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## Calcul formel pour la combinatoire

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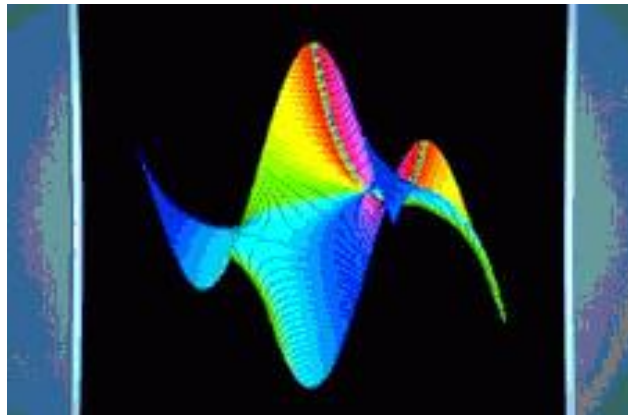
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# Computer algebra for Combinatorics

Alin Bostan & Bruno Salvy



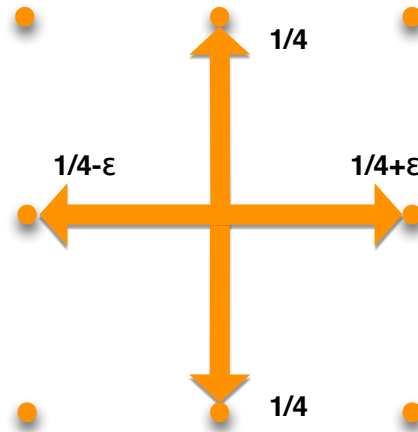
Algorithms Project, INRIA

ALEA 2012

# INTRODUCTION

## 1. Examples

# From the SIAM 100-Digit Challenge



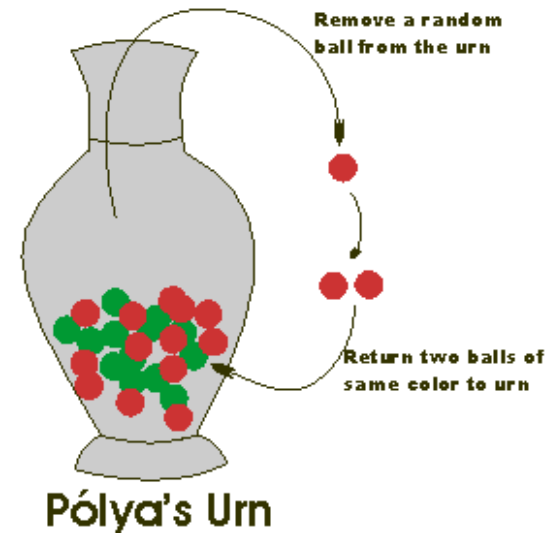
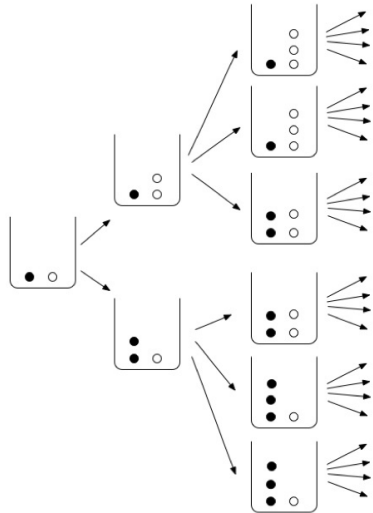
## Problem 6

*A flea starts at  $(0,0)$  on the infinite two-dimensional integer lattice and executes a biased random walk: At each step it hops north or south with probability  $1/4$ , east with probability  $1/4 + \epsilon$ , and west with probability  $1/4 - \epsilon$ . The probability that the flea returns to  $(0,0)$  sometime during its wanderings is  $1/2$ . What is  $\epsilon$ ?*

► Computer algebra [conjectures](#) and [proves](#)

$$p(\epsilon) = 1 - \sqrt{\frac{A}{2}} \cdot {}_2F_1 \left( \begin{matrix} \frac{1}{2}, \frac{1}{2} \\ 1 \end{matrix} \middle| \frac{2\sqrt{1-16\epsilon^2}}{A} \right)^{-1}, \quad \text{with } A = 1 + 8\epsilon^2 + \sqrt{1-16\epsilon^2}.$$

# Algebraic balanced urns



## Theorem [M.FI11]

The balanced urns class  $\begin{pmatrix} 2\alpha & \beta \\ \alpha & \alpha + \beta \end{pmatrix}$ , with  $\alpha > 0$ ,  $\beta \geq 0$ , has an **algebraic** bivariate generating function.

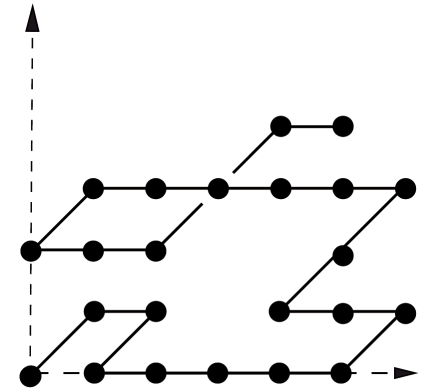
- ▶ Computer algebra **conjectures** and **proves** larger classes of algebraic balanced urns.
- ▶ More in Basile Morcrette's talk!

# Gessel's conjecture

- **Gessel walks**: walks in  $\mathbb{N}^2$  using only steps in  $\mathcal{S} = \{\nearrow, \swarrow, \leftarrow, \rightarrow\}$
- $g(i, j, n) =$  number of **walks** from  $(0, 0)$  to  $(i, j)$  with  $n$  steps in  $\mathcal{S}$

**Question:** Nature of the generating function

$$G(x, y, t) = \sum_{i, j, n=0}^{\infty} g(i, j, n) x^i y^j t^n \in \mathbb{Q}[[x, y, t]]$$



► Computer algebra **conjectures** and **proves**:

**Theorem** [B. & Kauers 2010]  $G(x, y, t)$  is an **algebraic function**<sup>†</sup> and

$$G(1, 1, t) = \frac{1}{2t} \cdot {}_2F_1 \left( \begin{matrix} -1/12 & 1/4 \\ 2/3 \end{matrix} \middle| -\frac{64t(4t+1)^2}{(4t-1)^4} \right) - \frac{1}{2t}.$$













► A simpler variant as **an exercise tomorrow**.

---

<sup>†</sup>Minimal polynomial  $P(x, y, t, G(x, y, t)) = 0$  has  $> 10^{11}$  monomials;  $\approx 30\text{Gb}$  (!)

## Inverse moment problem for walk sequences [B., Flajolet & Penson 2011]

**Question:** Given  $(f_n)$ , find  $I \subseteq \mathbb{R}$  and  $w : I \rightarrow \mathbb{R}$  s.t.  $f_n = \int_I w(t) t^n dt$  ( $n \geq 0$ ).

Step set and walks sequence		GF	Measure $(w(t))$ ;	$t$ :
A126087	 (1, 1, 3, 5, 15, 29, 87)	$\frac{2z - 1 + \sqrt{1 - 8z^2}}{2z(1 - 3z)}$	$\frac{1}{2\pi} \frac{\sqrt{8 - t^2}}{3 - t}$	$[-2\sqrt{2}, 2\sqrt{2}]$
A128386	 (1, 1, 4, 7, 28, 58, 232, 523, 2092)	$\frac{2z - 1 + \sqrt{1 - 12z^2}}{2z(1 - 4z)}$	$\frac{1}{2\pi} \frac{\sqrt{12 - t^2}}{4 - t}$	$[-2\sqrt{3}, 2\sqrt{3}]$
A151282	 (1, 2, 6, 18, 58, 190, 638)	$\frac{3z - 1 + \sqrt{1 - 2z - 7z^2}}{2z(1 - 4z)}$	$\frac{1}{2\pi} \frac{\sqrt{7 + 2t - t^2}}{4 - t}$	$[1 - 2\sqrt{2}, 1 + 2\sqrt{2}]$
A151292	 (1, 2, 7, 23, 85, 314, 1207, 4682)	$\frac{3z - 1 + \sqrt{1 - 2z - 11z^2}}{2z(1 - 5z)}$	$\frac{1}{2\pi} \frac{\sqrt{11 + 2t - t^2}}{5 - t}$	$[1 - 2\sqrt{3}, 1 + 2\sqrt{3}]$
A129400	 (1, 2, 8, 32, 144, 672, 3264)	$\frac{1 - 2z - \sqrt{1 - 4z - 12z^2}}{8z^2}$	$\frac{1}{8\pi} \sqrt{(t+2)(6-t)}$	$[-2, 6]$
A151318	 (1, 3, 13, 55, 249, 1131, 5253)	$\frac{5z - 1 + \sqrt{1 - 2z - 15z^2}}{4z(1 - 5z)}$	$\frac{1}{4\pi} \sqrt{\frac{3+t}{5-t}}$	$[-3, 5]$
A060899	 (1, 2, 8, 24, 96, 320, 1280, 4480)	$\frac{4z - 1 + \sqrt{1 - 16z^2}}{4z(1 - 4z)}$	$\frac{1}{4\pi} \sqrt{\frac{4+t}{4-t}}$	$[-4, 4]$
A005773	 (1, 2, 5, 13, 35, 96, 267, 750, 2123)	$\frac{3z - 1 + \sqrt{1 - 2z - 3z^2}}{2z(1 - 3z)}$	$\frac{1}{2\pi} \sqrt{\frac{1+t}{3-t}}$	$[-1, 3]$
A001405	 (1, 1, 2, 3, 6, 10, 20, 35, 70, 126)	$\frac{2z - 1 + \sqrt{1 - 4z^2}}{2z(1 - 2z)}$	$\frac{1}{2\pi} \sqrt{\frac{2+t}{2-t}}$	$[-2, 2]$
A151281	 (1, 2, 6, 16, 48, 136, 408, 1184)	$\frac{4z - 1 + \sqrt{1 - 8z^2}}{4z(1 - 3z)}$	$\frac{1}{4\pi} \frac{\sqrt{8 - t^2}}{3 - t}$	$[-2\sqrt{2}, 2\sqrt{2}]$
A129637	 (1, 3, 11, 41, 157, 607, 2367, 9277)	$\frac{5z - 1 + \sqrt{1 - 2z - 7z^2}}{4z(1 - 4z)}$	$\frac{1}{4\pi} \frac{\sqrt{7 + 2t - t^2}}{4 - t}$	$[1 - 2\sqrt{2}, 1 + 2\sqrt{2}]$
A151323	 (1, 3, 14, 67, 342, 1790, 9580)	$\frac{\sqrt[4]{\frac{1+2z}{1-6z}} - 1}{2z}$	$\frac{1}{2\sqrt{2}\pi} \sqrt[4]{\frac{2+t}{6-t}}$	$[-2, 6]$

# A SIAM Review combinatorial identity

Problem 87-8, by JOHN W. MOON (University of Alberta).

