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Sound and Complete Bisimilarities for Call-by-Name and Call-by-Value $\lambda\mu$ -Calculus

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Abstract: We propose the first sound and complete bisimilarities for the call-by-name and call-by-value untyped $\lambda\mu$ -calculus. We define applicative bisimilarities for both semantics and environmental bisimilarity for call-by-name. We give equivalence examples to illustrate how our relations can be used; in particular, we prove David and Py's counter-example, which cannot be proved with Lassen's preexisting normal form bisimilarities for the $\lambda\mu$ -calculus.

Key-words: Applicative bisimilarity, $\lambda\mu$ -calculus, call-by-name, call-by-value

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Bisimilarités correctes et complètes pour le $\lambda\mu$ -calcul en appel par nom et par valeur

Résumé : Nous proposons les premières définitions de bisimilarités correctes et complètes pour le $\lambda\mu$ -calcul non typé en appel par nom et en appel par valeur. Nous définissons une bisimilarité applicative pour chacune des sémantiques, et une bisimilarité environnementale en appel par nom. Nous donnons des exemples d'équivalences pour montrer comment ces relations peuvent être utilisées ; en particulier, nous prouvons le contre-exemple de David et Py, qui ne peut être démontré avec la bisimilarité de forme normale définie auparavant par Lassen.

Mots-clés : Bisimilarité applicative, $\lambda\mu$ -calcul, appel par nom, appel par valeur

1 Introduction

Contextual equivalence [11] is considered as the most natural behavioral equivalence in languages based on the λ -calculus. Two terms are contextually equivalent if an outside observer cannot tell them apart when they are evaluated within any *context* (a term with a hole). However, the quantification over contexts makes proving the equivalence of two given programs difficult. Consequently, characterizations of contextual equivalence are sought for, usually in the form of inductively defined *bisimilarities*.

Several kinds of bisimilarity have been proposed, such as, e.g., *applicative* bisimilarity [1], which relates terms by reducing them to values (if possible), and then compares these values by applying them to an arbitrary argument. The idea is the same for *environmental* bisimilarity [16], except the values are tested with arguments built from an environment, which represents the knowledge of an observer about the tested terms. Finally, *normal form* bisimilarity [9] (initially called *open* bisimilarity [15]) reduces open terms to normal forms and then compares their subterms. Applicative and environmental bisimilarities still contain some quantification over arguments, and usually coincide with contextual equivalence. In contrast, normal form bisimilarity is easier to use as its definition does not contain any quantification over arguments, but it is generally not *complete*, i.e., there exist equivalent terms that are not normal form bisimilar.

This article treats the behavioral theory of the untyped $\lambda\mu$ -calculus [13]. The $\lambda\mu$ -calculus provides a computational interpretation of the classical natural deduction and thus extends the Curry-Howard correspondence [6] from intuitionistic to classical logic. Operationally, the reduction rules of the calculus express not only function applications but also captures of the current context of evaluation. Therefore, when considered in the untyped setting, the calculus offers an approach to the semantics of abortive control operators such as *callcc* known from the Scheme programming language and it may be viewed as a closely related alternative to Felleisen and Hieb's syntactic theory of control [4].

So far no characterization of contextual equivalence have been proposed in either call-by-value or call-by-name $\lambda\mu$ -calculus. Lassen defines normal bisimilarities for call-by-name [8], for head reduction semantics [10], and, with Støvring, for call-by-value [18] that are not complete. However, normal form bisimilarity is complete for Saurin's $\Lambda\mu$ -calculus [17] with head reduction semantics [10], and also for call-by-value $\lambda\mu$ -calculus with store [18]. Lassen also defines a not complete applicative bisimilarity for call-by-name in [8]. A definition of applicative bisimilarity has also been proposed for a call-by-value typed μ PCF [12], but the resulting relation is neither sound nor complete.

In this work we propose the first characterizations of contextual equivalence for call-by-name and call-by-value $\lambda\mu$ -calculus. We define sound and complete applicative bisimilarities for both semantics, and we also discuss environmental bisimilarity for call-by-name. The relations we obtain are harder to use than Lassen's normal form bisimilarity to prove the equivalence of two given terms, but because they are complete, we can equate terms that cannot be related with normal form bisimilarity, such as David and Py's counter-example [3]. Even

though the applicative bisimilarities we define for call-by-name and call-by-value are built along the same principles, the relation we obtain in call-by-value is much more difficult to use than the one for call-by-name. However, we provide counter-examples showing that it cannot be simplified.

The rest of the paper is organized as follows. We first discuss the behavioral theory of the call-by-name $\lambda\mu$ -calculus in Section 2. We propose a notion of contextual equivalence (in Section 2.2) which observes top-level names, and we then characterize it with, respectively, applicative (Section 2.3) and environmental (Section 2.4) bisimilarities. In particular, we compare our definition of applicative bisimilarity with Lassen’s one and we prove David and Py’s counter-example using applicative bisimilarity. We then discuss call by value in Section 3. We propose a definition of applicative bisimilarity (Section 3.2) which coincides with contextual equivalence. We then show that the definition cannot be simplified by providing counter-examples (Section 3.3). Although the relation we obtain is harder to use than the one for call-by-name, we can still prove some equivalences of terms, as we demonstrate in Section 3.4. We conclude in Section 4, and the appendices contain some of the proofs missing from the main text (Appendix A for call-by-name and Appendix B for call-by-value).

2 Call-by-Name $\lambda\mu$ -calculus

In this section, we define applicative and environmental bisimilarities for the call-by-name (CBN) $\lambda\mu$ -calculus. We first recall its syntax and semantics, and we then discuss the definition of contextual equivalence for the calculus.

2.1 Syntax and Semantics

Syntax. The $\lambda\mu$ -calculus [13] extends the λ -calculus with named terms and a μ constructor that binds names in terms. We assume we have a set X of *variables*, ranged over by x, y , etc., and a distinct set A of *names*, ranged over by a, b , etc. The syntax of terms (T) and named terms (U) is given by the following grammar:

$$\begin{array}{l} \text{Terms:} \quad t ::= x \mid \lambda x.t \mid tt \mid \mu a.u \\ \text{Named terms:} \quad u ::= [a]t \end{array}$$

Values (V), ranged over by v , are terms of the form $\lambda x.t$. A λ -abstraction $\lambda x.t$ binds x in t and a μ -abstraction $\mu a.t$ binds a in t . We equate terms up to α -conversion of their bound variables and names. We write $\text{fv}(t)$ and $\text{fv}(u)$ for the set of free variables of, respectively, t and u , and we write $\text{fn}(t)$ and $\text{fn}(u)$ for their set of free names. A term t or named term u is said *closed* if, respectively, $\text{fv}(t) = \emptyset$ or $\text{fv}(u) = \emptyset$. Note that a closed (named) term may contain free names. The sets of closed terms, closed values, and named terms are T^0 , V^0 , and U^0 , respectively. In any discussion or proof, we say a variable or a name is *fresh* if it does not occur in any term under consideration.

We distinguish several kinds of contexts, represented outside-in, as follows:

$$\begin{aligned} \text{Contexts:} \quad C &::= \square \mid C t \mid t C \mid \lambda x.C \mid \mu a.C \\ \text{Named contexts:} \quad \mathbb{C} &::= [a]C \\ \text{CBN evaluation contexts:} \quad E &::= \square \mid E t \\ \text{Named evaluation contexts:} \quad \mathbb{E} &::= [a]E \end{aligned}$$

The syntax of (named) evaluation contexts reflects the chosen reduction strategy, here call by name. Contexts can be filled only with a term t , to produce either regular terms $C[t]$, $E[t]$, or named terms $\mathbb{C}[t]$, $\mathbb{E}[t]$; the free names and free variables of t may be captured in the process.

Reduction semantics. We write $t_0\{t_1/x\}$ and $u_0\{t_1/x\}$ for the usual capture-avoiding substitution of t_1 for x in respectively t_0 and u_0 . We write $t\langle\mathbb{E}/a\rangle$ and $u\langle\mathbb{E}/a\rangle$ for the capture-avoiding substitution of \mathbb{E} for a in, respectively, t or u (here, each subterm of t or u that has the form $[a]s$ is replaced by $\mathbb{E}[s]$). We define the call-by-name reduction relation \rightarrow_n inductively by the following rules:

$$\begin{aligned} (\beta_n) \quad [a](\lambda x.t_0) t_1 &\rightarrow_n [a]t_0\{t_1/x\} \\ (\mu) \quad [a]\mu b.u &\rightarrow_n u\langle [a]\square/b \rangle \\ (app) \quad [a]t_0 t_1 &\rightarrow_n u\langle [a]\square t_1/b \rangle \text{ if } [b]t_0 \rightarrow_n u \text{ and } b \notin \text{fn}([a]t_0 t_1) \end{aligned}$$

Reduction is defined on named terms only. The rule (β_n) is the usual call-by-name β -reduction. In rule (μ) , the current continuation, represented by a , is captured and substituted for b in u . In an application (cf. rule (app)), we reduce the term t_0 in function position by introducing a fresh name b which represents the top level. We then replace b with $[a]\square t_1$ in the result u of the reduction of $[b]t_0$. We can also express reduction with top level evaluation contexts as follows.

Lemma 1. $u \rightarrow_n u'$ iff $u = \mathbb{E}[(\lambda x.t_0) t_1]$ and $u' = \mathbb{E}[t_0\{t_1/x\}]$, or $u = \mathbb{E}[\mu a.u'']$ and $u' = u''\langle\mathbb{E}/a\rangle$.

Reduction is also compatible with evaluation contexts in the following sense.

Lemma 2. If $u \rightarrow_n u'$, then $u\langle\mathbb{E}/a\rangle \rightarrow_n u'\langle\mathbb{E}/a\rangle$.

We write \rightarrow_n^* for the transitive and reflexive closure of \rightarrow_n , and we define the evaluation relation of the calculus as follows.

Definition 1. We write $u \Downarrow_n u'$ if $u \rightarrow_n^* u'$ and $u' \not\rightarrow_n$.

The result of the evaluation of a closed named term, if it exists, is a named value. If a named term u admits an infinite reduction sequence, we say it *diverges*, written $u \Uparrow_n$. For example, let $\Omega \stackrel{\text{def}}{=} (\lambda x.x x) (\lambda x.x x)$; then $[a]\Omega \Uparrow_n$ for all a .

2.2 Contextual Equivalence

As in the λ -calculus, contextual equivalence in the $\lambda\mu$ -calculus is defined in terms of convergence. However, unlike previous definitions [8,10], we define contextual equivalence on named terms first, before extending it to any terms.

Definition 2. *Two closed named terms u_0, u_1 are contextually equivalent, written $u_0 \approx_c u_1$, if for all closed \mathbb{C} , $\mathbb{C}[\mu a.u_0] \Downarrow_n [b]v_0$ iff $\mathbb{C}[\mu a.u_1] \Downarrow_n [b]v_1$.*

Note that we can plug only terms in a context, therefore we prefix u_0 and u_1 with a μ -abstraction. Definition 2 is not as generic as it could be, because we require the resulting named values to have the same top-level name b ; a more general definition would simply say “ $\mathbb{C}[\mu a.u_0] \Downarrow_n$ iff $\mathbb{C}[\mu a.u_1] \Downarrow_n$.” Our definition is strictly finer than the general one, because contexts cannot discriminate upon top-level names in some cases, as we can see with the next example.

Example 1. Let $\Theta \stackrel{\text{def}}{=} (\lambda x.\lambda y.y (x x y)) (\lambda x.\lambda y.y (x x y))$ be Turing’s CBN fixed-point combinator, and let $v \stackrel{\text{def}}{=} \lambda x.\lambda y.x$. The terms $u_0 \stackrel{\text{def}}{=} [a]\lambda x.\mu c.[b]\lambda y.\Theta v$ and $u_1 \stackrel{\text{def}}{=} [b]\lambda y.\Theta v$ are related by the general contextual equivalence, but not by Definition 2 if $a \neq b$. We verify that $\mathbb{E}[\mu c.u_0] \Downarrow_n$ iff $\mathbb{E}[\mu c.u_1] \Downarrow_n$ holds for all \mathbb{E} and c , and we can then conclude that u_0 and u_1 are in the general equivalence with David and Py’s context lemma [3]. Let \mathbb{E} be such that $\mathbb{E} \neq [d]\square$. Then one can check that $\mathbb{E}[\mu a.u_0] \Downarrow_n [b]\lambda y.\Theta v$ and $\mathbb{E}[\mu a.u_1] \Downarrow_n [b]\lambda y.\Theta v$, $\mathbb{E}[\mu b.u_0] \Downarrow_n [a]\lambda x.\mu c.\mathbb{E}[\lambda y.\Theta v]$ and $\mathbb{E}[\mu b.u_1] \Downarrow_n [b]\lambda y.\Theta v$, and finally $\mathbb{E}[\mu c.u_0] \Downarrow_n u_0$ and $\mathbb{E}[\mu c.u_1] \Downarrow_n u_1$ for $c \notin \{a, b\}$. The case $\mathbb{E} = [d]\square$ is easy to check as well.

We choose Definition 2 because it is easier to characterize with a bisimilarity than the general equivalence. Besides, only very peculiar terms u_0 and u_1 are related by the general equivalence but not by Definition 2: these terms reduce (in some context \mathbb{C}) to values $[a]v_0$ and $[b]v_1$ with $a \neq b$ that never evaluate their arguments. Indeed, if $\mathbb{E} = [c]\square t_0 \dots t_n$, then $\mathbb{E}[\mu a.[a]v_0] \rightarrow_n \mathbb{E}[v_0 \langle \mathbb{E}/a \rangle]$, and $\mathbb{E}[\mu a.[b]v_1] \Downarrow_n [b]v_1 \langle \mathbb{E}/a \rangle$. Suppose that when evaluating $\mathbb{E}[v_0 \langle \mathbb{E}/a \rangle]$, we evaluate one of the t_i . Then by replacing t_i with Ω , we obtain a context \mathbb{E}' such that $\mathbb{E}'[\mu a.[a]v_0] \uparrow_n$ (because Ω will be evaluated), and $\mathbb{E}'[\mu a.[b]v_1] \Downarrow_n$, which is in contradiction with the fact that u_0 and u_1 are in the general equivalence (take $\mathbb{E}'[\mu a.\mathbb{C}]$).

We extend Definition 2 to any terms t_0, t_1 , by saying that $t_0 \approx_c t_1$ if $[a]t_0 \approx_c [a]t_1$ for any fresh a . Other versions of the extension are possible, for example by replacing “for any a ” by “for some a ”, or by dropping the freshness requirement; as can be shown using the results of Section 2.3, all these definitions are equivalent. We can also define contextual equivalence on open terms, using the notion of *open extension*, which extends any relation on closed (named) terms to open (named) terms.

Definition 3. *Let \mathcal{R} be a relation on closed (named) terms. Two terms t_0 and t_1 are in the open extension of \mathcal{R} , written $t_0 \mathcal{R}^\circ t_1$, if for all substitutions σ closing t_0 and t_1 , we have $t_0\sigma \mathcal{R} t_1\sigma$ (and similarly for $u_0 \mathcal{R}^\circ u_1$).*

2.3 Applicative bisimilarity

Definition and properties. We propose a notion of applicative bisimulation, which tests values by applying them to a random argument. As with contextual equivalence, we give the definitions for named terms, before extending it to regular terms.

Definition 4. *A relation \mathcal{R} on closed named terms is an applicative bisimulation if $u_0 \mathcal{R} u_1$ implies*

- if $u_0 \rightarrow_n u'_0$, then there exists u'_1 such that $u_1 \rightarrow_n^* u'_1$ and $u'_0 \mathcal{R} u'_1$;
- if $u_0 = [a]\lambda x.t_0$, then there exists t_1 such that $u_1 \rightarrow_n^* [a]\lambda x.t_1$ and for all t , we have $[a]t_0 \langle [a]\square t/a \rangle \{t/x\} \mathcal{R} [a]t_1 \langle [a]\square t/a \rangle \{t/x\}$;
- the converse of the above conditions on u_1 .

Applicative bisimilarity, written \approx , is the largest applicative bisimulation.

For regular terms, we write $t_0 \mathcal{R} t_1$ if $[a]t_0 \mathcal{R} [a]t_1$ for any $a \notin \text{fn}(t_0, t_1)$. The first item of Definition 4 plays the bisimulation game for named terms which are not named values. If u_0 is a named value $[a]\lambda x.t_0$, then u_1 has to reduce to a named value $[a]\lambda x.t_1$, and we compare the values by applying them to an argument t . However, we cannot simply apply t directly to the values, we have to prefix them first with μa . As a result, we consider the named terms $[a](\mu a.[a]\lambda x.t_0) t$ and $[a](\mu a.[a]\lambda x.t_1) t$, which reduce to, respectively, $[a](\lambda x.t_0 \langle [a]\square t/a \rangle) t$ and $[a](\lambda x.t_1 \langle [a]\square t/a \rangle) t$, and then to $[a]t_0 \langle [a]\square t/a \rangle \{t/x\}$ and $[a]t_1 \langle [a]\square t/a \rangle \{t/x\}$; we obtain the terms in the clause for values of Definition 4.

Remark 1. When considering $[a](\mu a.[a]\lambda x.t_0) t$ and $[a](\mu a.[a]\lambda x.t_1) t$, we use the same top level name a as the one of the named values $[a]\lambda x.t_0$ and $[a]\lambda x.t_1$. We could use a fresh name b instead; reusing the same name makes the bisimulation proofs easier (we do not have to introduce unnecessary fresh names).

We can also define a big-step version of the bisimulation, where we consider only evaluation to a value.

Definition 5. *A relation \mathcal{R} on closed named terms is a big-step applicative bisimulation if $u_0 \mathcal{R} u_1$ implies*

- if $u_0 \rightarrow_n^* [a]\lambda x.t_0$, then there exists t_1 such that $u_1 \rightarrow_n^* [a]\lambda x.t_1$ and for all t , we have $[a]t_0 \langle [a]\square t/a \rangle \{t/x\} \mathcal{R} [a]t_1 \langle [a]\square t/a \rangle \{t/x\}$;
- the converse of the above condition on u_1 .

Lemma 3. *If \mathcal{R} is a big-step applicative bisimulation, then $\mathcal{R} \subseteq \approx$.*

As a first property, we prove that reduction (and therefore, evaluation) is included in bisimilarity.

Lemma 4. *We have $\rightarrow_n^* \subseteq \approx$.*

Proof. By showing that $\{(u, u') \mid u \rightarrow_n^* u'\}$ is a big-step bisimulation.

We give a basic example to show how applicative bisimulation can be used.

Example 2. For all closed v and $b \notin \text{fn}(v)$, we prove that $[a]v \approx [a]\lambda x.\mu b.[a]v$ by showing that $\{([a]v, [a]\lambda x.\mu b.[a]v) \mid b \notin \text{fn}(v)\} \cup \approx$ is an applicative bisimulation. Indeed, if $v = \lambda x.t$, then for all t' , we have $[a]t\{t'/x\} \approx [a]\mu b.[a]v t'$, because $[a]\mu b.[a]v t' \rightarrow_n^* [a]t\{t'/x\}$ (and by Lemma 4).

