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Vincent Roca, Kazuhisa Matsuzono

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# Structured RLC codes: an update

### Vincent Roca (Inria, France) Kazuhisa Matsuzono (NICT, Japan) \*

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( A part of the work done while visiting Inria as post-doctorate)





#### Note well

- we, authors, didn't try to patent any of the material included in this presentation
- we, authors, are not reasonably aware of patents on the subject that may be applied for by our employer
- if you believe some aspects may infringe IPR you are aware of, then fill in an IPR disclosure and please, let us know
   <a href="http://irtf.org/ipr">http://irtf.org/ipr</a>

# Our proposal and some results in block mode... A reminder

For details, see:

http://www.ietf.org/proceedings/88/slides/slides-88-nwcrg-2.pdf





# **Goals (from IETF88)**

#### design codes that

Ocan be used indifferently as sliding/elastic/block codes

Ocan be used with encoding window/block sizes in 1-10,000s symbols range

Okeep high enc./decoding speeds and erasure recovery performance in all cases

Ocan be used as **small-rate** codes

Oit's not necessarily required, but it simplifies many things

ofocus only on use-cases that need end-to-end coding
Oe.g. for FLUTE/ALC, FECFRAME, or Tetrys

Oenable compact and robust signaling (essential!)
 Ovectors can help for tiny k values but it's unfeasible above
 Ouse a known function + key (e.g. PRNG + seed)

# Two key ideas

 idea 1: mix binary and non binary coefficients
 most equations are sparse and coefficients binary
 a limited number of columns are dense and use nonbinary coefficients on GF(2<sup>8</sup>)

idea 2: add a structure

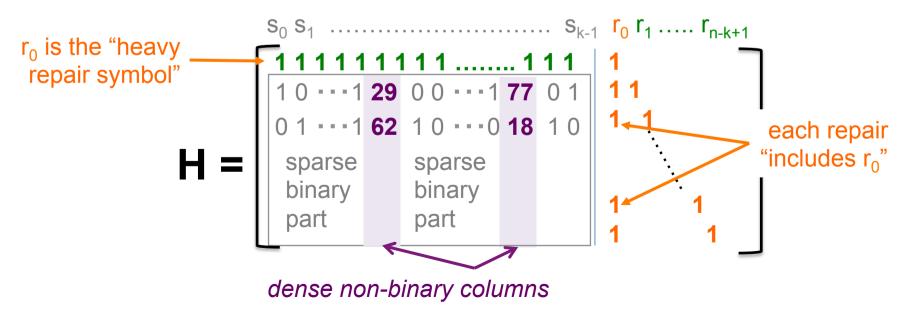
Oadd a single dense row (e.g. XOR sum of all source symbols) and make all repair symbols depend on it

# Let's put ideas 1 and 2 together

#### • 3 key parameters

- Ok block or encoding window size
- **OD\_bin** controls the density of the sparse sub-matrices
- **OD\_nonbin** controls number of dense non-binary columns
  - {D\_nonbin, D\_bin} depend on k and a target maximum average overhead

#### example: in block mode

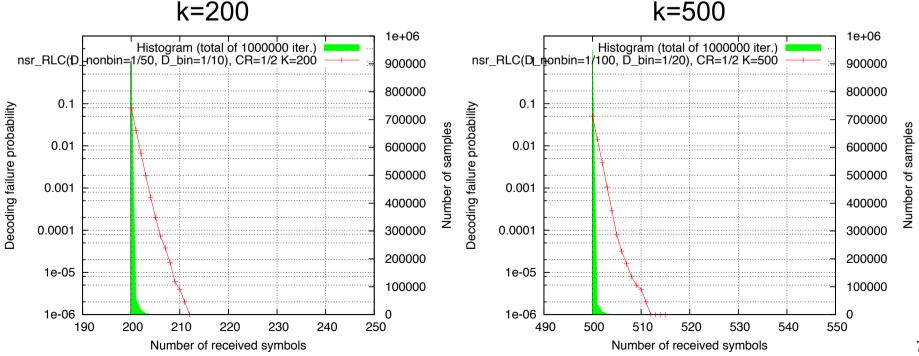


### It works well as a block AL-FEC code

#### • it works well on average...

Oparameters are chosen so that the average overhead is always below, say 10<sup>-3</sup> (meaning k\*10<sup>-3</sup> add. symbols needed)

and when looking at decoding failure proba. curves
 Ono visible error floor at 10<sup>-5</sup> failure probability ©



