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DHWM: A Scheme for Managing Watermarking Keys in the Aquarelle Multimedia Distributed System

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Abstract. This paper presents secure architecture and protocols for managing Intellectual Property Rights in distributed content databases in a close environment. The implementation of this architecture is currently being realized in the European project AQUARELLE. Registered users will access on the Internet to high value content through secure servers. The main interest of this paper is protocols and architecture developed for using watermarking technologies, with a clever and efficient key management based on the Diffie-Hellman (DH) protocol and Trusted Third Parties (TTP).

This paper presents a short survey of watermarking technologies. Next Aquarelle background is specified, along with the chosen watermarking algorithm, which is convenient for the project. Next the DHWM key exchange is presented, based on the simple idea that watermarking and verification can be separated. This scheme uses the Diffie-Hellman key-exchange protocol. Next some hints on the implementation of the scheme and on its correctness are given.

Keywords: IPR protection, watermarking, Key exchange, Aquarelle, multimedia distributed system.

1 Introduction

Watermarking (or fingerprinting or stamping) digital images is a new topic in the security domain. Roughly, this technology consists in hiding an invisible and robust mark into an image. This information should be sufficient to identify the copyright owner of the image: watermarked images can be traced to find their originator or their owner. This is clearly recognized of importance in the context of World Wild Web publishing. The given state-of-the-art does not allow copyright owners to protect their images after diffusion, and many services are blocked in their development because of the ease of reproduction of digital data.

Watermarking, or *embedding* is a very new and complex technology, which hardly seems by now scalable to the whole Internet. Indeed, the technology is

not as strong as classical cryptology, and reasonable attacks can be attempted with success.

This paper describes the solution devised for a closed environment designed for the access of the European Cultural Database. This work is done for the Aquarelle European Project. In this project, multimedia data is more clearly defined, and users are also well known and identified. In that context, a reasonable solution can be tailored.

A trusted third party *TTP* is introduced. Its role is to check the watermarked images. A first scheme is presented, and our scheme, named DHWM, improves that scheme by making use of the Diffie-Hellman protocol.

The paper is structured as follows: the second section briefly surveys techniques relatives to watermarking, and fixes some terminology. Third section presents Aquarelle and its context, giving user requirements and technicals constraints. In section four, the algorithm from UCL is sketched. Section five introduces the DHWM functional model, using the Diffie-Hellman protocol. Section six presents how the scheme is implemented in the Aquarelle prototype.

2 A short survey and a terminology

Many authors have proposed as many algorithms for “marking”, “fingerprinting”, “data hiding”, “steganography”, “label embeddings”, “watermarking” etc. We will first set a terminology with the most common concepts. We will then survey the functionalities claimed by different propositions, since the objectives are quite different depending on the authors. We will not discuss the technicals properties of the algorithms (robustness, invisibility, speed ...) but only their aims and objectives.

2.1 Methods from classical cryptography

It appears that the classical uses of cryptographic techniques are not able to prevent fraudulent use of the images:

encryption encryption can only protects the images during their transmission.

With encryption, an eavesdropper does not have access to the on-line image when it is transfered. But when the user has deciphered the image, then this image does not have any copyright protection anymore.

signature the owner of the image may electronically sign the image (with a hash function and a signature algorithm), but since the signature is added as a suffix to the image, it can easily be removed by anyone who gets the image.

It appears that additional tools are needed. Watermarking is one of these, it will be explained in next sections.

2.2 Terminology

steganography is a very generic concept that consists in hiding messages in a way that eavesdroppers or any monitors do not even know that there is a communication and a message is being sent. Many data hiding techniques were invented for this purpose. Those techniques inspired the development of watermarking algorithms for copyright protection [9].

watermarking is the robust embedding of a copyright information (e.g. time and date, copyright identifiers or simply a correlation pattern) into a content. This content may be a text [12] [1], an audio content [16], but most of the time watermarking is applied to still or moving images. This paper is focusing on images watermarking for both still pictures and motion pictures applications. In the current state-of-the-art, watermarking uses symmetric keys, in the sense that a secret key is used to hide data in a robust way and the same key is used to retrieve the data.

fingerprinting consists in uniquely marking and registering each copy of the data. This marking allows a distributor to detect any unauthorized copy and trace it back to the user. Fingerprinting englobes most often data hiding techniques and cryptographic protocols. Fingerprints have to resist to collusions attacks. It must be very difficult for a set of users to collaborate together and alter fingerprints by merging their copies. Data hiding techniques for fingerprinting can be for instance watermarking techniques but data can also be physically hidden in the media that support the data [14].