Soundness and completeness. We now prove that \approx coincides with \approx_c . We first show that \approx is a congruence using Howe's method [7,5], which is a classic proof method to show that an applicative bisimilarity is a congruence. As in [8], we need to slightly adapt the proof to the $\lambda\mu$ -calculus. Here we only sketch the application of the method, all the details can be found in Appendix A.1.

The principle of the method is to prove that a relation called the *Howe's closure* of \approx , which is a congruence by construction, is also a bisimulation. The definition of Howe's closure relies on an auxiliary relation, called the *compatible refinement* $\tilde{\mathcal{R}}$ of a relation \mathcal{R} , and inductively defined by the following rules:

$$\begin{array}{c} \frac{}{x \tilde{\mathcal{R}} x} \quad \frac{t_0 \mathcal{R} t_1}{\lambda x.t_0 \tilde{\mathcal{R}} \lambda x.t_1} \quad \frac{t_0 \mathcal{R} t_1 \quad t'_0 \mathcal{R} t'_1}{t_0 t'_0 \tilde{\mathcal{R}} t_1 t'_1} \quad \frac{u_0 \mathcal{R} u_1}{\mu a.u_0 \tilde{\mathcal{R}} \mu a.u_1} \\ \\ \frac{t_0 \mathcal{R} t_1}{[a]t_0 \tilde{\mathcal{R}} [a]t_1} \quad \frac{t_0 \mathcal{R} t_1 \quad \mathbb{E}_0 \tilde{\mathcal{R}} \mathbb{E}_1}{t_0 \langle \mathbb{E}_0/a \rangle \tilde{\mathcal{R}} t_1 \langle \mathbb{E}_1/a \rangle} \quad \frac{u_0 \mathcal{R} u_1 \quad \mathbb{E}_0 \tilde{\mathcal{R}} \mathbb{E}_1}{u_0 \langle \mathbb{E}_0/a \rangle \tilde{\mathcal{R}} u_1 \langle \mathbb{E}_1/a \rangle} \\ \\ \frac{}{\square \tilde{\mathcal{R}} \square} \quad \frac{E_0 \tilde{\mathcal{R}} E_1 \quad t_0 \mathcal{R} t_1}{E_0 t_0 \tilde{\mathcal{R}} E_1 t_1} \quad \frac{E_0 \tilde{\mathcal{R}} E_1}{[a]E_0 \tilde{\mathcal{R}} [a]E_1} \end{array}$$

In the original definition of compatible refinement [5], two terms are related by $\tilde{\mathcal{R}}$ if they have the same outer language constructor, and their direct subterms are related by \mathcal{R} . In the $\lambda\mu$ -calculus, compatible refinement is extended to (named) evaluation contexts, and we allow for the substitution of names with related named contexts.

We can now define Howe's closure of \approx , written \approx^\bullet , as follows.

Definition 6. *The Howe's closure \approx^\bullet is the smallest relation verifying:*

$$\approx^\circ \subseteq \approx^\bullet \quad \approx^\bullet \approx^\circ \subseteq \approx^\bullet \quad \tilde{\approx}^\bullet \subseteq \approx^\bullet$$

Howe's closure is defined on open (named) terms as well as on (named) evaluation contexts. Because it contains its compatible refinement, \approx^\bullet is a congruence. To prove it is a bisimulation, we need a stronger result, called a pseudo-simulation lemma, where we test named values not with the same argument, but with arguments t'_0, t'_1 related by \approx^\bullet .

Lemma 5. *Let $(\approx^\bullet)^c$ be \approx^\bullet restricted to closed terms, and let $u_0 (\approx^\bullet)^c u_1$.*

– *If $u_0 \rightarrow_n u'_0$, then $u_1 \rightarrow_n^* u'_1$ and $u'_0 (\approx^\bullet)^c u'_1$.*

- If $u_0 = [a]\lambda x.t_0$, then $u_1 \rightarrow_n^* [a]\lambda x.t_1$ and for all $t'_0 \approx^{\bullet c} t'_1$, we have $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \approx^{\bullet c} [a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$.

With this result, we can prove that $\approx^{\bullet c}$ is a bisimulation, and therefore included in \approx . Because it also contains \approx by definition, we have $\approx = (\approx^{\bullet c})^c$, and this implies that \approx is a congruence. As a result, \approx is sound w.r.t. to \approx_c .

Theorem 1. *We have $\approx \subseteq \approx_c$.*

To prove completeness (the reverse inclusion), we use an alternate definition of contextual equivalence, where we test terms with named evaluation contexts. By doing so, we prove a context lemma in the process.¹

Definition 7. *Let u_0, u_1 be closed terms. We write $u_0 \approx_c u_1$ if for all closed contexts \mathbb{E} and names a , we have $\mathbb{E}[\mu a.u_0] \Downarrow_n [b]v_0$ iff $\mathbb{E}[\mu a.u_1] \Downarrow_n [b]v_1$.*

Theorem 2. *We have $\approx_c \subseteq \approx_c \subseteq \approx$.*

The first inclusion is by definition, and the second one is by showing that \approx_c is a big-step applicative bisimulation.

Comparison with Lassen's work. In [8], Lassen also proposes a definition of applicative bisimilarity that he proves sound, but he conjectures that it is not complete. We discuss here the differences between the two approaches.

Lassen defines a notion of bisimulation for regular terms only, and not for named terms. The definition is as follows.

Definition 8. *A relation \mathcal{R} on closed terms is a Lassen applicative bisimulation if $t_0 \mathcal{R} t_1$ implies:*

- for all a , if $[a]t_0 \rightarrow_n^* [b]\lambda x.t'_0$, then there exists t'_1 such that $[a]t_1 \rightarrow_n^* [b]\lambda x.t'_1$, and for all t , we have $t_0 \langle [b]\square t/b \rangle \{t/x\} \mathcal{R} t_1 \langle [b]\square t/b \rangle \{t/x\}$;
- the converse of the above condition on t_1 .

Lassen's definition is quite similar to our definition of big-step applicative bisimulation (Definition 5), except it requires $t_0 \langle [b]\square t/b \rangle \{t/x\} \mathcal{R} t_1 \langle [b]\square t/b \rangle \{t/x\}$, which implies that these terms must be related when reduced with any top level name a . This is more restrictive than our definition, where we compare these terms only with the top level name b . To illustrate the difference, we consider Lassen's counter-example from [8].

Example 3. Let $t_0 \stackrel{\text{def}}{=} (\lambda x.\lambda y.x x) (\lambda x.\lambda y.x x)$, and $t_1 \stackrel{\text{def}}{=} \mu a.[a]\lambda y.\mu c.[a]t_0$ (with $c \neq a$). These terms are not bisimilar according to Lassen's definition. For all b , we have $[b]t_0 \rightarrow_n^* [b]\lambda y.t_0$ and $[b]t_1 \rightarrow_n^* [b]\lambda y.\mu c.[b]t_0$. With Lassen's definition, one has to relate t_0 and $\mu c.[b]t_0 t$ for any t , which means comparing $[d]t_0$ and $[d]\mu c.[b]t_0 t$ for all d . But these two terms are not equivalent if $d \neq b$.

¹ We cannot directly use David and Py's context lemma [3], because we use a different notion of contextual equivalence.

Lassen conjectures in [8] that these terms are contextually equivalent, and we can indeed prove that they are (big-step) applicative bisimilar with our definition: we just have to compare $[b]t_0$ and $[b]\mu a'.[b]t_0 t$ for any t , and both terms evaluate to $[b]\lambda x.t_0$ and are therefore equivalent.

By comparing primarily named terms, as we do in our definition, we can keep track of what happens to the top level, and especially of any connection between the top level and a subterm. In Example 3, we can see that it is essential to remember that b represents the top level in $\mu c.[b]t_0 t$, and therefore it does not make sense to compare $[d]t_0$ and $[d]\mu c.[b]t_0 t$ for any $d \neq b$, as we have to do with Lassen's definition. We believe comparing named terms is essential to obtain completeness w.r.t. contextual equivalence; note that the sound and complete normal form bisimilarity for the $\lambda\mu\rho$ -calculus [18] is also defined on named terms.

Example. In [3], David and Py give a counter-example showing that Böhm's theorem fails in CBN $\lambda\mu$ -calculus. They prove their terms are contextually equivalent with a context lemma. Here we slightly simplify their counter-example, and prove equivalence using applicative bisimilarity. Note that these terms cannot be proved equivalent with (a CBN variant of) eager normal form bisimilarity [8](papers defining enf for lambda mu).

Example 4. Let $0 \stackrel{\text{def}}{=} \lambda x.\lambda y.y$, $1 \stackrel{\text{def}}{=} \lambda x.\lambda y.x$, and $t_a \stackrel{\text{def}}{=} \mu c.[a]0$. Then we have $\lambda x.\mu a.[a]x \mu b.[a]x t_a 0 \approx \lambda x.\mu a.[a]x \mu b.[a]x t_a 1$.²

Proof. We fix a name c , and for all t , we want to relate $[c]\mu a.[a]t \mu b.[a]t t_a 0$ and $[c]\mu a.[a]t \mu b.[a]t t_a 1$, which reduce respectively to $[c]t \mu b.[c]t t_c 0$ and $[c]t \mu b.[c]t t_c 1$. Let $s_0^t \stackrel{\text{def}}{=} \mu b.[c]t t_c 0$, $s_1^t \stackrel{\text{def}}{=} \mu b.[c]t t_c 1$, $\mathbb{E}_0^t \stackrel{\text{def}}{=} [c]\square s_0^t$, and $\mathbb{E}_1^t \stackrel{\text{def}}{=} [c]\square s_1^t$. We define a relation $[d]t \rightsquigarrow^k u$ as $[d]t \rightarrow_n^* u$ if $k = 0$ and as $[d]t \Downarrow_n [d]\lambda x_1.t_1, [c]t_1 \Downarrow_n [d]\lambda x_2.t_2, \dots [c]t_k \rightarrow_n^* u$ for some $t_1 \dots t_k$ if $k > 0$. The rationale behind this relation appears in the proof. Note that if $[d]t \rightsquigarrow^k u$, then $\text{fv}(u) \subseteq \{x_1, \dots, x_k\}$, and $[d]t \langle \mathbb{E}/e \rangle \rightsquigarrow^k u \langle \mathbb{E}/e \rangle$ for all \mathbb{E} and $e \neq d$. We define \mathcal{R} as

$$\begin{aligned} & \{(u \langle \mathbb{E}_0^t/d \rangle \{s_0^t/x_i\}_{i=1}^k, u \langle \mathbb{E}_1^t/d \rangle \{s_1^t/x_i\}_{i=1}^k) \mid u \in U, t \in T^0, d \notin \text{fn}(t), [d]t \rightsquigarrow^k u\} \\ & \cup \{([c]t t_c 0, [c]t t_c 1) \mid t \in T^0, d \notin \text{fn}(t), [d]t \rightsquigarrow^k \mathbb{E}[x_i], k \geq 1, 1 \leq i \leq k\} \\ & \cup \{(u, u) \mid u \in U^0\} \end{aligned}$$

where $t\{t_i/x_i\}_{i=1}^k$ stands for $t\{t_1/x_1\} \dots \{t_k/x_k\}$, and we show that \mathcal{R} is an applicative bisimulation.

Let $u \langle \mathbb{E}_0^t/d \rangle \{s_0^t/x_i\}_{i=1}^k \mathcal{R} u \langle \mathbb{E}_1^t/d \rangle \{s_1^t/x_i\}_{i=1}^k$. If $u \rightarrow_n u'$, then it is easy to conclude. Otherwise, we distinguish several cases.

² The terms David and Py consider in their work are $\lambda x.\mu a.[a]x \mu b.[a](x t_a 0) t_a$ and $\lambda x.\mu a.[a]x \mu b.[a](x t_a 1) t_a$. However, the additional argument t_a would not come into play in the proof we present, so we have elided it.

If $u = [e]\lambda y.t_0$ with $e \neq d$, then $[e]t_0 \langle [e]\square t'/e \rangle \{t'/y\} \langle \mathbb{E}_0^{t''}/d \rangle \{s_0^{t''}/x_i\}_{i=1}^k \mathcal{R}$ $[e]t_0 \langle [e]\square t'/e \rangle \{t'/y\} \langle \mathbb{E}_1^{t''}/d \rangle \{s_1^{t''}/x_i\}_{i=1}^k$ (where $t'' = t \langle [e]\square t'/e \rangle$) for all t' , because the relation $[d]t \rightsquigarrow^k [e]\lambda y.t_0$ implies $[d]t'' \rightsquigarrow^k [e]t_0 \langle [e]\square t'/e \rangle \{t'/y\}$.

If $u = [d]\lambda x_{k+1}.t_{k+1}$, then we have the reduction $u \langle \mathbb{E}_0^t/d \rangle \{s_0^t/x_i\}_{i=1}^k \rightarrow_n [c]t_{k+1} \{s_0^t/x_{k+1}\} \langle \mathbb{E}_0^t/d \rangle \{s_0^t/x_i\}_{i=1}^k$ as well as the reduction $u \langle \mathbb{E}_1^t/d \rangle \{s_1^t/x_i\}_{i=1}^k \rightarrow_n [c]t_{k+1} \{s_1^t/x_{k+1}\} \langle \mathbb{E}_1^t/d \rangle \{s_1^t/x_i\}_{i=1}^k$. We obtain terms in \mathcal{R} because we can rewrite them into, respectively, $[c]t_{k+1} \langle \mathbb{E}_0^t/d \rangle \{s_0^t/x_i\}_{i=1}^{k+1}$ and $[c]t_{k+1} \langle \mathbb{E}_1^t/d \rangle \{s_1^t/x_i\}_{i=1}^{k+1}$, and $[d]t \rightsquigarrow^{k+1} [c]t_{k+1}$ holds.

Finally, if $u = \mathbb{E}[x_i]$, for $1 \leq i \leq k$ (assuming $k \geq 1$), then we have $u \langle \mathbb{E}_0^t/d \rangle \{s_0^t/x_i\}_{i=1}^k \rightarrow_n [c]t t_c 0$ and $u \langle \mathbb{E}_1^t/d \rangle \{s_1^t/x_i\}_{i=1}^k \rightarrow_n [c]t t_c 1$, and $[c]t t_c 0 \mathcal{R} [c]t t_c 1$ holds.

Let $[c]t t_c 0 \mathcal{R} [c]t t_c 1$ with $[d]t \rightsquigarrow^k \mathbb{E}[x_i]$. There exist t_1, \dots, t_k such that $[d]t \Downarrow_n [d]\lambda x_1.t_1$, $[c]t_1 \Downarrow_n [d]\lambda x_2.t_2$, \dots , $[c]t_{k-1} \Downarrow_n [d]\lambda x_k.t_k$, and $[c]t_k \rightarrow_n^* \mathbb{E}[x_i]$. Then $[c]t t_c 0 \rightarrow_n^* [c](\lambda x_1.t_1 \langle [c]\square t_c 0/d \rangle) t_c 0 \rightarrow_n [c]t_1 \langle [c]\square t_c 0/d \rangle \{t_c/x_1\} 0 \rightarrow_n^* [c]t_2 \langle [c]\square t_c 0/d \rangle \langle [c]\square 0/c \rangle \{t_c/x_i\}_{i=1}^2 0 \rightarrow_n^* [c]t_k \langle [c]\square t_c 0/d \rangle \langle [c]\square 0/c \rangle \{t_c/x_i\}_{i=1}^k 0 \rightarrow_n^* \mathbb{E}'[t_c] \rightarrow_n [c]0$ with $\mathbb{E}' \stackrel{\text{def}}{=} \mathbb{E} \langle [c]\square t_c 0/d \rangle \langle [c]\square 0/c \rangle \{t_c/x_i\}_{i=1}^k$. Similarly $[c]t t_c 1 \rightarrow_n^* \mathbb{E}''[t_c] \rightarrow_n [c]0$ with $\mathbb{E}'' \stackrel{\text{def}}{=} \mathbb{E} \langle [c]\square t_c 1/d \rangle \langle [c]\square 1/c \rangle \{t_c/x_i\}_{i=1}^k$. Because the two terms evaluate to $[c]0$, it is easy to conclude.

2.4 Environmental bisimilarity

Definition and basic properties. Environmental bisimilarity [16] uses environments (denoted by \mathcal{E}) to accumulate knowledge when comparing terms. Like applicative bisimilarity, it challenges values by passing them arguments, except the arguments are built from the environments instead of being completely arbitrary. In $\lambda\mu$ -calculus, we define an environment \mathcal{E} as a relation on named closed values; for example, the identity environment \mathcal{I} is defined as $\{([a]v, [a]v) \mid a \in A, v \in V^0\}$. To build arguments from \mathcal{E} , we use the notion of *context closure*, defined as follows: given a relation \mathcal{R} on (named) terms, its context closure $\widehat{\mathcal{R}}$ is the smallest relation verifying $\mathcal{R} \subseteq \widehat{\mathcal{R}}$ and $\widetilde{\mathcal{R}} \subseteq \widehat{\mathcal{R}}$. Like compatible refinement, context closure is defined on open (named) terms and evaluation contexts, but we restrict it to closed (named) terms and contexts, unless specified otherwise. Note that, unlike usual definitions of context closure, our definition requires context closure to be stable by substitution of a name by a named evaluation context.

An environmental relation \mathcal{X} is a set of environments \mathcal{E} and triples (\mathcal{E}, u_0, u_1) , where u_0 and u_1 are closed. We write $u_0 \mathcal{X}_{\mathcal{E}} u_1$ for $(\mathcal{E}, u_0, u_1) \in \mathcal{X}$; it means that u_0 and u_1 are compared within the current knowledge \mathcal{E} . We define environmental bisimulation as follows.

Definition 9. *An environmental relation \mathcal{X} is an environmental bisimulation if $u_0 \mathcal{X}_{\mathcal{E}} u_1$ implies*

- if $u_0 \rightarrow_n u'_0$, then there exists u'_1 such that $u_1 \rightarrow_n^* u'_1$ and $u'_0 \mathcal{X}_{\mathcal{E}} u'_1$;
- if $u_0 = [a]v_0$, then there exists v_1 such that $u_1 \rightarrow_n^* [a]v_1$ and $\{([a]v_0, [a]v_1)\} \cup \mathcal{E} \in \mathcal{X}$;

– the converse of the above condition on u_1 .

Furthermore, if $\mathcal{E} \in \mathcal{X}$ and $[a]\lambda x.t_0 \mathcal{E} [a]\lambda x.t_1$, then for all $t'_0 \widehat{\mathcal{E}} t'_1$, we have $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \mathcal{X}_{\mathcal{E}} [a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$.

Environmental bisimilarity \cong is the largest environmental bisimulation. To prove the equivalence between u_0 and u_1 , we want to relate them without any predefined knowledge, i.e., we want to prove that $u_0 \cong_{\emptyset} u_1$ holds; we also write \simeq for \cong_{\emptyset} . We extend these definitions to regular terms by saying that $t_0 \mathcal{X}_{\mathcal{E}} t_1$ holds if we have $[a]t_0 \mathcal{X}_{\mathcal{E}} [a]t_1$ for any $a \notin \text{fn}(t_0, t_1)$.

