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Pairwise Testing of Dynamic Composite Services

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ABSTRACT

Online services encapsulate enterprises, people, software systems and often operate in poorly understood environments. Using such services in tandem to predictably orchestrate a complex task is one of the principal challenges of service-oriented computing. A composite service orchestration soliciting multiple atomic services is plagued by a number of sources of variation. For instance, availability of an atomic service and its response time are two important sources of variation. Moreover, the number of possible variations in a composite service increases exponentially with increase in the number of atomic services. Testing such a composite service presents a crucial challenge as its often very expensive to exhaustively examine the variation space. Can we effectively test the dynamic behavior of a composite service using only a subset of these variations? This is the question that intrigues us. In this paper, we first model composite service variability as a feature diagram (FD) that captures all valid configurations of its orchestration. Second, we apply *pairwise testing* to sample the set of all possible configurations to obtain a concise subset. Finally, we test the composite service for selected pairwise configurations for a variety of QoS metrics such as response time, data quality, and availability. Using two case studies, *Car crash crisis management* and *eHealth management*, we demonstrate that pairwise generation effectively samples the full range of QoS variations in a dynamic orchestration. The pairwise sampling technique eliminates over 99% redundancy in configurations, while still calling all atomic services at least once. We rigorously evaluate pairwise testing for the criteria such as: a) ability to sample the extreme QoS metrics of the service b) stable behavior of the extracted configurations c) compact set of configurations that can help evaluate QoS tradeoffs and d) comparison with random sampling.

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General Terms

Performance, Experimentation, Measurement, Reliability

Keywords

Web Services, Pairwise Testing, QoS

1. INTRODUCTION

In a service-oriented world actors such as data sources, knowledge bases, people, processes, businesses, hardware sensors/actuators and software systems are all seen as services. In such a world, a *composite* service orchestrates a number of self-contained *atomic* services to perform complex tasks. The unpredictable and dynamic nature of each of these atomic services ultimately renders the functional and non-functional behavior of a composite service unpredictable and dynamic. For instance, the crisis management system in a large city orchestrates a number of atomic services such as the ambulance service, police service, GPS service, and phone service. The variable nature of each of these services renders the overall behavior of the crisis management system variable and dynamic.

Untested dynamic behavior of a composite service can have several critical consequences. For instance, a crisis management system dealing with an earthquake must mobilize a multitude of services within a predictable time frame and seldom deviate from it. An untested composite service may exhibit unreliable deviations from contractual agreements on Quality of Service (QoS) [22]. Service level agreements (SLAs) [16] are the industry standard to specify constraints on QoS for both service providers and consumers. Habitual deviations from SLAs are a result of ignoring QoS outliers and dynamic behavior of a composite service.

A key challenge in testing a composite service emerges from its inherent variability. We enlist three important dimensions to *composite service variability* (a) The variation in selection/non-selection of equivalent atomic services used in a composite service (b) The variation in QoS of each of these atomic services leads to variations in composite service QoS. For instance, in [19] we develop probabilistic models of QoS variability in atomic services (c) The variation in the way atomic services are called in a composite service such as in sequence or in parallel. In this paper, we are primarily concerned with the first two sources of variability. Changes in the orchestration can be triggered by these sources, enabling self-* composite services. We use the general term dynamic

composite service to encompass self-* service-oriented systems.

With an increase in number of equivalent atomic services there is an exponential increase in the invocations of a composite service. It is impractical and computationally expensive to test a composite service for all its possible variations. Therefore we ask, can we effectively test the dynamic behavior of a composite service using only a *subset of these variations*? Answering this question is the subject of this paper.

We present a methodology for combinatorial interaction testing (CIT) dynamic composite services. In particular, we perform *pairwise testing* of composite services. The methodology consists of three main phases: (1) Modeling variability in a composite service (2) Generation of composite service configurations satisfying pairwise interactions (3) Analyzing these composite service configurations to test composite service QoS. In our approach, we model the variability of a composite service as a feature diagram where each feature represents an atomic service. Inter-feature constraints represent dependencies between atomic services. Feature Diagrams (FD) [10] provide a formal framework to specify authorized variations in the configuration of a composite service. We transform the feature diagram and pairwise interactions between features (or atomic services) to a single constraint satisfaction problem in the formal specification language Alloy [8]. We solve the Alloy model to generate valid configurations of the composite service. The generation methodology is an extension of our previous work [18] to dynamic composite services. We empirically investigate the QoS of the resulting configurations. We demonstrate that combinatorial interaction testing (CIT) [5] to select a subset of configurations that covers *all valid pairwise interactions* of services is an efficient technique to sample configurations of an orchestration. Our premise is based on the observation that most software faults are triggered by interactions between a small number of variables [13]. For example, consider the *car crash crisis management system* case study [12] that we will examine in this paper. With 25 optional features that may / may not be invoked in a specific orchestration, the total number of exhaustive tests required will be 33, 554, 432. This is an extremely large number of tests that would considerable time and effort for QoS analysis. The number of tests satisfying pairwise interaction is just 185 reducing the number of required tests by **99.99%**.

Pairwise testing has been used to detect faults in software systems in extensive prior research [5]. Our main contribution is the application of pairwise testing to sample configurations in dynamic composite services: one form of efficient self-monitoring of variable behavior. This is based on the hypothesis that composite services' QoS behavior uncover faults in a service-oriented systems where choice of atomic services and the orchestration between them are primary artifacts. The extensive empirical studies, based on two case studies which are the car crash crisis management system (*C³MS*) [12] and a eHealth administration system, support our **claims about pairwise testing** dynamic composite services:

1. **C1:** Pairwise testing is an sufficient coverage strategy for dynamic composite service orchestrations
2. **C2:** Pairwise testing covers a wide range of QoS in dynamic composite services
3. **C3:** Pairwise testing is better than random testing
4. **C4:** Pairwise testing is a stable strategy to define global SLA for a dynamic composite service
5. **C5:** Pairwise testing is useful to generate families of orchestrations with differing SLAs

The paper is organized as follows. Section 2 provide foundational material to understand our paper. This includes feature diagrams in 2.1, the Orc language for specifying orchestrations in 2.2, pairwise configuration generation in 2.4, and formal description of QoS metrics in 2.5. The methodology followed in this paper is discussed in Section 3. The case studies for experiments are put forth in Section 4. The experiments related to QoS analysis are presented in 5. Comparison with respect to random generation and the stability of pairwise analysis are shown in 5.3 and 5.4, respectively. Further deliberation and perspectives of our analysis scheme are presented in Section 5.5. Threats to the validity of the empirical studies are discussed in Section 5.6. Related work in literature is put forth in Section 6. We conclude in Section 7.