Show that

$$\sum_{n=1}^{\infty} \frac{56n^2 + 33n - 8}{(n+2)(n+1)} f_n^2 = 1$$

where

$$f_n = \frac{4^{-n}}{n} \binom{2n-2}{n-1} \quad \text{for } n \geq 1.$$

*Background.* A *branch* of a rooted tree  $T_n$  is a maximal subtree that does not contain the root. A branch  $B$  with  $i$  nodes is a *primary* branch of  $T_n$  if  $n/2 \leq i \leq n-1$ ; if  $T_n$  has a primary branch  $B$  with  $i$  nodes, then a branch  $C$  with  $j$  nodes is a *secondary* branch if  $(n-i)/2 \leq j \leq n-1-i$ . For many families  $F$  of rooted trees, the fraction of trees  $T_n$  in  $F$  that have a primary branch tends to 1 as  $n \rightarrow \infty$ . (See A. Meir and J.W. Moon, *On major and minor branches of rooted trees*, *Canad. J. Math.*, 39 (1987) 673-693). It can be shown that the fraction of plane trees  $T_n$  that have a secondary branch tends to a limit  $p$  as  $n \rightarrow \infty$ , where

$$p = 3 - 12 \sum_{n=1}^{\infty} \frac{13n^2 + 5n - 2}{(n+1)(n+2)} f_n^2.$$

If we appeal to the proposed identity then we obtain the more rapidly converging expression

$$p = \frac{3}{14} + \frac{3}{14} \sum_{n=1}^{\infty} \frac{149n+8}{(n+1)(n+2)} f_n^2$$

from which we find that  $p = .59 \dots$

► Computer algebra [conjectures](#) and [proves](#)  $p = \frac{28}{15\pi}$ .



# Monthly (AMM) problems with a combinatorial flavor that can be solved using computer algebra

## Expansion of a Symmetric Determinant

E2297 [1971, 543]. *Proposed by Richard Stanley, Harvard University*

Let  $L(n)$  be the total number of distinct monomials appearing in the expansion of the determinant of an  $n \times n$  symmetric matrix  $A = (a_{ij})$ . For instance,  $L(3) = 5$ . Show that

$$\sum_{n=0}^{\infty} L(n)x^n/n! = (1-x)^{-1/2} \exp(\frac{1}{2}x + \frac{1}{4}x^2),$$

where  $|x| < 1$ , and where we define  $L(0) = 1$ .

## Units of Chains

6342 [1981, 294]. *Proposed by Richard Stanley, Massachusetts Institute of Technology.*

Let  $f(n)$  be the number of nonisomorphic  $n$ -element partially ordered sets  $P$  which do not contain three pairwise incomparable elements. (Equivalently,  $P$  is a union of two chains.) Let

$$F(x) = 1 + \sum_{n \geq 1} f(n)x^n = 1 + x + 2x^2 + 4x^3 + 10x^4 + \dots$$

Show that

$$F(x) = \frac{4}{2 - 2x + \sqrt{1 - 4x} + \sqrt{1 - 4x^2}}.$$

### Noncrossing Trees

E 3170 [1986, 650]. *Proposed by The Howard University Group, Washington, D.C.*

Construct a graph as follows: Put  $n + 1$  labeled vertices around a circle and let the edges be the straight line segments connecting any two vertices. A tree is noncrossing if no two edges intersect except at the vertices. Enumerate the number of noncrossing spanning trees for this graph. For  $n = 1, 2, 3$ , the numbers are 1, 3, 12, respectively.

### An Unexpected Appearance of the Catalan Numbers

**10905** [2001, 871]. *Proposed by Richard P. Stanley, Massachusetts Institute of Technology, Cambridge, MA.* Let  $f(n) = \sum_P (-1)^{w(P)}$ , where  $P$  ranges over all lattice paths in the plane with  $2n$  steps, starting and ending at the origin, with steps  $(1, 0)$ ,  $(0, 1)$ ,  $(-1, 0)$ ,  $(0, -1)$ , and where  $w(P)$  denotes the winding number of  $P$  with respect to the point  $(1/2, 1/2)$ . Show that  $f(n) = 4^n C_n$ , where  $C_n = \binom{2n}{n} / (n + 1)$ , the  $n$ th Catalan number.

### Three-dimensional Lattice Walks in the Upper Half-Space

**10795** [2000, 367]. *Proposed by Emeric Deutsch, Polytechnic University, Brooklyn, NY.* A 3-dimensional lattice walk of length  $n$  takes  $n$  successive unit steps, each in one of the six coordinate directions. How many 3-dimensional lattice walks of length  $n$  are there that begin at the origin and never go below the horizontal plane?

## Another Type of Lattice Path

**10658** [1998, 366]. *Proposed by Emeric Deutsch, Polytechnic University, Brooklyn, NY.* Consider walks on the integer lattice in the plane that start at  $(0, 0)$ , that stay in the first quadrant (they may touch the  $x$ -axis), and such that each step is either  $(2, 1)$ ,  $(1, 2)$ , or  $(1, -1)$ . For each nonnegative integer  $n$ , how many paths are there to  $(3n, 0)$ ?

### The First Third

**6637** [1990, 621]. *Proposed by Herbert S. Wilf, University of Pennsylvania, Philadelphia, PA.*

Let  $f(n)$  be the sum of the first one-third of the coefficients in the expansion of  $(1+x)^{3n}$ , i.e.,

$$f(n) = \sum_{k=0}^n \binom{3n}{k} \quad (n = 0, 1, 2, \dots).$$

Prove that

$$\sum_{n=0}^{\infty} f(n) \left( \frac{4u^2}{27} \right)^n = \frac{u}{u - 2 \sin\left(\frac{1}{3} \arcsin u\right)} - \frac{2u}{2u - 3 \sin\left(\frac{1}{3} \arcsin u\right)}.$$

**11501.** *Proposed by Finbarr Holland, University College Cork, Cork, Ireland. (Correction)* Let

$$g(z) = 1 - \frac{3}{\frac{1}{1-az} + \frac{1}{1-iz} + \frac{1}{1+iz}}.$$

Show that the coefficients in the Taylor series expansion of  $g$  about 0 are all nonnegative if and only if  $a \geq \sqrt{3}$ .

**11567.** *Proposed by David Callan, University of Wisconsin-Madison, Madison, WI.* How many arrangements  $(a_1, \dots, a_{2n})$  of the multiset  $\{1, 1, 2, 2, \dots, n, n\}$  satisfy the following two conditions: (i) All entries between the two occurrences of any given value  $i$  exceed  $i$ , and (ii) No three entries increase from left to right with the last two adjacent? (When  $n = 3$ , one such arrangement is 122133.)

**11573.** *Proposed by Rob Pratt, SAS Institute, Cary, NC.* A *Sudoku permutation matrix* (SPM) of order  $n^2$  is a permutation matrix of order  $n^2$  with exactly one 1 in each of the  $n^2$  submatrices of order  $n$  obtained by partitioning the original matrix into an  $n$ -by- $n$  array of submatrices. Thus, for  $n = 2$ , the permutation 1324 yields an SPM, but the identity permutation 1234 does not. Find the number of SPMs of order  $n^2$ .

**11610.** *Proposed by Richard P. Stanley, Massachusetts Institute of Technology, Cambridge, MA.* Let  $f(n)$  be the number of binary words  $a_1 \cdots a_n$  of length  $n$  that have the same number of pairs  $a_i a_{i+1}$  equal to 00 as equal to 01. Show that

$$\sum_{n=0}^{\infty} f(n)t^n = \frac{1}{2} \left( \frac{1}{1-t} + \frac{1+2t}{\sqrt{(1-t)(1-2t)(1+t+2t^2)}} \right).$$

► Last one as an exercise tomorrow.

# A money changing problem

**Question<sup>†</sup>**: The number of ways one can change any amount of banknotes of 10 €, 20 €, ... using coins of 50 cents, 1 € and 2 € is always a perfect square.



