

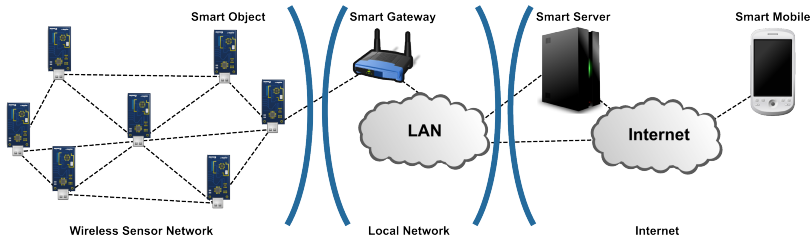
CCN Traffic Optimization for IoT

Jérôme François, Thibault Cholez, Thomas Engel

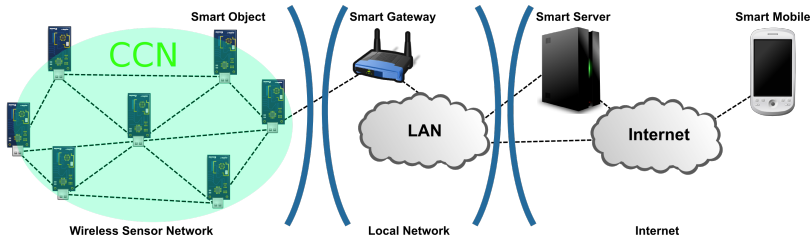
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- ▶ Internet of Things
 - ▶ multi sites network with sensors and gateways
 - ▶ **shared infrastructure** of multiple owners for multiple service providers
 - ▶ reliability, security, privacy... requirements



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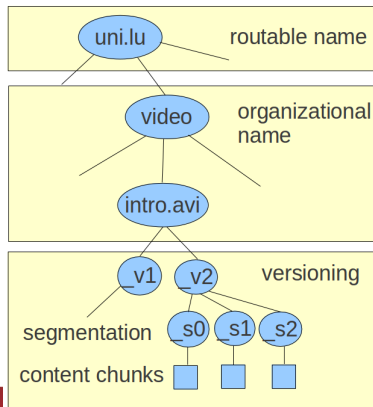
- ▶ Why?
 - ▶ CCN designed for distributed networking → scalability
 - ▶ ICN/CCN provides already many required basic functionalities (security, naming)
 - ▶ → alleviates the need of additional protocol, e.g. resource discovery protocols
 - ▶ Contiki implementation available
- ▶ but:
 - ▶ CCN: pull-based mode (on demand information retrieving)
 - ▶ many IoT use cases where: push based model for sensor (regular sensing or event based)
- ▶ → can we still use CCN for sensor communications on a push based manner?

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- ▶ Key ideas
 - ▶ Data as a name, not a location
 - ▶ **Data is directly requested at the network level** (no more DNS: more security; less delay)
 - ▶ Anybody with the data can answer
 - ▶ **Data is authenticated**, not the connections it traverses
 - ▶ + caching at intermediate routers
- ▶ 2 types of packets *Interest* and *Data*
 - ▶ Consumer driven: broadcast Interest, waits for Data
 - ▶ Data transmitted in response, consumes the Interest

Hierarchical naming

- ▶ Routable **prefix**, persistence of the names
- ▶ Hierarchical **aggregation** for fast routing and forwarding
- ▶ Dynamic generation of content
- ▶ Context-aware resolution (*/ThisRoom/projector*)



