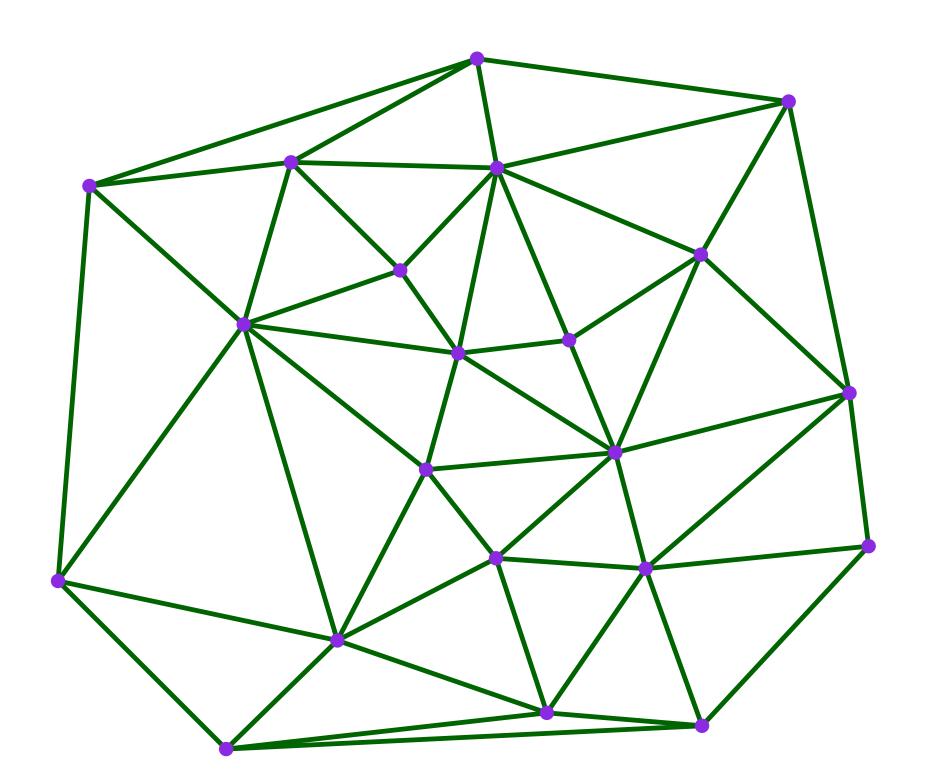
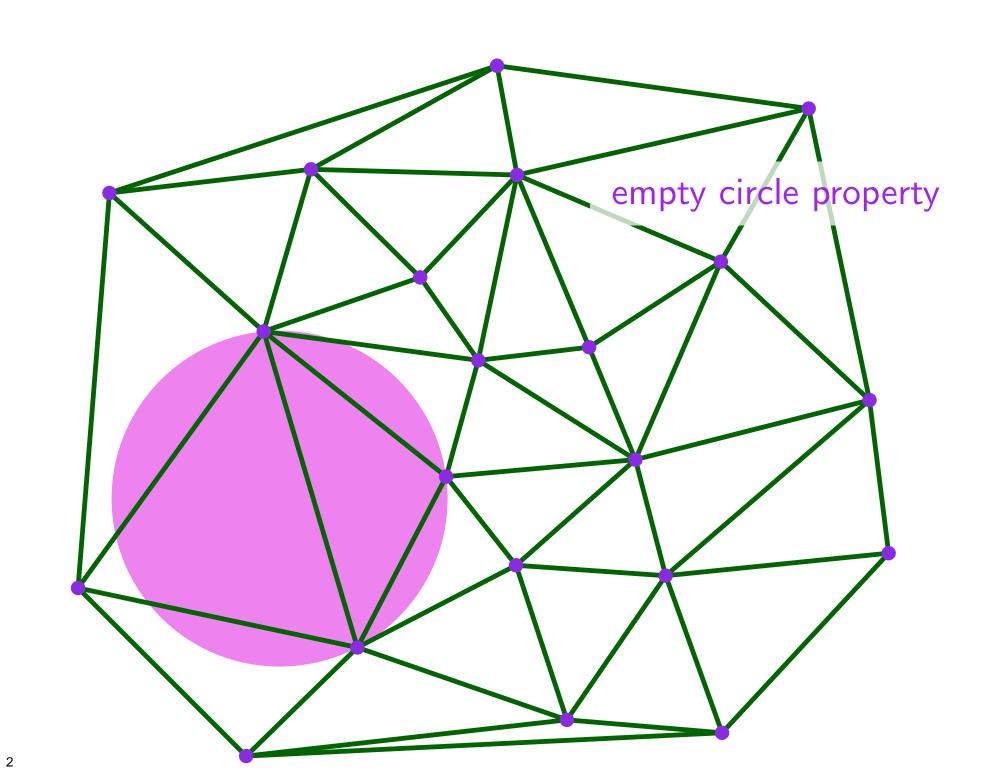
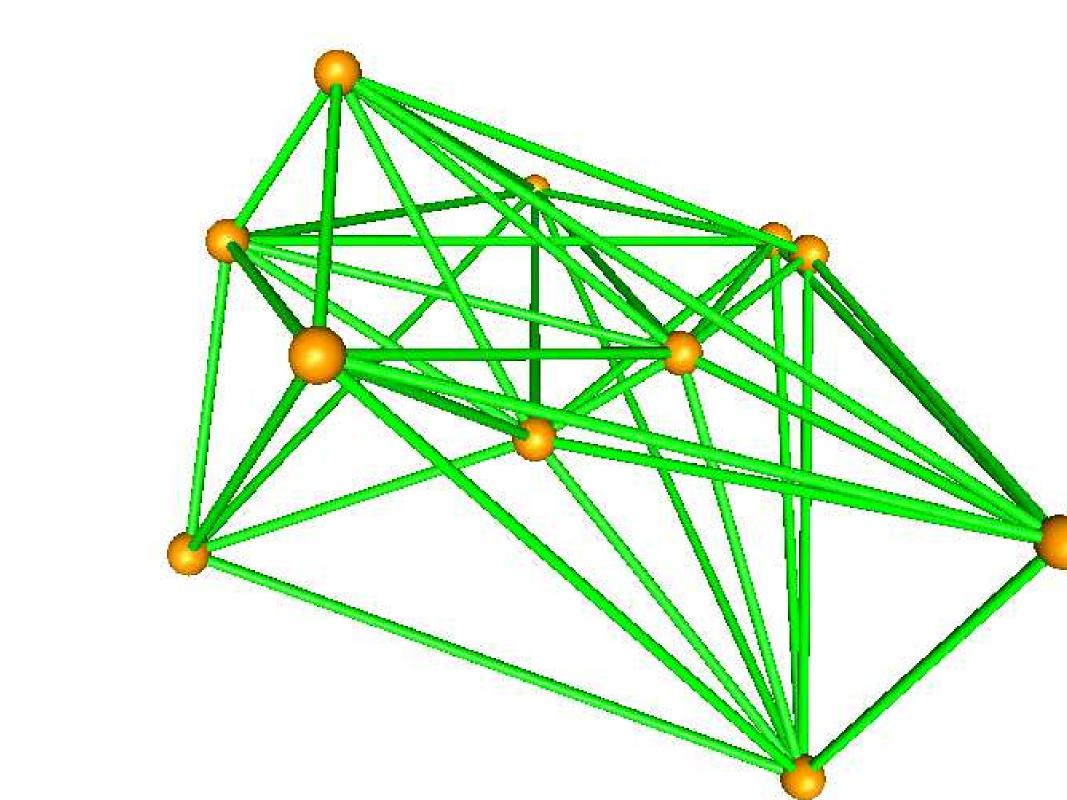
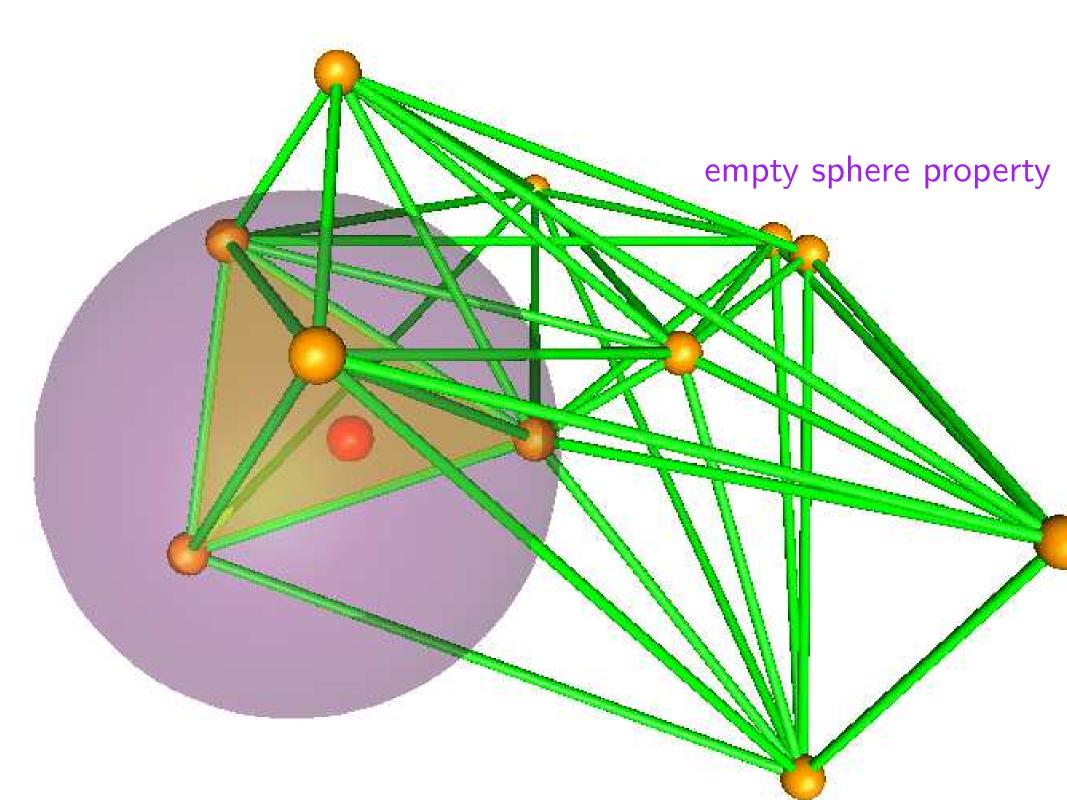


Everybody in this room know what Delaunay is!









A (partial) history of Delaunay algorithms (and of computational geometry)

A (partial) history of Delaunay algorithms (and of computational geometry)

Earlier algorithms [1970...]

(and of computational geometry)

Earlier algorithms [1970...]

Incremental 2D/3D

Gift wrapping 2D

(and of computational geometry)

Earlier algorithms

[1970...]

Incremental

2D/3D

Gift wrapping

2D

Non optimal

"simple"

actually coded

linked to applications: meshes, reconstruction

(and of computational geometry)

Earlier algorithms [1970...]

Worst case algorithms [1980...]

(and of computational geometry)

Earlier algorithms [1970...]

Worst case algorithms [1980...]

Divide & conquer 2D

Plane sweep 2D

(and of computational geometry)

Earlier algorithms [1970...]

Worst case algorithms [1980...]

Divide & conquer 2D

Plane sweep 2D optimal

'simple

actually coded

linked to applications: meshes, reconstruction

(and of computational geometry)

Earlier algorithms [1970...]

Worst case algorithms [1980...]

Randomized algorithms [1990...]

(and of computational geometry)

Earlier algorithms [1970...]

Worst case algorithms [1980...]

Randomized algorithms [1990...]

Delaunay tree

Clarkson & Shor

Dynamic updates (history graph)

Delaunay hierarchy

Spatial sorting (BRIO)

(and of computational geometry)

Earlier algorithms [1970...]

Worst case algorithms [1980...]

Randomized algorithms [1990...]

Robustness issues [1995...]

(and of computational geometry)

Earlier algorithms [1970...]

Worst case algorithms [1980...]

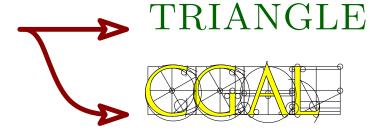
Randomized algorithms [1990...]

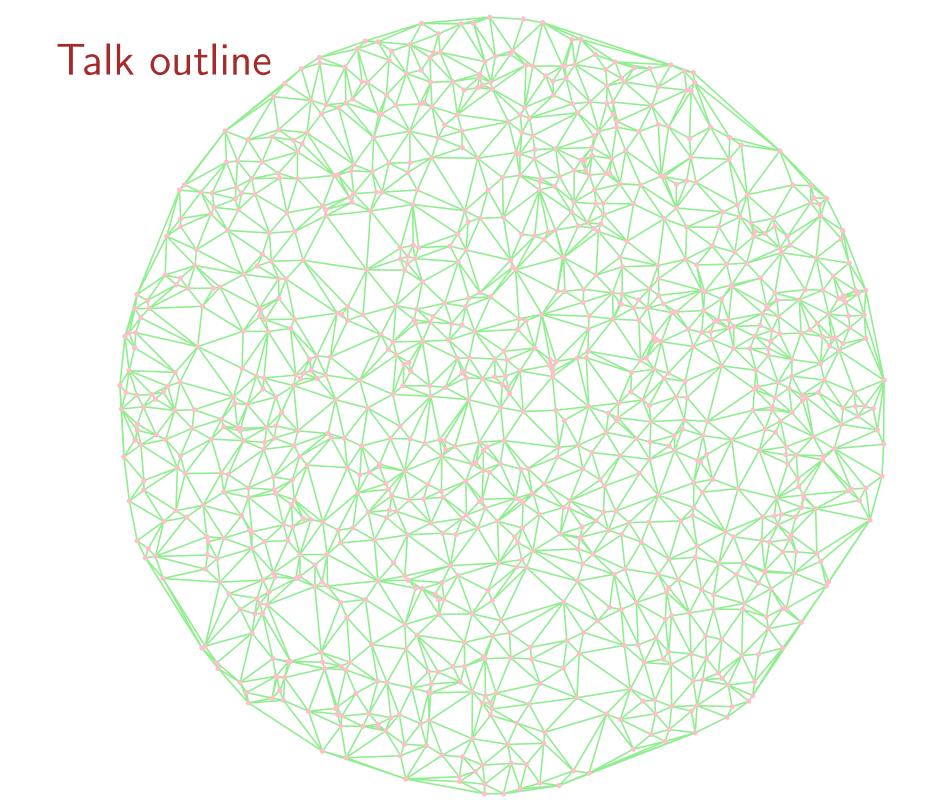
Robustness issues [1995...]

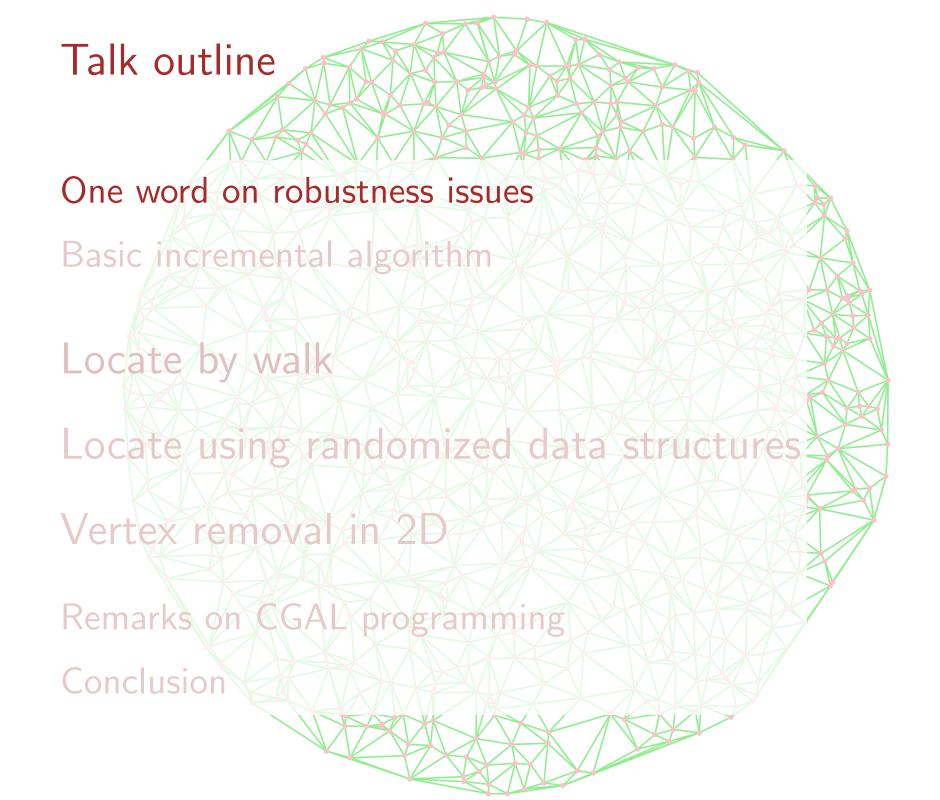
Properties checking

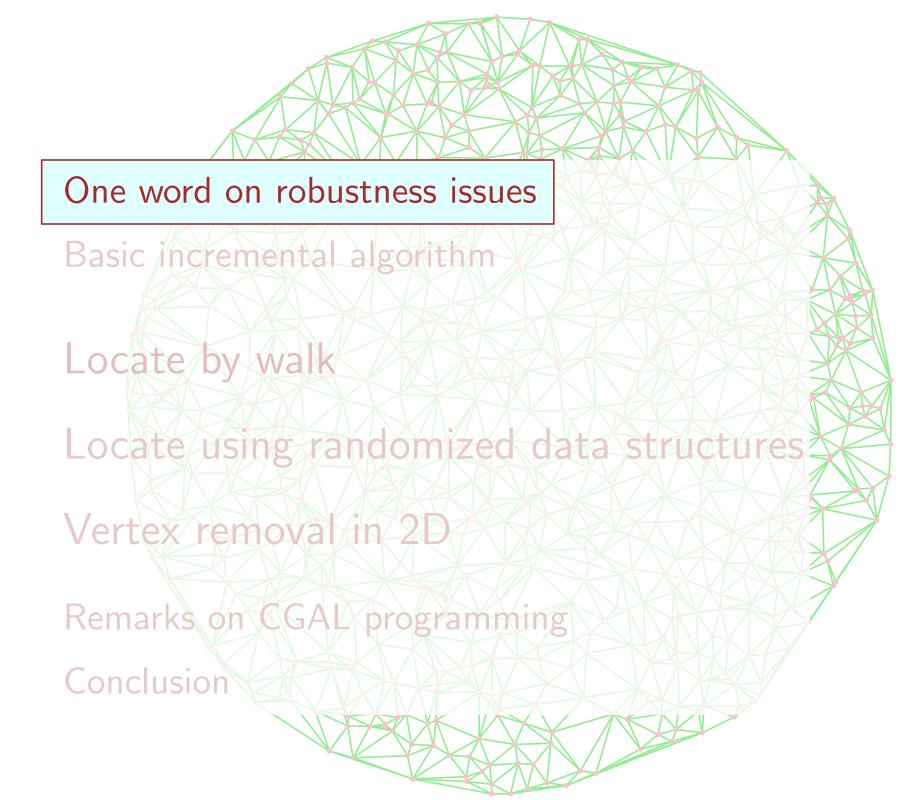
→ VRONI

Exact computation paradigm











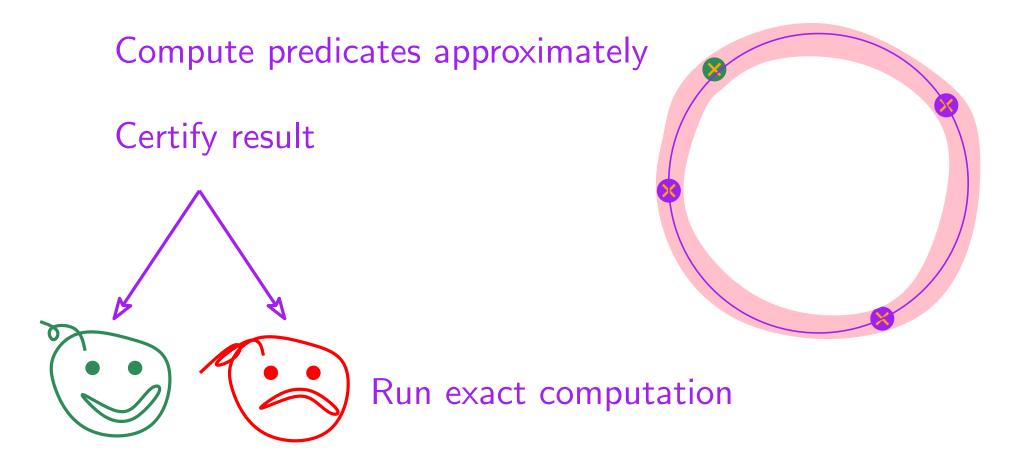


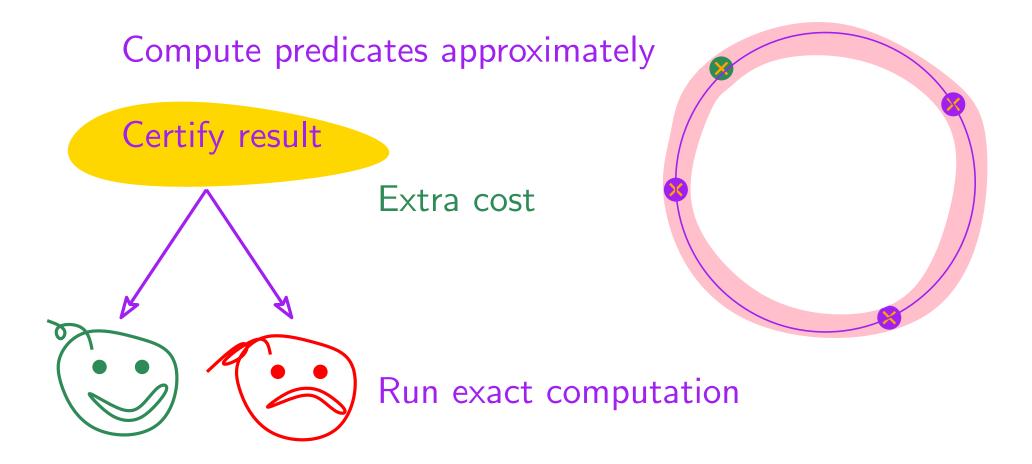
Compute predicates approximately

Compute predicates approximately

Certify result

Compute predicates approximately _____ Certify result







Deladitay 2D 10111 points

 ${\tt Exact_predicates_inexact_constructions_kernel}$

10.6 seconds

Cartesian < double >

9.7 seconds



Exact_predicates_inexact_constructions_kernel

Cartesian < double >

10.6 seconds 82 seconds

9.7 seconds 75 seconds



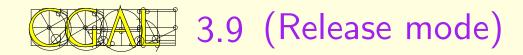
Exact_predicates_inexact_constructions_kernel

Cartesian < double >

10.6 seconds 82 seconds may loop (or crash)

All benchmarks

2.3 GHz, 16 GByte workstation



Exact_predicates_inexact_constructions_kernel

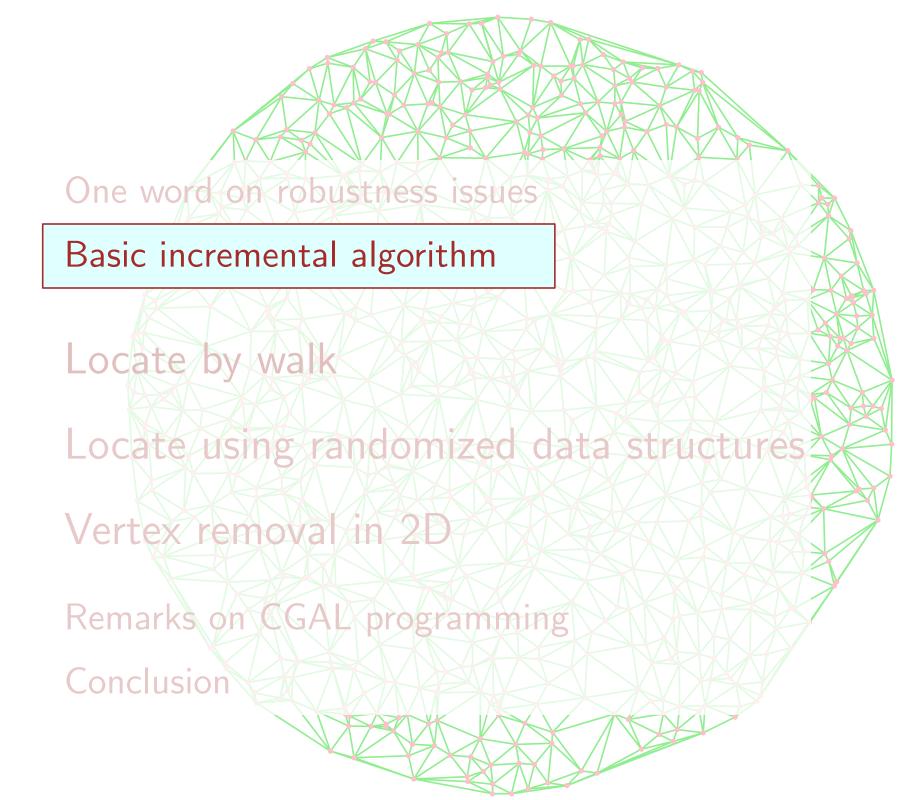
Exact_predicates_inexact_constructions_kernel