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<sup>†</sup>Free adaptation of Pb. 1, Ch. 1, p. 1, vol. 1 of Pólya and Szegő's Problems Book (1925)

This is equivalent to finding the number  $M_{20k}$  of solutions  $(a, b, c) \in \mathbb{N}^3$  of

$$a + 2b + 4c = 20k.$$

Euler-Comtet's denumerants:  $\sum_{n \geq 0} M_n x^n = \frac{1}{(1-x)(1-x^2)(1-x^4)}.$

```
> f:=1/(1-x)/(1-x^2)/(1-x^4):
> S:=series(f,x,201):
> [seq(coeff(S,x,20*k),k=1..10)];
```

[36, 121, 256, 441, 676, 961, 1296, 1681, 2116, 2601]

```
> subs(n=20*k,gfun[ratpolytocoeff](f,x,n)):
```

$$\frac{17}{32} + \frac{(20k+1)(20k+2)}{16} + 5k + \frac{(-1)^{-20k}(20k+1)}{16} + \frac{5(-1)^{-20k}}{32} + \sum_{\alpha^2+1=0} \left( -\frac{(\frac{1}{16} - \frac{1}{16}\alpha)\alpha^{-20k}}{\alpha} \right)$$

```
> value(subs(_alpha^(-20*k)=1,%)):
> simplify(%) assuming k::posint:
> factor(%);
```

2

(5 k + 1)

# INTRODUCTION

## 2. Computer Algebra

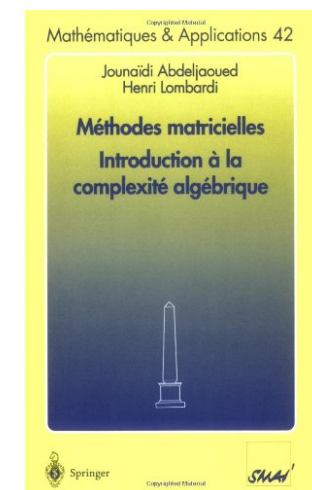
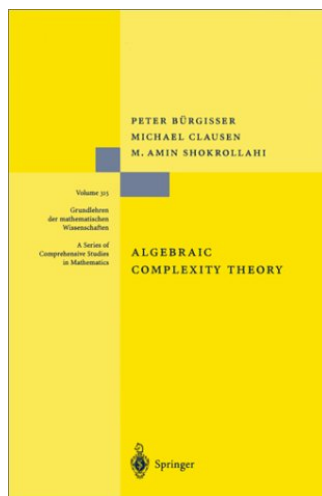
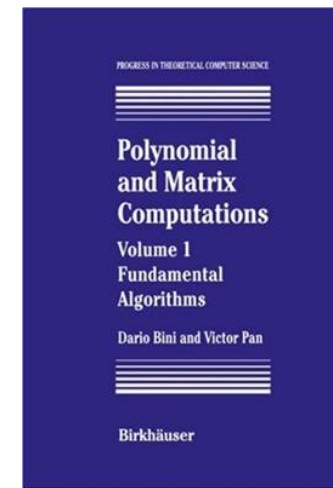
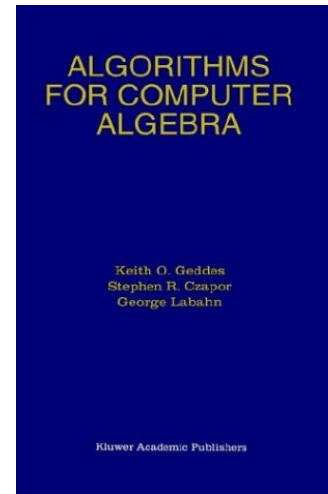
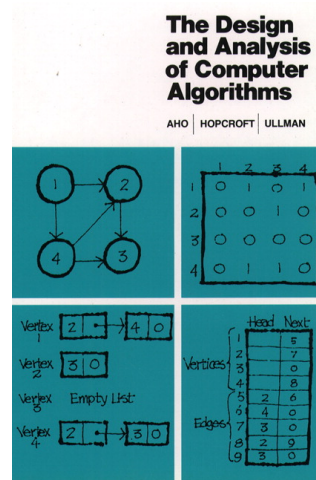
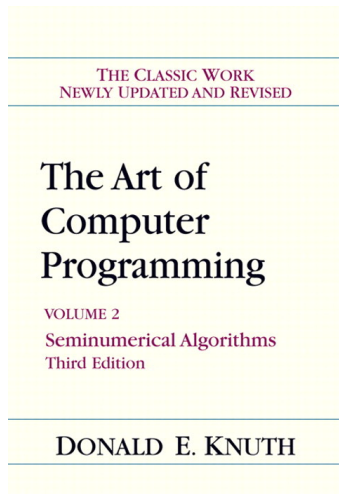
# General framework

Computeralgebra = effectivemathematics *and* algebraiccomplexity

- Effective mathematics: what can we compute?
- their complexity: how fast?



# Computer algebra books



# Mathematical Objects

- Main objects

- integers  $\mathbb{Z}$
- polynomials  $\mathbb{K}[x]$
- rational functions  $\mathbb{K}(x)$
- power series  $\mathbb{K}[[x]]$
- matrices  $\mathcal{M}_r(\mathbb{K})$
- linear recurrences with constant, or polynomial, coefficients  $\mathbb{K}[n]\langle S_n \rangle$
- linear differential equations with polynomial coefficients  $\mathbb{K}[x]\langle \partial_x \rangle$

where  $\mathbb{K}$  is a field (generally supposed of characteristic 0 or large)

- Secondary/auxiliary objects

- polynomial matrices  $\mathcal{M}_r(\mathbb{K}[x])$
- power series matrices  $\mathcal{M}_r(\mathbb{K}[[x]])$

# Overview

## Today

1. Introduction
2. High Precision **Approximations**
  - Fast multiplication, binary splitting, Newton iteration
3. Tools for **Conjectures**
  - Hermite-Padé approximants,  $p$ -curvature

## Tomorrow morning

4. Tools for **Proofs**
  - Symbolic method, resultants, D-finiteness, creative telescoping

## Tomorrow night

- Exercises with Maple

# **HIGH PRECISION**

## **1. Fast Multiplication**

# Complexity yardsticks

Important features:

- addition is easy: naive algorithm already optimal
- multiplication is the most basic (non-trivial) problem
- almost all problems can be reduced to multiplication

Are there quasi-optimal algorithms for:

- integer/polynomial/power series multiplication?
- matrix multiplication?

Yes!

Big open problem!

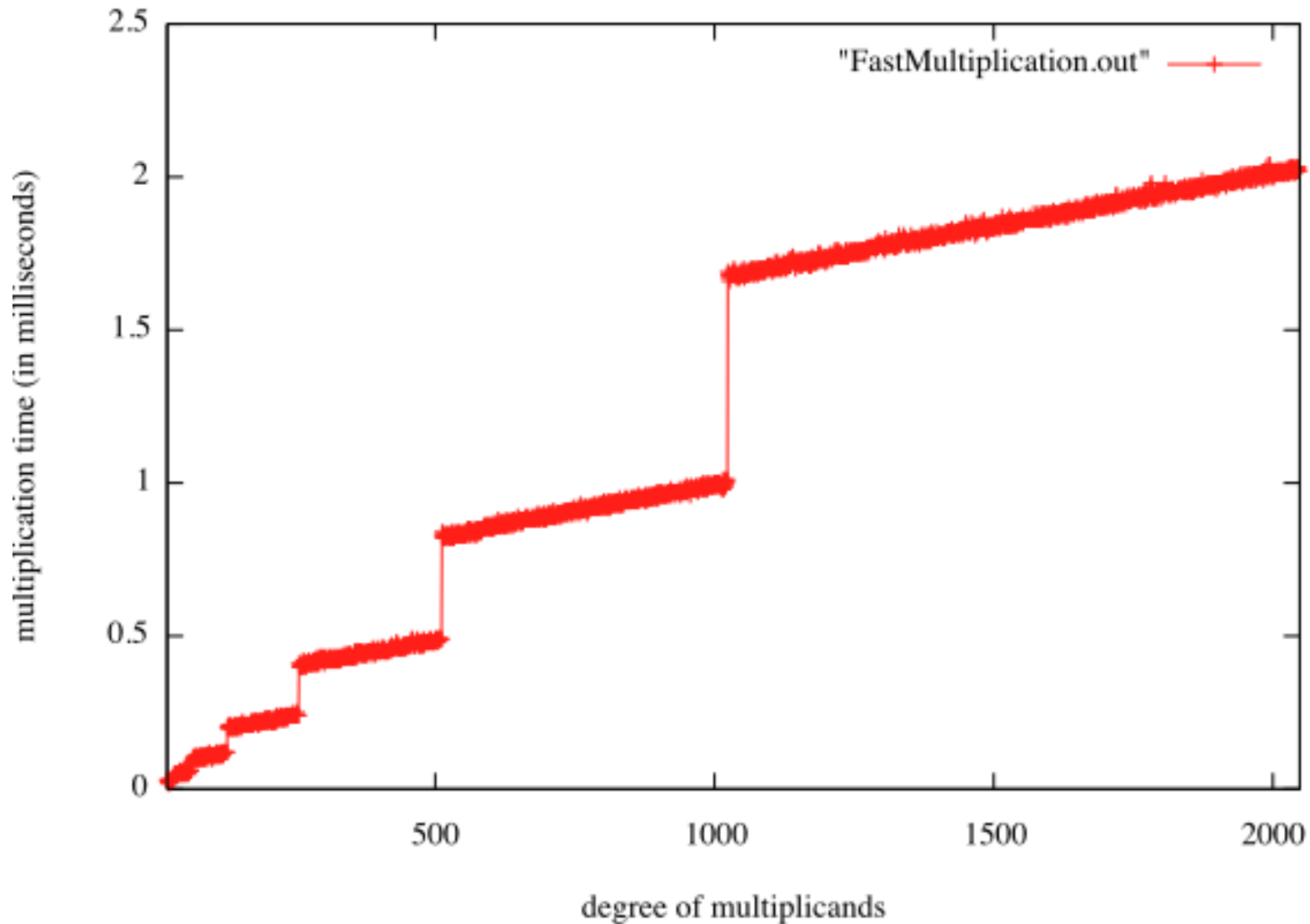
# Complexity yardsticks

$M(n)$  = complexity of **multiplication** in  $\mathbb{K}[x]_{<n}$ , and of  $n$ -bit integers  
=  $O(n^2)$  by the naive algorithm  
=  $O(n^{1.58})$  by **Karatsuba**'s algorithm  
=  $O(n^{\log_\alpha(2\alpha-1)})$  by the **Toom-Cook** algorithm ( $\alpha \geq 3$ )  
=  $O(n \log n \log \log n)$  by the **Schönhage-Strassen** algorithm

$MM(r)$  = complexity of **matrix product** in  $\mathcal{M}_r(\mathbb{K})$   
=  $O(r^3)$  by the naive algorithm  
=  $O(r^{2.81})$  by **Strassen**'s algorithm  
=  $O(r^{2.38})$  by the **Coppersmith-Winograd** algorithm

$MM(r, n)$  = complexity of **polynomial matrix product** in  $\mathcal{M}_r(\mathbb{K}[x]_{<n})$   
=  $O(r^3 M(n))$  by the naive algorithm  
=  $O(MM(r) n \log(n) + r^2 n \log n \log \log n)$  by the **Cantor-Kaltofen** algo  
=  $O(MM(r) n + r^2 M(n))$  by the **B-Schost** algorithm

# Fast polynomial multiplication in practice



Practical complexity of Magma's multiplication in  $\mathbb{F}_p[x]$ , for  $p = 29 \times 2^{57} + 1$ .