Message sequence

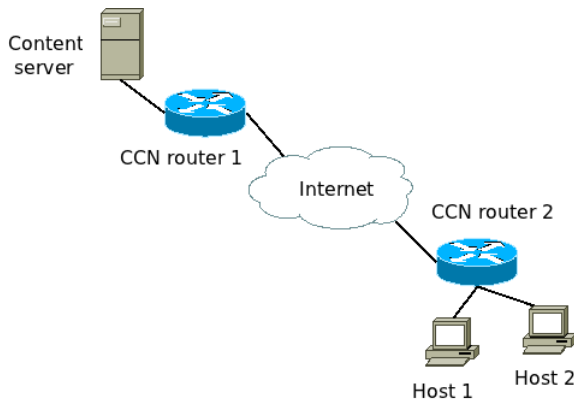


Figure: Scheme of a minimal Content-Centric network

Message sequence

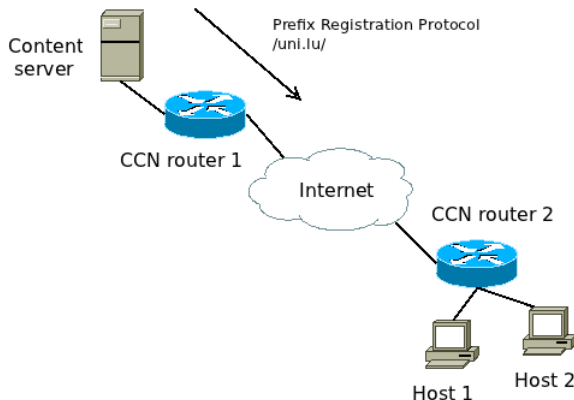


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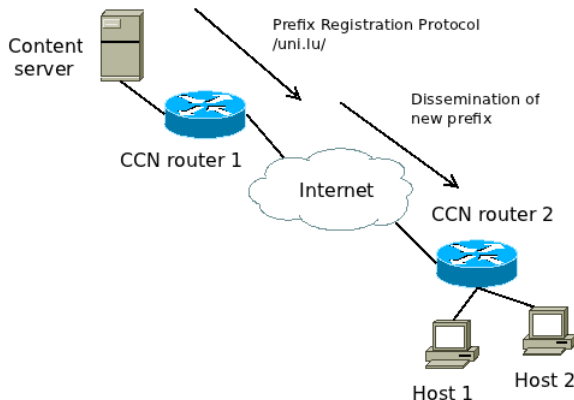


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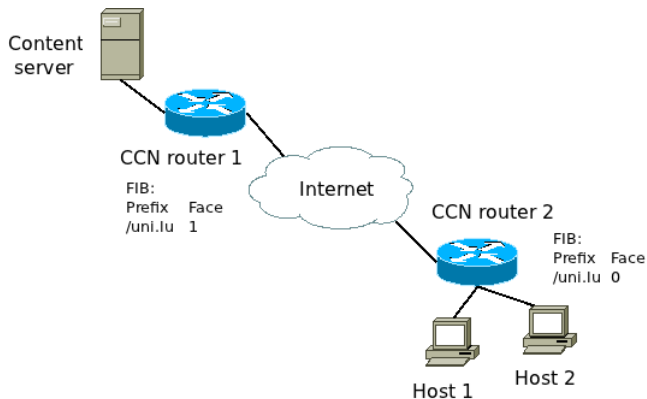


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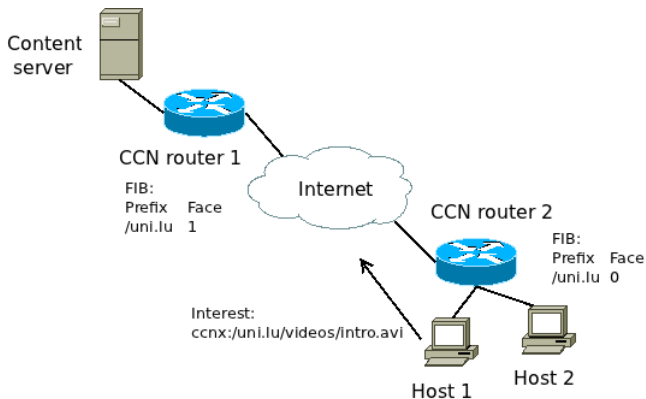


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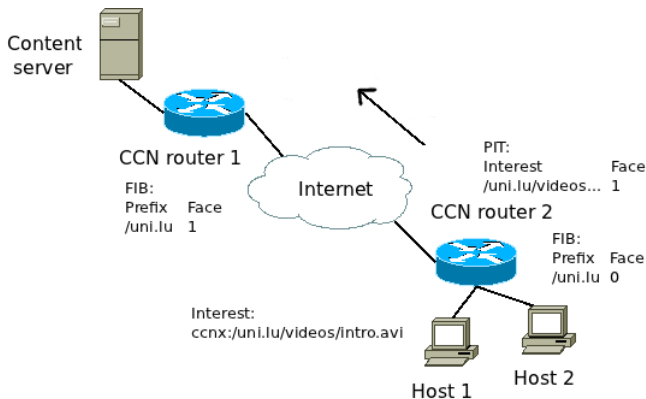