10.6 seconds

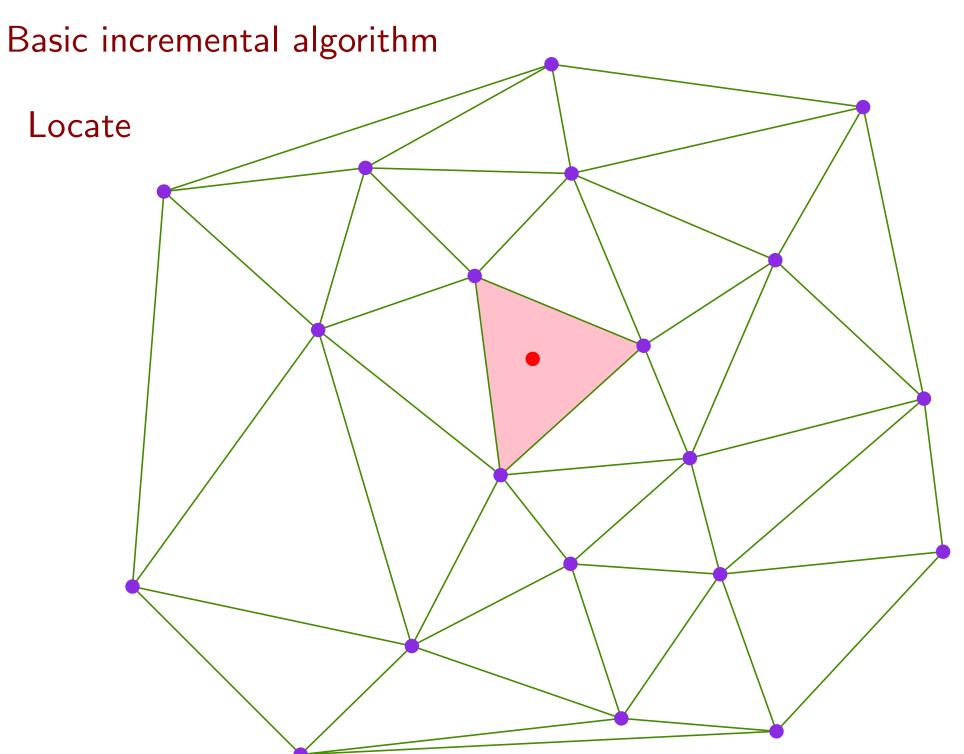
82 seconds

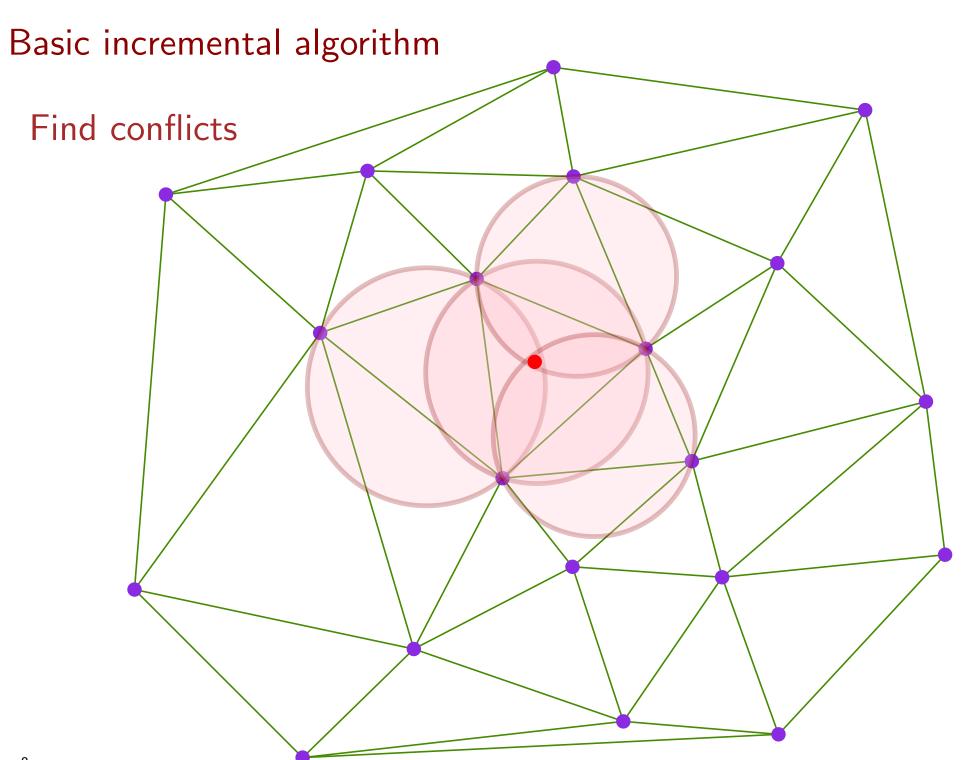
Cartesian < double >

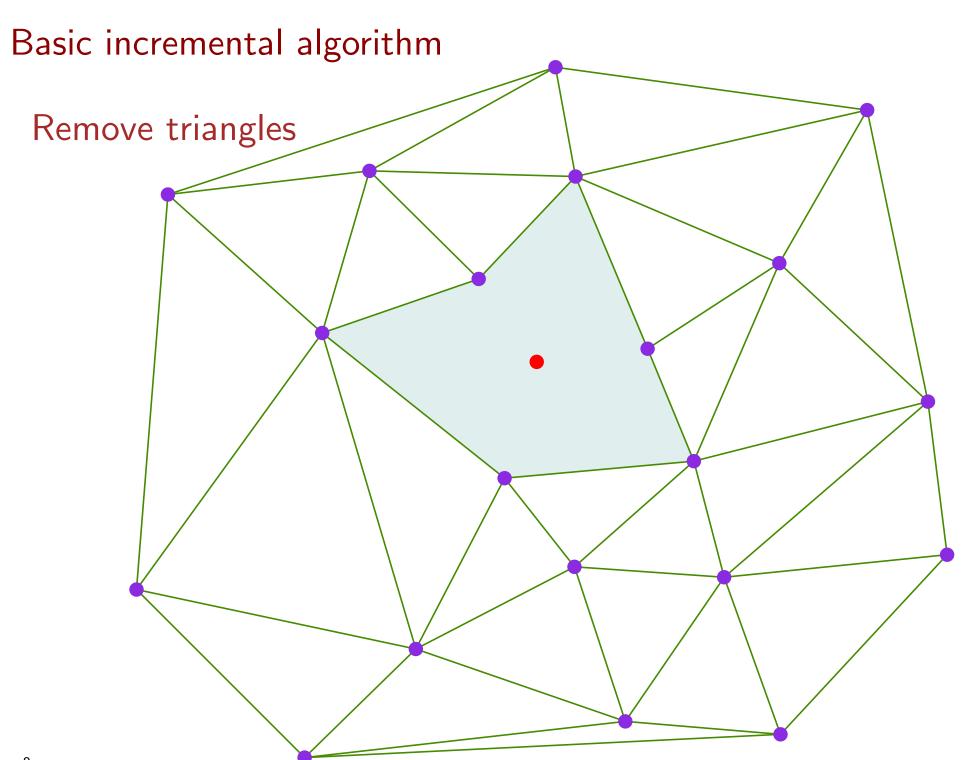
9.7 seconds 75 seconds

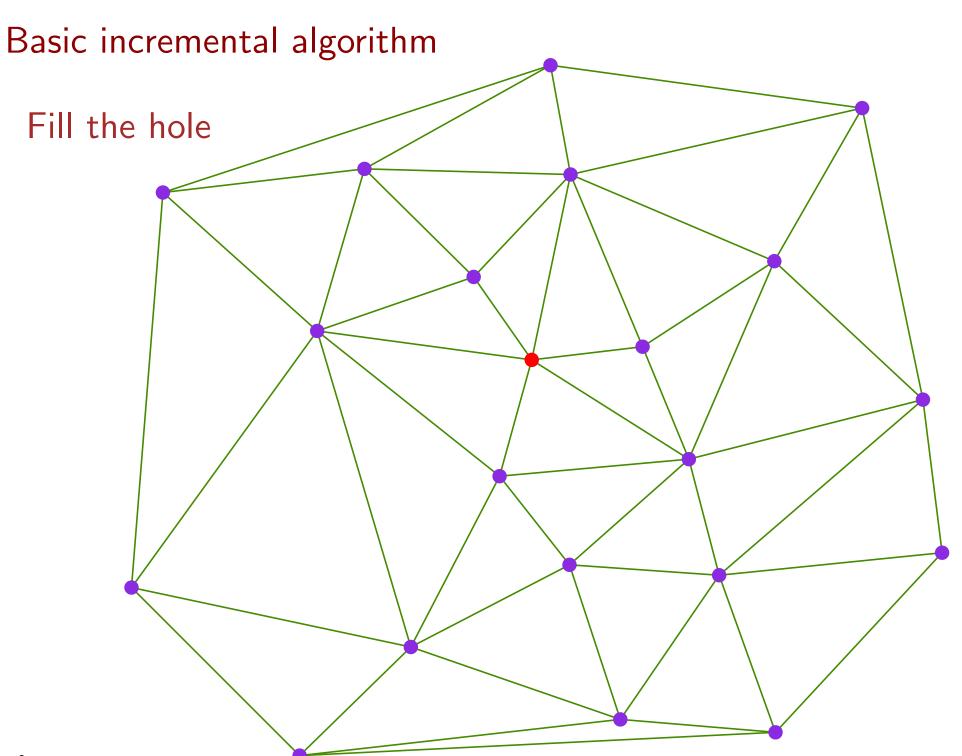


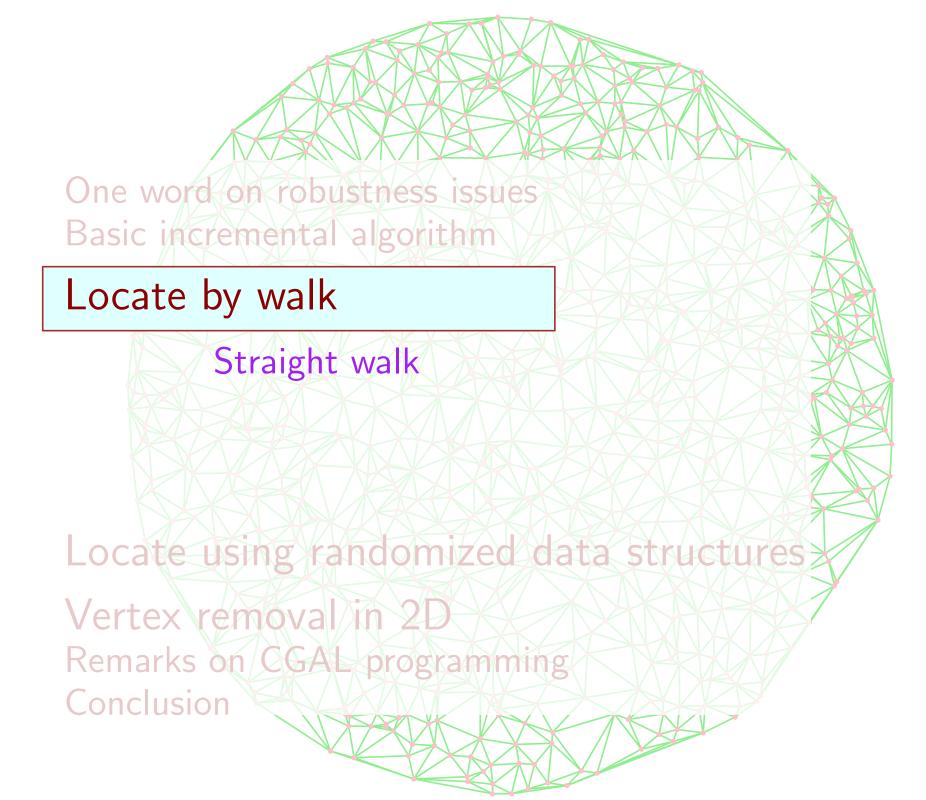
Basic incremental algorithm

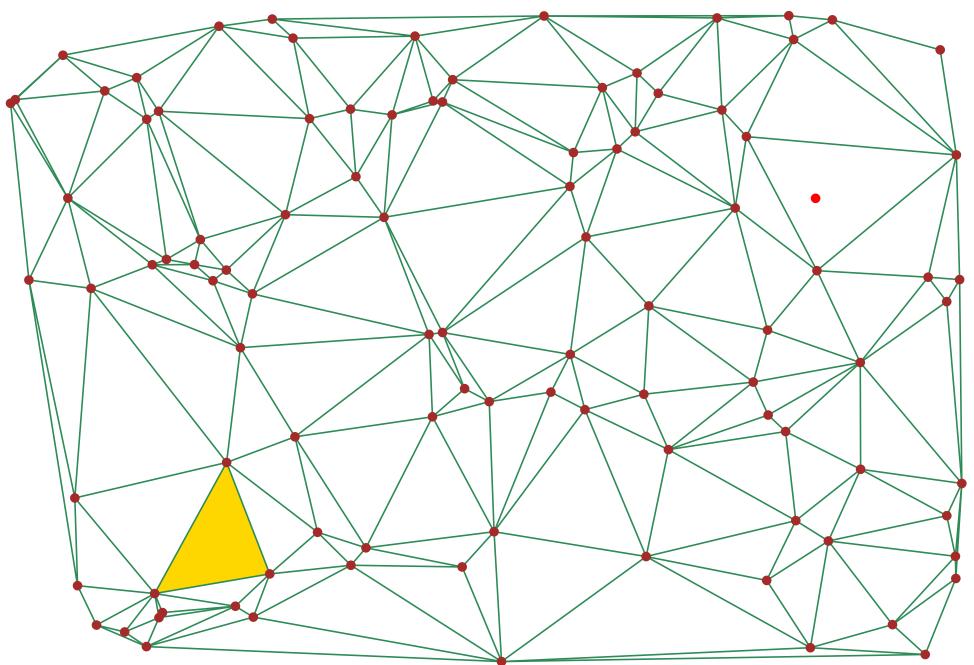


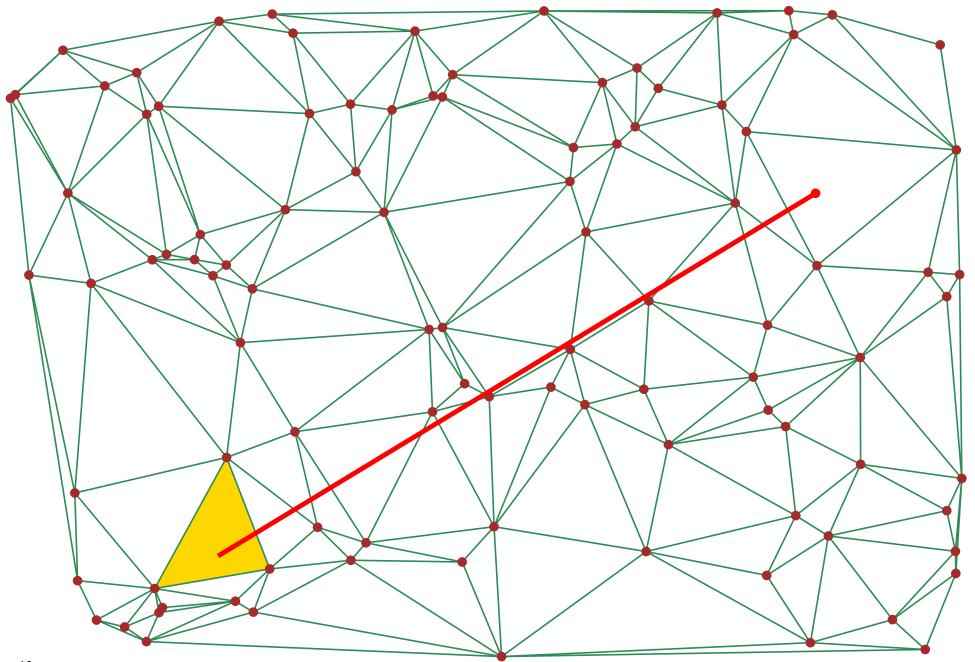


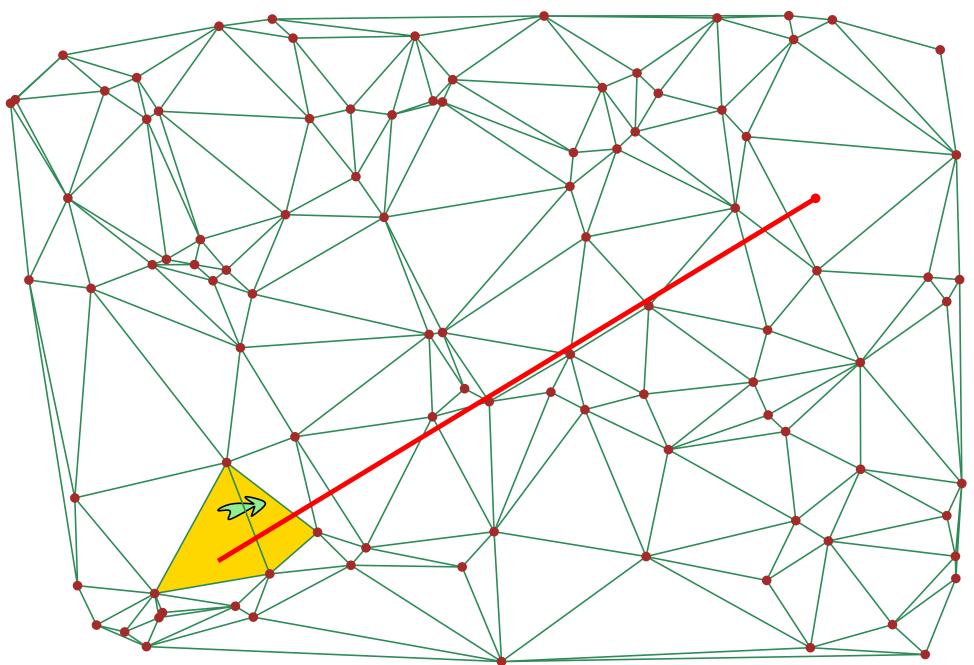


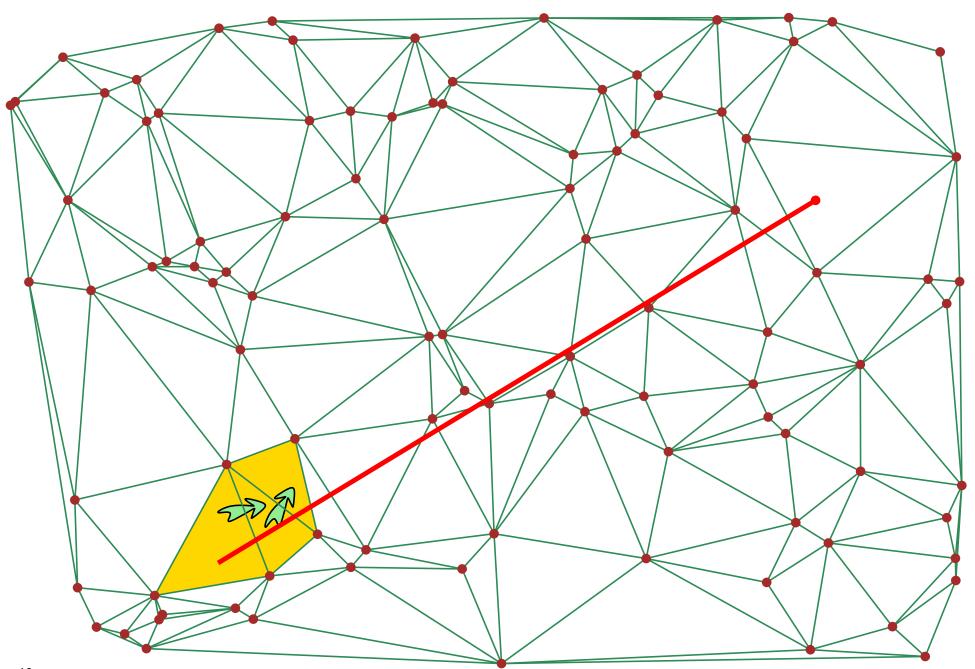


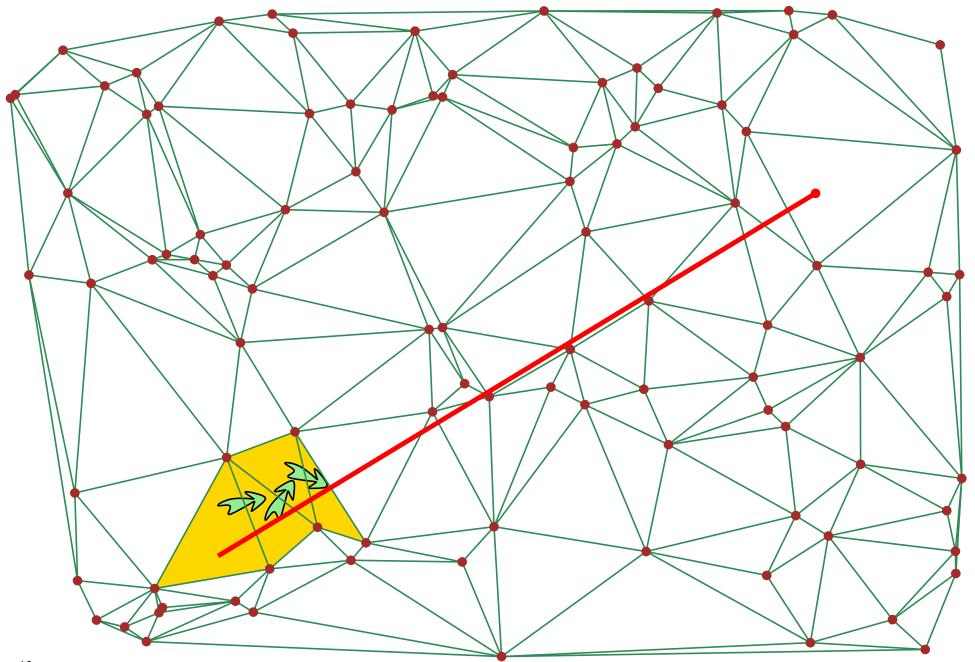


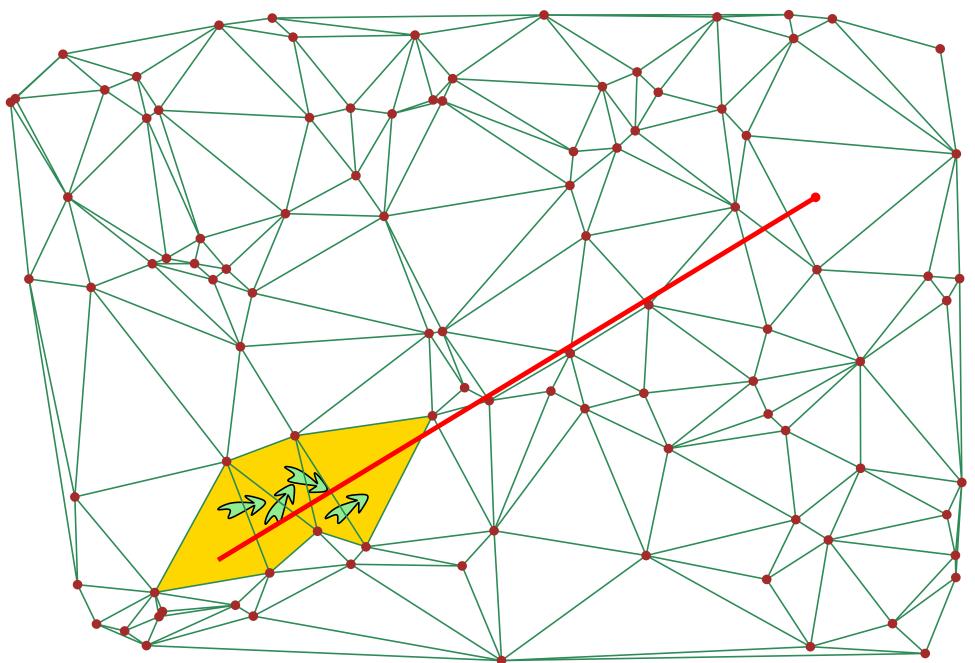


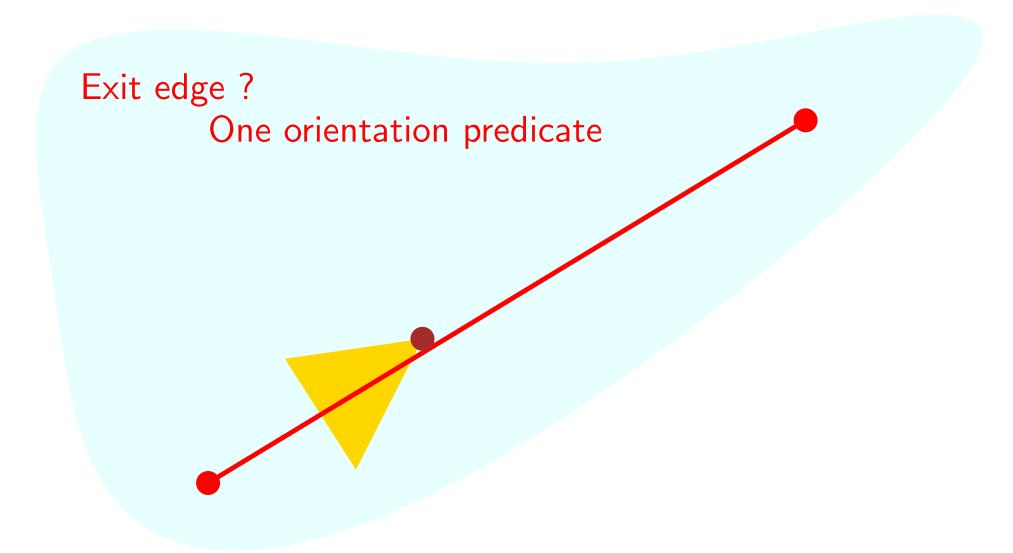


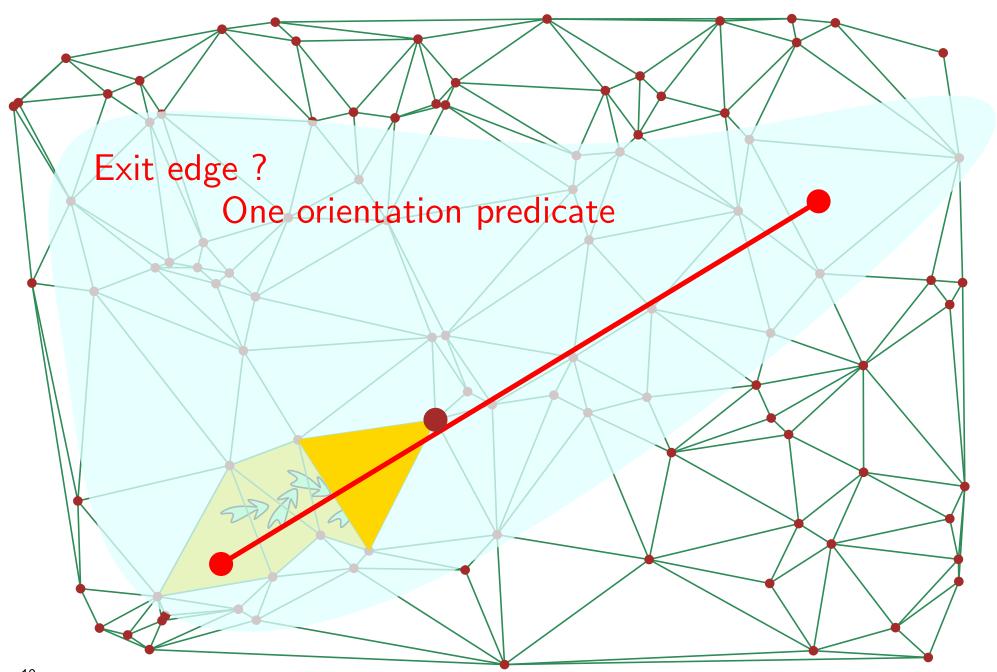


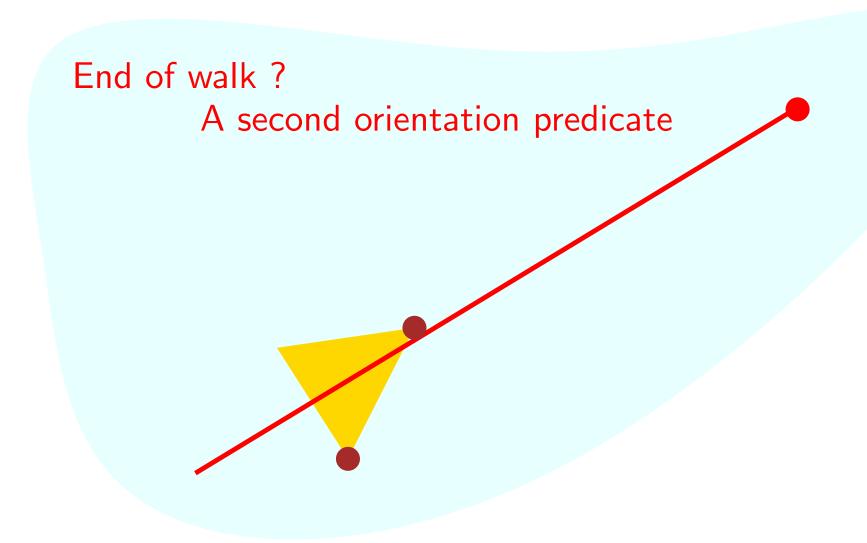


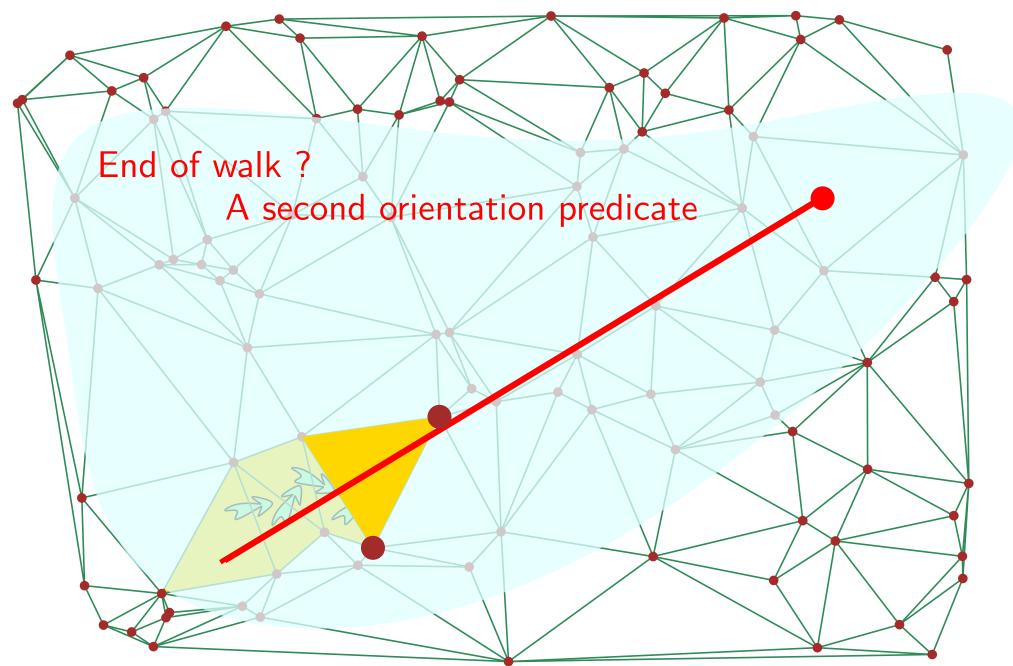


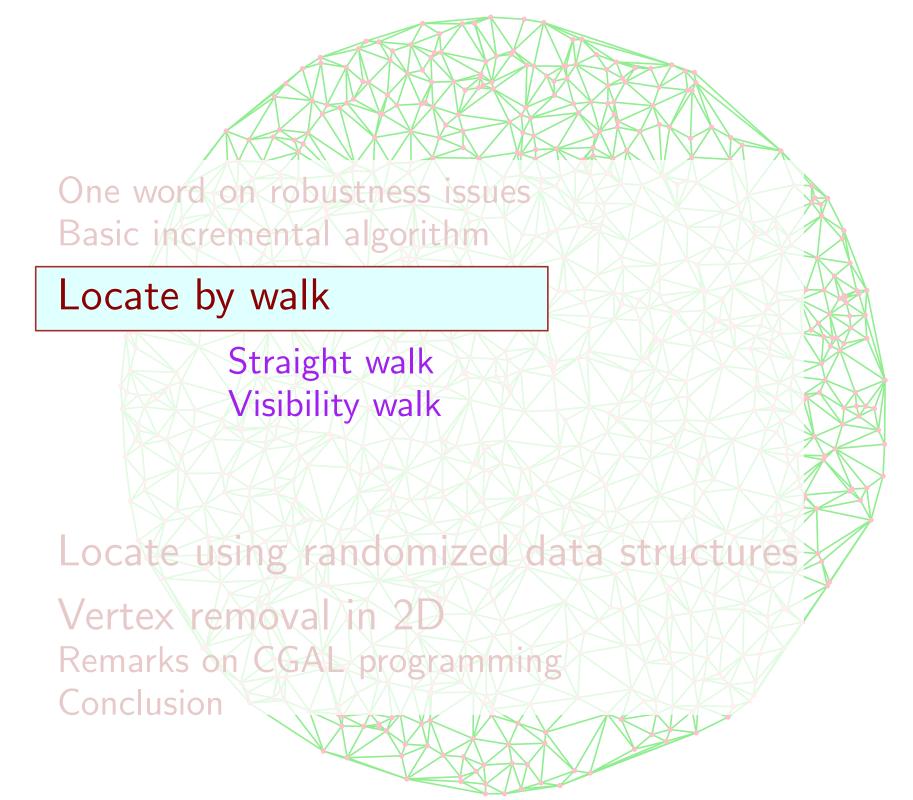


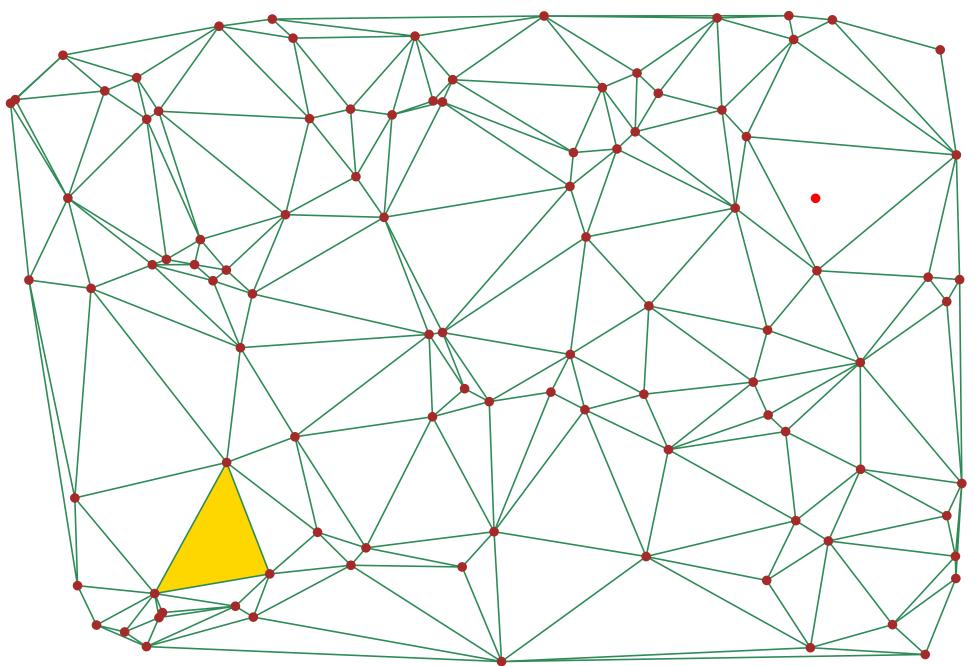


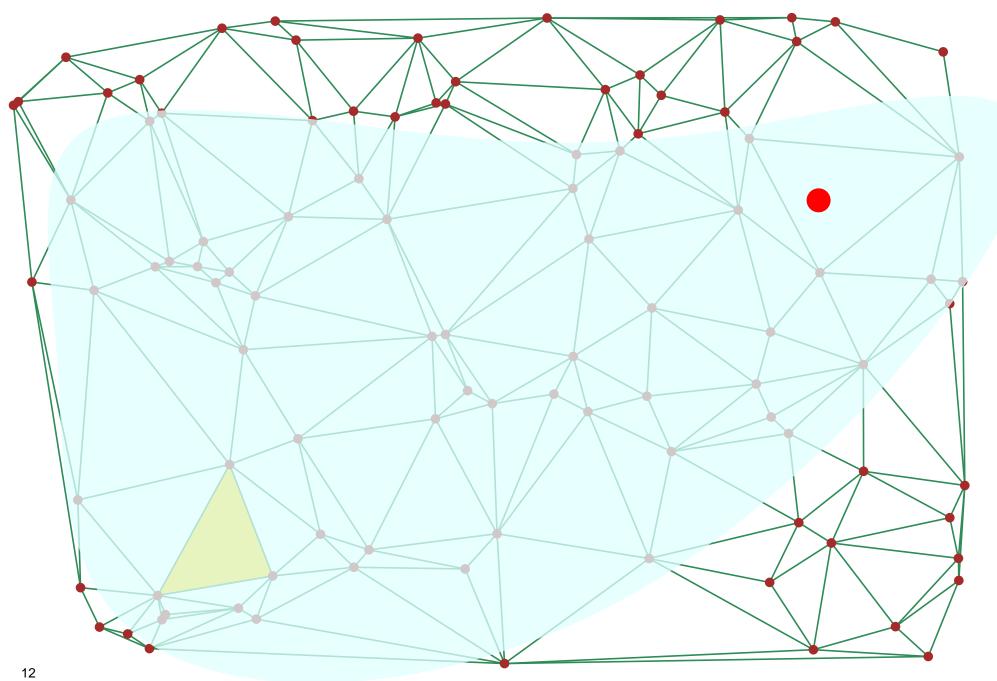


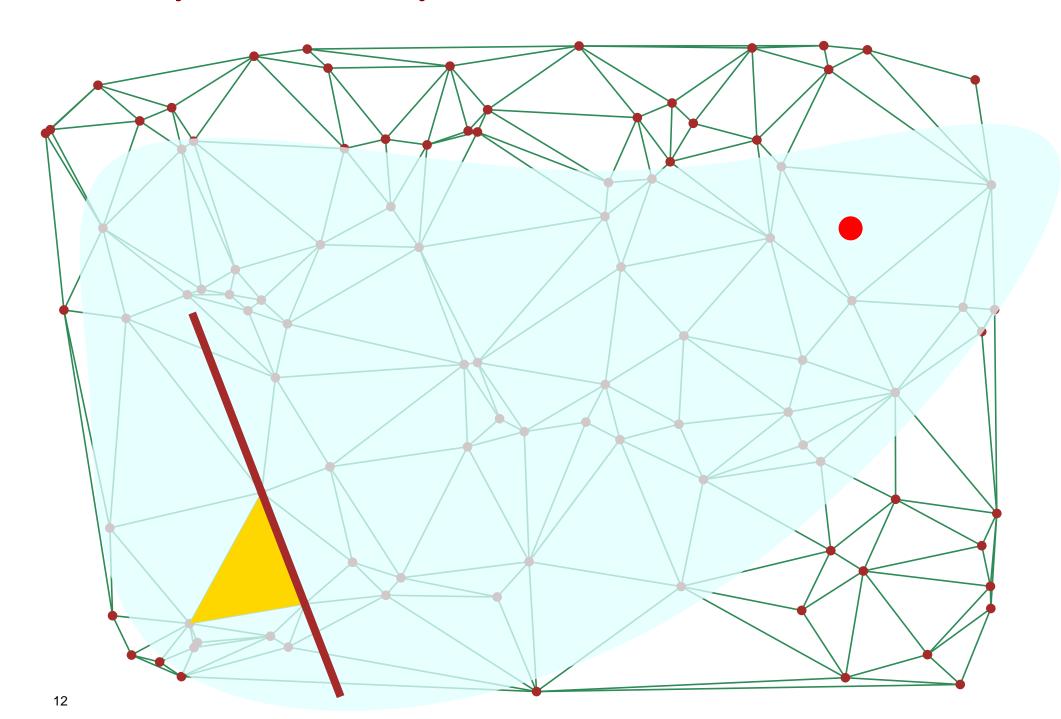


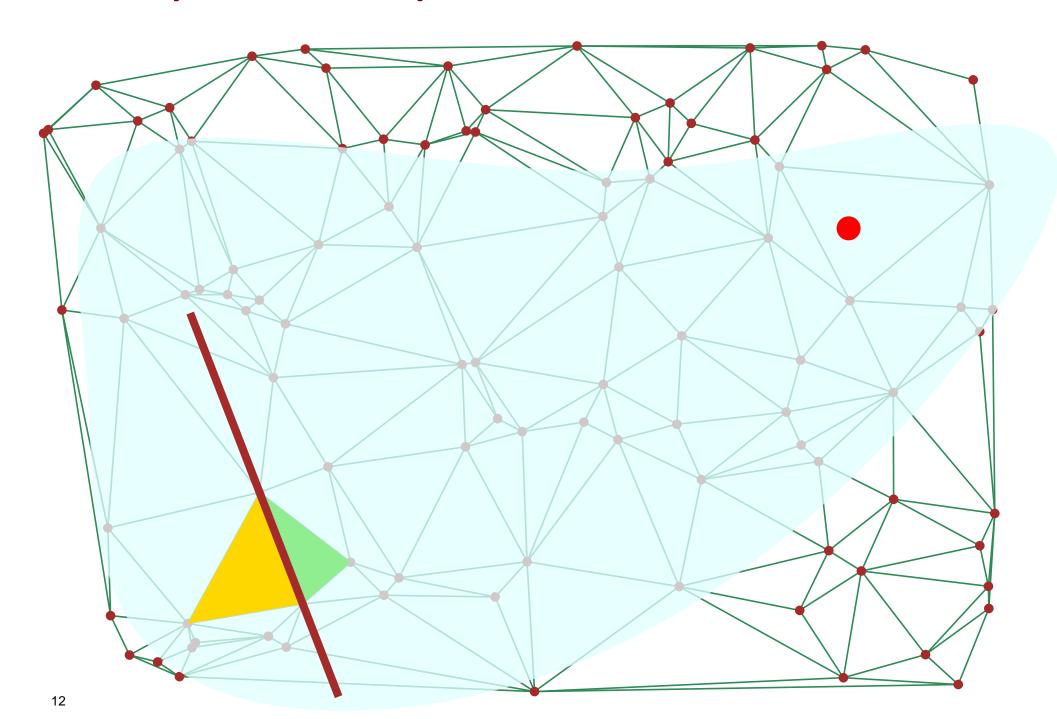


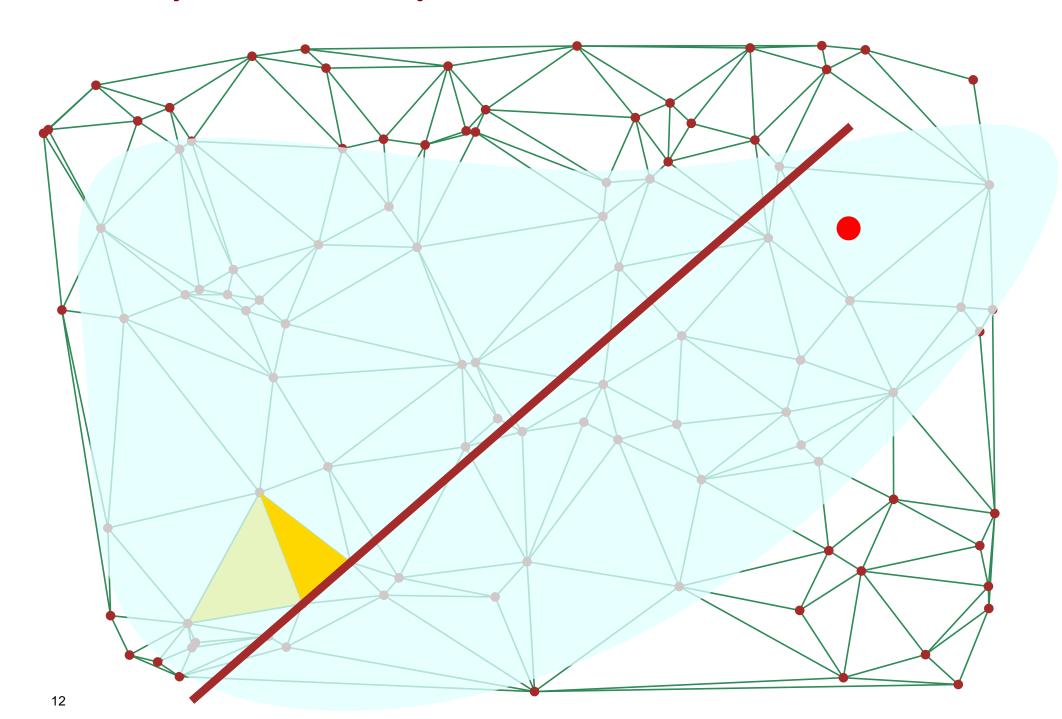


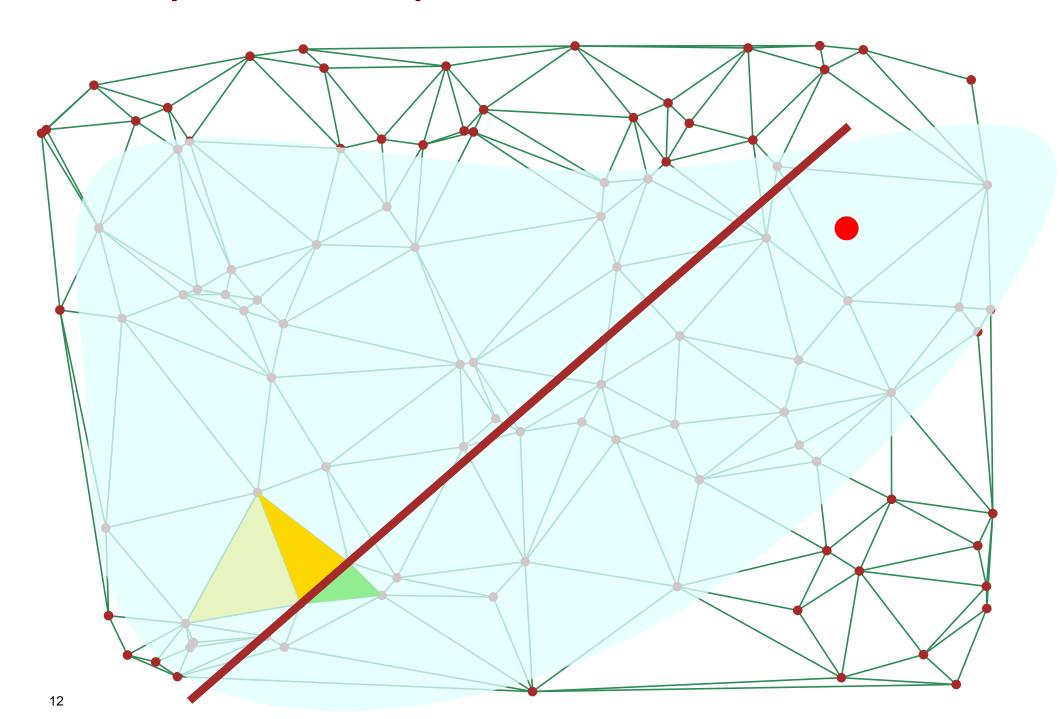


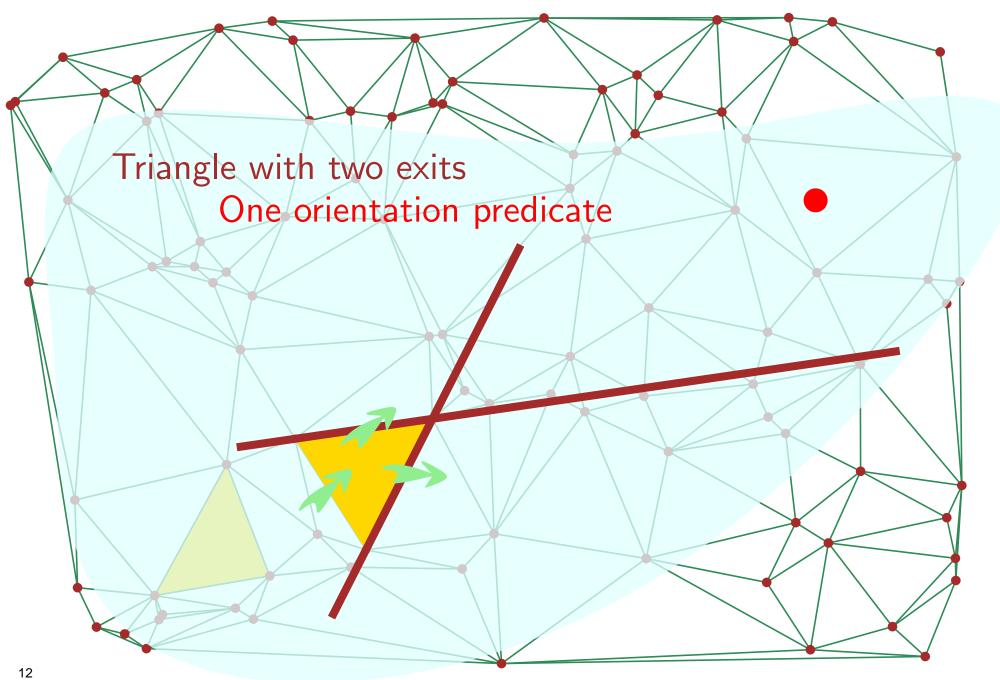


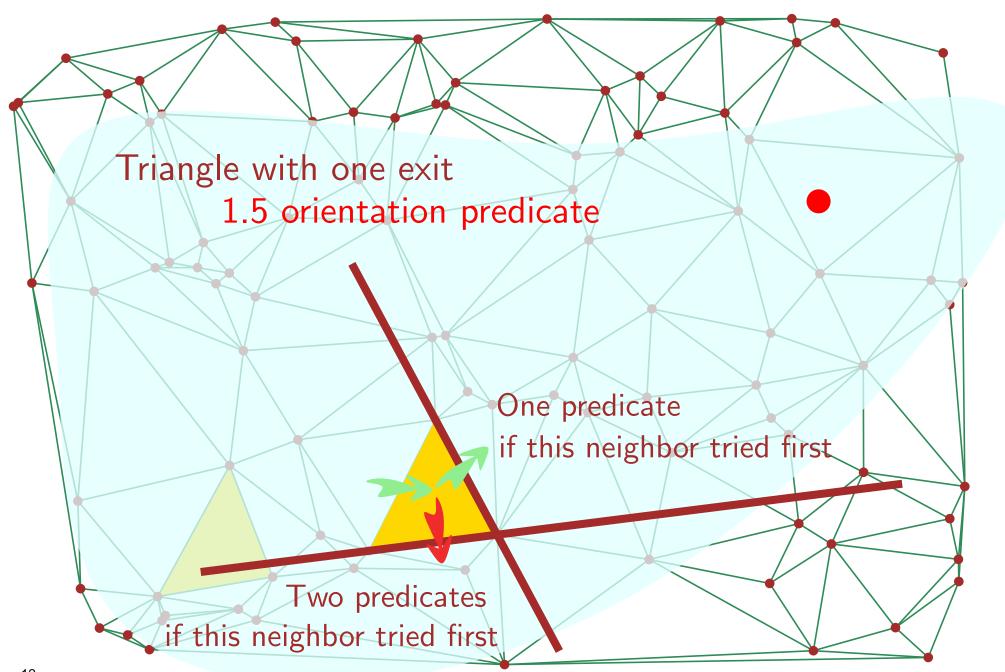


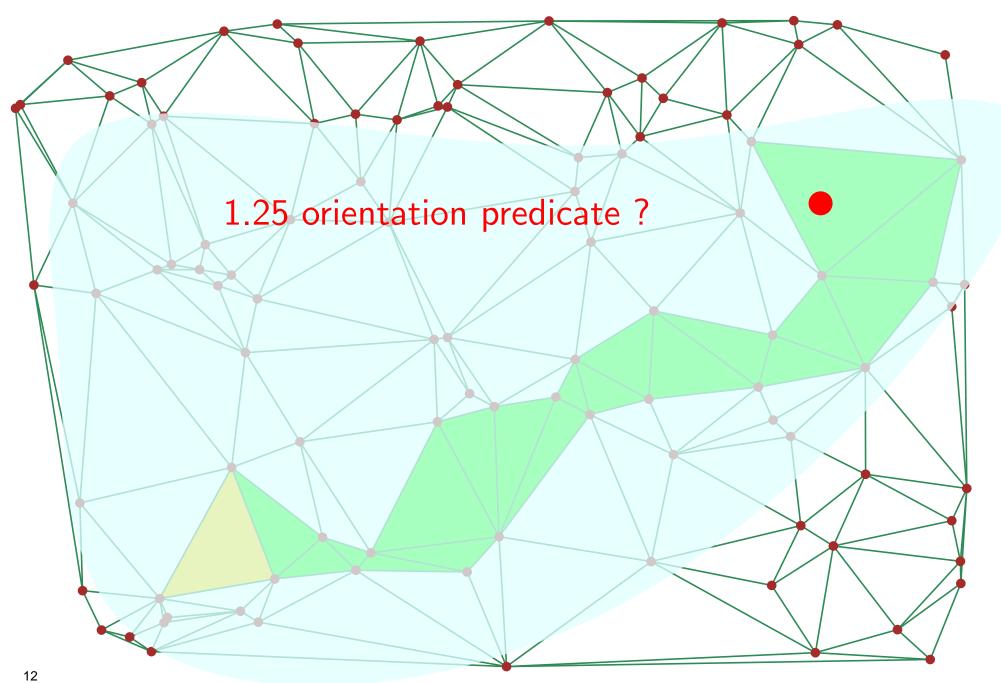


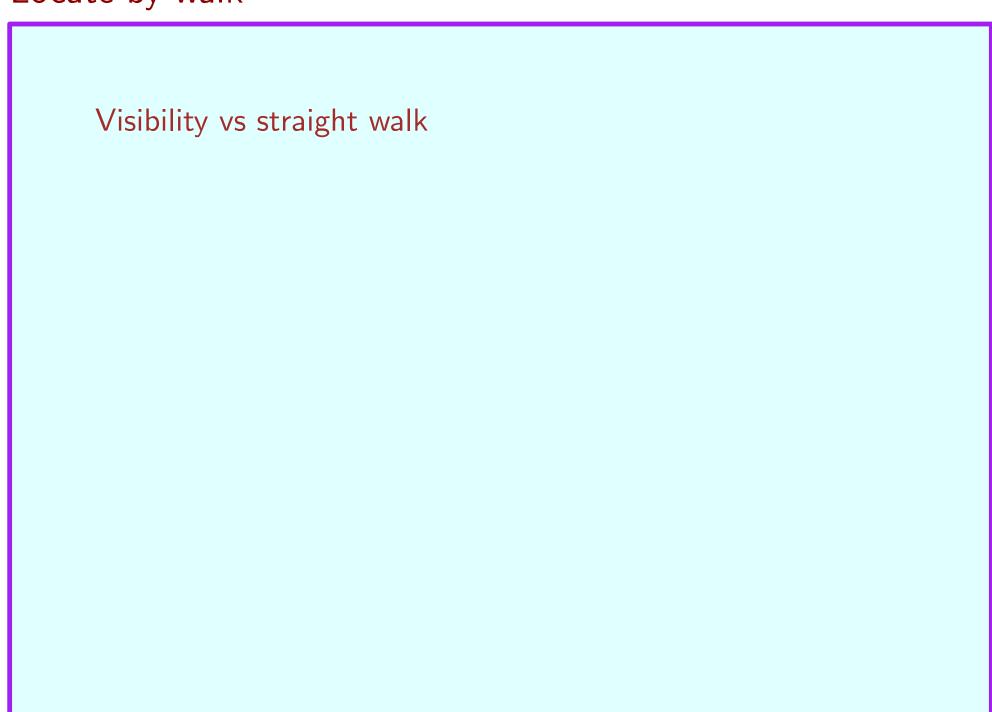












Visibility vs straight walk 2D and 3D

less predicates per crossed edge

similar number of crossed edges

experimental / theoretical

Visibility vs straight walk

Speed improvement?

