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The **DECODING** Library for List Decoding

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1 Introduction and motivation

Reed-Solomon (RS) codes form an important and well-studied family of codes. They were first proposed in 1960 by Reed and Solomon in their original paper [RS60]. They are widely used in practice [WB99]. RS codes can be efficiently unique decoded [Gao02] and [Jus76]. Sudan’s 1997 breakthrough on list decoding of RS codes [Sud97], further improved by Guruswami and Sudan in [GS98], showed that RS codes are list decodable up to the Johnson bound in polynomial time. **DECODING** is a C library whose main goal is to implement as efficiently as possible the Guruswami-Sudan algorithm. It is written in C89 and is stand-alone.

1.1 The Guruswami-Sudan algorithm

The **DECODING** library is devoted to algorithms concerning the Guruswami-Sudan list decoding scheme and does not limit itself to finite fields as it is often the case in coding theory. Let us fix a finite ring with identity $A$ not necessarily commutative. The Guruswami-Sudan algorithm has two main steps. The first one (interpolation step) consists in finding a “curve” of equation $Q(X, Y) = 0$ in $A^2$ which passes through given points with certain multiplicities. The second step (root-finding) finds the roots of $Q(X, Y)$ seen in $(A[X])[Y]$.

The interpolation step dominates the cost of the whole Guruswami-Sudan algorithm. Many methods have been proposed but without any available implementations or comparisons to other ones. A few timings of the Guruswami-Sudan algorithm can be found, but again without the corresponding implementation.

2 The implementation

To the knowledge of the author no implementation of the Guruswami-Sudan algorithm has been proposed. The only available implementation is constituted by a set of C++ functions, not directly accessible, inside PERCY++ [Gol07] whose purpose is not error correction and which does not use fast algorithms for dense bivariate polynomials.

2.1 The algorithms provided by decoding

The implemented algorithm for interpolation is a variant of the Koetter algorithm [McE03] in the *include/decoding/algos/koetter.c* file. It uses polynomial arithmetic with fast bivariate shifting (computation of $Q(X + x_0, Y + y_0)$ where $(x_0, y_0) \in A^2$) in *include/decoding/algos/dbpol_shift_fast.c* and fast univariate shifting (computation of $f(X + x_0)$ where $f \in A[X]$ and $x_0 \in A$) in *include/decoding/algos/upol_shift_fast.c*. Specific variants of these algorithms for commutative rings of characteristic 2 are also present in the same files.
The second step (root-finding) implemented in the library is a variant of the Roth and Ruckenstein algorithm \cite{RR98} and the naive algorithm of \cite{BLQ11}. It is in include/decoding/algos/dbpol_Xroots.c.

2.2 The design of decoding

The \textsc{decoding} library is designed to be easy to use in a C or C++ program. One of its particularities is to use the C preprocessor to generate algorithms for a ring (generally a finite field) which must be provided by the end-user. Therefore efficient libraries like MPFQ \cite{GT06} can be used by the end-user. For the sake of completeness, some finite fields are provided by default.

Error correcting code are often regarded over finite fields, in particular \( \mathbb{F}_2 \), together with the classical Hamming distance. But other distances, like the Lee distance, are better suited for some applications. Usually the Lee distance is needed for codes over Galois rings. Error correcting codes over the ring of matrices over a finite field or a finite commutative ring are also considered for example in \cite{OSB12}. Therefore \textsc{decoding} proposes generic algorithms whenever possible. This flexibility is needed when studying codes over Galois rings for example where the end-user needs to manipulate codes over a Galois ring and its residue field at the same time.

Although \textsc{decoding} proposes certain fast bivariate polynomial algorithms, it is not its goal to propose fast algorithms for univariate and bivariate polynomial multiplication. In fact, \textsc{decoding} is designed to be used in conjunction with other efficient libraries like \textsc{Gmp} \cite{Gra91}, \textsc{NTL} \cite{Sho90} or \textsc{Flint} \cite{Har10}. For the sake of completeness \textsc{decoding} provides these algorithms in their “schoolbook” form but it is recommended, for efficiency, to use external libraries.

A very simple mechanism using the C preprocessor allows one to override the default generic algorithms proposed by \textsc{decoding}. For example see at the include/decoding/rings/GF5.c file which implements the finite field \( \mathbb{F}_5 \). It shows how to replace the univariate polynomial root finding over \( \mathbb{F}_5 \). All C macros that control this mechanism are in the include/decoding/ring_reset.h file.

3 Presentation

The \textsc{decoding} library is the first library which proposes a flexible and efficient way to implement algorithms related to the Guruswami-Sudan decoding scheme. It can be used with efficient external libraries to obtain more efficient implementations of Guruswami-Sudan related algorithms.

I will first present quickly the history of the Guruswami-Sudan algorithm and show that it needs dense bivariate polynomials only available, not necessarily directly, in computer algebra systems such as MAGMA \cite{BCP97} or MATHEMAGIX \cite{H++02}. As error correcting codes are often used over binary fields, dedicated fast algorithms must be used. The bivariate polynomials appearing in the list decoding algorithms can have large degrees even when RS codes with small parameters are considered. Hence fast algorithms are needed.

I will then present the flexibility of \textsc{decoding}, needed to obtain efficient algorithms over several finite rings and fields.

• It is easy to replace a key algorithm, such as univariate polynomial multiplication or univariate polynomial root-finding, by a very efficient one provided by an external library such as \textsc{Flint} or \textsc{NTL} for example.

• It is easy to choose a finite ring or a finite field, or even to use different rings at the same time in order to implement algorithms related to RS codes over Galois rings. The \textsc{decoding} library is not restricted to mathematical object whose binary representation holds in a single machine word. It requires no supplementary efforts to use, for example, multiple precision integers from \textsc{Gmp} or large binary fields from \textsc{Mpfq}.
Finally, I will present the provided algorithms concerning dense bivariate polynomials and their applications to list decoding.

References


