

Online Sparse Bandits

David L. Saint-Pierre, Quentin Louveaux, Olivier Teytaud

▶ To cite this version:

David L. Saint-Pierre, Quentin Louveaux, Olivier Teytaud. Online Sparse Bandits. The 3rd Asian Conference on Machine Learning (ACML2011), Nov 2011, Taoyuan, Taiwan. hal-00642461

HAL Id: hal-00642461 https://inria.hal.science/hal-00642461

Submitted on 18 Nov 2011

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers. L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

David L. St-Pierre¹, Quentin Louveaux¹, and Olivier Teytaud^{2,3}

Online Sparse bandit

 $^{1}\,$ Department of Electrical Engineering and Computer Science, Faculty of Engineering, Liège University, Belgium

² TAO (Inria, Lri, Univ. Paris-Sud, UMR CNRS 8623), France

³ OASE Lab, National University of Tainan, Taiwan

What is a zero-sum Matrix Game (MG)?

M = KxK matrix, coefficients in [0,1] Player 1 chooses i in $\{1,2,3,...,K\}$ Player 2 chooses j in $\{1,2,3,...,K\}$ Player 1 gets reward Mij Player 2 gets reward 1-Mij

Is solving a MG hard?

= polynomial time (by linear programming) best known coef 3.5

How to find approximate solutions?

Grigoriadis & Khachiyan or Exp3 or Inf: time O(Klog K / ϵ 2) with proba ½ (Auer et al, Audibert et al, Grigoriadis et al)

Sparse versions?

Very often, (x^*,y^*) is very sparse; plenty of 0's. How to benefit from this?

Algo. in Flory et al.:

- 1. Approximate solving by t iterations of EXP3
- 2. Remove small components (keep only components $\geq (\max(tx))^{(4/5)} / t)$
- 3. Re-normalize ==> no proof

We propose the following online version ==>

(EXP3 recalled below...)

```
Algorithm 1 EXP3 algorithm for iteration t with K arms.
  Initialise \forall i, p(i) = \frac{1}{K}, n(i) = 0, S(i) = 0; t = 0
  \mathbf{while}\ t < T\ \mathbf{do}
      Arm i is chosen with probability p(i)
     n(i) \leftarrow n(i)+1
     Receive reward n
     S_i modified by the update formula S_i \leftarrow S_i + r/p(i) (and S_j for j \neq i is not
     modified)
      \forall i, p(i) = 1/(K\sqrt{t}) + (1 - 1/\sqrt{t}) \times \exp(S_i/\sqrt{t}) / \sum_i \exp(S_j/\sqrt{t})
  end while
  return n
```

What means "solving" a MG?

Strategy x: probability distribution on $\{1,2,3,...,K\}$ Strategy y: probability distribution on $\{1,2,3,...,K\}$ **Expected Reward:** R(x,y)=E Mij $|(i,j)\sim x,y$ Nash equilibrium: $(x^*,y^*) = Nash$ $\leq = >$ for all (x,y) $R(x,y^*) \ge R(x^*,y^*) \ge R(x^*,y)$

Approximate solving?

 (x^*,y^*) ε -approximate Nash equilibrium if for all (x,y), $R(x,y^*)+\varepsilon \ge R(x^*,y^*) \ge R(x^*,y)-\varepsilon$

So what?

There is a offline solution (i.e. sparsity used at the end). Can we use it online?

Algorithm 3 on EXP3, an online EXP3 algorithm with a cut solely based on

```
Initialise \forall i, p(i) = \frac{1}{K}, n(i) = 0, S(i) = 0; t = 0
while t < T do
   Arm i is chosen with probability p(i)
   n(i) \leftarrow n(i)+1
   t \leftarrow t+1
   S_i modified by the update formula S_i \leftarrow S_i + r/p(i)
   \forall i, p(i) = 1/(K\sqrt{t}) + (1 - 1/\sqrt{t}) \times \exp(S_i/\sqrt{t}) / \sum_i \exp(S_i/\sqrt{t})
   if x_i > \lceil \frac{t}{K} \rceil and x_i < (b_1 \times T^{\delta} \times (\frac{t}{T})^{\beta}) then
      Remove arm i
   end if
   if every arm has been pruned then
      Use plain EXP3
   Renormalize: p = p / \sum_{i} p(i)
end while
Execute the truncation TEXP3 as presented in 2.3
return n
```

Conclusions?

(i) it works (see numbers in paper) (ii) theory missing (iii) better (parameter-free?) versions



<== application: Urban Rivals (free, you can test!) Next application: Pokemon ==>