Definition 9 is divided into two parts, the first one making explicit the bisimulation game for named terms, and the second one focusing on environments. If u_0 is a named value $[a]v_0$, then u_1 has to evaluate to named value with the same top-level name $[a]v_1$, and \mathcal{E} extended with the newly acquired knowledge $([a]v_0, [a]v_1)$ must belong to \mathcal{X} . We then compare two named values $[a]\lambda x.t_0$ and $[a]\lambda x.t_1$ in \mathcal{E} by passing them two terms t'_0 and t'_1 built from \mathcal{E} . As in the applicative bisimulation definition, we have to prefix the named values with μa first, and we therefore have to relate $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\}$ and $[a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$.

It is possible to define a big-step variant of environmental bisimulation by removing the first item of Definition 9 and by changing equality in the value case by \rightarrow_n^* . We also use the following properties.

Lemma 6 (Weakening). *If $u_0 \cong_{\mathcal{E}} u_1$ and $\mathcal{E}' \subseteq \mathcal{E}$ then $u_0 \cong_{\mathcal{E}'} u_1$.*

A smaller environment is a weaker constraint, because we can build less arguments to test the named values in \mathcal{E} . The proof is as in [16]. Like with applicative bisimilarity, reduction and evaluation are included in \simeq .

Lemma 7. *If $u \rightarrow_n^* u'$, then $u \simeq u'$.*

Bisimulation proofs with environmental bisimulation are usually harder than with applicative bisimulation, as we can see with the next example.

Example 5. Following Example 2, we want to prove $[a]v \simeq [a]\lambda x.\mu b.[a]v$ for all closed v and $b \notin \text{fn}(v)$. To this end, we start with $\mathcal{X} \stackrel{\text{def}}{=} \{(\emptyset, [a]v, [a]\lambda x.\mu b.[a]v) \mid a \in A, v \in V^0, b \notin \text{fn}(v)\} \cup \{\mathcal{E} \mid \mathcal{E} \subseteq \mathcal{E}\}$, where $\mathcal{E} = \{([a]v, [a]\lambda x.\mu b.[a]v) \mid a \in A, v \in V^0, b \notin \text{fn}(v)\}$. Let $[a]v \mathcal{E} [a]\lambda x.\mu b.[a]v$, where $v = \lambda x.t$. For all $t_0 \widehat{\mathcal{E}} t_1$, we want $[a]t \{t_0/x\} \mathcal{X}_{\mathcal{E}} [a]\mu b.[a]v t_1$ to hold. In contrast with Example 2, $[a]\mu b.[a]v t_1$ does not reduce to $[a]t \{t_0/x\}$; to conclude the proof easily, we need some up-to techniques, introduced later (see Example 6).

Soundness and completeness. Here we sketch the soundness and completeness proofs for \simeq , which are adaptations of the proofs for the λ -calculus [16]. More details can be found in Appendix A.2. The first step is to define some basic up-to techniques, specifically bisimulation up to environment (which allows for bigger environments in the bisimulation clauses) and up to bisimilarity (where \simeq is used at some specific points in the bisimulation clauses); the definitions and proofs of soundness mimic the ones for the λ -calculus [16].

We then prove congruence w.r.t. evaluation contexts. Given a relation \mathcal{R} on named terms, we write \mathcal{R}^{nv} for its restriction to named values. Let \mathcal{Y} be an environmental bisimulation. Then the relation

$$\begin{aligned} & \{(\widehat{\mathcal{E}}^{\text{nv}}, u_0 \langle \mathbb{E}_0^0/a_0 \rangle \dots \langle \mathbb{E}_0^n/a_n \rangle, u_1 \langle \mathbb{E}_1^0/a_0 \rangle \dots \langle \mathbb{E}_1^n/a_n \rangle) \mid u_0 \mathcal{Y}_{\mathcal{E}} u_1, \mathbb{E}_0^i \widehat{\mathcal{E}} \mathbb{E}_1^i\} \\ & \cup \{(\widehat{\mathcal{E}}^{\text{nv}}, u_0, u_1) \mid \mathcal{E} \in \mathcal{Y}, u_0 \widehat{\mathcal{E}} u_1\} \cup \{\widehat{\mathcal{E}}^{\text{nv}} \mid \mathcal{E} \in \mathcal{Y}\} \end{aligned}$$

is an environmental bisimulation up to environment. From that, we can deduce the following lemma.

Lemma 8. *If $u_0 \cong_{\mathcal{E}} u_1$, then $u_0 \langle \mathbb{E}/a \rangle \cong_{\mathcal{E}} u_1 \langle \mathbb{E}/a \rangle$.*

Next, we prove that \simeq is a congruence by showing that $\{(\widehat{\simeq}^{\text{nv}}, u_0, u_1) \mid u_0 \widehat{\simeq} u_1\} \cup \{\widehat{\simeq}^{\text{nv}}\}$ is an environmental bisimulation up to bisimilarity.

Lemma 9. *If $u_0 \simeq u_1$, then $C[u_0] \widehat{\simeq}^{\text{nv}} C[u_1]$ and $\mathbb{C}[u_0] \widehat{\simeq}^{\text{nv}} \mathbb{C}[u_1]$.*

Using congruence, we can prove that \simeq is sound w.r.t. \approx_c , and completeness is proved using Definition 7, as in the case for applicative bisimulations.

Theorem 3. *We have $\simeq = \approx_c$.*

Up-to techniques. As witnessed by Example 5, proving the equivalence of two terms with Definition 9 alone is usually too difficult. To ease the proofs, it is common to rely on up-to techniques. We already mentioned bisimulation up to bisimilarity and up to environment; other techniques include bisimulation up to reduction, up to expansion, and up to context. We only discuss the last for space reasons, the others are defined and proved sound as in [16].

Bisimulation up to context allows to factor out a context common to both terms. Formally, we define the context closure of \mathcal{X} , written $\overline{\mathcal{X}}$, as follows: we have $u_0 \overline{\mathcal{X}}_{\mathcal{E}} u_1$ if

- either $u_0 = u'_0 \langle \mathbb{E}_0/a \rangle$, $u_1 = u'_1 \langle \mathbb{E}_1/a \rangle$ with $u'_0 \mathcal{X}_{\mathcal{E}} u'_1$ and $\mathbb{E}_0 \widehat{\mathcal{E}} \mathbb{E}_1$
- or $u_0 \widehat{\mathcal{E}} u_1$.

Terms related by $\overline{\mathcal{X}}_{\mathcal{E}}$ are either built from \mathcal{E} (second case), or we can decompose them into terms related by $\mathcal{X}_{\mathcal{E}}$ and evaluation contexts built from \mathcal{E} . We then define bisimulation up to context as follows.

Definition 10. *A relation \mathcal{X} is an environmental bisimulation up to context if $u_0 \mathcal{X}_{\mathcal{E}} u_1$ implies*

- if $u_0 \rightarrow_n u'_0$, then there exists u'_1 such that $u_1 \rightarrow_n^* u'_1$ and $u'_0 \overline{\mathcal{X}}_{\mathcal{E}} u'_1$;
- if $u_0 = [a]v_0$, then there exist v_1 and $\mathcal{E}' \subseteq \mathcal{X}$ such that $u_1 \rightarrow_n^* [a]v_1$ and $\{([a]v_0, [a]v_1)\} \cup \mathcal{E} \subseteq \widehat{\mathcal{E}}'$;
- the converse of the above condition on u_1 .

Furthermore, if $\mathcal{E} \in \mathcal{X}$ and $[a]\lambda x.t_0 \mathcal{E} [a]\lambda x.t_1$, then for all $t'_0 \widehat{\mathcal{E}} t'_1$, we have $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \overline{\mathcal{X}}_{\mathcal{E}} [a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$.

Lemma 10. *If \mathcal{X} is a bisimulation up to context, then $\mathcal{X} \subseteq \widehat{\cong}$.*

The proof of soundness is as in [16]. We now give some examples showing how bisimulation up to context (combined with other up-to techniques) can be used to simplify proofs.

Example 6. Continuing Example 5, we can prove that the candidate relation \mathcal{X} defined in this example is a bisimulation up to context up to reduction (meaning that we can do some reduction steps to get terms related by $\widehat{\mathcal{X}}$). Indeed, we remind that we want to relate $[a]t\{t_0/x\}$ and $[a]\mu b.[a]v t_1$ for all $t_0 \widehat{\mathcal{E}} t_1$, $v = \lambda x.t$, and $b \notin \text{fn}(v)$. But $[a]\mu b.[a]v t_1 \rightarrow_n^2 [a]t\{t_1/x\}$, and $[a]t\{t_0/x\} \widehat{\mathcal{E}} [a]t\{t_1/x\}$ holds, which implies $[a]t\{t_0/x\} \widehat{\mathcal{X}} [a]t\{t_1/x\}$, as wished.

3 Call-by-Value $\lambda\mu$ -calculus

We now discuss the behavioral theory of the call-by-value (CBV) $\lambda\mu$ -calculus and point out the differences with call-by-name.

3.1 Semantics and Contextual Equivalence

Semantics. In this section, we use CBV left-to-right evaluation, which is encoded in the syntax of the CBV evaluation contexts:

$$E ::= \square \mid E t \mid v E$$

CBV reduction \rightarrow_v is defined by the following rules.

$$\begin{aligned} (\beta_v) \quad & [a](\lambda x.t) v \rightarrow_v [a]t\{v/x\} \\ (\mu) \quad & [a]\mu b.u \rightarrow_v u\langle [a]\square/b \rangle \\ (app) \quad & [a]t_0 t_1 \rightarrow_v u\langle [a]\square t_1/b \rangle \text{ if } [b]t_0 \rightarrow_v u \text{ and } b \notin \text{fn}([a]t_0 t_1) \\ (app_v) \quad & [a]v t \rightarrow_v u\langle [a]v \square/b \rangle \text{ if } [b]t \rightarrow_v u \text{ and } b \notin \text{fn}([a]v t) \end{aligned}$$

With rule (app_v) , we reduce arguments to values, to be able to apply CBV β -reduction (rule (β_v)). The rules (μ) and (app) are the same as in CBN. We could also express reduction with top-level named evaluation contexts, as in Lemma 1. Furthermore, CBV reduction is compatible with CBV contexts, as in Lemma 2. We write \rightarrow_v^* for the reflexive and transitive closure of \rightarrow_v , \Downarrow_v for CBV evaluation, and \Uparrow_v for CBV divergence.

Contextual equivalence. We use the same definition as in CBN.

Definition 11. *Two closed named terms u_0 , u_1 are contextually equivalent, written $u_0 \approx_c u_1$, if for all closed \mathbb{C} , $\mathbb{C}[\mu a.u_0] \Downarrow_v [b]v_0$ iff $\mathbb{C}[\mu a.u_1] \Downarrow_v [b]v_1$.*

However, unlike in CBN, this definition (where we require the resulting values to have the same top-level names) coincides with the general definition where we simply say “ $\mathbb{C}[\mu a.u_0] \Downarrow_v$ iff $\mathbb{C}[\mu a.u_1] \Downarrow_v$.” Indeed, if $\mathbb{C}[\mu a.u_0] \Downarrow_v [b]v_0$ and $\mathbb{C}[\mu a.u_1] \Downarrow_v [c]v_1$ with $c \neq b$, then we can easily distinguish them, because $[b]\mu b.\mathbb{C}[\mu a.u_0] \Omega \rightarrow_v^* [b]v_0\langle [b]\square \Omega/b \rangle \Omega \Uparrow_v$, and $[b]\mu b.\mathbb{C}[\mu a.u_1] \Omega \Downarrow_v [c]v_1\langle [b]\square \Omega/b \rangle$.

3.2 Applicative Bisimilarity

We propose an applicative bisimilarity which equates contextual equivalence. We first give its definition and explain the rationale behind it, and then sketch the soundness proof.

Definition. We define applicative bisimulation in CBV as follows.

Definition 12. A relation \mathcal{R} on closed named terms is an applicative bisimulation if $u_0 \mathcal{R} u_1$ implies

- if $u_0 \rightarrow_v u'_0$, then there exists u'_1 such that $u_1 \rightarrow_v^* u'_1$ and $u'_0 \mathcal{R} u'_1$;
- if $u_0 = [a]\lambda x.t_0$, then there exists t_1 such that $u_1 \rightarrow_v^* [a]\lambda x.t_1$, and:
 1. for all t, b, v such that $[b]t \rightarrow_v^* [b]v$ and $b \notin \text{fn}(t)$, we have

$$[a]t_0 \langle [a]\square t/a \rangle \{v \langle [a]\lambda x.t_0 \langle [a]\square t/a \rangle \square/b \rangle /x\} \mathcal{R} [a]t_1 \langle [a]\square t/a \rangle \{v \langle [a]\lambda x.t_1 \langle [a]\square t/a \rangle \square/b \rangle /x\};$$

2. for all $v = \lambda x.t$, we have

$$[a]t \{ \lambda x.t_0 \langle [a]v \square/a \rangle /x \} \mathcal{R} [a]t \{ \lambda x.t_1 \langle [a]v \square/a \rangle /x \};$$

- the converse of the above conditions on u_1 .

Applicative bisimilarity, written \approx , is the largest applicative bisimulation.

The definition is extended to regular terms t_0, t_1 as in CBN, by using a fresh top-level name a . As we can see, we have two tests to perform in the value case (instead of one in CBN). In fact, each test corresponds to an elementary evaluation context: clause 1 checks contexts of the form $[a]\square t$, while clause 2 is about contexts of the form $[a]v \square$. As in CBN, we prefix the named values $[a]\lambda x.t_0$ and $[a]\lambda x.t_1$ with μa first to be able to test them.

When testing with $[a]\square t$, we consider only terms t such that $[b]t \rightarrow_v^* [b]v$ for a fresh b ; we explain why we make such a restriction. Testing with $[a]\square t$ requires us to look at the possible behaviors of $[a](\lambda x.t_0 \langle [a]\square t/a \rangle) t$ and $[a](\lambda x.t_1 \langle [a]\square t/a \rangle) t$. If $[b]t$ diverges, then these two terms diverge as well, and we gain no information on $[a]\lambda x.t_0$ and $[a]\lambda x.t_1$ themselves. If $[b]t \rightarrow_v^* [c]v$ with $b \neq c$, then $[a](\lambda x.t_0 \langle [a]\square t/a \rangle) t \rightarrow_v^* [c]v \langle [a]\lambda x.t_0 \langle [a]\square t/a \rangle \square/b \rangle$ (and similarly with $[a]\lambda x.t_1$). The values $[a]\lambda x.t_0$ and $[a]\lambda x.t_1$ are captured by $[b]t$, and no interaction between t and the two named values takes place in the process ($[a]\lambda x.t_0$ and $[a]\lambda x.t_1$ are not applied to any value); again, we do not gain any new knowledge on the behavior of $[a]\lambda x.t_0$ and $[a]\lambda x.t_1$. Finally, if $[b]t \rightarrow_v^* [b]v$, then $[a](\lambda x.t_0 \langle [a]\square t/a \rangle) t \rightarrow_v^* [b](\lambda x.t_0 \langle [a]\square t/a \rangle) v \langle [a]\lambda x.t_0 \langle [a]\square t/a \rangle \square/b \rangle$, and similarly with $[a]\lambda x.t_1$; in this case, a value is indeed passed to $[a]\lambda x.t_0$ and $[a]\lambda x.t_1$, and we can compare their respective behaviors. Therefore, an interaction happens between t and the tested

values only if $[b]t \rightarrow_v^* [b]v$, and the results of the interaction (after β -reduction) are the two terms we build in clause 1 of Definition 12.

In clause 2, we test $[a]\lambda x.t_0$ and $[a]\lambda x.t_1$ with $[a]v \square$, i.e., we look at the possible behaviors of $[a]v \lambda x.t_0 \langle [a]v \square / a \rangle$ and $[a]v \lambda x.t_1 \langle [a]v \square / a \rangle$. If $v = \lambda x.t$, then these two terms reduce in one β -reduction step into $[a]t \{ \lambda x.t_0 \langle [a]v \square / a \rangle / x \}$, and $[a]t \{ \lambda x.t_1 \langle [a]v \square / a \rangle / x \}$, which are the two terms we consider in clause 2. Note that this implies that a bisimulation \mathcal{R} is a congruence w.r.t. (regular) values; indeed, if $v_0 \mathcal{R} v_1$, then $[a]v_0 \mathcal{R} [a]v_1$ for a fresh a , and so we have $[a]t \{ v_0 / x \} \mathcal{R} [a]t \{ v_1 / x \}$ for all t (by clause 2). This is an unusual requirement on bisimulations which makes them harder to build, but we will provide a counter-example later on showing that clause 2 is necessary to have completeness of the bisimilarity.

As in CBN, we can define a big-step version of the bisimulation (where we use evaluation instead of reduction), and bisimilarity contains reduction.

Lemma 11. *We have $\rightarrow_v^* \subseteq \approx$.*

Example 7. We show that the applicative bisimulation for CBV is more difficult to use than the one for CBN by considering the terms of Example 2. Let $v = \lambda x.t$ be a closed value and $b \notin \text{fn}(v)$; then $[a]v \approx [a]\lambda x.\mu b.[a]v$. For clause 1, let t' be such that $[b]t' \rightarrow_v^* [b]v'$ for $b \notin \text{fn}(t')$; we have to compare $[a]t \{ v' \langle [a]v \square / b \rangle / x \}$ with $[a]\mu b.[a]v t'$. But $[a]\mu b.[a]v t' \rightarrow_v [a]v t' \rightarrow_v^* [a]t \{ v' \langle [a]v \square / b \rangle / x \}$, therefore we can conclude with Lemma 11.

For clause 2, we have to relate $[a]t' \{ v / y \}$ and $[a]t' \{ \lambda x.\mu b.[a]v' v / y \}$ for all $v' = \lambda y.t'$. We proceed by case analysis on t' ; the most interesting case is $t' = E[y v'']$. In this case, we have $[a]t' \{ \lambda x.\mu b.[a]v' v / y \} \rightarrow_v^* [a]v' v \rightarrow_v [a]t' \{ v / y \}$, therefore we can conclude with Lemma 11. To handle all the possible cases, we prove in Appendix B.1 that $\{ (u \{ v / y \}, u \{ \lambda x.\mu b.[a]t_0 / y \}) \mid [a]t_0 \rightarrow_v^* u \{ v / y \} \} \cup \approx$ is an applicative bisimulation.

Soundness and completeness. We briefly sketch the proof of soundness, more details can be found in Appendix B.2. The application of the Howe's method is easier than in CBN because, as already pointed out, an applicative bisimulation (and, therefore, the applicative bisimilarity) is already a congruence for regular values by definition. What is left to prove is congruence for (named) terms. We use the same definitions of compatible refinement and Howe's closure \approx^\bullet as in CBN. However, because \approx is a congruence for values, we can prove directly that the restriction of \approx^\bullet to closed terms (written $(\approx^\bullet)^c$) is an applicative bisimulation, without having to prove a pseudo-simulation lemma (similar to Lemma 5) beforehand.