2.3 Watermarking technologies

In this section, we make a short overview of most popular watermarking methods from different universities and companies. In the following “survey”, we describe roughly the methods, but we mainly focus on the functional aspects of the methods. Section 4 will describe the watermarking algorithm that has been chosen in the Aquarelle project.

- In [4], a “visible watermark” is introduced. It clearly identifies the ownership, and allows all image details to be seen through it. It is robust enough such that any attempt to remove it alters the image. However, the main drawback is of course that it reduces the quality of the picture.
- The algorithm introduced in [15] is envisaged to have application in image tagging, copyright enforcement, counterfeit access and controlled access, although the authors do not explain how to use their algorithm to perform these functionalities.
- In [1], electronic marking is applied to textual document, by word or line shifting. An indiscernible codeword is added to the document, and it identifies the registered user to whom the document has been delivered. This can be applied on non-ascii text representation, and not to images.
- In [10,17], the following requirements for an invisible copyright label are defined:

1. The image must contain a label or code, which marks it as the property of the copyright holder.
 2. The image data must contain a user code, which verifies that the user is in legal possession of the data.
 3. The image data is labeled in a manner which allows its distribution to be tracked. Unfortunately, the invisibility of the watermark is not totally guaranteed. The watermark is embedded in chosen DCT coefficients of 8x8 blocks. In order to be resistant against compression, the chosen DCT coefficients have to be quantified, which means that marked blocks are altered. This feature can be damageable for high quality pictures such as museum images.
- The authors of [5] consider watermarking using spread spectrum [8]. Their mark identifies ownership and the user who got the image. This method has good robustness properties, however, both the original image and the marked one are needed to check the mark.
 - In [2], an algorithm for invisibly marking an image is introduced. Here again, it is needed to have both the original image and the watermarked one to check the mark.
 - In [16], the authors focus on the notion of data hiding. They envisage the application to the problem of copyright proving and to the content integrity, but the paper does not describe the functional aspects clearly enough.
 - The paper [3] is cryptology oriented and makes abstraction of the marking algorithm. The authors discuss the main problem of the *fingerprinting* technique.
 - In [14], the author describes a global scheme to trace people who abuse broadcast encryption schemes and introduce the interest of fingerprinting. Nevertheless, this paper remains theoretical.
- Finally, there exist now quite a few companies involved in the area of watermarking, such as DIGIMARC, Signum Technologies, R3S, Mediasec Technologies or CRL. Some bigger companies are also currently developing watermarking techniques, such as IBM, AT&T, SONY, NTT, Matsushita, NEC, Philips. It is quite difficult to collect information about their technologies.

3 Aquarelle framework

3.1 Aquarelle aims and technical objectives

Documentation - in a broad sense - is becoming one of the major productions of museums and cultural organizations. To organize exhibitions and produce information products as books or CD-ROM, cultural organizations, museums, libraries, photo-agencies, research laboratories or publishers, have to share information. The aim of Aquarelle is to present a global system for accessing this information.

The Aquarelle users are museum curators, urban planners, commercial publishers and researchers. Aquarelle will provide the user with tools for searching

information and browsing in the folders available on the network of connected servers.

The main technical objectives of the Aquarelle project are the following: to develop a unified resource discovery system for the cultural heritage information available in archive and folder databases; to provide facilities supporting information access through hypertext navigation as well as information retrieval by querying.

The archive server databases contain the images that we want to protect. We consider that organizations running archive servers either own the images they contain, or distribute images belonging to another entity. In both cases, we consider that the organization running an archive server is willing to protect the images which are on the archive.

Serious threats result from the facilities of Aquarelle, and more generally of Internet-based information systems. The ease of getting copyrighted images through the system poses challenges for traditional intellectual property regimes. In the Aquarelle system, all users are registered, and have a legal access to the digital resources. So the main danger is further dissemination of the images by legal users, be it voluntary or not.