# What about sliding window mode?

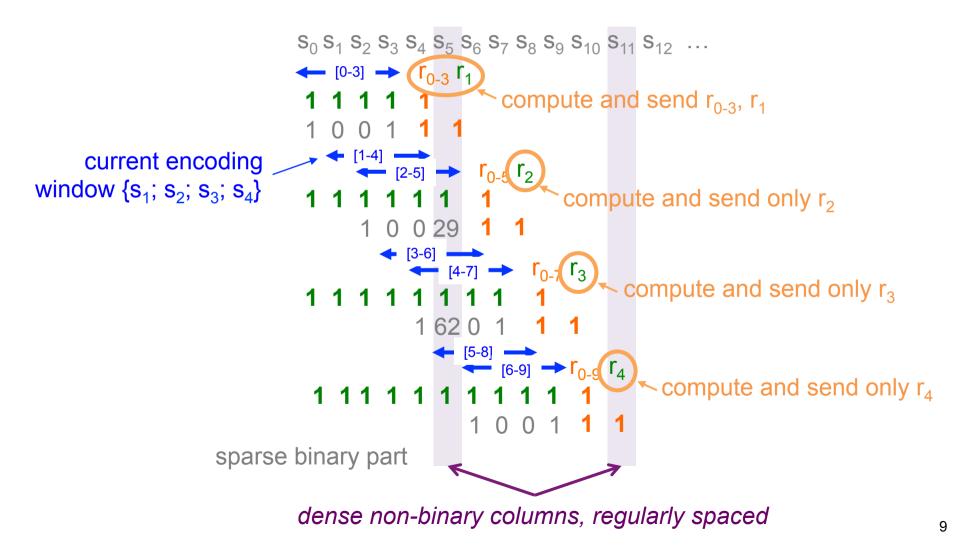




# Structured RLC in sliding window mode

### with a fixed length (k) sliding window

**Oexample:** k=4, CR=2/3  $\Rightarrow$  send one repair after 2 src symbols



# Struct. RLC in sliding window mode (cont')

#### about the previous example

- Oat session start, we wait k symbols to be available, and then compute and send a few repair symbols to match the target code rate
- Oafterwards we mix source and repair symbols in a periodic way
- Oeach repair that is not a heavy symbol "accumulates" the current heavy repair symbol
   Oi.e. the XOR sum from s₀ to the highest known symbol
   Othe current sum repair symbol is sent from time to time
- Othe D<sub>nonbin</sub>/D<sub>bin</sub> are set according to the fixed k value and desired average overhead, using pre-calculated tables

# A few experiments

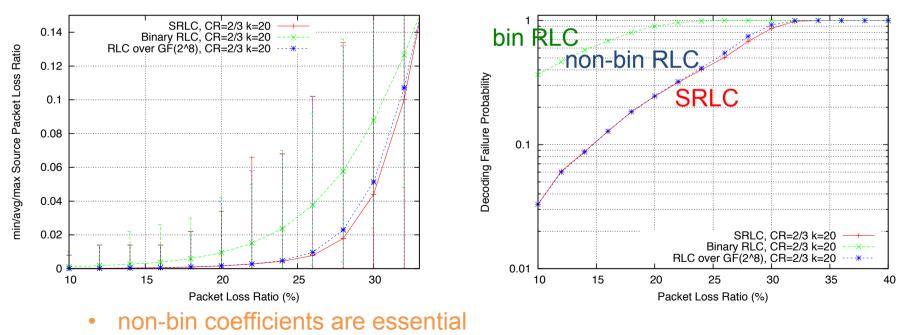
#### test conditions (small k=20)