Visibility vs straight walk

Speed improvement

Walk in Delaunay 1 Mpoints

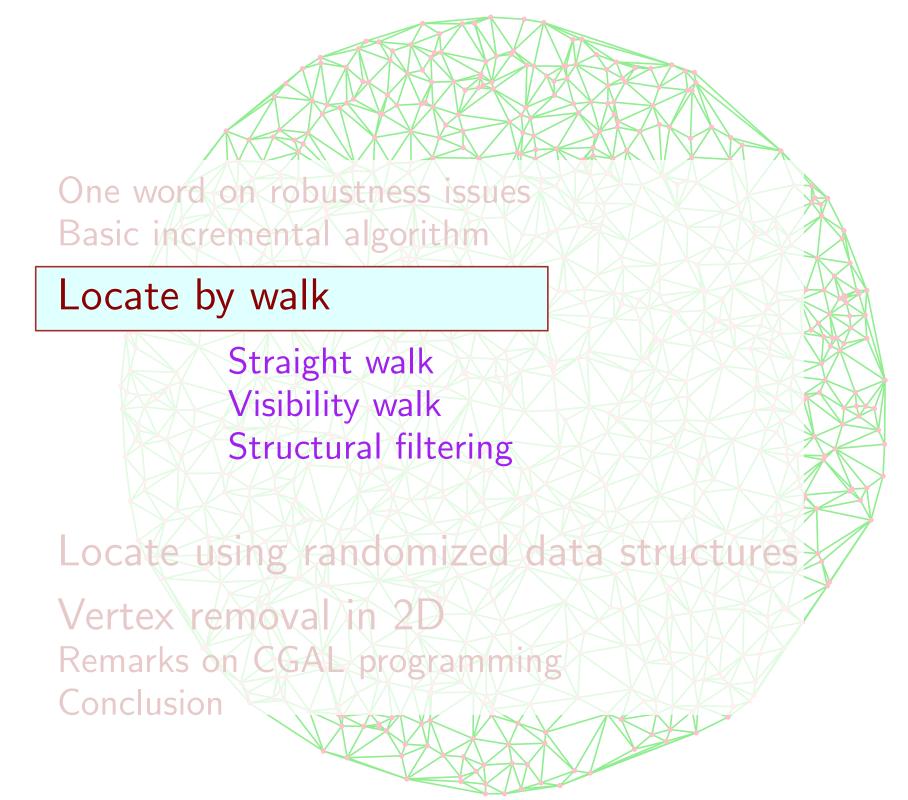
Straight: 324 μ s

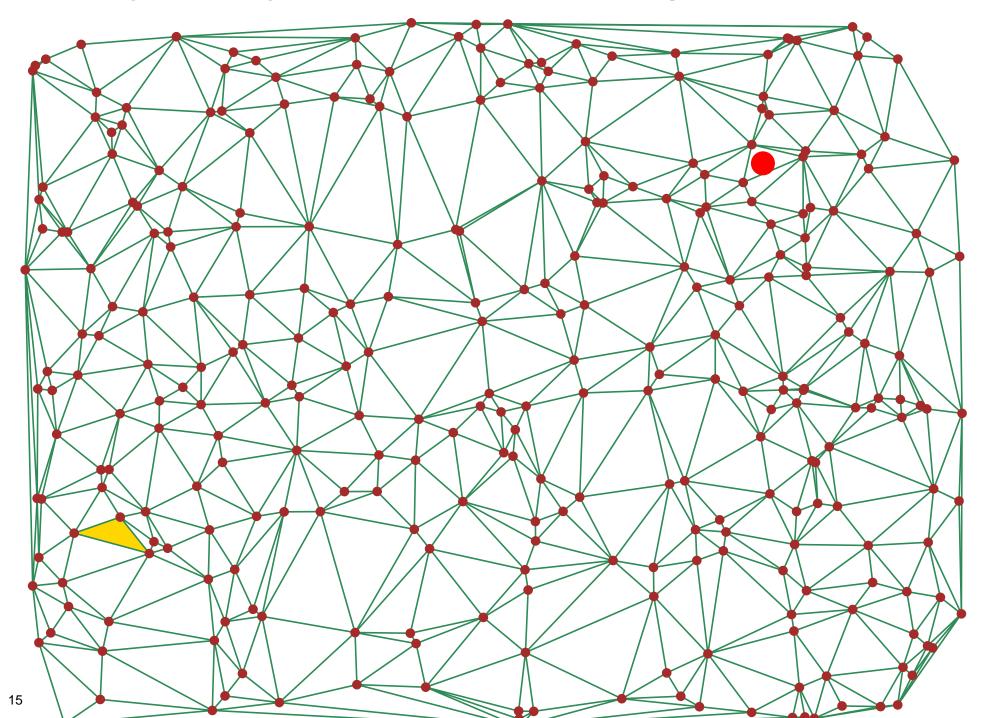
Visibility: 285 μ s

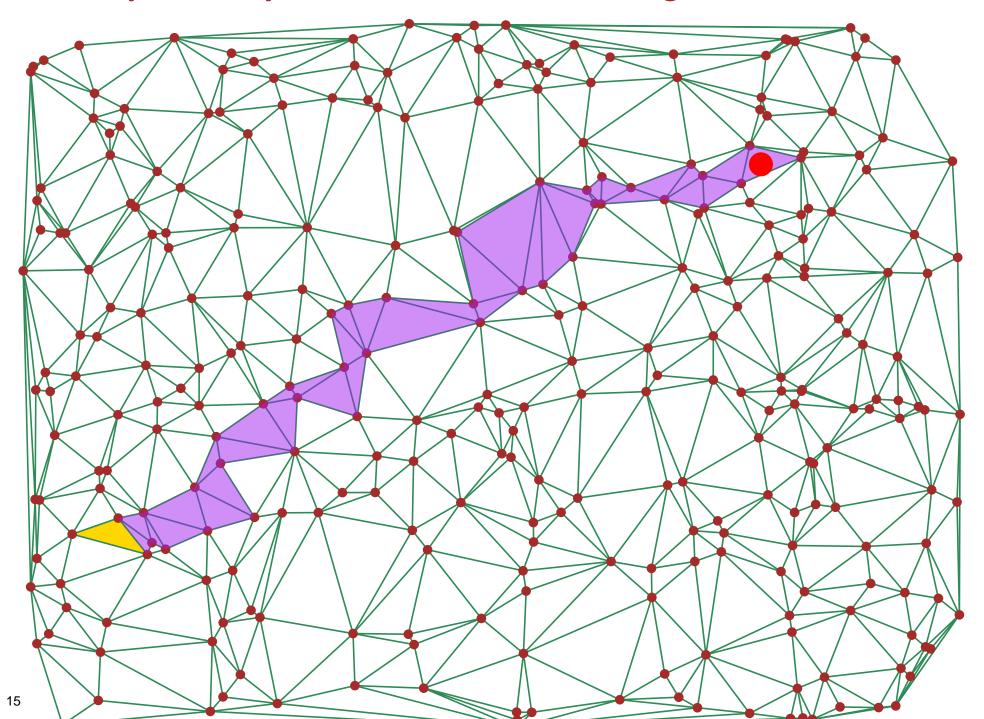
3D: 97 μ s

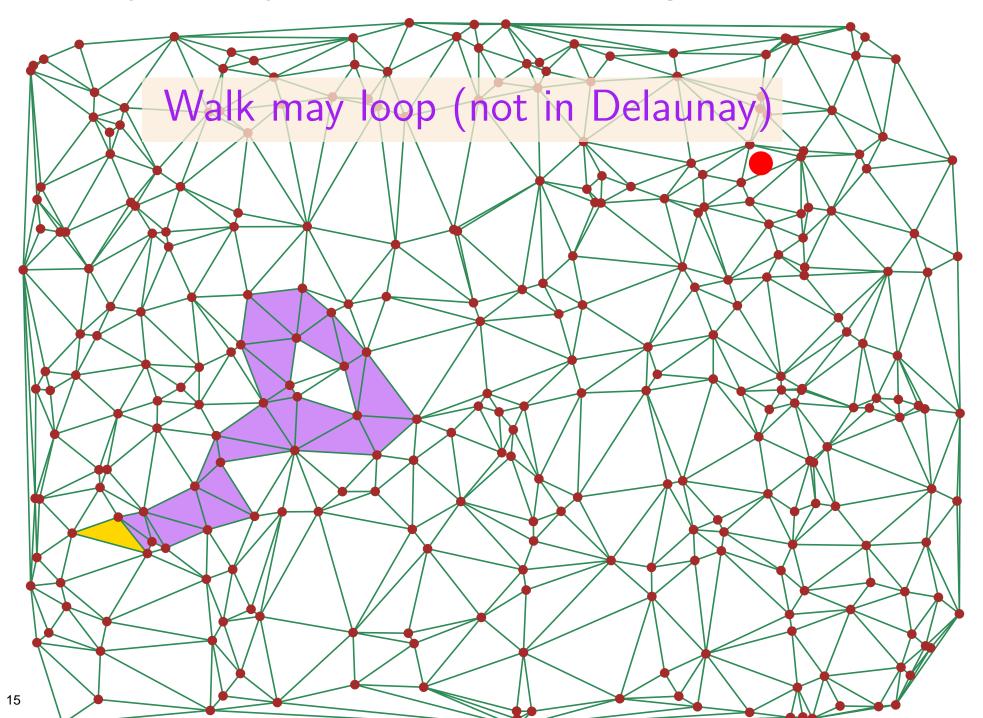
3D: 97 μ s

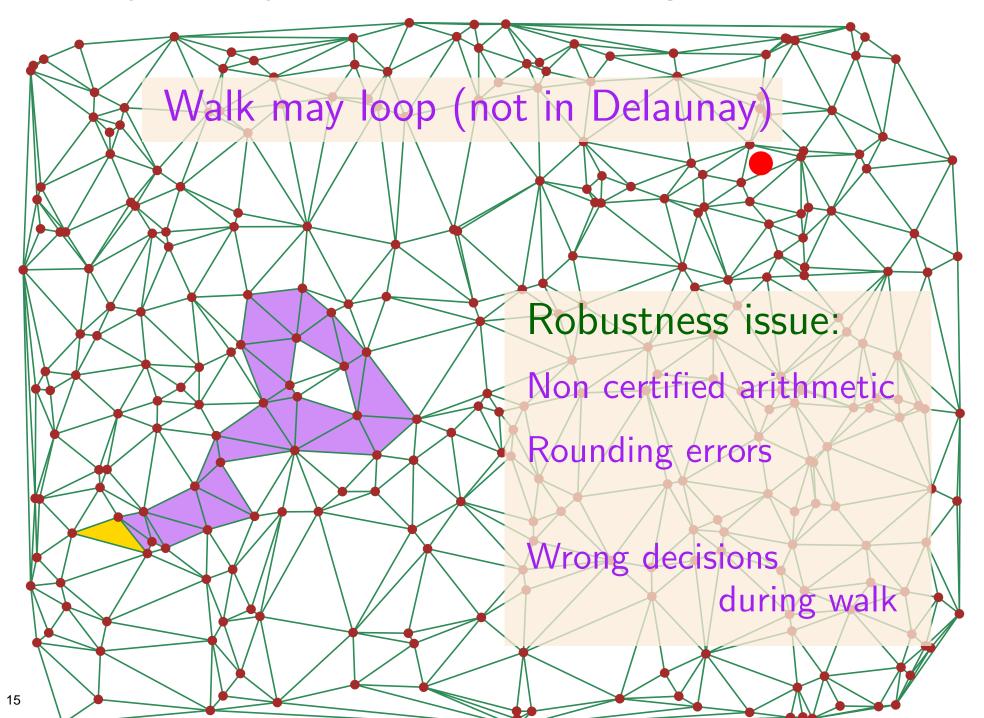
Much easier to code

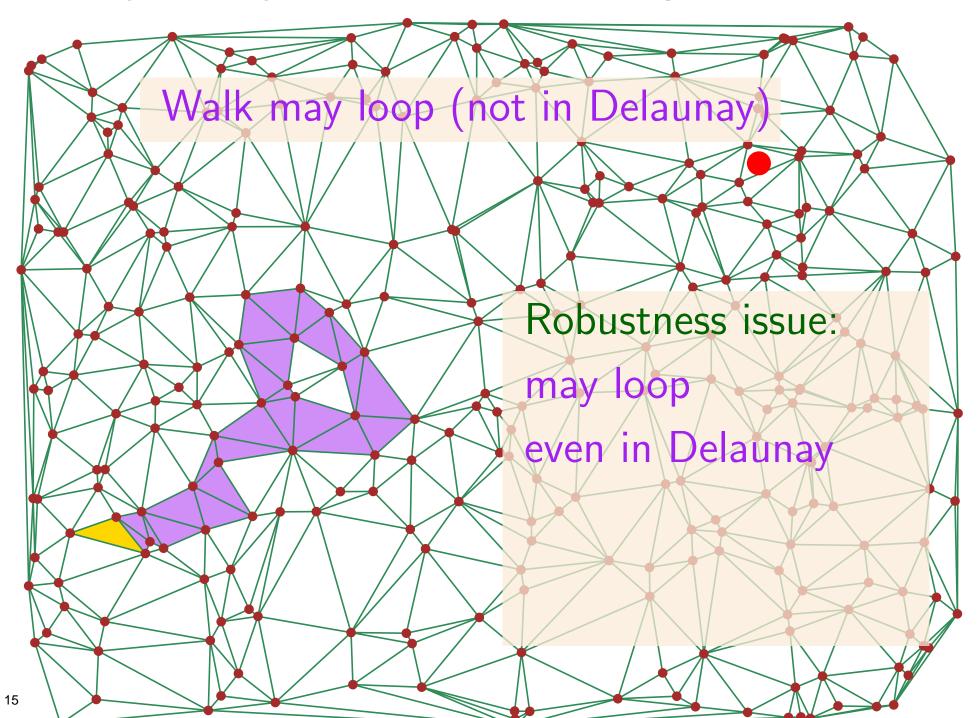


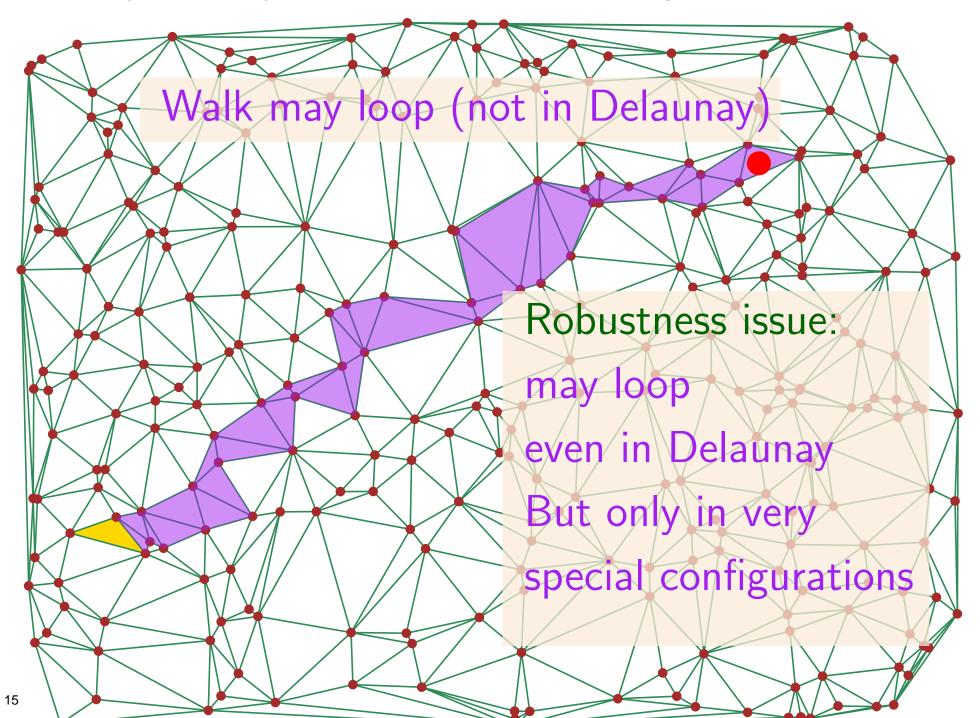


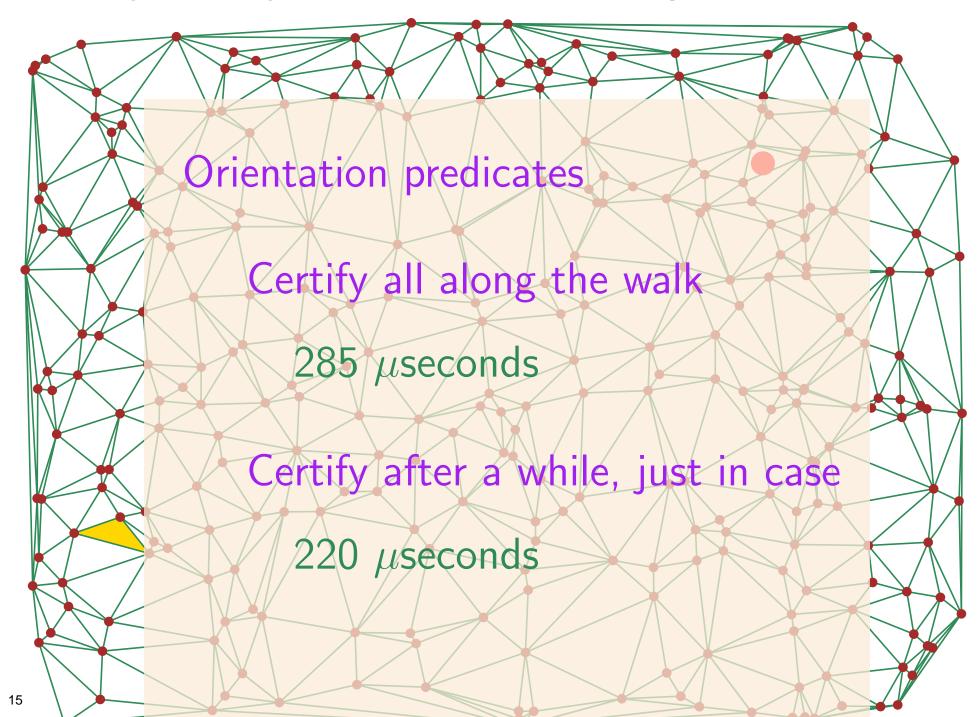




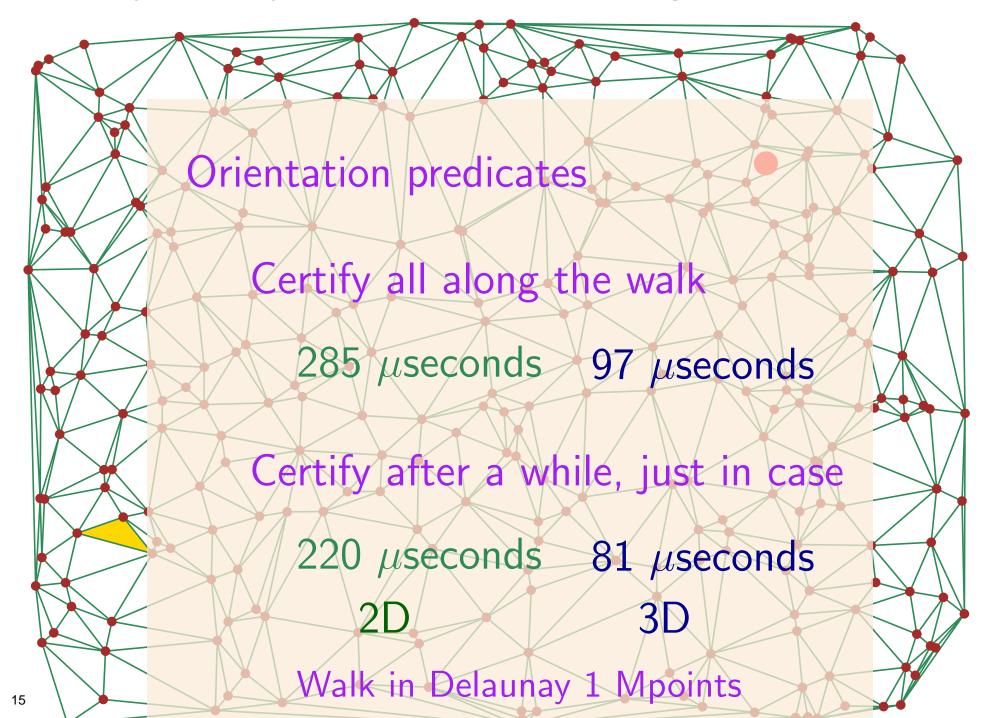


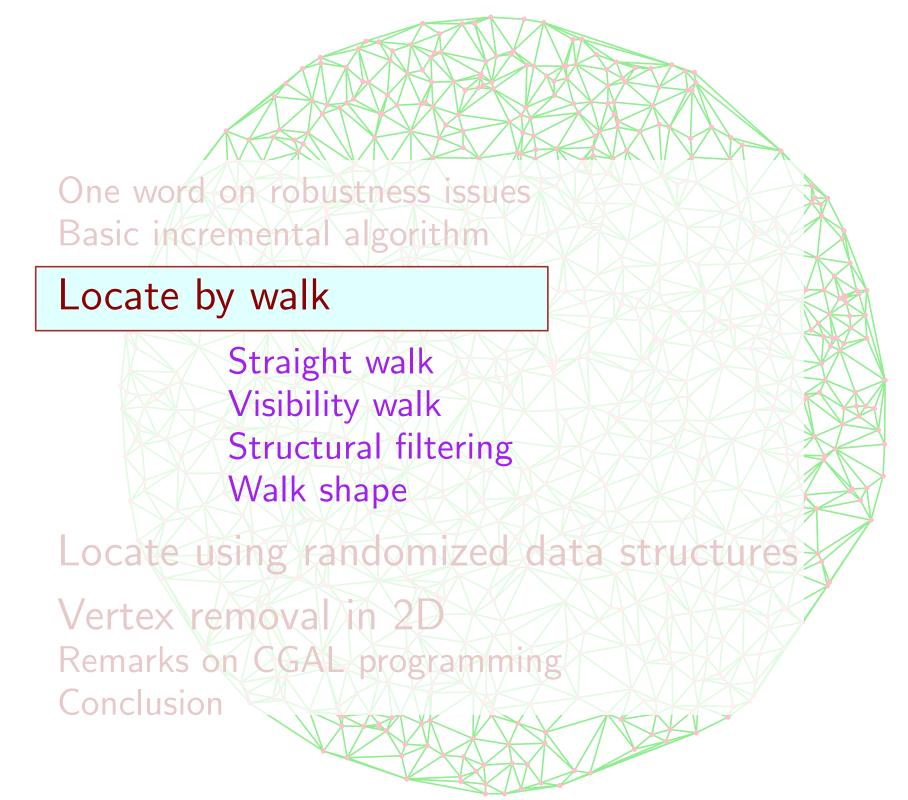


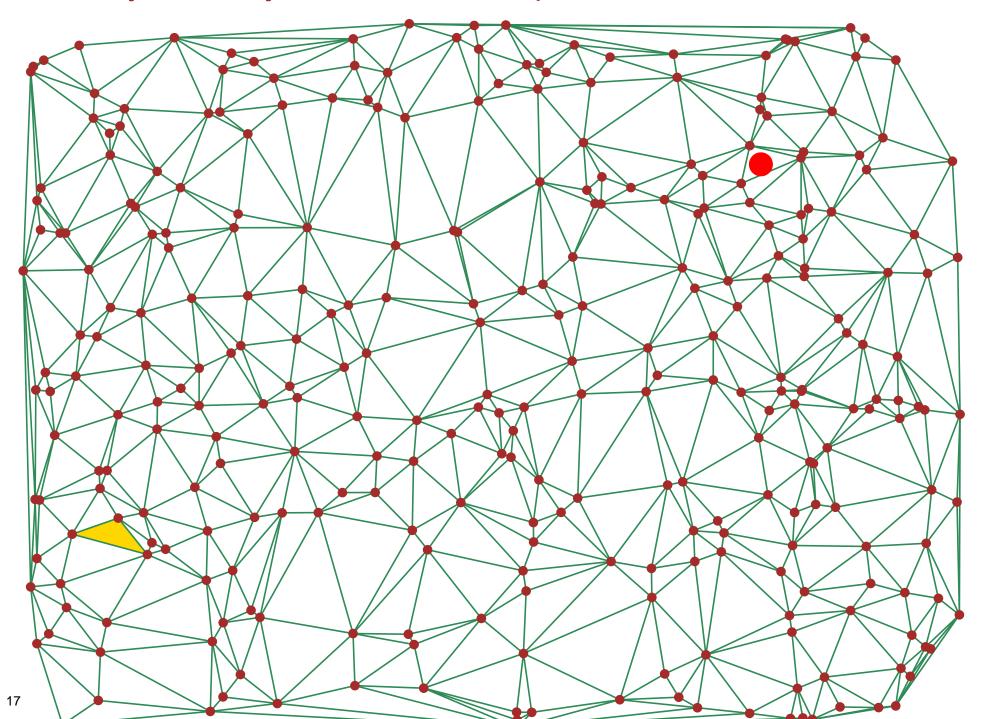


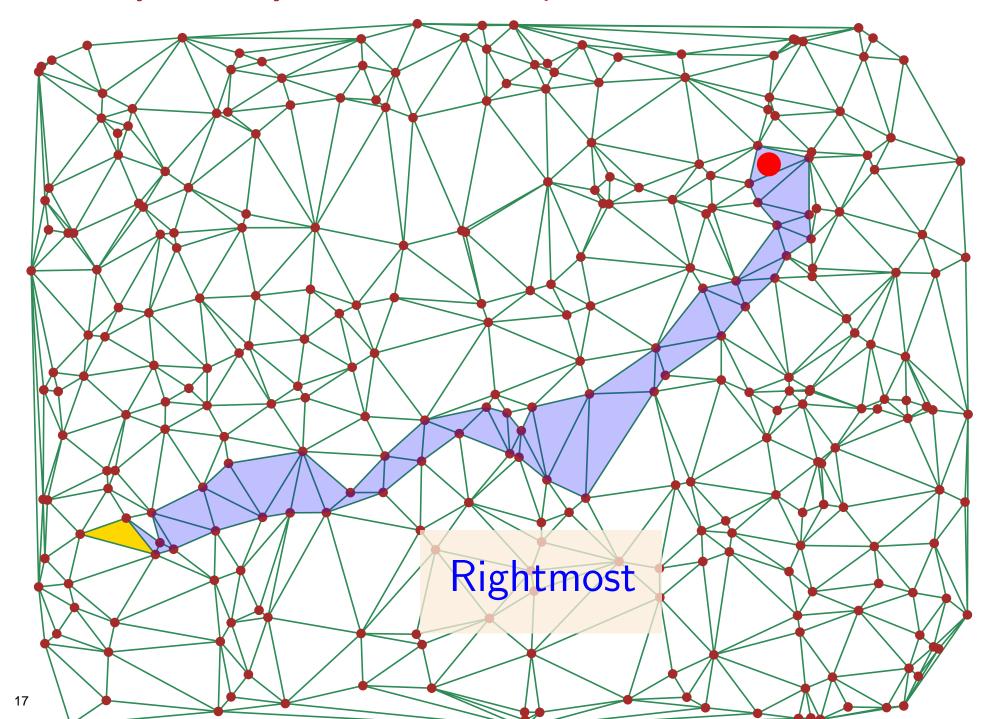


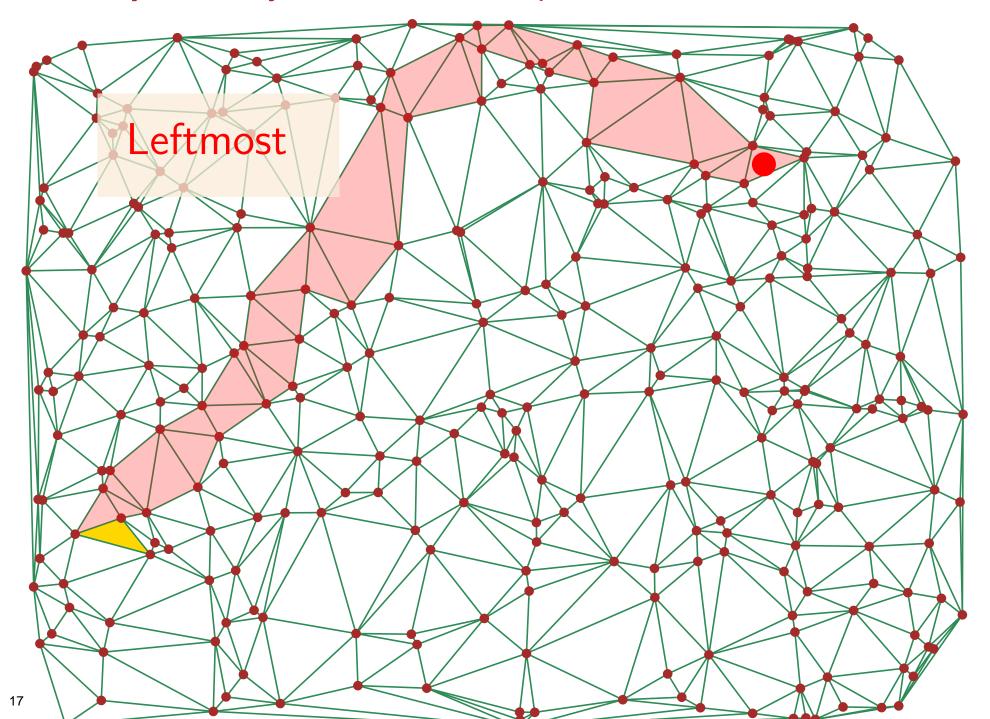
Locate by visibility walk - structural filtering