# What can be computed in 1 minute with a CA system\*

polynomial product<sup>†</sup> in degree 14,000,000 (>1 year with schoolbook)

product of two integers with 500,000,000 binary digits

factorial of  $N = 20,000,000$  (output of 140,000,000 digits)

gcd of two polynomials of degree 600,000

resultant of two polynomials of degree 40,000

factorization of a univariate polynomial of degree 4,000

factorization of a bivariate polynomial of total degree 500

resultant of two bivariate polynomials of total degree 100 (output 10,000)

product/sum of two algebraic numbers of degree 450 (output 200,000)

determinant (char. polynomial) of a matrix with 4,500 (2,000) rows

determinant of an integer matrix with 32-bit entries and 700 rows

---

\*on a PC, (Intel Xeon X5160, 3GHz processor, with 8GB RAM), running Magma V2.16-7

<sup>†</sup>in  $\mathbb{K}[x]$ , for  $\mathbb{K} = \mathbb{F}_{67108879}$



# Discrete Fourier Transform

[Gauss 1866, Cooley-Tukey 1965]

**DFT Problem:** Given  $n = 2^k$ ,  $f \in \mathbb{K}[x]_{<n}$ , and  $\omega \in \mathbb{K}$  a primitive  $n$ -th root of unity, compute  $(f(1), f(\omega), \dots, f(\omega^{n-1}))$

**Idea:** Write  $f = f_{\text{even}}(x^2) + x f_{\text{odd}}(x^2)$ , with  $\deg(f_{\text{even}}), \deg(f_{\text{odd}}) < n/2$ .

Then  $f(\omega^j) = f_{\text{even}}(\omega^{2j}) + \omega^j f_{\text{odd}}(\omega^{2j})$ , and  $(\omega^{2j})_{0 \leq j < n} = \frac{n}{2}$ -roots of 1.

**Complexity:**  $F(n) = 2 \cdot F(n/2) + O(n) \implies F(n) = O(n \log n)$

# Inverse DFT

**IDFT Problem:** Given  $n = 2^k$ ,  $v_0, \dots, v_{n-1} \in \mathbb{K}$  and  $\omega \in \mathbb{K}$  a primitive  $n$ -th root of unity, compute  $f \in \mathbb{K}[x]_{<n}$  such that  $f(1) = v_0, \dots, f(\omega^{n-1}) = v_{n-1}$

- $V_\omega \cdot V_{\omega^{-1}} = n \cdot I_n \rightarrow$  performing the **inverse DFT** in size  $n$  amounts to:
  - performing a DFT at

$$\frac{1}{1}, \frac{1}{\omega}, \dots, \frac{1}{\omega^{n-1}}$$

- dividing the results by  $n$ .
- this new DFT is the same as before:

$$\frac{1}{\omega^i} = \omega^{n-i},$$

so the outputs are just shuffled.

**Consequence:** the cost of the **inverse DFT** is  $O(n \log(n))$

# FFT polynomial multiplication

Suppose the basefield  $\mathbb{K}$  contains enough roots of unity

To multiply two polynomials  $f, g$  in  $\mathbb{K}[x]$ , of degrees  $< n$ :

- find  $N = 2^k$  such that  $h = fg$  has degree less than  $N$   $N \leq 4n$
- compute  $\text{DFT}(f, N)$  and  $\text{DFT}(g, N)$   $O(N \log(N))$
- multiply pointwise these values to get  $\text{DFT}(h, N)$   $O(N)$
- recover  $h$  by  $\text{inverse DFT}$   $O(N \log(N))$

**Complexity:**  $O(N \log(N)) = O(n \log(n))$

**General case:** Create artificial roots of unity  $O(n \log(n) \log \log n)$

# HIGH PRECISION

## 2. Binary Splitting

## Example: fast factorial

**Problem:** Compute  $N! = 1 \times \cdots \times N$

**Naive iterative way:** unbalanced multiplicands

$\tilde{O}(N^2)$

- **Binary Splitting:** balance computation sequence so as to take advantage of **fast** multiplication (operands of same sizes):

$$N! = \underbrace{(1 \times \cdots \times \lfloor N/2 \rfloor)}_{\text{size } \frac{1}{2} N \log N} \times \underbrace{((\lfloor N/2 \rfloor + 1) \times \cdots \times N)}_{\text{size } \frac{1}{2} N \log N}$$

and recurse. Complexity  $\tilde{O}(N)$ .

- Extends to **matrix factorials**  $A(N)A(N-1)\cdots A(1)$   
→ recurrences of arbitrary order.

$\tilde{O}(N)$

# Application to recurrences

**Problem:** Compute the  $N$ -th term  $u_N$  of a  $P$ -recursive sequence

$$p_r(n)u_{n+r} + \cdots + p_0(n)u_n = 0, \quad (n \in \mathbb{N})$$

**Naive algorithm:** unroll the recurrence  $\tilde{O}(N^2)$  bit ops.

**Binary splitting:**  $U_n = (u_n, \dots, u_{n+r-1})^T$  satisfies the 1st order recurrence

$$U_{n+1} = \frac{1}{p_r(n)} A(n) U_n \quad \text{with} \quad A(n) = \begin{bmatrix} & & & p_r(n) \\ & & & \vdots \\ & & & p_r(n) \\ -p_0(n) & -p_1(n) & \cdots & -p_{r-1}(n) \end{bmatrix}.$$

$\implies u_N$  reads off the **matrix factorial**  $A(N-1) \cdots A(0)$

**[Chudnovsky-Chudnovsky, 1987]:** Binary splitting strategy  $\tilde{O}(N)$  bit ops.

# Application: fast computation of $e = \exp(1)$

[Brent 1976]

$$e_n = \sum_{k=0}^n \frac{1}{k!} \quad \longrightarrow \quad \exp(1) = 2.7182818284590452 \dots$$

Recurrence  $e_n - e_{n-1} = 1/n! \iff n(e_n - e_{n-1}) = e_{n-1} - e_{n-2}$  rewrites

$$\begin{bmatrix} e_{N-1} \\ e_N \end{bmatrix} = \frac{1}{N} \underbrace{\begin{bmatrix} 0 & N \\ -1 & N+1 \end{bmatrix}}_{C(N)} \begin{bmatrix} e_{N-2} \\ e_{N-1} \end{bmatrix} = \frac{1}{N!} C(N)C(N-1) \dots C(1) \begin{bmatrix} 0 \\ 1 \end{bmatrix}.$$

►  $e_N$  in  $\tilde{O}(N)$  bit operations [Brent 1976]

► generalizes to the evaluation of any D-finite series at an algebraic number  
[Chudnovsky-Chudnovsky 1987]  $\tilde{O}(N)$  bit ops.

# Implementation in gfun

[Mezzarobba, S. 2010]

```
> rec:={n*(e(n) - e(n-1)) = e(n-1) - e(n-2), e(0)=1, e(1)=2};
```

```
> pro:=rectoproc(rec,e(n));
```

```
pro := proc(n::nonnegint)
```

```
local i1, loc0, loc1, loc2, tmp2, tmp1, i2;
```

```
  if n <= 22 then
```

```
    loc0 := 1;
```

```
    loc1 := 2;
```

```
    if n = 0 then return loc0
```

```
    else for i1 to n - 1 do
```

```
      loc2 := (-loc0 + loc1 + loc1*(i1 + 1))/(i1 + 1); loc0 := loc1; loc1 := loc2
```

```
    end do
```

```
    end if; loc1
```

```
  else
```

```
    tmp1 := 'gfun/rectoproc/binsplit'([  
      'ndmatrix'(Matrix([[0, i2 + 2], [-1, i2 + 3]]), i2 + 2), i2, 0, n,  
      matrix_ring(ad, pr, ze, ndmatrix(Matrix(2, 2, [[...],[...]]),  
      datatype = anything, storage = empty, shape = [identity]), 1)),  
      expected_entry_size], Vector(2, [...], datatype = anything));
```

```
    tmp1 := subs({e(0) = 1, e(1) = 2}, tmp1); tmp1
```

```
  end if
```

```
end proc
```

```
> tt:=time(): x:=pro(50000): time()-tt, evalf(x-exp(1), 200000);
```

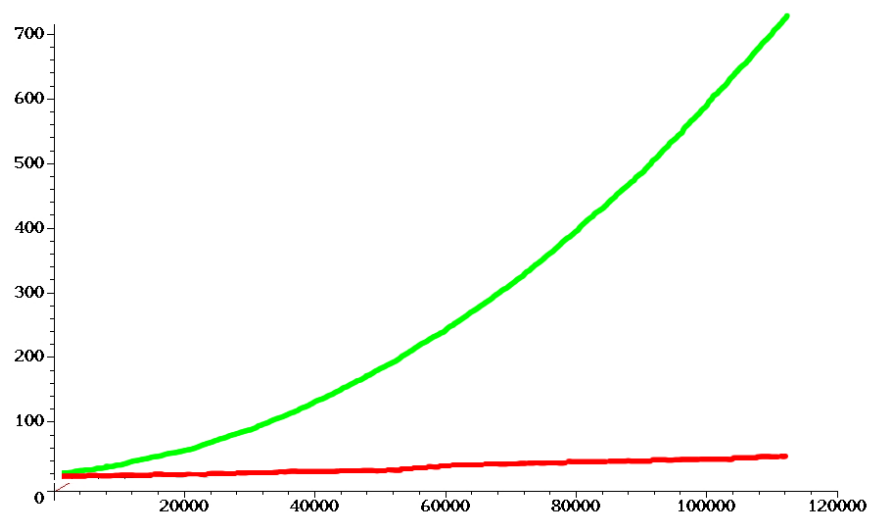
1.827, 0.



