $u_2 \equiv_n v_2$.
3. If $u_1 a u_2 \equiv_{n+1} v_1 a v_2$ for $a \in A \setminus (\alpha(u_2) \cup \alpha(v_2))$, then $u_1 \equiv_n v_1$ and $u_2 \equiv_n^{\mathcal{L}} v_2$.

Proof: “1”: We have $A^* / \equiv_n \in \mathbf{R} \vee \mathbf{L}$ since it is a submonoid of $(A^* / \equiv_n^{\mathcal{R}}) \times (A^* / \equiv_n^{\mathcal{L}})$, and $A^* / \equiv_n^{\mathcal{R}} \in \mathbf{R}$ and $A^* / \equiv_n^{\mathcal{L}} \in \mathbf{L}$ by Lemma 1 and its left-right dual. The properties “2” and “3” trivially follow from the definition of \equiv_n . \square

4 An Equation for the Join

The goal of this section is to prove $\mathbf{W} \subseteq \mathbf{R} \vee \mathbf{L}$. By Lemma 2 it suffices to show that for every A -generated monoid $M \in \mathbf{W}$ there exists an integer $n \geq 0$ such that M is a quotient of A^*/\equiv_n . The outline of the proof is as follows. First, in Lemma 3, we give a substitution rule valid in \mathbf{W} . Then, in Lemma 5, we show that \equiv_n -equivalence allows a factorization satisfying the premise for applying this substitution rule; this relies on a property of \mathbf{W} shown in Lemma 4. Finally, in Theorem 6, all the ingredients are put together.

Lemma 3 *Let $M \in \mathbf{W}$ and let $u, v, x \in M$. If $u \mathcal{R} ux$ and $v \mathcal{L} xv$, then $uxv = uv$.*

Proof: Since $u \mathcal{R} ux$ and $v \mathcal{L} xv$, there exist $y, z \in M$ with $u = uxy$ and $v = z xv$. In particular, we have $u = u(xy)^\omega$ and $v = (zx)^\omega v$. By $M \in \mathbf{W}$ we conclude $uxv = u(xy)^\omega x(zx)^\omega v = u(xy)^\omega (zx)^\omega v = uv$. \square

We will apply the previous lemma as follows. Let $M \in \mathbf{W}$ and $u, v, s, t \in M$ such that $u \mathcal{R} us \mathcal{R} ut$ and $v \mathcal{L} sv \mathcal{L} tv$. Then $usv = utv$ since $usv = uv$ and $utv = uv$ by Lemma 3. The \mathcal{R} -equivalences and \mathcal{L} -equivalences for being able to apply this substitution rule are established in Lemma 5. Before, we give a simple property of \mathbf{W} . It is the link between Green's relations and the congruence \equiv_n .

Lemma 4 *Let $M \in \mathbf{W}$ and let $u, v, a \in M$. If $u \mathcal{R} v \mathcal{R} va$, then $u \mathcal{R} ua$. If $u \mathcal{L} v \mathcal{L} av$, then $u \mathcal{L} au$.*

Proof: Since $u \mathcal{R} v$ and $u \mathcal{R} va$, there exist $x, y \in M$ with $v = ux$ and $u = vay$. Now, $u = uxa y = u(xay)^{2\omega+1} = u(xay)^\omega x(ayx)^\omega ay = u(xay)^\omega (ayx)^\omega ay = u(ayx)^\omega ay \in uaM$ where the fourth equality uses $M \in \mathbf{W}$. This shows $uM \subseteq uaM$ and thus $u \mathcal{R} ua$. The second implication is left-right symmetric. \square

The intuitive interpretation of the algebraic statement in Lemma 4 is the following: For $M \in \mathbf{W}$ it only depends on the element a and the \mathcal{R} -class of u whether $u \mathcal{R} ua$ or not (but not on the element u itself). The statement for \mathcal{L} -classes is analogous.

Lemma 5 *Let $M \in \mathbf{W}$ and let $\varphi : A^* \rightarrow M$ be a homomorphism. If $u \equiv_n v$ for $n \geq 2|M|$, then there exist factorizations $u = a_1 s_1 \cdots a_{\ell-1} s_{\ell-1} a_\ell$ and $v = a_1 t_1 \cdots a_{\ell-1} t_{\ell-1} a_\ell$ with $a_i \in A$ and $s_i, t_i \in A^*$ and with $\ell \leq 2|M|$ such that for all $i \in \{1, \dots, \ell-1\}$ we have:*

$$\begin{aligned} \varphi(a_1 s_1 \cdots a_{i-1} s_{i-1} a_i) \mathcal{R} \varphi(a_1 s_1 \cdots a_i s_i) \mathcal{R} \varphi(a_1 s_1 \cdots a_{i-1} s_{i-1} a_i t_i), \\ \varphi(a_{i+1} t_{i+1} \cdots a_{\ell-1} t_{\ell-1} a_\ell) \mathcal{L} \varphi(t_i a_{i+1} \cdots t_{\ell-1} a_\ell) \mathcal{L} \varphi(s_i a_{i+1} t_{i+1} \cdots a_{\ell-1} t_{\ell-1} a_\ell). \end{aligned}$$