The Aquarelle architecture is quite intricate and sophisticated. For our purpose, we only present the following simplification of the architecture in figure. The front-end user uses a standard browser or the Aquarelle advanced browser. He connects to the "User Client Server WEB", which provides front pages and various cgi-bin. The "access server" is the key entrance to the Aquarelle system. In that place, users are registered, and data is transmitted through that node. Through the Z39.50 protocol, data may come directly from the archive servers, or from the "folder server", where folders contain meta-data, and may be published by publishers or cultural organizations.

3.2 Users requirements and system constraints

Cultural partners were concerned with the possibility that their images could be re-used in an unauthorized way. For instance a pirate can make fraudulent copies of images for selling cheap CD-ROMs. Their main concern was to be able to prove, if presented with such fraudulent copies, their ownership of images. Such a possibility is a deterring threat to potential cheaters.

Because of such an objective, they wanted the watermarking system to present a high level of efficiency for protection. This means that the embedded mark must be very robust. Basic cryptographic commandments imply that the algorithm must be parameterized with some key. Following the strong requirements for cultural partners, a unique key will be used for each image.

Now we briefly describe the Aquarelle architecture. Users connect to the system through an *Aquarelle Access Server*. Users use their user terminal to connect on access servers, where they can log in. Once logged they are able to formulate their query. This query is broadcasted to folders and archive servers, and a result set is presented to the user. Connection between users and Access Servers is performed through the HTTP protocol. Connection between Access

Servers and core data servers are performed using the Z39.50 protocol. The login and password are managed by the Access Server.

For the Z39.50 connection between the Access Server and the core data Servers, a login and user password is provided. This enables to authenticate the Access Servers with respect to the core data servers. *There is no user authentication at this level.* This a simple security mechanism to ensure that only authorized connections have occurred to the core data servers.

Since users are unknown at the archive server level, there is no possibility for fingerprinting images here. Fingerprinting can not beat the access server level, since information cannot be cross-compared between users queries and delivered information. Furthermore access servers are seen as being too busy at logging users, managing connections, formulating queries, collecting and assembling results sets. They are not able to perform an expensive on-line operation as fingerprinting. It is also easier to implement off-line watermarking at the archive server using an easy to manage software piece.

4 Properties of the algorithm

4.1 Rough description

The watermarking technique used here is based on Human Visual System Model that guarantees that the watermarked picture has the same quality as the original. The watermark is a correlation pattern, which has strong correlation properties. The technique is additive since this watermark is added to the original to ensure protection.

What is added to the image ? Basically, the watermark is composed of 16 preprocessed replications of one basic correlation pattern. This pattern, that allows to identify copyright ownership, is an image of the same size as the original picture. The 16 replications are modulated and preprocessed before composing the watermark that will be added to the original. This process will be described more in detail in Section 4.1.

The basic pattern is composed of black and white rectangle of pixels. Each rectangle stands for one bit. Those bits form MLS sequences. MLS sequences [11] are binary sequences having very good correlation properties, since MLS sequences are nearly orthogonal to their shifted versions. This feature is taken into account during the retrieval process. The pattern is illustrated in Figure 1.

An additional security feature was added, the sequences bits are pseudo-randomly mixed before being mapped into the rectangles, with the use of a secret key.

How is it added? The 16 parts of the watermark are generated independently. Each part consists of a modulation of a basic pattern at a secret frequency and a secret orientation, determined by a secret key and a pseudo-random generator. For this frequency and this orientation, a perceptual mask is computed. Figure 2

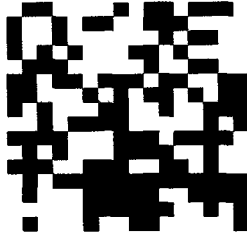


Fig. 1. Example of one basic pattern

shows one example of Lena filtered with a perceptual mask. It serves to adjust the level of the modulated pattern to have it invisible when added to the image. The white regions correspond to areas where the activity is high at the corresponding frequency.