2. FOUNDATIONS

In this section we present background or foundational ideas required to understand the rest of the paper. We present these concepts to make the paper as self-contained as possible.

2.1 Feature Diagrams

Feature Diagrams (FD) introduced by Kang et al. [10] compactly represent all the products (referred to as *configurations* in this paper) of a software product line (SPL) in terms of features which can be composed. Feature diagrams have been formalized to perform SPL analysis [21]. In [21], Schobbens et al. propose a generic formal definition of FD which subsumes many existing FD dialects. We define a FD as follows:

- A FD consists of k features f_1, f_2, \dots, f_k
- Each feature f_i may be associated with a software asset such as an atomic service.
- Features are organized in a parent-child relationship in a tree T . A feature with no children is called a leaf.
- A parent-child relationship between features f_p and f_c are categorized as follows:
 - *Mandatory* - child feature f_c is required if f_p is selected.
 - *Optional* - child feature f_c may be selected if f_p is selected.
 - *OR* - at least one of the child-features $f_{c1}, f_{c2}, \dots, f_{c3}$ of f_p must be selected.
 - *XOR* - one of the child-features $f_{c1}, f_{c2}, \dots, f_{ck}$ of f_p must be selected.
- Cross tree relationships between two features f_i and f_j in the tree T are categorized as follows:
 - f_i requires f_j - The selection of f_i in a product implies the selection of f_j .
 - f_i excludes f_j - f_i and f_j cannot be part of the same product and are *mutually exclusive*.

2.2 Service Orchestrations using Orc

A dynamic composite service is an orchestration of atomic services. We express the orchestration of atomic services available in an FD using the Orc language. Orc [15] serves as a simple yet powerful concurrent programming language to describe and execute service orchestrations.

The fundamental declaration used in the Orc language is a *site*. The type of a *site* is itself treated like a service - it is passed the types of its arguments, and responds with a

return type for those arguments. An Orc *expression* represents an execution and may call external services to publish some number of values (possibly zero).

Orc has the following combinators that are used on various examples as seen in [15]. The *Parallel* combinator $F|G$, where F and G are Orc expressions, runs by executing F and G concurrently. Whenever F or G communicates with a service or publishes a value, $F|G$ does so as well. The execution of the *Sequential* combinator $F >x> G$ starts by executing F . Sequential operators may also be written compactly as $F \gg G$. Values published by copies of G are published by the whole expression, but the values published by F are not published by the whole expression; they are consumed by the variable binding. If there is no response from either of the sites, the expression does not terminate. While the above two composition operators are for creating threads, Orc uses the following construct to prune operations. The *Pruning* combinator, written $F <x< G$, allows us to block a computation waiting for a result, or terminate a computation. The execution of $F <x< G$ starts by executing F and G in parallel. Whenever F publishes a value, that value is published by the entire execution. When G publishes its first value, that value is bound to x in F , and then the execution of G is immediately terminated. The *Otherwise* combinator, written $F;G$ has the following execution. First, F is executed. If F completes, and has not published any values, then G executes. If F did publish one or more values, then G is ignored. The publications of $F;G$ are those of F if F publishes, or those of G otherwise. In the *Fork-Join* combinator, two processes are invoked and run concurrently. The process waits until a response is obtained from both. This may be represented as (F,G) where the process waits for responses from both atomic services F and G .

2.3 Feature Diagrams with Orchestrations

The FD and the orchestration cover two dimensions that are complementary to each other. While the FD represents the variability in the configurations, the orchestration specifies the order in which the services are called. Making use of the terminology in [21], *primitive* features are “features” that are of interest and that will be incorporated in real-world services. On the contrary, *decomposable* features are just intermediate nodes used for decomposition. It is up to the modeler to determine such classification of features in the FD. We extend the semantics given in [21] to ensure compatibility of an orchestration with the feature model as follows:

- The set of available services S are the *primitive* nodes of the FD D ;
- For each orchestration, the set of corresponding services invoked (denoted N);
- $N \subseteq S$ in a configuration;
- A model of D is a subset of its (*primitive* and *decomposable*) nodes;
- There must exist a model of D ($[[D]]$) such that $[[D]] \cap S = N$ (a model of a FD is a subtree that is valid w.r.t. the operators and the dependence relation).

Drawing from the real-world services and the constraints shown in a FD, the composite service may be developed by an orchestrator.

2.4 Combinatorial Interaction Testing

We use combinatorial interaction testing (CIT) to synthesize a subset of configurations represented by the FD of a dynamic composite service. Originally, CIT was proposed by Cohen et al. [5] to select a subset of all combinations of

variables that define the input domain of a program, while still guaranteeing a certain level of coverage. This has led to the definition of pairwise interaction testing, or 2-wise testing. This samples the set of all combinations in such a way that all possible pairs of variable values are included in the set of test data. Pairwise testing has been generalized to t -wise testing which samples the input domain to cover all t -wise combinations. In this paper, a set of test data is often represented in the form of a *covering array* that contains all t -wise interaction of features in a FD.

DEFINITION. 1. *Covering Array* - A covering array $CA(N; t, k, v)$ is a $N \times k$ array on v symbols with the property that every $N \times t$ sub-array contains all ordered subsets of size t from v symbols at least once.

From the definition of a covering array, the strength t of the array is the parameter that allows achieving 2-wise (pairwise), 3-wise or t -wise combinations. The k columns on this array correspond to all the variables in the input domain which in our case are the features in a FD. For the generation of dynamic composite service configurations, k is the number of services, and v is 2 since we have only boolean variables (services may be present or absent in a configuration). The covering array is a set of configurations of features.

We demonstrate the concept of a minimal covering array using an example. Consider the set of four atomic services (A, B, C, D) with varying response times. The atomic services can be composed in 2^4 exhaustive combinations. However, if we consider the service combinations in pairs, we require fewer configurations. These can be subsumed by 6 sets of configurations that cover these pairs of interactions resulting in removal of 62.5% of redundancies. This is shown in Table 1 where, for example, interaction (A, B) refers to calling both service A and B while (A, \neg B) refers to calling only A with B explicitly not invoked.

