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On the Chromatic Number of some Flip Graphs

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This paper studies the chromatic number of the following four flip graphs (under suitable definitions of a flip):

- the flip graph of perfect matchings of a complete graph of even order,
- the flip graph of triangulations of a convex polygon (the associahedron),
- the flip graph of non-crossing Hamiltonian paths of a set of points in convex position, and
- the flip graph of triangles in a convex point set.

We give tight bounds for the latter two cases and upper bounds for the first two.

Keywords: chromatic number, flip graph, geometric graph, Hamiltonian path, perfect matching, triangulation

Mathematics Subject Classification: 05C15 (Coloring of graphs and hypergraphs)

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1 Introduction

Given a class \mathcal{C} of combinatorial objects of a given kind and a transformation (flip) between these objects, a flip graph is defined as the graph whose vertex set is \mathcal{C} , where two vertices are adjacent whenever the corresponding objects differ by a flip. Flip graphs have received considerable attention in the past. Properties such as Hamiltonicity, connectivity and diameter have been widely studied [4, 6, 13, 16]. This interest is very likely a consequence of the practical applications of these properties. For example, Hamiltonicity allows for rapid generation of the given combinatorial objects. We refer the interested reader to the survey [1].

A k -coloring of a graph G is a function that assigns one of k colors to each vertex of G , such that adjacent vertices receive distinct colors. The *chromatic number* $\chi(G)$ is the least integer k for which there is a k -coloring of G . The chromatic number can be formulated in terms of homomorphisms. A *homomorphism* from a graph G to a graph H is a mapping from $V(G)$ to $V(H)$ such that the images of any two adjacent vertices in G are adjacent in H . In this sense $\chi(G)$ is the minimum positive integer m such that there exists a homomorphism from G to the complete graph K_m . The composition of homomorphisms is again a homomorphism. Therefore if there exists a homomorphism from G to H then $\chi(G) \leq \chi(H)$.

The chromatic number of flip graphs has received little attention, with only a few papers concentrating on this parameter (see for example [5]). This paper studies the chromatic number of some flip graphs. We first consider three flip graphs defined on sets of points in convex position. (A set of points in the plane is in *convex position* if it forms the vertex-set of a convex polygon.) We then consider the flip graph on the perfect matchings of a complete graph of even order.

A *geometric non-crossing graph* is a planar graph in which the vertices are embedded as points in the plane, and where the edges are embedded as non-crossing straight line segments. Throughout the paper, S_n will denote a set of n points in convex position in the plane, and the term *convex hull* will refer to the boundary of the convex polygon having those points as vertex set. All of the considered graphs, except the flip graph of perfect matchings, have all of their vertices associated to non-crossing geometric graphs having S_n as vertex set. We remark that any underlying convex point set of the same size induces the same graphs.

This paper is organized as follows. In Section 2 we determine the exact chromatic number for non-crossing geometric Hamiltonian paths having S_n as their set of vertices. In Section 3 we define a certain flip graph on the triangles with vertices in S_n and determine its chromatic number up to a constant multiplicative factor. In Section 4 we give an upper bound on the chromatic number for the flip graph of the triangulations of S_n . In Section 5 we consider a flip graph on the perfect matchings of K_{2n} and give an upper bound on its chromatic number. (It should be stressed that in this graph no geometry is considered.) In Section 6 we conclude with some open problems and remarks.

2 Geometric Non-Crossing Hamiltonian Paths

Let G_n be the graph whose vertex set is the set of all geometric non-crossing paths with vertex set S_n . Two paths Γ and Γ' in $V(G_n)$ are adjacent if and only if Γ' is obtained from Γ by deleting some edge e and adding some edge f . (See Figure 1.) We say that Γ' is obtained from Γ by *flipping* the edge e . We point out that e and f may cross or they may share an endpoint. Rivera-Campo [15] and Urrutia-Galicia [17] proved that G_n is Hamiltonian. We determine $\chi(G_n)$.