Lemma 12. *The relation $(\approx^\bullet)^c$ is an applicative bisimulation.*

As in CBN, we can conclude that $(\approx^\bullet)^c = \approx$, and therefore \approx is a congruence. We can then deduce that \approx is sound w.r.t. \approx_c . For the reverse inclusion, we use an alternate definition of contextual equivalence where we test terms with evaluation contexts (see Definition 7), and we prove it is an applicative bisimulation. As a result, \approx coincides with \approx_c .

Theorem 4. *We have $\approx = \approx_c$.*

3.3 Counter-examples

As we can see with Example 7, applicative bisimulation is harder to use in CBV even for a basic example. We give counter-examples explaining why we cannot simplify its definition.

Example 8. In the CBV λ -calculus (and also with delimited control [2]), it is enough to test λ -abstractions with a value to have completeness. In the CBV $\lambda\mu$ -calculus, one could think that testing with contexts of the form $[a]\square v$ is also enough. It is not the case, as shown by the following counter-example.

Let $v_0 \stackrel{\text{def}}{=} \lambda x.\mu b.[a]w x$, $v_1 \stackrel{\text{def}}{=} \lambda x.w x x$, with $w \stackrel{\text{def}}{=} \lambda y.\lambda z.z y$. Then we have $([a]v_0)\langle [a]\square v/a \rangle = [a](\lambda x.\mu b.[a]w x v) v \rightarrow_v^* [a]w v v$ and $([a]v_1)\langle [a]\square v/a \rangle = [a](\lambda x.w x x) v \rightarrow_v^* [a]w v v$. Because they reduce to the same term, we can prove that $([a]v_0)\langle [a]\square v/a \rangle$ is contextually equivalent to $([a]v_1)\langle [a]\square v/a \rangle$.

Let $t \stackrel{\text{def}}{=} \mu d.[d]\lambda y.\mu c.[d]w'$ with $w' \stackrel{\text{def}}{=} \lambda x.\mu c.[d]x w''$ and $w'' \stackrel{\text{def}}{=} \lambda x.\lambda y.\lambda z.\Omega$. Then one can check that $([a]v_0)\langle [a]\square t/a \rangle \rightarrow_v^* \lambda z.\Omega$ and $([a]v_1)\langle [a]\square t/a \rangle \rightarrow_v^* \Omega$. Roughly, in $([a]v_1)\langle [a]\square t/a \rangle$, t is reduced to a value once and captures $[a]v_1$, while t is reduced twice to a value in $([a]v_0)\langle [a]\square t/a \rangle$ and each time captures a different term. Hence, we observe a discrepancy in their respective behaviors, and $[a]v_0$ and $[a]v_1$ are not contextually equivalent.

Example 9. This counter-example shows that we need clause 2 for \approx to be complete. Let $v_0 \stackrel{\text{def}}{=} \lambda x.\mu b.[a](\lambda y.\lambda z.w y) x$ and $v_1 \stackrel{\text{def}}{=} w$ with $w \stackrel{\text{def}}{=} \lambda x.w' (x \lambda y.y)$, and $w' \stackrel{\text{def}}{=} \lambda y.y \lambda z.\Omega$.

Let t such that $[b]t \Downarrow_v [b]v$ for $b \notin \text{fn}(t)$. Then we have $([a]v_0)\langle [a]\square t/a \rangle \rightarrow_v^* [a]w' (v\langle [a]v_0\langle [a]\square t/a \rangle \square/b \rangle \lambda x.x)$ and $([a]v_1)\langle [a]\square t/a \rangle \rightarrow_v^* [a]w' (v\langle [a]v_1 \square/b \rangle \lambda x.x)$. We can prove the two resulting terms are bisimilar by showing that the relation $\{(u\langle [a]E[v_0\langle [a]\square t/a \rangle \square/b \rangle], u\langle [a]E[v_1 \square/b \rangle]) \mid [b]t \Downarrow_v [b]v, b \notin \text{fn}(t)\}$ is an applicative bisimulation. The proof is by case analysis on u , and can be found in Appendix B.1. Because \approx is complete (Theorem 4), the terms $([a]v_0)\langle [a]\square t/a \rangle$ and $([a]v_1)\langle [a]\square t/a \rangle$ are contextually equivalent.

However, $[a]v_0$ and $[a]v_1$ themselves are not contextually equivalent. Indeed, we have $([a]v_0)\langle [a]\lambda x.x x \square/a \rangle \rightarrow_v^* \lambda z.\Omega$ and $([a]v_1)\langle [a]\lambda x.x x \square/a \rangle \rightarrow_v^* \Omega$. As a result, testing named values with clause 1 only is not enough; we need clause 2 as well.

3.4 Examples

Even if applicative bisimulation for CBV is difficult to use, we can still prove some equivalences with it. Here we give some examples inspired from Sabry and Felleisen's axiomatization of call/cc [14]. Given a name a , we write a^\dagger for the term $\lambda x.\mu b.[a]x$, and we encode call/cc into $\lambda x.\mu a.[a]x a^\dagger$. Given a named context \mathbb{E} , we also write \mathbb{E}^\dagger for $\lambda x.\mu b.\mathbb{E}[x]$, where $b \notin \text{fn}(\mathbb{E})$. The first example is the axiom C_{tail} of [14], where call/cc is exchanged with a λ -abstraction.

Example 10. If $y \notin \text{fv}(t_1)$ and b is fresh, then $[b](\lambda x.\mu a.[a]x a^\dagger) (\lambda y.(\lambda z.t_0) t_1) \approx [b](\lambda z.(\lambda x.\mu a.[a]x a^\dagger) (\lambda y.t_0)) t_1$.

Proof. Let $v_0 \stackrel{\text{def}}{=} \lambda z.t_0\{b^\dagger/y\}$ and $v_1 \stackrel{\text{def}}{=} \lambda z.(\lambda x.\mu a.[a]x a^\dagger) (\lambda y.t_0)$. The term on the left reduces to $[b]v_0 t_1$, so we relate this term to the one the right, i.e., $[b]v_1 t_1$. We distinguish several cases depending on t_1 . Let c be a fresh name. If $[c]t_1 \Downarrow_v [c]v$, then $[b]v_0 t_1 \rightarrow_v^* [b]t_0\{b^\dagger/y\}\{v\langle^{[b]v_0} \square/c\rangle/z\}$ and $[b]v_1 t_1 \rightarrow_v^* [b]t_0\{b^\dagger/y\}\{v\langle^{[b]v_1} \square/c\rangle/z\}$; because c is fresh, it does not occur in t_0 , and the previous terms can be written $u\langle^{[b]v_0} \square/c\rangle$ and $u\langle^{[b]v_1} \square/c\rangle$ with $u \stackrel{\text{def}}{=} [b]t_0\{b^\dagger/y\}\{v/z\}$.

Similarly, if $[c]t_1 \Downarrow_v [d]v$ with $c \neq d$, then $[b]v_0 t_1 \Downarrow_v [d]v\langle^{[b]v_0} \square/c\rangle$ and $[b]v_1 t_1 \Downarrow_v [d]v\langle^{[b]v_1} \square/c\rangle$. When testing these two values with clauses 1 and 2, we obtain each time terms of the form $u\langle^{[b]v_0} \square/c\rangle$ and $u\langle^{[b]v_1} \square/c\rangle$ for some u . With this reasoning, we can prove that $\{u\langle^{[b]v_0} \square/c\rangle, u\langle^{[b]v_1} \square/c\rangle \mid u \in U^0\}$ is an applicative bisimulation, by case analysis on u .

With the next example and congruence of \approx , we can prove the axiom C_{abort} .

Example 11. Let $a \neq b$; we have $[b]E[a^\dagger t] \approx [b]a^\dagger t$.

Proof. We prove that the relation $\mathcal{R} \stackrel{\text{def}}{=} \{u\langle^{[b]E[\mathbb{E}^\dagger] \square/c}\rangle, u\langle^{[b]\mathbb{E}^\dagger} \square/c\rangle \mid u \in U^0\}$ is an applicative bisimulation by case analysis on u . For example, if $u = [c]t$ and $u \Downarrow_v [c]v$, then $u\langle^{[b]E[\mathbb{E}^\dagger] \square/c}\rangle \rightarrow_v^* [b]E[\mathbb{E}^\dagger v\langle^{[b]E[\mathbb{E}^\dagger] \square/c}\rangle] \rightarrow_v \mathbb{E}[v\langle^{[b]E[\mathbb{E}^\dagger] \square/c}\rangle]$ and $u\langle^{[b]\mathbb{E}^\dagger} \square/c\rangle \rightarrow_v^* [b]\mathbb{E}^\dagger v\langle^{[b]\mathbb{E}^\dagger} \square/c\rangle \rightarrow_v \mathbb{E}[v\langle^{[b]\mathbb{E}^\dagger} \square/c\rangle]$. If $\mathbb{E} \neq [d]\square$, then the resulting terms are in \mathcal{R} , otherwise we get two named values; when checking clauses 1 and 2, we obtain terms of the form $u'\langle^{[b]E[\mathbb{E}'^\dagger] \square/c}\rangle$ and $u'\langle^{[b]\mathbb{E}'^\dagger} \square/c\rangle$ that are in \mathcal{R} . The remaining cases are similar.

Example 12 (axiom $C_{\text{if}t}$). We have $[b]E[(\lambda x.\mu a.[a]x a^\dagger) t] \approx [b]E[t(\lambda x.b^\dagger E[x])]$.

Proof. In this proof, we use an intermediary result, proved in Appendix B.1: if $\mathbb{E} = \mathbb{E}_0[E_I]$, then $\mathbb{E}^\dagger \approx \lambda x.\mathbb{E}_0^\dagger(E_I[x])$. The proof of the axiom itself is by case analysis on t . An interesting case is when $[d]t \Downarrow_v [d]\lambda y.t'$ where $d \notin \text{fn}(t)$. Then $[b]E[(\lambda x.\mu a.[a]x a^\dagger) t] \rightarrow_v^* [b]E[(\lambda x.\mu a.[a]x a^\dagger) \lambda y.t'\langle^{[b]E[(\lambda x.\mu a.[a]x a^\dagger) \square/d]}\rangle] \rightarrow_v^* [b]E[t'\langle^{[b]E[(\lambda x.\mu a.[a]x a^\dagger) \square/d]}\rangle\{\mathbb{E}^\dagger/y\}]$ (with $\mathbb{E} = [b]E$), and $[b]E[t(\lambda x.b^\dagger E[x])] \rightarrow_v^* [b]E[t'\langle^{[b]E[\square(\lambda x.b^\dagger E[x])]/d}\rangle\{\lambda x.b^\dagger E[x]/y\}]$. From the intermediary result, and because \approx is a congruence, we know that $[b]E[t'\{\mathbb{E}^\dagger/y\}] \approx [b]E[t'\{\lambda x.b^\dagger E[x]/y\}]$. Hence, to conclude the proof, one can show that

$$\{(u_0\langle^{[b]E[(\lambda x.\mu a.[a]x a^\dagger) \square/d]}\rangle, u_1\langle^{[b]E[\square(\lambda x.\mathbb{E}_0^\dagger E_I[x])]/d}\rangle) \mid u_0 \approx u_1, \mathbb{E} = \mathbb{E}_0[E_I]\}$$

is an applicative bisimulation.

4 Conclusion

In this work we propose a definition of applicative bisimilarity for CBN and CBV $\lambda\mu$ -calculus, and we also briefly discuss environmental bisimilarity for CBN (a

definition for CBV can be derived from the applicative one). Even if the two applicative definitions for CBN and CBV seem quite different, they follow the same principles. First, we believe it is essential for completeness to hold to relate primarily named terms, and then extend the definition to all terms, as we can see when comparing with Lassen’s definition of applicative bisimilarity. The top level names allow to keep track of how the top level is captured and manipulated in the compared terms.

Then, the idea is to test named values with elementary contexts, $[a]\square t$ for CBN, and $[a]\square t$ and $[a]v \square$ for CBV. In the CBV case, we slightly restrict the terms t tested when considering $[a]\square t$, but the resulting definition remains complex to use compared to CBN, as we can see with Examples 2 and 7. However, we provide counter-examples showing that we cannot simplify it further (see Examples 8 and 9). In CBV as well as in CBN, applicative bisimilarity is harder to use than eager normal form bisimilarity [18], but our relations are complete characterizations of contextual equivalence, and we can therefore prove equivalences of terms that cannot be related with normal form bisimilarity, such as David and Py’s example (see Example 4). To prove the equivalence between two given $\lambda\mu$ -terms, one should try with the bisimulation of [18], and if it fails, try next our applicative or environmental bisimulations.

We believe the relations we define remain complete w.r.t. contextual equivalence in other variants of the $\lambda\mu$ -calculus (perhaps with some slight variations), such as $\lambda\mu$ with different reduction semantics (like, e.g., in [3]), typed $\lambda\mu$ -calculus [13], or Saurin’s extended calculus ($\Lambda\mu$ -calculus [17]). However, any direct implications of this work for other calculi for abortive continuations such as the syntactic theory of control [4] are unclear and remain to be investigated. The reason is that our approach hinges on the syntactic notion of names, unique to the $\lambda\mu$ -calculus, that allows one to keep track of the whereabouts of the top level.

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A Call-by-Name $\lambda\mu$ -calculus

A.1 Soundness and Completeness of Applicative Bisimilarity

Let $(\approx^\bullet)^c$ be the restriction of \approx^\bullet to closed terms.

Lemma 13. *If $t_0 \approx^\bullet t_1$, then there exists a substitution σ which closes t_0 and t_1 such that $t_0\sigma (\approx^\bullet)^c t_1\sigma$, and the size of the derivation of $t_0\sigma (\approx^\bullet)^c t_1\sigma$ is equal to the size of the derivation of $t_0 \approx^\bullet t_1$. A similar result holds if $u_0 \approx^\bullet u_1$.*

Proof. As usual.

Lemma 14. *Let $t_0 (\approx^\bullet)^c t_1$, and $a \notin \text{fn}(t_0) \cup \text{fn}(t_1)$.*

- *If $[a]t_0 \rightarrow_n u_0$, then $[a]t_1 \rightarrow_n^* u_1$ and $u_0 (\approx^\bullet)^c u_1$.*
- *If $t_0 = \lambda x.t'_0$, then $[a]t_1 \rightarrow_n^* [a]\lambda x.t'_1$ and for all $t''_0 (\approx^\bullet)^c t''_1$, we have $[a]t'_0\{t''_0/x\} (\approx^\bullet)^c [a]t'_1\langle [a]\square t''_1/a \rangle\{t''_1/x\}$.*

Proof. By induction on $t_0 (\approx^\bullet)^c t_1$.

Suppose $t_0 \approx t_1$. Then $[a]t_0 \approx [a]t_1$. If $[a]t_0 \rightarrow_n u_0$, then by bisimilarity, we have $[a]t_1 \rightarrow_n^* u_1$ and $u_0 \approx u_1$, i.e., $u_0 (\approx^\bullet)^c u_1$, as required.

Suppose $t_0 = \lambda x.t'_0$, and let $t''_0 (\approx^\bullet)^c t''_1$. By the bisimilarity definition, we have $[a]t_1 \rightarrow_n^* [a]\lambda x.t'_1$ and $[a]t'_0\{t''_0/x\} \approx [a]t'_1\langle [a]\square t''_1/a \rangle\{t''_1/x\}$ (using the fact that $a \notin \text{fn}(t'_0)$). From $t''_0 (\approx^\bullet)^c t''_1$, we deduce $t'_0\{t''_0/x\} (\approx^\bullet)^c t'_0\{t''_1/x\}$, and then $[a]t'_0\{t''_0/x\} (\approx^\bullet)^c \approx [a]t'_1\langle [a]\square t''_1/a \rangle\{t''_1/x\}$, which implies $[a]t'_0\{t''_0/x\} (\approx^\bullet)^c [a]t'_1\langle [a]\square t''_1/a \rangle\{t''_1/x\}$, as required.

Suppose $t_0 \approx^\bullet t \approx^\circ t_1$, so that t is closed (using Lemma 13 if necessary). In fact, we have $t_0 (\approx^\bullet)^c t \approx t_1$. If $[a]t_0 \rightarrow_n u_0$, then by the induction hypothesis, there exists u such that $[a]t \rightarrow_n^* u$ and $u_0 (\approx^\bullet)^c u$. By bisimilarity, there exists u_1 such that $[a]t_1 \rightarrow_n^* u_1$ and $u \approx u_1$. From $u_0 (\approx^\bullet)^c u \approx u_1$, we deduce $u_0 (\approx^\bullet)^c u_1$, as wished.

Suppose $t_0 = \lambda x.t'_0$, and let $t''_0 (\approx^\bullet)^c t''_1$. By induction, there exists t' such that $[a]t \rightarrow_n^* [a]\lambda x.t'$ and $[a]t'_0\{t''_0/x\} (\approx^\bullet)^c [a]t'\langle [a]\square t''_1/a \rangle\{t''_1/x\}$. By bisimilarity, there exists t'_1 such that $[a]t_1 \rightarrow_n^* [a]\lambda x.t'_1$ and $[a]t'\langle [a]\square t''_1/a \rangle\{t''_1/x\} \approx [a]t'_1\langle [a]\square t''_1/a \rangle\{t''_1/x\}$. Hence, we have $[a]t'_0\{t''_0/x\} (\approx^\bullet)^c \approx [a]t'\langle [a]\square t''_1/a \rangle\{t''_1/x\}$, i.e., $[a]t'_0\{t''_0/x\} (\approx^\bullet)^c [a]t'_1\langle [a]\square t''_1/a \rangle\{t''_1/x\}$ as required.

If $t_0 \widetilde{(\approx^\bullet)^c} t_1$, then we have several cases to consider.

Suppose $t_0 = \lambda x.t'_0$ and $t_1 = \lambda x.t'_1$ with $t'_0 \approx^\bullet t'_1$. Let $t''_0 (\approx^\bullet)^c t''_1$. We have $[a]t'_0\{t''_0/x\} (\approx^\bullet)^c [a]t'_1\{t''_1/x\}$, hence the result holds (note that $a \notin \text{fv}(t'_1)$ because $a \notin \text{fv}(t_1)$).

Suppose $t_0 = \mu b.u_0$ and $t_1 = \mu b.u_1$ with $u_0 (\approx^\bullet)^c u_1$. We have $[a]t_0 \rightarrow_n u_0\langle [a]\square/b \rangle$, $[a]t_1 \rightarrow_n u_1\langle [a]\square/b \rangle$, and $u_0\langle [a]\square/b \rangle (\approx^\bullet)^c u_1\langle [a]\square/b \rangle$, hence the result holds.

Suppose $t_0 = t_0^1 t_0^2$, $t_1 = t_1^1 t_1^2$ with $t_0^1 (\approx^\bullet)^c t_1^1$ and $t_0^2 (\approx^\bullet)^c t_1^2$. We distinguish two cases.