Pairwise Interaction	Configurations
(A, B); (A, C); (A, D); (B, C); (C,D)	(A, B, C, D)
(A, \neg B); (A, \neg C); (A, \neg D)	(A)
(B, D); (B, \neg A); (B, \neg C); (D, \neg A)	(B, D)
(C, \neg A); (C, \neg B); (C, \neg D)	(C)
(D, \neg B); (D, \neg C)	(A, D)
(B, \neg D)	(A, B, C)

Table 1: Subsuming pairwise interactions in configurations

Essentially, the use of pairwise sampling reduces the number of cases needed to generate a range of outputs, a few of which that may be considered faulty. Consider a system S having a set of inputs p and a set of outputs q . With random testing, in which input vectors satisfying p are randomly generated, and the output of each execution is compared with the postcondition q as a set of tests. As structural features of system S are hidden, the efficacy of using manually designed test cases can be seen mainly through their cost effectiveness. In our case, we view this as the decrease in the number of samples needed to generate extreme output values (faults).

Let $\omega \in p$ be a set of tests for the system S . This produces a set of specifications $\omega \xrightarrow{S} q'$, where $q' \in q$. A successful set of tests is one that has a minimal cardinality of cases $|\omega|$ and maximal variance in the set of outputs q' . This generates a range of values as the system output. Empirical studies have shown pairwise sampling to be superior for precisely such a case - efficiently generating a minimal set of tests to generate all dual combinations of input values. This in turn produces a range of outputs q' that have higher variance than other comparative techniques of similar cardinality $|\omega|$.

The problem of generating a minimal covering array for a set of variables is a complex optimization problem that has been studied in extensive prior work for example [5]. It is important to notice that there exist very few studies that have tackled the automatic generation for CIT in the presence of constraints between variables [6]. In order to include properties that forbid combinations of values, CIT generation techniques have to allow the introduction of constraints in the algorithms that generate covering arrays. In recent work [18], we present a solution to generate t-wise configurations that satisfy all simultaneously constraints modeled in a feature diagram.

We transform the feature diagram to constraint satisfaction problem model in the language Alloy as described in [18]. The features in the FD are transformed to concepts in Alloy called *signatures*. Inter-feature constraints in the FD are transformed to Alloy *facts*. All pair-wise interactions between features are transformed to Alloy *predicates*. The goal of solving the Alloy model is to find the minimal set of configurations that cover conjunctions of all valid pair-wise predicates. The first step involves *detection* of all valid pairs that conform to the FD. In the second step, we construct conjunctions of pair-wise predicates and solve them via incrementally increasing the scope of the solution size. The result is a minimal set of configurations that cover conjunctions of all valid pairs. At times the SAT solver in Alloy is not scalable for a large FD. We apply divide-and-compose approaches as described in [18] to handle this scalability issue.

2.5 QoS Aspects of the Orchestration

In this paper, we test dynamic composite services for their probabilistic QoS behavior. In this section we summarize our work in [19], that presents the derivation of composite service QoS behavior from individual atomic service behaviors. Probabilistic analysis of QoS parameters as described in [19] [7] provide a more realistic study of actual services' behavior. The following QoS parameters have been chosen for experiments in this paper:

1. *Latency / Response Time* (T) - Denotes the overall delay due to the time taken by a service to respond. It is a discrete value that may be modeled as a long tailed distribution incorporating some *rare deviations*.
2. *Availability* (α) - The probability that a service is active and can respond to a service call. For a well managed service, this value is generally quite high.
3. *Cost* (χ) - Refers to the monetary cost associated with each invocation of a particular atomic service.
4. *Data Quality* (ξ) - A subjective measure of trade off to high Cost and Response times of services. It measures the "Quality" of the output of the service and the beneficial aspects of including a new atomic service into the composite orchestration.

These QoS metrics are normally defined for an atomic service. We derive these QoS metrics for a dynamic composite service by analyzing its orchestration. This analysis involves giving a semantic to a composite service QoS based on individual atomic service QoS and the Orc combinators (see Section 2.2) associating them. Taking two sites s_i and s_j , the QoS metrics may be computed as shown in Table 2 based on the Orc combinators in use. The cases of composing the service s_{ij} using the *sequential* and *fork-join* combinators have been considered. The latency, cost and availability metrics for the composite service s_{ij} are derived as shown in [4] with $Max(p, q)$ representing the maxima of the values p and q . For the sequential case, the latency and cost of the composite service is a sum of the atomic services' parameters while the

availability is a product of such parameters. Similarly, the maxima of the atomic services' response times contributes to the global response time under parallel invocation.

Orc Code	$s_{ij} \triangleq s_i \gg s_j$	$s_{ij} \triangleq (s_i, s_j)$
Latency	$T(s_{ij}) = T(s_i) + T(s_j)$	$T(s_{ij}) = Max(T(s_i), T(s_j))$
Cost	$\chi(s_{ij}) = \chi(s_i) + \chi(s_j)$	$\chi(s_{ij}) = \chi(s_i) + \chi(s_j)$
Availability	$\alpha(s_{ij}) = \alpha(s_i) \times \alpha(s_j)$	$\alpha(s_{ij}) = \alpha(s_i) \times \alpha(s_j)$

Table 2: QoS metrics extended to Orc combinators.

Some QoS metrics of an atomic service may be modeled as a random variable conforming to a probability distribution. We need to simulate the QoS metric by sampling from a probability distribution. For instance, we need to simulate the probabilistic response time distributions of each atomic service as done in [19]. By varying the degrees of freedom ν and non-centrality parameter δ in the t-distribution *dfittool* of MATLAB, it is possible to generate various heavy tailed distributions that mimic the response times of services. We sample these distributions to simulate the response times of actually invoked atomic services. In this paper, the t-distribution fitting was used to generate various distributions of services' response times.

3. METHODOLOGY

We present the methodology for pairwise testing and QoS analysis of dynamic composite services.