Theorem 2.1 *For every integer $n \geq 3$, $\chi(G_n) = n$.*

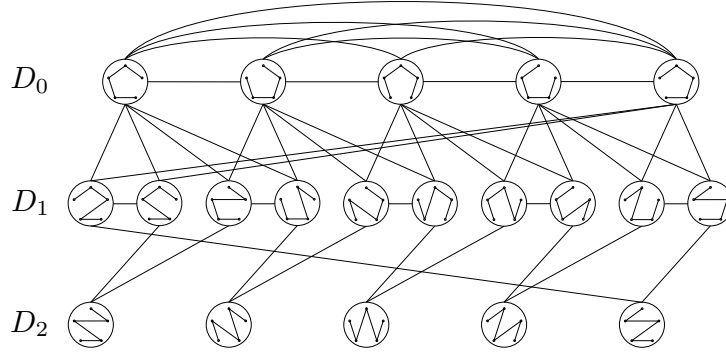


Fig. 1: G_5 drawn with vertices grouped by the number of non-convex hull edges.

Proof: Since $G_3 \simeq K_3$, assume $n \geq 4$. For all integers $i \in [0, n - 3]$, let D_i be the set of paths in $V(G_n)$ that contain exactly i edges not in the convex hull of S_n . (See Figure 1). Note that the set D_0 consists of n paths, all of whose edges are in the convex hull of S_n . Each element of D_0 is obtained by removing one edge of the convex hull of S_n . Clearly any two elements of D_0 are adjacent in G_n , therefore D_0 induces a complete subgraph of G_n . This proves that $\chi(G_n) \geq n$.

Let H_i be the subgraph induced in G_n by $D_0 \cup \dots \cup D_i$. We prove $\chi(G_n) \leq n$ by showing how to extend an n -coloring of H_{i-1} to an n -coloring of H_i for every $i \in [1, n - 3]$. For this, let Γ_i be a path in D_i .

For extending an n -coloring of H_0 to one of H_1 we use the next two properties of G_n . (1) The set D_1 induces a perfect matching in G_n . (2) Every adjacent pair in D_1 together with its neighborhood in H_1 induces a K_4 in G_n . Both properties follow because, on one hand, flipping an edge of Γ_1 that is in the convex hull of S_n yields a path in D_2 . On the other hand, flipping the edge of Γ_1 that is not in the convex hull of S_n yields either one path in D_1 , or one of two different paths in D_0 . Since $n \geq 4$, each pair of adjacent paths in D_1 can be colored with colors different to those of its neighbors in D_0 .

For extending an n -coloring of H_{i-1} to one of H_i when $i \geq 2$, we use the next two properties of G_n . (1) The set D_i is an independent set in G_n . (2) The vertex of G_n corresponding to the path Γ_i has degree 2 in H_i . The first property follows because, on one hand, flipping an edge of Γ_i that is in the convex hull of S_n yields a path in D_{i+1} . On the other hand, flipping an edge of Γ that is not in the convex hull of S_n yields a path in D_{i-1} . The latter property follows because the only edges of Γ that are not in the convex hull of S_n that can be flipped, are the first and last when walking from one leaf of Γ to the other.

Since n colors are needed to color D_0 , and any n -coloring of D_0 can be extended to an n -coloring of G_n , the result follows. \square

3 Triangle Graph

We introduce a new flip graph $G_\Delta(n)$, whose vertex set is the set of triangles with endpoints in S_n , two of which are adjacent if they share an edge and their interiors are disjoint. (See Figure 2.) We prove that this graph has chromatic number $\Theta(\log n)$.

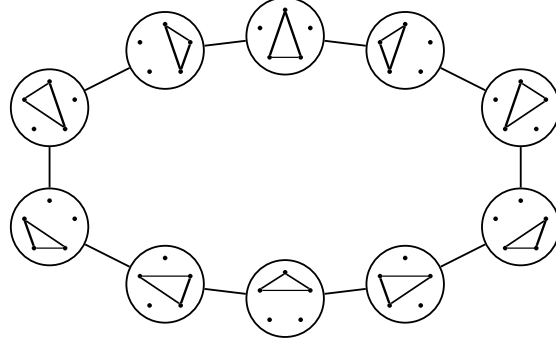


Fig. 2: The graph $G_{\Delta}(10)$.