- If $[b]t_0^1 \rightarrow_n u_0$ (for some fresh b), then $[a]t_0 \rightarrow_n u_0 \langle [a] \square t_0^2/b \rangle$. By the induction hypothesis, there exists u_1 such that $[b]t_1^1 \rightarrow_n^* u_1$ and $u_0 \approx^{\bullet} u_1$. Consequently, we have $[a]t_1 \rightarrow_n^* u_1 \langle [a] \square t_1^2/b \rangle$, and by definition of \approx^{\bullet} , we have $u_0 \langle [a] \square t_0^2/b \rangle \approx^{\bullet} u_1 \langle [a] \square t_1^2/b \rangle$, as required.
- If $t_0^1 = \lambda x.t'_0$, then we have $[a]t_0 \rightarrow_n [a]t'_0 \{t_0^2/x\}$. By the induction hypothesis, there exists t'_1 such that $[a]t_1^1 \rightarrow_n^* [a]\lambda x.t'_1$ and $[a]t'_0 \{t_0^2/x\} \approx^{\bullet} [a]t'_1 \langle [a] \square t'_1/a \rangle \{t_1^2/x\}$. From $[a]t_1^1 \rightarrow_n^* [a]\lambda x.t'_1$, we deduce

$$[a]t_1 \rightarrow_n^* [a] \langle \lambda x.t'_1 \langle [a] \square t'_1/a \rangle \rangle t_1^2 \rightarrow_n [a]t'_1 \langle [a] \square t'_1/a \rangle \{t_1^2/x\},$$

hence the result holds.

Suppose $t_0 = t'_0 \langle \mathbb{E}_0/b \rangle$, $t_1 = t'_1 \langle \mathbb{E}_1/b \rangle$ with $t'_0 \approx^{\bullet} t'_1$, $\mathbb{E}_0 \approx^{\bullet} \mathbb{E}_1$. If $[a]t_0 \rightarrow_n u_0$, then in fact $[a]t'_0 \rightarrow_n u'_0$ with $u_0 = u'_0 \langle \mathbb{E}_0/b \rangle$. By the induction hypothesis, there exists u'_1 such that $[a]t'_1 \rightarrow_n^* u'_1$ and $u'_0 \approx^{\bullet} u'_1$. Consequently, we have $[a]t_1 \rightarrow_n^* u'_1 \langle \mathbb{E}_1/b \rangle$, and $u'_0 \langle \mathbb{E}_0/b \rangle \approx^{\bullet} u'_1 \langle \mathbb{E}_1/b \rangle$ holds, as wished. If $[a]t_0$ is a named value, then in fact $t'_0 = \lambda x.t''_0$ and $t_0 = \lambda x.t''_0 \langle \mathbb{E}_0/b \rangle$. Let $t''_0 \approx^{\bullet} t''_1$. Then $t''_0 \langle [c] \square/b \rangle \approx^{\bullet} t''_1 \langle [c] \square/b \rangle$ for a fresh c . By the induction hypothesis, there exists t''_1 such that $[a]t'_1 \rightarrow_n^* [a]\lambda x.t''_1$ and

$$[a]t''_0 \{t''_0 \langle [c] \square/b \rangle / x\} \approx^{\bullet} [a]t''_1 \langle [a] \square t''_1 \langle [c] \square/b \rangle / a \rangle \{t''_1 \langle [c] \square/b \rangle / x\}.$$

Therefore, we have $[a]t_1 \rightarrow_n^* [a]\lambda x.t''_1 \langle \mathbb{E}_1/b \rangle$, and

$$[a]t''_0 \{t''_0 \langle [c] \square/b \rangle / x\} \langle \mathbb{E}_0/b \rangle \approx^{\bullet} [a]t''_1 \langle [a] \square t''_1 \langle [c] \square/b \rangle / a \rangle \{t''_1 \langle [c] \square/b \rangle / x\} \langle \mathbb{E}_1/b \rangle.$$

Because b does not occur in $t''_0 \langle [c] \square/b \rangle$, $t''_1 \langle [c] \square/b \rangle$ thanks to the renaming to a fresh c , we can switch the substitutions around, and in fact

$$[a]t''_0 \langle \mathbb{E}_0/b \rangle \{t''_0 \langle [c] \square/b \rangle / x\} \approx^{\bullet} [a]t''_1 \langle \mathbb{E}_1/b \rangle \langle [a] \square t''_1 \langle [c] \square/b \rangle / a \rangle \{t''_1 \langle [c] \square/b \rangle / x\}$$

holds. Renaming c back into b , we obtain

$$[a]t''_0 \langle \mathbb{E}_0/b \rangle \{t''_0/x\} \approx^{\bullet} [a]t''_1 \langle \mathbb{E}_1/b \rangle \langle [a] \square t''_1/a \rangle \{t''_1/x\},$$

which gives us the required result.

Lemma 15. *Let $u_0 \approx^{\bullet} u_1$.*

- If $u_0 \rightarrow_n u'_0$, then $u_1 \rightarrow_n^* u'_1$ and $u'_0 \approx^{\bullet} u'_1$.
- If $u_0 = [a]\lambda x.t_0$, then $u_1 \rightarrow_n^* [a]\lambda x.t_1$ and for all $t'_0 \approx^{\bullet} t'_1$, we have $[a]t_0 \langle [a] \square t'_0/a \rangle \{t'_0/x\} \approx^{\bullet} [a]t_1 \langle [a] \square t'_1/a \rangle \{t'_1/x\}$.

Proof. By induction on $u_0 \approx^{\bullet} u_1$.

Suppose $u_0 \approx u_1$. The first item holds by bisimilarity. Suppose $u_0 = [a]\lambda x.t_0$, and let $t'_0 \approx^{\bullet} t'_1$. By definition of the bisimilarity, we have $u_1 \rightarrow_n^* [a]\lambda x.t_1$ and $[a]t_0 \langle [a] \square t'_0/a \rangle \{t'_0/x\} \approx [a]t_1 \langle [a] \square t'_1/a \rangle \{t'_1/x\}$. From $t'_0 \approx^{\bullet} t'_1$, we deduce $t_0 \langle [a] \square t'_0/a \rangle \{t'_0/x\} \approx^{\bullet} t_0 \langle [a] \square t'_1/a \rangle \{t'_1/x\}$, which implies

$$t_0 \langle [a] \square t'_0/a \rangle \{t'_0/x\} \approx^{\bullet} [a]t_1 \langle [a] \square t'_1/a \rangle \{t'_1/x\},$$

which in turn implies

$$[a]t_0 \langle [a] \square t'_0/a \rangle \{t'_0/x\} (\approx^\bullet)^c [a]t_1 \langle [a] \square t'_1/a \rangle \{t'_1/x\},$$

as required.

Suppose $u_0 \approx^\bullet u \approx^\circ u_1$, so that u is closed (using Lemma 13 if necessary). In fact, we have $u_0 (\approx^\bullet)^c u \approx u_1$. If $u_0 \rightarrow_n u'_0$, then by the induction hypothesis, there exists u' such that $u \rightarrow_n^* u'$ and $u'_0 (\approx^\bullet)^c u'$. By bisimilarity, there exists u'_1 such that $u_1 \rightarrow_n^* u'_1$ and $u' \approx u'_1$. From $u'_0 (\approx^\bullet)^c u' \approx u'_1$, we deduce $u'_0 (\approx^\bullet)^c u'_1$, as wished.

Suppose $u_0 = [a]\lambda x.t_0$, and let $t'_0 (\approx^\bullet)^c t'_1$. By induction, there exists t such that $u \rightarrow_n^* [a]\lambda x.t$ and $[a]t_0 \langle [a] \square t'_0/a \rangle \{t'_0/x\} (\approx^\bullet)^c [a]t \langle [a] \square t'_1/a \rangle \{t'_1/x\}$. By bisimilarity, there exists t_1 such that $u_1 \rightarrow_n^* [a]\lambda x.t_1$ and $[a]t \langle [a] \square t'_1/a \rangle \{t'_1/x\} \approx [a]t_1 \langle [a] \square t'_1/a \rangle \{t'_1/x\}$. Consequently, we have

$$[a]t_0 \langle [a] \square t'_0/a \rangle \{t'_0/x\} (\approx^\bullet)^c \approx [a]t \langle [a] \square t'_1/a \rangle \{t'_1/x\},$$

i.e., $[a]t_0 \langle [a] \square t'_0/a \rangle \{t'_0/x\} (\approx^\bullet)^c [a]t_1 \langle [a] \square t'_1/a \rangle \{t'_1/x\}$ holds as required.

If $u_0 \widetilde{(\approx^\bullet)^c} u_1$, then we have two cases.

Suppose $u_0 = [a]t_0$ and $u_1 = [a]t_1$ with $t_0 (\approx^\bullet)^c t_1$. If $a \notin \text{fn}(t_0) \cup \text{fn}(t_1)$, then we can apply Lemma 14 directly to get the required result. Otherwise, let $b \notin \text{fn}(t_0) \cup \text{fn}(t_1)$. If $u_0 \rightarrow_n u'_0$, then $[b]t_0 \rightarrow_n u''_0$ and $u'_0 = u'_0 \langle [a] \square /b \rangle$. We can apply Lemma 14 to $[b]t_0$ and $[b]t_1$, and then rename b into a .

Suppose $u_0 = u'_0 \langle \mathbb{E}_0/a \rangle$, $u_1 = u'_1 \langle \mathbb{E}_1/a \rangle$ with $u'_0 (\approx^\bullet)^c u'_1$ and $\mathbb{E}_0 (\approx^\bullet)^c \mathbb{E}_1$.

If u_0 is a named value, then we distinguish two cases. First, we may have $u'_0 = [a]\lambda x.t_0$, $\mathbb{E}_0 = \mathbb{E}_1 = [b] \square$, and $u_0 = [b]\lambda x.t_0 \langle \mathbb{E}_0/a \rangle$. Let $t'_0 (\approx^\bullet)^c t'_1$. Let t''_0, t''_1 be t'_0 and t'_1 where a, b are renamed into fresh c, d to avoid some name clashes (we still have $t''_0 (\approx^\bullet)^c t''_1$). By the induction hypothesis, there exists t_1 such that $u'_1 \rightarrow_n^* [a]\lambda x.t_1$ and $[a]t_0 \langle [a] \square t''_0/a \rangle \{t''_0/x\} (\approx^\bullet)^c [a]t_1 \langle [a] \square t''_1/a \rangle \{t''_1/x\}$. This implies

$$[a]t_0 \langle [a] \square t''_0/a \rangle \{t''_0/x\} \langle [a] \square t''_0/b \rangle (\approx^\bullet)^c [a]t_1 \langle [a] \square t''_1/a \rangle \{t''_1/x\} \langle [a] \square t''_1/b \rangle,$$

which is the same as

$$[a]t_0 \langle [a] \square t''_0/a \rangle \langle [a] \square t''_0/b \rangle \{t''_0/x\} (\approx^\bullet)^c [a]t_1 \langle [a] \square t''_1/a \rangle \langle [a] \square t''_1/b \rangle \{t''_1/x\}$$

because b does not occur in t''_0, t''_1 . In turn, we have

$$[b]t_0 \langle [a] \square t''_0/a \rangle \langle [a] \square t''_0/b \rangle \{t''_0/x\} \langle \mathbb{E}_0/a \rangle (\approx^\bullet)^c [b]t_1 \langle [a] \square t''_1/a \rangle \langle [a] \square t''_1/b \rangle \{t''_1/x\} \langle \mathbb{E}_1/a \rangle,$$

which is equal to

$$[b]t_0 \langle \mathbb{E}_0/a \rangle \langle [b] \square t''_0/b \rangle \{t''_0/x\} (\approx^\bullet)^c [b]t_1 \langle \mathbb{E}_1/a \rangle \langle [b] \square t''_1/b \rangle \{t''_1/x\}$$

because a does not occur in t''_0, t''_1 . Renaming c, d back into a, b , we obtain

$$[b]t_0 \langle \mathbb{E}_0/a \rangle \langle [b] \square t'_0/b \rangle \{t''_0/x\} (\approx^\bullet)^c [b]t_1 \langle \mathbb{E}_1/a \rangle \langle [b] \square t'_1/b \rangle \{t''_1/x\},$$

and because $u_1 \rightarrow_n^* [b] \lambda x. t_1 \langle \mathbb{E}_1/a \rangle$, the result holds.

In the second case, we have $u'_0 = [b] \lambda x. t_0$ and $u_0 = [b] \lambda x. t_0 \langle \mathbb{E}_0/a \rangle$ with $b \neq a$. Let $t'_0 (\approx^\bullet)^c t'_1$. Let t''_0, t''_1 be t'_0 and t'_1 where a is renamed into a fresh c . By the induction hypothesis, there exists t_1 such that $u'_1 \rightarrow_n^* [b] \lambda x. t_1$ and

$$[b]t_0 \langle [b] \square t''_0/b \rangle \{t''_0/x\} (\approx^\bullet)^c [b]t_1 \langle [b] \square t''_1/b \rangle \{t''_1/x\}.$$

From $\mathbb{E}_0 \langle [b] \square t''_0/b \rangle (\approx^\bullet)^c \mathbb{E}_1 \langle [b] \square t''_1/b \rangle$ and the previous relation, we can deduce

$$[b]t_0 \langle [b] \square t''_0/b \rangle \{t''_0/x\} \langle \mathbb{E}_0 \langle [b] \square t''_0/b \rangle /a \rangle (\approx^\bullet)^c [b]t_1 \langle [b] \square t''_1/b \rangle \{t''_1/x\} \langle \mathbb{E}_1 \langle [b] \square t''_1/b \rangle /a \rangle.$$

Because a does not occur in t''_0, t''_1 , this can be rewritten into

$$[b]t_0 \langle \mathbb{E}_0/a \rangle \langle [b] \square t''_0/b \rangle \{t''_0/x\} (\approx^\bullet)^c [b]t_1 \langle \mathbb{E}_1/a \rangle \langle [b] \square t''_1/b \rangle \{t''_1/x\}.$$

By renaming c back into a , we obtain

$$[b]t_0 \langle \mathbb{E}_0/a \rangle \langle [b] \square t'_0/b \rangle \{t'_0/x\} (\approx^\bullet)^c [b]t_1 \langle \mathbb{E}_1/a \rangle \langle [b] \square t'_1/b \rangle \{t'_1/x\},$$

and because $u_1 \rightarrow_n^* [b] \lambda x. t_1 \langle \mathbb{E}_1/a \rangle$, the result holds.

If $u_0 \rightarrow_n$, then again we distinguish two cases. First, suppose $u'_0 \rightarrow_n u''_0$; then $u_0 \rightarrow_n u''_0 \langle \mathbb{E}_0/a \rangle$. By the induction hypothesis, there exists u'_1 such that $u'_1 \rightarrow_n^* u''_1$ and $u''_0 (\approx^\bullet)^c u''_1$. Then $u_1 \rightarrow_n^* u'_1 \langle \mathbb{E}_1/a \rangle$ and $u''_0 \langle \mathbb{E}_0/a \rangle (\approx^\bullet)^c u''_1 \langle \mathbb{E}_1/a \rangle$, hence the result holds.

Otherwise, $u'_0 = [a] \lambda x. t_0$, $\mathbb{E}_0 = \mathbb{E}'_0[\square t'_0]$, and $u_0 \rightarrow_n \mathbb{E}'_0[t_0 \langle \mathbb{E}_0/a \rangle \{t'_0/x\}]$. We can prove by induction on $\mathbb{E}_0 (\approx^\bullet)^c \mathbb{E}_1$ that $\mathbb{E}_1 = \mathbb{E}'_1[\square t'_1]$ with $\mathbb{E}'_0 (\approx^\bullet)^c \mathbb{E}'_1$ and $t'_0 (\approx^\bullet)^c t'_1$. Let \mathbb{E}''_i, t''_i be \mathbb{E}'_i, t'_i with a renamed into a fresh b . By the induction hypothesis, there exists t_1 such that $u'_1 \rightarrow_n^* [a] \lambda x. t_1$, and

$$[a]t_0 \langle [a] \square t''_0/a \rangle \{t''_0/x\} (\approx^\bullet)^c [a]t_1 \langle [a] \square t''_1/a \rangle \{t''_1/x\}.$$

This implies

$$([a]t_0 \langle [a] \square t''_0/a \rangle \{t''_0/x\}) \langle \mathbb{E}'_0/a \rangle (\approx^\bullet)^c ([a]t_1 \langle [a] \square t''_1/a \rangle \{t''_1/x\}) \langle \mathbb{E}'_1/a \rangle,$$

i.e., $\mathbb{E}'_0[t_0 \langle \mathbb{E}'_0[\square t''_0/a] \{t''_0/x\}] (\approx^\bullet)^c \mathbb{E}'_1[t_1 \langle \mathbb{E}'_1[\square t''_1/a] \{t''_1/x\}]$ because a does not occur in t''_0, t''_1 . Renaming b into a , we obtain

$$\mathbb{E}'_0[t_0 \langle \mathbb{E}_0/a \rangle \{t'_0/x\}] (\approx^\bullet)^c \mathbb{E}'_1[t_1 \langle \mathbb{E}_1/a \rangle \{t'_1/x\}],$$

and because $u_1 \rightarrow_n^* \mathbb{E}'_1[t_1 \langle \mathbb{E}_1/a \rangle \{t'_1/x\}]$, we have the required result.

From there, we can prove that $(\approx^\bullet)^c = \approx$ using the usual techniques [5], and then we deduce soundness of \approx . To prove completeness, we show that \approx_c is an applicative bisimulation.

Lemma 16. *The relation \approx_c is a big step applicative bisimulation.*

Proof. Suppose $u_0 \approx_c u_1$. If $u_0 \Downarrow_n [a] \lambda x. t_0$, then $u_1 \Downarrow_n [a] \lambda x. t_1$. We have $[a](\mu a. u_0) t \rightarrow_n^* [a]t_0 \langle [a] \square t/a \rangle \{t/x\}$, $[a](\mu a. u_1) t \rightarrow_n^* [a]t_1 \langle [a] \square t/a \rangle \{t/x\}$, but also $[a](\mu a. u_0) t \approx_c [a](\mu a. u_1) t$. From $\rightarrow_n^* \subseteq \approx \subseteq \approx_c$, we have $[a]t_0 \langle [a] \square t/a \rangle \{t/x\} \approx_c [a]t_1 \langle [a] \square t/a \rangle \{t/x\}$ as wished.

A.2 Soundness and Completeness of Environmental Bisimilarity

Lemma 17. *Let \mathcal{R} be a relation on closed terms and named terms. If $t_0 \widehat{\mathcal{R}} t_1$ and $t'_0 \widehat{\mathcal{R}} t'_1$, then $t_0\{t'_0/x\} \widehat{\mathcal{R}} t_1\{t'_1/x\}$ (and similarly if $u_0 \widehat{\mathcal{R}} u_1$)*

Proof. By induction on $t_0 \widehat{\mathcal{R}} t_1$.