1. **Inputs:** The inputs to our methodology is tuple (S , FD , O , $Strategy$):
 - (a) S is the set of all atomic services that can be used in a dynamic composite service.
 - (b) FD is a feature diagram that specifies various features in a dynamic composite service and the constraints between them. Primitive features in an FD are each associated with an atomic service S_i . A valid configuration C_k of a FD is the set of m features f_1, f_2, \dots, f_M that conform to the constraints in the FD . The features in valid configurations represents sets of atomic services S_1, S_2, \dots, S_N . The sets are subsets of S . See Section 2.1 for formal definition of a FD .
 - (c) O is the overall orchestration of the dynamic composite service. The orchestration is reconfigured based on valid configurations of the FD . The orchestration O may be reconfigured to orchestrations O_1, O_2, \dots, O_N for all valid configurations C_1, C_2, \dots, C_N of the FD . An orchestration only invokes the set of atomic services present in a valid configuration of the FD . In our paper, O is an Orc orchestration. See Section 2.2 for brief description of Orc.
 - (d) $Strategy$ is the strategy used to generate configurations. In this paper, we consider two strategies to guide generation of valid FD configurations:
 - i. Random Generation : We randomly select configurations conforming to FD by solving the Alloy model representing only the FD .
 - ii. Pairwise Generation : We generate a set of configurations that satisfy all pairwise interactions between features in FD . These configurations also satisfy the constraints in the FD .
2. **Configuration Generation:** We generate the configurations using the technique described in [18] and briefly outlined in Section 2.4. The process involves transformation of the FD to a constraint satisfaction

problem in Alloy. A chosen *Strategy* to generate configurations is also transformed in conjunction with the Alloy model. Solving the Alloy model gives valid configurations. Let the set of output configurations be C_1, C_2, \dots, C_N for a chosen strategy *Strategy*.

3. **Empirical Analysis of QoS:** The output configurations from the previous step C_1, C_2, \dots, C_N reconfigures O to orchestrations O_1, O_2, \dots, O_N by selecting only the atomic services that are present in each of the configurations. We compute QoS for each of the orchestrations invoking all atomic services in the configuration using the semantics described in Section 2.5. We use the experiments to address the questions motivated in Section 1.

4. CASE STUDIES

We consider two case studies for our experiments as described in Sections 4.1 and 4.2.

4.1 Car Crash Crisis Management System

The need for crisis management systems has grown significantly over time [12]. Crisis management involves identifying, assessing, and handling the crisis situation. A crisis management system facilitates this process by orchestrating the communication between all (distributed) parties involved in handling the crisis. The car crash crisis management system (C^3MS) [12] includes all the functionalities of a general crisis management systems, and some additional features specific to car crashes such as facilitating the rescuing of victims at the crisis scene and the use of tow trucks to remove damaged vehicles. As described in [12], the main goals of this system include: a) Facilitating the rescue mission carried out by the police / firemen and providing them with detailed information on the location of the crash. b) Managing the dispatch of ambulances or other alternate emergency vehicles to transport victims from the crisis scene to hospitals. c) Coordinating the first-aid missions by providing relevant medical history of identified victims by querying data bases of local hospitals. d) Ushering the medical treatment process of victims by providing important information about the crash to the concerned workers. e) Managing the use of tow trucks to remove obstacles and damaged vehicles from the crisis scene.

In Figure 1, we present the Car Crash Crisis Management System (C^3MS) FD [12]. The C^3MS FD contains several features that are associated with software assets represented by atomic services. For example, the *Paramedic* feature is represented by the *Paramedic* service. Some sets of features like *Police* and *PoliceMan* are subsumed by a single service *Police*. Constraints such as optional, requires and mutual exclusion (XOR) are also incorporated. For example, the *GPS* and *GSM* features are mutually exclusive while the *Doctor* feature requires the *PublicHospital* feature.

A host of services may be used to represent the C^3MS . The corresponding services referring to the features are shown in Fig. 2. Services are invoked either *sequentially* or in *parallel*(with synchronization merge) as shown in the workflow. It is assumed that the services perform the functions as generally specified by the nomenclature. For example, the *CommunicationManager* service manages the communication between parties while the *Ambulance* service regulates ambulances to the car crash sites. Their construction may be modified according to specifications to perform/subsume other associated tasks.

4.2 eHealth Management System

The need for efficient hospital management stems has been discussed in [20]. A hospital administration system is devised to remove some of the inefficiency plaguing current

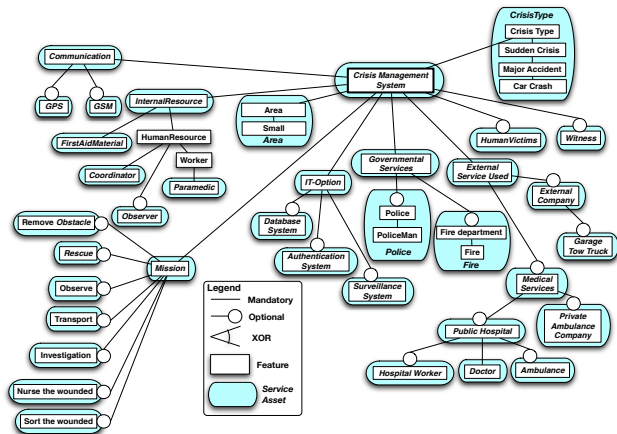


Figure 1: C^3MS Feature Diagram.

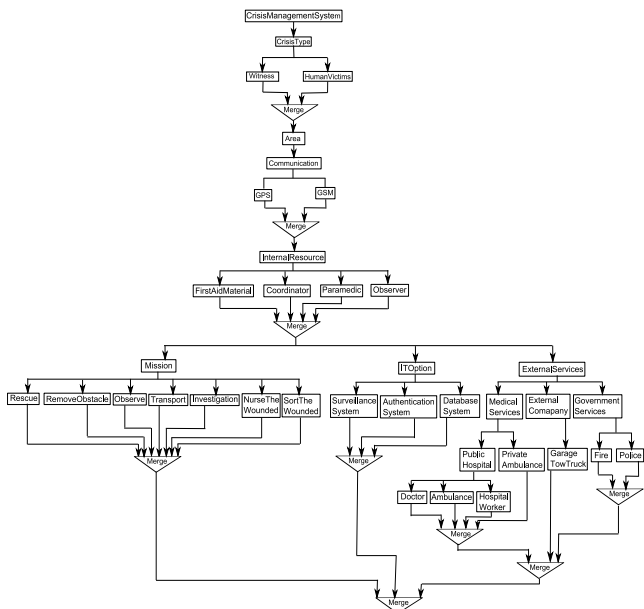


Figure 2: Composite Service Orchestration of the C^3MS .