Proof: To simplify notation, for some relation \mathcal{G} on M we write $u \mathcal{G} v$ for words $u, v \in A^*$ if $\varphi(u) \mathcal{G} \varphi(v)$. Consider the \mathcal{R} -factorization of u , i.e., let $u = b_1 u_1 \cdots b_k u_k$ with $b_i \in A$ such that

$$\begin{aligned} b_1 u_1 \cdots b_i \mathcal{R} b_1 u_1 \cdots b_i u_i & \quad \text{for all } i \in \{1, \dots, k\}, \\ b_1 u_1 \cdots b_i u_i >_{\mathcal{R}} b_1 u_1 \cdots b_i u_i b_{i+1} & \quad \text{for all } i \in \{1, \dots, k-1\}. \end{aligned}$$

Similarly, let $v = v_1 c_1 \cdots v_{k'} c_{k'}$ be the \mathcal{L} -factorization of v , i.e., we have $c_i \in A$ and

$$\begin{aligned} c_i \cdots v_{k'} c_{k'} \mathcal{L} v_i c_i \cdots v_{k'} c_{k'} & \quad \text{for all } i \in \{1, \dots, k'\}, \\ v_i c_i \cdots v_{k'} c_{k'} >_{\mathcal{L}} c_{i-1} v_i c_i \cdots v_{k'} c_{k'} & \quad \text{for all } i \in \{2, \dots, k'\}. \end{aligned}$$

We have $k, k' \leq |M|$ because neither the number of \mathcal{R} -classes nor the number of \mathcal{L} -classes can exceed $|M|$. By Lemma 4, we have $b_i \notin \alpha(u_{i-1})$ for all $i \in \{2, \dots, k\}$ and $c_i \notin \alpha(v_{i+1})$ for all $i \in \{1, \dots, k' - 1\}$. We use these properties to convert the \mathcal{R} -factorization of u to v and to convert the \mathcal{L} -factorization of v to u : Let $v = b_1 v'_1 \cdots b_k v'_k$ such that $b_i \notin \alpha(v'_{i-1})$, and let $u = u'_1 c_1 \cdots u'_{k'} c_{k'}$ with $c_i \notin \alpha(u'_{i+1})$. These factorizations exist because $u \equiv_n v$; in particular, by Lemma 2,

$$\begin{aligned} u_i b_{i+1} u_{i+1} \cdots b_k u_k &\equiv_{n-i} v'_i b_{i+1} v'_{i+1} \cdots b_k v'_k \\ v_1 c_1 \cdots v_{j-1} c_{j-1} v_j &\equiv_{n-k'-1+j} u'_1 c_1 \cdots u'_{j-1} c_{j-1} u'_j \end{aligned}$$

for all $i \in \{1, \dots, k\}$ and $j \in \{1, \dots, k'\}$. Moreover, we see that $\alpha(u_i) = \alpha(v'_i)$ and $\alpha(v_j) = \alpha(u'_j)$.

We now show that the relative positions of the b_i 's and c_j 's in the above factorizations are the same in u and v . Let p be the position of b_i in the \mathcal{R} -factorization of u and let q be the position of c_j in the above factorization of u . Similarly, let p' be the position of b_i in v and let q' be the position of c_j in v . First, suppose $p < q$. Let

$$u = b_1 u_1 \cdots b_{i-1} u_{i-1} b_i u' c_j u'_{j+1} c_{j+1} \cdots u'_{k'} c_{k'}.$$

By an i -fold application of property “2” in Lemma 2 with $a \in \{b_1, \dots, b_i\}$ (which is possible for u) we obtain $v = b_1 v'_1 \cdots b_{i-1} v'_{i-1} b_i z$ with $z \equiv_{n-i} u' c_j u'_{j+1} c_{j+1} \cdots u'_{k'} c_{k'}$. By a $(k' + 1 - j)$ -fold application of property “3” in Lemma 2 with $a \in \{c_{k'}, \dots, c_j\}$ (which is possible for the word $u' c_j u'_{j+1} c_{j+1} \cdots u'_{k'} c_{k'}$) we obtain $z = v' c_j v'_{j+1} c_{j+1} \cdots v'_{k'} c_{k'}$. Thus

$$v = b_1 v'_1 \cdots b_{i-1} v'_{i-1} b_i v' c_j v'_{j+1} c_{j+1} \cdots v'_{k'} c_{k'}$$

showing that $p' < q'$. Symmetrically, one shows that $p' < q'$ implies $p < q$. We conclude $p < q$ if and only if $p' < q'$. Similarly, we have $p = q$ if and only if $p' = q'$. It follows that the relative order of the b_i 's and c_j 's in u and v is the same. By factoring u and v at all b_i 's and c_j 's, we obtain $u = a_1 s_1 \cdots a_{\ell-1} s_{\ell-1} a_\ell$ and $v = a_1 t_1 \cdots a_{\ell-1} t_{\ell-1} a_\ell$ with $a_i \in A$ and $\ell \leq k + k' \leq 2|M|$.