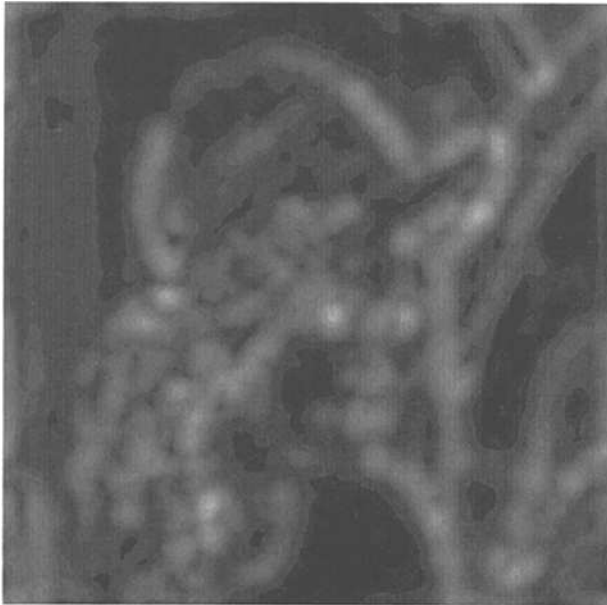


Fig. 2. Energy of Lena in middle frequencies

Figure 3 shows the original, the watermark and the watermarked image, which is the addition of these first two.

Finally, Figure 4 is the global scheme summarizing all these operations.

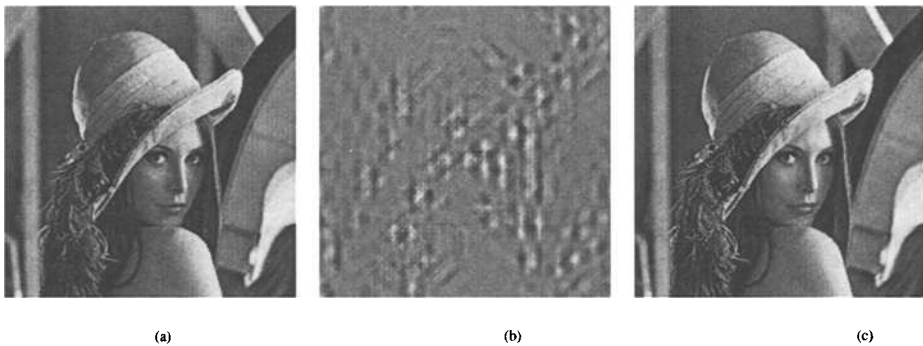


Fig. 3. (a) is the original image, (b) is the watermark, (c) is the watermarked image

How is it decoded? The retrieval procedure is simple. Each part of the watermark is extracted from the watermarked image, by demodulation and filtering, before being added together. The result is an image which is very correlated with the basic pattern if the watermark was present. Autocorrelations are compared to cross-correlations (correlations with shifted MLS) to determine whether an image has been watermarked.

4.2 Robustness

The robustness is provided by the use of MLS and the perceptual mask that allows to embed at a higher level in high activity regions of the image.

compression When the image is compressed at JPEG 10% the watermark can still be recovered, though the quality of the compressed image is very bad.

filtering The watermark is still recovered after low-pass filtering (e.g. blurring 7x7).

printing The watermark is still recovered after half-tone printing. After redigitizing by scanning, we still recover the watermark.

cropping and scanning In this case we need to know the size of the original to retrieve the watermark.

4.3 Invisibility

The quality of the watermarked image is the same as the original image, thanks to the use of perceptual masking.

4.4 Functional aspects

The algorithm, as described in the previous subsections, is well suitable for our purpose.