protocols such as cumbersome admission time, duplicate data entry, redundant lab tests, ineffective treatment coordination, and billing processes. Drawing inspiration from [20] composite health care applications are required to connect various parties and locations. The information flows seamlessly across organizational and system boundaries emitting from the use of such a centralized orchestration. This enhanced visibility gives everyone involved a unified view of relevant information and gives process owners the ability to improve existing methods and procedures. The eHealth system can be viewed as an extension of the C^3MS medical services to transport injured victims for speedy treatment of injuries. Examples of the utility of healthcare applications include: a) Healthcare providers can access the medical information of a prospective patient and use ambulance services to transfer the client to relevant healthcare facilities. b) Physicians can review a patient's medical history even though this data resides in several systems managed by diverse providers. c) Insurance claims and financial options can be updated and handled in a speedy way. d) Doctors

can use a composite application to determine the appropriate medication for a patient, order the drug, check the status of pharmacy approval, and monitor how the drug is dispensed. e) Special needs of the patient such as catering specific food items and lab tests can be coordinated in an effective way.

Fig. 3, presents the eHealth management system FD. Similar to the C^3MS FD, it contains several features that are associated with software assets represented by atomic services. Constraints such as optional, requires and mutual exclusion (XOR) are also incorporated. Two versions of the similar service $Ambulance_f$ and $Ambulance_s$ are in mutual exclusion. These atomic features or services can be set to varying QoS values resulting in interesting combinations of services.

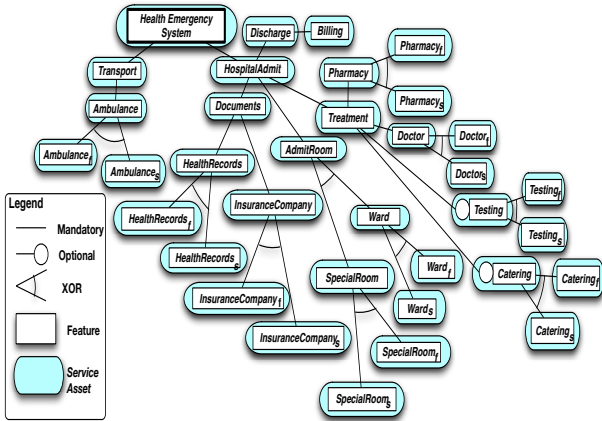


Figure 3: eHealth Feature Diagram.

The services used for the orchestration of the eHealth system are shown as follows. The operations are generic with services such as $HealthRecords$ and $InsuranceRecords$ used to request relevant medical history and insurance status of the patient, respectively.

The Orc pseudo code for the eHealth system is presented with the distinguishing feature being the choice of services that can be used to perform similar goals. For instance, either one of the mutually exclusive (MUX) services $Testing_f()$ or $Testing_s()$ services can be used to request for lab tests. However, the QoS associated with each of these services is different resulting in varying overall composite service QoS.

--eHealth Orchestration--

```
def HealthEmergencySystem() = Transport() >> HospitalAdmit()
  >> Billing() >> Discharge()

def Transport() = a >> Ambulance(a)
def HospitalAdmit() = Documents() >(hf,in) (HealthRecords(hf),
  InsuranceRecords(in)) >> AdmitRoom() >> Treatment()
def AdmitRoom() = (sr,w) >> (SpecialRoom(sr),Ward(w))
def Treatment() = (d,t,c,p) >> ((Doctor(d),Testing(t),
  Catering(c),Pharmacy(p))

def Ambulance(a) = let(Ambulance_f() | Ambulance_s())
def HealthRecords(hf) = let(HealthRecords_f() |
  HealthRecords_s())
def InsuranceCompany(in) = let(InsuranceCompany_f() |
  InsuranceCompany_s())
def SpecialRoom(sr) = let(SpecialRoom_f() | SpecialRoom_s())
def Ward(w) = let(Ward_f() | Ward_s())
def Doctor(d) = let(Doctor_f() | Doctor_s())
def Testing(t) = let(Testing_f() | Testing_s())
def Catering(c) = let(Catering_f() | Catering_s())
def Pharmacy(p) = let(Pharmacy_f() | Pharmacy_s())
```

5. EXPERIMENTS

Based on the methodology in Section 3 we perform experiments involving pairwise generation of configurations followed by simulations to obtain probabilistic QoS of dynamic composite services. We consider both case studies for these experiments.

5.1 Evaluating QoS of the Car Crash Crisis Management System

Configuration Generation: We first use the approach presented in [18] to generate a minimal set (given the resource constraints) of configurations that satisfy all valid pairwise interactions in the C^3MS case study. The input settings to the configuration generator are (a) Maximum scope for Alloy solver (b) Maximum time to solve (c) Divide-and-compose strategy for scalable generation. The maximum scope is set to 8 and maximum time to 2000 milliseconds with use of *incremental growth* strategy. Through this technique, **185** configurations for the C^3MS case study were generated. The 185 configurations satisfy all valid pairwise interactions between services in the C^3MS FD that originally specify 2^{25} configurations. All invalid pairs that do not conform to the FD are rejected by the approach. For instance, the not including the *Mission* feature in a configuration is invalid as it is a mandatory feature.

Computing Response Time: Second, we compute response times for these 185 configurations. We assign each atomic service in the dynamic composite service a t-distribution to model response time. The random settings for the atomic service t-distributions were degrees of freedom ν from 3 to 8 and non-centrality parameter δ from 5 to 15 seconds, respectively. We choose these values to provide diversity in atomic response times. For a chosen atomic service (in the current configuration), the individual timeout value was set to 95 percentile of the response time distribution. This largely ensures that the composite service obtains the result of the atomic service and not a timeout. For each of the 185 configuration, we obtain **10,000** Monte-Carlo samples of response times from all atomic services in a configuration. We compute the composite service response time from these atomic service response times. We collect the response times for the composite service for each configuration to create a t-distribution for the composite service. We set the global timeout of the composite service to a sufficiently high value (300 seconds) to allow capture of outliers in the distribution.

As seen in Fig. 4, the pairwise generated configurations cover a range of response time distributions. The distributions were sorted in increasing order of response time and are shown. The slowest and the fastest composite services are marked. Their median values are shown to be **113** and **201** seconds, respectively. This demonstrates the use of a few configurations to test significant changes of about **88** seconds response time in a composite service. These results support the claims **C1** and **C2** in Section 1, that pertain to the effectiveness of pairwise sampling to generate a wide range of orchestrations and output QoS values.