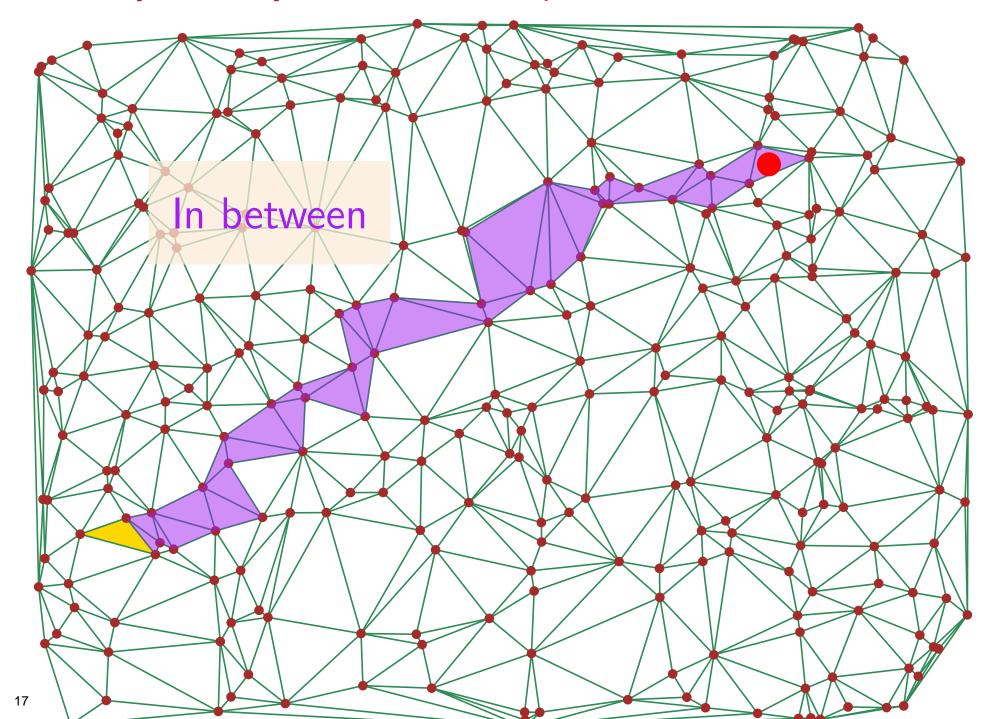


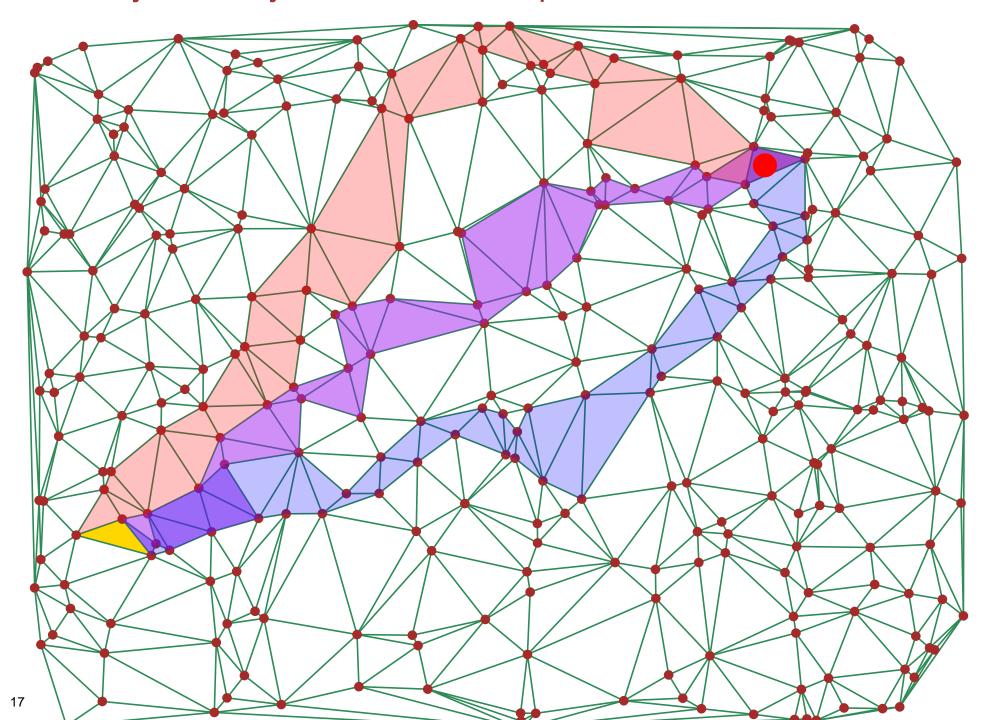




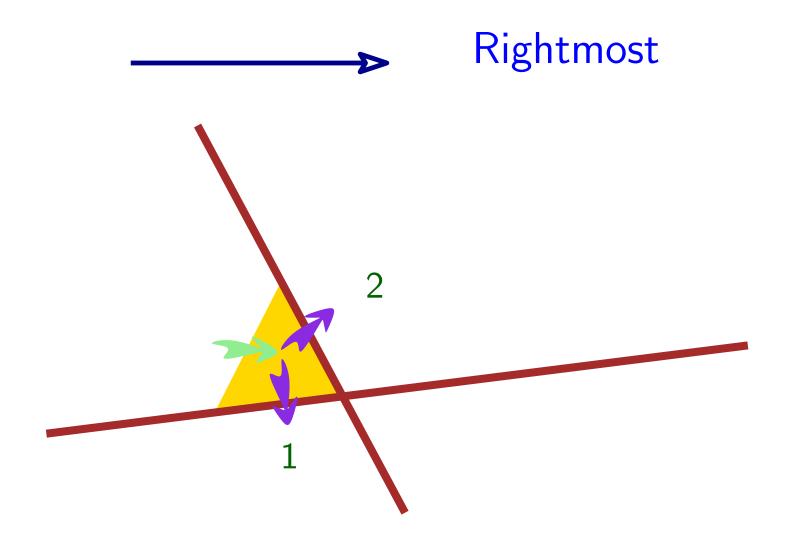




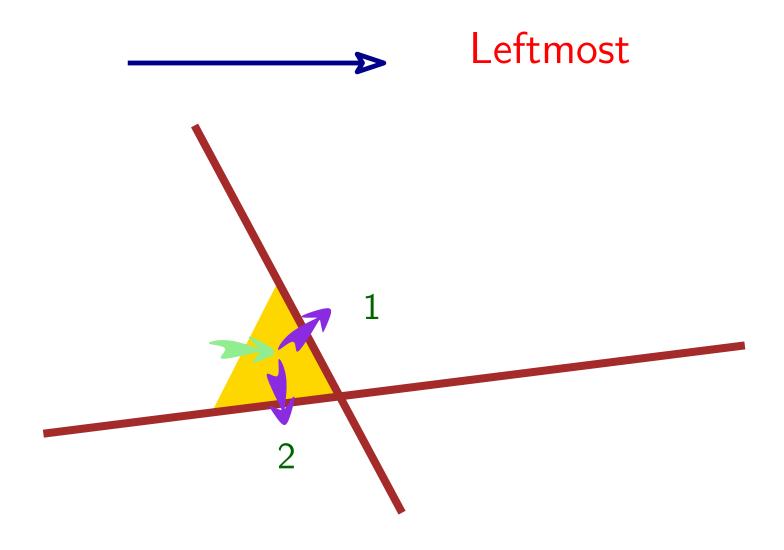




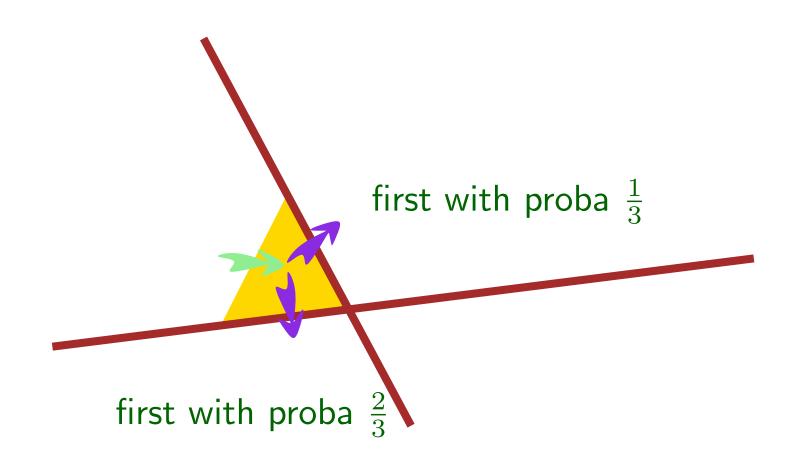
Turn counterclockwise from previous



Turn clockwise from previous



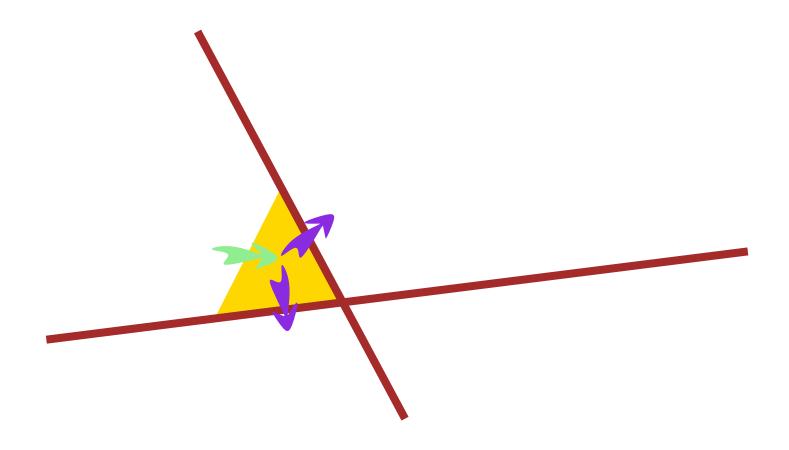
Balance left and right turns



Walk in Delaunay 1 Mpoints

Leftmost 220 μ seconds

Balanced 188 μ seconds



Locate by visibility walk

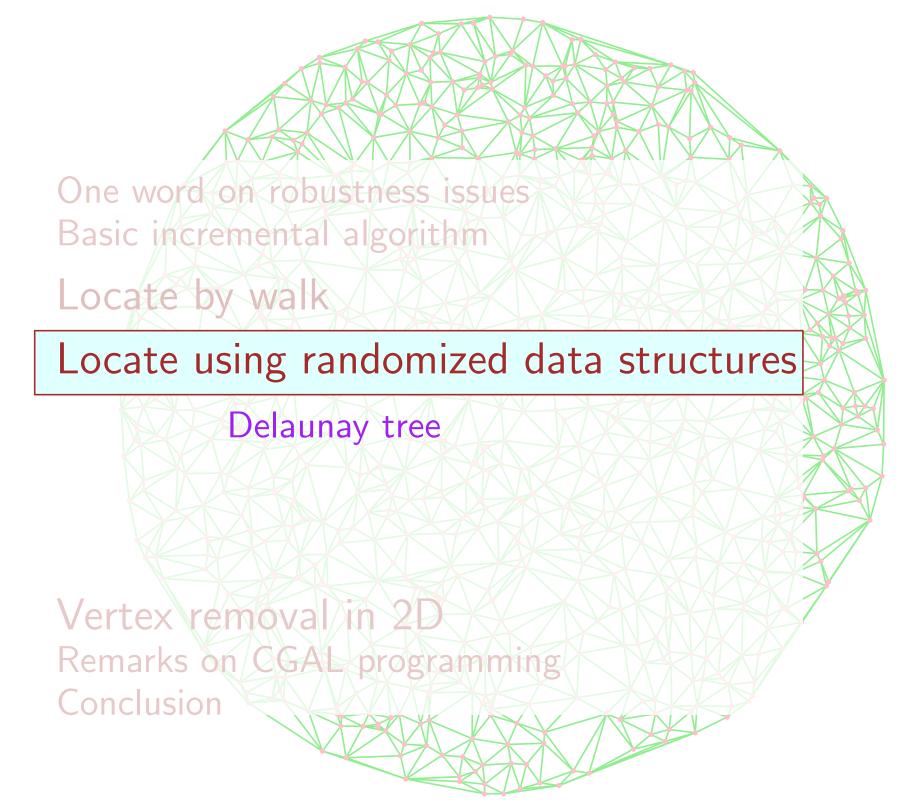
Walk in Delaunay 1 Mpoints

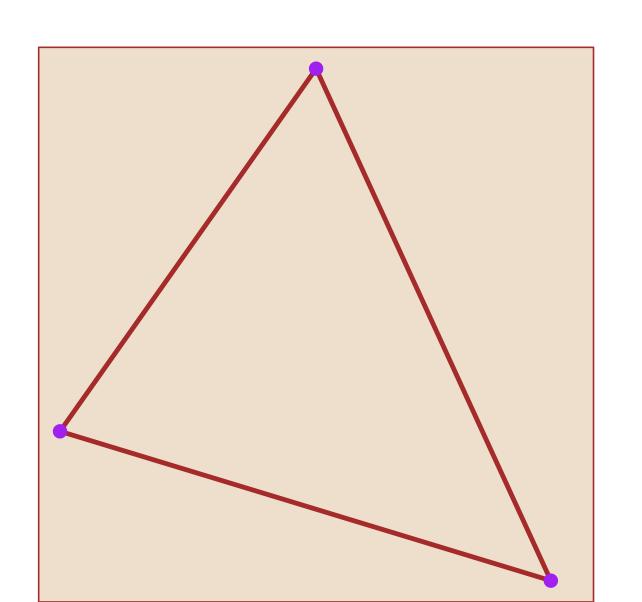
Straight walk 324 μ seconds

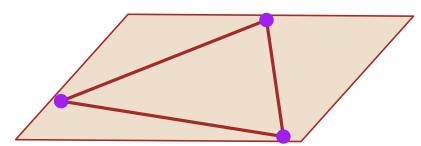
Visibility walk 285 μ seconds

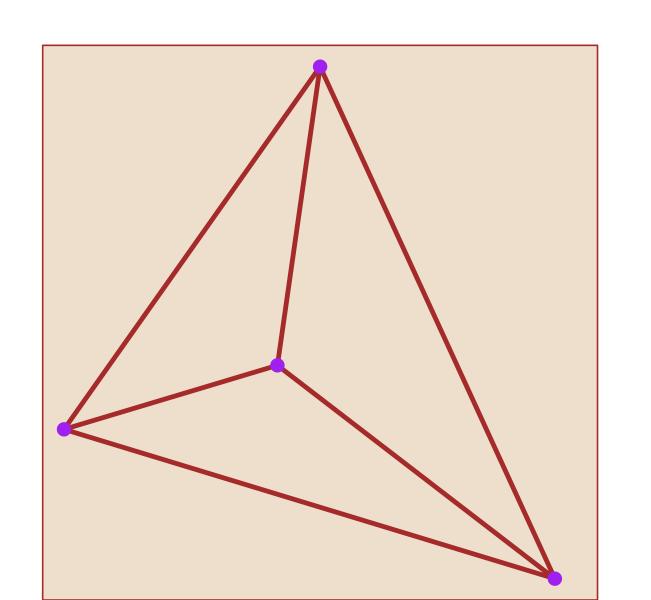
Structural filtering 220 μ seconds

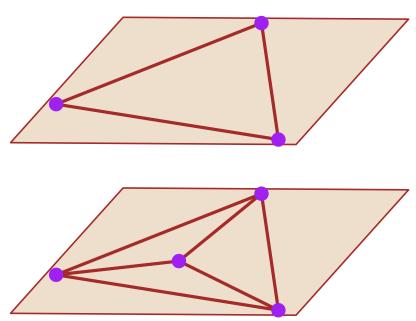
Balanced walk 188 μ seconds

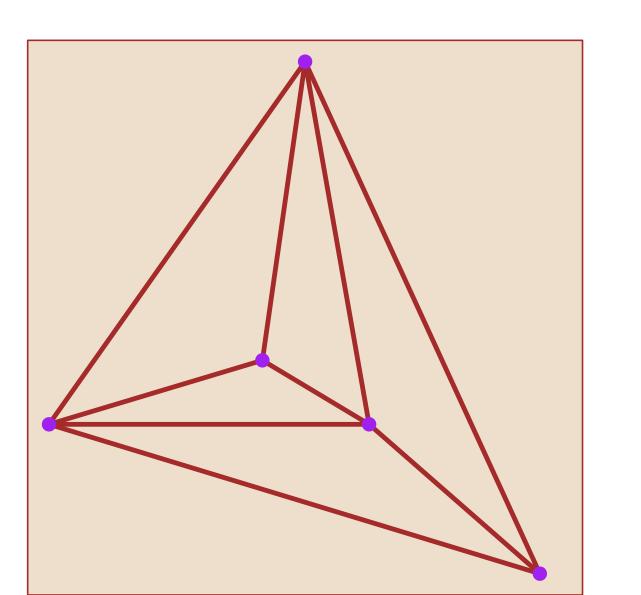


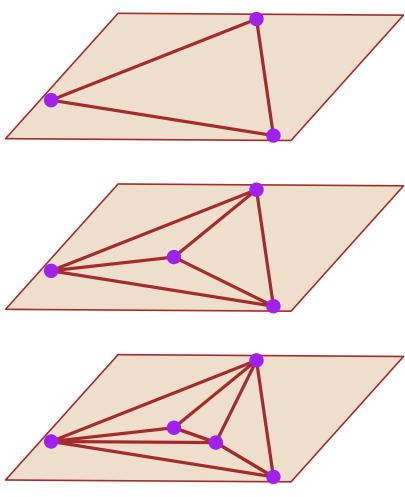


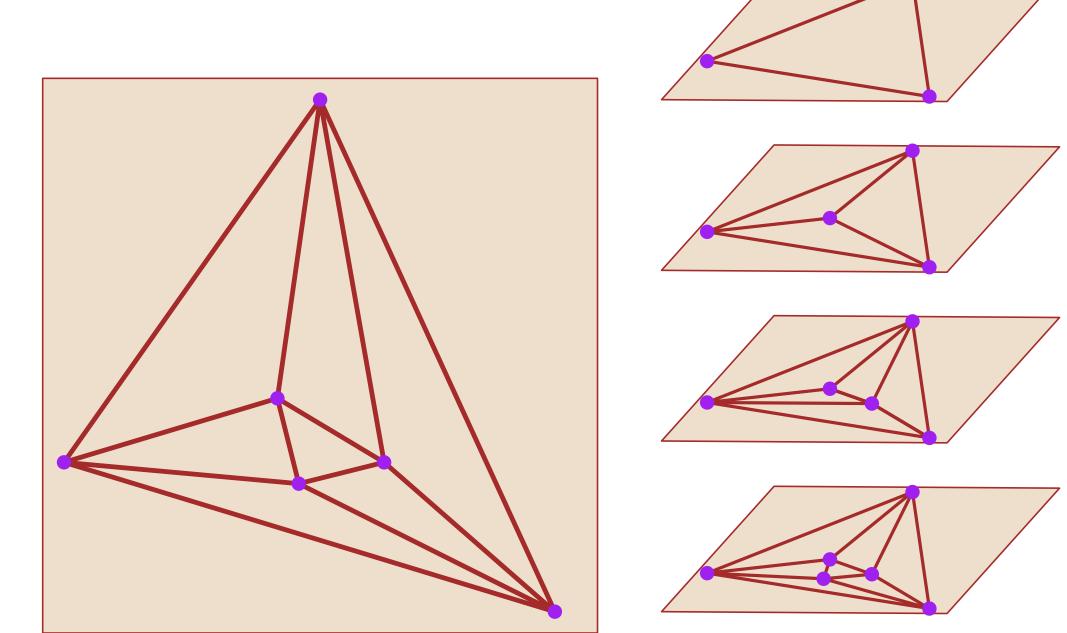


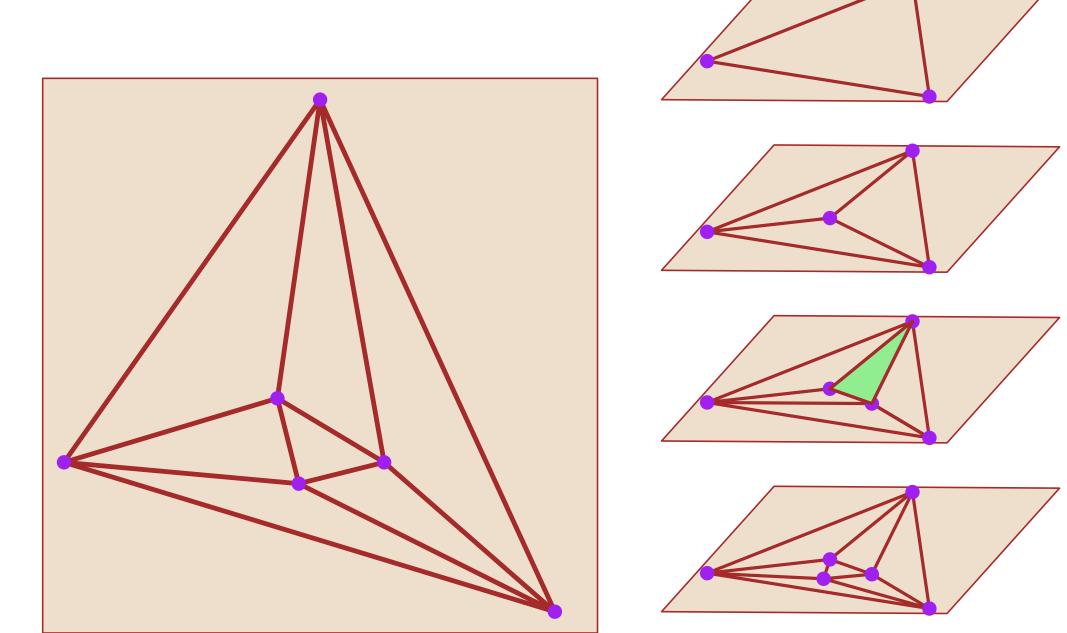


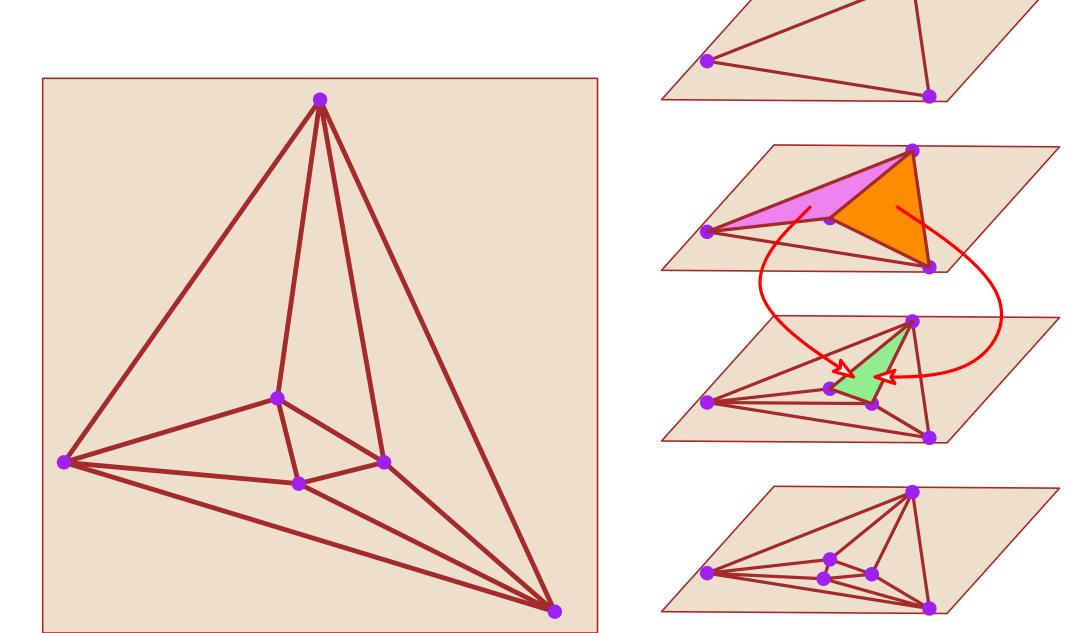


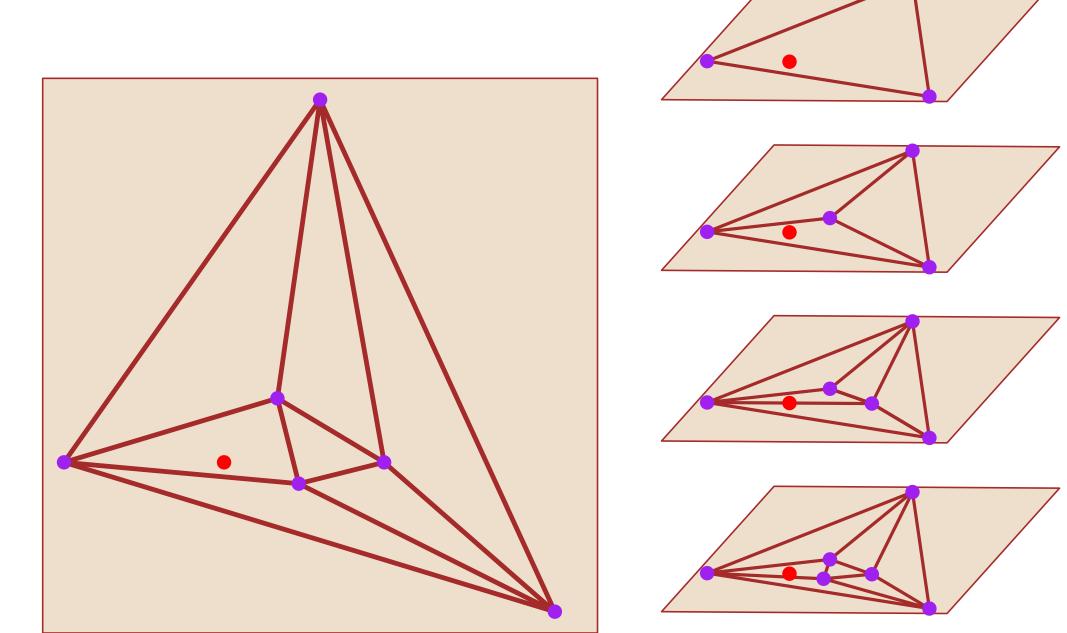


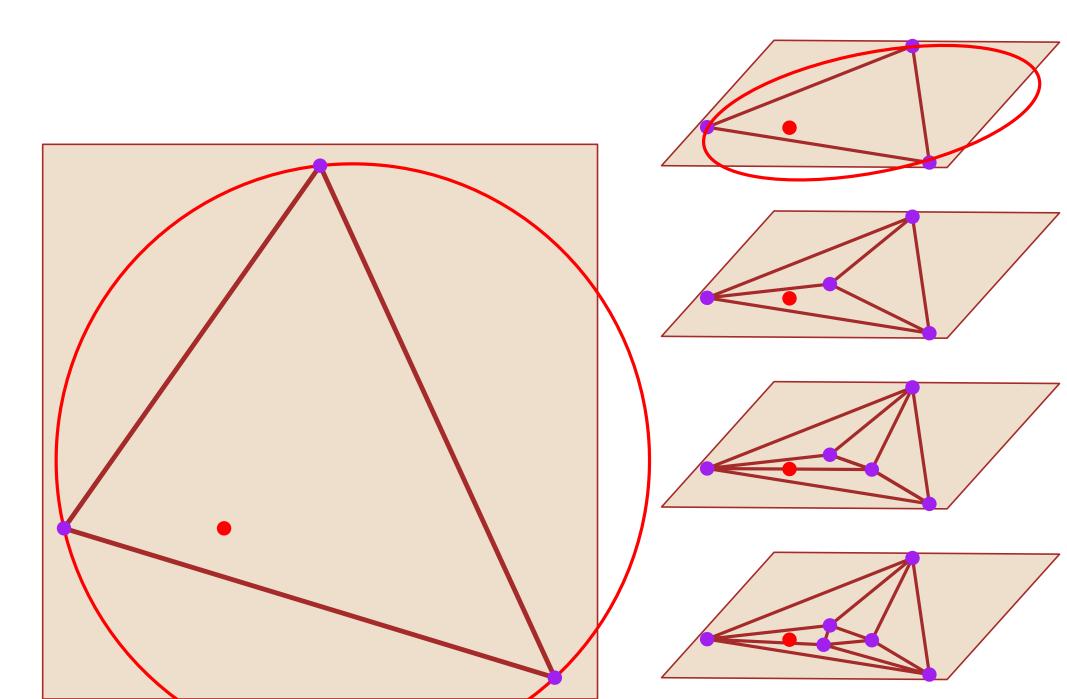


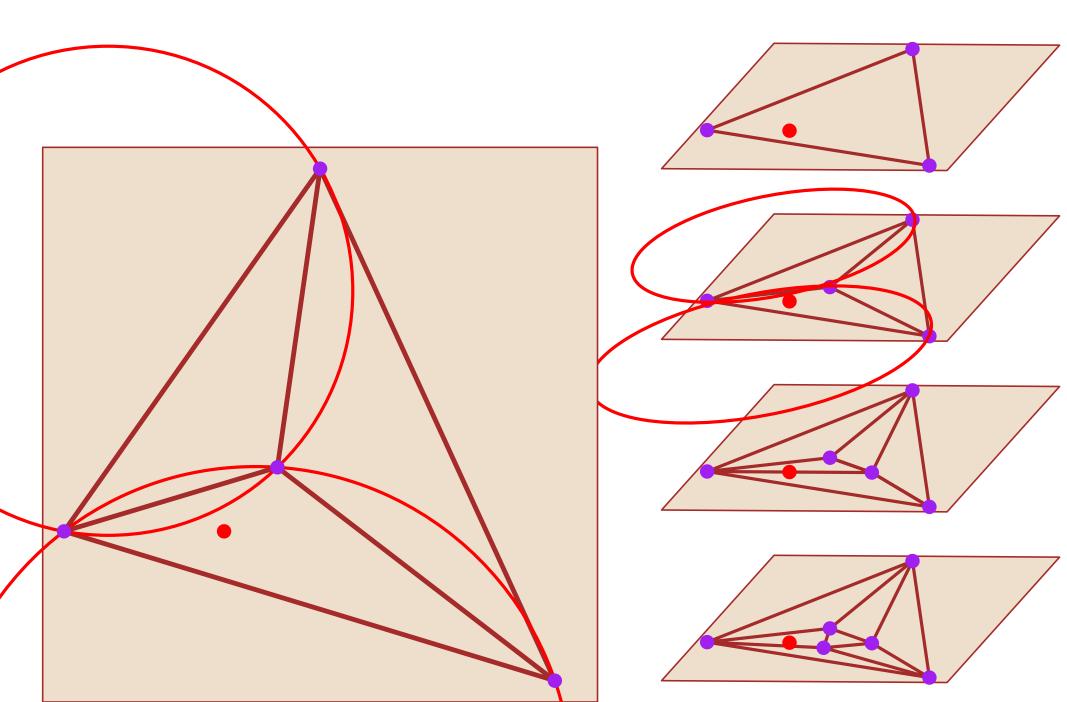


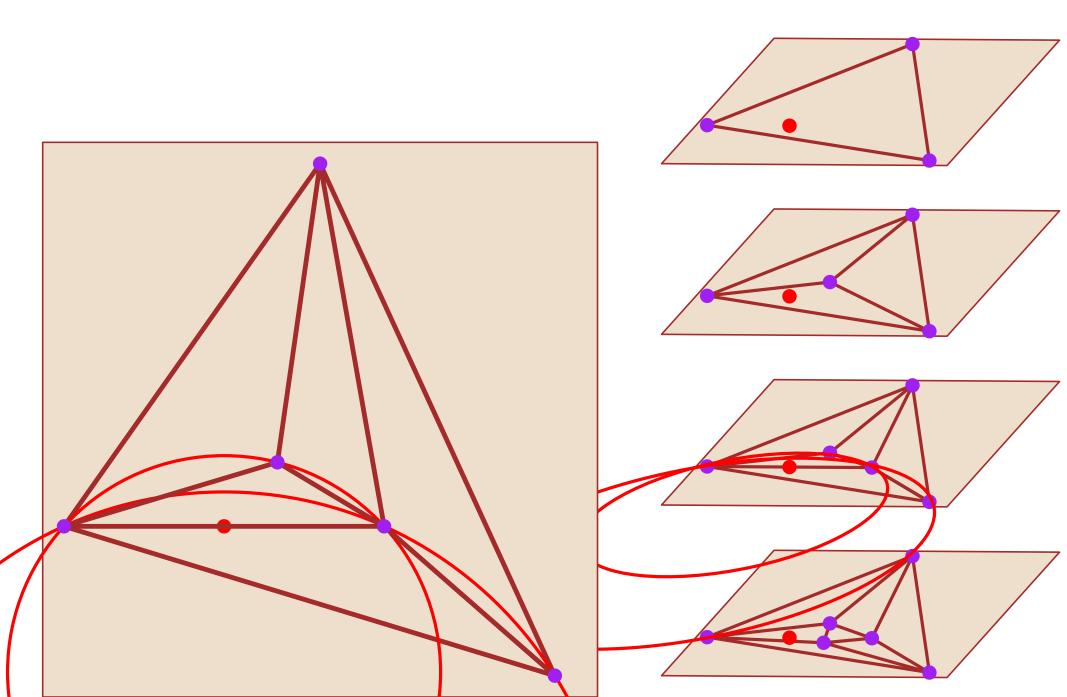


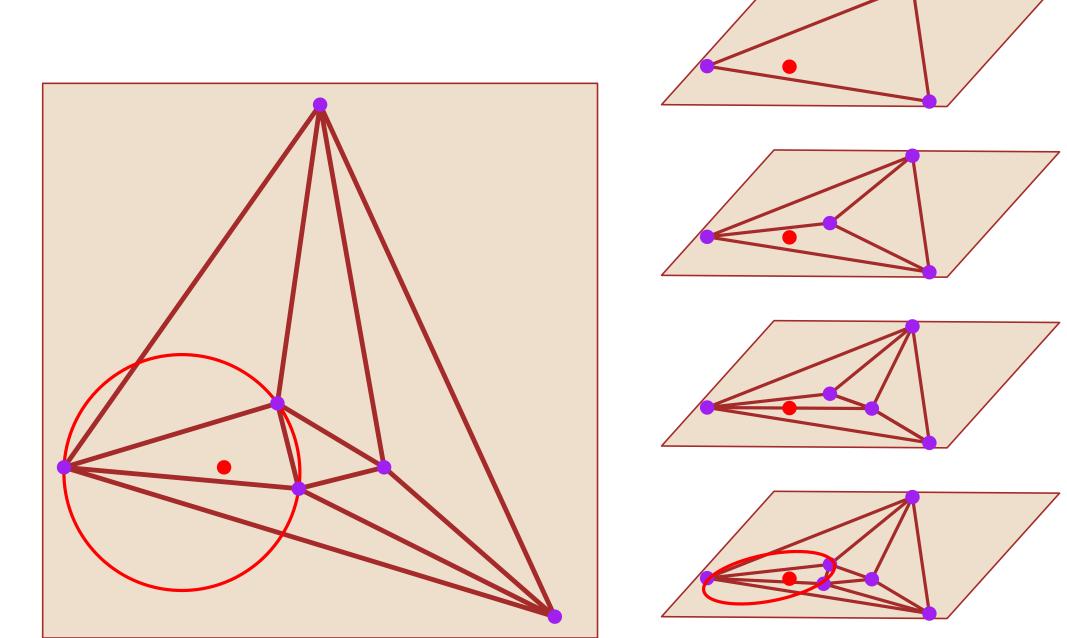






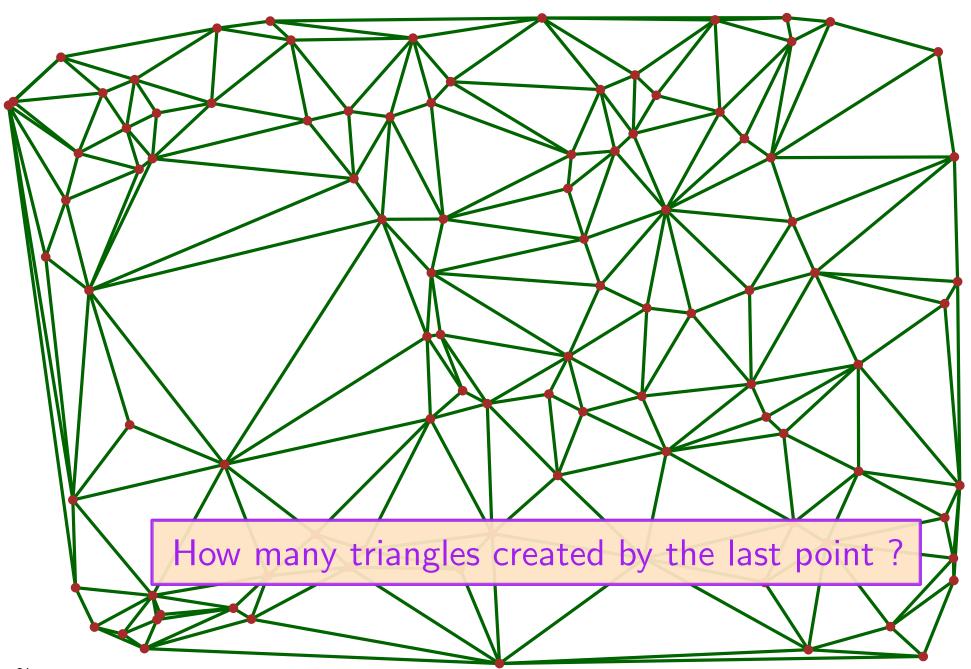


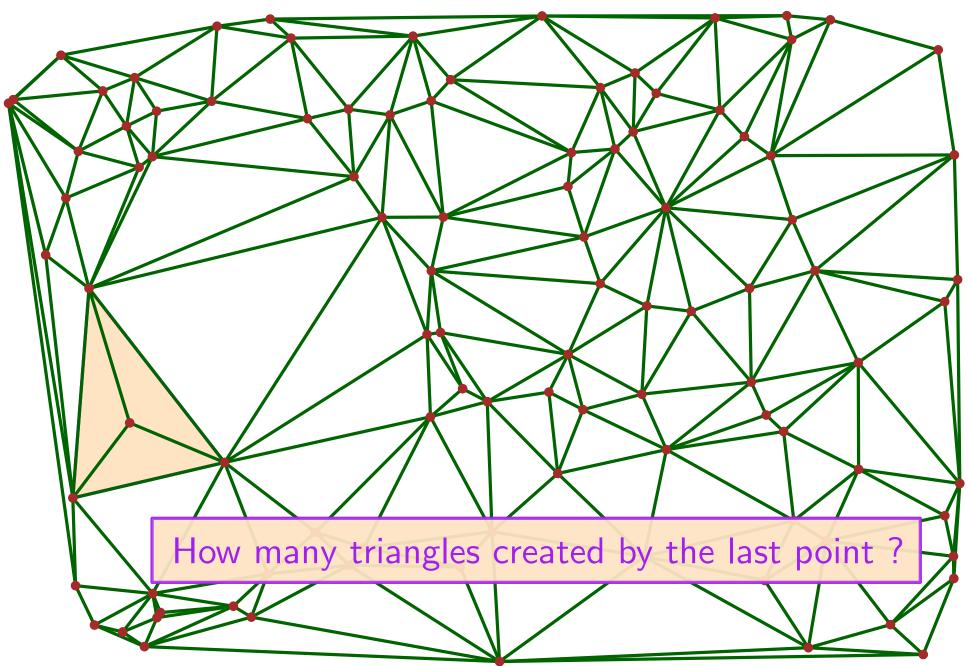


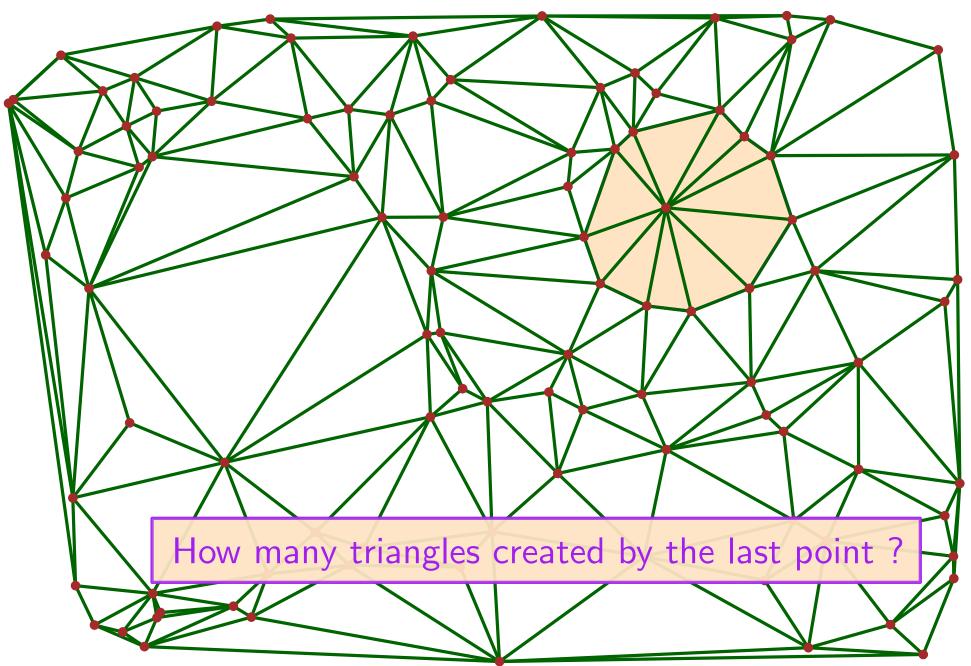


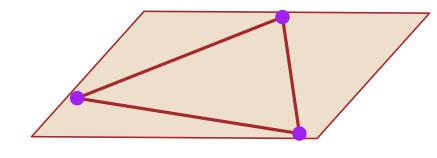
locate based on incircle predicate