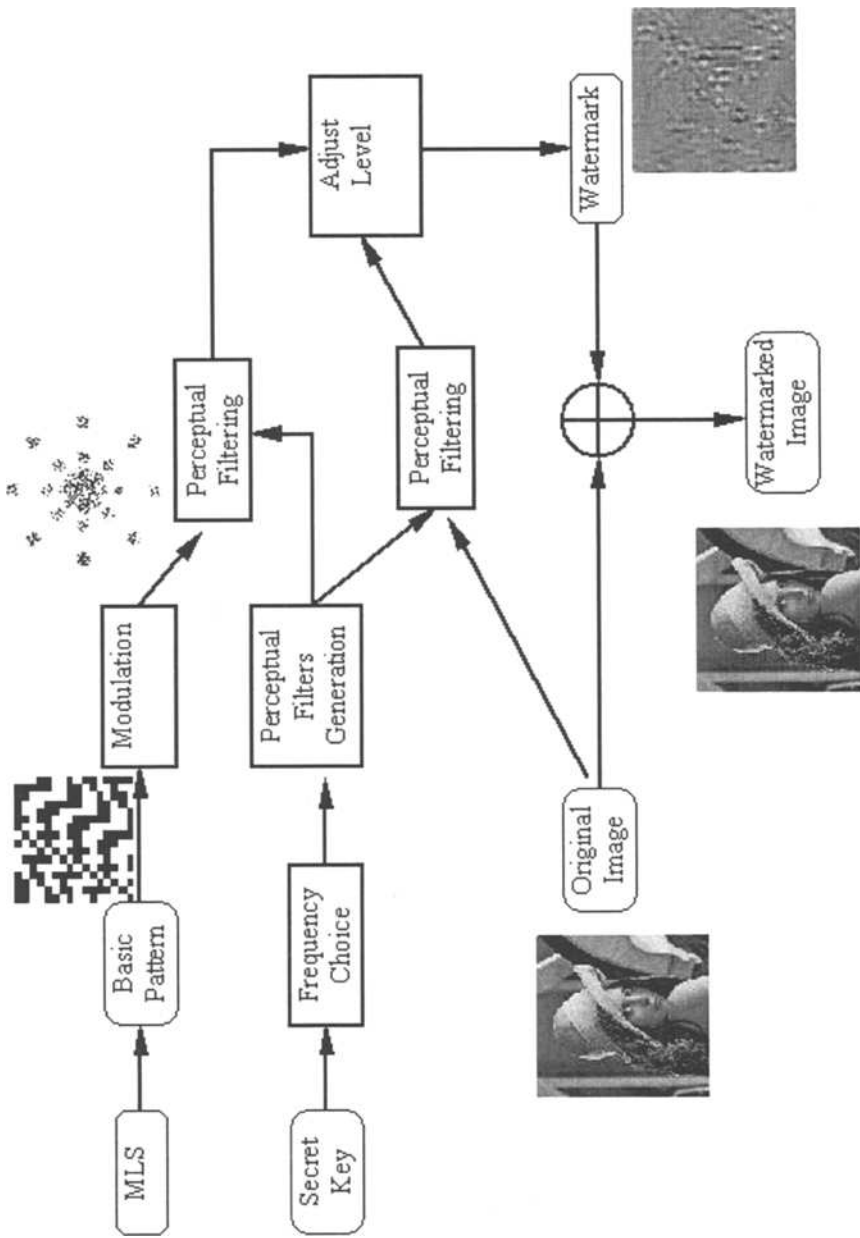


Fig. 4. Global embedding process

- It is secured by a secret key.
- It is optimized for high quality still pictures.
- The decoding procedure is a Yes or No decision that determines whether the image is watermarked or not.
- The watermark resists perfectly to classical image processing in image editing and distribution.

5 Functional models

5.1 The trusted third party

We think that any watermarking algorithm must be public, but parameterized by some key. It does not behave under a static way, since all images would then be embedded along the same scheme (even if the algorithm is *adaptive* to the image). Should the algorithm be delivered in compiled form, then reverse engineering applies to discover the principle of the algorithm. So we draw the conclusion that any watermarking algorithm must be parameterized with some key. In such a way the algorithm can be made public, and all the secrecy resides in the key.

So we consider keyed embedding algorithms such that the knowledge of the embedding key K is needed to verify the watermark. Such an algorithm can offer two modes of operation for verification:

1. The owner reveals the key K to a verifier. The verifier runs the decoding algorithm to check that the image has been marked with the key K .
2. The owner does not reveal the key K , and runs the algorithm for himself.

In the first case, the incusted image is not reusable, since the key K has been shown, and anyone knowing K is able to remove the mark.

In the second case, the owner may be a liar, since from an outside point of view, it seems only that the owner is running a black box which outputs YES. He can not be trusted.

We solve these issues by introducing a Trusted Third Party, the TTP, which plays the following role:

- The TTP knows the secret key K .
- The TTP will never reveal the key K .
- The TTP runs the decoding algorithm, outputs the answer and never lies.

Furthermore the TTP is highly secure, from many points of view (see Section 6.2). The secret K can not be violated, and there can be no impersonnification of the TTP.