Computing other QoS metrics: We compute additional QoS metrics such as availability of a service, the cost entailed in calling atomic services and output data quality for the 185 configurations. We compute QoS for a composite service based on rules given in Table 2 for different Orc combinators in an orchestration. For example, when we set atomic service availability to **0.99** (representing service availability in 99% of invocations) the composite availability of each configuration is shown in Fig. 5. The output data quality ξ is related to the cost χ by the constant κ given by $\xi = \chi/\kappa$ (assuming linear increase in data quality with each atomic service invocation). The output data quality ξ is can also be derived exponentially from the cost χ by

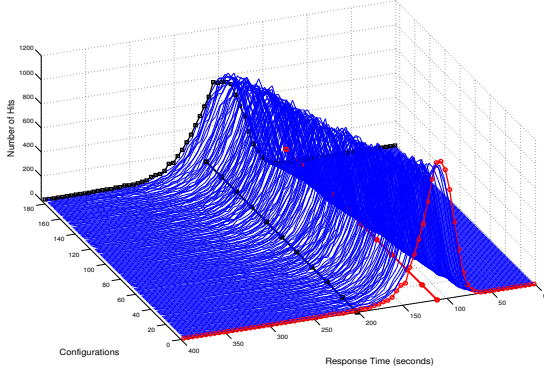


Figure 4: Response time distributions of the 185 pairwise configurations for C^3MS .

$\xi = e^{\chi/\kappa}$. For example, setting the $\chi = 5$ units for each invoked atomic service, the cost of each configuration is shown in Fig. 5. Furthermore, setting $\kappa = 20$, the linear and exponential output data quality of the configurations may also be derived. These variations in data-quality, response time and cost help analyze trade-offs between QoS parameters. These variations in QoS parameters substantiate the claim **C5** about pairwise testing in Section 1 referring to its use in generating families of composite services.

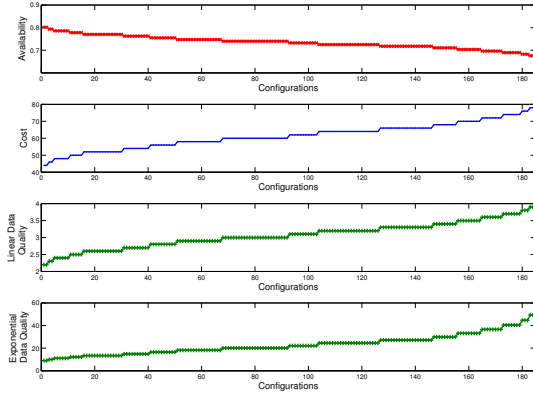


Figure 5: Availability, Data Quality and Cost of the pairwise configurations of C^3MS .

5.2 Evaluating QoS of the eHealth System

Configuration Generation: For the eHealth system, we generate 188 configurations that satisfy all valid pairwise interactions from a total set of 2^{12} configurations. The initial settings for configuration generation were exactly the same as in the C^3MS case study.

Computing Response Time: For each of the 188 configurations, we model atomic service QoS as t-distributions. The parameters of these distribution are chosen in random in certain bounds to ensure diversity. The parameter degrees of freedom ν was from 3 to 8 and non-centrality parameter δ from 5 to 15 seconds, respectively. For the faster services (marked with the subscript f), the δ parameter was set between 3 to 5 seconds, representing a faster response to a service call.

We obtain **10,000** Monte-Carlo samples of response times for each of the atomic services and compute the composite service response time distribution. As seen in Fig. 6, the pairwise generated configurations cover a wide range of response time distributions. The distributions are sorted in increasing order of response time. The slowest and the fastest composite services are marked with median values. In the case of eHealth, the 30 seconds range in response time values is due to the added diversity of *choice* in choosing a *fast* or *slow* atomic service.

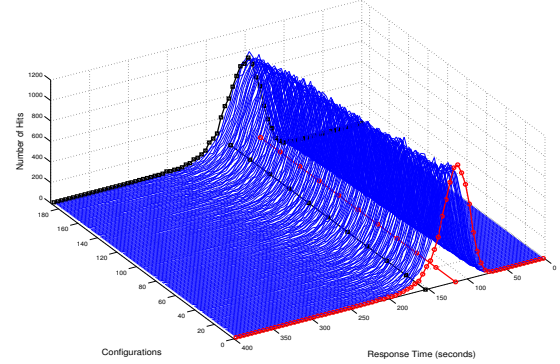


Figure 6: Response time distributions of the 188 pairwise configurations for eHealth.

Computing other QoS metrics: We use the rules for combinators described in Table 2 to compute QoS of composite service orchestrations. Setting atomic service availability to 0.99 the composite availability each configuration is shown in Fig. 7. We observe that the cost of the composite service varies with the choice of fast or slow services. A faster service (with subscript f) is set *double* the cost of its slower (with subscript s) counterpart. This changes the range of cost and data quality available for different configurations as seen in Fig. 7.

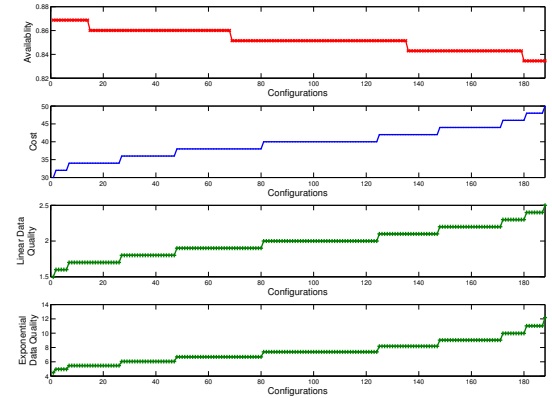


Figure 7: Availability, Data Quality and Cost of the pairwise configurations of eHealth.

5.3 Comparison with Random Sampling

It has been shown in [5] that pairwise interaction testing of such configurations is advantageous over random testing since its systematic and provides a better coverage. With

random runs, it is impossible to determine if all the atomic services have been invoked at least once. The configurations leading to extreme test case values need not be necessarily generated during random runs and there may be many redundant configurations invoked repeatedly. Setting SLAs based on random runs is both non-robust and can lead to habitual deviance. Generating families of configuration with accurately fixed bounds on QoS is also not possible. For these reasons, pairwise generation has comparative advantages over random runs.