Lemma 18. *If $u_0 \cong_{\mathcal{E}} u_1$, then $u_0\langle\mathbb{E}/a\rangle \cong_{\mathcal{E}} u_1\langle\mathbb{E}/a\rangle$.*

Let \mathcal{Y} be an environmental bisimulation. We define

$$\begin{aligned} \mathcal{X} &= \mathcal{X}_1 \cup \mathcal{X}_2 \cup \{\widehat{\mathcal{E}}^{\text{nv}}, \mathcal{E} \in \mathcal{Y}\} \\ \mathcal{X}_1 &= \{(\widehat{\mathcal{E}}^{\text{nv}}, u_0\langle\mathbb{E}_0/a_0\rangle \dots \langle\mathbb{E}_n/a_n\rangle, u_1\langle\mathbb{E}_1/a_0\rangle \dots \langle\mathbb{E}_1/a_n\rangle), u_0 \mathcal{Y}_{\mathcal{E}} u_1, \mathbb{E}_0 \widehat{\mathcal{E}} \mathbb{E}_1^i\} \\ \mathcal{X}_2 &= \{(\widehat{\mathcal{E}}^{\text{nv}}, u_0, u_1), \mathcal{E} \in \mathcal{Y}, u_0 \widehat{\mathcal{E}} u_1\} \end{aligned}$$

We first prove some preliminary lemmas about \mathcal{X} .

Lemma 19. *Let $\mathcal{E} \in \mathcal{Y}$.*

- *If $\lambda x.t_0 \widehat{\mathcal{E}} \lambda x.t_1$ and $t'_0 \widehat{\mathcal{E}} t'_1$, then $[a]t_0\{t'_0/x\} \widehat{\mathcal{E}} [a]t_1\{t'_1/x\}$.*
- *If $[a]\lambda x.t_0 \widehat{\mathcal{E}} [a]\lambda x.t_1$ and $t'_0 \widehat{\mathcal{E}} t'_1$, then we have $[a]t_0\langle[a]\square t'_0/a\rangle\{t'_0/x\} \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} [a]t_1\langle[a]\square t'_1/a\rangle\{t'_1/x\}$.*

Proof. The first item is proved by induction on $\lambda x.t_0 \widehat{\mathcal{E}} \lambda x.t_1$. If $t_0 \widehat{\mathcal{E}} t_1$, then $[a]t_0\{t'_0/x\} \widehat{\mathcal{E}} [a]t_1\{t'_1/x\}$ holds using Lemma 17. Otherwise, we have $\lambda x.t_0 = \lambda x.t''_0\langle\mathbb{E}_0/b\rangle$, $\lambda x.t_1 = \lambda x.t''_1\langle\mathbb{E}_1/b\rangle$ with $\lambda x.t''_0 \widehat{\mathcal{E}} \lambda x.t''_1$ and $\mathbb{E}_0 \widehat{\mathcal{E}} \mathbb{E}_1$. Let t''_0, t''_1 be t'_0, t'_1 where b is renamed to a fresh c (we still have $t''_0 \widehat{\mathcal{E}} t''_1$). By the induction hypothesis, we have $t''_0\{t''_0/x\} \widehat{\mathcal{E}} t''_1\{t''_1/x\}$, which implies $t''_0\{t''_0/x\}\langle\mathbb{E}_0/b\rangle \widehat{\mathcal{E}} t''_1\{t''_1/x\}\langle\mathbb{E}_1/b\rangle$, i.e., $t''_0\langle\mathbb{E}_0/b\rangle\{t''_0/x\} \widehat{\mathcal{E}} t''_1\langle\mathbb{E}_1/b\rangle\{t''_1/x\}$ because $b \notin \text{fn}(t''_0) \cup \text{fn}(t''_1)$. Renaming c back into b , we obtain $t''_0\langle\mathbb{E}_0/b\rangle\{t'_0/x\} \widehat{\mathcal{E}} t''_1\langle\mathbb{E}_1/b\rangle\{t'_1/x\}$ as wished.

The second item is proved by induction on $[a]\lambda x.t_0 \widehat{\mathcal{E}} [a]\lambda x.t_1$. If $[a]\lambda x.t_0 \mathcal{E} [a]\lambda x.t_1$, then we have $[a]t_0\langle[a]\square t'_0/a\rangle\{t'_0/x\} \mathcal{Y}_{\mathcal{E}} [a]t_1\langle[a]\square t'_1/a\rangle\{t'_1/x\}$ because \mathcal{Y} is a bisimulation, i.e., $[a]t_0\langle[a]\square t'_0/a\rangle\{t'_0/x\} \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} [a]t_1\langle[a]\square t'_1/a\rangle\{t'_1/x\}$ holds as wished.

Suppose $\lambda x.t_0 \widehat{\mathcal{E}} \lambda x.t_1$, and let t''_0, t''_1 be t'_0, t'_1 where a is renamed into a fresh b . By the first item, we have $[b]t_0\{t''_0/x\} \widehat{\mathcal{E}} [b]t_1\{t''_1/x\}$, which implies $[b]t_0\{t''_0/x\}\langle[a]\square t'_0/a\rangle \widehat{\mathcal{E}} [b]t_1\{t''_1/x\}\langle[a]\square t'_1/a\rangle$, which gives $[b]t_0\langle[a]\square t'_0/a\rangle\{t''_0/x\} \widehat{\mathcal{E}} [b]t_1\langle[a]\square t'_1/a\rangle\{t''_1/x\}$ because a does not occur in t''_0, t''_1 . Renaming b in a , we obtain $[a]t_0\langle[a]\square t'_0/a\rangle\{t'_0/x\} \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} [a]t_1\langle[a]\square t'_1/a\rangle\{t'_1/x\}$, as wished.

Suppose $[a]\lambda x.t_0 = [a]\lambda x.t''_0\langle\mathbb{E}_0/b\rangle$, $[a]\lambda x.t_1 = [a]\lambda x.t''_1\langle\mathbb{E}_1/b\rangle$ with $[a]\lambda x.t''_0 \widehat{\mathcal{E}} [a]\lambda x.t''_1$ and $\mathbb{E}_0 \widehat{\mathcal{E}} \mathbb{E}_1$. Let t''_0, t''_1 be t'_0, t'_1 where b is renamed to a fresh c . By the induction hypothesis, we have $[a]t''_0\langle[a]\square t''_0/a\rangle\{t''_0/x\} \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} [a]t''_1\langle[a]\square t''_1/a\rangle\{t''_1/x\}$, which implies $[a]t''_0\langle[a]\square t''_0/a\rangle\{t''_0/x\}\langle\mathbb{E}_0/b\rangle \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} [a]t''_1\langle[a]\square t''_1/a\rangle\{t''_1/x\}\langle\mathbb{E}_1/b\rangle$ (according to the definition of \mathcal{X}), which gives us $[a]t''_0\langle\mathbb{E}_0/b\rangle\langle[a]\square t''_0/a\rangle\{t''_0/x\} \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} [a]t''_1\langle\mathbb{E}_1/b\rangle\langle[a]\square t''_1/a\rangle\{t''_1/x\}$ (because a does not occur in t''_0, t''_1). Renaming b into a , we obtain $[a]t''_0\langle\mathbb{E}_0/b\rangle\langle[a]\square t'_0/a\rangle\{t'_0/x\} \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} [a]t''_1\langle\mathbb{E}_1/b\rangle\langle[a]\square t'_1/a\rangle\{t'_1/x\}$, as wished.

Lemma 20. *Let $\mathcal{E} \in \mathcal{Y}$, $t_0 \widehat{\mathcal{E}} t_1$, and $a \notin \text{fn}(t_0) \cup \text{fn}(t_1)$. If $[a]t_0 \rightarrow_n u_0$, then $[a]t_1 \rightarrow_n^* u_1$, and $u_0 \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1$.*

Proof. We proceed by induction on $t_0 \widehat{\mathcal{E}} t_1$.

Suppose $t_0 = \mu b.u_0$, $t_1 = \mu b.u_1$ with $u_0 \widehat{\mathcal{E}} u_1$. Then $[a]t_0 \rightarrow_n u_0 \langle [a]\square/b \rangle$, $[a]t_1 \rightarrow_n u_1 \langle [a]\square/b \rangle$, and $u_0 \langle [a]\square/b \rangle \widehat{\mathcal{E}} u_1 \langle [a]\square/b \rangle$ holds, so we have $u_0 \langle [a]\square/b \rangle \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1 \langle [a]\square/b \rangle$, as wished.

Suppose $t_0 = t_0^1 t_0^2$ and $t_1 = t_1^1 t_1^2$ with $t_0^1 \widehat{\mathcal{E}} t_1^1$ and $t_0^2 \widehat{\mathcal{E}} t_1^2$. We have two cases to consider.

- Assume $[a]t_0^1 \rightarrow_v u_0$, so that $[a]t_0 \rightarrow_v u_0 \langle [a]\square/t_0^2/a \rangle$. By the induction hypothesis, there exists u_1 such that $[a]t_1^1 \rightarrow_v^* u_1$ and $u_0 \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1$. From $t_0^2 \widehat{\mathcal{E}} t_1^2$ and $u_0 \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1$, we can deduce $u_0 \langle [a]\square/t_0^2/a \rangle \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1 \langle [a]\square/t_1^2/a \rangle$ by definition of \mathcal{X} . We also have $[a]t_1 \rightarrow_v^* u_1 \langle [a]\square/t_1^2/a \rangle$, hence the result holds.
- Assume $t_0^1 = \lambda x.t_0'$ so that $[a]t_0 \rightarrow_v [a]t_0' \{t_0^2/x\}$. Because \mathcal{E} relates only named values, t_1^1 must be value $\lambda x.t_1'$ as well. We have $[a]t_1 \rightarrow_v [a]t_1' \{t_1^2/x\}$, and by Lemma 19, we have $[a]t_0' \{t_0^2/x\} \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} [a]t_1' \{t_1^2/x\}$, hence the result holds.

Suppose $t_0 = t_0' \langle \mathbb{E}_0/b \rangle$, $t_1 = t_1' \langle \mathbb{E}_1/b \rangle$ with $t_0' \widehat{\mathcal{E}} t_1'$ and $\mathbb{E}_0 \widehat{\mathcal{E}} \mathbb{E}_1$. If $[a]t_0 \rightarrow_n$, then $[a]t_0' \rightarrow_n u_0$ and $[a]t_0 \rightarrow_n u_0 \langle \mathbb{E}_0/a \rangle$. By the induction hypothesis, there exists u_1 such that $[a]t_1' \rightarrow_n^* u_1$ and $u_0 \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1$. By definition of $\mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1$, we have $u_0 \langle \mathbb{E}_0/b \rangle \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1 \langle \mathbb{E}_1/b \rangle$, and we have $[a]t_1 \rightarrow_n^* u_1 \langle \mathbb{E}_1/b \rangle$, as wished.

We now prove Lemma 18 by showing that \mathcal{X} is a bisimulation up to environment.

Proof. We first prove the bisimulation for the elements in \mathcal{X}_2 (for these, we do not need the “up to environment”). Let $t_0 \widehat{\mathcal{E}} t_1$, with $\mathcal{E} \in \mathcal{Y}$. If u_0 is a named value, then one can check that so is t_1 (because \mathcal{E} relates only named values), and we have $\widehat{\mathcal{E}}^{\text{nv}} \cup \{(t_0, t_1)\} = \widehat{\mathcal{E}}^{\text{nv}} \in \mathcal{X}$.

Let $u_0 \widehat{\mathcal{E}} u_1$ such that $u_0 \rightarrow_n u_0'$. We proceed by induction on $u_0 \widehat{\mathcal{E}} u_1$. Suppose $u_0 = [a]t_0$, $u_1 = [a]t_1$, with $t_0 \widehat{\mathcal{E}} t_1$. If $a \notin \text{fn}(t_0) \cup \text{fn}(t_1)$, then we can apply Lemma 20 directly to get the required result. Otherwise, let $b \notin \text{fn}(t_0) \cup \text{fn}(t_1)$. If $u_0 \rightarrow_n u_0'$, then $[b]t_0 \rightarrow_n u_0'$ and $u_0' = u_0' \langle [a]\square/b \rangle$. We can apply Lemma 20 to $[b]t_0$ and $[b]t_1$, and then rename b into a .

Suppose $u_0 = u_0' \langle \mathbb{E}_0/a \rangle$, $u_1 = u_1' \langle \mathbb{E}_1/a \rangle$ with $u_0' \widehat{\mathcal{E}} u_1'$ and $\mathbb{E}_0 \widehat{\mathcal{E}} \mathbb{E}_1$. If u_0 is a named value, then $u_0' = [c]v_0$, and $u_0 = [b]v_0 \langle \mathbb{E}_0/a \rangle$. Because \mathcal{E} relates only named values, u_1' is also a named value $[b]v_1$, and we have $[b]v_0 \langle \mathbb{E}_0/a \rangle \widehat{\mathcal{E}} [b]v_1 \langle \mathbb{E}_1/a \rangle$, as wished.

If $u_0 \rightarrow_n$, then we distinguish two cases. First, suppose $u_0' \rightarrow_n u_0''$; then $u_0 \rightarrow_n u_0'' \langle \mathbb{E}_0/a \rangle$. By the induction hypothesis, there exists u_1'' such that $u_1' \rightarrow_n^* u_1''$ and $u_0'' \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1''$. Then $u_1 \rightarrow_n^* u_1'' \langle \mathbb{E}_1/a \rangle$ and $u_0'' \langle \mathbb{E}_0/a \rangle \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1'' \langle \mathbb{E}_1/a \rangle$ (by definition of \mathcal{X}), hence the result holds.

Otherwise, we have $u_0' = [a]\lambda x.t_0$, $\mathbb{E}_0 = \mathbb{E}_0'[\square t_0']$, which implies $u_0 \rightarrow_n \mathbb{E}_0'[\square t_0' \langle \mathbb{E}_0/a \rangle \{t_0'/x\}]$. We can prove by induction on $\mathbb{E}_0 \widehat{\mathcal{E}} \mathbb{E}_1$ that $\mathbb{E}_1 = \mathbb{E}_1'[\square t_1']$

with $\mathbb{E}'_0 \widehat{\mathcal{E}} \mathbb{E}'_1$ and $t'_0 \widehat{\mathcal{E}} t'_1$. Let \mathbb{E}''_i, t''_i be \mathbb{E}'_i, t'_i with a renamed into a fresh b . Because \mathcal{E} relates only named values, we must have $u'_1 = [a]\lambda x.t_1$. By Lemma 19, we have $[a]t_0 \langle [a]\square t'_0/a \rangle \{t''_0/x\} \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} [a]t_1 \langle [a]\square t'_1/a \rangle \{t''_1/x\}$, which in turn implies $([a]t_0 \langle [a]\square t'_0/a \rangle \{t''_0/x\}) \langle \mathbb{E}''_0/a \rangle \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} ([a]t_1 \langle [a]\square t'_1/a \rangle \{t''_1/x\}) \langle \mathbb{E}''_1/a \rangle$ by definition of \mathcal{X} . Because a does not occur in t''_0, t''_1 , the previous relation is equivalent to $\mathbb{E}''_0 [t_0 \langle \mathbb{E}''_0 \square t'_0/a \rangle \{t''_0/x\}] \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} \mathbb{E}''_1 [t_1 \langle \mathbb{E}''_1 \square t'_1/a \rangle \{t''_1/x\}]$. Renaming b into a , we obtain $\mathbb{E}'_0 [t_0 \langle \mathbb{E}'_0/a \rangle \{t'_0/x\}] \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} \mathbb{E}'_1 [t_1 \langle \mathbb{E}'_1/a \rangle \{t'_1/x\}]$, and because we have $u_1 \rightarrow_n^* \mathbb{E}'_1 [t_1 \langle \mathbb{E}'_1/a \rangle \{t'_1/x\}]$, the result holds.

We now prove bisimulation (up to environment) for elements in \mathcal{X}_1 . Let $u_0 \langle \mathbb{E}^0_0/a_0 \rangle \dots \langle \mathbb{E}^n_0/a_n \rangle \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u_1 \langle \mathbb{E}^0_1/a_0 \rangle \dots \langle \mathbb{E}^n_1/a_n \rangle$ so that $u_0 \mathcal{Y}_{\mathcal{E}} u_1$ and $\mathbb{E}^i_0 \widehat{\mathcal{E}} \mathbb{E}^i_1$. If u_0 is a named value $[a]v_0$, then because \mathcal{Y} is a bisimulation, there exists v_1 such that $u_1 \rightarrow_n^* [a]v_1$ and $\mathcal{E}' = \mathcal{E} \cup \{([a]v_0, [a]v_1)\} \in \mathcal{Y}$. We then have $u_1 \langle \mathbb{E}^0_1/a_0 \rangle \dots \langle \mathbb{E}^n_1/a_n \rangle \rightarrow_n^* ([a]v_1) \langle \mathbb{E}^0_1/a_0 \rangle \dots \langle \mathbb{E}^n_1/a_n \rangle$, and $([a]v_0) \langle \mathbb{E}^0_0/a_0 \rangle \dots \langle \mathbb{E}^n_0/a_n \rangle$ and $([a]v_1) \langle \mathbb{E}^0_1/a_0 \rangle \dots \langle \mathbb{E}^n_1/a_n \rangle$ are in \mathcal{X}_2 . We can prove the bisimulation property with these two named terms the same way we did with the named terms in \mathcal{X}_2 , except that we reason up to environment, because $\mathcal{E} \subseteq \mathcal{E}'$. Suppose u_0 is not a named value. There exists u'_0 such that $u_0 \rightarrow_v u'_0$, and so $u_0 \langle \mathbb{E}^0_0/a_0 \rangle \dots \langle \mathbb{E}^n_0/a_n \rangle \rightarrow_v u'_0 \langle \mathbb{E}^0_0/a_0 \rangle \dots \langle \mathbb{E}^n_0/a_n \rangle$. Because \mathcal{Y} is a bisimulation, there exists u'_1 such that $u_1 \rightarrow_n^* u'_1$ and $u'_0 \mathcal{Y}_{\mathcal{E}} u'_1$. We therefore have $u_1 \langle \mathbb{E}^0_1/a_0 \rangle \dots \langle \mathbb{E}^n_1/a_n \rangle \rightarrow_n^* u'_1 \langle \mathbb{E}^0_1/a_0 \rangle \dots \langle \mathbb{E}^n_1/a_n \rangle$ with $u'_0 \langle \mathbb{E}^0_0/a_0 \rangle \dots \langle \mathbb{E}^n_0/a_n \rangle \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} u'_1 \langle \mathbb{E}^0_1/a_0 \rangle \dots \langle \mathbb{E}^n_1/a_n \rangle$, as wished.

For the last bisimulation condition, let $[a]\lambda x.t_0 \widehat{\mathcal{E}}^{\text{nv}} [a]\lambda x.t_1$ and $t'_0 \widehat{\mathcal{E}} t'_1$. By Lemma 19, we have $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \mathcal{X}_{\widehat{\mathcal{E}}^{\text{nv}}} [a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$, hence the result holds.