For the rest of this section, assume that the elements of S_n are labeled $\{0, \dots, n-1\}$ in clockwise order, and let any triangle be denoted by the set of its vertices.

Lemma 3.1 For every integer $n \geq 3$, $\chi(G_{\Delta}(n)) \geq \log_2(n-1)$.

Proof: Let H be the subgraph of $G_{\Delta}(n)$ induced by all triangles containing 0 as a vertex. Suppose that H has a k -coloring. Let A_i (with $1 \leq i \leq n-2$), be the set of colors assigned to the triangles $\{0, i, j\}$ (with $i < j \leq n-1$); see Figure 3 (a). Observe that $A_r \neq A_s$ for $r \neq s$. This follows from the fact that if $r < s$, the triangle $\{0, r, s\}$ is adjacent to every triangle $\{0, s, k\}$ (with $s < k \leq n-1$); see Figure 3 (b). Thus, the color assigned to $\{0, r, s\}$ is in A_r but not in A_s . The number of possible non-empty color subsets must be at least $n-2$. Therefore $2^k \geq n-1$, implying $k \geq \log_2(n-1)$ and $\chi(G_{\Delta}(n)) \geq \log_2(n-1)$. \square

Note that the subgraph H of G_{Δ} considered in the proof of Lemma 3.1 is isomorphic to the so-called *shift graph* [7] of the set $\{1, 2, \dots, n-1\}$. It is folklore that its chromatic number is $\lceil \log_2(n-1) \rceil$.

Lemma 3.2 For every integer $m \geq 3$, let G be the subgraph of $G_{\Delta}(2m)$ induced by all the triangles without edges in the convex hull of S_{2m} . Then $\chi(G) = \chi(G_{\Delta}(m))$.

Proof: Color the vertices of S_{2m} red and blue such that consecutive vertices along the convex hull have distinct colors. Note that the subgraph H of G induced by the blue points is isomorphic to $G_{\Delta}(m)$. Map every triangle t in $V(G)$ to a triangle t' in $V(H)$ as follows: If i is a red vertex of t , substitute it by the blue vertex $i+1$, where addition is modulo $2m$. Observe that triangles whose vertices are all blue are mapped onto themselves. Since adjacent triangles in G are mapped to adjacent triangles in H , this mapping induces a homomorphism from G to H (see Figure 3). The result follows. \square

Lemma 3.3 For every integer $m \geq 3$, $\chi(G_{\Delta}(2m)) \leq \chi(G_{\Delta}(m)) + 3$.

Proof: We use a similar technique to that used in Theorem 2.1. Let G' be the subgraph of $G_{\Delta}(2m)$ induced by all those triangles of S_{2m} that have at least one edge in the convex hull of S_{2m} . Let $\tau = \{i, i+1, j\}$ be any such triangle. $S_{2m} - \{i, i+1, j\}$ can be partitioned into two maximal subsets of

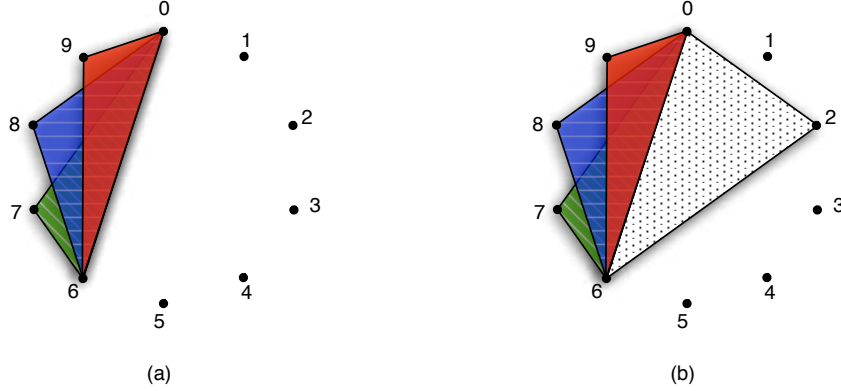


Fig. 3: Illustration of Lemma 3.1 (with $s = 6$). (a) The triangles considered in A_s are those having the segment $\overline{0s}$ as its right side. (b) The triangle $\{0, r, s\}$ (with $r < s$) is adjacent in H to all the triangles considered in A_s .