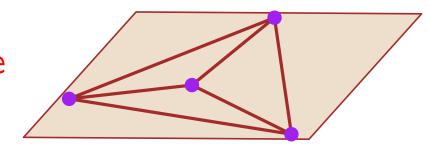






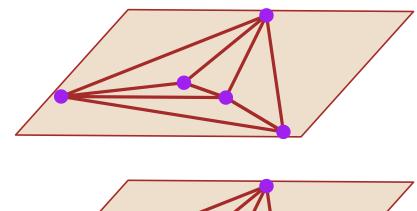


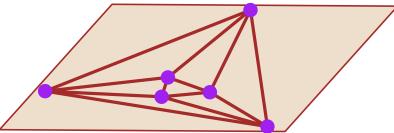
locate based on incircle predicate

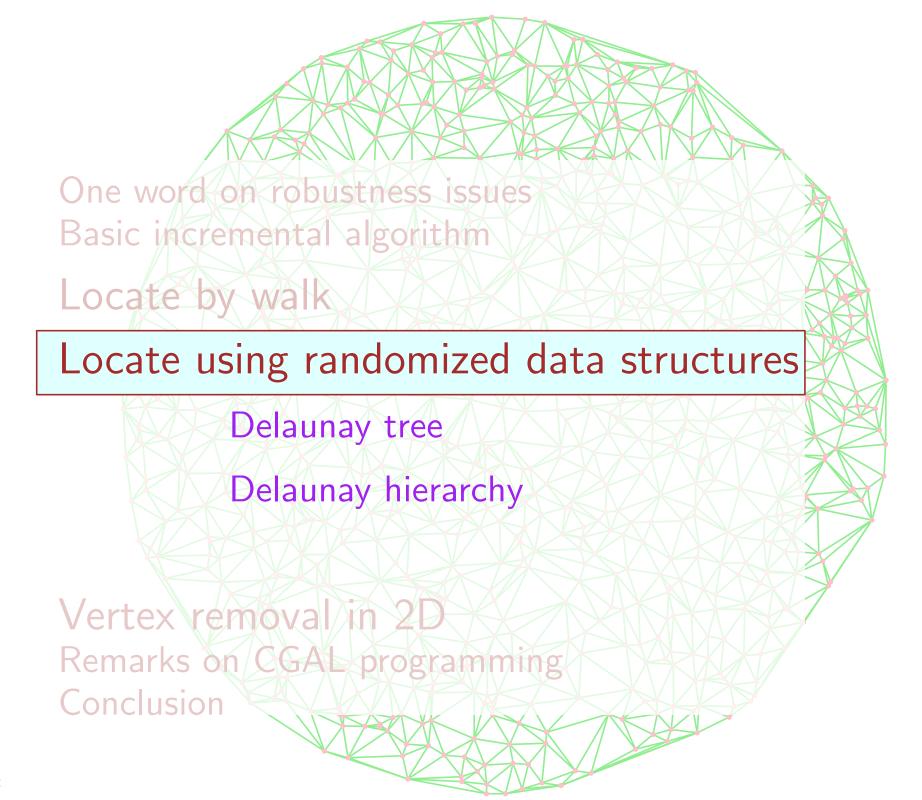


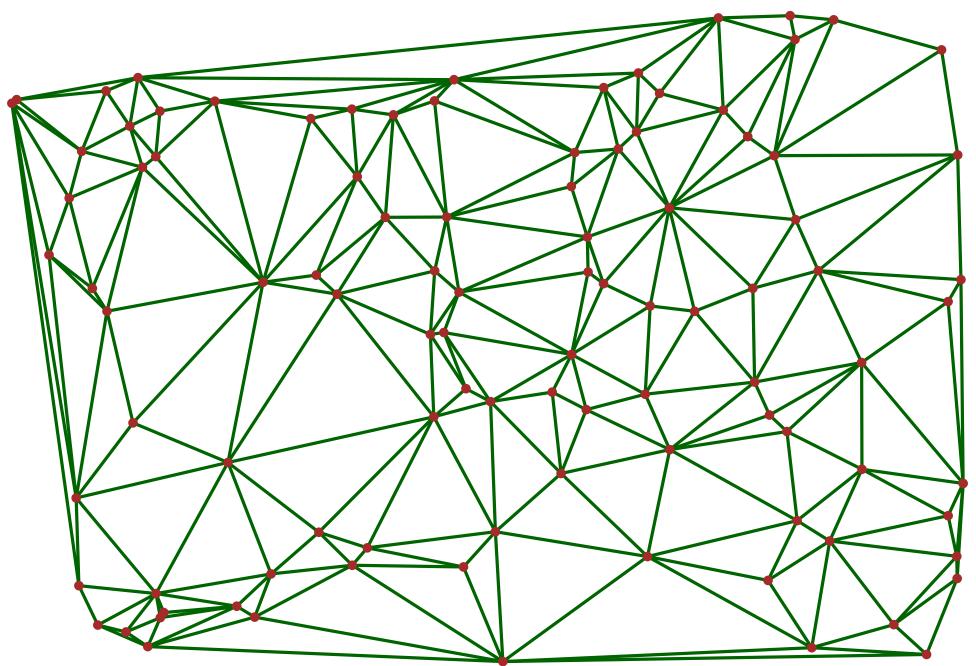
triangles in the Delaunay tree

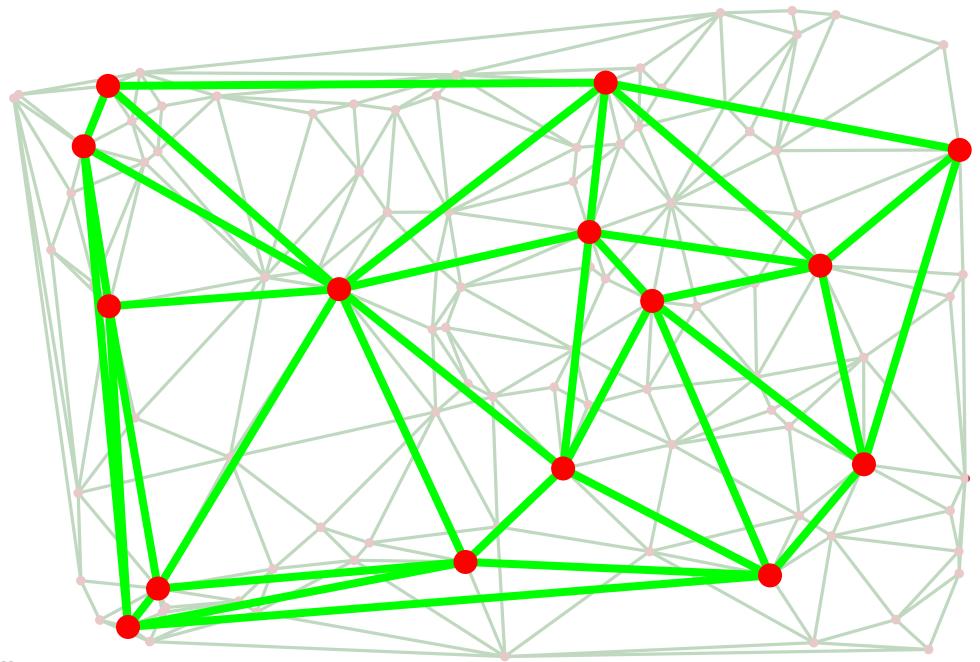
=6n (randomized)

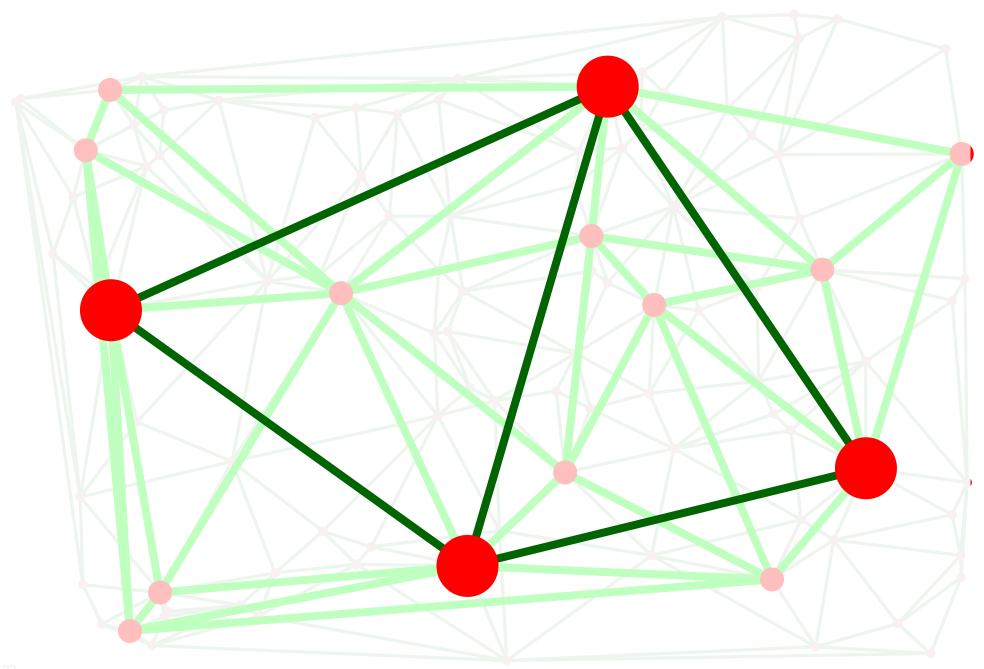


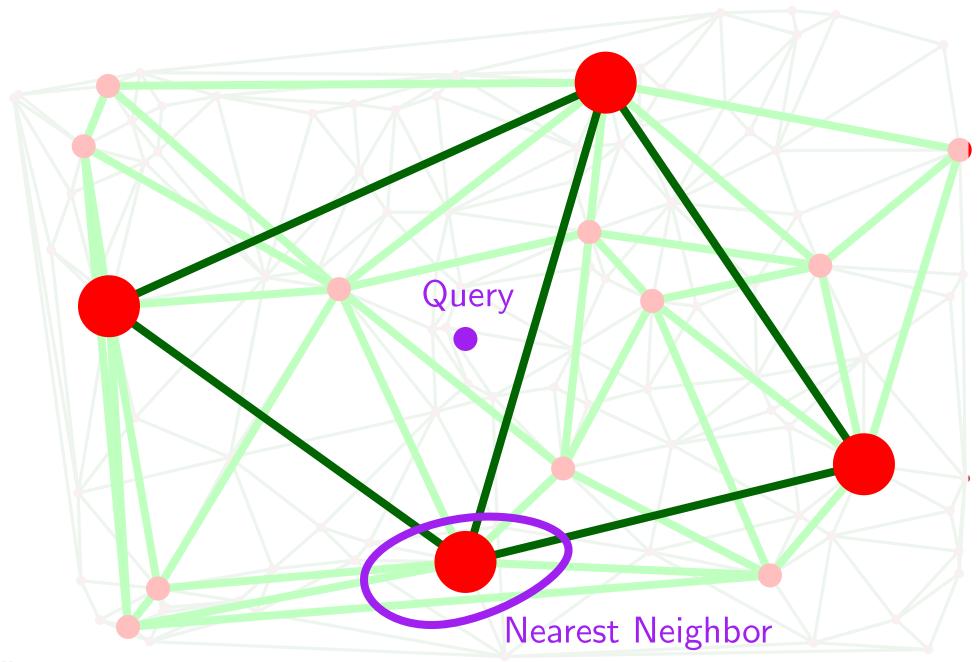


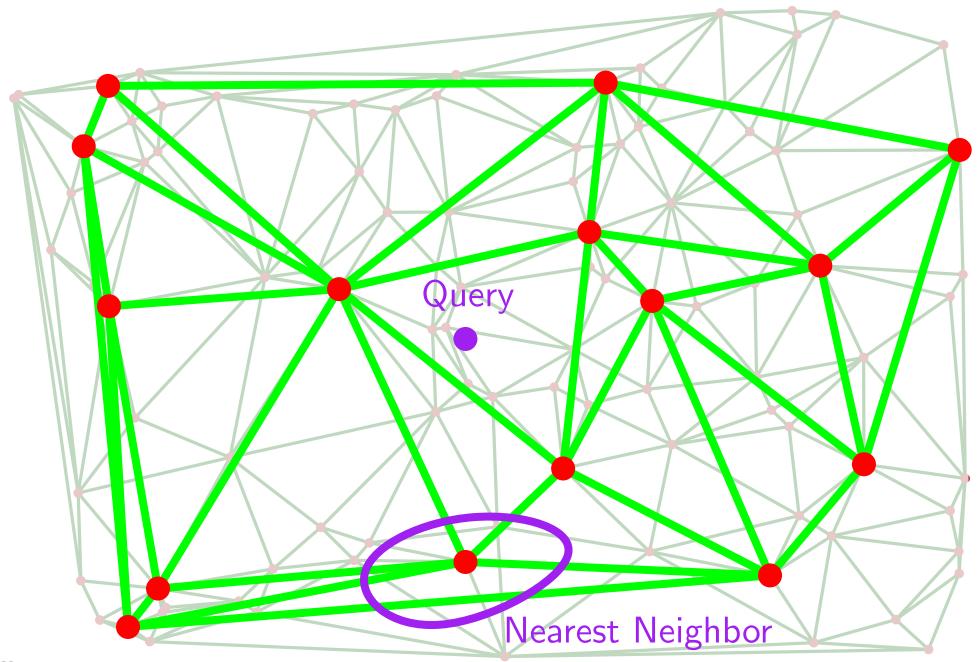


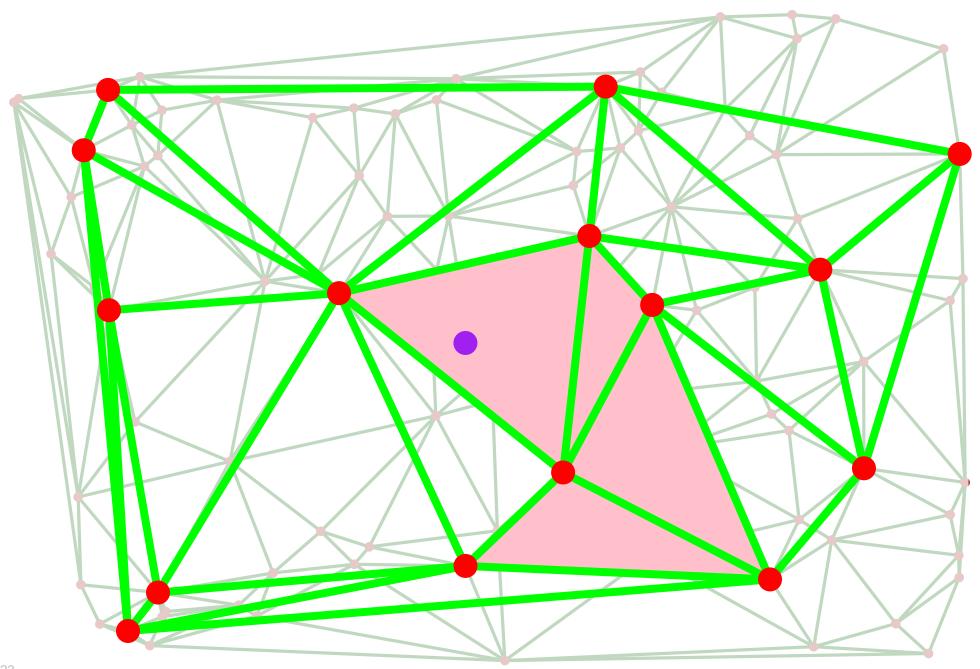


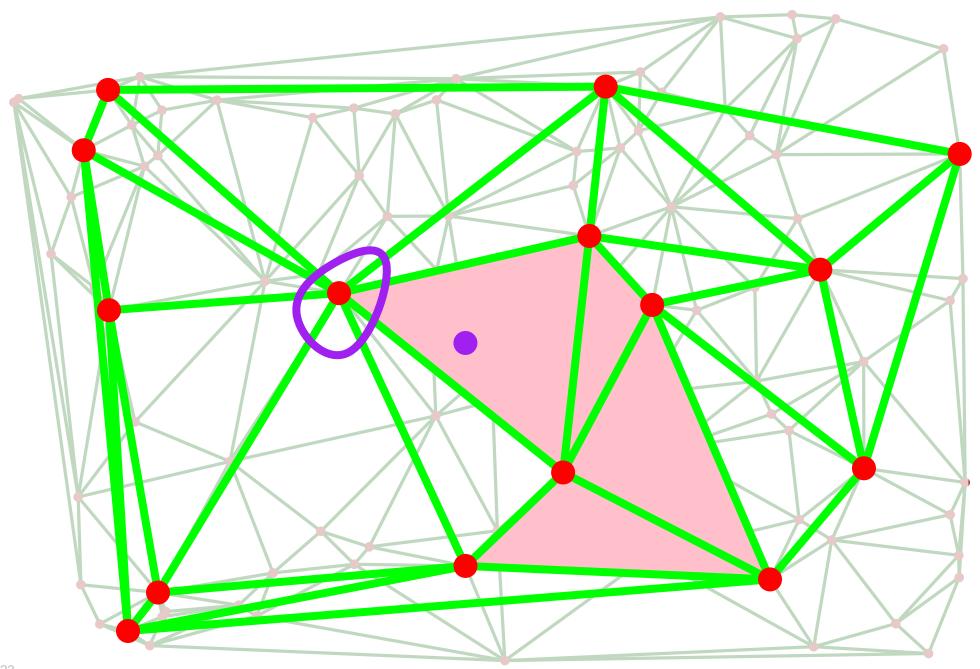


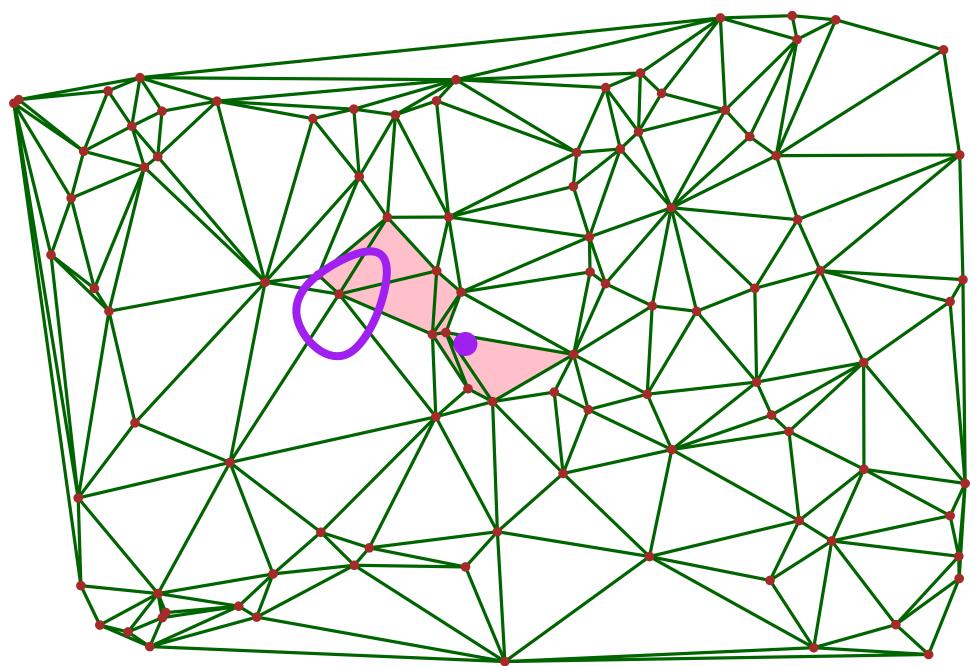












The Delaunay tree

locate based on incircle predicate

triangles in the Delaunay tree

=6n (randomized)

The Delaunay hierarchy

The Delaunay tree

based on orientation predicate

locate based on incircle predicate

triangles in the hierarchy

can be chosen

 $=1.03\times 2n$ (expected)

triangles in the Delaunay tree

=6n (randomized)

The Delaunay hierarchy

The Delaunay tree

based on orientation predicate

locate based on incircle predicate

triangles in the hierarchy

can be chosen

 $=1.03\times 2n$ (expected)

triangles in the Delaunay tree

=6n (randomized)

 $O(n \log n)$

The Delaunay hierarchy

The Delaunay tree

based on orientation predicate

locate based on incircle predicate

triangles in the hierarchy

can be chosen

 $=1.03\times 2n$ (expected)

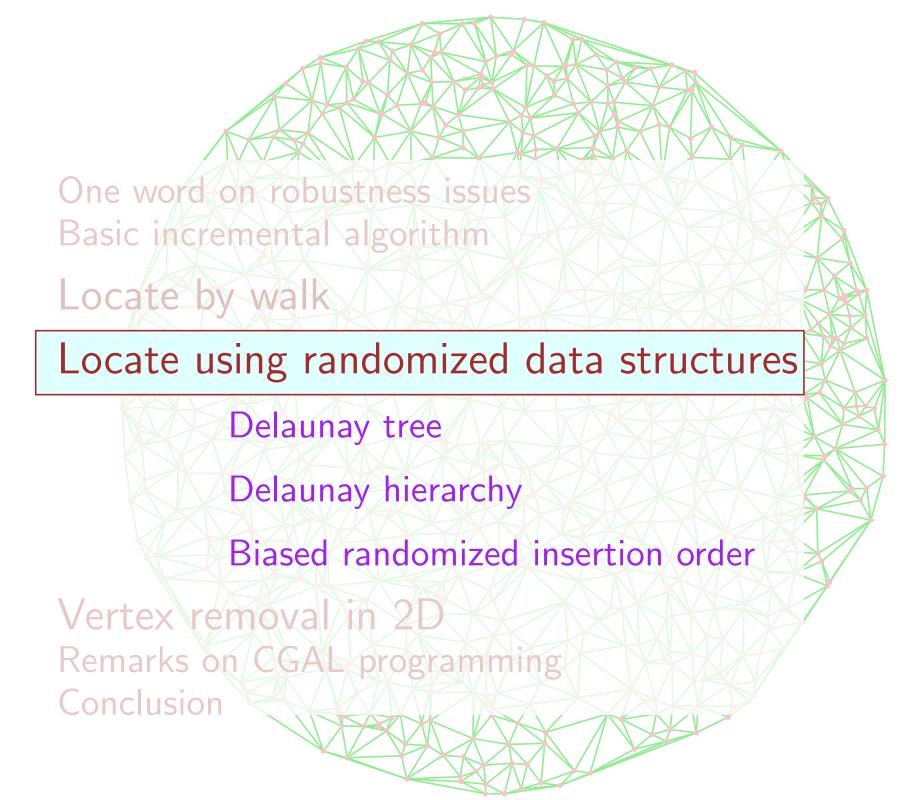
triangles in the Delaunay tree

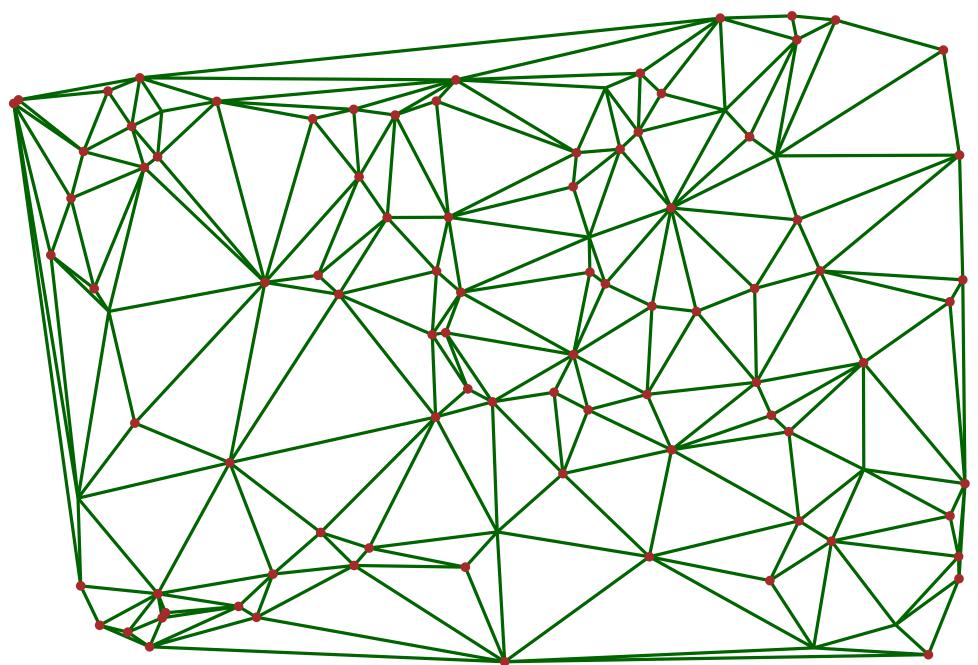
=6n (randomized)