Three sets of random configurations were generated as shown in Fig. 8, each with original configuration size 185. In each case, the number of valid configurations was found to be 17, 21 and 24 resulting in a maximum efficient generation percentage of 12.97%. Not only are there deviations in the number of valid configurations for each run (17, 21, 24), but also in the QoS metrics output in each run. SLA deviations are a result of resorting to such insufficient random runs of a composite service, which might generate invalid and redundant scenarios. To test the effectiveness of combinato-

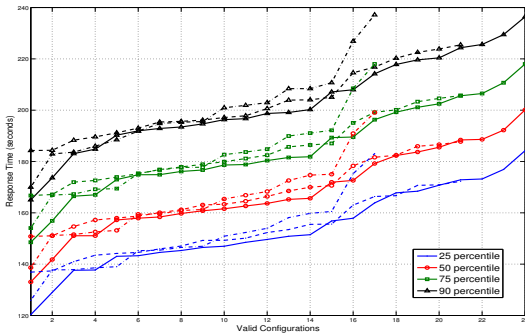


Figure 8: Three runs of random generation of configurations for C^3MS .

rial testing the 185 pairwise configurations were compared with random samples for the C^3MS . All the mandatory features were set to be invoked with the constrained and optional features randomized in invocation for the random case. This random sampling was performed by a Markov decision process of traversing features in the FD, which will always lead to generation of valid configurations (based on constraints). The comparison with pairwise is shown in Fig. 9 and it is seen that random generation can cover a large range of QoS values if sufficient number of configurations are generated. To determine that number, however, requires analysis of pairwise interactions. The random configurations are deficient as they cannot guarantee a) invocation of every possible service at least once; b) generating the extreme configurations for a particular composite service in every sample. When compared to the pairwise generation scheme that covered all pairs of services, the random generation covered only 8.8% of the service pairs. This shows that the same set of services are redundantly invoked in many configurations during random generation. Thus, for such orchestrations with numerous configurations, using pairwise interactions is a sufficient choice in order to examine the entire sample space. These results support our claim **C3** in Section 1, referring to the comparison between pairwise and random sampling.

5.4 Consistency of Pairwise Samples

Given one orchestration, there can be many different sets (or solutions) of configurations that cover pairwise services interactions. Thus, we compute QoS behavior over differ-

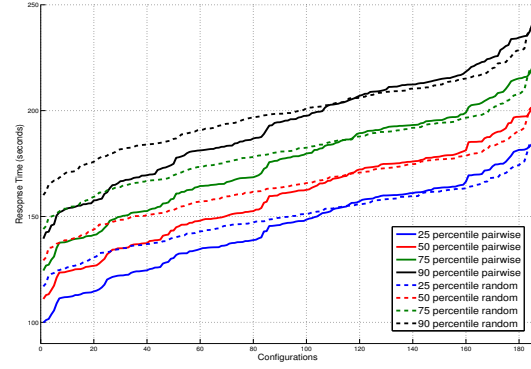


Figure 9: Comparison of pairwise and random response time (arranged in increasing order) of percentile values for 185 configurations of C^3MS .

ent samples of configurations. This aims at evaluating the stability of pairwise testing as a sampling technique to estimate the global QoS for a dynamic composite service. A collection of 10 samples that satisfy the pairwise interaction testing were generated for the eHealth case. The percentile statistics of the configurations in each sample was collected through 10,000 Monte-Carlo runs and is shown in Fig. 10. The lowest and highest percentile values of the configurations in each sample were collected. The mean inter-sample difference for the random case is 12.94 seconds compared to 6.44 seconds for the pairwise case. Further, these were compared with 10 samples of randomly generated configurations (with 300 configurations in each sample) in Fig. 10. Again, all the mandatory features were set to be invoked with the constrained and optional features randomized in invocation for the random configurations. The number of valid configurations for each sample ranged between 3.5% to 9% of the 300 configurations. Comparing the two cases, the stability of the pairwise generation is demonstrated through its consistently low standard deviation values in Table 3 when compared to random samples. Once again, the lowest and the highest percentile values of all the configurations in a particular sample are compared. These results support claim **C4** in Section 1, referring to the stability of pairwise sampling.

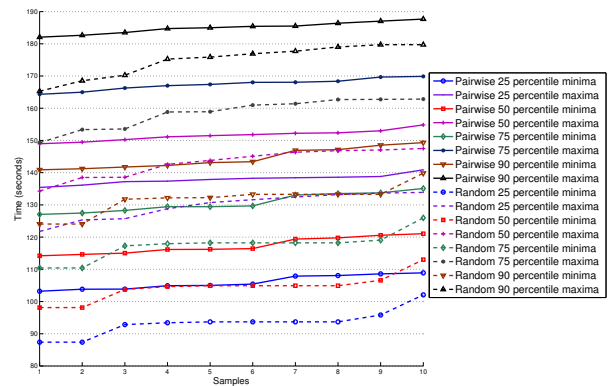


Figure 10: Comparing stability of pairwise and random samples for eHealth.

Percentile	25(min.)	25(max.)	50(min.)	50(max.)	75(min.)	75(max.)	90(min.)	90(max.)
Pairwise Std. Deviation(secs.)	2.18	1.52	2.59	1.73	2.90	1.82	3.19	1.83
Random Std. Deviation(secs.)	4.14	4.17	4.21	4.51	4.43	4.76	4.63	5.07

Table 3: Standard Deviation values for pairwise and random samples.

5.5 Perspectives due to Analysis

The methodology evaluated for the C^3MS and the eHealth orchestrations can lead to many possibilities for improving QoS metrics for composite services. This includes setting the SLA keeping into account the worst performing configuration. This will prevent contract deviation during actual deployment of the service.

A family of SLAs for a set of configurations taking into account trade-offs between QoS metrics and the output quality of configurations may be proposed. This leads to families of composite services with extensively analyzed SLAs. Configurations may be grouped along with their QoS behavior to develop an extended product line of composite services. For example, categories of services may be constructed for the C^3MS orchestration (based on Figs. 4 and 5) as shown in Table 4. Similarly, the two categories of service families for the eHealth case (Figs. 6 and 7) is shown in Table 5. In both cases, the family of services with higher data quality is traded-off by a slightly higher response time.