# Application: record computation of $\pi$

[Chudnovsky-Chudnovsky 1987] fast convergence hypergeometric identity

$$\frac{1}{\pi} = \frac{1}{53360\sqrt{640320}} \sum_{n \geq 0} \frac{(-1)^n (6n)! (13591409 + 545140134n)}{n!^3 (3n)! (8 \cdot 100100025 \cdot 327843840)^n}.$$



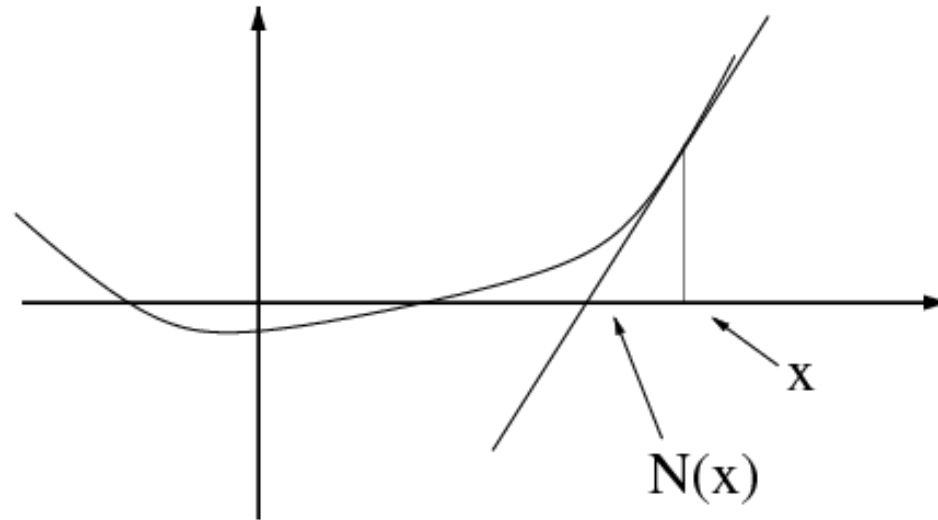
- ▶ **Used in Maple & Mathematica:** 1st order recurrence, yields 14 correct digits per iteration  $\longrightarrow$  4 billion digits [Chudnovsky-Chudnovsky 1994]
- ▶ **Current record on a PC:** 10000 billion digits [Kondo & Yee 2011]

# **HIGH PRECISION**

## **3. Newton Iteration**

# Newton's tangent method: real case

[Newton, 1671]



$$x_{\kappa+1} = \mathcal{N}(x_{\kappa}) = x_{\kappa} - (x_{\kappa}^2 - 2)/(2x_{\kappa}), \quad x_0 = 1$$

$$x_1 = 1.50000000000000000000000000000000$$

$$x_2 = 1.41666666666666666666666666666667$$

$$x_3 = 1.4142156862745098039215686274510$$

$$x_4 = 1.4142135623746899106262955788901$$

$$x_5 = 1.4142135623730950488016896235025$$

# Newton's tangent method: power series case

$$x_{\kappa+1} = \mathcal{N}(x_{\kappa}) = x_{\kappa} - (x_{\kappa}^2 - (1 - t))/(2x_{\kappa}), \quad x_0 = 1$$

$$x_1 = 1 - \frac{1}{2}t$$

$$x_2 = 1 - \frac{1}{2}t - \frac{1}{8}t^2 - \frac{1}{16}t^3 - \frac{1}{32}t^4 - \frac{1}{64}t^5 - \frac{1}{128}t^6 - \frac{1}{256}t^7 - \frac{1}{512}t^8 - \frac{1}{1024}t^9 + \dots$$

$$x_3 = 1 - \frac{1}{2}t - \frac{1}{8}t^2 - \frac{1}{16}t^3 - \frac{5}{128}t^4 - \frac{7}{256}t^5 - \frac{21}{1024}t^6 - \frac{33}{2048}t^7 - \frac{107}{8192}t^8 - \frac{177}{16384}t^9 + \dots$$

# Newton's tangent method: power series case

In order to solve  $\varphi(x, g) = 0$  in  $\mathbb{K}[[x]]$  (where  $\varphi \in \mathbb{K}[[x, y]]$ ,  $\varphi(0, 0) = 0$  and  $\varphi_y(0, 0) \neq 0$ ), iterate

$$g_{\kappa+1} = g_{\kappa} - \frac{\varphi(g_{\kappa})}{\varphi_y(g_{\kappa})} \pmod{x^{2^{\kappa+1}}}$$

$$g - g_{\kappa+1} = g - g_{\kappa} + \frac{\varphi(g) + (g_{\kappa} - g)\varphi_y(g) + O((g - g_{\kappa})^2)}{\varphi_y(g) + O(g - g_{\kappa})} = O((g - g_{\kappa})^2).$$

- ▶ The number of correct coefficients **doubles** after each iteration
- ▶ **Total cost** = **2** × (the cost of the **last** iteration)

**Theorem** [Cook 1966, Sieveking 1972 & Kung 1974, Brent 1975]

Division, logarithm and exponential of power series in  $\mathbb{K}[[x]]$  can be computed at precision  $N$  using  $O(M(N))$  operations in  $\mathbb{K}$

# Division, logarithm and exponential of power series

[Sieveking1972, Kung1974, Brent1975]

To compute the **reciprocal** of  $f \in \mathbb{K}[[x]]$  with  $f(0) \neq 0$ , choose  $\varphi(g) = 1/g - f$ :

$$g_0 = 1/f_0 \quad \text{and} \quad g_{\kappa+1} = g_{\kappa} + g_{\kappa}(1 - fg_{\kappa}) \quad \text{mod } x^{2^{\kappa+1}} \quad \text{for } \kappa \geq 0.$$

**Complexity:**  $C(N) = C(N/2) + O(M(N)) \quad \implies \quad C(N) = O(M(N))$

**Corollary:** division of power series at precision  $N$  in  $O(M(N))$

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**Corollary:** division of power series at precision  $N$  in  $O(M(N))$

**Corollary: Logarithm**  $\log(f) = -\sum_{i \geq 1} \frac{(1-f)^i}{i}$  of  $f \in 1 + x\mathbb{K}[[x]]$  in  $O(M(N))$ :

- compute the Taylor expansion of  $h = f'/f$  modulo  $x^{N-1}$   $O(M(N))$
- take the antiderivative of  $h$   $O(N)$

# Division, logarithm and exponential of power series

[Sieveking1972, Kung1974, Brent1975]

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- compute the Taylor expansion of  $h = f'/f$  modulo  $x^{N-1}$   $O(M(N))$
- take the antiderivative of  $h$   $O(N)$

**Corollary: Exponential**  $\exp(f) = \sum_{i \geq 0} \frac{f^i}{i!}$  of  $f \in x\mathbb{K}[[x]]$ . Use  $\phi(g) = \log(g) - f$ :

$$g_0 = 1 \quad \text{and} \quad g_{\kappa+1} = g_{\kappa} - g_{\kappa}(\log(g_{\kappa}) - f) \quad \text{mod } x^{2^{\kappa+1}} \quad \text{for } \kappa \geq 0.$$

**Complexity:**  $C(N) = C(N/2) + O(M(N)) \implies C(N) = O(M(N))$



# Application: Euclidean division for polynomials

[Strassen, 1973]

Pb: Given  $F, G \in \mathbb{K}[x]_{\leq N}$ , compute  $(Q, R)$  in **Euclidean division**  $F = QG + R$

Naive algorithm:

$O(N^2)$

Idea: look at  $F = QG + R$  **from infinity**:  $Q \sim_{+\infty} F/G$

Let  $N = \deg(F)$  and  $n = \deg(G)$ . Then  $\deg(Q) = N - n$ ,  $\deg(R) < n$  and

$$\underbrace{F(1/x)x^N}_{\text{rev}(F)} = \underbrace{G(1/x)x^n}_{\text{rev}(G)} \cdot \underbrace{Q(1/x)x^{N-n}}_{\text{rev}(Q)} + \underbrace{R(1/x)x^{\deg(R)}}_{\text{rev}(R)} \cdot x^{N-\deg(R)}$$

Algorithm:

- Compute  $\text{rev}(Q) = \text{rev}(F)/\text{rev}(G) \pmod{x^{N-n+1}}$   $O(M(N))$
- Recover  $Q$   $O(N)$
- Deduce  $R = F - QG$   $O(M(N))$

# Application: conversion coefficients $\leftrightarrow$ power sums

[Schönhage, 1982]

Any polynomial  $F = x^n + a_1x^{n-1} + \dots + a_n$  in  $\mathbb{K}[x]$  can be represented by its first  $n$  power sums  $S_i = \sum_{F(\alpha)=0} \alpha^i$

Conversions **coefficients  $\leftrightarrow$  power sums** can be performed

- either in  $O(n^2)$  using **Newton identities** (naive way):

$$ia_i + S_1a_{i-1} + \dots + S_i = 0, \quad 1 \leq i \leq n$$

- or in  $O(M(n))$  using **generating series**

$$\frac{\text{rev}(F)'}{\text{rev}(F)} = - \sum_{i \geq 0} S_{i+1} x^i \iff \text{rev}(F) = \exp \left( - \sum_{i \geq 1} \frac{S_i}{i} x^i \right)$$

# Application: special bivariate resultants

[B-Flajolet-S-Schost, 2006]

Composed products and sums: manipulation of algebraic numbers

$$F \otimes G = \prod_{F(\alpha)=0, G(\beta)=0} (x - \alpha\beta), \quad F \oplus G = \prod_{F(\alpha)=0, G(\beta)=0} (x - (\alpha + \beta))$$

Output size:

$$N = \deg(F) \deg(G)$$

Linear algebra:  $\chi_{xy}, \chi_{x+y}$  in  $\mathbb{K}[x, y]/(F(x), G(y))$

$$O(\text{MM}(N))$$

Resultants:  $\text{Res}_y (F(y), y^{\deg(G)} G(x/y))$ ,  $\text{Res}_y (F(y), G(x - y))$

$$O(N^{1.5})$$

Better:  $\otimes$  and  $\oplus$  are easy in Newton representation

$$O(\text{M}(N))$$

$$\sum \alpha^s \sum \beta^s = \sum (\alpha\beta)^s \quad \text{and}$$
$$\sum \frac{\sum (\alpha + \beta)^s}{s!} x^s = \left( \sum \frac{\sum \alpha^s}{s!} x^s \right) \left( \sum \frac{\sum \beta^s}{s!} x^s \right)$$

Corollary: Fast polynomial shift  $P(x + a) = P(x) \oplus (x + a)$

$$O(\text{M}(\deg(P)))$$

# Newton iteration on power series: operators and systems

In order to solve an equation  $\phi(Y) = 0$ , with  $\phi : (\mathbb{K}[[x]])^r \rightarrow (\mathbb{K}[[x]])^r$ ,

1. **Linearize**:  $\phi(Y_\kappa - U) = \phi(Y_\kappa) - D\phi|_{Y_\kappa} \cdot U + O(U^2)$ ,  
where  $D\phi|_Y$  is the differential of  $\phi$  at  $Y$ .
2. **Iterate**:  $Y_{\kappa+1} = Y_\kappa - U_{\kappa+1}$ , where  $U_{\kappa+1}$  is found by
3. **Solve linear** equation:  $D\phi|_{Y_\kappa} \cdot U = \phi(Y_\kappa)$  with  $\text{val } U > 0$ .