consecutive points along the convex hull, namely the points from $i + 1$ to j and the points from j to i , which we denote by l_τ and r_τ respectively. Let the order of τ be $\min\{|l_\tau|, |r_\tau|\}$. For $i \in [0, m - 2]$, let D_i be the subset of $V(G')$ of all triangles of order i . Note that the subgraph of G' induced by D_{m-2} has maximum degree 2 and therefore chromatic number at most 3. Note that every vertex of D_i is only adjacent to at most two vertices in $D_i \cup \dots \cup D_{m-2}$. Thus any 3-coloring of $D_{i+1} \cup \dots \cup D_{m-2}$ can be extended to a 3-coloring of $D_i \cup D_{i+1} \cup \dots \cup D_{m-2}$. Therefore $\chi(G') = 3$.

Let G be the subgraph of $G_\Delta(2m)$ induced by all the triangles without edges of the convex hull of S_{2m} . Observe that $V(G) \cup V(G') = V(G_\Delta(2m))$. Since there is a 3-coloring of G' and (by Lemma 3.2) there is a coloring of G with $\chi(G_\Delta(m))$ colors, there is a coloring of $G_\Delta(2m)$ with $\chi(G_\Delta(m)) + 3$ colors. \square

We now have:

Theorem 3.4 For every integer $n \geq 3$, $\log_2(n - 1) \leq \chi(G_\Delta(n)) \leq 3\lceil \log_2(n) \rceil - 4$.

Proof: Let $F(k) = \chi(G_\Delta(2^k))$. By Lemma 3.3, $F(k) \leq F(k - 1) + 3$. Observe that $\chi(G_\Delta(4)) = 2$. Thus $F(k) \leq 3k - 4$ and therefore $\chi(G_\Delta(2^k)) \leq 3k - 4$.

For $n \neq 2^k$, let m be the smallest power of 2 greater than n . Since $G_\Delta(n)$ is a subgraph of $G_\Delta(m)$, we can color the vertices of $G_\Delta(n)$ with $\chi(G_\Delta(m)) \leq 3 \log_2(m) - 4 = 3\lceil \log_2(n) \rceil - 4$ colors. \square

4 Triangulations of a Convex Polygon

We consider a flip graph on the triangulations of S_n . Let $G_T(n)$ be the graph whose vertex set is the set of triangulations of S_n , where two triangulations T_1 and T_2 are adjacent in $G_T(n)$ whenever they differ by one edge flip; that is, there exist edges $e \in T_1$ and $f \in T_2$ such that T_2 is obtained from T_1 by deleting e and adding f (see Figure 5).

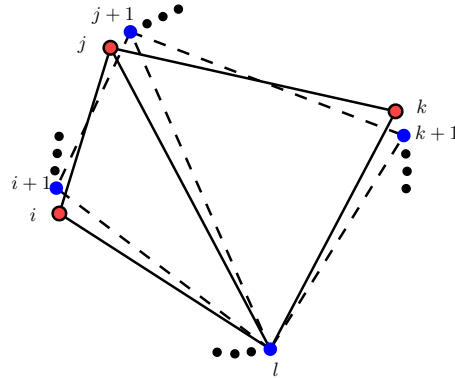


Fig. 4: The homomorphism between G and $G_\Delta(m)$.

Much is known about $G_T(n)$. This is probably because there is a bijection between triangulations of the n -gon and binary trees with $n - 2$ interior nodes. A flip in a triangulation corresponds to a rotation in its corresponding binary tree. Lucas [13] proved that $G_T(n)$ is Hamiltonian. Sleator et al. [16] proved that the diameter of $G_T(n)$ is $2n - 10$. Lee [11] proved that the automorphism group of $G_T(n)$ is the dihedral group of symmetries of a regular n -gon. Lee also proved that $G_T(n)$ can be realized as the skeleton of an $(n - 3)$ -dimensional convex polytope called the *associahedron*. Most of these results are proved again in [10] using a unifying framework called the *tree of triangulations*.