Note It is important to note that the TTP introduced here is not a registration authority of copyright-ownership. The TTP will trust the copyright-owners who wish to use its services, and will not check whether the image belongs or does not belong to the copyright-owner using its services. We shall see that the TTP can defeat image-owners trying to cheat and to use its services for watermarking *already protected images*.

5.2 Entities

In Aquarelle we will consider the following entities:

TTP	the Trusted Third Party
CO	the owner of the copyright of IM.
CO-ID	a string which is the unique image identifier of CO
IM	the original image
IM-ID	a string which is the unique image identifier of IM
D	the date
IM*	the watermarked image
IM**	the watermarked image, eventually modified by some hacker
K-IM	the secret used to perform the embedding for that particular image
User	a sample user of the Aquarelle system

From the Aquarelle point of view, we see the CO as an archive-server manager.

5.3 A first functional model

We present a first functional model for clarity purpose. It is NOT the one which is implemented, but it is useful to understand the next one and its advantages.

The protocol for watermarking using the above algorithm runs in 3 phases:

1. The Copyright-Owner sends IM, IM-ID and CO-ID to the TTP
2. The TTP generates a random key K-IM, watermarks the image with K-IM, and securely keeps $\boxed{\text{IM-ID, CO-ID, D, K-IM}}$ in a table.
3. The TTP sends the watermarked image IM* back to the Copyright-Owner, along with CO-ID, IM-ID.

The Copyright-Owner may now deliver the watermarked image IM* through the Aquarelle system. The verification phase is as follows:

1. A user submits an image IM**, IM-ID and CO-ID.
2. The TTP replies YES or NO.

The date field in the database of secret keys is introduced to prevent the following scenario. An image-owner CO1 wants to cheat: he picks an image that has been marked by CO at date D, and submits it to the TTP with the identifiers CO1 and IM-ID1 for watermarking. Both CO and CO1 are able to have their watermark checked by the TTP. But since CO1 submitted the image after CO, then the date field D1 related to CO1, IM-ID1 is bigger than the date D from the original query, the fraud can be detected.

But the above watermarking protocol has the two following disadvantages. First the image must be transmitted over a secure line for the first phase, since an eavesdropper may steal the unmarked image, which has no protection at that time. Secure line may mean encryption, which is a difficult issue because of

various regulations on that topic in European countries (notably France). The second disadvantage is that there are two exchanges of images between the CO and the TTP, which makes a large amount of data to be transmitted.

The improved watermarking protocol presented below solves these two problems.

5.4 Using the Diffie-Hellman protocol

The improved model uses the Diffie-Hellman key-exchange protocol [6]: it enables two persons to share a common secret, without any secure communication; it gets its security from the difficulty of calculating discrete logarithms in a finite field.

The protocol for watermarking runs in 3 phases (see figure 5).

1. The Copyright-Owner and the TTP share a common secret key K-IM using the Diffie-Hellman protocol (each of them sends to the other his Diffie-Hellman half public key, say K_A for the CO and K_B for the TTP).
2. The TTP securely keeps $\boxed{\text{IM-ID, CO-ID, D, K-IM}}$ secret.
3. The CO marks the image with the key K-IM.

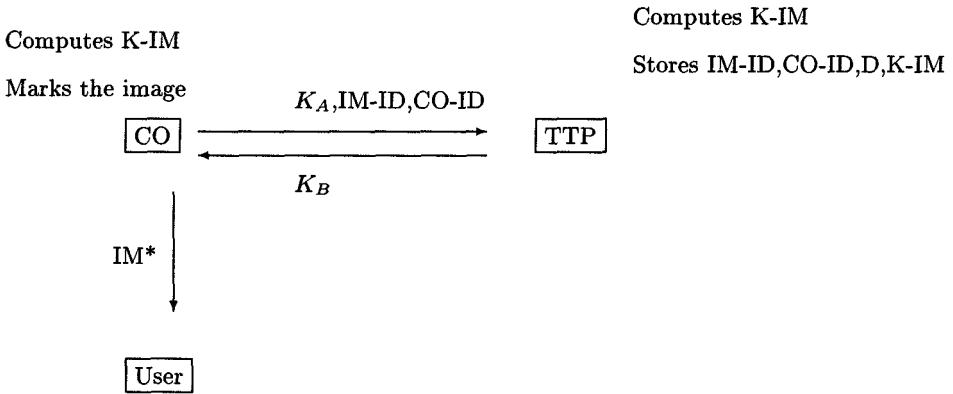


Fig. 5. A scheme for watermarking images with the Diffie-Hellman protocol

This protocol is an improvement of the previous one since no images are exchanged between the CO and the TTP, so there is no need for a secure communication. Second the data exchanged between the CO and the TTP is very small, a few thousands bits, say.