We have $a_1 s_1 \cdots a_{i-1} s_{i-1} a_i \mathcal{R} a_1 s_1 \cdots a_{i-1} s_{i-1} a_i s_i$ since the factorization $u = a_1 s_1 \cdots a_{\ell-1} s_{\ell-1} a_\ell$ is a refinement of the \mathcal{R} -factorization. Note that we cannot assume $\alpha(s_i) = \alpha(t_i)$. But each t_i is a factor of some v'_j , and at the same time s_i is a factor of u_j . More precisely, there exists $m \leq i$ such that

$$b_1 v'_1 \cdots b_{j-1} v'_{j-1} b_j = a_1 t_1 \cdots a_{m-1} t_{m-1} a_m \quad \text{and} \quad t_m a_{m+1} \cdots t_{i-1} a_i t_i \text{ is a prefix of } v'_j.$$

Furthermore, $s_m a_{m+1} \cdots s_{i-1} a_i s_i$ is a prefix of u_j . Now, $\alpha(t_i) \subseteq \alpha(v'_j) = \alpha(u_j)$ and, by Lemma 4, for all words z with $\alpha(z) \subseteq \alpha(u_j)$ we have $a_1 s_1 \cdots a_{i-1} s_{i-1} a_i \mathcal{R} a_1 s_1 \cdots a_{i-1} s_{i-1} a_i z$. Symmetrically we see $a_{i+1} t_{i+1} \cdots a_{\ell-1} t_{\ell-1} a_\ell \mathcal{L} t_i a_{i+1} \cdots t_{\ell-1} a_\ell \mathcal{L} s_i a_{i+1} t_{i+1} \cdots a_{\ell-1} t_{\ell-1} a_\ell$. \square

Theorem 6 (Almeida/Azevedo, 1989 [2])

$$\mathbf{R} \vee \mathbf{L} = \llbracket (xy)^\omega x (zx)^\omega = (xy)^\omega (zx)^\omega \rrbracket$$

Proof: The inclusion $\mathbf{R} \vee \mathbf{L} \subseteq \mathbf{W}$ is trivial since $\mathbf{R} \cup \mathbf{L} \subseteq \mathbf{W}$ and \mathbf{W} is a variety of finite monoids. Let $M \in \mathbf{W}$ be generated by A , and let $\varphi : A^* \rightarrow M$ be the homomorphism induced by $A \subseteq M$. Let $n = 2|M|$ and

suppose $u \equiv_n v$. Let $u = a_1 s_1 \cdots a_{\ell-1} s_{\ell-1} a_\ell$ and $v = a_1 t_1 \cdots a_{\ell-1} t_{\ell-1} a_\ell$ be the factorizations from Lemma 5. Applying Lemma 3 repeatedly, we get

$$\begin{aligned}
 \varphi(v) &= \varphi(a_1 t_1 a_2 t_2 \cdots a_{\ell-2} t_{\ell-2} a_{\ell-1} t_{\ell-1} a_\ell) \\
 &= \varphi(a_1 s_1 a_2 t_2 \cdots a_{\ell-2} t_{\ell-2} a_{\ell-1} t_{\ell-1} a_\ell) \\
 &= \varphi(a_1 s_1 a_2 s_2 \cdots a_{\ell-2} t_{\ell-2} a_{\ell-1} t_{\ell-1} a_\ell) \\
 &\quad \vdots \\
 &= \varphi(a_1 s_1 a_2 s_2 \cdots a_{\ell-2} s_{\ell-2} a_{\ell-1} t_{\ell-1} a_\ell) \\
 &= \varphi(a_1 s_1 a_2 s_2 \cdots a_{\ell-2} s_{\ell-2} a_{\ell-1} s_{\ell-1} a_\ell) = \varphi(u).
 \end{aligned}$$

Note that the substitution rules $t_i \rightarrow s_i$ are φ -invariant only when applied from left to right. This shows that M is a quotient of A^*/\equiv_n , and the latter is in $\mathbf{R} \vee \mathbf{L}$ by Lemma 2. Thus $M \in \mathbf{R} \vee \mathbf{L}$. \square

Acknowledgements

We thank the anonymous referees for several suggestions which helped to improve the presentation of the paper, and we gratefully acknowledge the support by the German Research Foundation (DFG) under grant DI 435/5-1.

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