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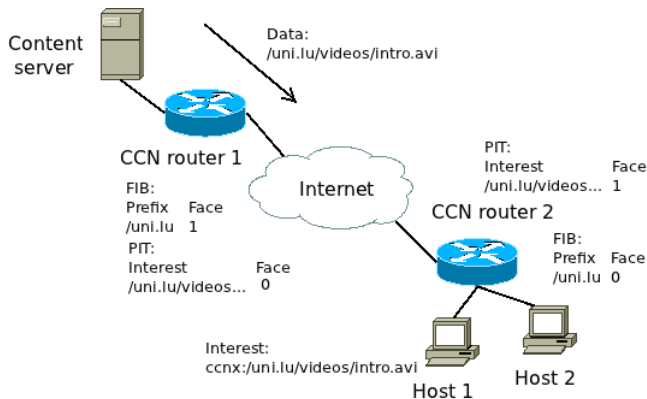


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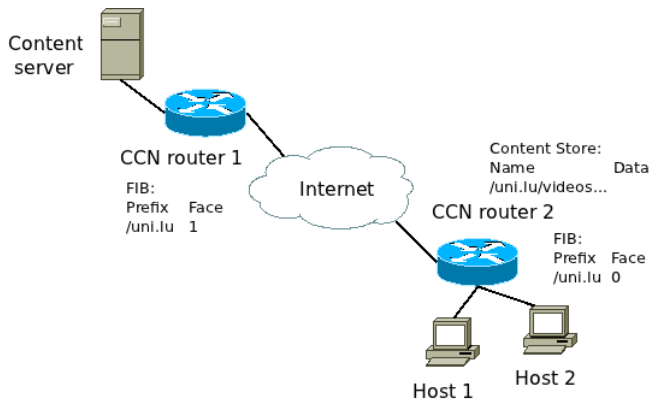


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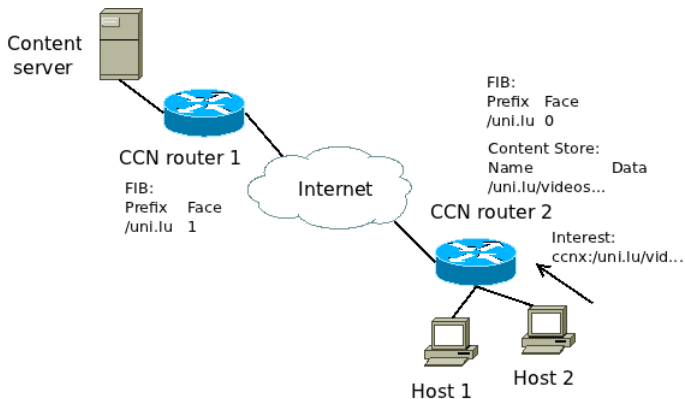


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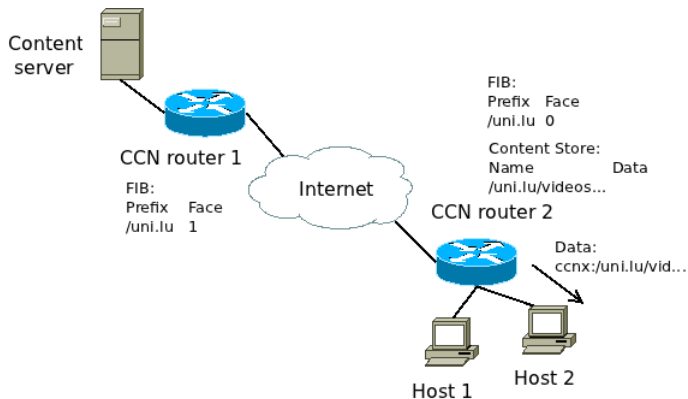
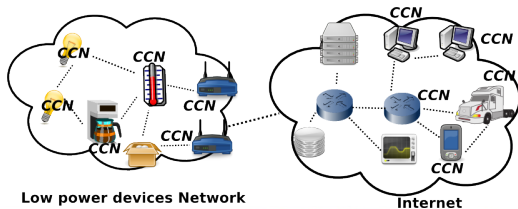


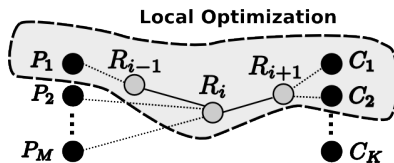
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- ▶ CCN faces → transparent IP network transversal
- ▶ CCN at all sides: sensors, servers, user devices
- ▶ context of this work = **sensor communication optimization** (low power devices)
- ▶ → retrieved sensed information
 - ▶ **push** based communication
 - ▶ **regular updates** (every 10 sec, 5 minutes)
 - ▶ **multiple consumers** (with different update frequencies)