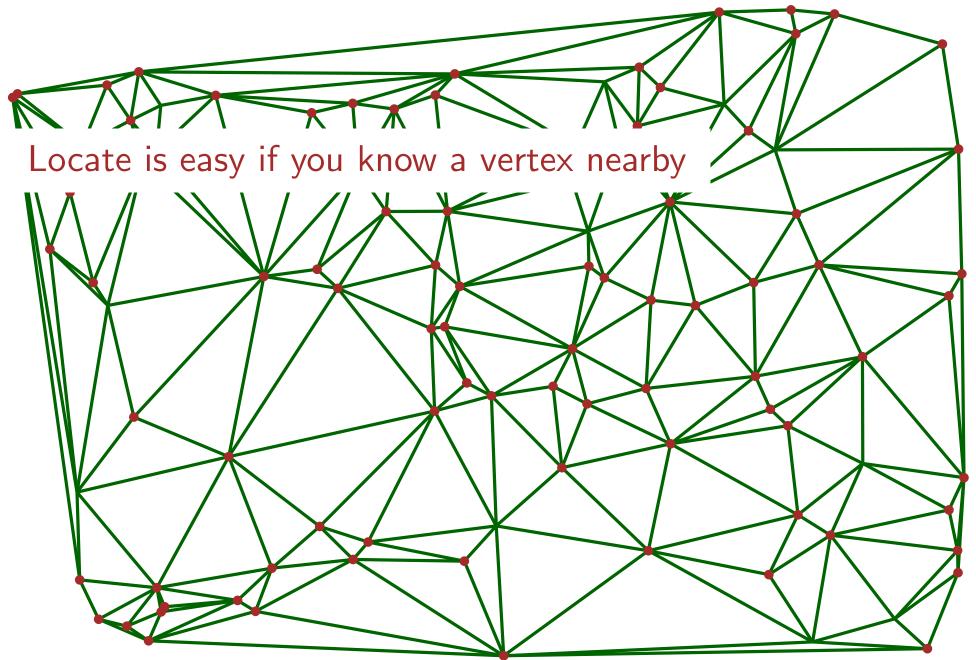
2.3 seconds

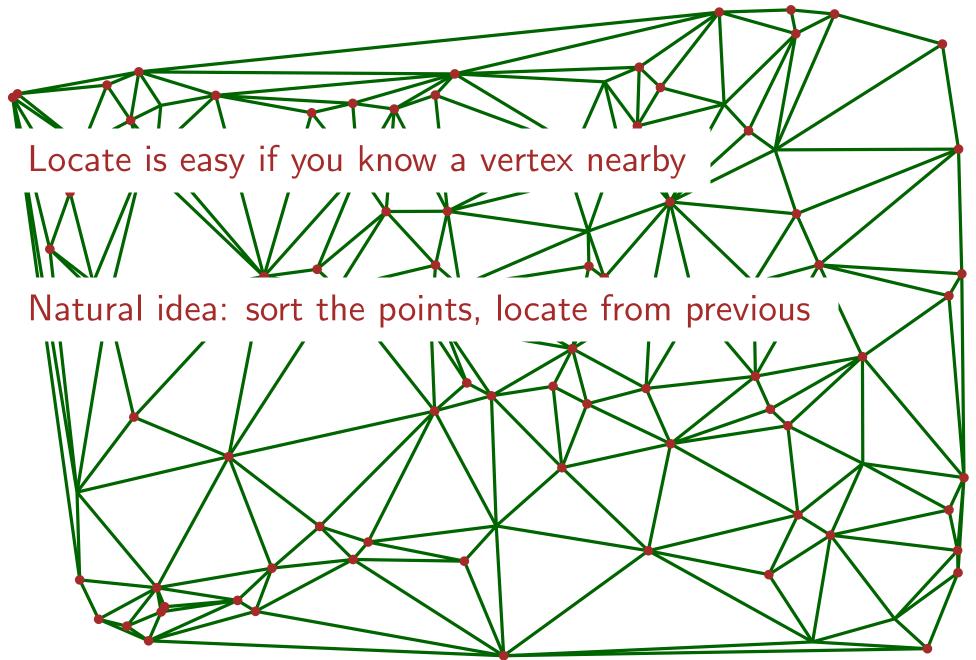
17 seconds

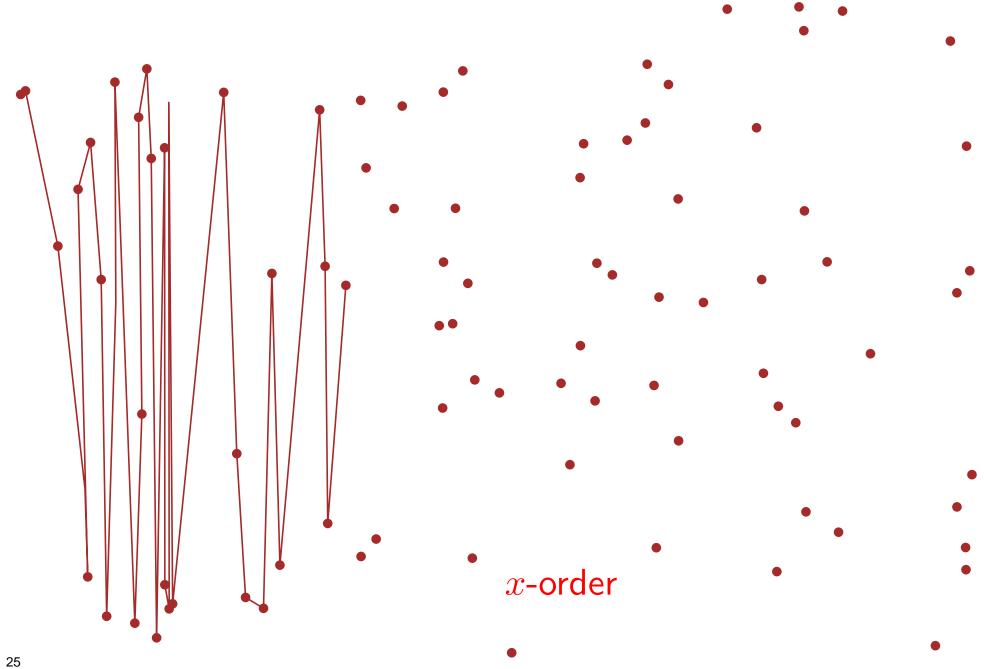
50000 random points (original benchmarks in 2000).

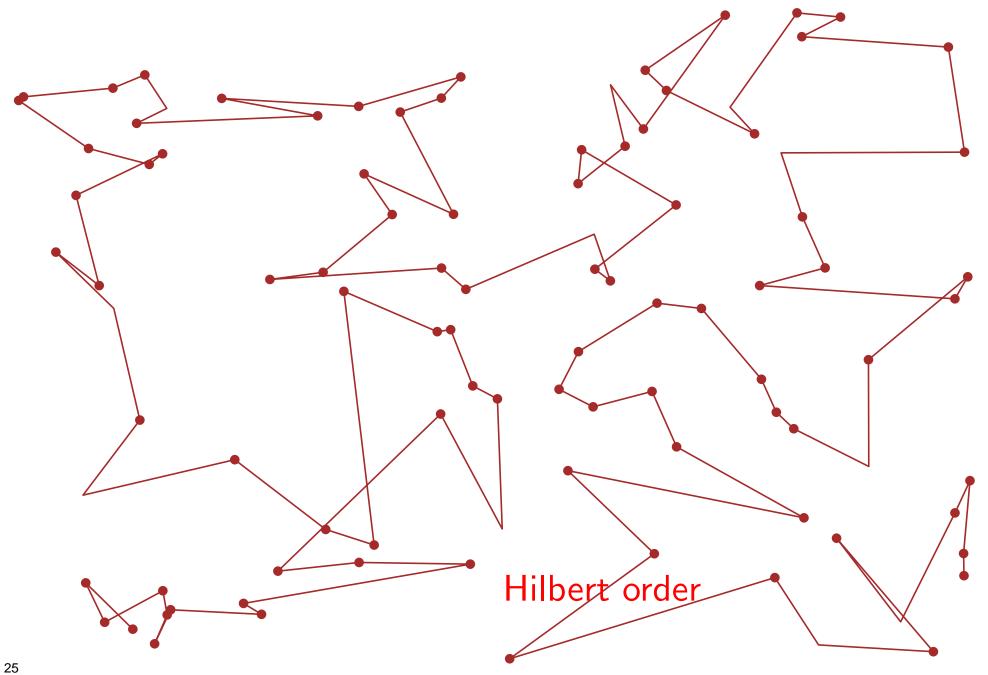


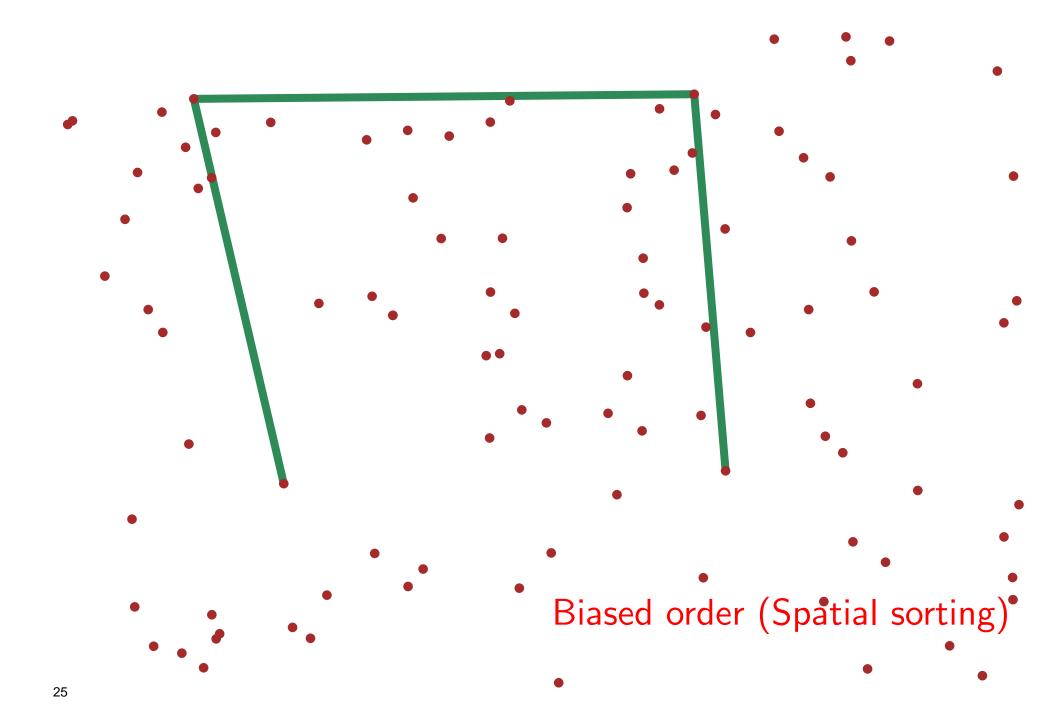


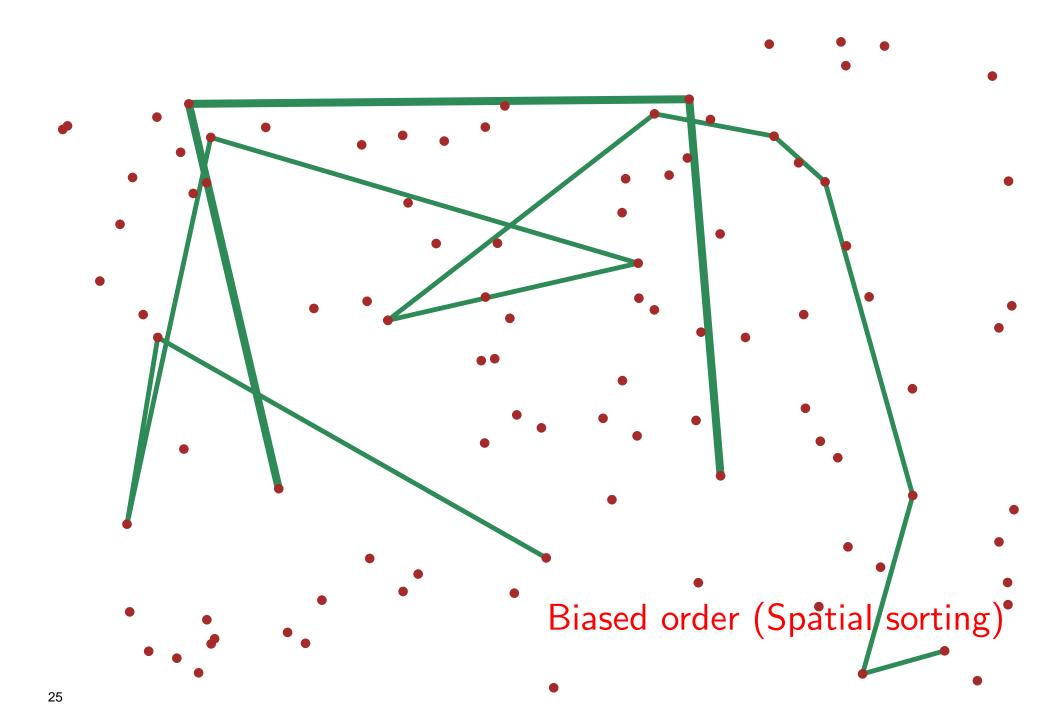


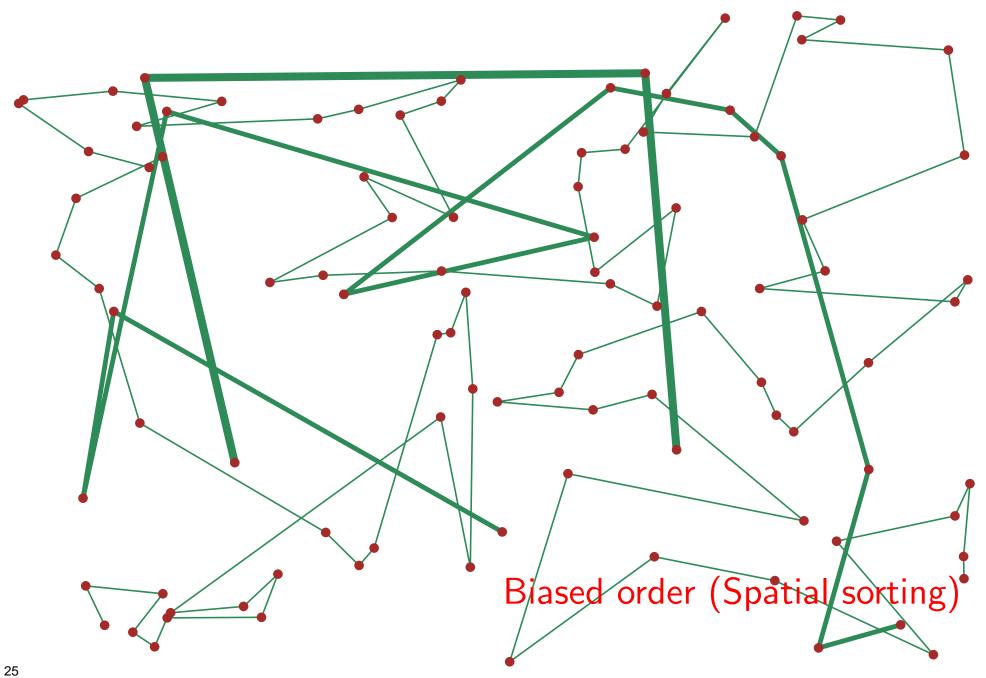












Delaunay 2D 1M random points locate using Delaunay hierarchy	6/seconds
random order	157 seconds
x-order	3 seconds
Hilbert order	0.8 seconds
Biased order (Spatial sorting)	0.7 seconds

Delaunay 2D 100K parabola points locate using Delaunay hierarchy 0.3 seconds 128 seconds random order 632 seconds x-order 46 seconds Hilbert order Biased order (Spatial sorting) 0.3 seconds

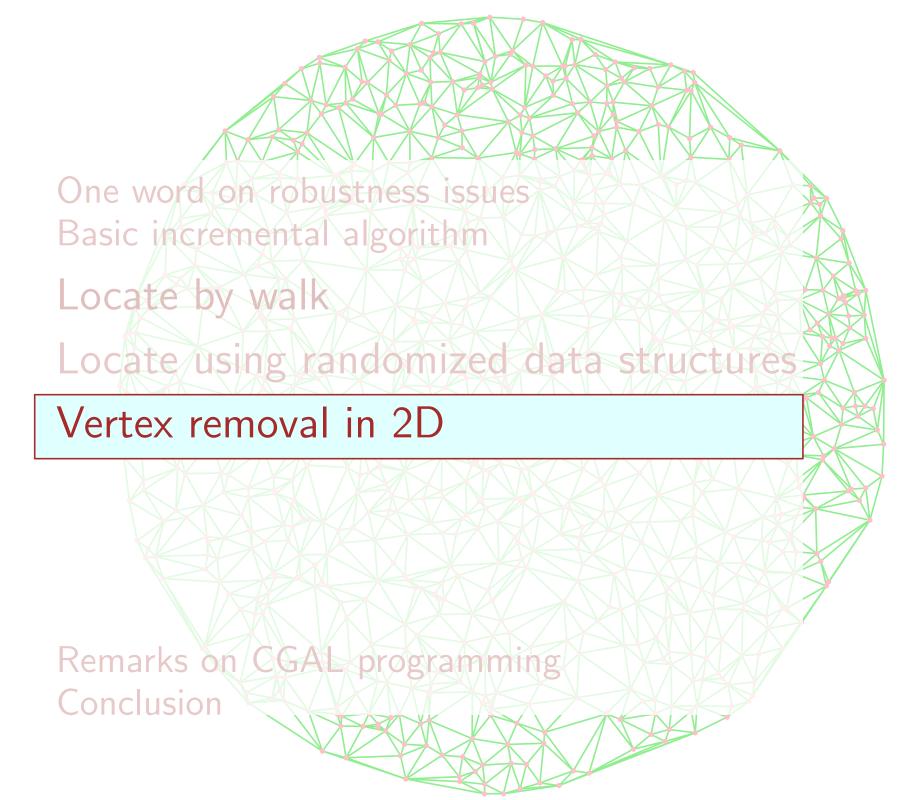
Data structures to locate

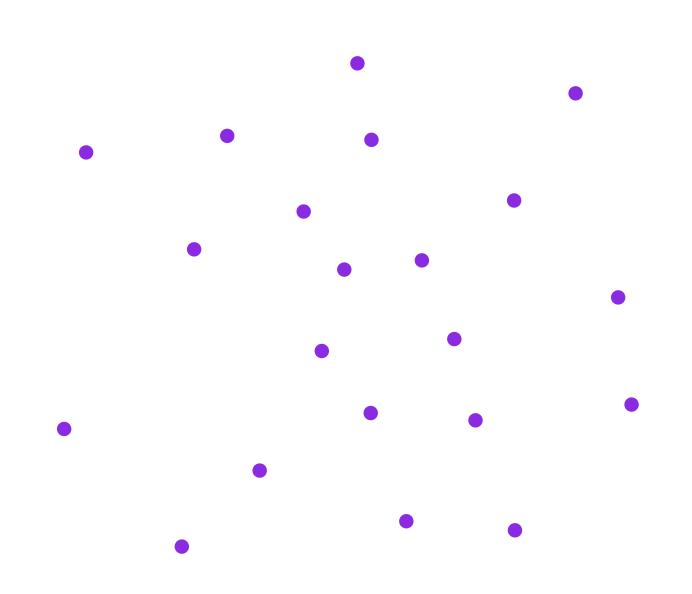
Construction of Delaunay 10 M random points

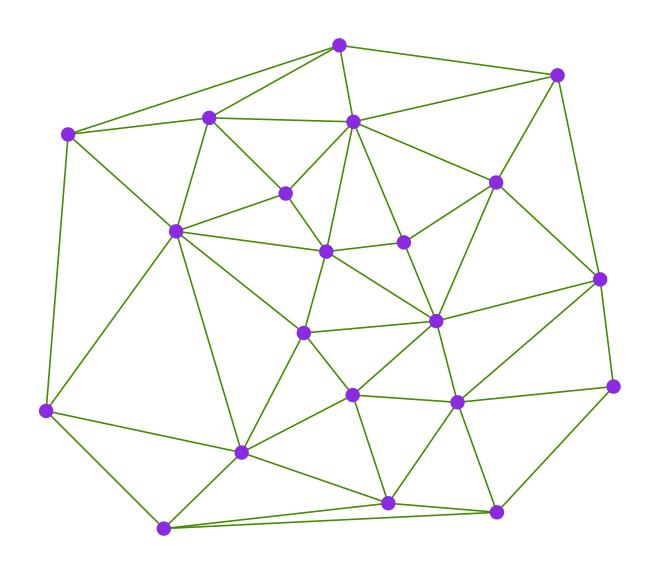
Delaunay tree ~ 10 mn (estimate)