Othe encoding window (size k = 20) slides over a flow of 25\*k = 500 source symbols

**OCR = 2/3**, send 1 repair after 2 source symbols

**Oplot Pr**<sub>fail</sub>(plr) post-repair curves for the whole transmission

does not catch the number of non recovered source symbols



• the heavy repair symbol improves performance WRT. RLC over GF(2<sup>8</sup>)

# A few experiments... (cont')

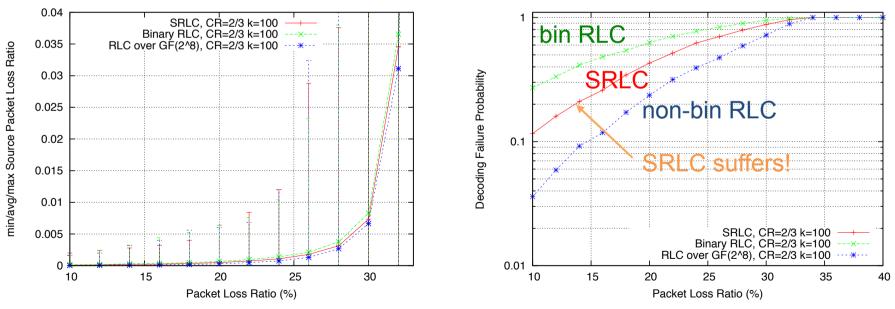
#### test conditions (medium k=100)

Othe encoding window (size k = 100) slides over a flow of 25\*k = 2500 source symbols

**OCR = 2/3, send 1 repair after 2 source symbols** 

**Oplot Pr**<sub>fail</sub>(plr) post-repair curves for the whole transmission

· does not catch the number of non recovered source symbols

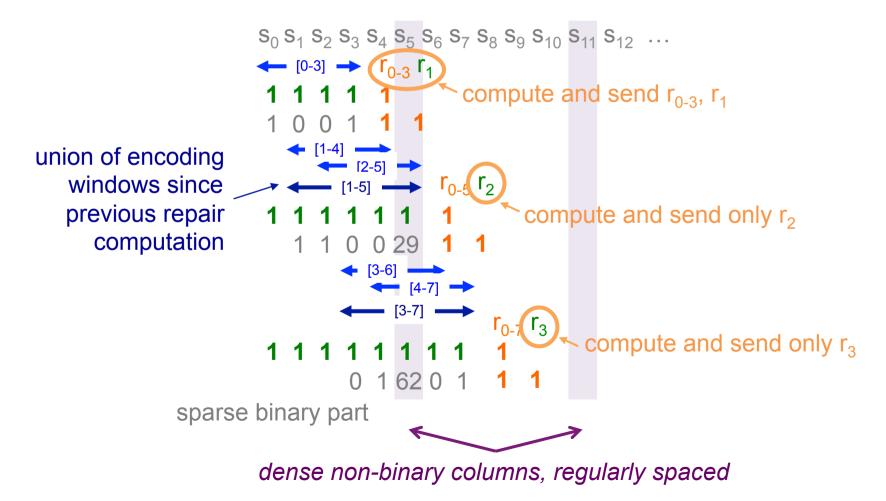


 we reused D\_bin/D\_nonbin values computed for the block mode, which is perhaps not appropriate here...

# An improvement (under progress)

 consider the union of encoding windows when computing new repair symbols...

Owill make a difference with small k and high CR values



# Conclusions





# **Conclusions**

#### our proposal tries to take the best of RLC

- fill in the gap between sliding/elastic window codes and block codes
- Ouse the right technique (bin vs. non-bin coefficients) at the right time, in the right way
  - find balance between erasure recovery perf. and complexity
- a lot remains to be done yet...
  - Ohow fast is it?
  - e.g., compared to our optimized LDPC-Staircase/RS codecs
     how does it scale with k?
    - e.g., compared to our optimized LDPC-Staircase codec

#### Odefine **signaling** aspects

· it's a critical practical topic

# Thank you!