Lemma 21. *If $[a]\lambda x.t_0 \simeq [a]\lambda x.t_1$ then $[a]t_0 \langle [a]\square t/a \rangle \{t/x\} \simeq [a]t_1 \langle [a]\square t/a \rangle \{t/x\}$*

Proof. Follows almost directly from the bisimilarity definition.

Lemma 22. *If $[a]\lambda x.t_0 \widehat{\mathcal{E}} [a]\lambda x.t_1$ and $t'_0 \widehat{\mathcal{E}} t'_1$ then $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \widehat{\mathcal{E}} [a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$.*

Proof. By case analysis on $[a]\lambda x.t_0 \widehat{\mathcal{E}} [a]\lambda x.t_1$. Suppose $[a]\lambda x.t_0 \simeq [a]\lambda x.t_1$. By Lemma 21, we have $[a]t_0 \langle [a]\square t'_1/a \rangle \{t'_1/x\} \simeq [a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$. From $t'_0 \widehat{\mathcal{E}} t'_1$, we deduce $t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \widehat{\mathcal{E}} t_0 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$, which implies $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \widehat{\mathcal{E}} [a]t_0 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$ (using Lemma 17). Finally, we get $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \widehat{\mathcal{E}} [a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$ as wished.

Suppose $\lambda x.t_0 \widehat{\mathcal{E}} \lambda x.t_1$. We distinguish two cases. If $\lambda x.t_0 \simeq \lambda x.t_1$, then $[a]\lambda x.t_0 \simeq [a]\lambda x.t_1$, and we proceed as before. Otherwise, we have $t_0 \widehat{\mathcal{E}} t_1$, which implies $t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \widehat{\mathcal{E}} t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$, which in turn implies $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \widehat{\mathcal{E}} [a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$ by Lemma 17, and finally we have $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \widehat{\mathcal{E}} [a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$, as wished.

Lemma 23. *Let $t_0 \widehat{\mathcal{E}} t_1$ and $a \notin \text{fn}(t_0) \cup \text{fn}(t_1)$.*

– *If $[a]t_0 \rightarrow_n u_0$, then $[a]t_1 \rightarrow_n^* u_1$, and $u_0 \widehat{\mathcal{E}} u_1$.*

- If $t_0 = \lambda x.t'_0$, then $[a]t_1 \rightarrow_n^* [a]\lambda x.t'_1$ and $[a]\lambda x.t'_0 \hat{\simeq} [a]v \simeq [a]\lambda x.t'_1$ for some v .

Proof. By induction on $t_0 \hat{\simeq} t_1$.

If $t_0 \simeq t_1$, then $[a]t_0 \simeq [a]t_1$, and the result holds by bisimilarity.

If $t_0 = \lambda x.t'_0$ and $t_1 = \lambda x.t'_1$ with $t'_0 \hat{\simeq} t'_1$, then $[a]t_0 \hat{\simeq} [a]t_1$, and the result holds.

Suppose $t_0 = \mu b.u_0$ and $t_1 = \mu b.u_1$ with $u_0 \hat{\simeq} u_1$. We have $[a]t_0 \rightarrow_n u_0 \langle [a]\square/b \rangle$, $[a]t_1 \rightarrow_n u_1 \langle [a]\square/b \rangle$, and $u_0 \langle [a]\square/b \rangle \hat{\simeq} u_1 \langle [a]\square/b \rangle$, hence the result holds.

Suppose $t_0 = t_0^1 t_0^2$, $t_1 = t_1^1 t_1^2$ with $t_0^1 \hat{\simeq} t_1^1$ and $t_0^2 \hat{\simeq} t_1^2$. We distinguish two cases.

- If $[b]t_0^1 \rightarrow_n u_0$ (for some fresh b), then $[a]t_0 \rightarrow_n u_0 \langle [a]\square/t_0^2/b \rangle$. By the induction hypothesis, there exist u_1, u such that $[b]t_1^1 \rightarrow_n^* u_1$ and $u_0 \hat{\simeq} u \simeq u_1$. Consequently, we have $[a]t_1 \rightarrow_n^* u_1 \langle [a]\square/t_1^2/b \rangle$. By congruence, we have $u_0 \langle [a]\square/t_0^2/b \rangle \hat{\simeq} u \langle [a]\square/t_1^2/b \rangle$, and by Lemma 18, we have $u \langle [a]\square/t_1^2/b \rangle \simeq u_1 \langle [a]\square/t_1^2/b \rangle$, hence the result holds.
- If $t_0^1 = \lambda x.t'_0$, then $[a]t_0 \rightarrow_n [a]t'_0 \{t_0^2/x\}$. By the induction hypothesis, there exist t, t'_1 such that $[a]t_1^1 \rightarrow_n^* [a]\lambda x.t'_1$ and $[a]\lambda x.t'_1 \hat{\simeq} [a]\lambda x.t \simeq [a]\lambda x.t'_1$. By Lemma 22, we have $[a]t'_0 \{t_0^2/x\} \hat{\simeq} [a]t \langle [a]\square/t_1^2/a \rangle \{t_1^2/x\}$, and by Lemma 21, we have $[a]t \langle [a]\square/t_1^2/a \rangle \{t_1^2/x\} \simeq [a]t'_1 \langle [a]\square/t_1^2/a \rangle \{t_1^2/x\}$. From $[a]t_1^1 \rightarrow_n^* [a]\lambda x.t'_1$, we deduce $[a]t_1 \rightarrow_n^* [a](\lambda x.t'_1 \langle [a]\square/t_1^2/a \rangle)t_1^2 \rightarrow_n [a]t'_1 \langle [a]\square/t_1^2/a \rangle \{t_1^2/x\}$, hence the result holds.

Suppose $t_0 = t'_0 \langle \mathbb{E}_0/b \rangle$, $t_1 = t'_1 \langle \mathbb{E}_1/b \rangle$ with $t'_0 \hat{\simeq} t'_1$, $\mathbb{E}_0 \hat{\simeq} \mathbb{E}_1$. If $[a]t_0 \rightarrow_n u_0$, then in fact $[a]t'_0 \rightarrow_n u'_0$ with $u_0 = u'_0 \langle \mathbb{E}_0/b \rangle$. By the induction hypothesis, there exists u'_1 such that $[a]t'_1 \rightarrow_n^* u'_1$ and $u'_0 \hat{\simeq} u'_1$. Consequently, we have $[a]t_1 \rightarrow_n^* u'_1 \langle \mathbb{E}_1/b \rangle$, and $u'_0 \langle \mathbb{E}_0/b \rangle \hat{\simeq} u'_1 \langle \mathbb{E}_1/b \rangle$ holds (using Lemma 18), as wished. If $[a]t_0$ is a named value, then in fact $t'_0 = \lambda x.t''_0$ and $t_0 = \lambda x.t''_0 \langle \mathbb{E}_0/b \rangle$. By induction, there exist t''_1 such that $[a]t'_1 \rightarrow_n^* [a]\lambda x.t''_1$ and $[a]\lambda x.t''_1 \hat{\simeq} [a]\lambda x.t''_1$. Hence, we have $[a]t_1 \rightarrow_n^* [a]\lambda x.t''_1 \langle \mathbb{E}_1/b \rangle$, and $[a]\lambda x.t''_0 \langle \mathbb{E}_0/b \rangle \hat{\simeq} [a]\lambda x.t''_1 \langle \mathbb{E}_1/b \rangle$ holds using Lemma 18.

Lemma 24. $u_0 \simeq u_1$ implies $C[u_0] \cong_{\hat{\simeq}^{nv}} C[u_1]$ and $\mathbb{C}[u_0] \cong_{\hat{\simeq}^{nv}} \mathbb{C}[u_1]$.

Proof. We prove that

$$\mathcal{X} = \{(\hat{\simeq}^{nv}, u_0, u_1), u_0 \hat{\simeq} u_1\} \cup \{\hat{\simeq}^{nv}\}$$

is a bisimulation up-to bisimilarity.

Let $u_0 \hat{\simeq} u_1$; we proceed by induction on $u_0 \hat{\simeq} u_1$.

If $u_0 \simeq u_1$, then the result holds by bisimilarity.

Suppose $u_0 = [a]t_0$, $u_1 = [a]t_1$, with $t_0 \hat{\simeq} t_1$. If $a \notin \text{fn}(t_0) \cup \text{fn}(t_1)$, then we can apply Lemma 23 directly to get the required result. Otherwise, let $b \notin \text{fn}(t_0) \cup \text{fn}(t_1)$. If $u_0 \rightarrow_n u'_0$, then $[b]t_0 \rightarrow_n u''_0$ and $u'_0 = u'_0 \langle [a]\square/b \rangle$. We can apply Lemma 23 to $[b]t_0$ and $[b]t_1$, and then rename b into a .

Suppose $u_0 = u'_0 \langle \mathbb{E}_0/a \rangle$, $u_1 = u'_1 \langle \mathbb{E}_1/a \rangle$ with $u'_0 \hat{\simeq} u'_1$ and $\mathbb{E}_0 \hat{\simeq} \mathbb{E}_1$.

If u_0 is a named value, then $u'_0 = [c]v_0$, and $u_0 = [b]v_0 \langle \mathbb{E}_0/a \rangle$. By the induction hypothesis, there exists v_1 such that $u'_1 \rightarrow_n^* [c]v_1$ and $[c]v_0 \hat{\simeq} [c]v_1$. We have $u_1 \rightarrow_n^* [b]v_1 \langle \mathbb{E}_1/a \rangle$, and also $[b]v_0 \langle \mathbb{E}_0/a \rangle \hat{\simeq} [b]v_1 \langle \mathbb{E}_1/a \rangle$ (using Lemma 18), hence the result holds.

If $u_0 \rightarrow_n$, then we distinguish two cases. First, suppose $u'_0 \rightarrow_n u''_0$; then $u_0 \rightarrow_n u''_0 \langle \mathbb{E}_0/a \rangle$. By the induction hypothesis, there exists u''_1 such that $u'_1 \rightarrow_n^* u''_1$ and $u''_0 \hat{\simeq} u''_1$. Then $u_1 \rightarrow_n^* u''_1 \langle \mathbb{E}_1/a \rangle$ and $u''_0 \langle \mathbb{E}_0/a \rangle \hat{\simeq} u''_1 \langle \mathbb{E}_1/a \rangle$ (using Lemma 18), hence the result holds.

Otherwise, we have $u'_0 = [a]\lambda x.t_0$, $\mathbb{E}_0 = \mathbb{E}'_0[\square t'_0]$, which implies $u_0 \rightarrow_n \mathbb{E}'_0[t_0 \langle \mathbb{E}_0/a \rangle \{t'_0/x\}]$. We can prove by induction on $\mathbb{E}_0 \hat{\simeq} \mathbb{E}_1$ that $\mathbb{E}_1 = \mathbb{E}'_1[\square t'_1]$ with $\mathbb{E}'_0 \hat{\simeq} \mathbb{E}'_1$ and $t'_0 \hat{\simeq} t'_1$. Let \mathbb{E}''_i, t''_i be \mathbb{E}'_i, t'_i with a renamed into a fresh b . By the induction hypothesis, there exist t, t_1 such that $u'_1 \rightarrow_n^* [a]\lambda x.t_1$, and $[a]\lambda x.t_0 \hat{\simeq} [a]\lambda x.t \simeq [a]\lambda x.t_1$. By Lemma 22 and Lemma 21, we have $[a]t_0 \langle [a]\square t'_0/a \rangle \{t''_0/x\} \hat{\simeq} [a]t \langle [a]\square t'_1/a \rangle \{t''_1/x\} \simeq [a]t_1 \langle [a]\square t'_1/a \rangle \{t''_1/x\}$, which in turn implies $([a]t_0 \langle [a]\square t'_0/a \rangle \{t''_0/x\}) \langle \mathbb{E}''_0/a \rangle \hat{\simeq} ([a]t_1 \langle [a]\square t'_1/a \rangle \{t''_1/x\}) \langle \mathbb{E}''_1/a \rangle$ using Lemma 18. Because a does not occurs in t''_0, t''_1 , the previous relation is equivalent to $\mathbb{E}''_0[t_0 \langle \mathbb{E}''_0[\square t'_0/a] \{t''_0/x\} \rangle] \hat{\simeq} \mathbb{E}''_1[t_1 \langle \mathbb{E}''_1[\square t'_1/a] \{t''_1/x\} \rangle]$. Renaming b into a , we obtain $\mathbb{E}'_0[t_0 \langle \mathbb{E}_0/a \rangle \{t'_0/x\}] \hat{\simeq} \mathbb{E}'_1[t_1 \langle \mathbb{E}_1/a \rangle \{t'_1/x\}]$, and because $u_1 \rightarrow_n^* \mathbb{E}'_1[t_1 \langle \mathbb{E}_1/a \rangle \{t'_1/x\}]$, we have the required result.

To prove the last item, we take $[a]\lambda x.t_0 \hat{\simeq} [a]\lambda x.t_1$ and $t'_0 \hat{\simeq} t'_1$. Then $[a]t_0 \langle [a]\square t'_0/a \rangle \{t'_0/x\} \hat{\simeq} [a]t_1 \langle [a]\square t'_1/a \rangle \{t'_1/x\}$ holds by Lemma 22.

B Call-by-Value $\lambda\mu$ -calculus

B.1 Equivalence Proofs

Lemma 25 (see Example 7). *The relation $\mathcal{R} \stackrel{def}{=} \{(u\{v/y\}, u\{\lambda x.\mu b.[a]t_0/y\}) \mid [a]t_0 \rightarrow_v^* u\{v/y\}\} \cup \approx$ is an applicative bisimulation.*

Proof. We proceed by case analysis on u . If $u \rightarrow_v u'$ or if $u = \mathbb{E}[v' y]$, then it is easy to conclude. If $u = \mathbb{E}[y v']$, then $u\{\lambda x.\mu b.[a]t_0/y\} \rightarrow_v [a]t_0$ and $[a]t_0 \rightarrow_v^* u\{v/y\}$ by definition, so it we can conclude with Lemma 11. Similarly, the transition from $u\{v/y\}$ is matched by $u\{\lambda x.\mu b.[a]t_0/y\}$.

If $u = [c]v'$, then we have to compare (the result of the reduction of) $u \langle \mathbb{E}/c \rangle \{v \langle \mathbb{E}/c \rangle / y\}$ and $u \langle \mathbb{E}/c \rangle \{\lambda x.\mu b.([a]t_0) \langle \mathbb{E}/c \rangle / y\}$ for some \mathbb{E} (which depends on the clause we check). The resulting terms are in \mathcal{R} , because $[a]t_0 \rightarrow_v^* u\{v/y\}$ implies $([a]t_0) \langle \mathbb{E}/c \rangle \rightarrow_v^* u \langle \mathbb{E}/c \rangle \{v \langle \mathbb{E}/c \rangle / y\}$.

For the next lemma and its proof, we use the same definitions of terms as in Example 9.

Lemma 26 (see Example 9). *The relation*

$$\mathcal{R} \stackrel{def}{=} \{(u \langle [a]E[v_0 \langle [a]E[\square t_1/a] \square] / b \rangle, u \langle [a]E[v_1 \square] / b \rangle) \mid [b]t \Downarrow_v [b]v, b \notin \text{fn}(t)\}$$

is an applicative bisimulation.

Proof. We proceed by case analysis on u . If $u \rightarrow_v u'$, then the result holds. If $u = [c]v'$ with $c \neq a$ and $c \neq b$, then when checking clauses 1 and 2, we obtain terms that can be written $u' \langle [a]E[v_0 \langle [a]E[\square \ t]/a \rangle \square] / b \rangle$ and $u' \langle [a]E[v_1 \square] / b \rangle$ for some u' , and are therefore in \mathcal{R} .

If $u = [a]v'$, then when checking clauses 1 and 2, we obtain terms that can be written $u' \langle [a]E'[E[v_0 \langle [a]E'[\square \ t]/a \rangle \square]] / b \rangle$ and $u' \langle [a]E'[E[v_1 \square]] / b \rangle$ for some u' and E' (depending on which clause we check). We obtain terms in \mathcal{R} .

If $u = [b]v'$, then

$$\begin{aligned} u \langle [a]E[v_0 \langle [a]E[\square \ t]/a \rangle \square] / b \rangle &= [a]E[v_0 \langle [a]E[\square \ t]/a \rangle v' \langle [a]E[v_0 \langle [a]E[\square \ t]/a \rangle \square] / b \rangle] \\ &\rightarrow_v^* [a]E[w' (v' \langle [a]E[v_0 \langle [a]E[\square \ t]/a \rangle \square] / b \rangle \lambda x.x)] \end{aligned}$$

and $u \langle [a]E[v_1 \square] / b \rangle = [a]E[v_1 v' \langle [a]E[v_1 \square] / b \rangle] \rightarrow_v^* [a]E[w' (v' \langle [a]E[v_1 \square] / b \rangle \lambda x.x)]$; we obtain terms in \mathcal{R} .

The next lemma uses the notations of Section 3.4

Lemma 27 (see Example 12). *Let \mathbb{E} , \mathbb{E}_0 , and E_1 be such that $\mathbb{E} = \mathbb{E}_0[E_1]$. Then $\mathbb{E}^\dagger \approx \lambda x. \mathbb{E}_0^\dagger E_1[x]$.*

Proof. Let a be a fresh name. Let t such that $[d]t \Downarrow_v [d]v$. To check clause 1, we have to relate $[a]\mu c. \mathbb{E}[v \langle [a]\mathbb{E}^\dagger \square / d \rangle]$ and $[a]\mathbb{E}_0^\dagger E_1[v \langle [a]\lambda x. \mathbb{E}_0^\dagger E_1[x] \square / d \rangle]$. These terms reduce respectively to $\mathbb{E}[v \langle [a]\mathbb{E}^\dagger \square / d \rangle]$ and $\mathbb{E}_0[E_1[v \langle [a]\lambda x. \mathbb{E}_0^\dagger E_1[x] \square / d \rangle]]$, which we can rewrite into respectively $u \langle [a]\mathbb{E}^\dagger \square / d \rangle$ and $u \langle [a]\lambda x. \mathbb{E}_0^\dagger E_1[x] \square / d \rangle$. To check clause 2, we have to relate $[a]t\{\mathbb{E}^\dagger / y\}$ and $[a]t\{\lambda x. \mathbb{E}_0^\dagger E_1[x] / y\}$ for all t . The most interesting case is when $t = E[yv]$; then these terms reduce to respectively $\mathbb{E}[v]\{\mathbb{E}^\dagger / y\}$ and $\mathbb{E}[v]\{\lambda x. \mathbb{E}_0^\dagger E_1[x] / y\}$. In fact, one can prove the result by showing that

$$\{(u \langle [a]\mathbb{E}^\dagger \square / d \rangle), u \langle [a]\lambda x. \mathbb{E}_0^\dagger E_1[x] \square / d \rangle, (u\{\mathbb{E}^\dagger / y\}), u\{\lambda x. \mathbb{E}_0^\dagger E_1[x] / y\}\} \mid \mathbb{E} = \mathbb{E}_0[E_1]$$

is an applicative bisimulation.