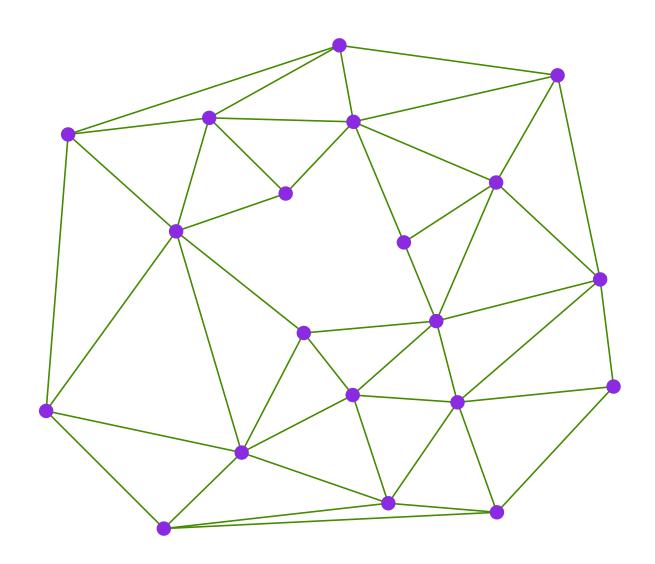
Delaunay hierarchy 90 seconds

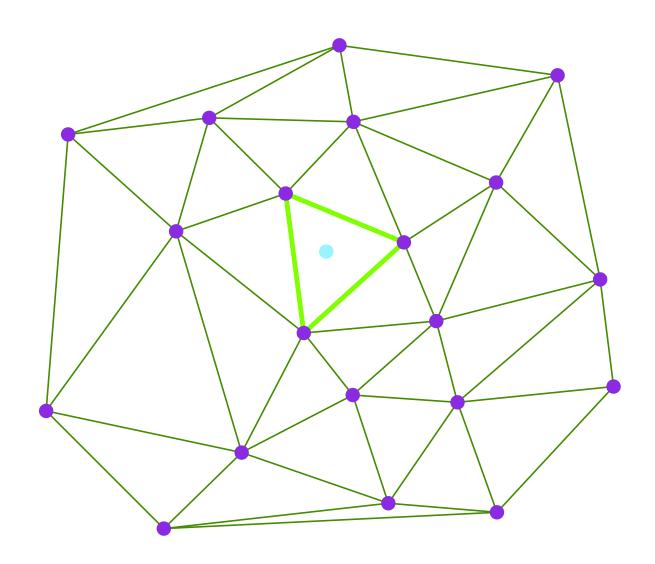
Biased random order 10.6 seconds

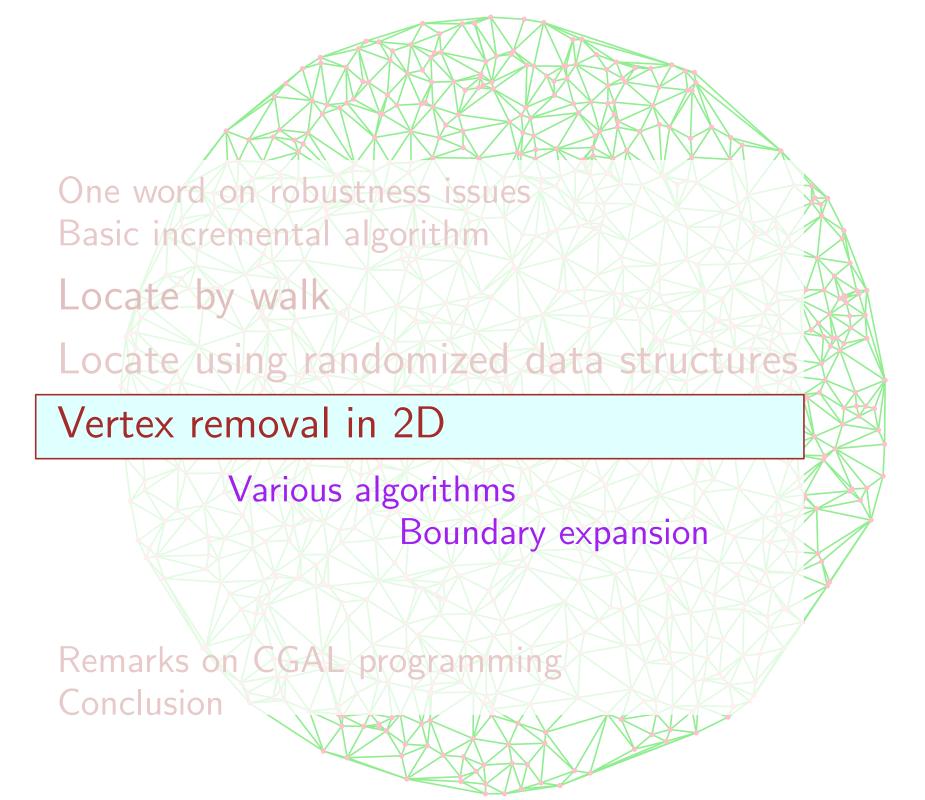




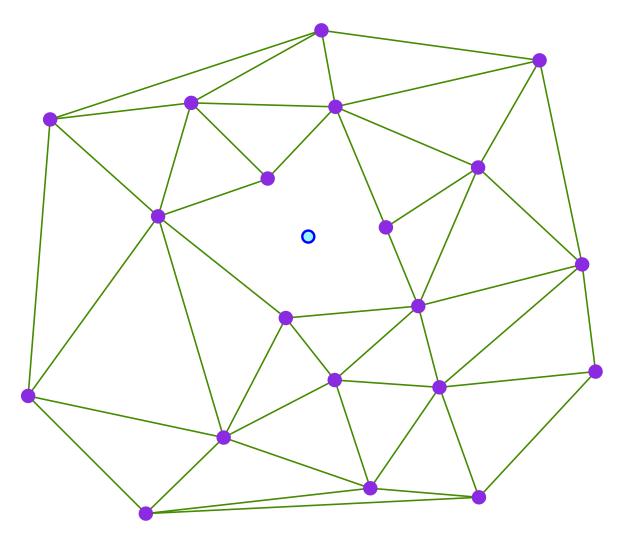




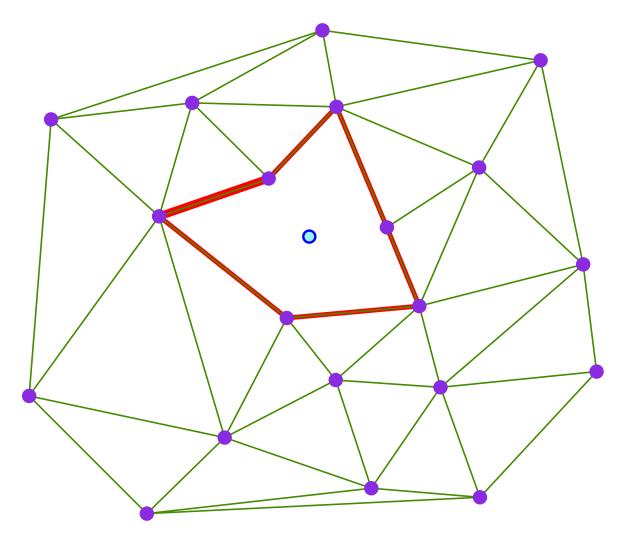




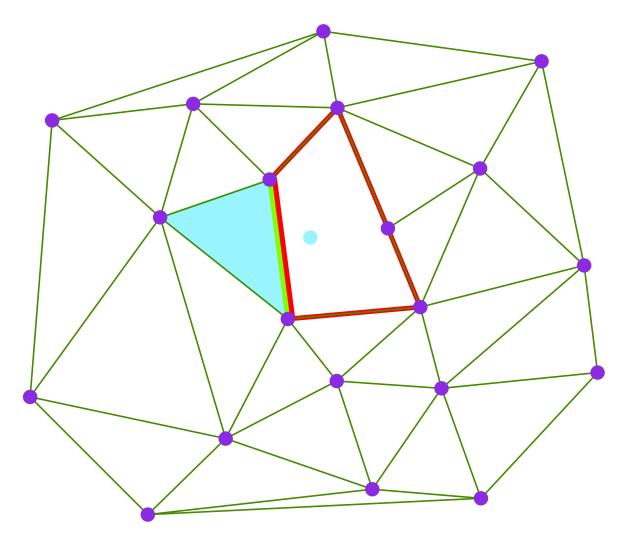




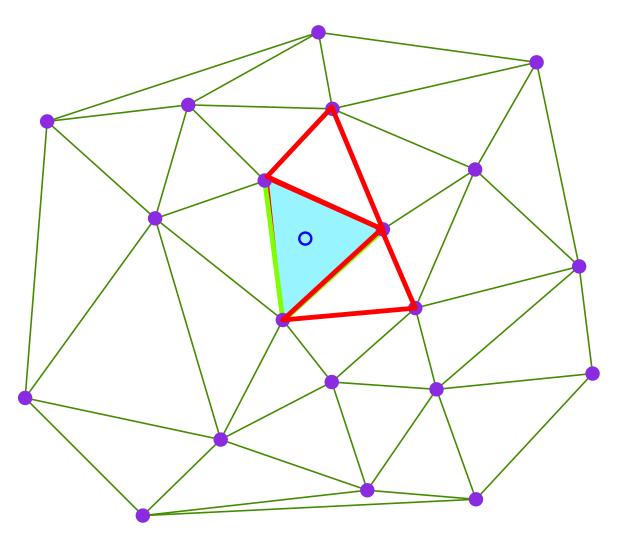






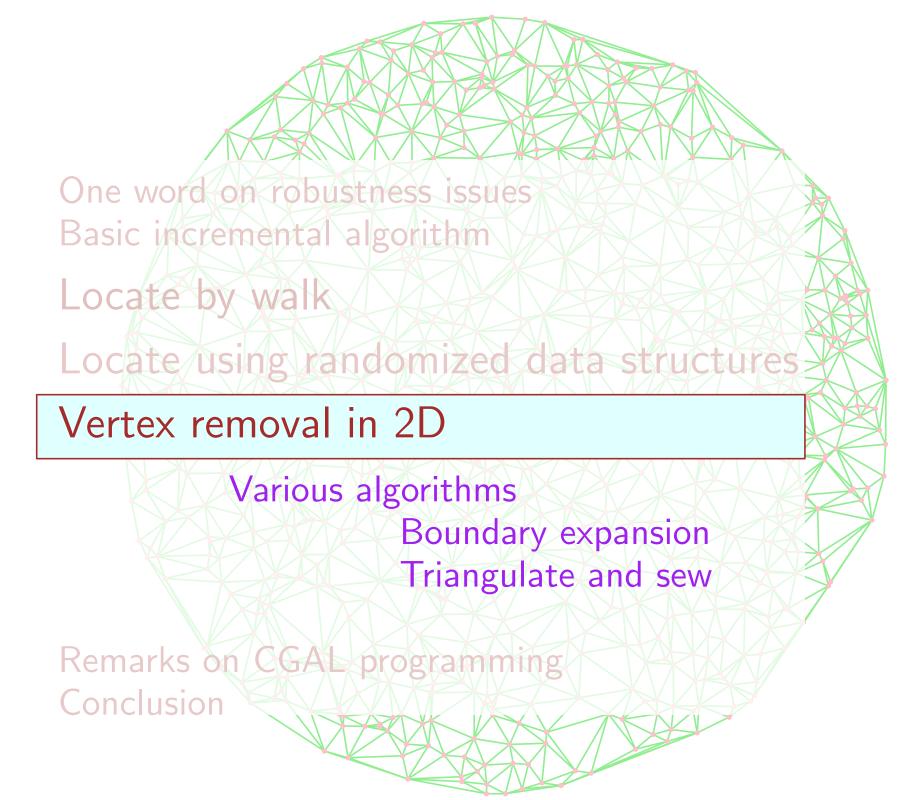




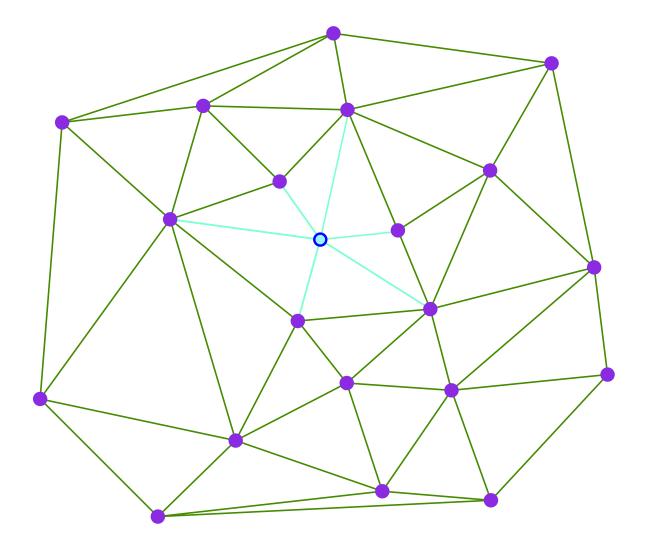




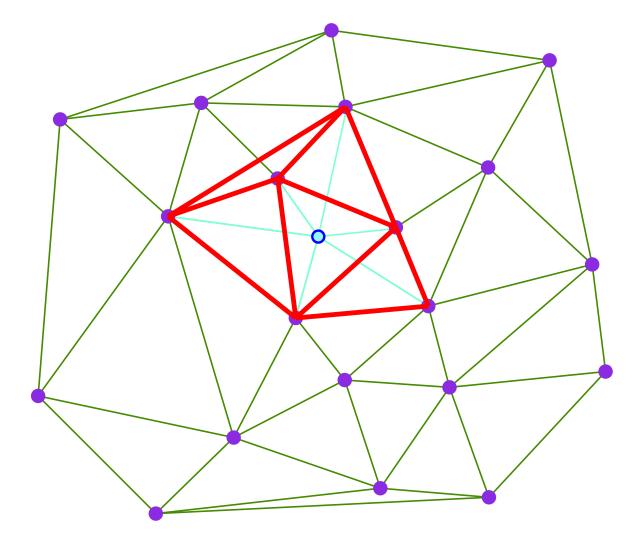




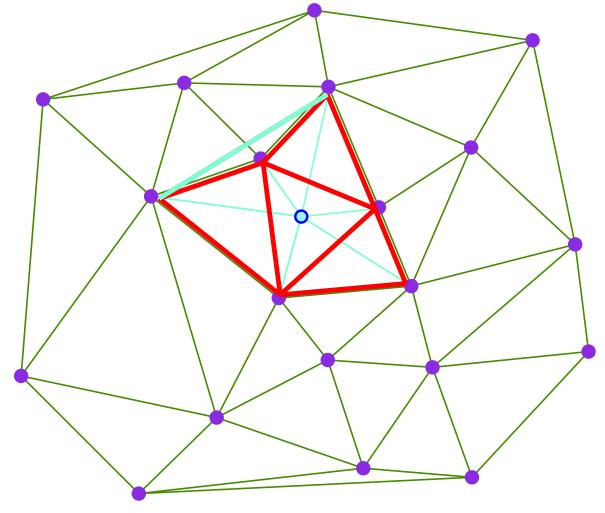
current 3D implementation



current 3D implementation

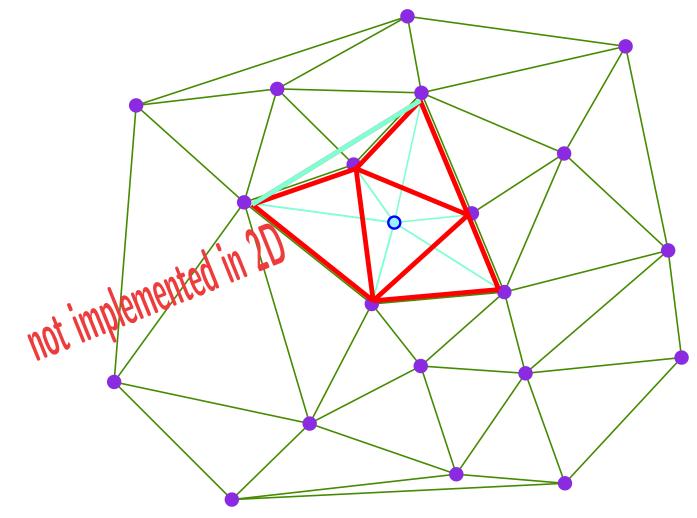


current 3D implementation

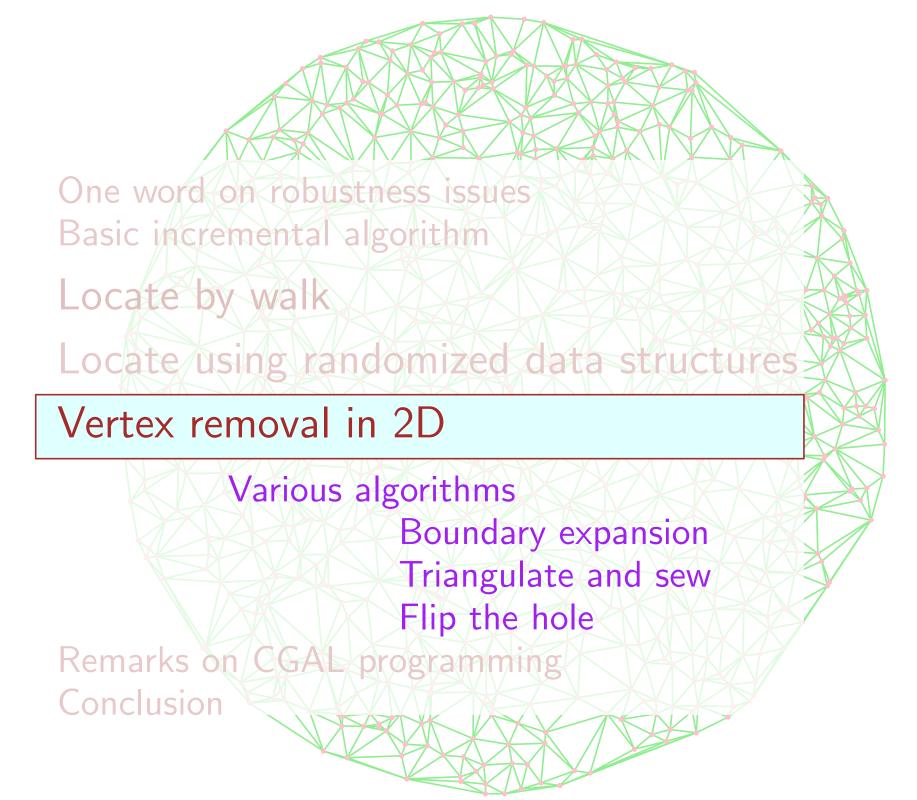


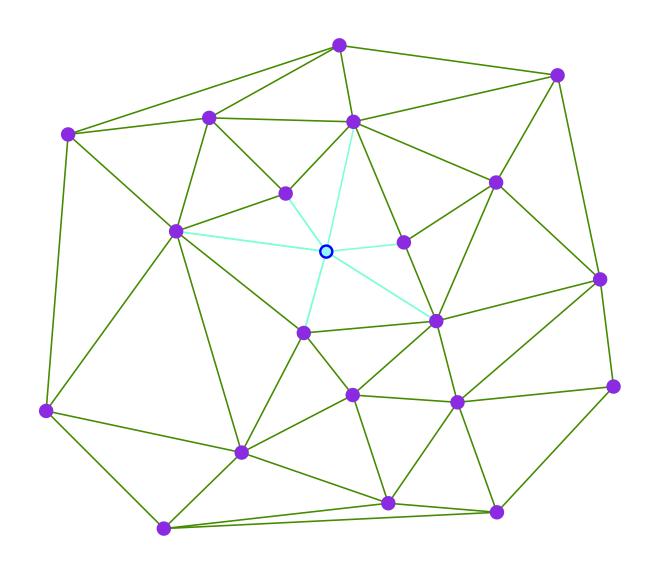
delete extra triangles and sew

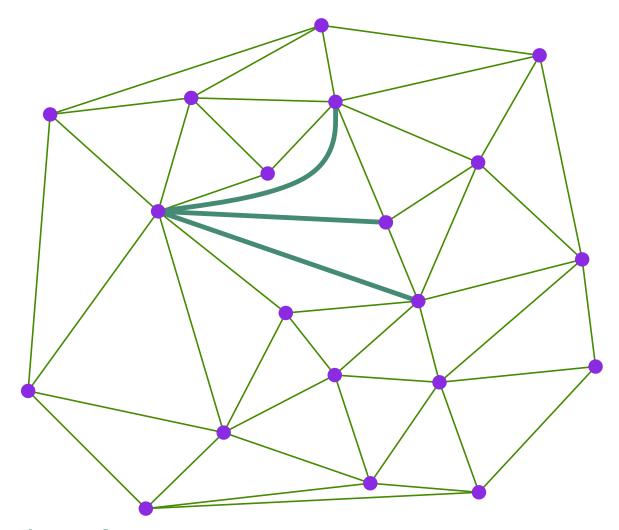




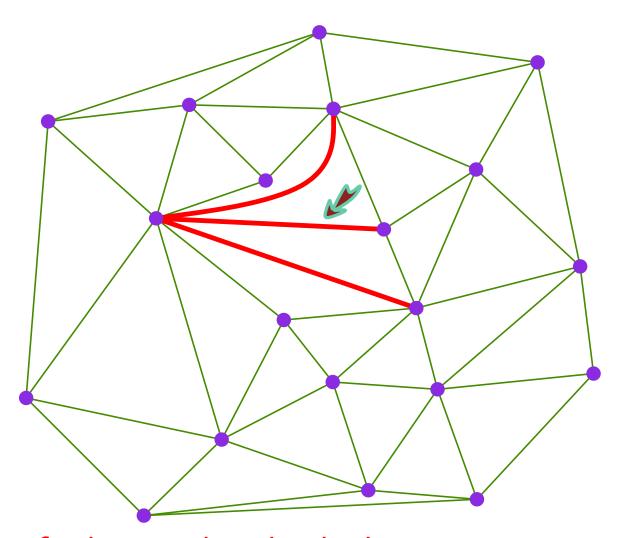
delete extra triangles and sew



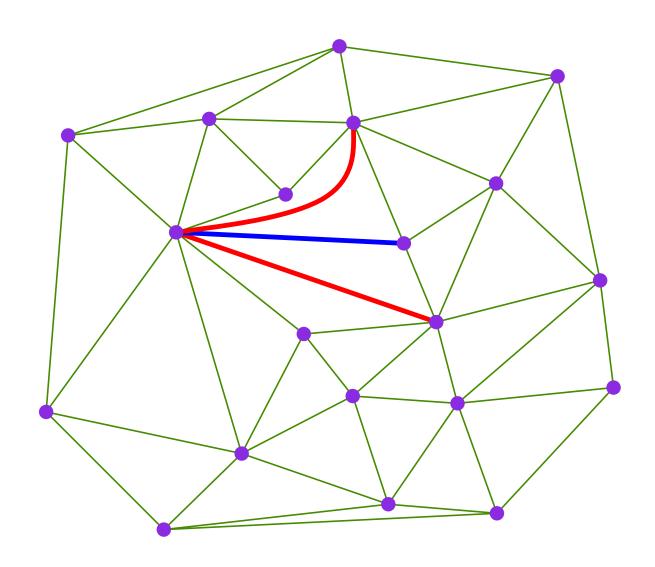


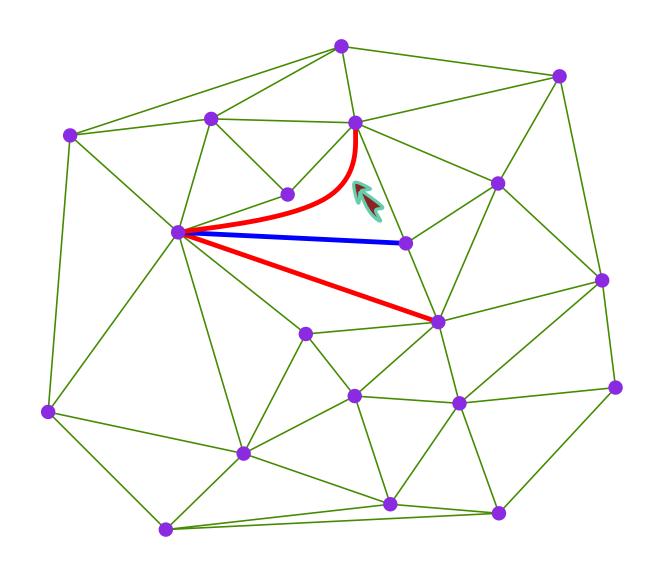


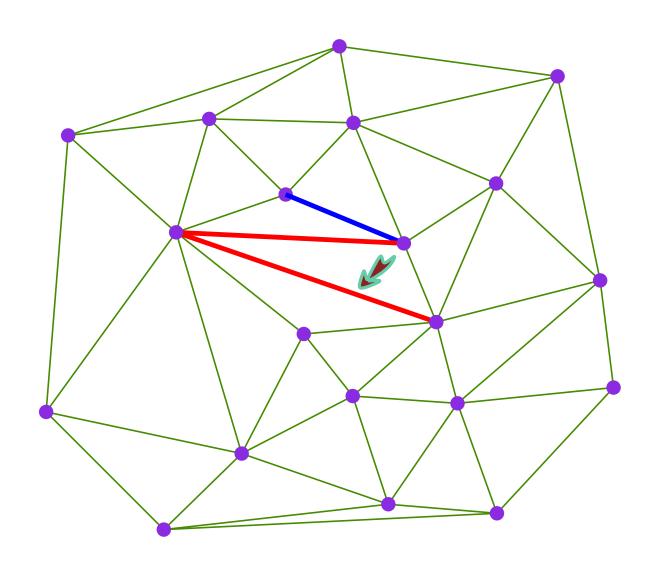
triangulate from any vertex

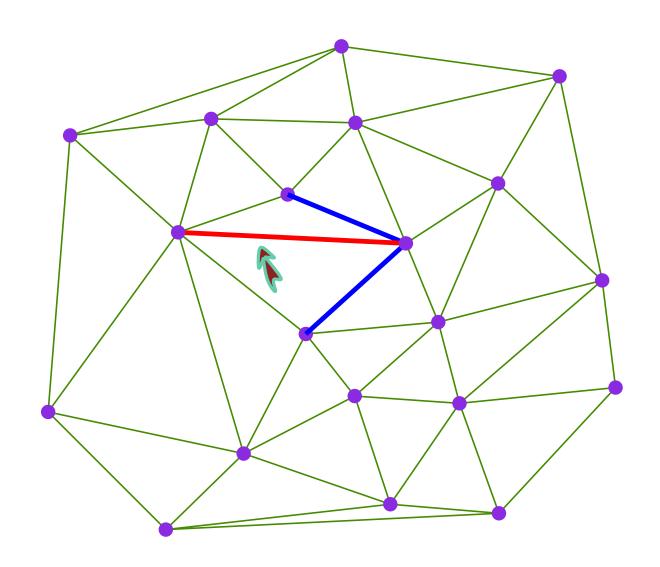


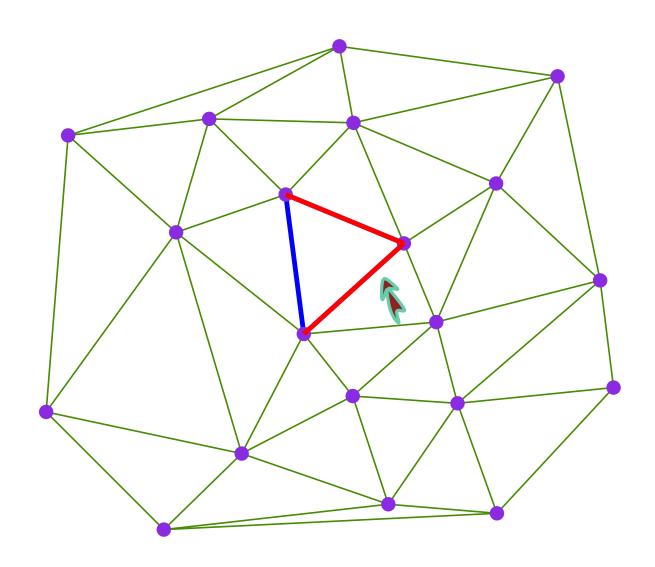
queue of edges to be checked

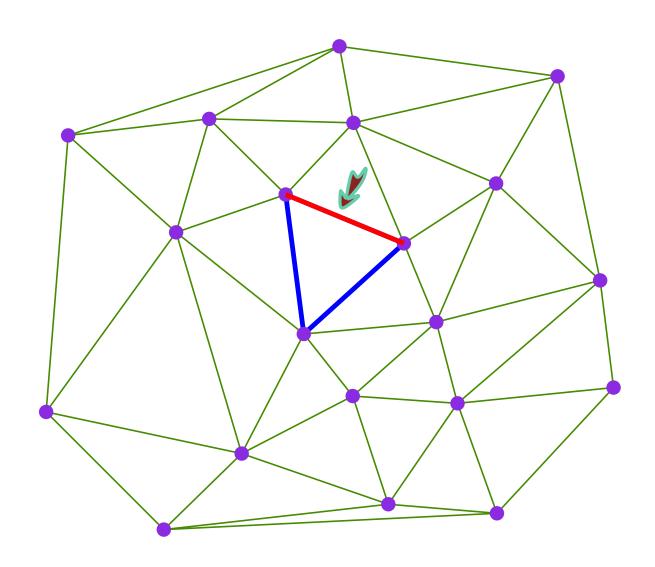


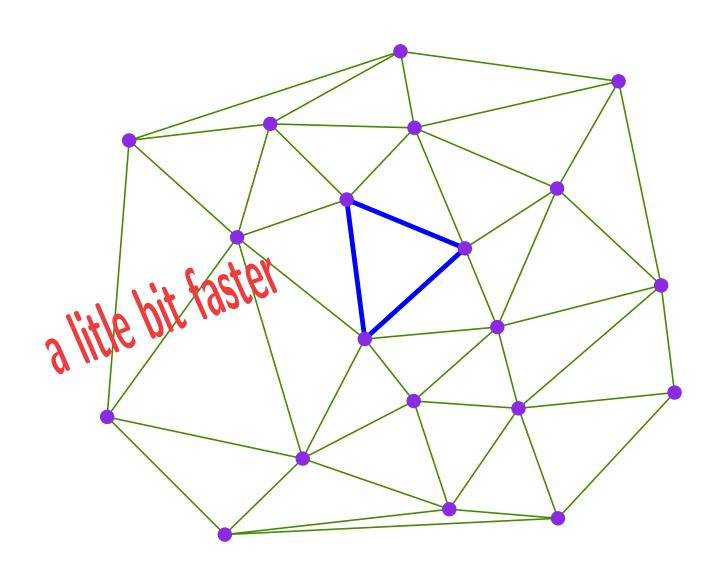


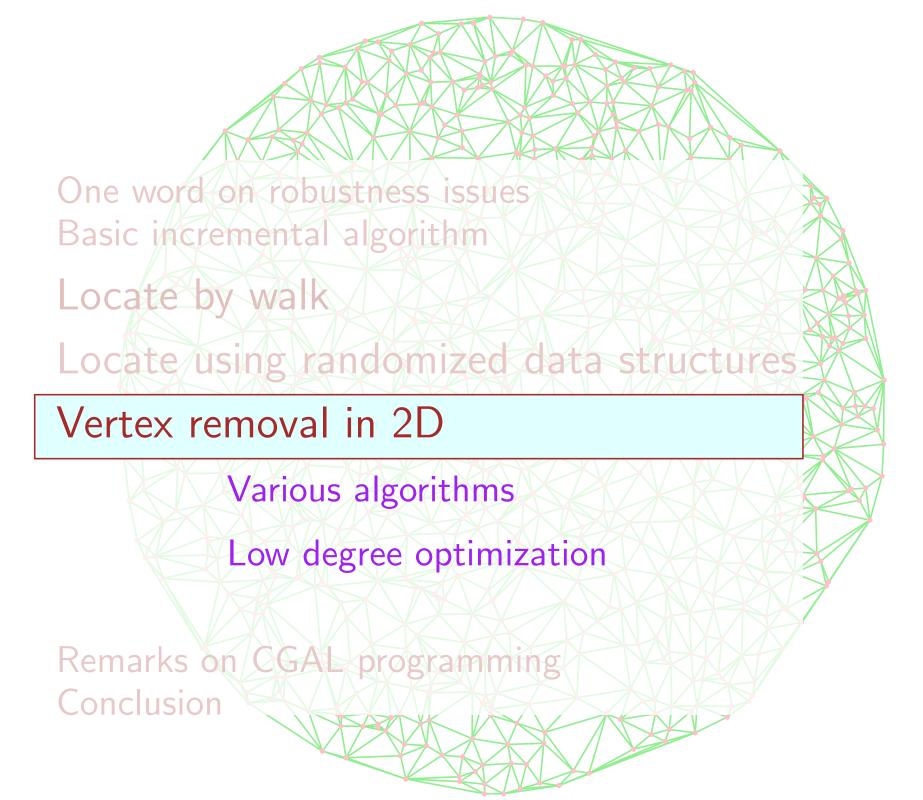




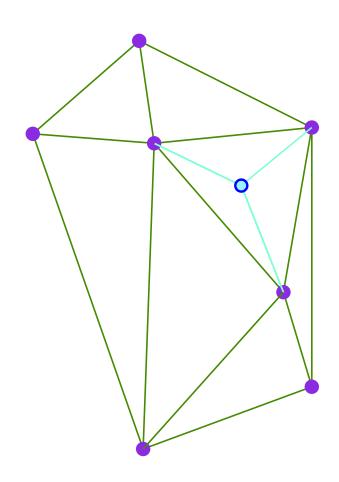




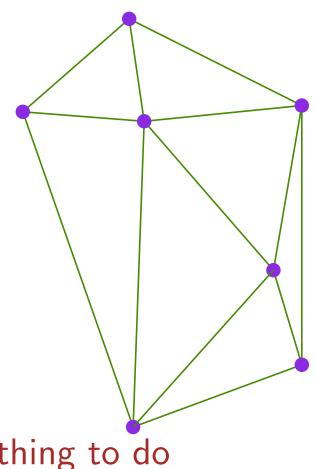




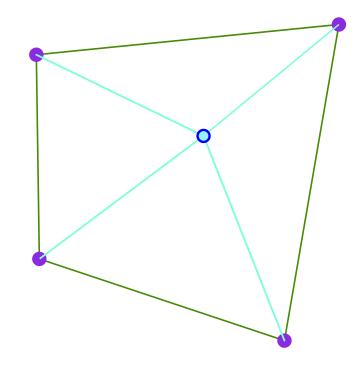
degree 3

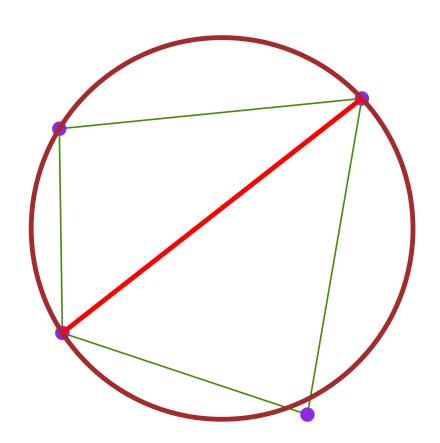


degree 3

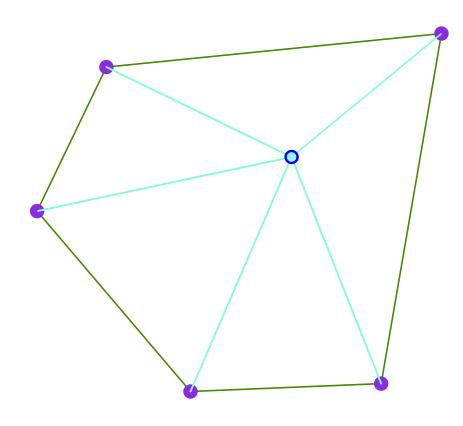


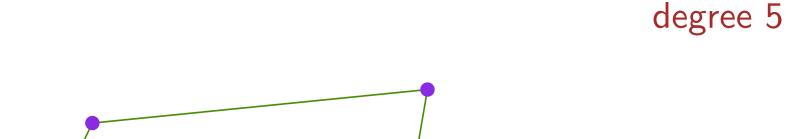
almost nothing to do





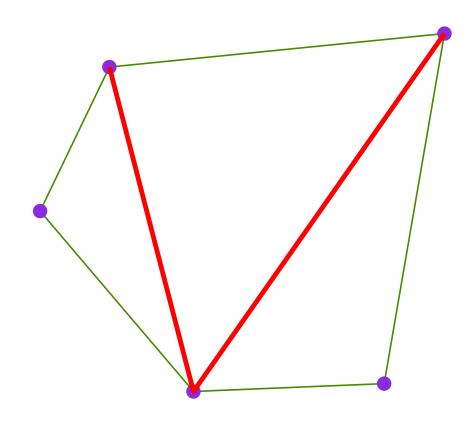
just one incircle test to decide



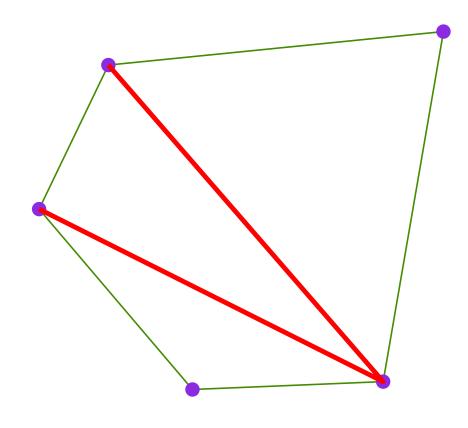




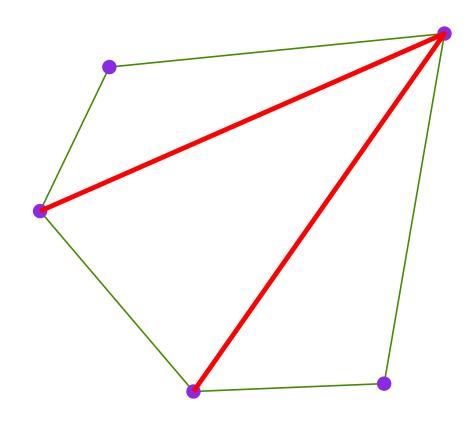




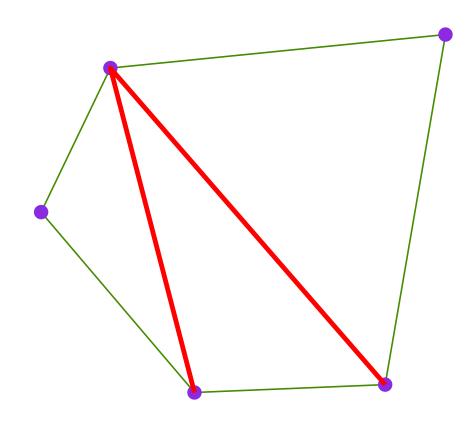




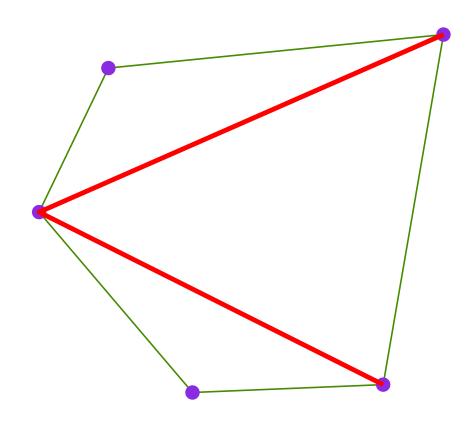
degree 5



degree 5

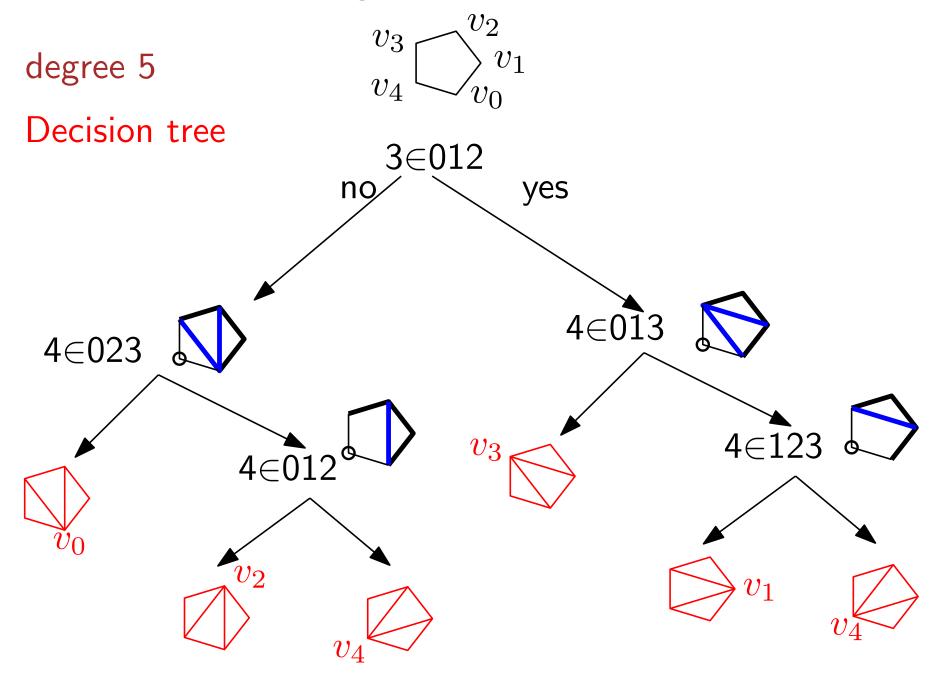


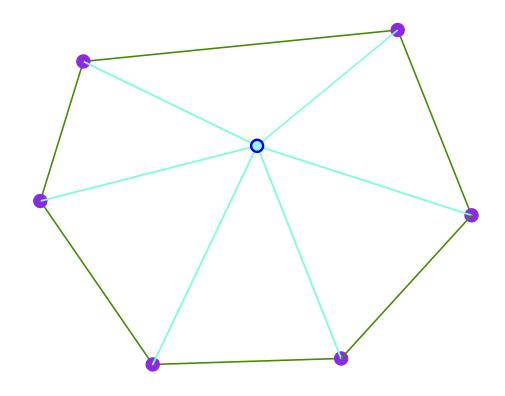
degree 5



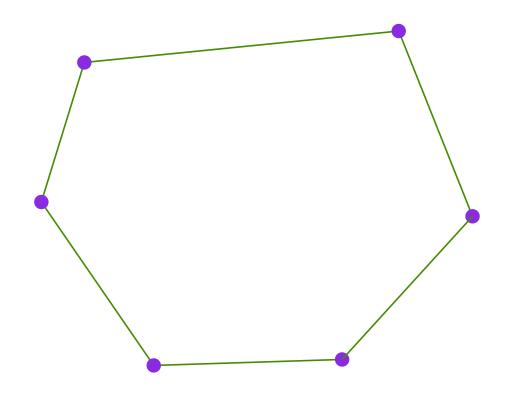
degree 5

Decision tree

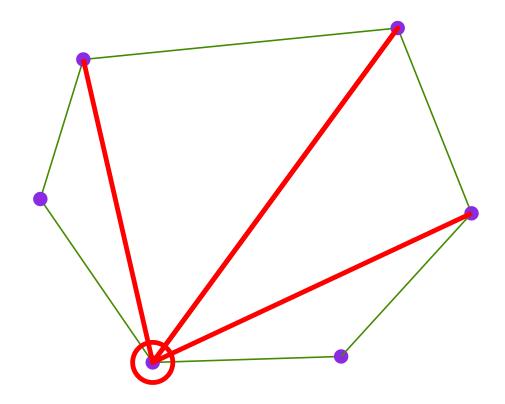




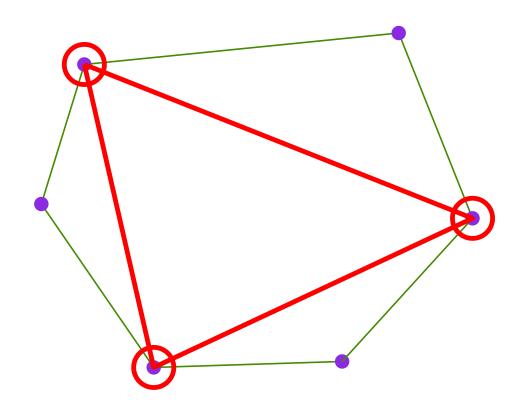




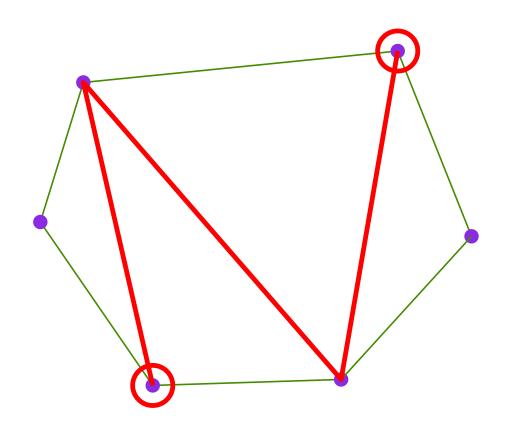
degree 6



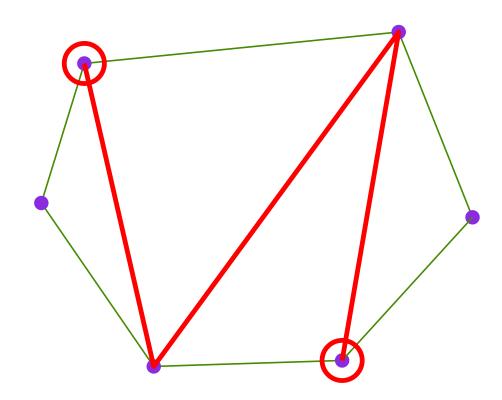
degree 6



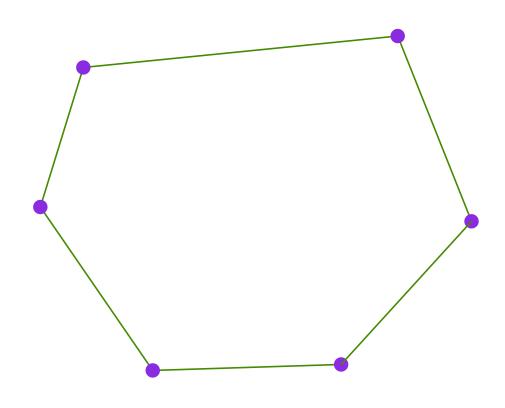
degree 6



degree 6

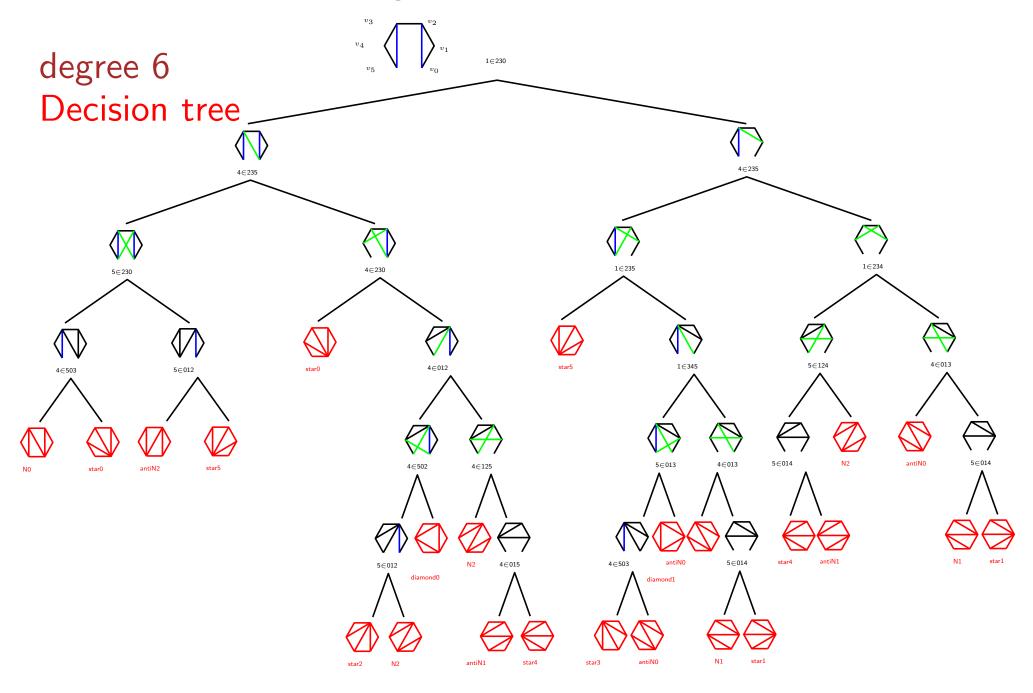


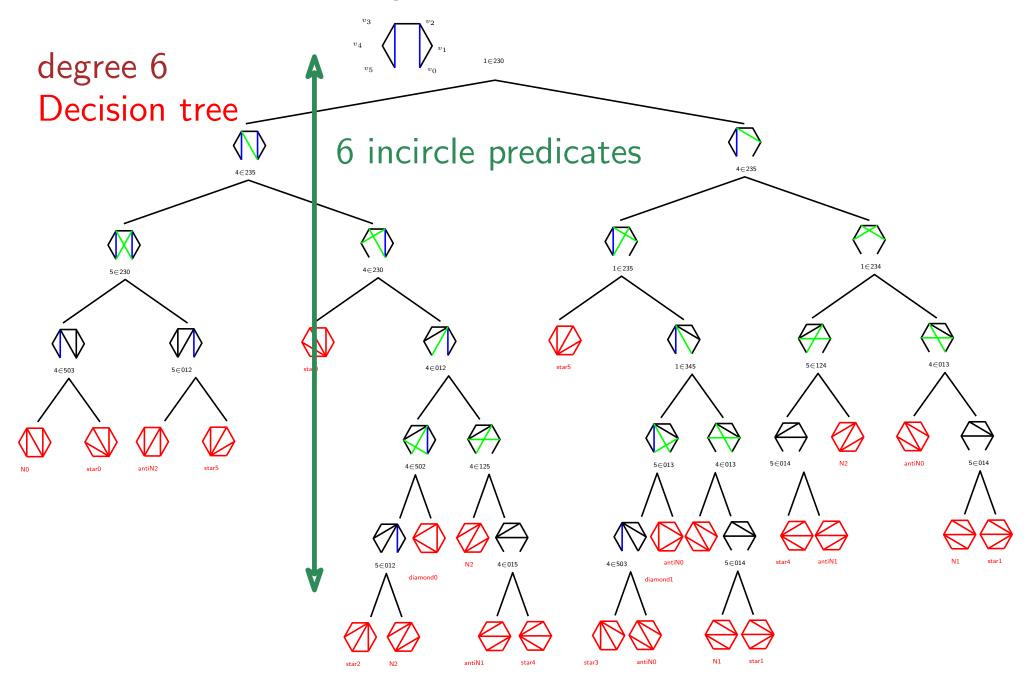


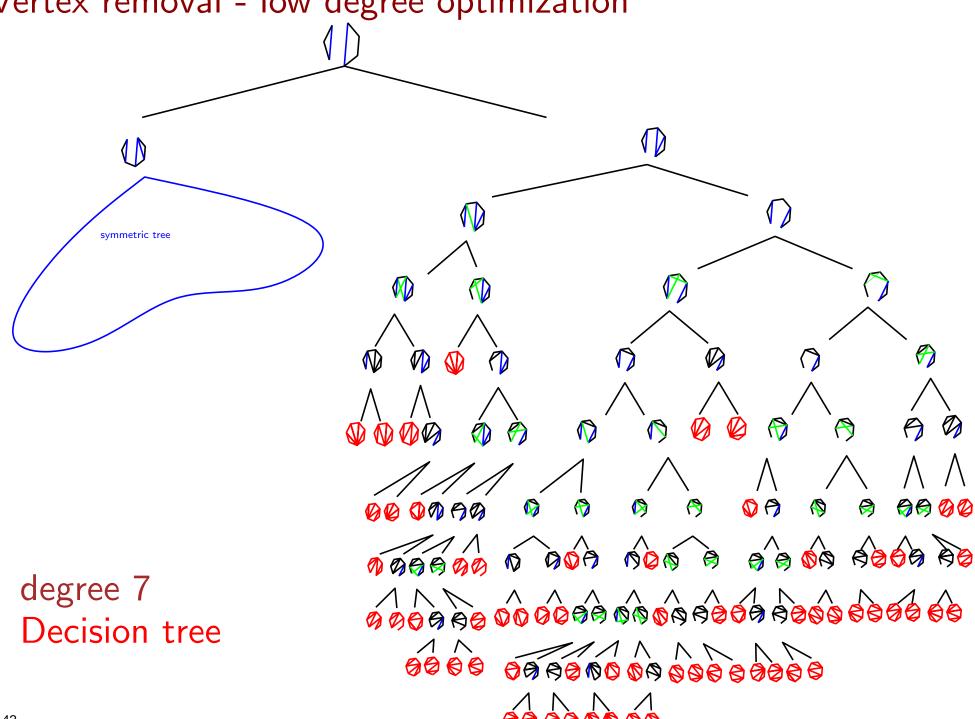


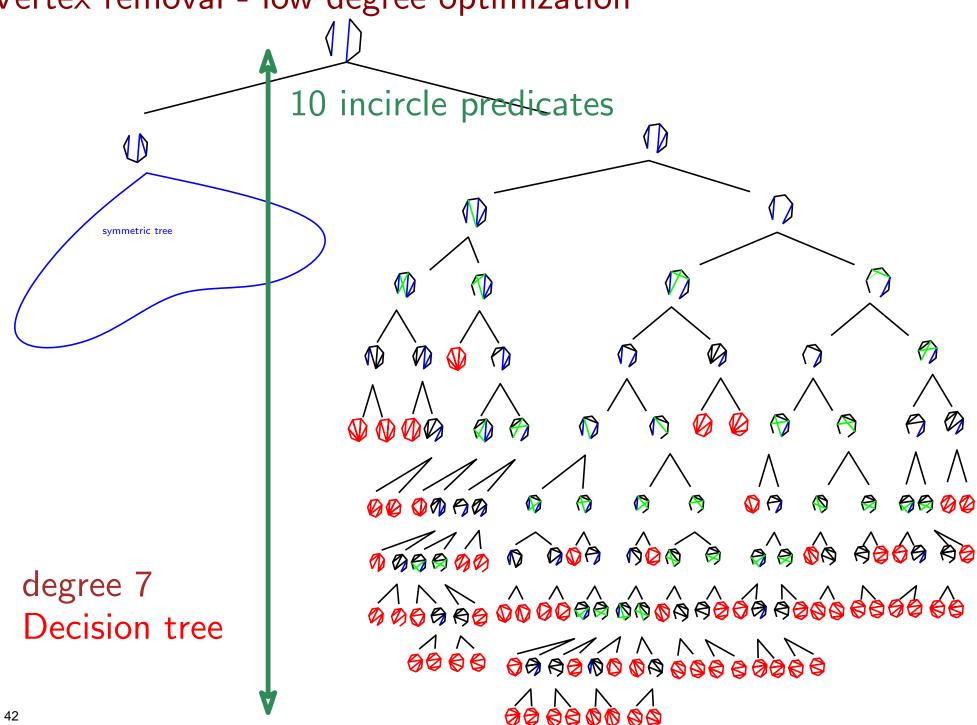
14 results

degree 6
Decision tree









degree	3	4	5	6	7	8*	9