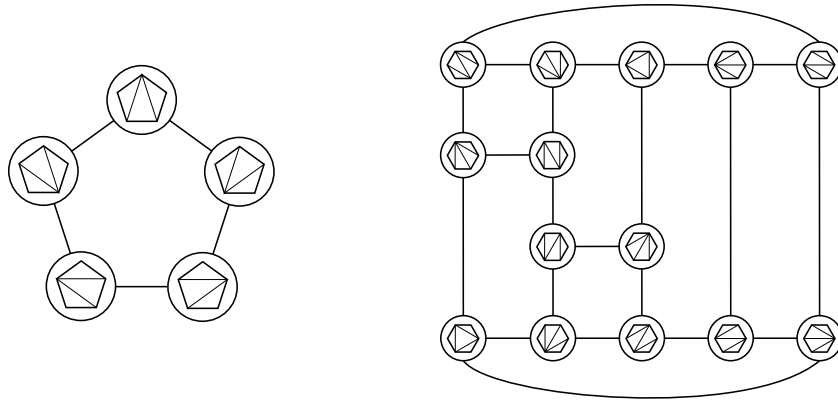


Fig. 5: The graphs $G_T(5)$ and $G_T(6)$.

We give an upper bound on the chromatic number of $G_T(n)$.

Theorem 4.1 For every integer $n \geq 3$, $\chi(G_T(n)) \leq \lceil \frac{n}{2} \rceil$.

Proof: It is well known that for n even, the set of $\binom{n}{2}$ edges between the vertices of S_n can be partitioned into $\frac{n}{2}$ edge-disjoint geometric non-crossing paths (see [12, 2] and Figure 6 for an example with $n = 6$).

Label these paths $P_1, \dots, P_{\frac{n}{2}}$. (Since $G_T(n)$ is a subgraph of $G_T(n+1)$, assume n is even for the time being.) If an edge e belongs to P_i , assign to it the weight $w(e) = i$. To every triangulation \mathcal{T} of S_n , assign the number $(\sum_{e \in \mathcal{T}} w(e)) \pmod{\frac{n}{2}}$. Let \mathcal{T}_1 and \mathcal{T}_2 be two adjacent triangulations of $G_T(n)$. That is, $\mathcal{T}_2 = \mathcal{T}_1 - e + f$ for some crossing edges e and f . Since e and f cross, $w(e) \neq w(f)$, and thus the numbers associated to \mathcal{T}_1 and \mathcal{T}_2 are different. This induces a coloring of $G_T(n)$ with $\frac{n}{2}$ colors for n even and $\lceil \frac{n}{2} \rceil$ in general. \square

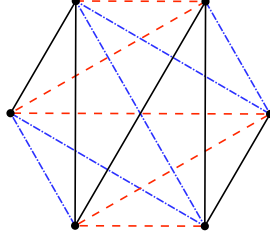


Fig. 6: Partition of the edges into disjoint non-crossing geometric Hamiltonian paths.

The upper bound on $\chi(G_T(n))$ in Theorem 4.1 is non-trivial since, for example, Brooks' Theorem [3] gives an upper bound of only $n - 3$. However, $\chi(G_T(n))$ is in fact sublinear, as we now show. Johansson proved that for sufficiently large Δ , every triangle-free graph with maximum degree Δ is $O(\frac{\Delta}{\log \Delta})$ -colorable; see [14]. The next theorem follows since $G_T(n)$ is $(n - 3)$ -regular and triangle-free.

Theorem 4.2 For every integer $n \geq 3$, $\chi(G_T(n)) = O(\frac{n}{\log n})$.

We remark that the proof of Theorem 4.1 is constructive, whereas Johansson's proof is probabilistic.

5 Perfect Matchings of K_{2n}

Given any graph G , let $\mathcal{M}(G)$ be the graph whose vertex set is the set of perfect matchings of G , where two vertices of $\mathcal{M}(G)$ are adjacent whenever the symmetric difference of the corresponding perfect matchings is a cycle of length 4. (See Figure 7).

We remark that (contrary to the other graphs we study in this paper) we do not consider geometry in the definition of $\mathcal{M}(G)$. This is because the geometric version of this graph⁽ⁱ⁾ is bipartite, as was shown by Hernando et al. [8]. Interestingly, the chromatic number turns out to be quite challenging when geometry is discarded.