B.2 Proof of Soundness

Lemma 28. *If $[a]v_0 \approx [a]v_1$, then for all $[b]t \rightarrow_v^* u$ ($b \notin \text{fn}(t)$), we have $u \langle [a]v_0 \langle [a]\square \ t/a \rangle \square / b \rangle \approx u \langle [a]v_1 \langle [a]\square \ t/a \rangle \square / b \rangle$.*

Proof. The relation $\{(u \langle [a]v_0 \langle [a]\square \ t/a \rangle \square / b \rangle), u \langle [a]v_1 \langle [a]\square \ t/a \rangle \square / b \rangle\} \mid \forall u, [b]t \rightarrow_v^* u\} \cup \approx$ is a bisimulation.

Lemma 29. *If $\lambda x. t_0 (\approx^\bullet)^c t_1$, then there exists t such that $t_0 \approx^\bullet t$, $\text{fv}(t) \subseteq \{x\}$ and $\lambda x. t \approx t_1$.*

Proof. By induction on $\lambda x. t_0 (\approx^\bullet)^c t_1$.

Lemma 30. *If $[a]t_0 (\approx^\bullet)^c u_1$, then there exists a closed t such that $t_0 (\approx^\bullet)^c t$ and $[a]t \approx u_1$.*

Proof. By induction on $[a]t_0 (\approx^\bullet)^c u_1$.

Lemma 31. *Let $t_0 (\approx^\bullet)^c t_1$, and $a \notin \text{fn}(t_0) \cup \text{fn}(t_1)$.*

- If $[a]t_0 \rightarrow_n u_0$, then $[a]t_1 \rightarrow_n^* u_1$ and $u_0 (\approx^\bullet)^c u_1$.
- If $t_0 = \lambda x.t'_0$, then $[a]t_1 \rightarrow_n^* [a]\lambda x.t'_1$ and for all t, b, v such that $[b]t \rightarrow_v^* [b]v$ and $b \notin \text{fn}(t)$, we have

$$[a]t'_0 \langle [a]\square t/a \rangle \{v \langle [a]\lambda x.t_0 \langle [a]\square t/a \rangle \square/b \rangle /x\} (\approx^\bullet)^c \\ [a]t'_1 \langle [a]\square t/a \rangle \{v \langle [a]\lambda x.t_1 \langle [a]\square t/a \rangle \square/b \rangle /x\}$$

and for all $v' = \lambda x.t'$, we have

$$[a]t' \langle \lambda x.t'_0 \langle [a]v \square/a \rangle /x \rangle (\approx^\bullet)^c [a]t' \langle \lambda x.t'_1 \langle [a]v \square/a \rangle /x \rangle.$$

Proof. By induction on $t_0 (\approx^\bullet)^c t_1$. If $t_0 \approx t_1$, then the result holds by bisimilarity. If $t_0 \approx^\bullet t \approx^\circ t_1$, so that t is closed (using Lemma 13 if necessary), then we can conclude with the induction hypothesis and the bisimilarity definition.

If $t_0 (\approx^\bullet)^c t_1$, then we have several cases to consider.

If $t_0 = \lambda x.t'_0$ and $t_1 = \lambda x.t'_1$ with $t'_0 \approx^\bullet t'_1$, we can prove the required result because \approx^\bullet is substitutive. Suppose $t_0 = \mu b.u_0$ and $t_1 = \mu b.u_1$ with $u_0 (\approx^\bullet)^c u_1$. We have $[a]t_0 \rightarrow_v u_0 \langle [a]\square/b \rangle$, $[a]t_1 \rightarrow_v u_1 \langle [a]\square/b \rangle$, and $u_0 \langle [a]\square/b \rangle (\approx^\bullet)^c u_1 \langle [a]\square/b \rangle$, hence the result holds.

Suppose $t_0 = t_0^1 t_0^2$, $t_1 = t_1^1 t_1^2$ with $t_0^1 (\approx^\bullet)^c t_1^1$ and $t_0^2 (\approx^\bullet)^c t_1^2$. We distinguish three cases.

- If $[b]t_0^1 \rightarrow_v u_0$ (for some fresh b), then $[a]t_0 \rightarrow_v u_0 \langle [a]\square t_0^2/b \rangle$. By the induction hypothesis, there exists u_1 such that $[b]t_1^1 \rightarrow_v^* u_1$ and $u_0 (\approx^\bullet)^c u_1$. Consequently, we have $[a]t_1 \rightarrow_v^* u_1 \langle [a]\square t_1^2/b \rangle$, and by definition of \approx^\bullet , we have $u_0 \langle [a]\square t_0^2/b \rangle (\approx^\bullet)^c u_1 \langle [a]\square t_1^2/b \rangle$, as required.
- Suppose $t_0^1 = \lambda x.t'_0$ and $[b]t_0^2 \rightarrow_v u_0$ for some fresh b ; then we have $[a]t_0 \rightarrow_v u_0 \langle [a]\lambda x.t'_0 \square/b \rangle$. By the induction hypothesis, there exists u_1 such that $[b]t_1^2 \rightarrow_v^* u_1$ and $u_0 (\approx^\bullet)^c u_1$. Because $\lambda x.t'_0 (\approx^\bullet)^c t_1^1$, there exists a t such that $t'_0 \approx^\bullet t$ and $\lambda x.t \approx t_1^1$ by Lemma 29. From $t'_0 \approx^\bullet t$, we deduce $[a]\lambda x.t'_0 \square (\approx^\bullet)^c [a]\lambda x.t \langle [a]\square t_1^1/a \rangle \square$, which implies $u_0 \langle [a]\lambda x.t'_0 \square/b \rangle (\approx^\bullet)^c u_1 \langle [a]\lambda x.t \langle [a]\square t_1^1/a \rangle \square/b \rangle$. Because $\lambda x.t \approx t_1^1$, there exists v_1 such that $[a]t_1^1 \rightarrow_v^* [a]v_1$, and we have $[a]\lambda x.t \approx [a]v_1$. By Lemma 28, $u_1 \langle [a]\lambda x.t \langle [a]\square t_1^1/a \rangle \square/b \rangle \approx u_1 \langle [a]v_1 \langle [a]\square t_1^1/a \rangle \square/b \rangle$ holds, so $u_0 \langle [a]\lambda x.t'_0 \square/b \rangle (\approx^\bullet)^c u_1 \langle [a]v_1 \langle [a]\square t_1^1/a \rangle \square/b \rangle$ holds as well. Besides, we have $[a]t_1 \rightarrow_v^* u_1 \langle [a]v_1 \langle [a]\square t_1^2/a \rangle \square/b \rangle$, hence we have the required result.
- If $t_0^1 = \lambda x.t'_0$ and $t_0^2 = \lambda y.t''_0$, then $[a]t_0 \rightarrow_v [a]t'_0 \langle \lambda x.t''_0/x \rangle$. By Lemma 29, there exist closed t, t' such that $t'_0 (\approx^\bullet)^c t$, $\lambda x.t \approx t_1^1$, $t''_0 (\approx^\bullet)^c t'$, and $\lambda y.t' \approx t_1^2$. From $t'_0 (\approx^\bullet)^c t$, we deduce $t'_0 (\approx^\bullet)^c t \langle [a]\square t_1^1/a \rangle$, and from $t''_0 (\approx^\bullet)^c t'$, we deduce $\lambda y.t''_0 (\approx^\bullet)^c \lambda y.t' \langle [a]\lambda x.t \langle [a]\square t_1^1/a \rangle \square/a \rangle$. Consequently, we have $[a]t'_0 \langle \lambda y.t''_0/x \rangle (\approx^\bullet)^c [a]t \langle [a]\square t_1^1/a \rangle \langle \lambda y.t' \langle [a]\lambda x.t \langle [a]\square t_1^1/a \rangle \square/a \rangle /x \rangle$.

From $\lambda y.t' \approx t_1^2$, by bisimilarity, there exists v_1 such that $[a]t_1^2 \rightarrow_v^* [a]v_1$, and

$$\begin{aligned} [a]t \langle [a] \square t_1^2/a \rangle \{ \lambda y.t' \langle [a] \lambda x.t \langle [a] \square t_1^2/a \rangle \square/a \rangle /x \} \\ \approx [a]t \langle [a] \square t_1^2/a \rangle \{ v_1 \langle [a] \lambda x.t \langle [a] \square t_1^2/a \rangle \square/a \rangle /x \} \end{aligned}$$

holds (clause 2). From $\lambda x.t \approx t_1^1$, by bisimilarity, there exists t_1 such that $[a]t_1^1 \rightarrow_v^* [a]\lambda x.t_1$, and

$$\begin{aligned} [a]t \langle [a] \square t_1^2/a \rangle \{ v_1 \langle [a] \lambda x.t \langle [a] \square t_1^2/a \rangle \square/a \rangle /x \} \\ \approx [a]t_1 \langle [a] \square t_1^2/a \rangle \{ v_1 \langle [a] \lambda x.t_1 \langle [a] \square t_1^2/a \rangle \square/a \rangle /x \} \end{aligned}$$

holds (clause 1). Finally, we have

$$[a]t_0 \{ \lambda y.t_0''/x \} (\approx^\bullet)^c [a]t_1 \langle [a] \square t_1^2/a \rangle \{ v_1 \langle [a] \lambda x.t_1 \langle [a] \square t_1^2/a \rangle \square/a \rangle /x \},$$

and because $[a]t_1 \rightarrow_v^* [a]t_1 \langle [a] \square t_1^2/a \rangle \{ v_1 \langle [a] \lambda x.t_1 \langle [a] \square t_1^2/a \rangle \square/a \rangle /x \}$, we have the required result.

Suppose $t_0 = t_0' \langle \mathbb{E}_0/b \rangle$, $t_1 = t_1' \langle \mathbb{E}_1/b \rangle$ with $t_0' (\approx^\bullet)^c t_1'$, $\mathbb{E}_0 (\approx^\bullet)^c \mathbb{E}_1$. If $[a]t_0 \rightarrow_v u_0$, then in fact $[a]t_0' \rightarrow_v u_0'$ with $u_0 = u_0' \langle \mathbb{E}_0/b \rangle$. By the induction hypothesis, there exists u_1' such that $[a]t_1' \rightarrow_v^* u_1'$ and $u_0' (\approx^\bullet)^c u_1'$. Consequently, we have $[a]t_1 \rightarrow_v^* u_1' \langle \mathbb{E}_1/b \rangle$, and $u_0' \langle \mathbb{E}_0/b \rangle (\approx^\bullet)^c u_1' \langle \mathbb{E}_1/b \rangle$ holds, as wished. If $[a]t_0$ is a named value, then in fact $t_0' = \lambda x.t_0'''$ and $t_0 = \lambda x.t_0''' \langle \mathbb{E}_0/b \rangle$. The result holds by using the induction hypothesis (and with some renaming of b into a fresh c to avoid name clashes).

Lemma 32. *The relation $(\approx^\bullet)^c$ is an applicative simulation.*

Proof. Let $u_0 (\approx^\bullet)^c u_1$; we prove the simulation clause by induction on $u_0 (\approx^\bullet)^c u_1$.

If $u_0 \approx u_1$, then the result holds by bisimilarity. If $u_0 \approx^\bullet u \approx^\circ u_1$, then we can conclude using the induction hypothesis and the bisimilarity definition.

If $u_0 (\approx^\bullet)^c u_1$, then we have two cases. Suppose $u_0 = [a]t_0$ and $u_1 = [a]t_1$ with $t_0 (\approx^\bullet)^c t_1$. If $a \notin \text{fn}(t_0) \cup \text{fn}(t_1)$, then we can apply Lemma 31 directly to get the required result. Otherwise, let $b \notin \text{fn}(t_0) \cup \text{fn}(t_1)$. If $u_0 \rightarrow_v u_0'$, then $[b]t_0 \rightarrow_v u_0'$ and $u_0' = u_0' \langle [a] \square/b \rangle$. We can apply Lemma 31 to $[b]t_0$ and $[b]t_1$, and then rename b into a .

Suppose $u_0 = u_0' \langle \mathbb{E}_0/a \rangle$, $u_1 = u_1' \langle \mathbb{E}_1/a \rangle$ with $u_0' (\approx^\bullet)^c u_1'$ and $\mathbb{E}_0 (\approx^\bullet)^c \mathbb{E}_1$.

If u_0 is a named value, then we use the induction hypothesis, with some renaming (as in the call-by-name case) to avoid some name clashes.

If $u_0 \rightarrow_v$, then again we distinguish several cases. First, suppose $u_0' \rightarrow_v u_0''$; then $u_0 \rightarrow_v u_0'' \langle \mathbb{E}_0/a \rangle$. By the induction hypothesis, there exists u_1'' such that $u_1' \rightarrow_v^* u_1''$ and $u_0'' (\approx^\bullet)^c u_1''$. Then $u_1 \rightarrow_v^* u_1'' \langle \mathbb{E}_1/a \rangle$ and $u_0'' \langle \mathbb{E}_0/a \rangle (\approx^\bullet)^c u_1'' \langle \mathbb{E}_1/a \rangle$, hence the result holds.

Second, assume $u'_0 = [a]\lambda x.t_0$, $\mathbb{E}_0 = \mathbb{E}'_0[\square v_0]$ and $u_0 \rightarrow_v \mathbb{E}'_0[t_0\langle\mathbb{E}_0/a\rangle\{v_0/x\}]$. By Lemma 30, there exists t such that $\lambda x.t_0 (\approx^\bullet)^c t$ and $[a]t \approx u'_1$. By Lemma 29, there exists t' such that $t_0 \approx^\bullet t'$ and $\lambda x.t' \approx t$. We can prove by induction on $\mathbb{E}_0 (\approx^\bullet)^c \mathbb{E}_1$ that $\mathbb{E}_1 = \mathbb{E}'_1[\square t'_1]$ with $\mathbb{E}'_0 (\approx^\bullet)^c \mathbb{E}'_1$ and $v_0 (\approx^\bullet)^c t'_1$. Suppose $v_0 = \lambda x.t'_0$. By Lemma 29, there exists s such that $t'_0 \approx^\bullet s$ and $\lambda x.s \approx t'_1$. Let $\mathbb{E}''_i, s', t''_i$ be $\mathbb{E}'_i, s,$ and t'_i with a renamed into a fresh c . From $t_0 (\approx^\bullet)^c t', \lambda x.t''_0 (\approx^\bullet)^c t'_1,$ and $\lambda x.t''_0 (\approx^\bullet)^c \lambda x.s',$ we deduce

$$[a]t_0\langle[a]\square \lambda x.t''_0/a\rangle\{\lambda x.t''_0/x\} (\approx^\bullet)^c [a]t'\langle[a]\square t'_1/a\rangle\{\lambda x.s'/x\}.$$

Because $\lambda x.s' \approx t'_1$, there exists v'_1 such that $[b]t''_1 \rightarrow_v^* [b]v'_1$ (for a fresh b), and $[a]t'\langle[a]\square t'_1/a\rangle\{\lambda x.s'/x\} \approx [a]t'\langle[a]\square t'_1/a\rangle\{v'_1\langle[a]\lambda x.t'\langle[a]\square t'_1/a\rangle\square/b\rangle/x\}$ (using the second item of the value case of the bisimulation definition). Because $\lambda x.t' \approx t$, there exists t'' such that $[a]t \rightarrow_v^* [a]\lambda x.t''$ and

$$\begin{aligned} [a]t'\langle[a]\square t'_1/a\rangle\{v'_1\langle[a]\lambda x.t'\langle[a]\square t'_1/a\rangle\square/b\rangle/x\} \\ \approx [a]t''\langle[a]\square t'_1/a\rangle\{v'_1\langle[a]\lambda x.t''\langle[a]\square t'_1/a\rangle\square/b\rangle/x\} \end{aligned}$$

(by clause 1). Because $[a]\lambda x.t' \approx u'_1$, there exists t_1 such that $u'_1 \rightarrow_v^* [a]\lambda x.t_1$ and we have also

$$\begin{aligned} [a]t''\langle[a]\square t'_1/a\rangle\{v'_1\langle[a]\lambda x.t''\langle[a]\square t'_1/a\rangle\square/b\rangle/x\} \\ \approx [a]t_1\langle[a]\square t'_1/a\rangle\{v'_1\langle[a]\lambda x.t_1\langle[a]\square t'_1/a\rangle\square/b\rangle/x\} \end{aligned}$$

by clause 1. Finally, we obtain

$$[a]t_0\langle[a]\square \lambda x.t''_0/a\rangle\{\lambda x.t''_0/x\} (\approx^\bullet)^c [a]t_1\langle[a]\square t'_1/a\rangle\{v'_1\langle[a]\lambda x.t_1\langle[a]\square t'_1/a\rangle\square/b\rangle/x\}$$

by transitivity of \approx and definition of \approx^\bullet , from which we can deduce

$$\begin{aligned} \mathbb{E}''_0[t_0\langle\mathbb{E}''_0[\square \lambda x.t''_0/a]\rangle\{\lambda x.t''_0/x\}] \\ (\approx^\bullet)^c \mathbb{E}''_1[t_1\langle\mathbb{E}''_1[\square t'_1/a]\rangle\{v'_1\langle\mathbb{E}''_1[\lambda x.t_1\langle\mathbb{E}''_1[\square t'_1/a]\rangle\square/b\rangle/x\}]. \end{aligned}$$

Renaming c back into a , we get

$$\begin{aligned} \mathbb{E}'_0[t_0\langle\mathbb{E}'_0[\square \lambda x.t'_0/a]\rangle\{\lambda x.t'_0/x\}] \\ (\approx^\bullet)^c \mathbb{E}'_1[t_1\langle\mathbb{E}'_1[\square t'_1/a]\rangle\{v_1\langle\mathbb{E}'_1[\lambda x.t_1\langle\mathbb{E}'_1[\square t'_1/a]\rangle\square/b\rangle/x\}] \end{aligned}$$

(assuming v_1 is the result of renaming c into a in v'_1), which is the same as

$$\mathbb{E}'_0[t_0\langle\mathbb{E}_0/a\rangle\{\lambda x.t'_0/x\}] (\approx^\bullet)^c \mathbb{E}'_1[t_1\langle\mathbb{E}_1/a\rangle\{v_1\langle\mathbb{E}'_1[\lambda x.t_1\langle\mathbb{E}_1/a\rangle\square/b\rangle/x\}].$$

One can check that $u_1 \rightarrow_v^* \mathbb{E}'_1[t_1\langle\mathbb{E}_1/a\rangle\{v_1\langle\mathbb{E}'_1[\lambda x.t_1\langle\mathbb{E}_1/a\rangle\square/b\rangle/x\}]$, hence the result holds.

Finally, the last case is $u'_0 = [a]v_0$, $\mathbb{E}_0 = \mathbb{E}'_0[\lambda x.t_0 \square]$, which give $u_0 \rightarrow_v \mathbb{E}'_0[t_0\{v_0\langle\mathbb{E}_0/a\rangle/x\}]$. This case is similar to the previous one and is left to the reader.



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