# results	1	2	5	14	42	132	429
# leaves	1	2	6	24	130	<i>≃500</i>	
$\lceil \log_2 \sharp results \rceil$	0	1	3	4	6	8	9
tree height	0	1	3	6	10	\simeq 14	
# lines of code	30	40	90	280	700	<i>≃2500</i>	

^{*} not implemented. The sizes of the tree and the code are estimated

Remarks on implementation

limited memory allocation, use old faces "in place"

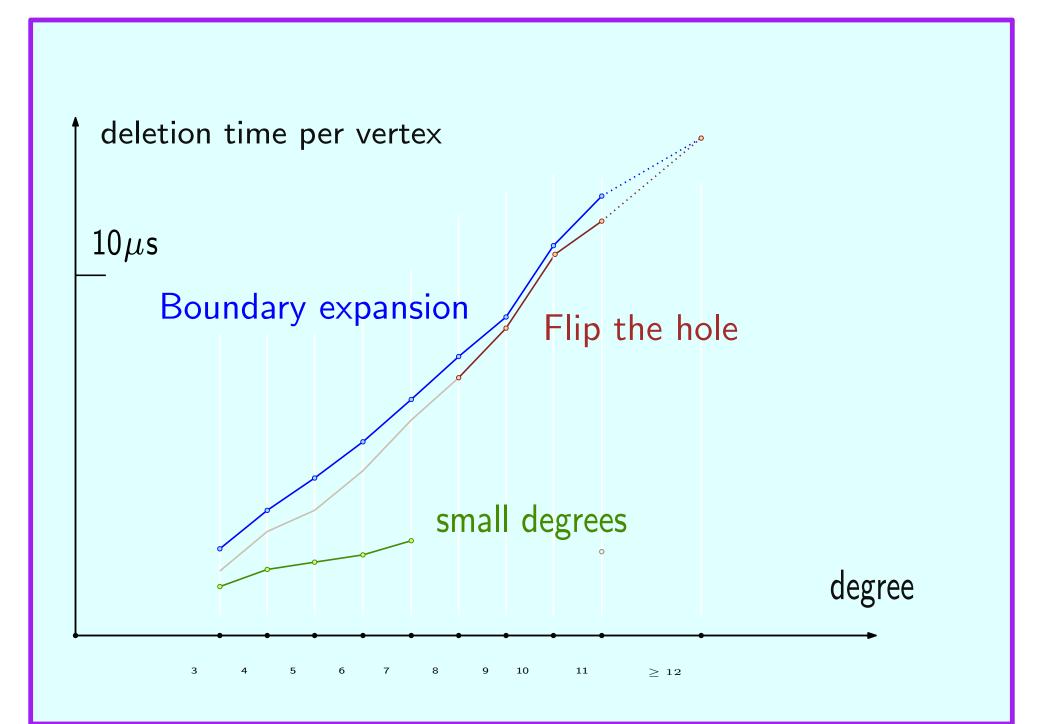
re-use as many neighbor links as possible

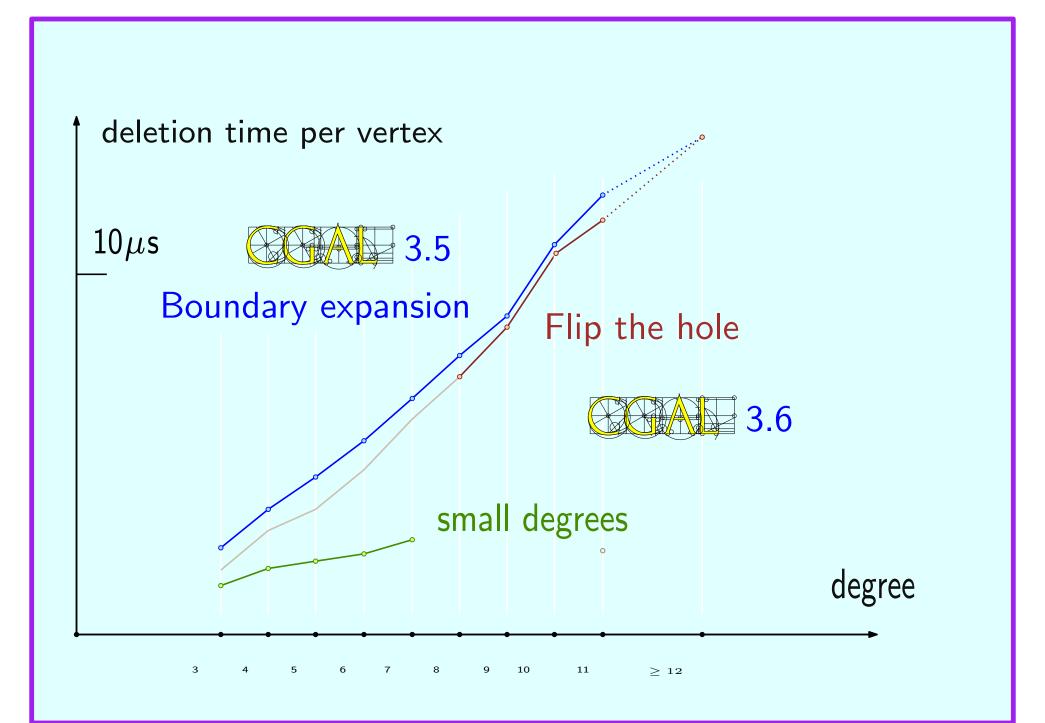
Remarks on implementation

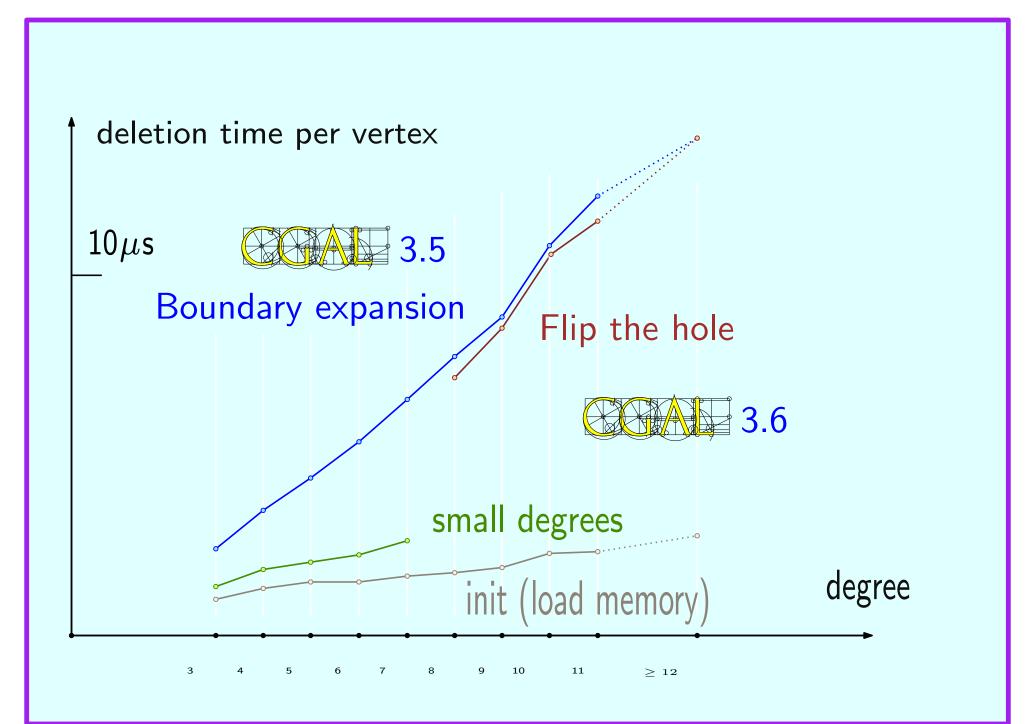
limited memory allocation, use old faces "in place"

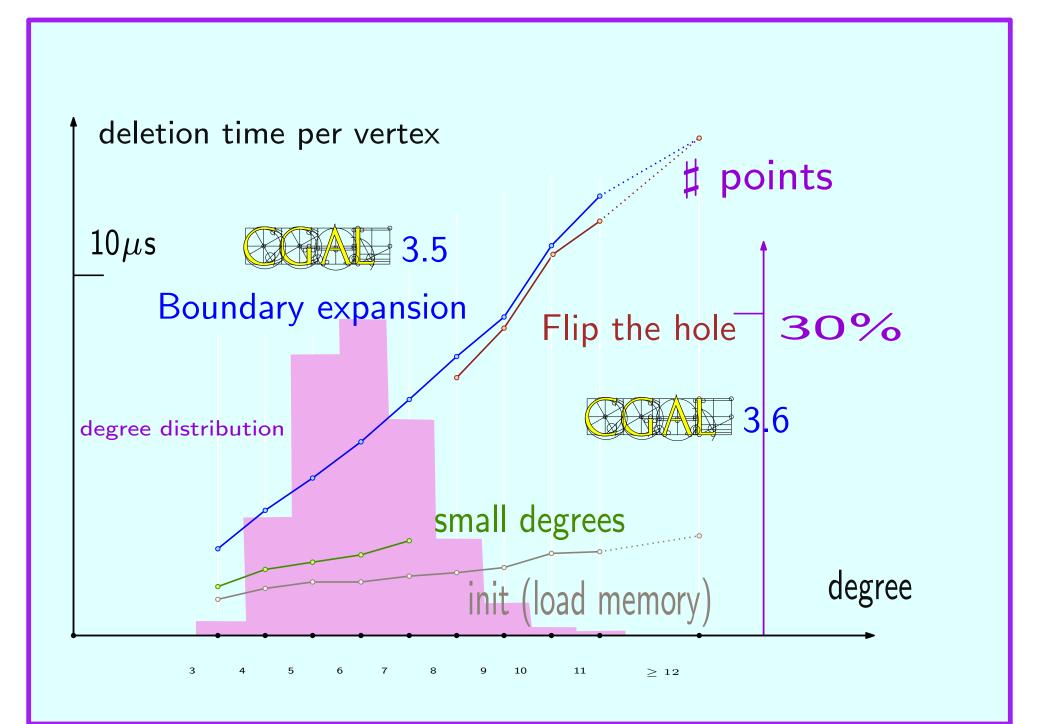
re-use as many neighbor links as possible

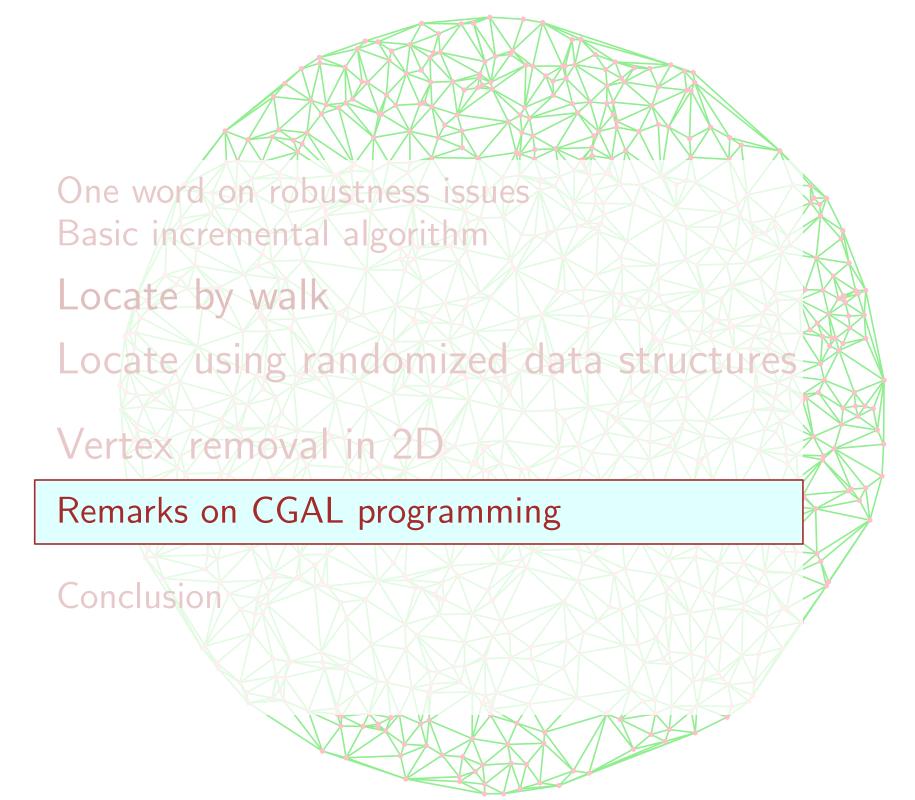
tree implementation













Algorithmic choices

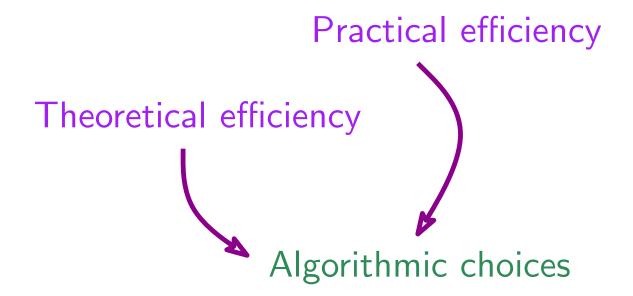


Theoretical efficiency

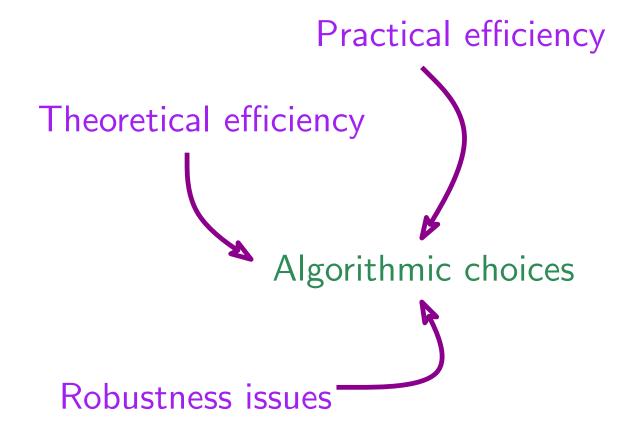


Algorithmic choices

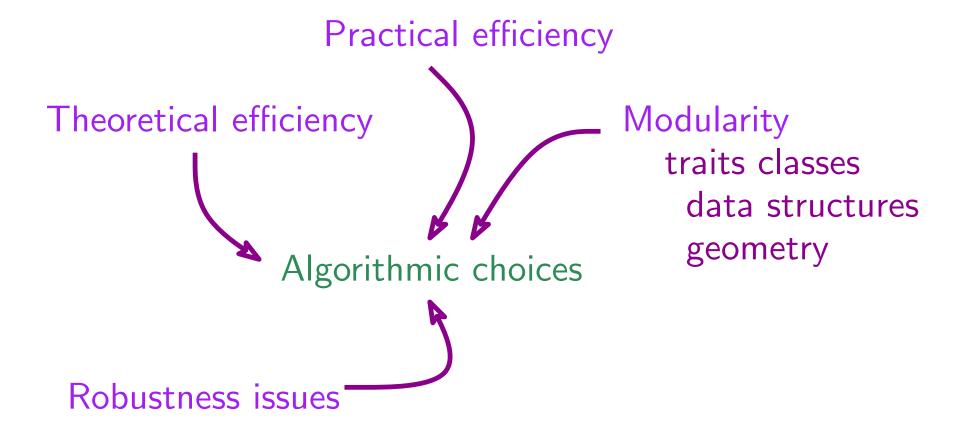






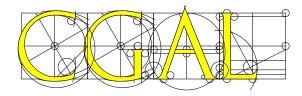


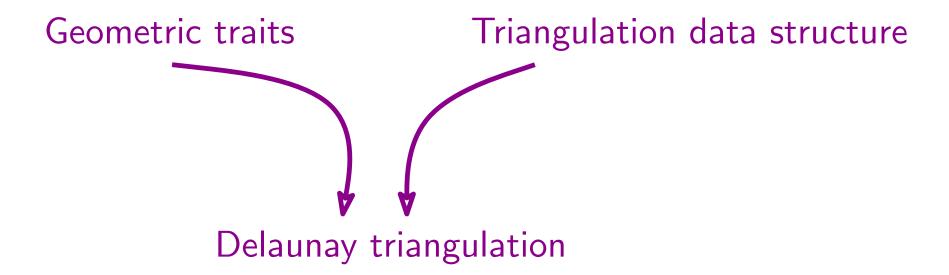




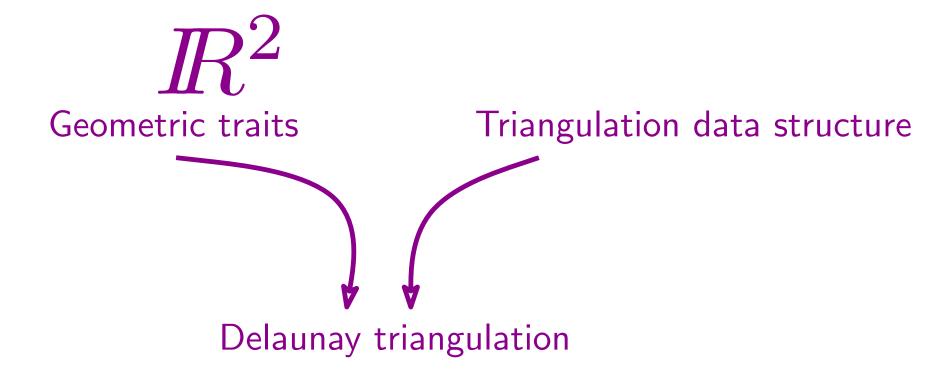


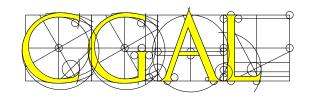
Practical efficiency Theoretical efficiency Modularity traits classes data structures geometry Algorithmic choices Robustness issues Minimal requirements e.g. do not use strange predicates



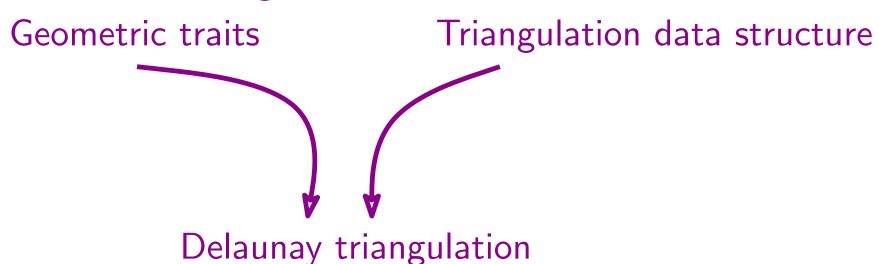


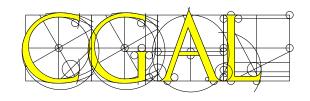




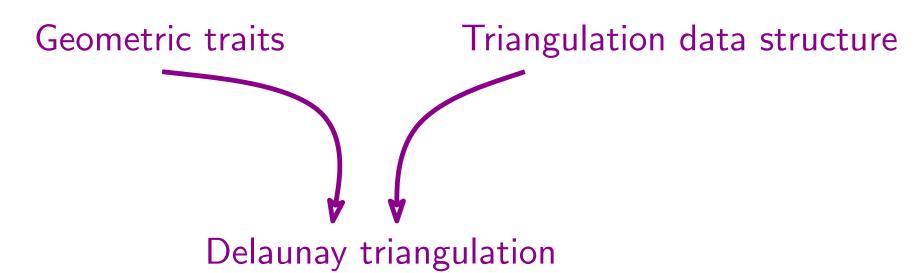


$I\!\!R^3$ projection

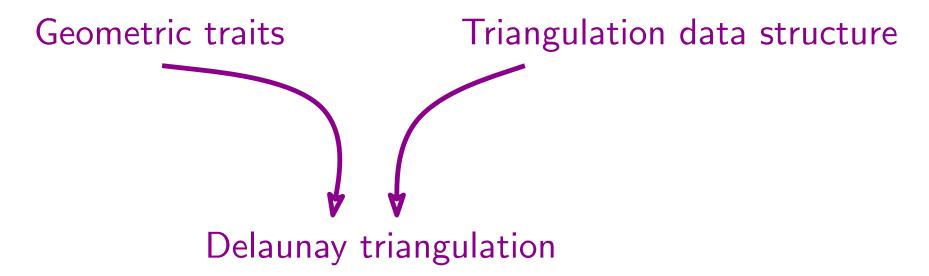




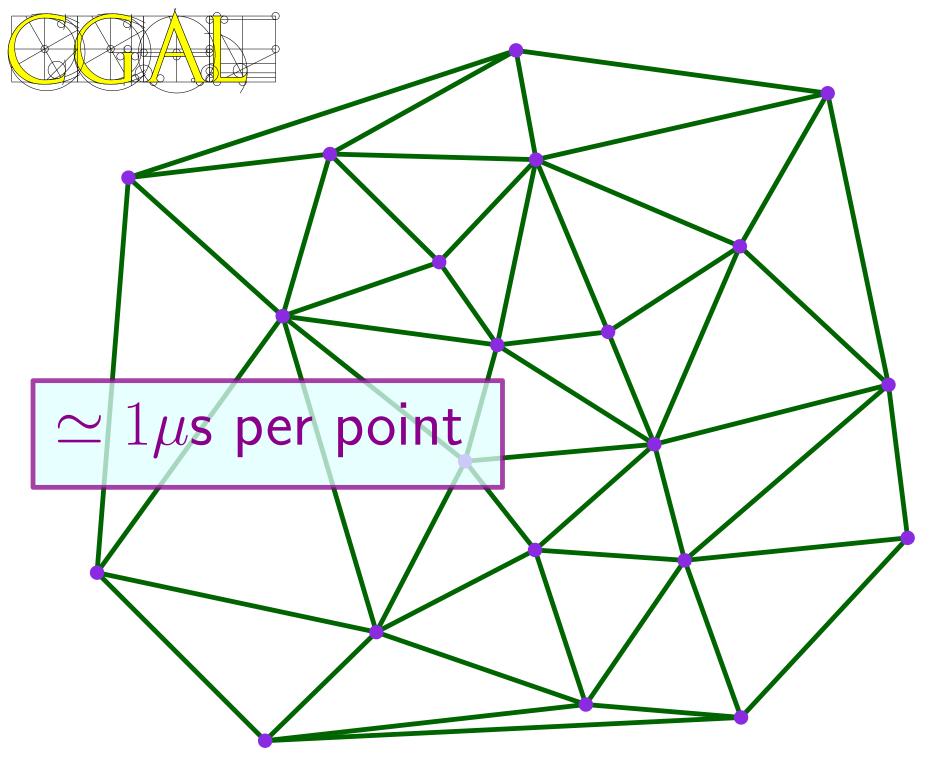
$I\!\!R^2$ other metric

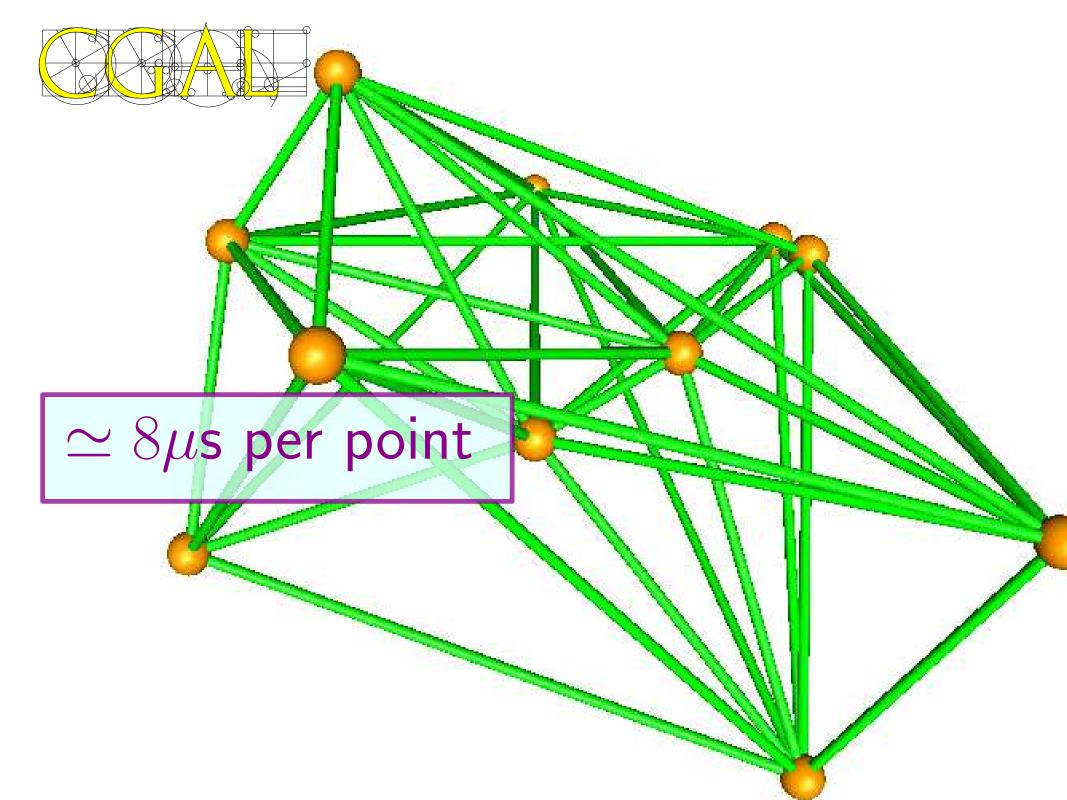


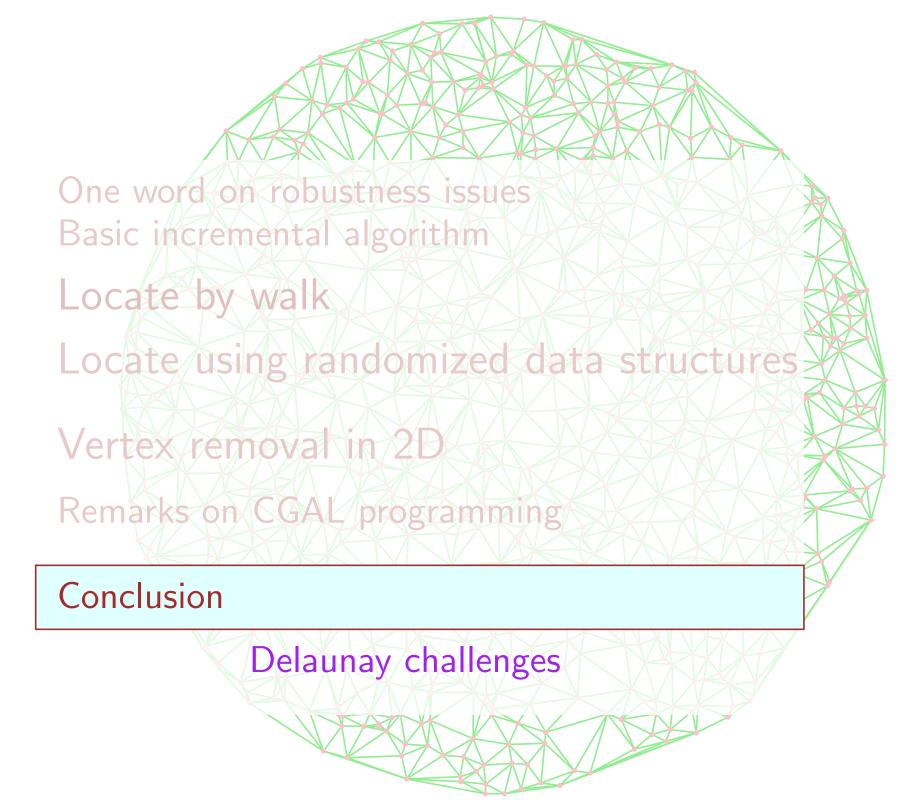




Periodic triangulation







Practical vs worst case size of Delaunay 3D

Practical vs worst case size of Delaunay 3D

Known results

 $\Theta(n^2)$ worst case

 $\Theta(n)$ random in ball

 $\Omega(n)O(n\log n)$ random on polyhedron

 $O(n \log n)$ good sample of smooth generic surface

 $\Theta(n \log n)$ random on cylinder

Practical vs worst case size of Delaunay 3D

Known results

 $\Theta(n^2)$ worst case

Find good models of practical data

 $\Theta(n)$ random in ball

(Smooth analysis)

 $\Omega(n)O(n\log n)$ random on polyhedron

 $O(n \log n)$ good sample of smooth generic surface

 $\Theta(n \log n)$ random on cylinder

Practical vs worst case size of Delaunay 3D

Better algorithm for 3D deletion

10 μ s to insert

100 μ s to delete

Practical vs worst case size of Delaunay 3D

Better algorithm for 3D deletion

One billion points

Needs memory efficient algorithms

Cache effects are already important