While the diversity in QoS families for the C^3MS is due to optional services that may / may not be included, the variability in the eHealth case is mainly due to other factors. An inherent choice in replacing a *slow* atomic service with a *fast* counterpart can lead to a range of QoS values. Generated configuration families can use of combination of these options of optimally compose atomic services to specific QoS bounds. These service families can have associated contracts (albeit in the soft-sense as in [19]) to monitor deviations from specifications. These instances support our claim C5 in Section 1, that pertains to developing families of composite service orchestration with significantly different QoS behavior. With numerous possible combinations

Configuration Families	Bronze	Silver	Gold
90 percentile Response Time (T)	< 183 s	< 216 s	> 216 s
Median Response Time (T)	< 150 s	< 179 s	> 179 s
Availability (α)	> 0.75	> 0.71	> 0.71
Cost (χ)	< 60	< 70	> 70
Linear Data Quality (ξ)	< 3	< 3.5	> 3.5
Exponential Data Quality (ξ)	< 20	< 30	> 30

Table 4: Configuration families for C^3MS .

Configuration Families	Standard	Premium
90 percentile Response Time (T)	< 171 s	> 171 s
Median Response Time (T)	< 139 s	> 139 s
Availability (α)	> 0.85	> 0.85
Cost (χ)	< 40	> 40
Linear Data Quality (ξ)	< 2	> 2
Exponential Data Quality (ξ)	< 8	> 8

Table 5: Configuration families for eHealth.

of atomic services, such a dedicated families of services with significantly different QoS outputs enable accurate monitoring of services provided. The pairwise scheme is both a robust and compact representation of the behavior space of the set of orchestrations. This provides an effective pre-SLA technique to enunciate the QoS metrics and threshold levels.

5.6 Threats to Validity

This section considers the threats to the validity of the experimental results. These may be *internal* (whether there is a bias/error in the experimental design which could affect

the causal relationship) or *external* (ability to generalize the results of the experiment to industrial practice).

The hypothesis studied in this paper concerns the use of pairwise sampling to evaluate QoS of large orchestrations. Sources of internal error can be a result of the MiniSAT solver used to generate the pairwise configurations or the MATLAB statistical tools used for QoS evaluation. These tools have not been compared with available alternatives for consistency of results. Furthermore, the assumption is that for each sample of configurations, the pairwise analysis scheme can provide consistently large range of QoS values. Systematic bias in QoS may be introduced in samples when extreme cases are not generated.

To ensure scalability to large industry level FDs, the pairwise generation in [18] makes use of incremental growth / binary splitting schemes. Redundancies in the number of configurations can be seen due to these schemes. For generating more than one sample of solutions, the symmetry breaking scheme in Alloy was used. This introduces more constraints with each proceeding sample, which increases the time required to generate such samples.

6. RELATED WORK

The combinatorial testing framework described by Cohen et al. [5] has been applied extensively to efficient testing for fault detection. In the work of Cohen et al. [6], this technique is extended to software product lines with highly configurable systems. Modeling variability in SPLs using feature models is the work of Jaring and Boschet [9] where they show that the robustness of a SPL architecture is related to the type of variability. To ensure that constraints in the FD are incorporated in the efficient sampling of t-wise tests, the scalable solver proposed by Perrouin et al. [18] is used. In [14], variability in software as a service applications are modeled using the orthogonal variability model to study the customization choices in such workflows.

Pre-deployment testing of SLAs has been studied by Di Penta et al. [17], where they make use of genetic algorithms to generate test data causing SLA violations. Analysis of white and black box approaches are provided in the paper. In [2], Bruno et al. make use of regression testing to ensure that an evolving service maintains the functional and QoS assumptions. The service consistency verification due to evolution is done by executing test suites contained in a XML encoded facet attached to the service.

The use of probabilistic QoS and soft contracts was introduced by Rosario et. al [19] and Bistarelli et al. [1]. Instead of using fixed hard bound values for parameters such as response time, the authors proposed a soft contract monitoring approach to model the QoS measurement. The composite service QoS was modeled using probabilistic processes by Hwang et al. [7] where the authors combine orchestration constructs to derive global probability distributions.

In our paper, we extend these two notions to analyze the QoS of a composite orchestration under various configurations. Effective sampling of orchestrations is necessary specially in conjunction with exceedingly flexible and large configuration spaces. When combined with the probabilistic behavior QoS behavior of services, this provides an accurate portrayal of the composite service's end-to-end QoS. In a recent submission [11], similar methodology is used to compare pairwise and exhaustive analysis of configuration spaces in smaller orchestrations. In this paper, that no-

tion is extended to comparison with random runs of larger configuration spaces (where exhaustive analysis is impossible). This entails a scalable approach for robust pairwise interaction generation that is not required for the smaller examples. The case studies and corresponding experiments are much larger in this paper and study the effect of not only orchestration variability, but also choice in compatible atomic service counterparts. Correspondingly, this requires further experiments on the sampling robustness and comparison with random generation, which is not included in [11].

Though formal analysis of end-to-end QoS has been studied in Cardoso et al. [4], there are no practical testing tools available for the composite service provider. The pairwise testing procedure has been shown to outperform other testing techniques in [5]. We extend this testing tool to develop a generic testing methodology to query end-to-end QoS of a web service. Related empirical studies of optimal QoS compositions make use of genetic programming in Canfora et al. [3] and linear programming in Zeng et al. [23]. These are dynamic techniques to choose the best possible atomic services and configurations for SLAs. The goal in our paper is to analyze the dynamic configurations that may result due to invocation/non-invocation of particular web services when atomic SLAs have already been established.

7. CONCLUSION

We demonstrate that combinatorial interaction testing and in particular pairwise testing effectively portrays the overall behavior of a dynamic composite service. Pairwise testing drastically reduces the number of composite service configurations while successfully analyzing a wide range of QoS values. It provides good coverage for two large case studies (C^3MS and eHealth). We also observe that the analysis remains stable over multiple solutions for the same case study. Pairwise testing is superior to random generation of configurations in terms of coverage and stability of results. Pairwise testing helps specify SLAs based on a deterministic and systematic sampling scheme rather than random sampling. We use our approach to create many families of composite services which can be seen as products with varying costs and SLAs. We largely augment the predictability of a dynamic composite service by performing offline pairwise testing in advance.

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