We name this protocol the DHWM protocol, standing for “Diffie-Hellman protocol for Water-Marking”.

6 Integration into the Aquarelle prototype

6.1 Consequences on the Aquarelle architecture

The TTP is implemented here as an HTTP server, and the protocols between the Aquarelle's users and the TTP are built on the HTTP/1.1 protocol.

We note that, ideally, an image is watermarked once for its whole lifetime, and that the process can be done "off-line". So the TTP running the DHWM protocol suite may not be a part of the architecture, and since images are stored on the archive servers, the DHWM protocol needs only to be run between the TTP and the CO of the archive servers.

The watermarking algorithm from UCL focuses on the invisibility of the mark, and, from that point of view, is quite performant. This means that high quality images will not lose their quality after being watermarked. Nevertheless, we think that the COs running the archive servers wish to keep an unmarked copy of their image. In that case, this means that there is a duplication of images: watermarked images to be retrieved through the Aquarelle links and unwatermarked originals that the COs wish to keep by themselves, such that they remain unreachable through the network.

6.2 Security considerations

There remains three main issues to address in the DHWM protocol.

Random numbers Running the Diffie-Hellman protocol, we need random numbers for generating the Diffie-Hellman half public keys. We chose here to use a self-shrinking generator [13]. Known attacks against this kind of pseudo-random generator only apply when the opponent is able to look at a very long string of bits. Here, we only need small pseudo-random strings for our protocol, and we are then protected against the attacks on our generator. Moreover, it is very fast, and this is an important point since the TTP could perform the random numbers generator very often. The weak point is that the state of the machine must be stored in a file, and attacks on this file may be considered. So this file (DH_seed in our implementation) must be protected.

Security of the TTP database It is more obvious that the secrets maintained by the TTP must not be discovered by anyone. A cryptographic solution may consist in encrypting the IM-ID field in the database with a key only known by the TTP, but since the TTP acts automatically, this key must be stored somewhere. So the problem of protecting the file where the key is stored still remains.

We believe that this problem is more related to computer security than to cryptology. In our implementation, the file is simply protected by usual Unix rights, and only the HTTP server is able to read this file.

We leave the problem to computer security specialists, and suggest to use specialized software for this issue. We also suggest to limit the Internet Protocols that are used by the TTP.

Authentication While the Diffie-Hellman protocol is designed to be protected against an eavesdropper (i.e. a passive attack), it does not offer protection against active attacks. We mainly think of authentication: it is a main issue that the CO and the TTP can be ABSOLUTELY assured of each other identity when they run the DHWM protocol to share a common secret key.

Since the protocol is built onto the HTTP/1.1 protocol, any security tool or software for authentication for the World Wide Web is convenient here. We think that the COs must be registered by the TTP, and the TTP must not be subject to an impersonnification attack.

This issue is not covered by the DHWM scheme. Many authentication schemes are proposed for the WEB and for the HTTP protocol. HTTP/1.1 provides a better authentication tool than HTTP/1.0 (login,password). SSL protocol also enables authentication, with heavier tools.

Nevertheless, the Diffie-Hellman protocol can integrate authentication, using a three-round protocol instead of a two-round protocol [7]. This protocol combines Diffie-Hellman key exchange and authentication, and is the basis of the Photuris protocol for IP security.

7 Conclusion

In a system like Aquarelle, where images are distributed to registered users, we protect the images with a watermarking solution. The chosen algorithm from the catholic University of Louvain has very good properties with respect to robustness, invisibility, resistance to JPEG compression. It offers a low functionality, since it actually embeds a single bit of information in an image. Using the DHWM scheme, this small amount of information is turned into a copyright protection system, using a Trusted Third Party.

A simple implementation using an HTTP server has been made. This enables to use any emerging security tool for the WEB to improve the security of the model.

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