Then, the sequence  $Y_\kappa$  **converges quadratically** to the solution  $Y$ .

# Application: inversion of power series matrices

[Schulz, 1933]

To compute the inverse  $Z$  of a matrix of power series  $Y \in \mathcal{M}_r(\mathbb{K}[[x]])$ :

- Choose the map  $\phi : Z \mapsto I - YZ$  with differential  $D\phi|_Y : U \mapsto -YU$
- Equation for  $U$ :  $-YU = I - YZ_\kappa \pmod{x^{2^{\kappa+1}}}$
- Solution:  $U = -Y^{-1}(I - YZ_\kappa) = -Z_\kappa(I - YZ_\kappa) \pmod{x^{2^{\kappa+1}}}$

This yields the following Newton-type iteration for  $Y^{-1}$

$$Z_{\kappa+1} = Z_\kappa + Z_\kappa(I_r - YZ_\kappa) \pmod{x^{2^{\kappa+1}}}$$

Complexity:

$$C_{\text{inv}}(N) = C_{\text{inv}}(N/2) + O(\text{MM}(r, N)) \quad \Longrightarrow \quad C_{\text{inv}}(N) = O(\text{MM}(r, N))$$

# Application: non-linear systems

In order to solve a system  $Y = H(Y) = \phi(Y) + Y$ , with  $H : (\mathbb{K}[[x]])^r \rightarrow (\mathbb{K}[[x]])^r$ , such that  $I_r - \partial H/\partial Y$  is invertible at 0.

1. **Linearize:**  $\phi(Y_\kappa - U) - \phi(Y_\kappa) = U - \partial H/\partial Y(Y_\kappa) \cdot U + O(U^2)$ .
2. **Iterate**  $Y_{\kappa+1} = Y_\kappa - U_{\kappa+1}$ , where  $U_{\kappa+1}$  is found by
3. **Solve linear** equation:  $(I_r - \partial H/\partial Y(Y_\kappa)) \cdot U = H(Y_\kappa) - Y_\kappa$  with  $\text{val } U > 0$ .

This yields the following Newton-type iteration:

$$\begin{cases} Z_{\kappa+1} &= Z_\kappa + Z_\kappa(I_r - (I_r - \partial H/\partial Y(Y_\kappa))Z_\kappa) \pmod{x^{2^{\kappa+1}}} \\ Y_{\kappa+1} &= Y_\kappa - Z_{\kappa+1}(H(Y_\kappa) - Y_\kappa) \pmod{x^{2^{\kappa+1}}} \end{cases}$$

computing simultaneously a matrix and a vector.

## Example: Mappings

- > mappings := {M=Set(Cycle(Tree)), Tree=Prod(Z, Set(Tree))}:
- > combstruct [gfeqns] (mappings, labeled, x);

$$[M(x) = \frac{1}{1 - Tree(x)}, \quad Tree(x) = x \exp(Tree(x))]$$

- > countmappings := SeriesNewtonIteration(mappings, labelled, x):
- > countmappings(10);

$$\left[ M = 1 + x + 2x^2 + \frac{9}{2}x^3 + \frac{32}{3}x^4 + \frac{625}{24}x^5 + \frac{324}{5}x^6 \right. \\ \left. + \frac{117649}{720}x^7 + \frac{131072}{315}x^8 + \frac{4782969}{4480}x^9 + O(x^{10}), \right. \\ Tree = x + x^2 + \frac{3}{2}x^3 + \frac{8}{3}x^4 + \frac{125}{24}x^5 + \frac{54}{5}x^6 + \\ \left. \frac{16807}{720}x^7 + \frac{16384}{315}x^8 + \frac{531441}{4480}x^9 + O(x^{10}) \right]$$

Code Pivoteau-S-Soria, should end up in combstruct

# Application: quasi-exponential of power series matrices

[B-Chyzak-Ollivier-Salvy-Schost-Sedoglavic 2007]

To compute the solution  $Y \in \mathcal{M}_r(\mathbb{K}[[x]])$  of the system  $Y' = AY$

- choose the map  $\phi : Y \mapsto Y' - AY$ , with differential  $\phi$ .
- the equation for  $U$  is  $U' - AU = Y'_\kappa - AY_\kappa \pmod{x^{2^{\kappa+1}}}$
- the method of variation of constants yields the solution  $U = Y_\kappa V_\kappa \pmod{x^{2^{\kappa+1}}}$ ,  $Y'_\kappa - AY_\kappa = Y_\kappa V'_\kappa \pmod{x^{2^{\kappa+1}}}$

This yields the following Newton-type iteration for  $Y$ :

$$Y_{\kappa+1} = Y_\kappa - Y_\kappa \int Y_\kappa^{-1} (Y'_\kappa - AY_\kappa) \pmod{x^{2^{\kappa+1}}}$$

Complexity:

$$C_{\text{solve}}(N) = C_{\text{solve}}(N/2) + O(\text{MM}(r, N)) \quad \Longrightarrow \quad C_{\text{solve}}(N) = O(\text{MM}(r, N))$$



# **TOOLS FOR CONJECTURES**

## **1. Hermite-Padé Approximants**

# Definition

**Definition:** Given a column vector  $\mathbf{F} = (f_1, \dots, f_n)^T \in \mathbb{K}[[x]]^n$  and an  $n$ -tuple  $\mathbf{d} = (d_1, \dots, d_n) \in \mathbb{N}^n$ , a **Hermite-Padé approximant of type  $\mathbf{d}$  for  $\mathbf{F}$**  is a row vector  $\mathbf{P} = (P_1, \dots, P_n) \in \mathbb{K}[x]^n$ , ( $\mathbf{P} \neq 0$ ), such that:

- (1)  $\mathbf{P} \cdot \mathbf{F} = P_1 f_1 + \dots + P_n f_n = O(x^\sigma)$  with  $\sigma = \sum_i (d_i + 1) - 1$ ,
- (2)  $\deg(P_i) \leq d_i$  for all  $i$ .

$\sigma$  is called the **order** of the approximant  $\mathbf{P}$ .

- Very useful concept in number theory (transcendence theory):
- [[Hermite, 1873](#)]:  $e$  is transcendent.
  - [[Lindemann, 1882](#)]:  $\pi$  is transcendent, and so does  $e^\alpha$  for any  $\alpha \in \overline{\mathbb{Q}} \setminus \{0\}$ .
  - [[Beukers, 1981](#)]: reformulate Apéry's proof that  $\zeta(3) = \sum_n \frac{1}{n^3}$  is irrational.
  - [[Rivoal, 2000](#)]: there exist an infinite number of  $k$  such that  $\zeta(2k + 1) \notin \mathbb{Q}$ .

## Worked example

Let us compute a Hermite-Padé approximant of **type (1, 1, 1)** for  $(1, C, C^2)$ , where  $C(x) = 1 + x + 2x^2 + 5x^3 + 14x^4 + 42x^5 + O(x^6)$ .

This boils down to finding  $\alpha_0, \alpha_1, \beta_0, \beta_1, \gamma_0, \gamma_1$  such that

$$\alpha_0 + \alpha_1 x + (\beta_0 + \beta_1 x)(1 + x + 2x^2 + 5x^3 + 14x^4) + (\gamma_0 + \gamma_1 x)(1 + 2x + 5x^2 + 14x^3 + 42x^4) = O(x^5).$$

By identifying coefficients, this is equivalent to a homogeneous linear system:

$$\begin{bmatrix} 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 1 & 2 & 1 \\ 0 & 0 & 2 & 1 & 5 & 2 \\ 0 & 0 & 5 & 2 & 14 & 5 \\ 0 & 0 & 14 & 5 & 42 & 14 \end{bmatrix} \times \begin{bmatrix} \alpha_0 \\ \alpha_1 \\ \beta_0 \\ \beta_1 \\ \gamma_0 \\ \gamma_1 \end{bmatrix} = 0 \iff \begin{bmatrix} 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 1 & 2 \\ 0 & 0 & 2 & 1 & 5 \\ 0 & 0 & 5 & 2 & 14 \\ 0 & 0 & 14 & 5 & 42 \end{bmatrix} \times \begin{bmatrix} \alpha_0 \\ \alpha_1 \\ \beta_0 \\ \beta_1 \\ \gamma_0 \end{bmatrix} = -\gamma_1 \begin{bmatrix} 0 \\ 1 \\ 2 \\ 5 \\ 14 \end{bmatrix}.$$

By homogeneity, one can choose  $\gamma_1 = 1$ . Then, the **violet minor** shows that one can take  $(\beta_0, \beta_1, \gamma_0) = (-1, 0, 0)$ . The other values are  $\alpha_0 = 1, \alpha_1 = 0$ .