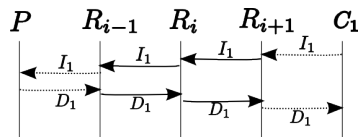
- ▶ **Limited resources** (CPU, memory, energy...)
 - ▶ no broadcast / no multicast to every consumer at every time unit
 - ▶ example: temperature update every 10 sec, service need update every 5 minutes \rightarrow overhead = 29 messages
- ▶ **Problem definition**
 - ▶ M producers, K destinations, CCN routers
 - ▶ Disjoint sensed information \rightarrow local problem
 - ▶ **1 producer, N consumers** $C = c_1, \dots, c_N$ with sampling periods i_j
 - ▶ R_i : Cap_i (available resources)
 - ▶ **objective: $min(msg(R_i))$**



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Push-based communication

- ▶ CCN data forwarded when an interest is emitted = pull model

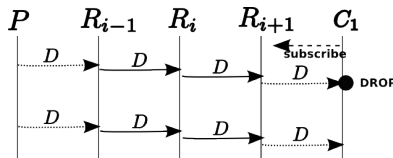


- ▶ SotA propositions
 - ▶ Jacobson *et. al*, ACM ReArch 2009: creation of a **rendez-vous point** → transaction model, individual demand from the producer (no subscription mechanism)
 - ▶ Carzaniga *et. al*, ACM SIGCOMM ICN 2011: dissemination of ephemeral messages → **a new unicast table** (no multicast) based on **IP routing** for avoiding routing loops

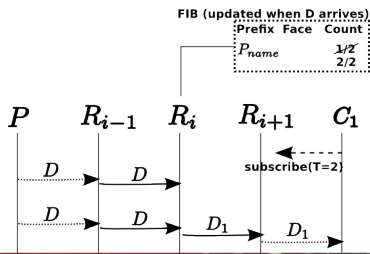
Push-based communication

► Approach overview

- **subscription = content publications**
- sensor value disseminations
 - → **use the standard FIB to disseminate** sensor value
 - interest message towards the consumer content name prefix + value as its suffix
 - */roomA/thermometer/value/20*
 - → but does not work as there is no content in the other way
 - use FIB for forwarding not requested content → **no PIT update**
 - → routing loop (interests are not consumed anymore → ***last_seen* flag in FIB**)



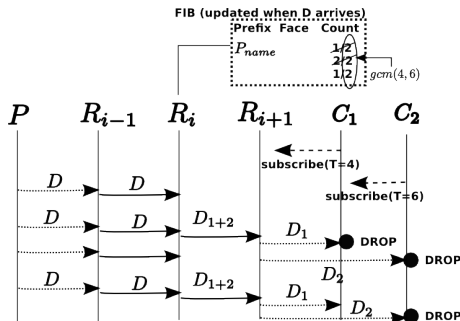
- ▶ Avoid sending useless dropped messages
 - ▶ simple filtering
 - ▶ the router only forwards data to C_i at the period t only if the consumer C_i has expressed its interest regarding this period
 - ▶ $\rightarrow t$ is a multiple of the sampling period i ;
 - ▶ \rightarrow each C_i has to specify its sampling period to R_i (smart router)
 - ▶ need a counter in the FIB, one counter per subscriber = optimal strategy



Smart forwarding – gcd

- ▶ Use a limited number of counters Cap_i
- ▶ **Extreme case $Cap_i = 1$**
 - ▶ apply partially simple filtering
 - ▶ **some packets still dropped** → minimize it

- ▶ remove **identical subscribed periods** i_i
- ▶ if **multiples in subscribed periods** and keep the smallest, e.g. $i_1 = 2$ and $i_2 = 4$ → forward packets every 2
- ▶ extend the previous concept using the **GCD** (Greatest Common Divisor)