We first determine the number of vertices of $\mathcal{M}(K_{2n})$. Let $f(n) := |V(\mathcal{M}(K_{2n}))|$ be the number of perfect matchings of K_{2n} . Say $V(K_{2n}) = \{v_1, \dots, v_{2n}\}$. In every such matching, v_1 is matched with some vertex in $\{v_2, \dots, v_{2n}\}$, and the remaining matching is isomorphic to a perfect matching of K_{2n-2} . Every matching obtained in this way is distinct. Thus

$$f(n) = (2n - 1) \cdot f(n - 1),$$

⁽ⁱ⁾ In the geometric version of this graph, the vertices are the non-crossing perfect matchings of a set of $2n$ points in convex position; two of which are adjacent whenever their symmetric difference is a cycle of length four.

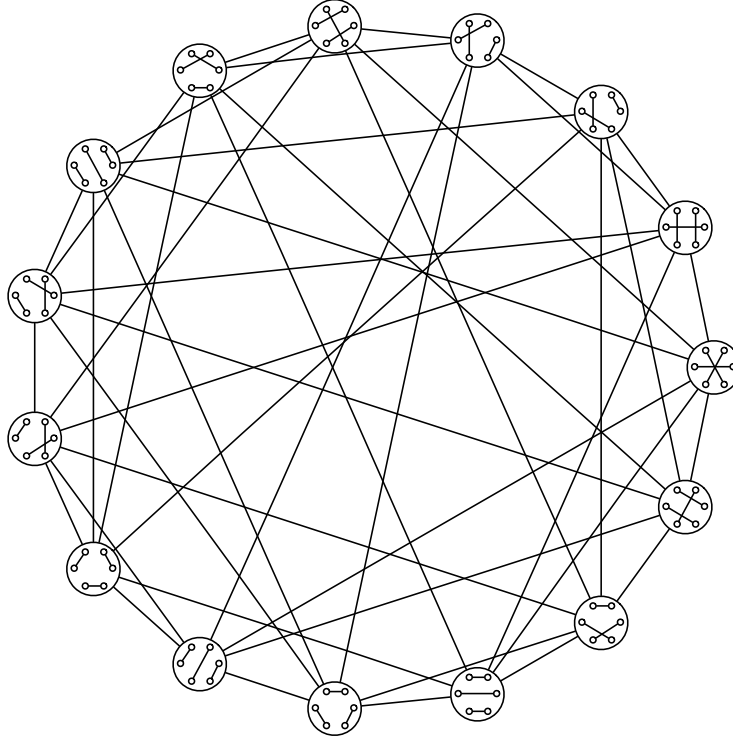


Fig. 7: The graph $\mathcal{M}(K_6)$.

where $f(1) = 1$. Hence⁽ⁱⁱ⁾

$$f(n) = (2n - 1)!! = (2n - 1)(2n - 3)(2n - 5) \dots 1 = \frac{(2n)!}{2^n n!}.$$

By Stirling's Approximation,

$$f(n) = \frac{(2n)!}{2^n n!} \approx \sqrt{2} \left(\frac{2n}{e} \right)^n.$$

Also observe that $\mathcal{M}(K_{2n})$ is $n(n - 1)$ -regular. Furthermore, the open neighborhood of each vertex of $\mathcal{M}(K_{2n})$ induces a perfect matching; thus $\mathcal{M}(K_{2n})$ has no K_4 subgraph.

We now give an upper bound on $\chi(\mathcal{M}(K_{2n}))$. To do so, we use the fact that $\chi(\mathcal{M}(K_{n,n})) = 2$. In [8], the flip graph of the non-crossing geometric matchings of a set of $2n$ points in convex position is shown to be bipartite. This graph is actually a subgraph of $\mathcal{M}(K_{n,n})$. Using the same arguments as in [8] we show that $\mathcal{M}(K_{n,n})$ is bipartite: Let $\{1, \dots, n\}$ and $\{1', \dots, n'\}$ be the color classes of $K_{n,n}$. Then each perfect matching of $\mathcal{M}(K_{n,n})$ can be interpreted as a permutation σ of $\{1, \dots, n\}$, where

⁽ⁱⁱ⁾ Note that the operator $!!$ is called *double factorial* in the literature.