Thus the approximant is  $(1, -1, x)$ , which corresponds to  $P = 1 - y + xy^2$  such that  $P(x, C(x)) = 0 \pmod{x^5}$ .

# Algebraic and differential approximation = guessing

- Hermite-Padé approximants of  $n = 2$  power series are related to Padé approximants, i.e. to approximation of series by rational functions
- algebraic approximants = Hermite-Padé approximants for  $f_\ell = A^{\ell-1}$ , where  $A \in \mathbb{K}[[x]]$  seriestoalgeq, listtoalgeq
- differential approximants = Hermite-Padé approximants for  $f_\ell = A^{(\ell-1)}$ , where  $A \in \mathbb{K}[[x]]$  seriestodiffeq, listtodiffeq

> listtoalgeq([1,1,2,5,14,42,132,429],y(x));

$$[1 - y(x) + x y(x)^2, \text{ogf}]$$

> listtodiffeq([1,1,2,5,14,42,132,429],y(x));

$$[\{-2 y(x) + (2 - 4 x) \frac{d}{dx} y(x) + x \frac{d^2}{dx^2} y(x)\}, y(0) = 1, D(y)(0) = 1], \text{egf}]$$

# Existence and naive computation

**Theorem** For any vector  $\mathbf{F} = (f_1, \dots, f_n)^T \in \mathbb{K}[[x]]^n$  and for any  $n$ -tuple  $\mathbf{d} = (d_1, \dots, d_n) \in \mathbb{N}^n$ , there exists a **Hermite-Padé approximant of type  $\mathbf{d}$  for  $\mathbf{F}$** .

**Proof:** The undetermined coefficients of  $P_i = \sum_{j=0}^{d_i} p_{i,j} x^j$  satisfy a linear homogeneous system with  $\sigma = \sum_i (d_i + 1) - 1$  equations and  $\sigma + 1$  unknowns.

**Corollary** Computation in  $O(\text{MM}(\sigma)) = O(\sigma^\theta)$ , for  $2 \leq \theta \leq 3$ .

► There are better algorithms:

- The linear system is **structured** (Sylvester-like / quasi-Toeplitz)
- **Derksen's algorithm** (Gaussian-like elimination)  $O(\sigma^2)$
- **Beckermann-Labahn's algorithm** (DAC)  $\tilde{O}(\sigma) = O(\sigma \log^2 \sigma)$

# Quasi-optimal computation

**Theorem** [Beckermann-Labahn, 1994] One can compute a Hermite-Padé approximant of type  $(d, \dots, d)$  for  $\mathbf{F} = (f_1, \dots, f_n)$  in  $O(\text{MM}(n, d) \log(nd))$ .

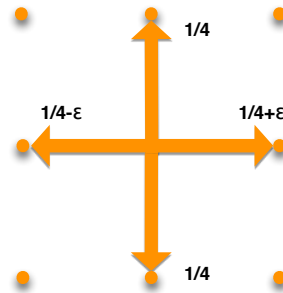
**Ideas:**

- Compute a whole matrix of approximants
- Exploit divide-and-conquer

**Algorithm:**

1. If  $\sigma = n(d + 1) - 1 \leq \text{threshold}$ , call the naive algorithm
  2. Else:
    - (a) recursively compute  $\mathbf{P}_1 \in \mathbb{K}[x]^{n \times n}$  s.t.  $\mathbf{P}_1 \cdot \mathbf{F} = O(x^{\sigma/2})$ ,  $\deg(\mathbf{P}_1) \approx \frac{d}{2}$
    - (b) compute “residue”  $\mathbf{R}$  such that  $\mathbf{P}_1 \cdot \mathbf{F} = x^{\sigma/2} \cdot (\mathbf{R} + O(x^{\sigma/2}))$
    - (c) recursively compute  $\mathbf{P}_2 \in \mathbb{K}[x]^{n \times n}$  s.t.  $\mathbf{P}_2 \cdot \mathbf{R} = O(x^{\sigma/2})$ ,  $\deg(\mathbf{P}_2) \approx \frac{d}{2}$
    - (d) return  $\mathbf{P} := \mathbf{P}_2 \cdot \mathbf{P}_1$
- The precise choices of degrees is a delicate issue
- Gcd, extended gcd, Padé approximants in  $O(\text{M}(n) \log n)$

# Example: Flea from SIAM 100-Digit Challenge



```
> proba:=proc(i,j,n,c)
option remember;
  if abs(i)+abs(j)>n then 0
  elif n=0 then 1
  else
    expand(proba(i-1,j,n-1,c)*(1/4+c)+proba(i+1,j,n-1,c)*(1/4-c)
    +proba(i,j+1,n-1,c)*1/4+proba(i,j-1,n-1,c)*1/4)
  fi
end:
> seq(proba(0,0,k,c),k=0..6);
1, 0,  $\frac{1}{4} - 2c^2$ , 0,  $\frac{9}{64} - \frac{9}{4}c^2 + 6c^4$ , 0,  $\frac{25}{256} - \frac{75}{32}c^2 + 15c^4 - 20c^6$ 
> gfun:-listtodiffeq([seq(proba(0,0,2*k,c),k=0..20)],y(x));
```

$$\begin{aligned}
& [\{ (-1 + 8c^2 + 48xc^4) y(x) + (4 - 8x + 64xc^2 + 192x^2c^4) \frac{d}{dx} y(x) \\
& \quad + (4x + 64x^3c^4 - 4x^2 + 32x^2c^2) \frac{d^2}{dx^2} y(x), \\
& \quad y(0) = 1, D(y)(0) = 1/4 - 2c^2 \}, ogf]
\end{aligned}$$

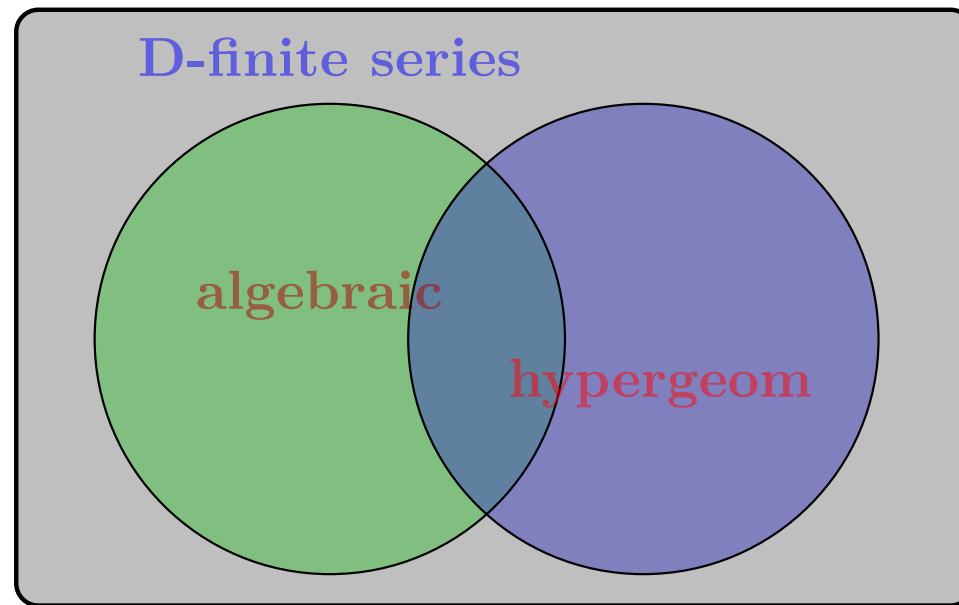
Next steps: `dsolve (+ help)` and evaluation at  $x = 1$ .



# TOOLS FOR CONJECTURES

## 2. $p$ -Curvature of Differential Operators

# Important classes of power series



**Algebraic:**  $S(x) \in \mathbb{K}[[x]]$  root of a polynomial  $P \in \mathbb{K}[x, y]$ .

**D-finite:**  $S(x) \in \mathbb{K}[[x]]$  satisfying a **linear differential equation with polynomial (or rational function) coefficients**  $c_r(x)S^{(r)}(x) + \cdots + c_0(x)S(x) = 0$ .

**Hypergeometric:**  $S(x) = \sum_n s_n x^n$  such that  $\frac{s_{n+1}}{s_n} \in \mathbb{K}(n)$ . E.g.

$${}_2F_1 \left( \begin{matrix} a & b \\ c \end{matrix} \middle| x \right) = \sum_{n=0}^{\infty} \frac{(a)_n (b)_n}{(c)_n} \frac{x^n}{n!}, \quad (a)_n = a(a+1) \cdots (a+n-1).$$

# Linear differential operators

**Definition:** If  $\mathbb{K}$  is a field,  $\mathbb{K}\langle x, \partial; \partial x = x\partial + 1 \rangle$ , or simply  $\mathbb{K}(x)\langle \partial \rangle$ , denotes the associative algebra of linear differential operators with coefficients in  $\mathbb{K}(x)$ .

$\mathbb{K}[x]\langle \partial \rangle$  is called the **(rational) Weyl algebra**. It is the **algebraic formalization** of the notion of linear differential equation with rational function coefficients:

$$a_r(x)y^{(r)}(x) + \cdots + a_1(x)y'(x) + a_0(x)y(x) = 0$$

$$\iff$$

$$L(y) = 0, \quad \text{where} \quad L = a_r(x)\partial^r + \cdots + a_1(x)\partial + a_0(x)$$

The commutation rule  $\partial x = x\partial + 1$  formalizes Leibniz's rule  $(fg)' = f'g + fg'$ .

► Implementation in the DEtools package: **diffop2de, de2diffop, mult**

```
DEtools[mult](Dx, x, [Dx, x]);
```

```
x Dx + 1
```

# Weyl algebra is Euclidean

**Theorem** [Libri 1833, Brassinne 1864, Wedderburn 1932, Ore 1932]

$\mathbb{K}(x)\langle\partial\rangle$  is a non-commutative (left and right) **Euclidean domain**: for any  $A, B \in \mathbb{K}(x)\langle\partial\rangle$ , there exist unique operators  $Q, R \in \mathbb{K}(x)\langle\partial\rangle$  such that

$$A = QB + R, \quad \text{and} \quad \deg_{\partial}(R) < \deg_{\partial}(B).$$

This is called the **Euclidean right division** of  $A$  by  $B$ .