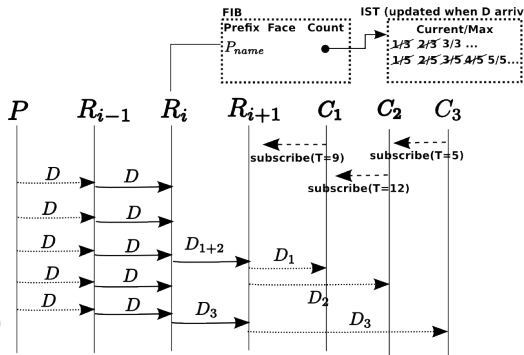
Smart forwarding

- ▶ More counters, $Cap_i > 1$
 - ▶ IST (Information Sampling Table) for tracking multiple counters
 - ▶ guarantee that each C_i receives the packets as subscribed or more
 - ▶ \rightarrow select $G = \{g_1, \dots, g_{Cap_i}\}$ such that:

$$\forall C_i \in C, \exists g_k \in G, \alpha \in \mathbb{N}, g_k \times \alpha = i$$

- ▶ incremental IST updates
- ▶ heuristic: maximal GCD between two i_j

- 1: T : new registered interval
- 2: $G.append(T)$
- 3: **if** $|G| > Cap_i$ **then**
- 4: $g_select \leftarrow g_i \in G, \forall g_k \in G, gcd(g_i, T) \geq gcd(g_k, T)$
- 5: $G.replace(g_select, gcd(g_select, T))$
- 6: **end if**

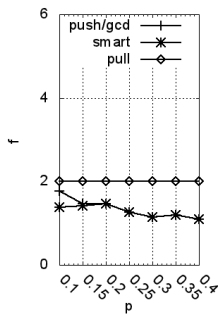
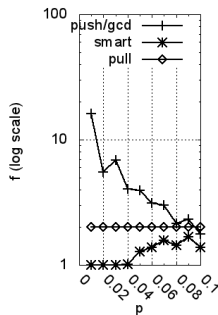


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- ▶ **Basic simulator**
 - ▶ limited functions necessary to assess our approach
 - ▶ no PIT or Content Store, no physical layer...
- ▶ **Parameters**
 - ▶ MAX_T : maximal sampling interval that a consumer can subscribe to,
 - ▶ p : number of consumers with a unique sampling interval expressed as a proportion of MAX_T
 - ▶ $\#counters$: number of available counters,
 - ▶ $\#steps = 5 \times MAX_T$ the number of simulation step, 1 step = the initial sampling period of the sensor
- ▶ **Strategies**
 - ▶ $s \in \{pull, push, optimal, gcd, smart\}$
 - ▶ **message overhead**: $f_s = msg_s / msg_{optimal}$
- ▶ i_i chosen randomly
- ▶ simulation executed 10 times

Number of consumers

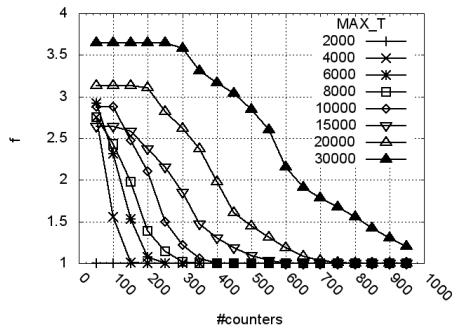
- ▶ $MAX_T = 3000$, $\#counters = 100$
- ▶ *pull* strategy \rightarrow 2 messages (1 interest + 1 data)
- ▶ prime numbers $\rightarrow gcd = 1 \rightarrow gcd \sim push$ (forwarding every packet)
- ▶ when p increases
 - ▶ more consumers \rightarrow higher chance to have low i_i including 1
 - ▶ \rightarrow all strategies tends to 1
- ▶ smart forwarding benefit
 - ▶ $1 < f_{smart} < 1.67$
 - ▶ $p < 0.04$ a unique counter per i_i (equivalent to *optimal*)
 - ▶ worst case: 240 unique subscribed periods ($p = 0.08$)



Number of consumers

- ▶ $p = 0.05 \times MAX_T$
- ▶ smart forwarding strategy
- ▶ *#counters* and *MAX_T* vary

- ▶ when *#counters* increases \rightarrow overhead reduces faster with less subscribers (lower *MAX_T*)
- ▶ **fine tuning of *Cap_i*** during the design phase of the sensor
 - ▶ assumption: a maximal overhead, a number of subscribers
 - ▶ \rightarrow **number of counters**



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- ▶ **CCN for distributed IoT communication**
- ▶ Context: producers (sensors) regularly updating consumers
- ▶ A new **push based mechanism to limit the message overhead** based on subscription requests
- ▶ **Smart forwarding is close to the optimal forwarding and always better than the standard pull based communication of CCN**
- ▶ Future work: multiple-content producers

Acknowledgement: This work was partly funded by IoT6 FP7 EU project under the grant agreement 288445.

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