$\sigma(i) = j'$ if i is matched to j' . Now, two perfect matchings of $\mathcal{M}(K_{n,n})$ are adjacent if and only if their associated permutations differ by a transposition. Thus, these associated permutations have different parity and consequently $\mathcal{M}(K_{n,n})$ is bipartite.

Theorem 5.1 For every integer $n \geq 2$, $\chi(\mathcal{M}(K_{2n})) \leq 4n - 4$.

Proof: Let $\{1, \dots, n\}$ and $\{1', \dots, n'\}$ be the color classes of $K_{n,n}$, and let $\{1, \dots, 2n\}$ be the set of vertices of K_{2n} . For every perfect matching M of K_{2n} , let $U(M) := \{i \in V(K_{2n}) \mid (i, j) \in M \text{ and } i > j\}$ and $D(M) := \{i \in V(K_{2n}) \mid (i, j) \in M \text{ and } i < j\}$. Assign to every matching M the number $c(M) := \left(\sum_{i \in D(M)} i \right) \bmod (2n - 2)$. Given two sets U and D , let $\mathbb{M}(U, D)$ be the set of matchings M such that $U(M) = U$ and $D(M) = D$. The subgraph $H_{U,D}$ of $\mathcal{M}(K_{2n})$ induced by $\mathbb{M}(U, D)$ is a subgraph of $\mathcal{M}(K_{n,n})$, and is thus 2-colorable. Color the vertices of $H_{U,D}$ with colors $c(M)$ and $c'(M) := c(M) + 2n - 2$.

We show now that if two matchings M and M' differ by a flip, then they receive different colors. Two cases arise: Either $U(M) = U(M')$ and $D(M) = D(M')$, or for some $i \neq j$, $U(M) = U(M') - \{i\} + \{j\}$ and $D(M) = D(M') - \{j\} + \{i\}$. In the first case, M and M' belong to $H_{U(M), D(M)}$ and thus receive different colors. In the second case, $c(M) = c(M') - j + i$. Since $1 \in D(M) \cap D(M')$ and $2n \in U(M) \cap U(M')$, then $|i - j| \leq 2n - 3$. Therefore $c(M) \neq c(M')$, and consequently the colors assigned to M and M' are different.

Thus we obtain a coloring of $\mathcal{M}(K_{2n})$ using $2(2n - 2) = 4n - 4$ colors. \square

6 Open Problems

The most challenging open problems arising from this work are to determine the chromatic numbers of $G_T(n)$ and $\mathcal{M}(K_{2n})$. Despite our efforts, we have not been able to obtain non-trivial lower bounds for these graphs. We can formulate the following educated guess on the chromatic number of $\mathcal{M}(K_{2n})$.

Conjecture 6.1 For every integer $n \geq 2$, $\chi(\mathcal{M}(K_{2n})) = n + 1$.

We have verified this conjecture with the aid of a computer for $n = 2, 3, 4$.

We have studied $\mathcal{M}(K_{2n})$, but it would be very interesting to study $\mathcal{M}(G)$ for other graphs.

There is another flip graph, $\mathcal{M}'(G)$, related to $\mathcal{M}(G)$. $\mathcal{M}'(G)$ has as its vertex set the perfect matchings of G , where two matchings are now adjacent if and only if their symmetric difference is a cycle of arbitrary length. We have yet to study its chromatic number.

With respect to $\chi(G_T(n))$, the problem seems far more intriguing since we not only do not have a non-trivial lower bound, but we also believe that our upper bound is far from being tight:

Conjecture 6.2 $\lim_{n \rightarrow \infty} \chi(G_T(n)) = \infty$ and $\chi(G_T(n)) = O(\log(n))$.

Let us mention that we do not have any strong evidence on the conjectured upper bound, and the tight value might be $\Theta(\log \log(n))$ or $\Theta(\text{polylog}(n))$, say, but we have chosen a precise value as more stimulating for research.

It would be interesting to see which techniques will be capable of improving the lower bounds of both of these graphs. It is likely that such techniques will turn out to be useful in determining the chromatic number of other graphs as well.

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