Moreover, any  $A, B \in \mathbb{K}(x)\langle\partial\rangle$  admit a **greatest common right divisor (GCRD)** and a **least common left multiple (LCLM)**. They can be computed by a non-commutative version of the extended Euclidean algorithm.

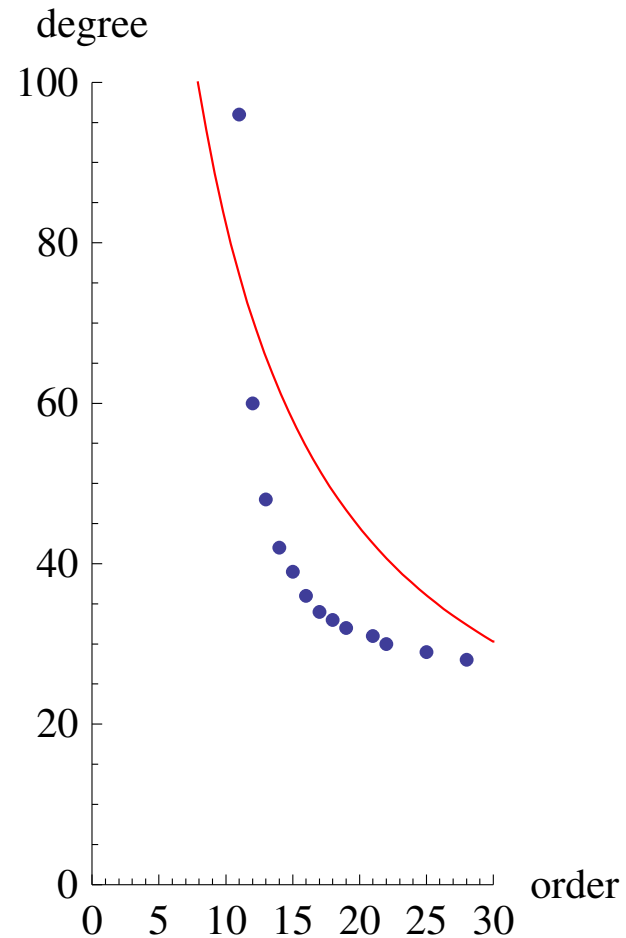
► **rightdivision, GCRD, LCLM** from the DEtools package

```
> rightdivision(Dx^10, Dx^2-x, [Dx, x]) [2];
```

$$(20x^3 + 80) Dx + 100x^2 + x^5$$

proves that  $Ai^{(10)}(x) = (20x^3 + 80)Ai'(x) + (100x^2 + x^5)Ai(x)$

# Application to differential guessing



1000 terms of a series are enough to guess candidate differential equations below the red curve. GCRD of candidates could jump above the red curve.

# The Grothendieck–Katz $p$ -curvatures conjecture

**Q:** when does a differential equation possess a basis of **algebraic solutions**?

E.g. for the Gauss hypergeometric equation  $x(1-x)\partial^2 + (\gamma - (\alpha + \beta + 1)x)\partial - \alpha\beta x$ , **Schwarz's list (1873)** classifies algebraic  ${}_2F_1$ 's in terms of  $\alpha, \beta, \gamma$

**Conjecture** [Grothendieck, 1960's, unpublished; Katz, 1972]

Let  $A \in \mathbb{Q}(x)^{n \times n}$ . The system **(S) :  $y' = Ay$**  has a full set of algebraic solutions if and only if, for almost all prime numbers  $p$ , the system **(S<sub>p</sub>)** defined by reduction of **(S)** modulo  $p$  has a full set of algebraic solutions over  $\mathbb{F}_p(x)$ .

**Definition:** The  **$p$ -curvature of (S)** is the matrix  $A_p$ , where

$$A_0 = I_n, \quad \text{and} \quad A_{\ell+1} = A'_\ell + A_\ell A \quad \text{for} \quad \ell \geq 0.$$

**Theorem** [Cartier, 1957]

The sufficient condition of the **G.-K. Conjecture** is equivalent to  $A_p = 0 \pmod{p}$ .

► For each  $p$ , this can be checked algorithmically.

# Grothendieck's conjecture

**Q:** when does a differential equation possess a basis of algebraic solutions?

For a scalar differential equation, the **G.-K. Conjecture** can be reformulated:

**Grothendieck's Conjecture:** Suppose  $L \in \mathbb{K}(x)\langle\partial\rangle$  is irreducible. The equation **(E) :  $L(y) = 0$**  has a basis of **algebraic solutions** if and only if, for almost all prime numbers  $p$ , the operator  $L$  **right-divides  $\partial^p$  modulo  $p$** .

- ▶ For each  $p$ , this can be checked algorithmically.
- ▶ Conjecture is proved for **Picard-Fuchs equations** [Katz 1972] (in particular, for **diagonals** [Christol 1984]), for  ${}_nF_{n-1}$  equations [Beukers & Heckman 1989].

# Grothendieck's conjecture for combinatorics

Suppose that we have guessed a linear differential equation  $L(f) = 0$  (by differential Hermite-Padé approximation) for some power series  $f \in \mathbb{Q}[[x]]$ , and that we want to recognize whether  $f$  is algebraic or not.

**Recipe 1:** try algebraic guessing.

**Recipe 2:** For several primes  $p$ , compute  $p$ -curvatures mod  $p$ , and check whether they are zero; equivalently, test if  $\partial^p \bmod L = 0 \pmod{p}$ .

► For many power series coming from counting problems (diagonals, constant terms, integrals of algebraic functions, ...) Grothendieck's conjecture is true.



# Grothendieck's conjecture at work

Chebychev in his work on the distribution of primes numbers used the following fact

$$u_n := \frac{(30n)!n!}{(15n)!(10n)!(6n)!} \in \mathbb{Z}, \quad n = 0, 1, 2, \dots$$

This is not immediately obvious (for example, this ratio of factorials is not a product of multinomial coefficients) but it is not hard to prove. The only proof I know proceeds by checking that the valuations  $v_p(u_n)$  are non-negative for every prime  $p$ ; an interpretation of  $u_n$  as counting natural objects or being dimensions of natural vector spaces is far from clear.

As it turns out, the generating function

$$u := \sum_{\nu \geq 1} u_n \lambda^n$$

is algebraic over  $\mathbb{Q}(\lambda)$ ; i.e. there is a polynomial  $F \in \mathbb{Z}[x, y]$  such that

$$F(\lambda, u(\lambda)) = 0.$$

However, we are not likely to see this polynomial explicitly any time soon as its degree is 483,840 (!)

(excerpt from Rodriguez-Villegas's "Integral ratios of factorials")

- ▶ Algebraicity of  $u$  can be however guessed using any prior knowledge, by computing  $p$ -curvatures of the (minimal) order-8 operator  $L$  s.t.  $L(u) = 0$
- ▶ For  $p < 300$ , they are all zero, except when  $p \in \{11, 13, 17, 19, 23\}$

# $G$ -series and global nilpotence

**Definition:** A power series  $\sum_{n \geq 0} \frac{a_n}{b_n} x^n$  in  $\mathbb{Q}[[x]]$  is called a  $G$ -series if it is  
(a) D-finite; (b) analytic at  $x = 0$ ; (c)  $\exists C > 0, \text{lcm}(b_0, \dots, b_n) \leq C^n$ .

**Basic examples:** (1) algebraic functions [Eisenstein 1852]

(2)  $-\log(1 - x) = \sum_{n \geq 1} x^n/n$  ([Chebyshev 1852]  $\text{lcm}(1, 2, \dots, n) \leq 4^n$ )

(3)  ${}_2F_1 \left( \begin{matrix} \alpha & \beta \\ \gamma \end{matrix} \middle| x \right), \alpha, \beta, \gamma \in \mathbb{Q}$

(4) OGF of any  $P$ -recursive, integer-valued, exponentially bounded, sequence

**Theorem** [Chudnovsky 1985] The minimal-order linear differential operator annihilating a  $G$ -series is **globally nilpotent**: for almost all prime numbers  $p$ , it right-divides  $\partial^{p\mu}$  modulo  $p$ , for some  $\mu \leq \deg_{\partial} L$ .

(this condition is equivalent to the **nilpotence** mod  $p$  of the  $p$ -curvature matrix)

**Examples:** algebraic resolvents; Gauss's  $x(1 - x)\partial^2 + (\gamma - (\alpha + \beta + 1)x)\partial - \alpha\beta x$ .

# Global nilpotence for combinatorics

Suppose we have guessed (by differential approximation) a linear differential equation  $L(f) = 0$  for a power series  $f \in \mathbb{Q}[[x]]$  which is a  $G$ -series (typically, the OGF of a  $P$ -recursive, integer-valued, exponentially bounded, sequence).

A way to **empirically certify** that  $L$  is very plausible:

**Recipe:** compute  $p$ -curvatures mod  $p$ , and check whether they are nilpotent; equivalently, test if  $\partial^{pr} \bmod L = 0 \pmod{p}$ , where  $r = \deg_{\partial} L$

**Example:**

```
> L:=x^2*(64*x^4+40*x^3-30*x^2-5*x+1)*Dx^3+
  x*(576*x^4+200*x^3-252*x^2-33*x+5)*Dx^2+
  4*(1+288*x^4+22*x^3-117*x^2-12*x)*Dx+384*x^3-12-144*x-72*x^2:
> p:=7; for j to 3 do N:=rightdivision(Dx^(3*p),L,[Dx,x])[2] mod p;
  p:=nextprime(p); print(p, N); od:
```

11, 0

13, 0

17, 0

# Overview

## Today

1. Introduction
2. High Precision **Approximations**
  - Fast multiplication, binary splitting, Newton iteration
3. Tools for **Conjectures**
  - Hermite-Padé approximants,  $p$ -curvature

## Tomorrow morning

4. Tools for **Proofs**
  - Symbolic method, resultants, D-finiteness, creative telescoping

## Tomorrow night

- Exercises with Maple